

# Types for Information Flow Control: Labeling Granularity and Semantic Models (Technical appendix)

Vineet Rajani and Deepak Garg

Max Planck Institute for Software Systems

April 30, 2018

## Contents

<b>1</b>	<b>Part I: Details of what's in the paper</b>	<b>8</b>
1.1	Fine-grained IFC enforcement (FG)	8
1.1.1	FG type system	8
1.1.2	FG semantics	10
1.1.3	Logical relation for $FG$	10
1.1.4	Soundness proof for FG	12
1.2	Coarse-grained IFC enforcement (CG)	73
1.2.1	CG type system	73
1.2.2	CG semantics	73
1.2.3	Logical relation for CG	73
1.2.4	Soundness proof for CG	76
1.3	CG to FG translation	145
1.3.1	Type directed translation from CG to FG	145
1.3.2	Type preservation for CG to FG translation	145
1.3.3	Logical relation for CG to FG translation	151
1.3.4	Soundness proof for CG to FG translation	152
1.4	FG to CG translation	188
1.4.1	Type directed (direct) translation from FG to CG	188
1.4.2	Type preservation for FG to CG translation	189
1.4.3	Logical relation for FG to CG translation	204
1.4.4	Soundness proof for FG to CG translation	205
<b>2</b>	<b>Part II: Alternate development with original HLIO in place of CG</b>	<b>236</b>
2.1	Fine-grained IFC enforcement (FG)	236
2.1.1	FG type system	236
2.1.2	FG semantics	239
2.1.3	Logical relation for $FG$	239
2.1.4	Soundness proof for FG	242
2.2	Coarse-grained IFC enforcement (CG)	329
2.2.1	CG type system	329
2.2.2	CG semantics	329

2.2.3	Logical relation for CG . . . . .	331
2.2.4	Soundness proof for CG . . . . .	333
2.3	CG to FG translation . . . . .	414
2.3.1	Type directed translation from CG to FG . . . . .	414
2.3.2	Type preservation for CG to FG translation . . . . .	415
2.3.3	Logical relation for CG to FG translation . . . . .	424
2.3.4	Soundness proof for CG to FG translation . . . . .	425
2.4	Translation from FG to $FG^-$ . . . . .	469
2.4.1	$FG^-$ typesystem . . . . .	469
2.4.2	Type translation . . . . .	471
2.4.3	Type preservation: FG to $FG^-$ . . . . .	472
2.5	FG to CG translation . . . . .	487
2.5.1	Type directed (direct) translation from FG to CG . . . . .	487
2.5.2	Type preservation for FG to CG translation . . . . .	488
2.5.3	Logical relation for FG to CG translation . . . . .	513
2.5.4	Soundness proof for FG to CG translation . . . . .	514

## List of Figures

1	Type system for FG . . . . .	8
2	FG subtyping . . . . .	9
3	FG semantics . . . . .	10
4	Type system of CG. . . . .	73
5	CG subtyping . . . . .	74
6	CG semantics . . . . .	74
7	Expression translation from CG to FG . . . . .	145
8	Type system for FG . . . . .	237
9	FG subtyping . . . . .	238
10	Well-formedness relation for FG . . . . .	238
11	FG semantics . . . . .	240
12	Type system for CG . . . . .	329
13	CG subtyping . . . . .	330
14	Well-formedness relation for CG . . . . .	330
15	Type system for $FG^-$ . . . . .	470
16	$FG^-$ subtyping . . . . .	471

## List of Theorems

1.1	Lemma (Reflexivity of subtyping)	8
1.2	Definition ( $\theta_2$ extends $\theta_1$ )	10
1.3	Definition ( $W_2$ extends $W_1$ )	11
1.4	Definition (Binary value relation)	11
1.5	Definition (Binary expression relation)	11
1.6	Definition (Unary value relation)	11
1.7	Definition (Unary expression relation)	11
1.8	Definition (Unary heap well formedness)	12
1.9	Definition (Binary heap well formedness)	12
1.10	Definition (Binary substitution)	12
1.11	Definition (Unary substitution)	12
1.12	Definition (Unary interpretation of $\Gamma$ )	12
1.13	Definition (Binary interpretation of $\Gamma$ )	12
1.14	Lemma (Binary value relation subsumes unary value relation)	12
1.15	Lemma (Monotonicity Unary)	15
1.16	Lemma (Monotonicity binary)	16
1.17	Lemma (Unary monotonicity for $\Gamma$ )	18
1.18	Lemma (Binary monotonicity for $\Gamma$ )	18
1.19	Lemma (Unary monotonicity for $H$ )	19
1.20	Lemma (Binary monotonicity for heaps)	19
1.21	Theorem (Fundamental theorem unary)	20
1.22	Lemma (Expression subtyping)	32
1.23	Lemma (Subtyping unary)	34
1.24	Lemma (Binary interpretation of $\Gamma$ implies Unary interpretation of $\Gamma$ )	38
1.25	Theorem (Fundamental theorem binary)	38
1.26	Lemma (Binary heap well formedness implies unary heap well formedness)	67
1.27	Lemma (Subtyping binary)	67
1.28	Theorem (NI for FG)	71
1.29	Definition ( $\theta_2$ extends $\theta_1$ )	73
1.30	Definition ( $W_2$ extends $W_1$ )	75
1.31	Definition (Value Equivalence)	75
1.32	Definition (Binary value relation)	75
1.33	Definition (Binary expression relation)	75
1.34	Definition (Unary value relation)	75
1.35	Definition (Unary expression relation)	76
1.36	Definition (Unary heap well formedness)	76
1.37	Definition (Binary heap well formedness)	76
1.38	Definition (Binary substitution)	76
1.39	Definition (Unary substitution)	76
1.40	Definition (Unary interpretation of $\Gamma$ )	76
1.41	Definition (Binary interpretation of $\Gamma$ )	76
1.42	Lemma (Binary value relation subsumes unary value relation)	76
1.43	Lemma (Monotonicity Unary)	80
1.44	Lemma (Monotonicity binary)	82
1.45	Lemma (Unary monotonicity for $\Gamma$ )	84
1.46	Lemma (Binary monotonicity for $\Gamma$ )	85
1.47	Lemma (Unary monotonicity for $H$ )	85

1.48	Lemma (Binary monotonicity for heaps)	86
1.49	Theorem (Fundamental theorem unary)	86
1.50	Lemma (Subtyping unary)	102
1.51	Lemma (Binary interpretation of $\Gamma$ implies Unary interpretation of $\Gamma$ )	106
1.52	Theorem (Fundamental theorem binary)	107
1.53	Lemma	137
1.54	Lemma (Subtyping binary)	138
1.55	Theorem (NI for CG)	143
1.56	Theorem (Type preservation, $CG \rightsquigarrow FG$ )	145
1.57	Lemma (Subtyping type preservation: CG to FG)	149
1.58	Definition ( ${}^s\theta_2$ extends ${}^s\theta_1$ )	151
1.59	Definition ( $\hat{\beta}_2$ extends $\hat{\beta}_1$ )	151
1.60	Definition (Unary value relation)	151
1.61	Definition (Unary expression relation)	151
1.62	Definition (Unary heap well formedness)	151
1.63	Definition (Value substitution)	151
1.64	Definition (Unary interpretation of $\Gamma$ )	151
1.65	Lemma (Monotonicity)	152
1.66	Lemma (Unary monotonicity for $\Gamma$ )	155
1.67	Lemma (Unary monotonicity for $H$ )	155
1.68	Theorem (Fundamental theorem)	156
1.69	Lemma (Subtyping)	179
1.70	Theorem (Deriving CG NI via compilation)	183
1.71	Definition	188
1.72	Theorem (Type preservation: FG to CG)	189
1.73	Lemma (Subtyping - Type preservation)	202
1.74	Definition ( ${}^s\theta_2$ extends ${}^s\theta_1$ )	204
1.75	Definition ( $\hat{\beta}_2$ extends $\hat{\beta}_1$ )	204
1.76	Definition (Unary value relation)	204
1.77	Definition (Unary expression relation)	204
1.78	Definition (Unary heap well formedness)	204
1.79	Definition (Value substitution)	205
1.80	Definition (Unary interpretation of $\Gamma$ )	205
1.81	Lemma (Monotonicity)	205
1.82	Lemma (Unary monotonicity for $\Gamma$ )	208
1.83	Lemma (Unary monotonicity for $H$ )	208
1.84	Lemma (Coercion lemma)	209
1.85	Theorem (Fundamental theorem)	209
1.86	Lemma (Subtyping lemma)	227
1.87	Theorem (Deriving FG NI via compilation)	232
2.1	Lemma (FG: Reflexivity of subtyping)	236
2.2	Definition (FG: $\theta_2$ extends $\theta_1$ )	239
2.3	Definition (FG: $W_2$ extends $W_1$ )	239
2.4	Definition (FG: Binary value relation)	241
2.5	Definition (FG: Binary expression relation)	241
2.6	Definition (FG: Unary value relation)	241
2.7	Definition (FG: Unary expression relation)	242
2.8	Definition (FG: Unary heap well formedness)	242
2.9	Definition (FG: Binary heap well formedness)	242

2.10	Definition (FG: Label substitution)	242
2.11	Definition (FG: Value substitution to value pairs)	242
2.12	Definition (FG: Value substitution to values)	242
2.13	Definition (FG: Unary interpretation of $\Gamma$ )	242
2.14	Definition (FG: Binary interpretation of $\Gamma$ )	242
2.15	Lemma (FG: Binary value relation subsumes unary value relation)	242
2.16	Lemma (FG: Monotonicity Unary)	246
2.17	Lemma (FG: Monotonicity binary)	248
2.18	Lemma (FG: Unary monotonicity for $\Gamma$ )	251
2.19	Lemma (FG: Binary monotonicity for $\Gamma$ )	252
2.20	Lemma (FG: Unary monotonicity for $H$ )	252
2.21	Lemma (FG: Binary monotonicity for heaps)	253
2.22	Theorem (FG: Fundamental theorem unary)	253
2.23	Lemma (FG: Expression subtyping with closed labels and types)	271
2.24	Lemma (FG: Subtyping unary)	272
2.25	Lemma (FG: Binary interpretation of $\Gamma$ implies Unary interpretation of $\Gamma$ )	277
2.26	Theorem (FG: Fundamental theorem binary)	278
2.27	Lemma (FG: Binary heap well formedness implies unary heap well formedness)	320
2.28	Lemma (FG: Subtyping binary)	321
2.29	Theorem (FG: NI)	327
2.30	Definition (CG: $\theta_2$ extends $\theta_1$ )	331
2.31	Definition (CG: $W_2$ extends $W_1$ )	331
2.32	Definition (CG: Value Equivalence)	331
2.33	Definition (CG: Binary value relation)	332
2.34	Definition (CG: Binary expression relation)	332
2.35	Definition (CG: Unary value relation)	332
2.36	Definition (CG: Unary expression relation)	333
2.37	Definition (CG: Unary heap well formedness)	333
2.38	Definition (CG: Binary heap well formedness)	333
2.39	Definition (CG: Label substitution)	333
2.40	Definition (CG: Value substitution to value pairs)	333
2.41	Definition (CG: Value substitution to values)	333
2.42	Definition (CG: Unary interpretation of $\Gamma$ )	333
2.43	Definition (CG: Binary interpretation of $\Gamma$ )	333
2.44	Lemma (CG: Binary value relation subsumes unary value relation)	333
2.45	Lemma (CG: Monotonicity Unary)	338
2.46	Lemma (CG: Monotonicity binary)	340
2.47	Lemma (CG: Unary monotonicity for $\Gamma$ )	345
2.48	Lemma (CG: Binary monotonicity for $\Gamma$ )	345
2.49	Lemma (CG: Unary monotonicity for $H$ )	345
2.50	Lemma (CG: Binary monotonicity for heaps)	346
2.51	Theorem (CG: Fundamental theorem unary)	347
2.52	Lemma (CG: Subtyping unary)	365
2.53	Lemma (CG: Binary interpretation of $\Gamma$ implies Unary interpretation of $\Gamma$ )	370
2.54	Theorem (CG: Fundamental theorem binary)	371
2.55	Lemma (CG: Equivalence of values)	404
2.56	Lemma (CG: Subtyping binary)	405
2.57	Theorem (CG: NI)	412
2.58	Assumption	415

2.59	Theorem (CG $\rightsquigarrow$ FG: Type preservation)	415
2.60	Lemma (CG $\rightsquigarrow$ FG: Subtyping)	419
2.61	Lemma (CG $\rightsquigarrow$ FG: Preservation of well-formedness)	421
2.62	Lemma (CG $\rightsquigarrow$ FG: Free variable lemma)	423
2.63	Definition (CG $\rightsquigarrow$ FG: ${}^s\theta_2$ extends ${}^s\theta_1$ )	424
2.64	Definition (CG $\rightsquigarrow$ FG: $\hat{\beta}_2$ extends $\hat{\beta}_1$ )	424
2.65	Definition (CG $\rightsquigarrow$ FG: Unary value relation)	424
2.66	Definition (CG $\rightsquigarrow$ FG: Unary expression relation)	425
2.67	Definition (CG $\rightsquigarrow$ FG: Unary heap well formedness)	425
2.68	Definition (CG $\rightsquigarrow$ FG: Label substitution)	425
2.69	Definition (CG $\rightsquigarrow$ FG: Value substitution to values)	425
2.70	Definition (CG $\rightsquigarrow$ FG: Unary interpretation of $\Gamma$ )	425
2.71	Lemma (CG $\rightsquigarrow$ FG: Monotonicity)	425
2.72	Lemma (CG $\rightsquigarrow$ FG: Unary monotonicity for $\Gamma$ )	430
2.73	Lemma (CG $\rightsquigarrow$ FG: Unary monotonicity for $H$ )	430
2.74	Theorem (CG $\rightsquigarrow$ FG: Fundamental theorem)	431
2.75	Lemma (CG $\rightsquigarrow$ FG: Subtyping)	459
2.76	Theorem (CG $\rightsquigarrow$ FG: Deriving CG NI via compilation)	464
2.77	Lemma (FG $^-$ : Reflexivity of subtyping)	469
2.78	Definition (FG $\rightsquigarrow$ FG $^-$ : Type translation)	471
2.79	Theorem (FG $\rightsquigarrow$ FG $^-$ : Type preservation)	472
2.80	Lemma (FG $\rightsquigarrow$ FG $^-$ : Subtyping)	479
2.81	Lemma (FG $\rightsquigarrow$ FG $^-$ : Subtyping with label)	484
2.82	Lemma (FG $\rightsquigarrow$ FG $^-$ : Subtyping for $\tau \searrow \ell$ )	484
2.83	Lemma (FG $\rightsquigarrow$ FG $^-$ : Lemma for protection relation)	484
2.84	Lemma (FG $\rightsquigarrow$ FG $^-$ : Substitution lemma)	484
2.85	Lemma (FG $\rightsquigarrow$ FG $^-$ : Preservation of protection relation)	486
2.86	Definition (FG $\rightsquigarrow$ CG: Type translation)	487
2.87	Lemma (Coercion lemma - typing)	488
2.88	Theorem (FG $\rightsquigarrow$ CG: Type preservation)	489
2.89	Lemma (FG $\rightsquigarrow$ CG: Subtyping preservation)	505
2.90	Lemma (FG $\rightsquigarrow$ CG: Preservation of well-formedness)	509
2.91	Lemma (FG $\rightsquigarrow$ CG: Free variable lemma)	511
2.92	Definition (FG $\rightsquigarrow$ CG: ${}^s\theta_2$ extends ${}^s\theta_1$ )	513
2.93	Definition (FG $\rightsquigarrow$ CG: $\hat{\beta}_2$ extends $\hat{\beta}_1$ )	513
2.94	Definition (FG $\rightsquigarrow$ CG: Unary value relation)	513
2.95	Definition (FG $\rightsquigarrow$ CG: Unary expression relation)	513
2.96	Definition (FG $\rightsquigarrow$ CG: Unary heap well formedness)	513
2.97	Definition (FG $\rightsquigarrow$ CG: Label substitution)	513
2.98	Definition (FG $\rightsquigarrow$ CG: Value substitution to values)	513
2.99	Definition (FG $\rightsquigarrow$ CG: Unary interpretation of $\Gamma$ )	513
2.100	Lemma (FG $\rightsquigarrow$ CG: Monotonicity)	514
2.101	Lemma (FG $\rightsquigarrow$ CG: Unary monotonicity for $\Gamma$ )	517
2.102	Lemma (FG $\rightsquigarrow$ CG: Unary monotonicity for $H$ )	518
2.103	Lemma (Coercion lemma)	518
2.104	Theorem (FG $\rightsquigarrow$ CG: Fundamental theorem)	519
2.105	Lemma (FG $\rightsquigarrow$ CG: Semantic Subtyping lemma)	543
2.106	Theorem (FG $\rightsquigarrow$ CG: Deriving FG NI via compilation)	549

# 1 Part I: Details of what's in the paper

## 1.1 Fine-grained IFC enforcement (FG)

### 1.1.1 FG type system

**Syntax, types, constraints:**

Expressions	$e$	::=	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, x.e) \mid \text{new } e \mid !e \mid e := e$
Labels	$\ell, pc$	::=	$\perp \mid \top \mid l \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
(Labeled) Types	$\tau$	::=	$A^\ell$
Unlabeled types	$A$	::=	$\mathbf{b} \mid \text{unit} \mid \tau \xrightarrow{\ell_e} \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \tau$

**Type system:**  $\boxed{\Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var} \qquad \frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp} \text{FG-lam} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG-app} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp} \text{FG-prod} \\
\\
\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG-fst} \qquad \frac{\Gamma \vdash_{pc} e : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp} \text{FG-inl} \\
\\
\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau} \text{FG-case} \\
\\
\frac{\Gamma \vdash_{pc'} e : \tau' \quad \mathcal{L} \vdash pc \sqsubseteq pc' \quad \mathcal{L} \vdash \tau' <: \tau}{\Gamma \vdash_{pc} e : \tau} \text{FG-sub} \qquad \frac{\Gamma \vdash_{pc} e : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp} \text{FG-ref} \\
\\
\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau'} \text{FG-deref} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit}} \text{FG-assign} \qquad \frac{}{\Gamma \vdash_{pc} () : \text{unit}^\perp} \text{FG-unitI}
\end{array}$$

Figure 1: Type system for FG

**Lemma 1.1** (Reflexivity of subtyping). *The following hold:*

1. For all  $\tau$ :  $\mathcal{L} \vdash \tau <: \tau$
2. For all  $A$ :  $\mathcal{L} \vdash A <: A$



$$\begin{array}{c}
\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^\ell <: A^{\ell'}} \text{FGsub-label} \qquad \frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base} \\
\\
\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref} \qquad \frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod} \\
\\
\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum} \\
\\
\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow} \qquad \frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}
\end{array}$$

Figure 2: FG subtyping.

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$ .

Proof of statement (1)

Let  $\tau = A^\ell$ . Then, we have:

$$\frac{\frac{}{\mathcal{L} \vdash A <: A} \text{IH(2)} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on  $A$ .

1.  $A = \mathbf{b}$ :

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2.  $A = \text{ref } \tau$ :

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}$$

4.  $A = \tau_1 + \tau_2$ :

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\mathcal{L} \vdash \ell_e \sqsubseteq \ell_e}}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6.  $A = \text{unit}$ :

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}}$$

□

### 1.1.2 FG semantics

Judgement:  $(H, e) \Downarrow_i (H', v)$

The semantics are described in Figure 3

$$\begin{array}{c} \frac{(H, e_1) \Downarrow_i (H', \lambda x. e_i) \quad (H', e_2) \Downarrow_j (H'', v_2) \quad (H'', e_i[v_2/x]) \Downarrow_k (H''', v_3)}{(H, e_1 e_2) \Downarrow_{i+j+k+1} (H''', v_3)} \text{fg-app} \\ \\ \frac{(H, e_1) \Downarrow_i (H', v_1) \quad (H', e_2) \Downarrow_j (H'', v_2)}{(H, (e_1, e_2)) \Downarrow_{i+j+1} (H'', (v_1, v_2))} \text{fg-prod} \quad \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{fst}(e)) \Downarrow_{i+1} (H', v_1)} \text{fg-fst} \\ \\ \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{snd}(e)) \Downarrow_{i+1} (H', v_2)} \text{fg-snd} \quad \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inl}(e)) \Downarrow_{i+1} (H', \text{inl}(v))} \text{fg-inl} \\ \\ \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inr}(e)) \Downarrow_{i+1} (H', \text{inr}(v))} \text{fg-inr} \quad \frac{(H, e) \Downarrow_i (H', \text{inl } v) \quad (H', e_1[v/x]) \Downarrow_j (H'', v_1)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_1)} \text{fg-case1} \\ \\ \frac{(H, e) \Downarrow_i (H', \text{inr } v) \quad (H', e_2[v/x]) \Downarrow_j (H'', v_2)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_2)} \text{fg-case2} \\ \\ \frac{(H, e) \Downarrow_i (H', v) \quad a \notin \text{dom}(H)}{(H, \text{new}(e)) \Downarrow_{i+1} (H'[a \mapsto v], a)} \text{fg-ref} \quad \frac{(H, e) \Downarrow_i (H', a)}{(H, !e) \Downarrow_{i+1} (H', H(a))} \text{fg-deref} \\ \\ \frac{(H, e_1) \Downarrow_i (H', a) \quad (H', e_2) \Downarrow_j (H'', v)}{(H, e_1 := e_2) \Downarrow_{i+j+1} (H''[a \mapsto v], ())} \text{fg-assign} \quad \frac{e \in \{x, \lambda y. -\}}{(H, e) \Downarrow_0 (H, e)} \text{fg-val} \end{array}$$

Figure 3: FG semantics

### 1.1.3 Logical relation for FG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \leftrightarrow \text{Loc}))$

**Definition 1.2** ( $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 1.3** ( $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

**Definition 1.4** (Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \mathbf{inl} \ v, \mathbf{inl} \ v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \mathbf{inr} \ v, \mathbf{inr} \ v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})\} \\
[\mathbf{ref} \ \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W.\hat{\beta} \wedge W.\theta_1(a_1) = W.\theta_2(a_2) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V^A \triangleq \begin{cases} \{(W, n, v_1, v_2) \mid (W, n, v_1, v_2) \in [\mathbf{A}]_V^A\} & \ell' \sqsubseteq \mathbf{A} \\ \{(W, n, v_1, v_2) \mid \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbf{A}]_V\} & \ell' \not\sqsubseteq \mathbf{A} \end{cases}$$

**Definition 1.5** (Binary expression relation).

$$\begin{aligned}
[\tau]_E^A &\triangleq \{(W, n, e_1, e_2) \mid \\
&\quad \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge \\
&\quad (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\
&\quad \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_V^A\}
\end{aligned}$$

**Definition 1.6** (Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V\} \\
[\tau_1 + \tau_2]_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in [\tau_1]_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in [\tau_2]_V\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\
&\quad (\theta', j, e[v/x]) \in [\tau_2]_E^{\ell_e}\} \\
[\mathbf{ref} \ \tau]_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V \triangleq [\mathbf{A}]_V$$

**Definition 1.7** (Unary expression relation).

$$\begin{aligned}
[\tau]_E^{pc} &\triangleq \{(\theta, n, e) \mid \forall H. (n, H) \triangleright \theta \wedge \forall j < n. (H, e) \Downarrow_j (H', v') \implies \\
&\quad \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)\}
\end{aligned}$$

**Definition 1.8** (Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n-1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

**Definition 1.9** (Binary heap well formedness).

$$\begin{aligned} (n, H_1, H_2) \triangleright^A W \triangleq & \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ & (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ & \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ & (W, n-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ & \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

**Definition 1.10** (Binary substitution).  $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

**Definition 1.11** (Unary substitution).  $\delta : \text{Var} \mapsto \text{Val}$

**Definition 1.12** (Unary interpretation of  $\Gamma$ ).

$$\lfloor \Gamma \rfloor_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \lfloor \Gamma(x) \rfloor_V\}$$

**Definition 1.13** (Binary interpretation of  $\Gamma$ ).

$$\lfloor \Gamma \rfloor_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A\}$$

### 1.1.4 Soundness proof for FG

**Lemma 1.14** (Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n.$

*The following holds:*

1.  $\forall \mathcal{A}.$

$$(W, n, v_1, v_2) \in \lceil \mathcal{A} \rceil_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \lfloor \mathcal{A} \rfloor_V$$

2.  $\forall \tau.$

$$(W, n, v_1, v_2) \in \lceil \tau \rceil_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \lfloor \tau \rfloor_V$$

*Proof.* Proof by simultaneous induction on  $\mathcal{A}$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $\mathcal{A}$  in the last step:

1. Case **b**, unit:

From Definition 1.6

2. Case  $\tau_1 \times \tau_2$ :

$$\underline{\text{Given:}} (W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in \lceil \tau_1 \times \tau_2 \rceil_V^A$$

To prove:

$$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V \quad (\text{P01})$$

and

$$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V \quad (\text{P02})$$

From Definition 1.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A \quad (\text{P1})$$

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in \lfloor \tau_1 \rfloor_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in \lfloor \tau_2 \rfloor_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in \lfloor \tau_2 \rfloor_V$

From (P01) we know that given some  $m$  we need to prove

$$(W.\theta_1, m, (v_{i1}, v_{i2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$$

Similarly from (P02) we know that given some  $m$  we need to prove

$$(W.\theta_2, m, (v_{j1}, v_{j2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V \text{ and } (W.\theta_1, m, v_{i2}) \in \lfloor \tau_2 \rfloor_V$$

Then from Definition 1.6, we get

$$(W.\theta_1, m, (v_{i1}, v_{i2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V \text{ and } (W.\theta_2, m, v_{j2}) \in \lfloor \tau_2 \rfloor_V$$

Then from Definition 1.6, we get

$$(W.\theta_2, m, (v_{j1}, v_{j2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$$

### 3. Case $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V \quad (\text{S01})$$

and

$$\forall m. (W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V \quad (\text{S02})$$

From Definition 1.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

IH1:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in \lfloor \tau_1 \rfloor_V$  and

IH2:  $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some  $m$  and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

Also from (S02) we know that given some  $m$  and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

We instantiate IH1 with  $m$  from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with  $m$  from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{j2})$

Symmetric case as (a)

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$$

This means from Definition 1.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A &\implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V &\implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V &\implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e}) \quad (\text{L0}) \end{aligned}$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$ :

This means from Definition 1.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

This further means that we have some  $\theta', j$  and  $v$  s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

(b)  $\forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$ :

Similar reasoning with  $e_2$

5. Case ref  $\tau$ :

From Definition 1.4 and 1.6

Proof of statement (2)

Let  $\tau = A^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement(1))

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 1.4

□

**Lemma 1.15** (Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m'.$

1.  $\forall A. (\theta, m, v) \in [A]_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in [A]_V$
2.  $\forall \tau. (\theta, m, v) \in [\tau]_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in [\tau]_V$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $A$  in the last step:

1. case **b, unit**:

Directly from Definition 1.6

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

To prove:  $(\theta', m', (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V$

IH1 :  $(\theta', m', v_1) \in [\tau_1]_V$

IH2 :  $(\theta', m', v_2) \in [\tau_2]_V$

We get the desired from IH1, IH2 and Definition 1.6

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

- (a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in [\tau_1 + \tau_2]_V$

To prove:  $(\theta', m', \text{inl } v_1) \in [\tau_1 + \tau_2]_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in [\tau_1]_V$

IH :  $(\theta', m', v_1) \in [\tau_1]_V$

Therefore from IH and Definition 1.6 we get the desired

- (b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(\theta, m, (\lambda x. e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

To prove:  $(\theta', m', (\lambda x. e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

This means from Definition 1.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e} \quad (1)$$

Similarly from Definition 1.6 we know that we are required to prove

$$\forall \theta''' . \theta' \sqsubseteq \theta''' \wedge \forall k < m' . \forall v_1 . (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating Equation 57 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

5. case ref  $\tau$ :

From Definition 1.6 and Definition 1.2

Proof of statement (2)

Let  $\tau = A^\ell$

Since  $[A^\ell]_V = [A]_V$ , therefore from IH (statement 1) □

**Lemma 1.16** (Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, A, n, n'$ .

1.  $\forall A$ .  $(W, n, v_1, v_2) \in [A]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [A]_V^A$
2.  $\forall \tau$ .  $(W, n, v_1, v_2) \in [\tau]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_V^A$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We analyze the different cases of  $A$  in the last step:

1. Case **b**, unit:

From Definition 1.4

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 :  $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 :  $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 1.4 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:



(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 1.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

This means from Definition 1.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A2)

Similarly from Definition 1.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

5. Case ref  $\tau$ :

From Definition 1.4 and Definition 1.3

Proof of statement (2)

Let  $\tau = \mathbf{A}^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement 1)

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

From Lemma 1.15 and Definition 1.4

□

**Lemma 1.17** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .  
 $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

*Proof.* Given:  $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove:  $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 1.12 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$$

And again from Definition 1.12 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

•  $\text{dom}(\Gamma) \subseteq \text{dom}(\delta)$ :

Given

•  $\forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$ :

Since we know that  $\forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$  (given)

Therefore from Lemma 1.15 we get

$$\forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

□

**Lemma 1.18** (Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .

$$(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in [\Gamma]_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(W', n', \gamma) \in [\Gamma]_V$

From Definition 1.13 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And again from Definition 1.12 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

- $dom(\Gamma) \subseteq dom(\gamma)$ :

Given

- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$ :

Since we know that  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$  (given)

Therefore from Lemma 1.16 we get

$$\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

□

**Lemma 1.19** (Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .

$$(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 1.8 it is given that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$$

And again from Definition 1.12 we are required to prove that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

- $dom(\theta) \subseteq dom(H)$ :

Given

- $\forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$ :

Since we know that  $\forall a \in dom(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$  (given)

Therefore from Lemma 1.15 we get

$$\forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

□

**Lemma 1.20** (Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 1.9 it is given that

$$dom(W.\theta_1) \subseteq dom(H_1) \wedge dom(W.\theta_2) \subseteq dom(H_2) \wedge$$

$$(W.\hat{\beta}) \subseteq (dom(W.\theta_1) \times dom(W.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge$$

$$(W, n - 1, H_1(a_1), H_2(a_2)) \in [W.\theta_1(a_1)]_V^A \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in dom(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in [W.\theta_i(a_i)]_V$$

And again from Definition 1.9 we are required to prove:

- $dom(W.\theta_1) \subseteq dom(H_1) \wedge dom(W.\theta_2) \subseteq dom(H_2)$ :

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :  
Given
- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2) \text{ and } (W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A)$ :  
 $\forall (a_1, a_2) \in (W.\hat{\beta})$ .
  - $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given
  - $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
Given and from Lemma 1.16
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :  
Given

□

**Theorem 1.21** (Fundamental theorem unary).  $\forall \Gamma, pc, \theta, e, \tau, \delta, n$ .

$$\begin{aligned} & \Gamma \vdash_{pc} e : \tau \wedge \\ & (\theta, n, \delta) \in \lfloor \Gamma \rfloor_V \implies \\ & (\theta, n, e \delta) \in \lfloor \tau \rfloor_E^{pc} \end{aligned}$$

*Proof.* Proof by induction on  $FG$  typing derivation

1. FG-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove:  $(\theta, n, x \delta) \in \lfloor \tau \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n. (H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, x \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-V0}) \end{aligned}$$

In order to prove FU-V0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $v' = \delta(x)$  and  $j = 1$

We need to prove the following:

- (a)  $\theta \sqsubseteq \theta \wedge (n - 1, H) \triangleright \theta \wedge (\theta, n - 1, v') \in \lfloor \tau \rfloor_V$ :
  - $\theta \sqsubseteq \theta$ : From Definition 1.2
  - $(n - 1, H) \triangleright \theta$ : From Lemma 1.19

- $(\theta, n-1, v') \in \lfloor \tau \rfloor_V$ :  
 Since we are given that  $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$  and  $v' = \delta(x)$   
 Therefore  $(\theta, n, v') \in \lfloor \Gamma(x) \rfloor_V$ , where  $\Gamma(x) = \tau$   
 And finally from Lemma 1.15 we get  $(\theta, n-1, v') \in \lfloor \tau \rfloor_V$
- (b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :  
 Since  $H' = H$ , so we are done
- (c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :  
 Since  $\theta' = \theta$ , so we are done

## 2. FG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(\theta, \lambda x. e_i \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall j < n. (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H') \triangleright \theta' \wedge (\theta', n-j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H') \triangleright \theta' \wedge (\theta', n-j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-L0}) \end{aligned}$$

IH1:

$$\forall \theta_i, v_x, n. (\theta_i, n, e_i \delta \cup \{x \mapsto v_x\}) \in \lfloor \tau_2 \rfloor_E^{\ell_e}, \text{ s.t } (\theta_i, n, v_x) \in \lfloor \tau_1 \rfloor_V$$

In order to prove FU-L0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $j = 0$  and  $v' = \lambda x. e_i \delta$

$$(a) \theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_V:$$

- $\theta \sqsubseteq \theta$ : From Definition 1.2
- $(n, H) \triangleright \theta$ : Given
- $(\theta, n, (\lambda x. e_i) \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_V$ :

From Definition 1.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_i[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

This means given some  $\theta''$ ,  $j$  and  $v$  such that  $\theta \sqsubseteq \theta''$ ,  $j < n$  and  $(\theta'', j, v) \in \lfloor \tau_1 \rfloor_V$

It suffices to prove that  $(\theta'', j, e_i[v/x] \delta) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$

Since  $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$  and  $j < n$  therefore from Lemma 1.17 we have

$$(\theta, j, \delta) \in \lfloor \Gamma \rfloor_V$$

So we can apply IH1 instantiated with  $\theta''$ ,  $v$  and  $j$  we get

$$(\theta'', j, e_i[v/x] \delta) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

(b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $H' = H$  so we are done

(c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :

Since  $\theta' = \theta$  so we are done

### 3. FG-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  s.t.  $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-P0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall n_1, H_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $n, H$  and since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-P1}) \end{aligned}$$

From evaluation rule we know that  $v'_1 = \lambda x. e_i$ . Since from FU-P1 we know that

$$(\theta'_1, n - i, \lambda x. e_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V$$

This means from Definition 1.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < (n - i). \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2]_E^{\ell_e} \quad (2)$$

IH2:

$$\begin{aligned} & \forall n_2, \forall H_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall k < n_2. (H_2, (e_2) \delta) \Downarrow_k (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - k, v'_2) \in [(\tau_1)]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with  $n-i$ ,  $H'_1$  and since we know that  $(n-i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1 \ e_2) \ \delta) \Downarrow_{n'}$   $(H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n-i-k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n-i-k, v'_2) \in [(\tau_1)]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-P2})
\end{aligned}$$

Instantiating  $\theta''$ ,  $j$  and  $v$  in Equation 2 with  $\theta'_2$ ,  $n-i-k$  and  $v'_2$  from FU-P2 respectively, we get

$$(\theta'_2, n-i-k, e_i[v'_2/x]) \in [\tau_2]_E^{\ell_e}$$

This means from Definition 1.7 we have

$$\begin{aligned}
& \forall H_3. (n-i-k, H_3) \triangleright \theta'_2 \wedge \forall l < (n-i-k). (H_3, e_i[v'_2/x]) \Downarrow_l (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e)
\end{aligned}$$

Instantiating  $H_3$  with  $H'_2$  from FU-P2 and since we know that  $((n-i-k), H'_2) \triangleright \theta'_2$  and since the reduction happens therefore we have

$$\begin{aligned}
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e) \quad (\text{FU-P3})
\end{aligned}$$

In order to prove FU-P0 we choose  $\theta'$  as  $\theta'_3$  from FU-P3. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i+k+l$ . Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V:$$

- $\theta \sqsubseteq \theta'_3$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-P1,  $\theta'_1 \sqsubseteq \theta'_2$  from FU-P2 and  $\theta'_2 \sqsubseteq \theta'_3$  from FU-P3 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_3$
- $((n-i-k-l), H'_3) \triangleright \theta'_3$ :  
From FU-P3 we get  $((n-i-k-l), H'_3) \triangleright \theta'_3$
- $(\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V$ :  
From FU-P3 we get  $(\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

Since  $pc \sqsubseteq \ell_e$  therefore we get the desired from FU-P1, FU-P2 and FU-P3

$$(c) \ (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta). \theta'_3(a) \searrow pc)$$

Since  $pc \sqsubseteq \ell_e$  therefore we get the desired from FU-P1, FU-P2 and FU-P3

#### 4. FG-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(\theta, n, (e_1, e_2) \ \delta) \in [(\tau_1 \times \tau_2)^\perp]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall n' < n.(H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 \times \tau_2)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  s.t  $H \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 \times \tau_2)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-PA0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1.(n_1, H_1) \triangleright \theta \wedge \forall i < n_1.(H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau_1]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

We instantiate IH1 with  $H$  and  $n$ . And since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau_1]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-PA1}) \end{aligned}$$

IH2:

$$\begin{aligned} & \forall H_2, n_2.(n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2.(H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau_2)]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \end{aligned}$$

We instantiate IH2 with  $H'_1$  and  $n - i$ . And since we know that  $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau_2)]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-PA2}) \end{aligned}$$

In order to prove FU-PA0 we choose  $\theta'$  as  $\theta'_2$  from FU-PA2. Also we know from the evaluation rule, that let  $v' = (v'_1, v'_2)$ ,  $H' = H'_2$  and  $n' = i + j + 1$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v') \in [(\tau_1 \times \tau_2)^\perp]_V:$$

- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-PA1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-PA2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$
- $(n - i - j - 1, H') \triangleright \theta'_2$ :  
From FU-PA2 we get  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.19 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$



- $(\theta'_2, n - i - j, v') \in [(\tau_1 \times \tau_2)^\perp]_V$ :

From Definition 1.6 it suffices to show

- i.  $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1)]_V$ :

Since from FU-PA1 we know that  $(\theta'_1, n - i, v'_1) \in [(\tau_1)]_V$  and since  $\theta'_1 \sqsubseteq \theta'_2$

(from FU-PA2) therefore from Lemma 1.15 we get

$$(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1)]_V$$

- ii.  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2)]_V$ :

From FU-PA2 we know that  $(\theta'_2, n - i - j, v'_2) \in [(\tau_2)]_V$  therefore from

Lemma 1.15 we get  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2)]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

From FU-PA1 and FU-PA2

$$(c) \ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$$

From FU-PA1 and FU-PA2

5. FG-fst:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove:  $(\theta, n, \text{fst}(e_i) \delta) \in [\tau_1]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v') \implies$$

$$\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$$

This means that given some heap  $H$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-F0})$$

IH1:

$$\forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies$$

$$\exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell]_V \wedge$$

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow (H', v')$  therefore we have

$$\exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell]_V \wedge$$

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-F1})$$

From evaluation rule we know that  $v'_1 = (v''_1, v''_2)$

In order to prove FU-F0 we choose  $\theta'$  as  $\theta'_1$  from FU-P1. Also we know that  $H' = H'_1$  and  $v' = v''_1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v'_1) \in \lfloor \tau_1 \rfloor_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-F1
  - $(n - i - 1, H'_1) \triangleright \theta'_1$ :  
From FU-F1 we know  $(n - i, H'_1) \triangleright \theta'_1$  therefore from Lemma 1.19 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$
  - $(\theta'_1, n - i, v''_1) \in \lfloor \tau_1 \rfloor_V$ :  
Since from FU-F1 we know that  $(\theta'_1, n - i, (v''_1, v''_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$   
Therefore from Definition 1.6 we know that  $(\theta'_1, n - i, v''_1) \in \lfloor \tau_1 \rfloor_V$   
From Lemma 1.15 we get  $(\theta'_1, n - i - 1, v''_1) \in \lfloor \tau_1 \rfloor_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
From FU-F1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$   
From FU-F1

## 6. FG-snd:

Symmetric case to FG-fst

## 7. FG-inl:

$$\frac{\Gamma \vdash_{pc} e_i : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(\theta, n, \text{inl}(e_i) \delta) \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-LE0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor \tau_1 \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor \tau_1 \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-LE1}) \end{aligned}$$

In order to prove FU-LE0 we choose  $\theta'$  as  $\theta'_1$  from FU-LE1. Also we know from the evaluation rule, that let  $v' = \text{inl}(v'_1)$ ,  $H' = H'_1$  and  $n' = i + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-LE1
  - $(n - i - 1, H') \triangleright \theta'_1$ :  
From FU-LE1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 1.19 we get  $(n - i - 1, H') \triangleright \theta'_1$
  - $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$ :  
Since  $v' = \text{inl}(v'_1)$  and from FU-LE1 we know that  $(\theta'_1, n - i, v'_1) \in [\tau_1]_V$   
Therefore from Definition 1.6 we get  $(\theta'_1, n - i, v') \in [(\tau_1 + \tau_2)]_V$   
From Lemma 1.15 we get  $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
From FU-LE1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$   
From FU-LE1

8. FG-inr:

Symmetric case to FG-inl

9. FG-case:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau}$$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t.  $(n, H) \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-C0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_c) \delta) \Downarrow_i (H'_1, v'_c) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_c) \in [(\tau_1 + \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_c) \in \llbracket (\tau_1 + \tau_2)^\ell \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-C1}) \end{aligned}$$

2 cases arise:

(a)  $v'_c = \text{inl}(v_{ci})$ :

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_1) \delta \cup \{x \mapsto v_{ci}\}) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Instantiating IH2 with  $H'_1$  and  $n - i$  since we know that  $H'_1 \triangleright \theta'_1 \wedge (H', (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \quad (\text{FU-C2}) \end{aligned}$$

In order to prove FU-C0 we choose  $\theta'$  as  $\theta'_2$  from FU-C2. Also we know that  $H' = H'_2$ ,  $v' = v'_2$  and  $n' = i + j + 1$ . Now we are required to show

i.  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$ :

•  $\theta \sqsubseteq \theta'_2$ :

Since  $\theta \sqsubseteq \theta'_1$  from FU-C1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-C2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$

•  $(n - i - j - 1, H'_2) \triangleright \theta'_2$ :

From FU-C2 we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.19 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$

•  $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$ :

From FU-C2 we know that  $(\theta'_2, n - i - j, v'_2) \in \llbracket \tau \rrbracket_V$  therefore from Lemma 1.15 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$

ii.  $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since from FU-C2 we know that

$$(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell')$$

therefore we also have

$$(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc) \sqsubseteq \ell')$$

and from FU-C1 we know that

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge (pc) \sqsubseteq \ell')$$

Combining the two we get

$$(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

iii.  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$ :

Since from FU-C2 we know that

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell))$$

therefore we also have

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc))$$

and from FU-C1 we know that  
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow (pc \sqcup \ell))$

Combining the two we get  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$

- (b)  $v'_c = \text{inr}(v_{ci})$ :  
 Symmetric case as  $v'_c = \text{inl}(v_{ci})$

10. FG-ref:

$$\frac{\Gamma \vdash_{pc} e_i : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(\theta, n, \text{new } (e_i) \delta) \in [(\text{ref } \tau)^\perp]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\text{ref } \tau)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\text{ref } \tau)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-R0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-R1}) \end{aligned}$$

From the evaluation rule we know that  $H' = H'_1[a \mapsto v'_1]$  where  $a \notin \text{dom}(H'_1)$ ,  $v' = a$  and  $n' = i + 1$ . In order to prove FU-R0 we choose  $\theta'$  as  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau\})$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_2 \wedge (n - i - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - 1, v') \in [(\text{ref } \tau)^\perp]_V$ :

- $\theta \sqsubseteq \theta'_2$ :

From FU-R1 we know that  $\theta \sqsubseteq \theta'_1$  therefore from Definition 1.2  $\theta \sqsubseteq \theta'_2$

- $(n - i - 1, H') \triangleright \theta'_2$ :  
 From FU-R1 we know that  $(n - i, H'_1) \triangleright \theta'_1$ . Therefore from Lemma 1.19 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$ .  
 We also know that  $(\theta'_1, n - i, v'_1) \in \lfloor \tau \rfloor_V$  (from FU-R1) therefore from Lemma 1.15 we get  $(\theta'_1, n - i - 1, v'_1) \in \lfloor \tau \rfloor_V$   
 Since  $H' = H'_1[a \mapsto v'_1]$  and  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau\})$  therefore from Definition 1.8 we get  $(n - i - 1, H') \triangleright \theta'_2$
  - $(\theta'_2, n - i - 1, a) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V$ :  
 Since  $\theta'_2(a) = \tau$  therefore from Definition 1.6 we get  $(\theta'_2, n - i - 1, a) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
 From FU-R1
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$ :  
 We get this from FU-R1 and  $\tau \searrow pc$  (given)

11. FG-deref:

$$\frac{\Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(\theta, n, (!e_i) \delta) \in \lfloor \tau' \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (!e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (!e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-D0}) \end{aligned}$$

**IH1:**

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (!e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-D1}) \end{aligned}$$

In order to prove FU-D0 we choose  $\theta'$  as  $\theta'_1$  from FU-D1. Also we know from the evaluation rule, that  $H' = H'_1$ ,  $v' = H'_1(a)$ ,  $v'_1 = a$  and  $n' = i + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in \lfloor \tau \rfloor_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-D1
  - $(n - i - 1, H') \triangleright \theta'_1$ :  
From FU-D1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 1.19 we get  $(n - i - 1, H') \triangleright \theta'_1$
  - $(\theta'_1, n - i - 1, v') \in \lfloor \tau' \rfloor_V$ :  
Since from FU-D1 we know that  $(n - i, H'_1) \triangleright \theta'_1$  therefore from the Definition 1.8 we get  $(\theta'_1, n - i, H'_1(a)) \in \lfloor \tau \rfloor_V$   
From Lemma 1.15 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in \lfloor \tau \rfloor_V$   
Since  $\tau <: \tau'$  therefore from Lemma 1.23 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in \lfloor \tau' \rfloor_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
From FU-D1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$   
From FU-D1

12. FG-assign:

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit}}$$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in \lfloor \text{unit} \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \text{unit} \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \text{unit} \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-A0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-A1}) \end{aligned}$$

IH2:

$$\begin{aligned}
& \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq (n_2 - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with  $H'_1$  and since we know that  $H'_1 \triangleright \theta'_1 \wedge (H', (e_1 := e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq (n - i - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad \text{(FU-A2)}
\end{aligned}$$

In order to prove FU-A0 we choose  $\theta'$  as  $\theta'_2$  from FU-A2. Also we know from the evaluation rule, assign, that let  $v'_1 = a_1$ ,  $H' = H'_2[a_1 \mapsto v'_2]$ ,  $v' = ()$  and  $n' = i + j + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$ :
- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-A1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-A2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$
  - $(n - i - j - 1, H') \triangleright \theta'_2$ :  
From Definition 1.8 it suffices to prove that
    - i.  $\text{dom}(\theta'_2) \subseteq \text{dom}(H')$ : From FU-A2
    - ii.  $\forall a \in \text{dom}(\theta'_2). (\theta'_2, n - i - j - 1, H'(a)) \in \llbracket \theta'_2(a) \rrbracket_V$ :  
 $\forall a \in \text{dom}(\theta'_2)$ .
      - $a = a_1$ :  
From FU-A2 (since we know that  $(\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V$ )  
Therefore from Lemma 1.15 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket (\tau) \rrbracket_V$
      - $a \neq a_1$ :  
From FU-A2 (since we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.19 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$ )
  - $(\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$ :  
From Definition 1.6
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell')$   
 $\forall a \in \text{dom}(H)$ .
- $a = a_1$ :  
Since we know that  $H(a_1) \neq H'(a_1)$  and  $\theta(a_1) = \tau = A^{\ell_i}$  (given)  
It is given that  $\tau \searrow pc$  therefore  $pc \sqsubseteq \ell_i$
  - $a \neq a_1$ :  
From FU-A2
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$   
From FU-A2

□

**Lemma 1.22** (Expression subtyping).  $\forall pc, pc', \tau$ .

$$\mathcal{L} \models pc \sqsubseteq pc' \implies \llbracket \tau \rrbracket_E^{pc'} \subseteq \llbracket \tau \rrbracket_E^{pc}$$



*Proof.* Given:  $\mathcal{L} \models pc \sqsubseteq pc'$

To prove:  $\llbracket(\tau)\rrbracket_E^{pc'} \subseteq \llbracket(\tau)\rrbracket_E^{pc}$

This means we need to prove that

$$\forall(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc'} . (\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc}$$

This means given  $\forall(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc'}$

It suffices to prove that  $(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc}$

From Definition 1.7 for the chosen  $\theta, n, e$  we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc') \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall k < n.(H_1, e) \Downarrow_k (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some  $H_1$  and  $k$  such that  $(n, H_1) \triangleright \theta$ ,  $k < n$  and  $(H_1, e) \Downarrow_k (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned} \quad (\text{B})$$

Instantiate  $H$  with  $H_1$  and  $j$  with  $k$  in (A) to get

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc') \end{aligned} \quad (\text{C})$$

In order to prove (B) we choose  $\theta'_1$  as  $\theta'$  and we need to prove

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \llbracket\tau\rrbracket_V$ :

We get this directly from (C)

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $pc \sqsubseteq pc'$  and we are given

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell')$$

Therefore

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$ :

We are given

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc')$$

and since  $pc \sqsubseteq pc'$  Therefore

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$$

□

**Lemma 1.23** (Subtyping unary). *The following holds:*

1.  $\forall A, A', \mathcal{L}$ .

$$(a) \mathcal{L} \vdash A <: A' \implies \llbracket (A) \rrbracket_V \subseteq \llbracket (A') \rrbracket_V$$

2.  $\forall \tau, \tau', \mathcal{L}$ .

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$$

$$(b) \forall pc. \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_E^{pc} \subseteq \llbracket (\tau') \rrbracket_E^{pc}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

IH1:  $\llbracket (\tau'_1) \rrbracket_V \subseteq \llbracket (\tau_1) \rrbracket_V$  (Statement 2(a))

IH2:  $\forall pc. \llbracket (\tau_2) \rrbracket_E^{pc} \subseteq \llbracket (\tau'_2) \rrbracket_E^{pc}$  (Statement 2(b))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V$

Therefore from Definition 1.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E^{\ell_e} \quad (3)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

Again from Definition 1.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E^{\ell'_e}$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E^{\ell'_e}$

Since  $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in \llbracket \tau_1 \rrbracket_V$

As a result from Equation 3 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E^{\ell_e}$$

From IH2, we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2]_E^{\ell_e}$$

Since  $\mathcal{L} \models \ell'_e \sqsubseteq \ell_e$  therefore from Lemma 1.22 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2]_E^{\ell'_e}$$

## 2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $[\!(\tau_1 \times \tau_2)\!]_V \subseteq [\!(\tau'_1 \times \tau'_2)\!]_V$

IH1:  $[\!(\tau_1)\!]_V \subseteq [\!(\tau'_1)\!]_V$  (Statement 2(a))

IH2:  $[\!(\tau_2)\!]_V \subseteq [\!(\tau'_2)\!]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, (v_1, v_2)) \in [\!(\tau_1 \times \tau_2)\!]_V. (\theta, n, (v_1, v_2)) \in [\!(\tau'_1 \times \tau'_2)\!]_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [\!(\tau_1 \times \tau_2)\!]_V$

Therefore from Definition 1.6 we are given:

$$(\theta, n, v_1) \in [\tau_1]_V \wedge (\theta, n, v_2) \in [\tau_2]_V \tag{4}$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [\!(\tau'_1 \times \tau'_2)\!]_V$

Again from Definition 1.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1]_V \wedge (\theta, n, v_2) \in [\tau'_2]_V$$

Since from Equation 4 we know that  $(\theta, n, v_1) \in [\tau_1]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2]_V$  from Equation 4 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2]_V$

## 3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[\!(\tau_1 + \tau_2)\!]_V \subseteq [\!(\tau'_1 + \tau'_2)\!]_V$

IH1:  $[\!(\tau_1)\!]_V \subseteq [\!(\tau'_1)\!]_V$  (Statement 2(a))

IH2:  $[\!(\tau_2)\!]_V \subseteq [\!(\tau'_2)\!]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, v_s) \in [\!(\tau_1 + \tau_2)\!]_V. (\theta, v_s) \in [\!(\tau'_1 + \tau'_2)\!]_V$

This means that given:  $(\theta, n, v_s) \in [\!(\tau_1 + \tau_2)\!]_V$

And it suffices to prove:  $(\theta, n, v_s) \in [\!(\tau'_1 + \tau'_2)\!]_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_1 \rrbracket_V \quad (5)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_1 \rrbracket_V$$

From Equation 5 and IH1 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_1 \rrbracket_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_2 \rrbracket_V \quad (6)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_2 \rrbracket_V$$

From Equation 6 and IH2 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_2 \rrbracket_V$$

#### 4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\llbracket ((\text{ref } \tau)) \rrbracket_V \subseteq \llbracket ((\text{ref } \tau)) \rrbracket_V$

It suffices to prove:  $\forall (\theta, n, a) \in \llbracket ((\text{ref } \tau)) \rrbracket_V. (\theta, n, a) \in \llbracket ((\text{ref } \tau)) \rrbracket_V$

Trivial

#### 5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b})) \rrbracket_V \subseteq \llbracket ((\mathbf{b})) \rrbracket_V$

Directly from Definition 1.6

#### 6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit})) \rrbracket_V \subseteq \llbracket ((\text{unit})) \rrbracket_V$

Directly from Definition 1.6

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell' \quad \mathcal{L} \vdash \mathbf{A} <: \mathbf{A}'}{\mathcal{L} \vdash \mathbf{A}^\ell <: \mathbf{A}^{\ell'}} \text{FGsub-label}$$

To prove:  $\llbracket (\mathbf{A}^\ell) \rrbracket_V \subseteq \llbracket (\mathbf{A}^{\ell'}) \rrbracket_V$

From Definition 1.6 it suffices to prove:  $\llbracket (\mathbf{A}) \rrbracket_V \subseteq \llbracket (\mathbf{A}') \rrbracket_V$

This we get directly from IH (Statement 1(a))

Proof of statement 2(b)

Given:  $\mathcal{L} \vdash \tau <: \tau'$

To prove:  $\llbracket (\tau) \rrbracket_E^{pc} \subseteq \llbracket (\tau') \rrbracket_E^{pc}$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}. (\theta, n, e) \in \llbracket (\tau') \rrbracket_E^{pc}$$

This means given  $(\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}$

It suffices to prove that  $(\theta, n, e) \in \llbracket (\tau') \rrbracket_E^{pc}$

From Definition 1.7 we know we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall i < n.(H, e) \Downarrow_i (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - i, H') \triangleright \theta' \wedge (\theta', n - i, v') \in \llbracket \tau \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall j < n.(H_1, e) \Downarrow_j (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some  $H_1$  and  $j < n$  s.t  $(n, H_1) \triangleright \theta \wedge (H_1, e) \Downarrow_j (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate  $H$  in (A) with  $H_1$  and  $i$  with  $j$  then we choose  $\theta'_1$  as  $\theta'$

Also we have IH1 as  $\llbracket \tau \rrbracket_V \subseteq \llbracket \tau' \rrbracket_V$  (Statement 2(a))

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau' \rrbracket_V$ :

We are given  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau \rrbracket_V$

From IH1 we know that  $\llbracket \tau \rrbracket_V \subseteq \llbracket \tau' \rrbracket_V$

Therefore,  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau' \rrbracket_V$

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Given

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$ :

Given

□

**Lemma 1.24** (Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$   
 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 1.13 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case  $i = 1$

Given some  $m$  we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_1)$ :

$$\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_1)$$

Therefore,  $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_1))$  (Given)

- $\forall x \in \text{dom}(\Gamma). (W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$ :

$$\text{We are given: } \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

Therefore from Lemma 1.14 we know that

$$\forall m'. (W.\theta_1, m', \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Instantiating  $m'$  with  $m$  we get

$$(W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Case  $i = 2$

Symmetric case as  $i = 1$

□

**Theorem 1.25** (Fundamental theorem binary).  $\forall \Gamma, pc, W, \mathcal{A}, e, \tau, \gamma, n.$

$$\Gamma \vdash_{pc} e : \tau \wedge (W, n, \gamma) \in [\Gamma]_V^A \implies$$

$$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau]_E^A$$

*Proof.* Proof by induction on the typing derivation

1. FG-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

$$\text{To prove: } (W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau]_E^A$$

$$\text{Say } e_1 = x(\gamma \downarrow_1) \text{ and } e_2 = x(\gamma \downarrow_2)$$

From Definition of  $[\tau]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall j < n. (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^A$$

This means given some  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

We are required to prove:  $\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^A$

Here

- $H'_1 = H_1$  and  $H'_2 = H_2$
- $e_1 = v'_1 = \gamma(x) \Downarrow_1$
- $e_2 = v'_2 = \gamma(x) \Downarrow_2$
- $j = 1$

We choose  $W' = W$ .

- $W \sqsubseteq W$ : From Definition 1.3
- $(n - 1, H_1, H_2) \stackrel{A}{\triangleright} W$ :  
Since we know that  $(n, H_1, H_2) \stackrel{A}{\triangleright} W$  therefore from Lemma 1.20 we get  $(n - 1, H_1, H_2) \stackrel{A}{\triangleright} W$
- $(W, n - 1, \gamma(x) \Downarrow_1, \gamma(x) \Downarrow_2) \in [\tau]_{\mathcal{V}}^A$ :  
We are given that  $(W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^A$  therefore from Lemma 1.18 we get  $(W, n - 1, \gamma) \in [\Gamma]_{\mathcal{V}}^A$   
which means from Definition 1.13 we have  $(W, n - 1, \gamma(x) \Downarrow_1, \gamma(x) \Downarrow_2) \in [\tau]_{\mathcal{V}}^A$

## 2. FG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(W, n, \lambda x. e (\gamma \Downarrow_1), \lambda x. e (\gamma \Downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)]_E^A$

Say  $e_1 = \lambda x. e (\gamma \Downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \Downarrow_2)$

From Definition of  $[(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_E^A$  it suffices to prove that

$$\forall H_1, H_2, j < n. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\mathcal{V}}^A$$

This means that given  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

It suffices to prove:

$$\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\mathcal{V}}^A \quad (\text{FB-L0})$$

IH1:

$\forall W, n. (W, n, e (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2]_E^A$

s.t

$(W, n, (v_1, v_2)) \in [\tau_1]_V^A$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \lambda x.e (\gamma \downarrow_1)$ ,  $v'_2 = e_2 = \lambda x.e (\gamma \downarrow_2)$  and  $j = 0$ . In order to prove FB-L0 we choose  $W' = W$  and we need to prove the following:

- $W \sqsubseteq W$ : From Definition 1.3
- $(n, H_1, H_2) \triangleright^A W$ : Given
- $(W, n, \lambda x.e (\gamma \downarrow_1), \lambda x.e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^A$

From Definition 1.4 it suffices to prove that:

$\forall W'' \sqsupseteq W, k < n, v_1, v_2.$

$((W'', k, v_1, v_2) \in [\tau_1]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, k, v_c.$

$((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, k, v_c.$

$((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_2)[v_c/x]) \in [\tau_2]_E^{\ell_e})$

This means that we need to prove the following:

- $\forall W'' \sqsupseteq W, k < n, v_1, v_2. ((W'', k, v_1, v_2) \in [\tau_1]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A):$

This means given  $W'' \sqsupseteq W, k < n, v_1, v_2$  s.t  $((W'', k, v_1, v_2) \in [\tau_1]_V^A$

We need to prove:  $(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A$

We instantiate IH1 with  $W''$  and  $k$

And since  $(W'', k, v_1, v_2) \in [\tau_1]_V^A$  therefore we get

$(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A$

- $\forall \theta_l \sqsupseteq W.\theta_1, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}):$

This means that we are given  $\theta_l, k$  and  $v_c$  s.t

$\theta_l \sqsupseteq W.\theta_1$  and  $(\theta_l, k, v_c) \in [\tau_1]_V$

And we are required to prove:

$(\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}$

It is given to us that

$\forall v_1, v_2. (W, n, \gamma \in [\Gamma]_V^A$

Therefore from Lemma 1.24 we know that

$\forall m. (W.\theta_1, m, (\gamma \downarrow_1) \in [\Gamma]_V$

Therefore, we can apply Theorem 1.21 to obtain

$\forall m. (W.\theta_1, m, \lambda x.e \gamma \downarrow_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V$

From Definition 1.6 it means that we have



$$\forall m. \forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\ (\theta', j, e[v/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$$

We instantiate  $m$  with some  $l > k$ ,  $\theta'$  with  $\theta_l$ ,  $j$  with  $k$  and  $v$  with  $v_c$  to get  $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$

Since we thow that  $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1]_V$  therefore we get  $(\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$

$$- \forall \theta_l \sqsupseteq W.\theta_2, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies \\ (\theta_l, k, e(\gamma \downarrow_2)[v_c/x]) \in [\tau_2]_E^{\ell_e}): \\ \text{Symmetric case as above}$$

### 3. FG-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2)]_E^A$

This means from Definition 1.5 we need to prove:

$$\forall H_1, H_2, n' < n. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A$$

This further means that given  $H_1, H_2, n' < n$  s.t

$$(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A \quad (\text{FB-A0})$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}, i < n. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge (H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$ .  $(H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V^A \quad (7)$$

$$\underline{\text{IH2}}: (W'_1, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [(\tau_1)]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{j1}, H_{j2}, j < (n - i). (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge (H_{j1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \\ \exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Instantiating  $H_{j_1}$  with  $H'_1$  and  $H_{j_2}$  with  $H'_2$  in IH2. Since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Also,  $e_1$  reduces to value  $\gamma \downarrow_1$  in  $i < n'$  steps. Therefore  $\exists j < n' - i < n - i$  s.t  $(H_{i_1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j_1}, v'_{j_1})$ .  $(H_{i_2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j_2}, v'_{j_2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in \lceil (\tau_1) \rceil_V^A \quad (8)$$

We case analyze on  $(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$  from Equation 7

- Case  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$$

Let  $v'_1 = \lambda x.e_{h1}$  and  $v'_2 = \lambda x.e_{h2}$

Again from Definition 1.4 it means that

$$\forall W'_{h1} \sqsupseteq W'_1, j_1 < (n - i), v_1, v_2.$$

$$((W'_{h1}, j_1, v_1, v_2) \in \lceil \tau_1 \rceil_V^A \implies (W'_{h1}, j_1, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in \lceil \tau_2 \rceil_E^A) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_1, m_1, v_c.$$

$$\wedge ((\theta_{l1}, m_1, v_1) \in \lceil \tau_1 \rceil_V \implies (W'_{h1}.\theta_1, e_{h1}[v_1/x]) \in \lceil \tau_2 \rceil_E^{\ell_e}) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_2, m_1, v_c.$$

$$\wedge (\theta_{l1}, m_1, v_2) \in \lceil \tau_1 \rceil_V \implies (W'_{h1}.\theta_2, e_{h2}[v_2/x]) \in \lceil \tau_2 \rceil_E^{\ell_e}$$

We instantiate  $W'_{h1}$  with  $W'_2$  obtained from Equation 8. Similarly we also instantiate  $v_1$  and  $v_2$  with  $v'_{j_1}$  and  $v'_{j_2}$  respectively from Equation 8, and  $j_1$  with  $n - i - j$ . And we get

$$(W'_2, n - i - j, e_{h1}[v'_{j_1}/x], e_{h2}[v'_{j_2}/x]) \in \lceil \tau_2 \rceil_E^A$$

From Definition 1.5 we get

$$\forall H_1, H_2, k_e < (n - i - j).(n - i - j, H_1, H_2) \stackrel{A}{\triangleright} W'_2 \wedge$$

$$(H_1, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1}) \wedge (H_2, e_{h2}[v'_{j_2}/x]) \Downarrow (H'_{f2}, v_{f2}) \implies$$

$$\exists W' \sqsupseteq W'_2.(n - i - j - k_e, H'_{f1}, H'_{f2}) \stackrel{A}{\triangleright} W' \wedge (W', n - i - j - k_e, v_{f1}, v_{f2}) \in \lceil \tau_2 \rceil_V^A$$

Instantiating  $H_1$  with  $H'_{j_1}$  and  $H_2$  with  $H'_{j_2}$  obtained from Equation 8. And since we know that  $e_1 e_2$  reduces with  $\gamma \downarrow_1$  in  $n' < n$  steps. And  $e_2$  reduces to value  $\gamma \downarrow_1$  in  $j < n' - 1 < n - i$  steps. Therefore  $\exists k_e = n' - i - j < n - i - j$  s.t  $(H_1, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1})$ .  $(H_2, e_{h2}[v'_{j_2}/x]) \Downarrow (H'_{f2}, v_{f2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W' \sqsupseteq W'_2.((n - i - j - k_e), H'_{f1}, H'_{f2}) \stackrel{A}{\triangleright} W' \wedge (W', (n - i - j - k_e), v_{f1}, v_{f2}) \in \lceil \tau_2 \rceil_V^A \quad (9)$$

This concludes the proof in this case.

- Case  $\ell \not\sqsubseteq \mathcal{A}$ :

From FB-A0 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A$$

In this case since we know that  $\ell \not\sqsubseteq \mathcal{A}$ . Let  $\tau_2 = A^{\ell_i}$  and since  $\tau_2 \searrow \ell$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V \quad (10)$$

In this case from Definition 1.6 we know that

$$\forall m. (W'_1.\theta_1, m, \lambda x. e_{h1}) \in [(\tau_1 \xrightarrow{\ell_\xi} \tau_2)]_V \quad (11)$$

$$\forall m. (W'_1.\theta_2, m, \lambda x. e_{h2}) \in [(\tau_1 \xrightarrow{\ell_\xi} \tau_2)]_V \quad (12)$$

Applying Definition 1.6 on Equation 11 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2]_E^{\ell_e} \text{ where } \theta = W'_1.\theta_1$$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 1 + t_1). \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2]_E^{\ell_e} \quad (\text{FB-AC1})$$

Since from Equation 8 we have

$$(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Therefore from Lemma 1.14 we get

$$\forall m. (W'_2.\theta_1, m, v'_{j1}) \in [\tau_1]_V$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(W'_2.\theta_1, m_1 + 1 + t_1, v'_{j1}) \in [\tau_1]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_1$ ,  $j_1$  with  $m_1 + t_1$  and  $v$  with  $v'_{j1}$  from Equation 8.

Therefore we get  $(W'_2.\theta_1, m_1 + 1 + t_1, e_{h1}[v'_{j1}/x]) \in [\tau_2]_E^{\ell_e}$

From Definition 1.7, we get

$$\begin{aligned}
& \forall H.(m_1 + 1 + t_1, H) \triangleright W'_2.\theta_1 \wedge \forall k_c < (m_1 + 1 + t_1).(H, e_{h1}[v'_{j1}/x]) \Downarrow_{k_c} (H'_1, v'_1) \implies \\
& \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1 + t_1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1 + t_1 - k_c), v'_1) \in [\tau_2]_V \wedge \\
& (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e))
\end{aligned}$$

Since from Equation 8 we have  $(n - i - j, H'_{j1}, H'_{j1}) \triangleright W'_2$

Therefore from Lemma 1.26 we get  $\forall m.(m, H'_{j1}) \triangleright W'_2.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get  $(m_1 + 1 + t_1, H'_{j1}) \triangleright W'_2.\theta_1$

Now instantiating  $H$  with  $H'_{j1}$  from Equation 8 and  $k_c$  with  $t_1$  we get

$$\begin{aligned}
& \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2]_V \wedge \\
& (\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e)) \quad (\text{R1})
\end{aligned}$$

Similarly we can apply Definition 1.6 on Equation 12 to get

$$\begin{aligned}
& \forall m. \forall \theta'_2.(m, W'_1.\theta_2) \sqsubseteq \theta'_2 \wedge \forall j_2 < m. \forall v. (\theta'_2, j_2, v) \in [\tau_1]_V \implies \\
& (\theta'_2, j_2, e_{h2}[v/x]) \in [\tau_2]_E^{\ell_e}
\end{aligned}$$

We instantiate  $m$  with  $m_2 + 2 + t_2$  where  $t_2$  is the number of steps in which  $e_{h2}$  reduces

$$\begin{aligned}
& \forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies \\
& (\theta', j_1, e_{h2}[v/x]) \in [\tau_2]_E^{\ell_e} \quad (\text{FB-AC2})
\end{aligned}$$

Since from Equation 8 we have

$$(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Therefore from Lemma 1.14 we get

$$\forall m. (W'_2.\theta_2, m, v'_{j2}) \in [\tau_1]_V$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, v'_{j2}) \in [\tau_1]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_2$ ,  $j_1$  with  $m_2 + 1 + t_2$  and  $v$  with  $v'_{j2}$  from Equation 8 in FB-AC2 we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, e_{h2}[v'_{j2}/x]) \in [\tau_2]_E^{\ell_e}$$

From Definition 1.7, we get

$$\begin{aligned}
& \forall H.(m_2 + 1 + t_2, H) \triangleright W'_2.\theta_2 \wedge \forall k_c < (m_2 + 1 + t_2).(H, e_{h2}[v'_{j1}/x]) \Downarrow_{k_c} (H'_2, v'_2) \implies \\
& \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1 + t_2 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1 + t_2 - k_c), v'_2) \in [\tau_2]_V \wedge \\
& (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e))
\end{aligned}$$

Since from Equation 8 we have  $(n - i - j, H'_{j1}, H'_{j1}) \triangleright W'_2$

Therefore from Lemma 1.26 we get  $\forall m.(m, H'_{j2}) \triangleright W'_2.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get  $(m_2 + 1 + t_2, H'_{j2}) \triangleright W'_2.\theta_2$

Now Instantiating  $H$  with  $H'_{j2}$  from Equation 8 and  $k_c$  with  $t_2$ .

$$\begin{aligned}
& \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2]_V \wedge \\
& (\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e)) \quad (\text{R2})
\end{aligned}$$

In order to prove FB-A0 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_2.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

–  $dom(W'.\theta_1) \subseteq dom(H'_1) \wedge dom(W.\theta_2) \subseteq dom(H'_2)$ :

From R1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get

$dom(W'.\theta_1) \subseteq dom(H'_1)$

Similarly, from R2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $dom(W'.\theta_2) \subseteq dom(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (dom(W'.\theta_1) \times dom(W'.\theta_2))$ :

Since from Equation 8 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$  therefore from

Definition 1.9 we know that  $(W'_2.\hat{\beta}) \subseteq (dom(W'_2.\theta_1) \times dom(W'_2.\theta_2))$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  therefore

$(W'_2.\hat{\beta}) \subseteq (dom(\theta'_1) \times dom(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$

$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

4 cases arise for each  $(a_1, a_2) \in W'_2.\hat{\beta}$

i.  $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

We know from Equation 8 that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$

Therefore from Definition 1.9 we have

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

Since  $W'.\hat{\beta} = W'_2.\hat{\beta}$  by construction therefore

$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From Equation 8 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$

This means from Definition 1.9 that

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_2 \sqsubseteq W'$  and  $n - n' - 1 < n - i - j - 1$  (since  $n' = i + j + t_1$  where  $t_1$  is the number of steps taken by  $e_{h_1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce and  $j$  is the number of steps taken by  $e_2 \gamma \downarrow_1$  to reduce) therefore from Lemma 1.16 we get

$(W', n - n' - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$

ii.  $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From R1 and R2 we know that

$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell')$

$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell')$

This means we have

$$\begin{aligned}\exists \ell'. W'_2.\theta_1(a_1) &= A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell' \text{ and} \\ \exists \ell'. W'_2.\theta_2(a_2) &= A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell'\end{aligned}$$

Since  $pc \sqcup \ell \sqsubseteq \ell_e$  (given) and  $\ell \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from R1 and R2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ . Therefore from Definition 1.8 we have

$$\begin{aligned}(\theta'_1, m_1, H'_1(a_1)) &\in [\theta'_1(a_1)]_V \text{ and} \\ (\theta'_2, m_2, H'_2(a_1)) &\in [\theta'_2(a_2)]_V\end{aligned}$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iii. } H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same reasoning as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A:$$

From R2 we know that

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e$  in the world before the modification. Since  $pc \sqcup \ell \sqsubseteq \ell_e$  (given) and  $\ell \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 8 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$  that means from Definition 1.9 that  $(W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in [W'_2.\theta_1(a_1)]_V^A$ . Since  $(\ell_e) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H'_{j_1}(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_2.\theta_1, m, H'_{j_1}(a_1)) \in W'_2.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_2.\theta_2, m, H'_{j_2}(a_2)) \in W'_2.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.15 we get

$$(\theta'_1, m_1, H'_{j_1}(a_1)) \in \theta'_1(a_1)$$

Since from R2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$  therefore from Definition 1.8

$$\text{we know that } (\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W'.\theta_i(a_i)]_V:$$

$$\underline{i = 1}$$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W'.\theta_i(a_i)]_V$$

Like before we instantiate Equation 11 and Equation 12 with  $m + 2 + t_1$  and  $m + 2 + t_2$  respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_2. \theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2]_V \wedge \\ & (\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2. \theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2. \theta_1). \theta'_1(a) \searrow (\ell_e)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_2. \theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2]_V \wedge \\ & (\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2. \theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2. \theta_2). \theta'_2(a) \searrow (\ell_e)) \end{aligned}$$

Since we have  $(m + 1, H'_1) \triangleright \theta'_1$  and  $(m + 1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 1.8

$$\underline{i = 2}$$

Symmetric to  $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \searrow \ell$  and since  $\ell \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

From R1 and R2 we and Definition 1.4 we get the desired.

#### 4. FG-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\perp]_E^{\mathcal{A}}$

Say  $e_1 = (e_1, e_2) (\gamma \downarrow_1)$  and  $e_2 = (e_1, e_2) (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \times \tau_2)^\perp]_E^{\mathcal{A}}$  it suffices to prove that

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \stackrel{\mathcal{A}}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ & \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^{\mathcal{A}} \end{aligned}$$

This means that given some  $H_1, H_2$  and  $n' < n$  s.t

$$(n, H_1, H_2) \stackrel{\mathcal{A}}{\triangleright} W \wedge (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^{\mathcal{A}} \quad (13)$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [\tau_1]_E^{\mathcal{A}}$$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{p11}, H_{p12}. (n, H_{p11}, H_{p12}) \stackrel{\mathcal{A}}{\triangleright} W \wedge \forall i < n. (H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11}) \wedge (H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12}) \implies \\ & \exists W'_1 \sqsupseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{\mathcal{A}}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \end{aligned}$$

Instantiating  $H_{p11}$  with  $H_1$  and  $H_{p22}$  with  $H_2$  in IH1 and since the  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12})$ . Hence we get

$$\exists W'_1 \sqsubseteq W.(n-i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n-i, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \quad (14)$$

$$\underline{\text{IH2}} (W, n-i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_2]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{p21}, H_{p22}.(n-i, H_{p21}, H_{p22}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n-i.(H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22}) \implies$$

$$\exists W'_2 \sqsubseteq W'_1.(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

Instantiating  $H_{p21}$  with  $H'_{p11}$  and  $H_{p22}$  with  $H'_{p21}$  and in IH2. Since  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  has reduced with  $i < n'$  steps. Therefore we know that  $\exists j < n' - i < n - i$  s.t  $(H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{p21}, v'_{p21})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p22}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22})$ . Hence we get

since the  $(e_1, e_2)$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22})$ . Hence we get

$$\exists W'_2 \sqsubseteq W'_1.(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A \quad (15)$$

In order to prove Equation 13 we instantiate  $W'$  in Equation 13 with  $W'_2$  we are required to show the following:

- $W \sqsubseteq W'_2$ :

Since  $W \sqsubseteq W'_1$  from Equation 14 and  $W'_1 \sqsubseteq W'_2$  from Equation 15

Therefore,  $W \sqsubseteq W'_2$  from Definition 1.3

- $(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W'$ :

Here  $n' = i + j + 1$

From evaluation rule of products we know that  $H'_1 = H'_{p21}$  and  $H'_2 = H'_{p22}$

From Equation 15 we know that  $(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.20 we get  $(n-i-j-1, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

- $(W', n-i-j-1, v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^A$ :

From evaluation rule of products we know that  $v'_1 = (v'_{p11}, v'_{p21})$  and  $v'_2 = (v'_{p12}, v'_{p22})$

We are required to show

$$- (W'_2, n-i-j-1, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j-1, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A:$$

From Equation 14 and Equation 15 we know that

$$(W'_2, n-i-j, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

Therefore from Lemma 1.16 we get

$$(W'_2, n-i-j-1, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j-1, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

5. FG-fst:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$



To prove:  $(W, n, (\text{fst}(e_i)) (\gamma \downarrow_1), (\text{fst}(e_i)) (\gamma \downarrow_2)) \in [\tau_1]_E^A$

Say  $e_1 = (\text{fst}(e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{fst}(e_i)) (\gamma \downarrow_2)$

From Definition 1.5 it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1]_V^A \quad (16)$$

### IH1

$$(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\ell]_E^A$$

This means from Definition 1.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow \\ (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{fst}(e_i)$  reduces to value reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{fst}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A \quad (17)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A$  from Equation 17

- Case  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V^A$$

Let  $v'_{i1} = (v_{i1}, v_{j1})$  and  $v'_{i2} = (v_{i2}, v_{j2})$

Again from Definition 1.4 it means that

$$(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W'_1, n - i, v_{i2}, v_{j2}) \in [\tau_2]_V^A \quad (\text{F1})$$

In order to prove Equation 16 we choose  $W'$  as  $W'_1$  and from the evaluation rule of  $\text{fst}$  we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . Also, from reduction rules we know that  $n' = i + 1$ . And then we need to show:

- $W \sqsubseteq W'_1$ :

Directly from Equation 17

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

Since from Equation 17 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 1.20 we get  $(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

–  $(W'_1, n - n', v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A$ :

From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$

From F1 we know that  $(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^A$

Therefore from Lemma 1.16 we get  $(W'_1, n - i - 1, v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^A$

• Case  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 1.6 we know that

(a)  $\forall m. (W'_1.\theta_1, m, v'_{i1}) \in [(\tau_1 \times \tau_2)]_V$  and

(b)  $\forall m. (W'_1.\theta_2, m, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V$

where

$v'_{i1} = (v_{i1}, v_{i2})$  and  $v'_{i2} = (v_{j1}, v_{j2})$

In order to prove Equation 16 we choose  $W'$  as  $W'_1$  and from the evaluation rule of fst we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . And then we need to show:

–  $W \sqsubseteq W'_1$ :

Directly from Equation 17

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 17 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 1.20 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

–  $(W'_1, n - n', v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A$ :

From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$

Let  $\tau_1 = \mathbf{A}^{\ell_i}$  Since  $\tau_1 \searrow \ell$  and since  $\ell \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$\forall m_1. (W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$

and

$\forall m_2. (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \tag{18}$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V \tag{19}$$

Instantiating (a) with  $m_1$

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \wedge (W'_1.\theta_1, m_1, v_{i2}) \in [\tau_2]_V \tag{20}$$

Instantiating (b) with  $m_2$

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V \wedge (W'_1.\theta_2, m_2, v_{j2}) \in [\tau_2]_V \tag{21}$$

From Equation 20 and Equation 21 we get

$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and  $(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$

6. FG-snd:

Symmetric case as FG-fst

7. FG-inl:

$$\frac{\Gamma \vdash_{pc} e_i : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(W, n, (\text{inl}(e_i)) (\gamma \downarrow_1), (\text{inl}(e_i)) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\perp]_E^A$

Say  $e_1 = (\text{inl}(e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{inl}(e_i)) (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 + \tau_2)^\perp]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp]_V^A \quad (22)$$

$$\text{IH1 } (W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]_V^A \quad (23)$$

Instantiating  $W'$  in Equation 22 with  $W'_1$ . Also from reduction relation we know that  $n' = i + 1$  we are required to show the following:

- $W \sqsubseteq W'_1$ :

Directly from Equation 23

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

From Equation 23 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

Therefore from Lemma 1.20 we get

$$(n - n', H'_1, H'_2) \triangleright^A W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^+]^A_{\mathcal{V}}$ :

From evaluation rule of  $\text{inl}$  we know that  $v'_1 = \text{inl}(v'_{i1})$  and  $v'_2 = \text{inl}(v'_{i2})$

We are required to show

- $(W'_1, n - n', v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$ :

From Equation 23 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$

Therefore from Lemma 1.16 we get

$$(W'_1, n - i - 1, v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$$

#### 8. FG-inr:

Symmetric case to FG-inl.

#### 9. FG-case:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{i1} : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_{i2} : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e_i, x.e_{i1}, y.e_{i2}) : \tau}$$

To prove:  $(W, (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1), (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)) \in [(\tau)]^A_E$

Say  $e_1 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1)$  and  $e_2 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]^A_{\mathcal{V}} \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]^A_{\mathcal{V}} \quad (24)$$

$$\underline{\text{IH1}} \quad (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\ell]^A_E$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell]^A_{\mathcal{V}}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell]^A_{\mathcal{V}} \quad (25)$$

IH2:

$$(W'_1, n - i, (e_{i1}) (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\}), (e_{i1}) (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \in [(\tau)]^A_E$$

This means from Definition 1.5 we get

$$\forall H_{j_1}, H_{j_2}. (n - i, H_{j_1}, H_{j_2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_1, e_{i1} (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\})) \Downarrow_j (H'_{j_1}, v'_{j_1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \Downarrow (H'_{j_2}, v'_{j_2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$$

Instantiating  $H_{j_1}$  with  $H'_1$  and  $H_{j_2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{j_1}, v'_{j_1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{j_2}, v'_{j_2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A \quad (26)$$

IH3:

$$(W'_1, n - i, (e_{i2}) (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\}), (e_{i2}) (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \in [(\tau)]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{k_1}, H_{k_2}. (n - i, H_{k_1}, H_{k_2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < n - i. (H_1, e_{i2} (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\})) \Downarrow_k (H'_{k_1}, v'_{k_1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \Downarrow (H'_{k_2}, v'_{k_2}) \implies$$

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k_1}, H'_{k_2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [(\tau)]_V^A$$

Instantiating  $H_{k_1}$  with  $H'_1$  and  $H_{k_2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i2}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{k_1}, v'_{k_1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{k_2}, v'_{k_2})$ . Hence we get

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k_1}, H'_{k_2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [(\tau)]_V^A \quad (27)$$

We case analyze  $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 + \tau_2)^\ell]_V^A$  from Equation 25

- Case  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.4 2 further cases arise:

- $v'_1 = \text{inl}(v_{i1})$  and  $v'_2 = \text{inl}(v_{i2})$ :

In this case from Definition 1.4 we know that  $(W, n - i, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

In order to prove Equation 24 we choose  $W'$  as  $W'_2$  from Equation 26 and from the first evaluation rule of case we know that  $H'_1 = H'_{j_1}$  and  $H'_2 = H'_{j_2}$ . Also we know from the evaluation rule that  $n' = i + j + 1$ . And then we need to show:

- \*  $W \sqsubseteq W'_2$ :

Since  $W \sqsubseteq W'_1$  from Equation 25 and  $W'_1 \sqsubseteq W'_2$  from Equation 26

Therefore,  $W \sqsubseteq W'_2$  from Definition 1.3

- \*  $(n - n', H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$ :

From Equation 26 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.20 we get

$$(n - i - j - 1, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$$

\*  $(W'_2, n - n', v'_1, v'_2) \in [\tau]_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{j_1}$  and  $v'_2 = v'_{j_2}$

From Equation 26 we know that  $(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [\tau]_V^A$

Therefore from Lemma 1.16 we get

$(W'_2, n - i - j - 1, v'_{j_1}, v'_{j_2}) \in [\tau]_V^A$

-  $v'_1 = \text{inr}(v_{i_1})$  and  $v'_2 = \text{inr}(v_{i_2})$ :

In this case from Definition 1.4 we know that  $(W, v_{i_1}, v_{i_2}) \in [\tau_2]_V^A$

In order to prove Equation 24 we choose  $W'$  as  $W'_3$  from Equation 27 and from the second evaluation rule of case we know that  $H'_1 = H'_{k_1}$  and  $H'_2 = H'_{k_2}$ . Also we know from the evaluation rule that  $n' = i + k + 1$ . And then we need to show:

\*  $W \sqsubseteq W'_3$ :

Since  $W \sqsubseteq W'_1$  from Equation 25 and  $W'_1 \sqsubseteq W'_3$  from Equation 27

Therefore,  $W \sqsubseteq W'_3$  from Definition 1.3

\*  $(n - n', H'_1, H'_2) \triangleright^A W'_3$ :

From Equation 27 we know that  $(n - i - k, H'_{k_1}, H'_{k_2}) \triangleright^A W'_3$

Therefore from Lemma 1.20 we get

$(n - i - k - 1, H'_{k_1}, H'_{k_2}) \triangleright^A W'_3$

\*  $(W'_3, n - n', v'_1, v'_2) \in [\tau]_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{k_1}$  and  $v'_2 = v'_{k_2}$

From Equation 27 we know that  $(W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [\tau]_V^A$

Therefore from Lemma 1.16 we get

$(W'_3, n - i - k - 1, v'_{k_1}, v'_{k_2}) \in [\tau]_V^A$

• Case  $\ell \not\sqsubseteq \mathcal{A}$ :

The following cases arise:

- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case1 :  
Exactly the same reasoning as in the  $v'_1 = \text{inl}(v_{i_1})$  and  $v'_2 = \text{inl}(v_{i_2})$  subcase of the  $\ell \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case2 :  
Exactly the same reasoning as in the  $v'_1 = \text{inr}(v_{i_1})$  and  $v'_2 = \text{inr}(v_{i_2})$  subcase of the  $\ell \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case2 :

From Equation 24 we know that we need to prove

$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]_V^A$

In this case since we know that  $\ell \not\sqsubseteq \mathcal{A}$ . Let  $\tau = A^{\ell_i}$  and since  $\tau \searrow \ell$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]_V^A$

From Definition 1.4 it will suffice to prove

$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1.(W'.\theta_1, m_1, v'_1) \in [(\tau)]_V) \wedge (\forall m_2.(W'.\theta_1, m_2, v'_2) \in [(\tau)]_V)$

This means it suffices to prove

$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau)]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau)]_V)$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau)]_V \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau)]_V) \quad (28)$$

Since we know that  $(W, n, \gamma) \in [\Gamma]_V^A$  (given) therefore from Lemma 1.24 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore by instantiating it at  $m_1 + 1 + j$  we know that

$$(W.\theta_1, m_1 + 1 + j, \gamma \downarrow_1) \in [\Gamma]_V \quad (29)$$

Next we apply Theorem 1.21 on  $e_{i1} \gamma \downarrow_1$ . Here  $j$  is the number of steps in which  $e_{i1} \gamma \downarrow_1$  reduces. We use  $\gamma \downarrow_1 \cup \{x \mapsto v'_{s1}\}$  as the unary substitution to get  $(W.\theta_1, m_1 + 1 + j, e_{i1} \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \in [(\tau)]_E^{pc}$

This means from Definition 1.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_1 + 1 + j, H_{c1}) \triangleright W_1.\theta_1 \wedge \forall l_c < (m_1 + 1 + j). (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1 + j - l_c, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1 + j - l_c, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W_1.\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since from Equation 25 we know that  $(n-i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.26 we get  $\forall m. (m, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + j$  we get  $(m_1 + 1 + j, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $H_{c2}$  with  $H'_1$  from Equation 25 and  $l_c$  with  $j$  we get

$$\begin{aligned} & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \quad (\text{CC1}) \end{aligned}$$

Similarly we apply Theorem 1.21 on  $e_{i2} \gamma \downarrow_2$ . Here  $j_2$  is the number of steps in which  $e_{i2} \gamma \downarrow_2$  reduces. We use  $\gamma \downarrow_2 \cup \{y \mapsto v'_{s2}\}$  as the unary substitution to get  $(W_1.\theta_2, m_2 + 1 + j_2, e_{i2} \gamma \downarrow_2 \cup \{y \mapsto v'_c\}) \in [(\tau)]_E^{pc}$

This means from Definition 1.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_2 + 1 + j_2, H_{c1}) \triangleright W_1.\theta_2 \wedge \forall l_c < m_2 + 1 + j_2. (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1 + j_2 - l_c, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1 + j_2 - l_c, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since from Equation 25 we know that  $(n-i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.26 we get  $\forall m. (m, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + j_2$  we get  $(m_2 + 1 + j_2, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $H_{c2}$  with  $H'_2$  (from Equation 25) and  $l_c$  with  $j_2$  to get

$$\begin{aligned} & \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \quad (\text{CC2}) \end{aligned}$$

We choose

$$W_n.\theta_1 = \theta'_1 \text{ (from CC1)}$$

$$W_n.\theta_2 = \theta'_2 \text{ (from CC2)}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \text{ (from Equation 25)}$$

In order to prove Equation 24 we choose  $W'$  as  $W_n$

i.  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H'_2):$$

From (CC1) we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from (CC2) we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since from Equation 25 we have  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Definition 1.9 we get  $(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore  $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

$$- \forall(a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge \\ (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

4 cases arise for each  $a_1$  and  $a_2$

A.  $H'_{j1}(a_1) = H'_1(a_1) \wedge H'_{j2}(a_2) = H'_2(a_2):$

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2):}$$

We know from Equation 25 that  $(n - i, H'_1, H'_2) \triangleright W'_1$

Therefore from Definition 1.9 we have

$$\forall(a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:}$$

From Equation 25 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

This means from Definition 1.9 that

$$\forall(a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) \wedge (W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in \lceil W'_1.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$  (since  $n' = i + t_1 + 1$  where  $t_1$  is the number of steps taken by  $e_{i1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce) therefore from Lemma 1.16 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

B.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2):$

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2):}$$



Same as before

$$\frac{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^{\mathcal{A}}}{\text{From (CC1) and (CC2) we know that}}$$

$$(\forall a. H'_1(a) \neq H'_{c1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell)) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell)) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell)) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell)) \sqsubseteq \ell'$$

Since  $\ell \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from (CC1) and (CC2),  $(m_1 + 1, H'_{c1}) \triangleright \theta'_1$  and  $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$ .

Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_{c1}(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V \text{ and}$$

$$(\theta'_2, m_2, H'_{c2}(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we

get (here  $H'_1 = H'_{c1}$  and  $H'_2 = H'_{c2}$ )

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^{\mathcal{A}}$$

C.  $H'_{j1}(a_1) = H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2)$ :

$$\frac{W'.\theta_1(a_1) = W'.\theta_2(a_2):}{\text{Same as before}}$$

Same as before

$$\frac{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^{\mathcal{A}}}{\text{From (CC2) we know that}}$$

From (CC2) we know that

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell)) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $(pc \sqcup \ell)$  in the world before the modification. Since  $\ell \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 25 we know that  $(n - i, H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W'_1$  that means from Definition 1.9 that  $(W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^{\mathcal{A}}$ . Since  $((pc \sqcup \ell)) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H'_1(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_1(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_2(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.15 we get

$$(\theta'_1, m_1, H'_1(a_1)) \in \theta'_1(a_1)$$

Since from (CC2) we know that  $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$  therefore from Definition 1.8 we know that  $(\theta'_2, m_2, H'_{c2}(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_{c1}(a_1), H'_{c2}(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^{\mathcal{A}}$$

D.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2)$ :

Symmetric case as above

–  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V$ :

$i = 1$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i).(W'.\theta_i, m, H'_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$$

Like before we apply Theorem 1.21 on  $e_{i1}$   $\gamma_1$  and  $e_{i2}$   $\gamma_2$  but this time using  $m + 1 + i$  and  $m + 1 + j$  where  $i$  and  $j$  are the number of steps in which  $e_{i1}$   $\gamma_1$  and  $e_{i2}$   $\gamma_2$  reduces respectively. This will give us

$$\begin{aligned} \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m + 1, v'_c) \in \llbracket (\tau) \rrbracket_V \wedge \\ (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

and

$$\begin{aligned} \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m + 1, v'_c) \in \llbracket (\tau) \rrbracket_V \wedge \\ (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_2). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since we have  $(m + 1, H'_{c1}) \triangleright \theta'_1$  and  $(m + 1, H'_{c2}) \triangleright \theta'_2$  therefore we get the desired from Definition 1.8

$$i = 2$$

Symmetric to  $i = 1$

$$\text{ii. } (W', n - n' - 1, v'_1, v'_2) \in \llbracket \tau_2 \rrbracket_V^A:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \searrow \ell$  and since  $\ell \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

From CC1 and CC2 we and Definition 1.4 we get the desired.

- (d) Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case1 :  
Symmetric case as before

10. FG-ref:

$$\frac{\Gamma \vdash_{pc} e_i : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(W, (\text{new } (e_i)) (\gamma \downarrow_1), (\text{new } (e_i)) (\gamma \downarrow_2)) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_E^A$

Say  $e_1 = (\text{new } (e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{new } (e_i)) (\gamma \downarrow_2)$

From Definition of  $\llbracket (\text{ref } \tau)^\perp \rrbracket_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A \quad (30)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in \llbracket \tau \rrbracket_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n.(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsubseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{ref}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$ . s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since  $\text{ref}(e_i)$  reduces with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsubseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A \quad (31)$$

From the evaluation rule of  $\text{ref}$  we know that  $H'_1 = H'_{i1} \cup \{a_{n1} \mapsto v_{i1}\}$  and  $H'_2 = H'_{i2} \cup \{a_{n2} \mapsto v_{i2}\}$

Inorder to prove Equation 30 we instantiate  $W'$  with  $W_n$  where  $W_n$  is

$$W_n.\theta_1 = W'_1.\theta_1 \cup \{a_{n1} \mapsto \tau\}$$

$$W_n.\theta_2 = W'_1.\theta_2 \cup \{a_{n2} \mapsto \tau\}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$$

Also we know that  $n' = i + 1$

We are now required to prove

- $W \sqsubseteq W_n$ :

From Equation 31 we know that  $W \sqsubseteq W'_1$  and  $W'_1 \sqsubseteq W_n$  by construction. Therefore from Definition 1.3,  $W \sqsubseteq W_n$

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W_n$ :

From Definition 1.9 it suffices to show that

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2):$$

From Equation 31 and by construction of  $W_n$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

From Equation 31 and by construction of  $W_n$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, n - n', H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A:$$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

From Equation 31 and by construction of  $W_n$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A:$$

From Equation 31 since we know that  $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$  that means

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$$

Therefore from Lemma 1.16 we get  $(n - i - 2 = n - n' - 1, \text{ since } n' = i + 1)$

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 2, H'_1(a_1), H'_2(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$$

Since  $W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$  and from Equation 31 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A$

Therefore combining the two we get

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

–  $\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$ :

From Equation 31 we have  $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$  that means from Definition 1.9 we have

$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W'_1.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$

Also from Equation 31 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in \llbracket \tau \rrbracket_V^A$

Therefore from Lemma 1.14 and Lemma 1.15 we get

$\forall m. (W'_1.\theta_1, m, v'_{i1}) \in \llbracket \tau \rrbracket_V$

and

$\forall m. (W'_1.\theta_2, m, v'_{i2}) \in \llbracket \tau \rrbracket_V$

Combining the two we get

$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$

•  $(W_n, n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A$ :

Here  $v'_1 = a_{n1}$  and  $v'_2 = a_{n2}$

Since  $(a_{n1}, a_{n2}) \in W_n$  and also  $W_n.\theta_1(a_{n1}) = W_n.\theta_1(a_{n1}) = \tau$

Therefore from Definition 1.4  $(W_n, v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A$

11. FG-deref:

$$\frac{\Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(W, n, !(e_i) (\gamma \downarrow_1), !(e_i) (\gamma \downarrow_2)) \in \llbracket (\tau') \rrbracket_E^A$

Say  $e_1 = !(e_i) (\gamma \downarrow_1)$  and  $e_2 = !(e_i) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$

$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\tau') \rrbracket_V^A$

This further means that given

$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\tau') \rrbracket_V^A \quad (32)$$

IH1  $(W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in \llbracket (\text{ref } \tau)^\ell \rrbracket_E^A$

This means from Definition 1.5 we get

$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$

$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\ell \rrbracket_V^A$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $!(e_i)$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$ . Similarly since  $!e_i$  reduces to value with  $\gamma \downarrow_2$  therefore  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell]_V^A \quad (33)$$

We case analyze on  $(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell]_V^A$  from Equation 33

- Case  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)]_V^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } (\tau))]_V^A$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{D1})$$

In order to prove Equation 32 we instantiate  $W'$  with  $W'_1$ . Also we know that  $n' = i + 1$

- $W'_1 \sqsupseteq W$ :

From Equation 33

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 33 we know that

$$(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 1.20 we get

$$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau')]_V^A$ :

From the evaluation rule of deref we know that  $v'_1 = H'_1(a_{i1})$  and  $v'_2 = H'_2(a_{i2})$

Since from Equation 33 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ , therefore from Definition 1.9 we know that

$$(W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in [W'_1.\theta_1(a_{i1})]_V^A$$

And from D1 we know that  $W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau$

Therefore  $(W'_1, v'_1, v'_2) \in [(\tau)]_V^A$

Since  $\tau <: \tau'$  Therefore from Lemma 1.27, we get

$$(W'_1, n - i - 1, v'_1, v'_2) \in [(\tau')]_V^A$$

- Case  $\ell \not\sqsubseteq \mathcal{A}$ :

From the evaluation rule of deref we know that  $v'_{i1} = a_1$  and  $v'_{i2} = a_2$

In this case from Definition 1.4 we know that

$$\forall m_1.(W'_1.\theta_1, m_1, a_1) \in [(\text{ref } \tau)]_V \quad (34)$$

and

$$\forall m_2.(W'_1.\theta_2, m_2, a_2) \in [(\text{ref } \tau)]_V \quad (35)$$

In order to prove Equation 32 we choose  $W'$  as  $W'_1$ . And then we need to show:

- $W \sqsubseteq W'_1$ :  
Directly from Equation 33
- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :  
From Equation 33 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$   
Therefore from Lemma 1.20 we get  
 $(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$
- $(W'_1, n - n', v'_1, v'_2) \in [\tau']_V^A$ :  
Let  $\tau' = A^{\ell_i}$  Since  $\tau' \searrow \ell$  and since  $\ell \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v'_1) \in [\tau']_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v'_2) \in [\tau']_V$$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau']_V \quad (36)$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau']_V \quad (37)$$

Since from Equation 33 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.26 we get

$$\forall m_{h1}. (m_{h1}, H'_1) \triangleright W'_1.\theta_1 \quad (38)$$

$$\forall m_{h2}. (m_{h2}, H'_2) \triangleright W'_1.\theta_2 \quad (39)$$

Instantiating  $m_{h1}$  in Equation 38 with  $m_1 + 1$  we get  $(m_1, H'_1) \triangleright W'_1.\theta_1$

Therefore from Definition 1.8, we get

$$\forall a \in \text{dom}(W'_1.\theta_1). (W'_1.\theta_1, m_1, H'_1(a)) \in [W'_1.\theta_1(a)]_V$$

Instantiating  $a$  with  $a_1$  we get  $(W'_1.\theta_1, m_1, H'_1(a_1)) \in [W'_1.\theta_1(a)]_V$

Since  $W'_1.\theta_1(a_{i1}) = \tau$  therefore we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau]_V$$

and since  $\tau <: \tau'$  therefore from Lemma 1.23 we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau']_V$$

Similarly we also get

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau']_V$$

Finally from Definition 1.4 we get

$$(W'_1, v'_1, v'_2) \in [(\tau')]_V^A$$

12. FG-assign:

$$\frac{\Gamma \vdash_{pc} e_{i1} : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_{i2} : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_{i1} := e_{i2} : \text{unit}}$$

To prove:  $(W, n, (e_{i1} := e_{i2}) (\gamma \downarrow_1), (e_{i1} := e_{i2}) (\gamma \downarrow_2)) \in [(\mathbf{unit})]_E^A$

Say  $e_1 = (e_{i1} := e_{i2}) (\gamma \downarrow_1)$  and  $e_2 = (e_{i1} := e_{i2}) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ & \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\mathbf{unit})]_V^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\mathbf{unit})]_V^A \quad (40)$$

IH1  $(W, n, (e_{i1}) (\gamma \downarrow_1), (e_{i1}) (\gamma \downarrow_2)) \in [(\mathbf{ref} \ \tau)^\ell]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \\ & \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_{i1} := e_{i2})$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{i1}, v'_{i1})$ . Similarly since  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_2$  therefore we also have  $(H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A \quad (41)$$

IH2  $(W, n - i, (e_{i2}) (\gamma \downarrow_1), (e_{i2}) (\gamma \downarrow_2)) \in [(\tau)]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{j1}, H_{j2}. (n - i, H_{j1}, H_{j2}) \triangleright^A W'_1 \wedge \forall j < n - i. (H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \\ & \exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \triangleright^A W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A \end{aligned}$$

Instantiating  $H_{j1}$  with  $H'_{i1}$  and  $H_{j2}$  with  $H'_{i2}$  in IH2 and since the  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  reduces  $\gamma \downarrow_1$  with  $i < n'$  steps therefore  $\exists j < (n' - i) < (n - i)$  s.t  $(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1})$ . Similarly we also have  $(H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \triangleright^A W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A \quad (42)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A$  from Equation 41

- Case  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)]_V^{\mathcal{A}}$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau))]_V^{\mathcal{A}}$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{A1})$$

In order to prove Equation 40 we instantiate  $W'$  with  $W'_2$

- $W'_2 \sqsupseteq W$ :

Since  $W'_1 \sqsupseteq W$  from Equation 41 and  $W'_2 \sqsupseteq W'_1$  from Equation 42

Therefore from Definition 1.3 we get  $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W'_2$ :

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

In order to prove  $(n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W'_2$  we need to show:

$$* \text{ dom}(W'_2.\theta_1) \subseteq \text{ dom}(H'_1) \wedge \text{ dom}(W'_2.\theta_2) \subseteq \text{ dom}(H'_2):$$

Directly from Equation 42

$$* W'_2.\hat{\beta} \subseteq (\text{ dom}(W'_2.\theta_1) \times \text{ dom}(W'_2.\theta_2)):$$

Directly from Equation 42

$$* \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge \\ (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}:$$

$$(a) \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2):$$

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :

From A1 we know that  $W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$

and since  $W'_1 \sqsubseteq W'_2$  therefore from Lemma 1.15 we get  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

- ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise

- iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise

- iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 42 and Lemma 1.16

$$(b) \forall (a_1, a_2) \in (W'_2.\hat{\beta}). (W'_2, n - n', H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}:$$

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :

Since  $H'_1(a_{i1}) = v'_{j1}$  and  $H'_1(a_{i2}) = v'_{j2}$

From A1 we know that  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

And since from Equation 42 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^{\mathcal{A}}$

Therefore from Lemma 1.16 we get

$$(W'_2, n - j - i - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}$$

- ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise

- iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise

- iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 42 and from Lemma 1.16



\*  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \llbracket W'_2.\theta_i(a_i) \rrbracket_V$ :

When  $i = 1$

Given some  $m$

$\forall a_1 \in \text{dom}(W'_2.\theta_1)$ .

· when  $a_1 = a_{i1}$ :

From Equation 42 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \llbracket (\tau) \rrbracket_V^A$  thus from Lemma 1.14 we know that

$\forall m_1. (W'_2.\theta_1, m_1, H'_1(a_1)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V$

Instantiating with  $m$  we get

$(W'_2.\theta_1, m, H'_1(a_1)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V$

· Otherwise:

From Equation 42 and Lemma 1.26

When  $i = 2$

Similar reasoning as with  $i = 1$

–  $(W'_1, n - n', val'_1, v'_2) \in \llbracket (\text{unit}) \rrbracket_V^A$ :

From evaluation rule assign we know that  $v'_1 = v'_2 = ()$

Directly from Definition 1.4

• Case  $\ell \not\sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$\forall m_1. (W'_1.\theta_1, m_1, a_{i1}) \in \llbracket (\text{ref } \tau) \rrbracket_V \quad (43)$$

$$\forall m_2. (W'_1.\theta_2, m_2, a_{i2}) \in \llbracket (\text{ref } \tau) \rrbracket_V \quad (44)$$

In order to prove Equation 40 we instantiate  $W'$  with  $W'_2$  and then we need to show that:

–  $W'_2 \sqsupseteq W$ :

Since  $W'_1 \sqsupseteq W$  from Equation 41 and  $W'_2 \sqsupseteq W'_1$  from Equation 42

Therefore from Definition 1.3 we get  $W'_2 \sqsupseteq W$

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$ :

From the evaluation rule assign we know that

$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}]$  and  $H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$

In order to prove  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$  we need to show:

\*  $\text{dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2)$ :

Directly from Equation 42

\*  $W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$ :

Directly from Equation 42

\*  $\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V^A$ :

(a) When  $(a_{i1}, a_{i2}) \in W'_2.\hat{\beta}$ :

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$ .

i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :

Instantiating Equation 43 and Equation 44 with  $n - n' - 1$  we get

$W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$

and since  $W'_1 \sqsubseteq W'_2$  therefore from Definition 1.3 we get  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

From Equation 42 we know that  $(W'_2, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$

Therefore  $(W'_2, H'_1(a_{i_1}), H'_2(a_{i_2})) \in [(\tau)]_V^A$

- ii. When  $a_1 = a_{i_1}$  and  $a_2 \neq a_{i_2}$ : This case cannot arise
  - iii. When  $a_1 \neq a_{i_1}$  and  $a_2 = a_{i_2}$ : This case cannot arise
  - iv. When  $a_1 \neq a_{i_1}$  and  $a_2 \neq a_{i_2}$ : From Equation 42
- (b) When  $(a_{i_1}, a_{i_2}) \notin W'_2.\hat{\beta}$ :

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$ .

i. When  $a_1 = a_{i_1}$  and  $a_2 = a_{i_2}$ : This case cannot arise

ii. When  $a_1 = a_{i_1}$  and  $a_2 \neq a_{i_2}$ :

From Equation 42 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$  and since  $(a_{i_1}, a_2) \in W'_2.\hat{\beta}$  therefore from Definition 1.9 we know that

$$(W'_2.\theta_1(a_{i_1}) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in [W'_2.\theta_1(a_{i_1})]_V^A) \quad (45)$$

Instantiating Equation 43 and Equation 44 with  $n - i - j - 1$  we get  $W'_1.\theta_1(a_{i_1}) = \tau$  therefore from monotonicity we also have  $W'_2.\theta_1(a_{i_1}) = \tau$ . As a result from Equation 45 we get  $W'_2.\theta_2(a_2) = \tau$

Also since from Equation 45  $(W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in [(\tau)]_V^A$  and  $\tau \searrow \ell, \ell \not\sqsubseteq \mathcal{A}$  therefore from Lemma 1.14 we know that

$$\forall m.(W'_2.\theta_1, m, H'_{j_1}(a_{i_1})) \in [(\tau)]_V \quad (46)$$

$$\forall m.(W'_2.\theta_2, m, H'_{j_2}(a_2)) \in [(\tau)]_V \quad (47)$$

Instantiating  $m$  with  $n - i - j - 1$  in Equation 46 and Equation 47 to get

$$(W'_2.\theta_1, n - i - j - 1, H'_{j_1}(a_{i_1})) \in [(\tau)]_V$$

and

$$(W'_2.\theta_2, n - i - j - 1, H'_{j_2}(a_2)) \in [(\tau)]_V$$

Since  $H'_1(a_{i_1}) = v'_{j_1}$  and  $H'_2(a_2) = H'_{j_2}(a_2)$

Again from Equation 42 we know that  $(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$ . This means from Lemma 1.14 and instantiating it with  $n - i - j - 1$  we get

$$(W'_2.\theta_1, n - i - j - 1, v'_{j_1}) \in [(\tau)]_V \quad (48)$$

Therefore from Equation 47 and Equation 48 we have  $(W'_2, n - i - j - 1, H'_1(a_{i_1}), H'_2(a_2)) \in [(\tau)]_V^A$

iii. When  $a_1 \neq a_{i_1}$  and  $a_2 = a_{i_2}$ :

Symmetric case as (ii)

iv. When  $a_1 \neq a_{i_1}$  and  $a_2 \neq a_{i_2}$ :

From Equation 42 and Definition 1.9

\*  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in [W'_2.\theta_i(a_i)]_V$ :  
When  $i = 1$   
 Given some  $m$   
 $\forall a_1 \in \text{dom}(W'_2.\theta_1)$ .  
 · when  $a_1 = a_{i1}$ :  
 From Equation 42 we know that  $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A$  thus from Lemma 1.14 we know that  
 $(W'_2.\theta_1, H'_1(a_1)) \in [W'_2.\theta_1(a_1)]_V$   
 · Otherwise:  
 From Equation 42 and Lemma 1.26  
When  $i = 2$   
 Similar reasoning as with  $i = 1$   
 –  $(W'_1, n - n', v'_1, v'_2) \in [(\text{unit})]_V^A$ :  
 From evaluation rule assign we know that  $v'_1 = v'_2 = ()$   
 Directly from Definition 1.4

□

**Lemma 1.26** (Binary heap well formedness implies unary heap well formedness).  $\forall H_1, H_2, W. (n, H_1, H_2) \triangleright W \implies \forall i \in \{1, 2\}. \forall m. (m, H_i) \triangleright W.\theta_i$

*Proof.* Directly from Definition 1.9

□

**Lemma 1.27** (Subtyping binary). *The following holds:*

1.  $\forall A, A'$ .

$$(a) \mathcal{L} \vdash A <: A' \implies [(A)]_V^A \subseteq [(A')]_V^A$$

2.  $\forall \tau, \tau'$ .

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V^A \subseteq [(\tau')]_V^A$$

$$(b) \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^A \subseteq [(\tau')]_E^A$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $[((\tau_1 \xrightarrow{\ell_e} \tau_2))]_V^A \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_V^A$

$$\text{IH1: } [(\tau'_1)]_V^A \subseteq [(\tau_1)]_V^A$$

$$\text{IH2: } [(\tau_2)]_E^A \subseteq [(\tau'_2)]_E^A$$

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2))]_{\mathcal{V}}^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_{\mathcal{V}}^A$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2))]_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_{\mathcal{V}}^A$

From Definition 1.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^A &\implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_{\mathcal{E}}^A &\wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 1.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A &\implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2]_{\mathcal{E}}^A &\wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) \end{aligned}$$

This means given some  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$  we need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A) :$$

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_{\mathcal{E}}^A) \quad (49)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A$

Thus from Equation 49 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_{\mathcal{E}}^A$

Finally using IH2 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A$

$$(b) \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) :$$

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}}$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}}$  and since  $\tau'_1 <: \tau_1$  therefore from Lemma 1.23 we get

$$(\theta'_l, k, v'_c) \in [\tau_1]_{\mathcal{V}} \quad (50)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \quad (51)$$

Therefore from Equation 50 and 51 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}$

Since  $\tau_2 <: \tau'_2$  and  $\ell'_e \sqsubseteq \ell_e$  therefore from Lemma 1.23 and 1.22 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}$$

(c)  $\forall \theta'_i \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_i, k, v'_c) \in \lceil \tau'_1 \rceil_V \implies (\theta'_i, k, e_2[v'_c/x]) \in \lceil \tau'_2 \rceil_E^{\ell'_e})$ :

Similar reasoning as in the previous case

## 2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $\lceil ((\tau_1 \times \tau_2)) \rceil_V^A \subseteq \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

IH1:  $\lceil (\tau_1) \rceil_V^A \subseteq \lceil (\tau'_1) \rceil_V^A$

IH2:  $\lceil (\tau_2) \rceil_V^A \subseteq \lceil (\tau'_2) \rceil_V^A$

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2)) \rceil_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2)) \rceil_V^A$

Therefore from Definition 1.4 we are given:

$$(W, n, v_1, v'_1) \in \lceil \tau_1 \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau_2 \rceil_V^A \quad (52)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

Again from Definition 1.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in \lceil \tau'_1 \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau'_2 \rceil_V^A$$

Since from Equation 52 we know that  $(W, n, v_1, v'_1) \in \lceil \tau_1 \rceil_V^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in \lceil \tau'_1 \rceil_V^A$

Similarly since  $(W, n, v_2, v'_2) \in \lceil \tau_2 \rceil_V^A$  from Equation 52 therefore from IH2 we have  $(W, n, v_2, v'_2) \in \lceil \tau'_2 \rceil_V^A$

## 3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $\lceil ((\tau_1 + \tau_2)) \rceil_V^A \subseteq \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

IH1:  $\lceil (\tau_1) \rceil_V^A \subseteq \lceil (\tau'_1) \rceil_V^A$

IH2:  $\lceil (\tau_2) \rceil_V^A \subseteq \lceil (\tau'_2) \rceil_V^A$

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2)) \rceil_V^A. (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2)) \rceil_V^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

2 cases arise

- (a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s1} = \text{inl } v_{i2}$ :  
From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1 \rceil_{\mathcal{V}}^A \quad (53)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1' \rceil_{\mathcal{V}}^A$$

From Equation 53 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1' \rceil_{\mathcal{V}}^A$$

- (b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :  
From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2 \rceil_{\mathcal{V}}^A \quad (54)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2' \rceil_{\mathcal{V}}^A$$

From Equation 54 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2' \rceil_{\mathcal{V}}^A$$

4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\lceil ((\text{ref } \tau)) \rceil_{\mathcal{V}}^A \subseteq \lceil ((\text{ref } \tau)) \rceil_{\mathcal{V}}^A$

Directly from Definition 1.4

5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\lceil ((\mathbf{b})) \rceil_{\mathcal{V}}^A \subseteq \lceil ((\mathbf{b})) \rceil_{\mathcal{V}}^A$

Directly from Definition 1.4

6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\lceil ((\text{unit})) \rceil_{\mathcal{V}}^A \subseteq \lceil ((\text{unit})) \rceil_{\mathcal{V}}^A$

Directly from Definition 1.4

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell \subseteq \ell' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^\ell <: A^{\ell'}} \text{FGsub-label}$$

To prove:  $\lceil ((A^\ell)) \rceil_{\mathcal{V}}^A \subseteq \lceil ((A^{\ell'})) \rceil_{\mathcal{V}}^A$

2 cases arise

1.  $\ell \sqsubseteq \ell'$ :

From Definition 1.4 it suffices to prove:  $\lceil ((A)) \rceil_V^A \subseteq \lceil ((A')) \rceil_V^A$

This we get directly from IH (Statement (1))

2.  $\ell \not\sqsubseteq \ell'$ :

We need to prove that

$$\forall (W, n, v_1, v_2) \in \lceil A \rceil_V^A. (W, n, v_1, v_2) \in \lceil A' \rceil_V^A$$

From Definition 1.4 it suffices to prove:

$$\forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in \lceil A \rceil_V. (W(n). \theta_i, m, v_i) \in \lceil A \rceil_V \in \lceil A' \rceil_V$$

Since  $A <: A'$  therefore from Lemma 1.23 we get the desired

Proof of statement 2(b)

Given:  $\mathcal{L} \vdash \tau <: \tau'$

To prove:  $\lceil (\tau) \rceil_E^A \subseteq \lceil (\tau') \rceil_E^A$

This means we need to prove that

$$\forall (W, n, e_1, e_2) \in \lceil (\tau) \rceil_E^A. (W, n, e_1, e_2) \in \lceil (\tau') \rceil_E^A$$

This means given  $\forall (W, n, e_1, e_2) \in \lceil (\tau) \rceil_E^A$

It suffices to prove that  $(W, n, e_1, e_2) \in \lceil (\tau') \rceil_E^A$

From Definition 1.5 we know we are given:

$$\begin{aligned} & \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in \lceil \tau \rceil_V^A \quad (\text{Sub-exp1}) \end{aligned}$$

And we need prove that

$$\begin{aligned} & \forall H_{21}, H_{22}, k < n. (n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22}) \implies \\ \exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \end{aligned}$$

This means that we are given some  $H_{21}, H_{22}$  and  $k < n$  such that  $(n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22})$

It suffices to prove:

$$\exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \quad (55)$$

Instantiating (Sub-exp1) with  $H_{21}, H_{22}$  and  $k$  we get

$$\exists W' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W' \wedge (W', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \quad (56)$$

We choose  $W''$  in Equation 55 as  $W'$  from Equation 56 and we are done

□

**Theorem 1.28** (NI for FG). *Say*  $\text{bool} = (\text{unit} + \text{unit})$

$\forall v_1, v_2, e, \tau, n_1.$

$\emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_{-} (-, v'_2) \implies$

$v'_1 = v'_2$

*Proof.* Given some

$$\begin{aligned} & \emptyset \vdash_{\perp} v_1 : \mathbf{bool}^{\top} \wedge \emptyset \vdash_{\perp} v_2 : \mathbf{bool}^{\top} \\ & x : \mathbf{bool}^{\top} \vdash_{\perp} e : \mathbf{bool}^{\perp} \wedge \\ & (\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2) \end{aligned}$$

We need to prove

$$\overline{v'_1 = v'_2}$$

From Theorem 1.25 we have

$$\forall n. (\emptyset, n, v_1, v_2) \in \lceil \mathbf{bool}^{\top} \rceil_E^{\perp}$$

Therefore from Theorem 1.25 and from Definition 1.13 we have

$$\forall n. (\emptyset, n, e[v_1/x], e[v_2/x]) \in \lceil \mathbf{bool}^{\perp} \rceil_E^{\perp}$$

Therefore from Definition 1.5 we know that

$$\begin{aligned} & \forall n. (\forall H_1, H_2, j < n. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W' \sqsupseteq \\ & W.(n - j, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A) \end{aligned}$$

Instantiating with  $n_1 + 1$  and then with  $\emptyset, \emptyset, n_1$  we get

$$\exists W' \sqsupseteq W.(1, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', 1, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A$$

Since we have  $(W', 1, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A$  therefore from Definition 1.4 we get  $v'_1 = v'_2$   $\square$



## 1.2 Coarse-grained IFC enforcement (CG)

### 1.2.1 CG type system

**Term, type, constraint syntax:**

Expressions	$e ::= x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid \text{new } e \mid !e \mid e := e \mid () \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid \text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::= \perp \mid \top \mid l \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::= \mathbf{b} \mid \text{unit} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{Labeled } \ell \tau \mid \mathbb{C} \ell_1 \ell_2 \tau$

**Type system:**  $\boxed{\Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types  $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau, \text{unit}$  are included.)

$$\begin{array}{c}
\frac{\Gamma \vdash e : \tau}{\Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{CG-label} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau} \text{CG-unlabel} \\
\\
\frac{\Gamma \vdash e : \mathbb{C} \ell \ell' \tau}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell \perp (\text{Labeled } \ell' \tau)} \text{CG-toLabeled} \qquad \frac{\Gamma \vdash e : \tau}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell \ell' \tau} \text{CG-ret} \\
\\
\frac{\Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \quad \Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau'} \text{CG-bind} \\
\\
\frac{\Gamma \vdash e : \tau' \quad \mathcal{L} \vdash \tau' <: \tau}{\Gamma \vdash e : \tau} \text{CG-sub} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)} \text{CG-ref} \\
\\
\frac{\Gamma \vdash e : \text{ref } \ell' \tau}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell' \tau)} \text{CG-deref} \\
\\
\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}} \text{CG-assign}
\end{array}$$

Figure 4: Type system of CG.

### 1.2.2 CG semantics

Judgement:  $e \Downarrow_i v$  and  $(H, e) \Downarrow_i^f (H', v)$

### 1.2.3 Logical relation for CG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type})) \times (\text{Loc} \leftrightarrow \text{Loc})$

**Definition 1.29** ( $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

$$\begin{array}{c}
\frac{}{\mathcal{L} \vdash \tau <: \tau} \text{CGsub-refl} \qquad \frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{CGsub-arrow} \\
\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{CGsub-prod} \qquad \frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{CGsub-sum} \\
\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{CGsub-labeled} \\
\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'} \text{CGsub-monad}
\end{array}$$

Figure 5: CG subtyping

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x. e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 e_2 \Downarrow_{i+j+k+1} v_3} \text{cg-app} \qquad \frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{cg-prod} \\
\frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{cg-fst} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{cg-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{cg-inl} \\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{cg-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{cg-case1} \\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{cg-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{cg-Lb} \qquad \frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-ret} \\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{cg-bind} \\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-unlabel} \qquad \frac{e \Downarrow_i v \quad (H, v) \Downarrow_j^f (H', v')}{(H, \text{toLabeled}(e)) \Downarrow_{i+j+1}^f (H', \text{Lb}(v'))} \text{cg-toLabeled} \\
\frac{e \Downarrow_i \text{Lb}v \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1}^f (H[a \mapsto \text{Lb}v], a)} \text{cg-ref} \qquad \frac{e \Downarrow_i a}{(H, !e) \Downarrow_{i+1}^f (H, H(a))} \text{cg-deref} \\
\frac{e_1 \Downarrow_i a \quad e_2 \Downarrow_j \text{Lb}v}{(H, e_1 := e_2) \Downarrow_{i+j+1}^f (H[a \mapsto \text{Lb}v], ())} \text{cg-assign} \\
\frac{e \in \{x, \lambda y. -, \text{ret } -, \text{bind}(-, -, -), \text{unlabel}(-), \text{toLabeled}(-), \text{new } (-), !-, - := -\}}{e \Downarrow_0 e} \text{cg-val}
\end{array}$$

Figure 6: CG semantics

**Definition 1.30** ( $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

**Definition 1.31** (Value Equivalence).

$$\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_V^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_V \wedge (W.\theta_2, j, v_2) \in [\tau]_V & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

**Definition 1.32** (Binary value relation).

$$\begin{aligned} [\mathbf{b}]_V^{\mathcal{A}} &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\ [\mathbf{unit}]_V^{\mathcal{A}} &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\ [\tau_1 \times \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^{\mathcal{A}} \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^{\mathcal{A}}\} \\ [\tau_1 + \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^{\mathcal{A}}\} \cup \\ &\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^{\mathcal{A}}\} \\ [\tau_1 \rightarrow \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, \lambda x.e_1, \lambda x.e_2) \mid \\ &\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\ &\quad ((W', j, v_1, v_2) \in [\tau_1]_V^{\mathcal{A}} \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^{\mathcal{A}}) \wedge \\ &\quad \forall \theta_l \sqsupseteq W.\theta_1, v_c, j. \\ &\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\ &\quad \forall \theta_l \sqsupseteq W.\theta_2, v_c, j. \\ &\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\ [\text{ref } \ell \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, a_1, a_2) \mid \\ &\quad (a_1, a_2) \in W.\hat{\beta} \wedge W.\theta_1(a_1) = W.\theta_2(a_2) = \text{Labeled } \ell \tau\} \\ [\text{Labeled } \ell \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\ [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, v_1, v_2) \mid \\ &\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\ &\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ &\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ &\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ &\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell')) \wedge \\ &\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\} \end{aligned}$$

**Definition 1.33** (Binary expression relation).

$$[\tau]_E^{\mathcal{A}} \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}\}$$

**Definition 1.34** (Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta' \sqsupseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbb{C} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

**Definition 1.35** (Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

**Definition 1.36** (Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

**Definition 1.37** (Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

**Definition 1.38** (Binary substitution).  $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

**Definition 1.39** (Unary substitution).  $\delta : \text{Var} \mapsto \text{Val}$

**Definition 1.40** (Unary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

**Definition 1.41** (Binary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

### 1.2.4 Soundness proof for CG

**Lemma 1.42** (Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

*Proof.* Proof by induction on  $\tau$

1. Case  $\mathbf{b}$ ,  $\mathbf{unit}$ :

From Definition 1.34

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 1.32 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 1.34, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 1.34, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 1.32 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A$  (S0)

IH1:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH2:  $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some  $m$  and we are required to prove:

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

Also from (S02) we know that given some  $m$  and we are required to prove:

$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH1 with  $m$  from (S01) to get

$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 1.34, we get

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH2 with  $m$  from (S02) to get

$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 1.34, we get

$(W.\theta_2, m, \text{inl}(v_{j1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{j2})$

Symmetric reasoning as in the (a) case above

#### 4. Case $\tau_1 \rightarrow \tau_2$ :

Given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in \lceil \tau_1 \rightarrow \tau_2 \rceil_V^A$

This means from Definition 1.32 we know that

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \lceil \tau_1 \rceil_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in \lceil \tau_2 \rceil_E^A)$   
 $\wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, i, e_1[v_c/x]) \in \lfloor \tau_2 \rfloor_E)$   
 $\wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, k, e_2[v_c/x]) \in \lfloor \tau_2 \rfloor_E)$  (L0)

To prove:

(a)  $\forall m. (W.\theta_1, m, \lambda x.e_1) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$ :

This means from Definition 1.34 we need to prove:

$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

This further means that we have some  $\theta', j$  and  $v$  s.t

$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

And we need to prove:  $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

Therefore we get  $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

(b)  $\forall m. (W.\theta_2, m, \lambda x.e_2) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$ :

Similar reasoning with  $e_2$

#### 5. Case ref $\ell \tau$ :

From Definition 1.32 and 1.34

6. Case Labeled  $\ell \tau$ :

Given  $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 1.31 we know that

$$(W, n, v_1, v_2) \in [\tau]_V^A$$

Therefore from IH we get  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 1.31

7. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

This means from Definition 1.32 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 1.34 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t. } (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case  $l = 1$

And given some  $m$  and  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with  $l = 1$  and the given  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  we get the desired.

Case  $l = 2$

Symmetric reasoning as in the previous case above

□

**Lemma 1.43** (Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$(\theta, m, v) \in \lfloor \tau \rfloor_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \lfloor \tau \rfloor_V$

*Proof.* Proof by induction on  $\tau$

1. case **b**, unit:

Directly from Definition 1.34

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

This means from Definition 1.34 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V \wedge (\theta, m, v_2) \in \lfloor \tau_2 \rfloor_V$

IH1 :  $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

IH2 :  $(\theta', m', v_2) \in \lfloor \tau_2 \rfloor_V$

We get the desired from IH1, IH2 and Definition 1.34

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

This means from Definition 1.34 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V$

IH :  $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

Therefore from IH and Definition 1.34 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(\theta, m, (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

To prove:  $(\theta', m', (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

This means from Definition 1.34 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E \quad (57)$$

Similarly from Definition 1.34 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V \implies (\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$



Instantiating Equation 57 with  $\theta'''$ ,  $k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

5. case ref  $\ell \tau$ :

From Definition 1.34 and Definition 1.29

6. case Labeled  $\ell \tau$ :

Given:  $(\theta, m, (\text{Lb } v)) \in [\text{Labeled } \ell \tau]_V$

To prove:  $(\theta', m', (\text{Lb } v)) \in [\text{Labeled } \ell \tau]_V$

This means from Definition 1.34 we know that  $(\theta, m, v) \in [\tau]_V$

IH:  $(\theta', m', v) \in [\tau]_V$

Therefore from IH and Definition 1.34 we get the desired

7. case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(\theta, m, e) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

To prove:  $(\theta', m', e) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 1.34 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 1.34 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in [\tau]_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t } (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in [\tau]_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0),  $k$  with  $k_1$ ,  $\theta_e$  with  $\theta_{e1}$ ,  $H$  with  $H_1$  and  $j$  with  $j_1$ . We know that  $k_1 < m' < m$ ,  $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$ ,  $(k_1, H_1) \triangleright \theta_{e1}$ ,  $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$  and  $j_1 < k_1$ . Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in [\tau]_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

**Lemma 1.44** (Monotonicity binary). *The following holds:*

$$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$$

$$(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**, unit:

From Definition 1.32

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH1} : (W', n', v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH2} : (W', n', v_{i2}, v_{j2}) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

From IH1, IH2 and Definition 1.32 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 we know that we are given

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH} : (W', n', v_{i1}, v_{i2}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

Therefore from Definition 1.32 we get

$$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

(b)  $v_1 = \text{inr}(v_{i1})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \rightarrow \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

This means from Definition 1.32 we know that the following holds

$$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_{\mathcal{E}}^{\mathcal{A}}) \quad (\text{BM-A0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_{\mathcal{E}}) \quad (\text{BM-A1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_{\mathcal{E}}) \quad (\text{BM-A2})$$

Similarly from Definition 1.32 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$$

(b)  $\forall \theta'_i \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau_1]_V \implies (\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$$(\theta'_i, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_i, k$  and  $v'_c$  we get

$$(\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E$$

(c)  $\forall \theta'_i \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau_1]_V \implies (\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$$(\theta'_i, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_i, k$  and  $v'_c$  we get

$$(\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E$$

5. Case ref  $\ell \tau$ :

From Definition 1.32 and Definition 1.30

6. Case Labeled  $\ell \tau$ :

Given:  $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

To prove:  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

From Definition 1.32 2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, n, v_1, v_2) \in [\tau]_V^A$

Therefore from IH we know that  $(W', n', v_1, v_2) \in [\tau]_V^A$

Hence from Definition 1.32 we get  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case we know that  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $(W.\theta_2, m, v_2) \in [\tau]_V$

Since  $W.\theta_1 \sqsubseteq W'.\theta_1$  (from Definition 1.30). Therefore from Lemma 1.43 we know that  $\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since  $W.\theta_2 \sqsubseteq W'.\theta_2$  (from Definition 1.30). Therefore from Lemma 1.43 we know that

$$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$$

Finally from Definition 1.32 we get  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

7. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

To prove:  $(W', n', v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

From Definition 1.32 we are given that

$$\begin{aligned}
& (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \\
& \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\
& \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
& \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{BM-M0})
\end{aligned}$$

Similarly from Definition 1.32 it suffices to prove that

$$\begin{aligned}
& \text{(a)} \quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \\
& \quad \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \quad \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)): \\
& \quad \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\
& \quad (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k
\end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given  $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  and since we know that  $n' \leq n$  and  $W \sqsubseteq W'$  we get the desired

$$\begin{aligned}
& \text{(b)} \quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
& \quad \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
& \quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& \quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)):
\end{aligned}$$

Similar reasoning as in the previous case but using Lemma 1.43

□

**Lemma 1.45** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'.$   
 $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

*Proof.* Given:  $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove:  $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 1.40 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$$

And again from Definition 1.40 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

- $dom(\Gamma) \subseteq dom(\delta)$ :

Given

- $\forall x \in dom(\Gamma).(\theta', n', \delta(x)) \in [\Gamma(x)]_V$ :

Since we know that  $\forall x \in dom(\Gamma).(\theta, n, \delta(x)) \in [\Gamma(x)]_V$  (given)

Therefore from Lemma 1.43 we get

$$\forall x \in dom(\Gamma).(\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

□

**Lemma 1.46** (Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .

$$(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in [\Gamma]_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(W', n', \gamma) \in [\Gamma]_V$

From Definition 1.41 it is given that

$$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And again from Definition 1.40 we are required to prove that

$$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

- $dom(\Gamma) \subseteq dom(\gamma)$ :

Given

- $\forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$ :

Since we know that  $\forall x \in dom(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$  (given)

Therefore from Lemma 1.44 we get

$$\forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

□

**Lemma 1.47** (Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .

$$(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 1.36 it is given that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta).(\theta, n - 1, H(a)) \in [\theta(a)]_V$$

And again from Definition 1.40 we are required to prove that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta).(\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

- $dom(\theta) \subseteq dom(H)$ :

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :

Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)

Therefore from Lemma 1.43 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

**Lemma 1.48** (Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 1.37 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W.\hat{\beta}). &(W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 1.37 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :  
Given
- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :  
Given
- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
 $\forall (a_1, a_2) \in (W.\hat{\beta})$ .
  - $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given
  - $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
Given and from Lemma 1.44
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :  
Given

□

**Theorem 1.49** (Fundamental theorem unary).  $\forall \Gamma, \theta, e, \tau, \delta, n$ .

$$\begin{aligned} \Gamma \vdash e : \tau \wedge \\ (\theta, n, \delta) \in \lfloor \Gamma \rfloor_V \implies \\ (\theta, n, e \delta) \in \lfloor \tau \rfloor_E \end{aligned}$$

*Proof.* Proof by induction on *CG* typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, x \delta) \in [\tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some  $i < n$  s.t  $x \delta \Downarrow_i v$

(from cg-val we know that  $v = x \delta$  and  $i = 0$ )

It suffices to prove  $(\theta, n, x \delta) \in [\tau]_V$  (FU-V0)

Since  $(\theta, n, \delta) \in [\Gamma']_V$  where  $\Gamma' = \Gamma \cup \{x : \tau\}$ . Therefore from Definition 1.40 we know that  $(\theta, n, \delta(x)) \in [\Gamma'(x)]_V$

So we are done.

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e' : \tau_2}{\Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2)]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from cg-val we know that  $v = \lambda x. e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2)]_V \quad (\text{FU-L0})$$

From Definition 1.34 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

This means given some  $\theta'', v', j$  s.t  $\theta'' \sqsupseteq \theta, j < n$  and  $(\theta'', j, v') \in [\tau_1]_V$  (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

Since  $(\theta, n, \delta) \in [\Gamma]_V$  therefore from Lemma 1.45 we know that  $(\theta, j, \delta) \in [\Gamma]_V$  where  $j < n$  (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2]_E, \text{ s.t } (\theta_h, j, v_x) \in [\tau_1]_V$$

Instantiating IH with  $\theta''$  and  $v'$  from (FU-L1) we get  $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$

3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2]_V$$

This means that given some  $i < n$  s.t  $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2)]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2)]_V$$

From cg-app we know that  $v_1 = \lambda x. e'$ . Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2)]_V \quad (\text{FU-P1})$$

This means from Definition 1.34 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2]_E \quad (58)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1]_V \quad (\text{FU-P2})$$

Instantiating Equation 58 with  $\theta, (n - j - k), v_2$  and since we know that  $(\theta, n - j - k, v_2) \in [\tau_1]_V$  therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in [\tau_2]_E$$

This means from Definition 1.35 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_f) \in [\tau_2]_E$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore we know that  $\exists J < i < n$  s.t  $i = j + k + J$  (since  $j + k + J < n$  therefore  $J < n - j - k$ ) and  $e'[v_2/x] \Downarrow_J v_f$

Therefore we have  $(\theta, n - j - k - J, v_f) \in [\tau_2]_E$

Since we know that  $i = j + k + J$  and  $v = v_f$  therefore we get  $(\theta, n - i, v_f) \in [\tau_2]_E$  (so FU-P0 is proved)



4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in [(\tau_1 \times \tau_2)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \times \tau_2)]_V$$

This means that given some  $i < n$  s.t.  $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [(\tau_1 \times \tau_2)]_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [\tau_1]_V$$

Since we know that  $(e_1, e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in [\tau_1]_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_2]_V$$

Since we know that  $(e_1, e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t.  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in [\tau_2]_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from cg-prod we know that  $i = j + k + 1$  and  $v = (v_1, v_2)$  therefore from Definition 1.34 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in [\tau_1]_V \text{ and } (\theta, n - j - k - 1, v_2) \in [\tau_2]_V$$

We get this from (FU-PA1) and Lemma 1.43 and from (FU-PA2) and Lemma 1.43

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \text{fst}(e') \delta) \in [\tau_1]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_1]_V$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor \tau_1 \rfloor_V \quad (\text{FU-F0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$$

Since we know that  $\text{fst}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j (v_1, v_2)$ . This means we have

$$(\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$$

From Definition 1.34 we know the following holds

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V \text{ and } (\theta, n - j, v_2) \in \lfloor \tau_2 \rfloor_V \quad (\text{FU-F1})$$

From cg-fst we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-F0), we are required to prove

$$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \rfloor_V$$

We get this from (FU-F1) and Lemma 1.43

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove:  $(\theta, n, \text{inl}(e') \delta) \in \lfloor (\tau_1 + \tau_2) \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \rfloor_V$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \rfloor_V \quad (\text{FU-LE0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V$$

Since we know that  $\text{inl}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V \quad (\text{FU-LE1})$$

From cg-inl we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-LE0) w we are required to prove

$$(\theta, n - j - 1, v_1) \in [(\tau_1 + \tau_2)]_V$$

From Definition 1.34 it suffices to prove

$$(\theta, n - j - 1, v_1) \in [\tau_1]_V$$

We get this from (FU-LE1) and Lemma 1.43

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some  $i < n$  s.t  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau]_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in [(\tau_1 + \tau_2)]_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_c \delta \Downarrow_j v_c$ . This means we have

$$(\theta, n - j, v_c) \in [(\tau_1 + \tau_2)]_V \quad (\text{FU-C1})$$

2 cases arise:

(a)  $v_c = \text{inl}(v_l)$ :

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in [\tau]_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$ . This means we have

$$(\theta, n - j - k, v_1) \in [\tau]_V \quad (\text{FU-C2})$$

From cg-case1 we know that  $i = j + k + 1$  and  $v = v_1$ . Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in [\tau]_V$$

We get this from (FU-C2) and Lemma 1.43

(b)  $v_c = \text{inr}(v_r)$ :

Symmetric reasoning as in the previous case

10. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } (e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V$$

This means that given some  $i < n$  s.t  $\text{new } (e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{new } (e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

Also from cg-ref we know that  $v' = a$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in [(\text{ref } \ell' \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell' \tau)]_E$$

From Definition 1.35 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in [(\text{Labeled } \ell' \tau)]_V$$

Since we know that  $(H, \text{new } (e')) \Downarrow_j^f (H', a)$  therefore from cg-ref we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h$$

Therefore we have

$$(\theta_e, n - l, v_h) \in [(\text{Labeled } \ell' \tau)]_V \quad (\text{FU-R2})$$

In order to prove (FU-R0) we choose  $\theta'$  as  $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_n$ :

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in [\theta_n(a)]_V$$

- $dom(\theta_n) \subseteq dom(H')$ :  
 We know that  $dom(H') = dom(H) \cup \{a\}$   
 We know that  $dom(\theta_n) = dom(\theta_e) \cup \{a\}$   
 And  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that  $dom(\theta_e) \subseteq dom(H)$   
 So we are done
  - $\forall a \in dom(\theta_n).(\theta_n, (k-j)-1, H'(a)) \in [\theta_n(a)]_V$ :  
 Since from (FU-R2) we know that  $(\theta_h, n-l, v_h) \in [(\text{Labeled } \ell' \tau)]_V$   
 Since  $\theta_h \sqsubseteq \theta_n$  and  $k-j-1 < n-l$  (since  $k < n$  and  $l < j$ ) therefore from Lemma 1.43 we know that  $(\theta_n, k-j-1, v_h) \in [(\text{Labeled } \ell' \tau)]_V$
- (b)  $(\theta_n, k-j-1, a) \in [(\text{ref } \ell' \tau)]_V$ :  
 From Definition 1.34 it suffices to prove that  $\theta_n(a) = \text{Labeled } \ell' \tau$   
 We get this by construction of  $\theta_n$
- (c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :  
 Holds vacuously
- (d)  $(\forall a \in dom(\theta_n) \setminus dom(\theta_e). \theta_n(a) \searrow \ell)$ :  
 From CG-ref we know that  $\ell \sqsubseteq \ell'$

11. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \tau)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, (!e') \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (!e') \delta \Downarrow_i v \implies (\theta, n-i, v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_V$$

(From cg-val we know that  $v = !e' \delta$  and  $i = 0$ )

This means that given some  $i < n$  s.t  $!e' \delta \Downarrow_i !e' \delta$

It suffices to prove

$$(\theta, n, !e' \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-D0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau)]_E$$

From Definition 1.35 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau)]_V$$

Since we know that  $(H, !(e')) \Downarrow_j^f (H', a)$  therefore from cg-deref we know that

$$\exists l < j < k \text{ s.t. } e' \delta \Downarrow_l v_h, v_h = a$$

Therefore we have

$$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau)]_V \quad (\text{FU-D1})$$

In order to prove (FU-D0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

And  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

And since  $H' = H$  (from cg-deref) so we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since  $H' = H$  and from Lemma 1.43 we get

$$\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b)  $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau)]_V$ :

From cg-deref we know that  $H = H'$  and  $v' = H(a)$

From (FU-D1) and Definition 1.34 we know that  $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from cg-deref we know that  $j \geq 1$ . Therefore from Lemma 1.43 we get  $(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau)]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \top)$ :

Holds vacuously

12. CG-assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [(\mathbb{C} \ell \perp \text{unit})]_E^{pc}$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

This means that given some  $i < n$  s.t  $(e_1 := e_2) \delta \Downarrow_i v$ .

It suffices to prove

$$(\theta, n - i, ()) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-assgn we know that  $v' = ()$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, ()) \in [\text{unit}]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0}) \end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists l < j < k$  s.t  $e_1 \delta \Downarrow_l a$ . This means we have

$$(\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k - l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k - l - m, v_2) \in [\text{Labeled } \ell' \tau]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists m < j - l$  (since  $j < k$  therefore  $j - l < k - l$ ) s.t  $e_2 \delta \Downarrow_m v_2$ . This means we have

$$(\theta, k - l - m, v_2) \in [(\text{Labeled } \ell' \tau)]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

•  $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H)$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$ :  
 $\forall a \in \text{dom}(\theta_e)$ .

i.  $H(a) = H'(a)$ :

Since  $(k, H) \triangleright \theta_e$  therefore from Definition 1.36 we know that

$$(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Therefore from Lemma 1.43 we get

$$(\theta_e, k - 1 - j, H(a)) \in [\theta_e(a)]_V$$

ii.  $H(a) \neq H'(a)$ :

From cg-assign we know that  $H'(a) = v_2$

From (FU-A1) we know that  $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that  $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k - l - m, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

Therefore we get

$$(\theta, k - j + 1, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

Therefore from Lemma 1.43 we get

$$(\theta, k - j - 1, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

(b)  $(\theta_e, k - j - 1, ()) \in [\text{unit}]_V$ :

From Definition 1.34

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$ :

From CG-assign we know that  $\ell \sqsubseteq \ell'$

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

13. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \text{Lb}(e') \delta) \in [\text{Labeled } \ell \tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{Labeled } \ell \tau]_V$$

This means we are given some  $i < n$  s.t  $\text{Lb}(e') \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\text{Labeled } \ell \tau]_V$$

Let  $v = \text{Lb}(v_i)$ . This means from Definition 1.34 we are required to prove

$$(\theta, n - i, v_i) \in [\tau]_V$$

IH:  $(\theta, n, e' \delta) \in [\tau]_E$

This means from Definition 1.35 we have

$$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau]_V$$

Since we know that  $\text{Lb}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_i$



Therefore we have  $(\theta, n - j, v_i) \in \lfloor \tau \rfloor_V$

From cg-label we know that  $i = j + 1$  therefore from Lemma 1.43 we have

$$(\theta, n - i, v_i) \in \lfloor \tau \rfloor_V$$

14. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove:  $(\theta, n, \text{unlabel}(e') \delta) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_V$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{unlabel}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-unlabel we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-U0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in \lfloor (\text{Labeled } \ell \tau) \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in \lfloor (\text{Labeled } \ell \tau) \rfloor_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from cg-unlabel we know that

$$\exists h_1 < j < k \text{ s.t } e' \delta \Downarrow_{h_1} \text{Lb } v'$$

This means we have

$$(\theta_e, k - h_1, \text{Lb } v') \in \lfloor (\text{Labeled } \ell \tau) \rfloor_V$$

This means from Definition 1.34 we have

$$(\theta_e, k - h_1, v') \in [\tau]_V \quad (\text{FU-U1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since we have  $(k, H) \triangleright \theta_e$  therefore from Lemma 1.47 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau]_V$ :

Since from (FU-U1) we know that  $(\theta_e, k - h_1, v') \in [\tau]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 1.43 we get  $(\theta_e, k - j, v') \in [\tau]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)$ :

Holds vacuously

15. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell \ell' \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau]_V$$

This means we are given some  $i < n$  s.t.  $\text{ret}(e') \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau]_V$$

(from cg-val we know that  $v = \text{ret}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

Also from cg-ret we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in \lfloor \tau \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in \lfloor \tau \rfloor_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from cg-ret we know that  $\exists h_1 < j < k$  s.t  $e' \delta \Downarrow_{h_1} v'$

This means we have

$$(\theta_e, k - h_1, v') \in \lfloor \tau \rfloor_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 1.47 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in \lfloor \tau \rfloor_V$ :

Since from (FU-R1) we know that  $(\theta_e, k - h_1, v') \in \lfloor \tau \rfloor_V$

And since  $j = h_1 + 1$ , therefore from Lemma 1.43 we get  $(\theta_e, k - j, v') \in \lfloor \tau \rfloor_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :

Holds vacuously

16. CG-bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau'}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove:  $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

This means we are given some  $i < n$  s.t  $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

(from cg-val we know that  $v = \text{bind}(e_1, x.e_2) \delta$  and  $i = 0$ )

Therefore we need to prove

$$(\theta, n, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

From Definition 1.34 it suffices to prove

$$\forall k \leq n. \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore from cg-bind we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_V$$

From Definition 1.34 we know that

$$\begin{aligned} & \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h_1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ & \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1, \theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 1.47 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} & \exists \theta'' \sqsupseteq \theta_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-B1}) \end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_2 < k - h_1 - J. e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$  therefore from cg-bind we know that

$$\exists h_2 < j - h_1 - J < k - h_1 - J \text{ s.t } e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_V$$

From Definition 1.34 we know that

$$\begin{aligned} \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\ \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau']_V \wedge \\ (\forall a. H(a) \neq H''(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell_3) \end{aligned}$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists v_{h_2}, i$  s.t  $(v'' \Downarrow_i v_{h_2})$ . From cg-val we know that  $v_{h_2} = v''$  and  $i = 0$ . Instantiating  $k_{h_2}$  with  $k - h_1 - J - h_2$ ,  $\theta'_e$  with  $\theta''$ ,  $H$  with  $H'$  (from FU-B1) and  $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$  s.t  $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$ . And since we already know that  $(k - h_1, H') \triangleright \theta''$  therefore from Lemma 1.47 we get  $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned} \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H''(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell_3) \quad (\text{FU-B2}) \end{aligned}$$

We get (FU-B0) by choosing  $\theta'$  as  $\theta'''$  (from FU-B2)

17. CG-toLabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma]_V$

To prove:  $(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{toLabeled}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

And given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-tolabeled we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \quad (\text{FU-TL0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$$

Since  $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$  therefore from cg-tolabeled we know that  $\exists h_1 < j < k$  s.t  $e' \delta \Downarrow_{h_1} v_1$

Therefore we get  $(\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

From Definition 1.34 we know that

$$\begin{aligned} \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J.(k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in [\tau]_V \wedge \\ (\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1$ ,  $H_h$  with  $H$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{toLabeled}(e')) \Downarrow_j^f (H', v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 1.47 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} \exists \theta'' \sqsupseteq \theta'_e.(k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-TL1}) \end{aligned}$$

In order to prove (FU-TL0) we choose  $\theta'$  as  $\theta''$ . Now we need to prove the following

(a)  $(k - j, H') \triangleright \theta''$ :

Since  $(k - h_1 - J, H') \triangleright \theta''$  and  $j = h_1 + J + 1$  therefore from Lemma 1.47 we get  $(k - j, H') \triangleright \theta''$

(b)  $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau)]_V$ :

From cg-tolabeled we know that  $v' = \text{toLabeled}(v_1)$

From Definition 1.32 it suffices to prove that  $(\theta'', k - j - 1, v_1) \in [\tau]_V$

We get this from (FU-TL1) and Lemma 1.43

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Directly from (FU-TL1)

(d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :

Directly from (FU-TL1)

□

**Lemma 1.50** (Subtyping unary). *The following holds:*

$$\forall \mathcal{L}, \tau, \tau'.$$

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V \subseteq [(\tau')]_V$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E \subseteq [(\tau')]_E$$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $\llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

IH1:  $\llbracket (\tau'_1) \rrbracket_V \subseteq \llbracket (\tau_1) \rrbracket_V$  (Statement (1))

$\llbracket (\tau_2) \rrbracket_E \subseteq \llbracket (\tau'_2) \rrbracket_E$  (Sub-A0, From Statement (2))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E \quad (59)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

Again from Definition 1.34, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$

Since  $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in \llbracket \tau_1 \rrbracket_V$

As a result from Equation 59 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$$

2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V$

IH1:  $\llbracket (\tau_1) \rrbracket_V \subseteq \llbracket (\tau'_1) \rrbracket_V$  (Statement (1))

IH2:  $\llbracket (\tau_2) \rrbracket_V \subseteq \llbracket (\tau'_2) \rrbracket_V$  (Statement (1))

It suffices to prove:  $\forall (\theta, n, (v_1, v_2)) \in \llbracket ((\tau_1 \times \tau_2)) \rrbracket_V. (\theta, n, (v_1, v_2)) \in \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [((\tau_1 \times \tau_2))]_V$

Therefore from Definition 1.34 we are given:

$$(\theta, n, v_1) \in [\tau_1]_V \wedge (\theta, n, v_2) \in [\tau_2]_V \quad (60)$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [((\tau'_1 \times \tau'_2))]_V$

Again from Definition 1.34, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1]_V \wedge (\theta, n, v_2) \in [\tau'_2]_V$$

Since from Equation 60 we know that  $(\theta, n, v_1) \in [\tau_1]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2]_V$  from Equation 60 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2]_V$

### 3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $[((\tau_1 + \tau_2))]_V \subseteq [((\tau'_1 + \tau'_2))]_V$

IH1:  $[(\tau_1)]_V \subseteq [(\tau'_1)]_V$  (Statement (1))

IH2:  $[(\tau_2)]_V \subseteq [(\tau'_2)]_V$  (Statement (1))

It suffices to prove:  $\forall(\theta, n, v_s) \in [((\tau_1 + \tau_2))]_V. (\theta, v_s) \in [((\tau'_1 + \tau'_2))]_V$

This means that given:  $(\theta, n, v_s) \in [((\tau_1 + \tau_2))]_V$

And it suffices to prove:  $(\theta, n, v_s) \in [((\tau'_1 + \tau'_2))]_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 1.34 we are given:

$$(\theta, n, v_i) \in [\tau_1]_V \quad (61)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_1]_V$$

From Equation 61 and IH1 we know that

$$(\theta, n, v_i) \in [\tau'_1]_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 1.34 we are given:

$$(\theta, n, v_i) \in [\tau_2]_V \quad (62)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_2]_V$$

From Equation 62 and IH2 we know that

$$(\theta, n, v_i) \in [\tau'_2]_V$$



4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V \subseteq \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$

IH:  $\llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$  (Statement (1))

It suffices to prove:

$\forall (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V. (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$

This means that given some  $\theta, n$  and  $\text{Lb}(e_i)$  s.t  $(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$(\theta, n, v_i) \in \llbracket (\tau) \rrbracket_V \quad (\text{SL})$$

And we are required to prove that

$$(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$$

From Definition 1.34 it suffices to prove

$$(\theta, n, v_i) \in \llbracket (\tau') \rrbracket_V$$

We get this directly from (SL) and IH

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $\llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V \subseteq \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$

IH:  $\llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$  (Statement (1))

It suffices to prove:

$\forall (\theta, n, e) \in \llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V. (\theta, n, e) \in \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$

This means that given some  $\theta, n$  and  $e$  s.t  $(\theta, n, e) \in \llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) & \quad (\text{SC0}) \end{aligned}$$

And we are required to prove

$$(\theta, n, e) \in \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$$

So again from Definition 1.34 we need to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau' \rrbracket_V \wedge & \end{aligned}$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$   
(SC1)

And we need to prove

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

We instantiate (SC0) with  $k, \theta_e, H, j$  from (SC1) and we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)$$

Since  $\tau <: \tau'$  therefore from IH we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V$$

And since  $\ell'_i \sqsubseteq \ell_i$  therefore we also have

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

## 6. CGsub-base:

Trivial

### Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E. (\theta, n, e) \in \lfloor (\tau') \rfloor_E$$

This means that we are given  $(\theta, n, e) \in \lfloor (\tau) \rfloor_E$

From Definition 1.35 it means we have

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau') \rfloor_E$$

From Definition 1.35 we need to prove

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \rfloor_V$$

This further means that given some  $i < n$  s.t  $e \Downarrow_i v$ , it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \rfloor_V$$

Instantiating (Sub-E0) with the given  $i$  we get  $(\theta, n - i, v) \in \lfloor \tau \rfloor_V$

Finally from Statement(1) we get  $(\theta, n - i, v) \in \lfloor \tau' \rfloor_V$

□

**Lemma 1.51** (Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in \lfloor \Gamma \rfloor_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \Downarrow_i) \in \lfloor \Gamma \rfloor_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 1.41 we know that we are given:

$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

And we are required to prove:

$\forall i \in \{1, 2\}. \forall m.$

$dom(\Gamma) \subseteq dom(\gamma \downarrow_i) \wedge \forall x \in dom(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 1$

Given some  $m$  we need to show:

- $dom(\Gamma) \subseteq dom(\gamma \downarrow_1)$ :

$$dom(\gamma) = dom(\gamma \downarrow_1)$$

Therefore,  $dom(\Gamma) \subseteq (dom(\gamma) = dom(\gamma \downarrow_1))$  (Given)

- $\forall x \in dom(\Gamma). (W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$ :

We are given:  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

Therefore from Lemma 1.42 we know that

$$\forall m'. (W.\theta_1, m', \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Instantiating  $m'$  with  $m$  we get

$$(W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Case  $i = 2$

Symmetric reasoning as in the  $i = 1$  case above

□

**Theorem 1.52** (Fundamental theorem binary).  $\forall \Gamma, pc, W, \mathcal{A}, e, \tau, \gamma, n.$

$\Gamma \vdash e : \tau \wedge$

$(W, n, \gamma) \in [\Gamma]_V^A \implies$

$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau]_E^A$

*Proof.* Proof by induction on the typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

To prove:  $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau]_E^A$

Say  $e_1 = x(\gamma \downarrow_1)$  and  $e_2 = x(\gamma \downarrow_2)$

From Definition 1.33 it suffices to prove that

$$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some  $i < n$  s.t  $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2$

We are required to prove:  $(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

From cg-val we know that  $x (\gamma \downarrow_1) \Downarrow x (\gamma \downarrow_1)$  and  $x (\gamma \downarrow_2) \Downarrow x (\gamma \downarrow_2)$

This means  $v'_1 = x (\gamma \downarrow_1)$  and  $v'_2 = x (\gamma \downarrow_2)$

Since  $(W, n, \gamma) \in [\tau]_V^A$ . Therefore from Definition 1.41 we know that

$(W, n, v'_1, v'_2) \in [\tau]_V^A$

From Lemma 1.44 we get

$(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove:  $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2)]_E^A$

Say  $e_1 = \lambda x. e (\gamma \downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \rightarrow \tau_2)]_E^A$  it suffices to prove that

$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$

This means given some  $i < n$  s.t  $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From cg-val we know that  $v'_1 = (\lambda x. e_i) \gamma \downarrow_1$  and  $v'_2 = (\lambda x. e_i) \gamma \downarrow_2$

We are required to prove:

$(W, n - i, (\lambda x. e_i) \gamma \downarrow_1, (\lambda x. e_i) \gamma \downarrow_2) \in [\tau]_V^A$

From Definition 1.32 it suffices to prove

$\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in [\tau_2]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E) \wedge$

$\forall \theta_l \sqsupseteq W. \theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2]_E) \quad \text{(FB-L0)}$

IH:

$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2]_E^A$

s.t

$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in [\Gamma]_V^A$

In order to prove (FB-L0) we need to prove the following:

- (a)  $\forall W' \sqsupseteq W, j < n, v_1, v_2.$   
 $((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2]_E^A):$

This means given some  $W' \sqsupseteq W, j < n, v_1, v_2$  s.t.  $(W', j, v_1, v_2) \in [\tau_1]_V^A$   
We need to prove  $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2]_E^A$

We get this by instantiating IH with  $W'$  and  $j$

- (b)  $\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$   
 $((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E):$   
This means given some  $\theta_l \sqsupseteq W.\theta_1, v_c, j$  s.t.  $(\theta_l, j, v_c) \in [\tau_1]_V$   
We need to prove:  $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E$

It is given to us that

$$(W, n, \gamma) \in [\Gamma]_V^A$$

Therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating  $m$  with  $j$  we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

From Lemma 1.46 we know that

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Since we know that  $(\theta_l, j, v_c) \in [\tau_1]_V$

Therefore we also have

$$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1\}]_V$$

Therefore, we can apply Theorem 1.49 to obtain

$$(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_V$$

- (c)  $\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$   
 $((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2]_E):$   
Similar reasoning as in the previous case

### 3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2]_V^A$$

This further means that given some  $i < n$  s.t.  $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2]_V^A$$

IH1:  $(W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2)]_E^A$

This means from Definition 1.33 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h_1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h_2} \implies (W, n - j, v_{h_1}, v_{h_2}) \in [(\tau_1 \rightarrow \tau_2)]_V^A$$

Since we know that  $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \downarrow_j v_{h_1}$ . Similarly since  $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_1 \gamma \downarrow_2 \downarrow v_{h_2}$

This means we have  $(W, n - j, v_{h_1}, v_{h_2}) \in [(\tau_1 \rightarrow \tau_2)]_V^A$

From cg-app we know that  $val_{h_1} = \lambda x. e_{h_1}$  and  $val_{h_2} = \lambda x. e_{h_2}$

From Definition 1.32 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, J < (n - j), v_1, v_2. \\ ((W', J, v_1, v_2) \in [\tau_1]_V^A \implies (W', J, e_{h_1}[v_1/x], e_{h_2}[v_2/x]) \in [\tau_2]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, v_c, j. \\ ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, v_c, j. \\ ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E) \end{aligned} \quad (\text{FB-A1})$$

$$\underline{\text{IH2}}: (W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.33 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \downarrow_j v_{h_1'} \wedge e_2 \gamma \downarrow_2 \downarrow v_{h_2'} \implies (W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1]_V^A$$

Since we know that  $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \downarrow_k v_{h_1'}$ . Similarly since  $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \downarrow v_{h_2'}$

This means we have  $(W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1]_V^A$  (FB-A2)

Instantiating the first conjunct of (FB-A1) as follows  $W'$  with  $W$ ,  $J$  with  $n - j - k$ ,  $v_1$  and  $v_2$  with  $v'_{h_1}$  and  $v'_{h_2}$  respectively, we obtain

$$(W, n - j - k, e_{h_1}[v'_{h_1}/x], e_{h_2}[v'_{h_2}/x]) \in [\tau_2]_E^A$$

From Definition 1.33

$$\forall l < n - j - k. (e_{h_1}[v'_{h_1}/x]) \gamma \downarrow_l v_{f_1} \wedge e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2} \implies (W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2]_V^A$$

Since we know that  $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$ . Therefore  $\exists l < i - j - k < n - j - k$  s.t  $e_{h_1}[v'_{h_1}/x] \downarrow_l v_{f_1}$ . Similarly since  $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2}$

Therefore we have  $(W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2]_V^A$

Since  $i = j + k + l$  therefore we are done

#### 4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

To prove:  $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. (e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2}) \implies \\ (W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2})$

We are required to prove

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A \quad (\text{FB-P0})$$

$$\underline{\text{IH1}}: (W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.33 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge e_1 \gamma \downarrow_2 \downarrow v'_{f_1} \implies (W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \downarrow_j v_{f_1}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_1 \gamma \downarrow_2 \downarrow v'_{f_1}$

This means we have

$$(W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A \quad (\text{FB-P1})$$

$$\underline{\text{IH2}}: (W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2]_E^A$$

This means from Definition 1.33 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \downarrow_j v_{f_2}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 1.32 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A$$

Since  $i = j + k + 1$  therefore from (FB-P1) and (FB-P2) and from Lemma 1.44 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove:  $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [(\tau_1)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1]_V^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies \\ (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

Since we know that  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$  therefore  $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

From Definition 1.32 we know that

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

Since from cg-fst  $i = j + 1$  therefore from Lemma 1.44 we get

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove:  $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1}) \wedge \text{inl}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1}) \implies \\ (W, n - i, \text{inl}(v_{f_1}), \text{inl}(v'_{f_1})) \in [(\tau_1 + \tau_2)]_V^A$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1}) \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f_1}), \text{inl}(v'_{f_1})) \in [(\tau_1 + \tau_2)]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f_1} \implies \\ (W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

Since we know that  $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j v_{f_1}$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1})$  therefore  $e' \gamma \downarrow_2 \Downarrow v'_{f_1}$

This means we have



$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 1.32 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1]_V^A$$

From cg-inl since  $i = j + 1$  therefore from (FB-IL1) and Lemma 1.44 we get (FB-IL0)

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove:  $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau]_V^A$$

This means that given some  $i < n$  s.t  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2}$

We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau]_V^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_j v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2)]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$  therefore  $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2)]_V^A \quad (\text{FB-C1})$$

2 cases arise

(a)  $v_{h1} = \text{inl}(v_1)$  and  $v'_{h1} = \text{inl}(v'_1)$ :

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_i v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau]_V^A$$

From cg-case1 we know that  $i = j + k + 1$  therefore from Lemma 1.44 we get (FB-C0)

(b)  $v_{h1} = \text{inr}(v_1)$  and  $v'_{h1} = \text{inr}(v'_1)$ :

Symmetric case

10. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove:  $(W, n, \text{Lb}(e') (\gamma \downarrow_1), \text{Lb}(e') (\gamma \downarrow_2)) \in [\text{Labeled } \ell \tau]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1}) \implies (W, n - i, \text{Lb}(v_{f1}), \text{Lb}(v'_{f1})) \in [\text{Labeled } \ell \tau]_V^A$$

This means that given some  $i < n$  s.t  $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1})$

We are required to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\text{Labeled } \ell \tau]_V^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, v_{f1}, v'_{f1}) \in [\tau]_V^A$$

Since we know that  $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $\text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1})$  therefore  $e' \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau]_V^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 1.32 it suffices to prove that

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau]_V^A$$

From cg-label we know that  $i = j + 1$ . Therefore we get the desired from (FB-LB1) and Lemma 1.44

11. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

To prove:  $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \top \ell \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \top \ell \tau)]_V^A$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{unlabel}(e') \gamma \downarrow_1$  and  $v'_{f1} = \text{unlabel}(e') \gamma \downarrow_2$ . Also  $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\mathbb{C} \top \ell \tau)]_V^A$$

This means from Definition 1.32 we need to prove

Let  $e_1 = \text{unlabel}(e') \gamma \downarrow_1$  and  $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right) \end{aligned}$$

We need to show

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau): \end{aligned}$$

Also given is some  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau) \quad (\text{FB-U0})$$

$$\underline{\text{IH}}: (W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{Labeled } \ell \tau)]_E^A$$

This means from Definition 1.33 we are given

$$\forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) \implies (W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau)]_V^A$$

Since we know that

$(H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k$  therefore  
 $\exists I < j < k$  s.t  $e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1})$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau)]_V^A$$

This means from Definition 1.32 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose  $W'$  as  $W_e$  and from cg-unlabel we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . And we already know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore from Lemma 1.48 we get  $(k - j, H_1, H_2) \triangleright W_e$

From cg-unlabel we know that  $v'_1, v'_2$  in (FB-U0) is  $v_{h1}, v'_{h1}$  respectively. And since from (FB-U1) we know that  $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau)$ . Therefore from Lemma 1.53 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, \ell, v_{h1}, v'_{h1}, \tau)$$

$$(b) \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right):$$

Case  $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t } (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k$$

We need to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

$$\text{Instantiating } m \text{ with } k \text{ we get } (W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (\text{unlabel } e') \gamma \downarrow_1) \in [(\mathbb{C} \top \ell \tau)]_E$$

This means from Definition 1.35 we get

$$\forall c < k. (\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \top \ell \tau)]_V$$

This further means that given some  $c < k$  s.t  $(\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v$ . From cg-val we know that  $c = 0$  and  $v = (\text{unlabel } e') \gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, (\text{unlabel } e') \gamma \downarrow_1) \in [(\mathbb{C} \top \ell \tau)]_V$$

From Definition 1.34 we have

$$\begin{aligned} & \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e') \gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies \\ & \exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \top) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

12. CG-tolabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)}$$

To prove:  $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} &\implies \\ (W, n - i, v_{f_1}, v'_{f_1}) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

From cg-val we know that  $v_{f_1} = \text{toLabeled}(e') \gamma \downarrow_1$ ,  $v_{f_2} = \text{toLabeled}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_V^A$$

Let  $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$  and  $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 1.32 we are required to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau))) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled} \ell_o \tau)]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

We need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau)): \end{aligned}$$

This means that we are given some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2, v'_1, v'_2, j < k$  s.t

$$(k, H_1, H_2) \triangleright W_e \text{ and } (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau))$$

From Definition 1.31 it suffices to prove that

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled} \ell_2 \tau)]_V^A$$

Further from Definition 1.32 it suffices to prove

$$\begin{aligned} \exists W' \sqsupseteq W_e.(k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-TL0}) \\ \text{where } v''_1 = \text{Lb } v''_1 \text{ and } v''_2 = \text{Lb } v''_2 \end{aligned}$$

IH:

$$(W_e, k, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v'_{h1} \implies (W_e, n-J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j (H'_1, v'_1)$  and  $(H_2, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j (H'_2, v'_2)$ . Therefore from cg-val we know that  $\exists J < j < k \leq n$  s.t  $e' \gamma \downarrow_1 \downarrow_J v_{h1}$  and similarly we also know that  $e' \gamma \downarrow_2 \downarrow v'_{h1}$

This means we have

$$(W_e, k-J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

From Definition 1.32 we know that

$$\begin{aligned} (\forall k_1 \leq (k-J), W_e'' \sqsupseteq W_e. \forall H''_1, H''_2. (k_1, H''_1, H''_2) \triangleright W_e'' \wedge \forall v''_1, v''_2, m. \\ (H''_1, v_{h1}) \downarrow_m^f (H''_1, v''_1) \wedge (H''_2, v'_{h1}) \downarrow^f (H''_2, v''_2) \wedge m < k_1 \implies \\ \exists W' \sqsupseteq W_e''. (k_1 - m, H''_1, H''_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1 - m, \ell_2, v''_1, v''_2, \tau)) \wedge \\ \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_l) \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{FB-TL1}) \end{aligned}$$

We instantiate  $W_e''$  with  $W_e$ ,  $H''_1$  with  $H_1$ ,  $H''_2$  with  $H_2$  and  $k_1$  with  $k$  in (FB-TL1).

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \downarrow_j^f (H'_2, v'_2)$ , therefore  $\exists m < j < k \leq n$  s.t  $(H_1, v_{h1}) \downarrow_m^f (H'_1, v'_1) \wedge (H_2, v'_{h1}) \downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W' \sqsupseteq W_e.(k-m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-m, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-TL2})$$

In order to prove (FB-TL0) we choose  $W'$  as  $W'$  from (FB-TL2). Since from cg-tolabeled we know that  $v''_1 = \text{Lb}(v''_1)$ ,  $v''_2 = \text{Lb}(v''_2)$  and  $j = m+1$  (therefore from Lemma 1.48 we get  $(k-j, H'_1, H'_2) \triangleright W'$ ) and from (FB-TL2) and Lemma 1.53 we get  $\text{ValEq}(\mathcal{A}, W', k-j, \ell_2, v''_1, v''_2, \tau)$

$$\begin{aligned} \text{(b) } \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_l) \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))): \end{aligned}$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_l) \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that  
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_E$$

This means from Definition 1.35 we get

$$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

Instantiating  $c$  with 0 and from cg-val we know  $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$

From Definition 1.34 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_j^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\text{Labeled } \ell_2 \tau]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

13. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_1 \ell_2 \tau}$$

To prove:  $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. \text{ret}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{ret}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{ret}(e')\gamma \downarrow_1$ ,  $v_{f2} = \text{ret}(e')\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{ret}(e')\gamma \downarrow_1, \text{ret}(e')\gamma \downarrow_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

Let  $v_1 = \text{ret}(e')\gamma \downarrow_1$  and  $v_2 = \text{ret}(e')\gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \end{aligned}$$

$$\begin{aligned}
& \forall l \in \{1, 2\}. \left( \forall v, i. (e_l \Downarrow_i v_l) \implies \right. \\
& \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
& \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right)
\end{aligned}$$

It suffices to prove:

$$\begin{aligned}
\text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\
& (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau):
\end{aligned}$$

We are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t.  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

From cg-ret we know that  $H'_1 = H_1$  and  $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau]_V^A$$

Since we know that  $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow_j^f (H_2, v'_2)$ , therefore  $\exists J < j < k$  s.t.  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly  $e' \gamma \downarrow_2 \Downarrow v'_{h1}$ .

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau]_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_e$  and from cg-ret we know that  $v'_1 = v_{h1}$  and  $v'_2 = v'_{h1}$ . We need to prove the following:

i.  $(k - j, H_1, H_2) \triangleright W_e$ :

Since we have  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 1.48 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_2, v'_1, v'_2, \tau)$ :

2 cases arise:

A.  $\ell_2 \sqsubseteq \mathcal{A}$ :

In this case from Definition 1.31 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in [\tau]_V^A$$

Since  $j = J + 1$  therefore we get this from (FB-R1) and Lemma 1.44

B.  $\ell_2 \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 1.31 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in [\tau]_V \text{ and } \forall m. (W_e, m, v'_2) \in [\tau]_V$$

We get this From (FB-R1) and Lemma 1.42



(b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_o \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_o)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E$

This means from Definition 1.35 we get

$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

Instantiating  $c$  with 0 and from cg-val we know that  $v = (\text{ret } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

From Definition 1.34 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$   
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

14. CG-bind:

$$\frac{\Gamma \vdash e_l : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_b : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_l, x.e_b) : \mathbb{C} \ell \ell' \tau'}$$

To prove:  $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \ell' \tau']_E^A$

This means from Definition 1.33 we need to prove:

$\forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} \implies$   
 $(W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \ell' \tau']_V^A$

This means that given some  $i < n$  s.t.  $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f_1} = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$ ,  $v_{f_2} = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b)\gamma \downarrow_1, \text{bind}(e_l, x.e_b)\gamma \downarrow_2) \in [\mathbb{C} \ell \ell' \tau']_V^A$$

Let  $v_1 = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$  and  $v_2 = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$

This means from Definition 1.32 we need to prove

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau): \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} & \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} \implies \\ & (W_e, k - f, v_{h_1}, v'_{h_1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

This means from Definition 1.32 we have

$$\begin{aligned} & (\forall K \leq (k - f), W_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W_e \wedge \forall v''_1, v''_2, J. \\ & (H''_1, v_{h_1}) \Downarrow_J^f (H''_1, v''_1) \wedge (H''_2, v'_{h_1}) \Downarrow_J^f (H''_2, v''_2) \wedge J < K \implies \\ & \exists W'' \sqsupseteq W_e. (K - J, H''_1, H''_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell_2, v''_1, v''_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \end{aligned}$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)$$

Instantiating  $K$  with  $(k - f)$ ,  $W'_e$  with  $W_e$ ,  $H_1''$  with  $H_1$  and  $H_2''$  with  $H_2$  in the first conjunct of the above equation. Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 1.48 we also have  $(k - f, H_1, H_2) \triangleright W_e$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists J < j - f < k - f$  s.t.  $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v''_1) \wedge (H_2, v'_{h1}) \Downarrow_J^f (H'_2, v''_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e. (k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-B1})$$

From Definition 1.31 two cases arise:

i.  $\ell_2 \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W'', k - f - J, v''_1, v''_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

IH2:

$$(W'', k - f - J, e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}), e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\})) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{E}}^{\mathcal{A}}$$

This means from Definition 1.33 we need to prove:

$$\forall s < k - f - J. e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2} \implies \\ (W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{V}}^{\mathcal{A}}$$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists s < j - f - J < k - f - J$  s.t.  $e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{V}}^{\mathcal{A}}$$

This means from Definition 1.32 we know that

$$\left( \forall K_s \leq (k - f - J - s), W_s \sqsupseteq W'' . \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right.$$

$$(H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow_{J_s}^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies$$

$$\left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_4, v'_1, v'_2, \tau') \right) \wedge$$

$$\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$$

$$\left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_{\mathcal{V}} \wedge \right.$$

$$\left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_3 \sqsubseteq \ell') \wedge \right.$$

$$\left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_3) \right)$$

Instantiating  $K_s$  with  $(k - f - J - s)$ ,  $W_s$  with  $W''$ ,  $H_1$  with  $H'_1$  and  $H'_2$  with  $H_2$ . Since we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Lemma 1.48 we also have  $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists J_s < j - f - J - s < k - f - J - s$  s.t.  $(H'_1, v'_1) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_2) \Downarrow_{J_s}^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s.(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau') \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose  $W'$  as  $W'_s$ . From cg-bind we know that  $H'_1 = H'_{s1}$ ,  $H'_2 = H'_{s2}$ ,  $v'_1 = v'_{s1}$ ,  $v'_2 = v'_{s2}$  and  $j = f + J + s + J_s + 1$ . And we need to prove:

A.  $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$ :

Since from (FB-B2) we know that  $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$  therefore from Lemma 1.48 we get

$$(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$$

B.  $\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau')$ :

Since from (FB-B2) we know that  $\text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau')$  therefore from Lemma 1.53 we get

$$\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau')$$

ii.  $\ell_2 \not\sqsubseteq \mathcal{A}$ :

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau')$$

Since  $\ell_2 \sqsubseteq \ell_4 \sqsubseteq \ell'$  and  $\ell \not\sqsubseteq \mathcal{A}$  therefore we have  $\ell_4 \not\sqsubseteq \mathcal{A}$  and  $\ell' \not\sqsubseteq \mathcal{A}$

This means that from Definition 1.31 it suffices to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'.\theta_1, m_{u1}, v'_1) \in [\tau']_V \wedge \forall m_{u2}. (W'.\theta_2, m_{u2}, v'_2) \in [\tau']_V$$

This means given some  $m_{u1}, m_{u2}$  and we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in [\tau']_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau']_V \quad (\text{FB-B01})$$

In this case from (FB-B1) and Definition 1.31 we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in [\tau]_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in [\tau]_V \quad (\text{FB-B3})$$

Since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_1 \downarrow_j v'_1$  therefore  $\exists J_1 < j - f - J < k - f - J$  s.t  $(e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{J_1} v'_1$ . Similarly,  $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$  s.t  $(H'_1, v'_1) \downarrow_{J'_1}^f -$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  in the first conjunct of (FB-B3)

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in [\tau]_V$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  we get  $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 1.45 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V \quad (\text{FB-B4})$$

Now we can apply Theorem 1.49 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E$$

This means from Definition 1.35 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V \quad (\text{FB-B5})$$

Instantiating  $c_1$  with  $J_1$  in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

From Definition 1.34 we have

$$\forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \Downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies$$

$$\begin{aligned} \exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in [\tau']_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \end{aligned}$$

Instantiating  $K$  with  $m_{u1} + 1 + J'_1$ ,  $\theta'_e$  with  $W''.\theta_1$ ,  $H_1$  with  $H'_1$  (from FB-B1) and  $J_2$  with  $J'_1$  we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1. (m_{u1} + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau']_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \Downarrow v'_2$ . Say this reduction happens in  $t$  steps. Therefore  $\exists t_1 < t < k \leq n$  s.t.  $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{t_1} v_{l2}$  and similarly  $\exists t_2 < t - t_1 < k - t_1$  s.t.  $(H, v_{l2})\gamma \downarrow_2 \Downarrow_{t_2}^f (H''_2, v''_2)$

Again since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \Downarrow v'_2$  therefore  $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$  s.t.  $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{J_2} v'_2$ . Similarly  $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$  s.t.  $(H'_2, v'_2) \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with  $m_{u2} + 1 + J_2 + J'_2$  we get  $(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v''_2) \in [\tau]_V$

Again since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that  $\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $m_{u2} + 1 + J_2 + J'_2$  we get  $(W.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 1.45 we know that

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 1.49 to get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E$$

This means from Definition 1.35 we get

$$\forall c_2 < (m_{u2} + 1 + J_2 + J'_2). (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{c_2} v_{o2} \implies (W''.\theta_2, m_{u2} + 1 + J_2 - c_2, v_{o2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V \quad (\text{FB-B8})$$

Instantiating  $c_2$  with  $J_2$  in (FB-B8) we get

$$(W''.\theta_2, m_{u2} + 1 + J'_2, v_{o2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

From Definition 1.34 we have

$$\forall K \leq (m_{u2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3. (K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o2}) \Downarrow_{J_3}^f (H''_2, v'_2) \wedge J_3 < K \implies$$

$$\begin{aligned} \exists \theta'_2 \sqsupseteq \theta'_e. (K - J_3, H''_2) \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v'_2) \in [\tau']_V \wedge \\ (\forall a. H_2(a) \neq H''_2(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \end{aligned}$$

Instantiating  $K$  with  $m_{u2} + 1 + J'_2$ ,  $\theta'_e$  with  $W''.\theta_2$ ,  $H_2$  with  $H'_2$  (from FB-B1) and  $J_3$  with  $J'_2$ , we get

$$\begin{aligned} & \exists \theta'_2 \sqsupseteq W''.\theta_2.(m_{u2} + 1, H''_2) \triangleright \theta'_2 \wedge (\theta'_2, m_{u2} + 1, v'_2) \in [\tau']_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose  $W'$  as  $W_n$  where  $W_n$  is defined as follows:

$$\begin{aligned} W_n.\theta_1 &= \theta'_1 \quad (\text{From (FB-B6)}) \\ W_n.\theta_2 &= \theta'_2 \quad (\text{From (FB-B9)}) \\ W_n.\hat{\beta} &= W''.\hat{\beta} \quad (\text{From (FB-B1)}) \end{aligned}$$

It suffices to prove

- $(k - j, H'_1, H'_2) \triangleright W_n$ :

From Definition 1.37 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2):$$

From (FB-B6) we know that  $(m_{u1} + 1, H''_1) \triangleright \theta'_1$  therefore from Definition 1.36 we know that  $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly from (FB-B9) we know that  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$  therefore from Definition 1.36 we know that  $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Definition 1.37 we know that  $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq W_n.\theta_1$  and  $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell']_V^A):$$

4 cases arise for each  $(a_1, a_2) \in W_n.\hat{\beta}$

$$A. H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

We know from that  $(k - f - J, H'_1, H'_2) \triangleright W''$

Therefore from Definition 1.37 we have

$$\forall (a'_1, a'_2) \in (W''.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

Since  $W_n.\hat{\beta} = W''.\hat{\beta}$  by construction therefore

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

From (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq \theta'_1$  and  $W''.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.29

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). \theta'_1(a'_1) = \theta'_2(a'_2)$$

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

From (FB-B1) we know that  $(k - f - J, H_1', H_2') \triangleright^A W''$

This means from Definition 1.37 we know that

$$\forall (a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) = W''.\theta_2(a_{i2}) \wedge \\ (W'', k - f - J - 1, H_1'(a_{i1}), H_2'(a_{i2})) \in \lceil W''.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'' \sqsubseteq W_n$  and  $k - j - 1 < k - f - J - 1$  (since  $j = f + J + J_1 + 1$  therefore from Lemma 1.44 we get

$$(W_n, k - j - 1, H_1'(a_1), H_2'(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

B.  $H_1'(a_1) \neq H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$ :

To prove:

$$\overline{W_n.\theta_1(a_1) = W_n.\theta_2(a_2)}$$

Same reasoning as in the previous case

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

From (FB-B6) and (FB-B9) we know that

$$(\forall a. H_1'(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell') \\ (\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell' \text{ and} \\ \exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell'$$

Since  $\ell_2 \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_3 \not\sqsubseteq \mathcal{A}$ .

Also from (FB-B6) and (FB-B9),  $(m_{u1} + 1, H_1'') \triangleright \theta_1'$  and  $(m_{u2} + 1, H_2'') \triangleright \theta_2'$ .

Therefore from Definition 1.36 we have

$$(\theta_1', m_{u1}, H_1''(a_1)) \in \lceil \theta_1'(a_1) \rceil_V \text{ and} \\ (\theta_2', m_{u2}, H_2''(a_1)) \in \lceil \theta_2'(a_2) \rceil_V$$

Since  $m_{u1}$  and  $m_{u2}$  are arbitrary indices therefore from Definition 1.32 we get

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \lceil \theta_1'(a_1) \rceil_V^A$$

C.  $H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$ :

To prove:

$$\overline{W_n.\theta_1(a_1) = W_n.\theta_2(a_2)}$$

Same reasoning as in the previous case

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

From (FB-B9) we know that

$$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell')$$

This means we have

$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell'$   
 Since  $\ell_2 \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_3 \not\sqsubseteq \mathcal{A}$ .

Since from (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W''$  that means from Definition 1.37 that  $(W'', k - f - J - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W''.\theta_1(a_1) \rceil_V^{\mathcal{A}}$ . Since  $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$  and since  $\ell' \not\sqsubseteq \mathcal{A}$  therefore from Definition 1.32 and Definition 1.31 we know that

Therefore  
 $\forall m. (W''.\theta_1, m, H'_1(a_1)) \in W''.\theta_1(a_1)$  (F)

Instantiating the (F) with  $m_{u1}$  and using Lemma 1.43 we get  
 $(\theta'_1, m_{u1}, H'_1(a_1)) \in \theta'_1(a_1)$

Since from (FB-B9) we know that  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$  therefore from Definition 1.36 we know that  $(\theta'_2, m_{u2}, H''_2(a_2)) \in \theta'_2(a_2)$   
 Therefore from Definition 1.32 we get  
 $(W', k - j - 1, H''_1(a_1), H''_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^{\mathcal{A}}$

D.  $H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2)$ :  
 Symmetric reasoning as in the previous case

–  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$ :

Case  $i = 1$

Given some  $m$  we need to prove

$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$

This further means that given some  $a_1 \in \text{dom}(W_n.\theta_1)$  we need to show  
 $(W_n.\theta_1, m, H''_1(a_1)) \in \lfloor W_n.\theta_1(a_1) \rfloor_V$

Since  $W_n.\theta_1 = \theta'_1$ , it suffices to prove  
 $(\theta'_1, m, H''_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$

Like before we apply Theorem 1.49 on  $e_b \gamma \downarrow_1 \cup \{x \mapsto v''_1\}$  but this time at  $m + 1 + J_1 + J'_1$  to get

$\exists \theta'_1 \sqsupseteq W''.\theta_1.(m + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in \lfloor \tau' \rfloor_V \wedge$   
 $(\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3)$

Since we have  $\ell \sqsubseteq \ell_3$  and  $(m + 1, H''_1) \triangleright \theta'_1$  therefore from Definition 1.36 we get the desired.

Case  $i = 2$

Similar reasoning as in the  $i = 1$  case

•  $(W'.\theta_1, m_{u1}, v'_1) \in \lfloor \tau' \rfloor_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in \lfloor \tau' \rfloor_V$ :

We get this from (FB-B6), (FB-B9) and Lemma 1.43 we get the desired

15. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new}(e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

To prove:  $(W, n, \text{new}(e')(\gamma \downarrow_1), \text{new}(e')(\gamma \downarrow_2)) \in \lceil (\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \rceil_E^{\mathcal{A}}$



This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{new}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $\text{new}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{new}(e') \gamma \downarrow_1$ ,  $v_{f2} = \text{new}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{new}(e') \gamma \downarrow_1, \text{new}(e') \gamma \downarrow_2) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_{\mathcal{V}}^A$$

Let  $v_1 = \text{new}(e') \gamma \downarrow_1$  and  $v_2 = \text{new}(e') \gamma \downarrow_2$

From Definition 1.32 we are required to prove

$$\left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ \left. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \right. \\ \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau)) \right) \wedge \\ \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_{\mathcal{V}} \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right)$$

This means we need to prove the following:

$$(a) \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau)):$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also we are given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau))$$

Further from Definition 1.31 it suffices to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau)]_{\mathcal{V}}^A \quad (\text{FB-R0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau]_{\mathcal{V}}^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e' \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau]_{\mathcal{V}}^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_n$  where

$$W_n.\theta_1 = W_e.\theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\theta_2 = W_e.\theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\hat{\beta} = W_e.\hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

i.  $(k - j, H'_1, H'_2) \triangleright W_n$ :

From Definition 1.37 it suffices to prove:

$$\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2)) \wedge$$

$$(W_n, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$$

This means we need to prove

- $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$ :

We know that  $\text{dom}(W_n.\theta_1) = \text{dom}(W_e.\theta_1) \cup \{a_1\}$  and  $\text{dom}(W_n.\theta_2) = \text{dom}(W_e.\theta_2) \cup \{a_2\}$

Also  $\text{dom}(H'_1) = \text{dom}(H_1) \cup \{a_1\}$  and  $\text{dom}(H'_2) = \text{dom}(H_2) \cup \{a_2\}$

Therefore from  $(k, H_1, H_2) \triangleright W_e$  and from construction of  $W_n$  we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2)) \wedge (W_n, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_n.\theta_1(a'_1) \rceil_V^A$ :

$\forall (a'_1, a'_2) \in (W_n.\hat{\beta})$ .

A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

Since from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$

And since from cg-ref we know that  $H'_1(a_1) = v_{h1}$ ,  $H'_2(a_2) = v'_{h1}$  and  $j = f + 1$  thfore from Lemma 1.44 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise

C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise

D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 1.37

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in \lfloor W_n.\theta_i(a'_i) \rfloor_V$ :

When  $i = 1$

Given some  $m$

$\forall a'_1 \in \text{dom}(W_n.\theta_1)$ .

– when  $a'_1 = a_1$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

And from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$

Therefore from Lemma 1.42 get the desired

– Otherwise:

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 1.37

When  $i = 2$

Similar reasoning as with  $i = 1$

ii.  $(W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau)]_V^A$ :

From cg-ref we know that  $v'_1 = a_1$  and  $v'_2 = a_2$

From Definition 1.32 it suffices to prove

$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau)$

This holds from construction of  $W_n$

(b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge$

$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_E$

This means from Definition 1.35 we get

$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$

This further means that given some  $c < k$  s.t  $\text{ref } (e')\gamma \downarrow_1 \Downarrow_c v$ . From cg-val we know that  $c = 0$  and  $v = \text{ref } (e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$

From Definition 1.34 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$

$\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau)]_V \wedge$

$(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

16. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \ \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)}$$

To prove:  $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. !e' \ \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \ \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) &\in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $!e' \ \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \ \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = !e' \ \gamma \downarrow_1$ ,  $v_{f2} = !e' \ \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, !e' \ \gamma \downarrow_1, !e' \ \gamma \downarrow_2) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_V^A$$

Let  $v_1 = !e' \ \gamma \downarrow_1$  and  $v_2 = !e' \ \gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau))) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \ \tau)]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \ \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau)): \end{aligned}$$

This means we are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau))$$

This means from Definition 1.31 it suffices to prove  $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \ \tau)]_V^A$  (FB-D0)

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{ref } \ell \ \tau)]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} \implies \\ (W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A \quad (\text{FB-D1})$$

In order to prove (FB-D0) we choose  $W'$  as  $W_e$ . Also from cg-deref we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . Also we know that  $v_{h_1} = a_1$  and  $v'_{h_1} = a_2$ .

- $(k - j, H_1, H_2) \triangleright W_e$ :

Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 1.48 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

- $(W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \tau)]_V^A$ :

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau]_V^A$

Therefore from Definition 1.32 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that

$$(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau]_V^A$$

Also from cg-ref we know that  $v'_1 = H_1(a_1)$  and  $v'_2 = H_2(a_2)$

From Lemma 1.44 we get  $(W', k - j, H_1(a_1), H_2(a_2)) \in [(\text{Labeled } \ell \tau)]_V^A$

- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (!e' \gamma \downarrow_1)) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_E$$

This means from Definition 1.35 we get

$$\forall c < k. !e' \gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$$

Instantiating  $c$  with 0 and from cg-val we know that  $v = !e' \gamma \downarrow_1$

And we have  $(W.\theta_1, k, !e' \gamma \downarrow_1) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$

From Definition 1.34 we have

$$\begin{aligned} & \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ & \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \top) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

17. CG-assign:

$$\frac{\Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_l := e_r : \mathbb{C} \ell \perp \text{unit}}$$

To prove:  $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \perp \text{unit}]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} & \forall i < n. (e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ & (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \perp \text{unit}]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t.  $(e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = (e_l := e_r) \gamma \downarrow_1$ ,  $v_{f2} = (e_l := e_r) \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in [\mathbb{C} \ell \ell \text{unit}]_V^A$$

Let  $e_1 = (e_l := e_r) \gamma \downarrow_1$  and  $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{unit}]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} & \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}): \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2$  s.t.  $(k, H_1, H_2) \triangleright W_e$

And finally given some  $v'_1, v'_2, j < k$  s.t.  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit})$   
(FB-A0)

IH1:

$(W_e, k, e_l(\gamma \downarrow_1), e_l(\gamma \downarrow_2)) \in [\text{ref } \ell' \tau]_E^A$

This means from Definition 1.33 we need to prove:

$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies$   
 $(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  
 $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$  (FB-A1)

IH2:

$(W_e, k - f, e_r(\gamma \downarrow_1), e_r(\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau]_E^A$

This means from Definition 1.33 we need to prove:

$\forall s < k - f. e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2} \implies$   
 $(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f < k - f$  s.t  $e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$  (FB-A2)

In order to prove (FB-A0) we choose  $W'$  as  $W_e$ . Also from cg-assign we know that  $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$  and  $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$ , and  $j = f + s + 1$

We need to prove the following:

i.  $(k - j, H'_1, H'_2) \triangleright W_e$ :

Say  $v_{h1} = a_1$  and  $v'_{h1} = a_2$

From Definition 1.37 it suffices to prove:

$\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge$

$(W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge$

$\forall (a_1, a_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a_1) = W_e.\theta_2(a_2)) \wedge$

$(W_e, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [\text{ref } \ell' \tau]_V^A \wedge$

$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [W_e.\theta_i]_V$

This means we need to prove

- $\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2))$ :

Since  $\text{dom}(H_1) = \text{dom}(H'_1)$  and  $\text{dom}(H_2) = \text{dom}(H'_2)$ , and also we know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore we obtain the desired directly from Definition 1.37

- $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).(W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_e.\theta_1(a'_1) \rceil_V^A)$   
 $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).$ 
  - A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :  
From (FB-A1) and from Definition 1.32 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$   
Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$   
And since from cg-assign we know that  $H'_1(a_1) = v_{h2}$ ,  $H'_2(a_2) = v'_{h2}$  and  $j = f + s + 1$  therefore from Lemma 1.44 we get  
 $(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_e.\theta_1(a_1) \rceil_V^A$
  - B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise
  - C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise
  - D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :  
Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 1.37
- $\forall i \in \{1, 2\}.\forall m.\forall a'_i \in \text{dom}(W_e.\theta_i).(W_e.\theta_i, m, H_i(a'_i)) \in \lfloor W_e.\theta_i(a'_i) \rfloor_V$ :  
When  $i = 1$   
Given some  $m$   
 $\forall a'_1 \in \text{dom}(W_e.\theta_1).$ 
  - when  $a'_1 = a_1$ :  
From (FB-A1) and from Definition 1.32 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$   
Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$   
Therefore from Lemma 1.42 get the desired
  - Otherwise:  
Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 1.37
- When  $i = 2$   
Similar reasoning as with  $i = 1$

ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \perp, (), (), \text{unit})$ :

Holds directly from Definition 1.31 and Definition 1.32

- (b)  $\forall l \in \{1, 2\}.\left(\forall k, \theta_e \sqsupseteq \theta, H, j, (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \text{unit} \rfloor_V \wedge$   
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$ :

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor (\text{unit}) \rfloor_V \wedge$   
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$

Since  $(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$  therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in \lfloor \Gamma \rfloor_V$



Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{unit}))]_E$$

This means from Definition 1.35 we get

$$\forall c < k.(e_l := e_r)\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{unit}))]_V$$

Instantiating  $c$  with 0 and from cg-val we know that  $v = (e_l := e_r)\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit}))]_V$

From Definition 1.34 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell') \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

□

**Lemma 1.53.**  $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies$$

$$\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$$

*Proof.* Given that  $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$ . From Definition 1.31 two cases arise

1.  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

2 cases arise

(a)  $\ell' \sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 1.44 we know that  $(W', j, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

And thus from Definition 1.31 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 1.42 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 1.43 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 1.31 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Given is  $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 1.43 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 1.31 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

**Lemma 1.54** (Subtyping binary). *The following holds:*

$\forall, \tau, \tau'$ .

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_{\mathcal{V}}^A \subseteq [(\tau')]_{\mathcal{V}}^A$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^A \subseteq [(\tau')]_E^A$$

*Proof.* Proof of statement (1)

Proof by induction on the  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A \subseteq [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$

IH1:  $[(\tau'_1)]_{\mathcal{V}}^A \subseteq [(\tau_1)]_{\mathcal{V}}^A$  (Statement 1)

$[(\tau_2)]_E^A \subseteq [(\tau'_2)]_E^A$  (Sub-A0 From Statement 2)

It suffices to prove:

$$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$

From Definition 1.32 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^A &\implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A &\wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 1.32 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A &\implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2]_E^A &\wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_E) \end{aligned}$$

This means need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A) :$$

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A) \quad (63)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

Thus from Equation 63 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Finally using (Sub-A0) we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E)$ :

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1]_V$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1]_V$  and since  $\tau'_1 <: \tau_1$  therefore from Lemma 1.50 we get

$$(\theta'_l, k, v'_c) \in [\tau_1]_V \tag{64}$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2]_E) \tag{65}$$

Therefore from Equation 64 and 65 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Since  $\tau_2 <: \tau'_2$  therefore from Lemma 1.50 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_E)$ :

Similar reasoning as in the previous case

## 2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[\tau_1 \times \tau_2]_V^A \subseteq [(\tau'_1 \times \tau'_2)]_V^A$

IH1:  $[\tau_1]_V^A \subseteq [\tau'_1]_V^A$  (Statement (1))

IH2:  $[\tau_2]_V^A \subseteq [\tau'_2]_V^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau_1 \times \tau_2)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau'_1 \times \tau'_2)]_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau_1 \times \tau_2)]_V^A$

Therefore from Definition 1.32 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A \tag{66}$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau'_1 \times \tau'_2)]_V^A$

Again from Definition 1.32, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau'_2]_V^A$$

Since from Equation 66 we know that  $(W, n, v_1, v'_1) \in [\tau_1]_V^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in [\tau'_1]_V^A$

Similarly since  $(W, n, v_2, v'_2) \in [\tau_2]_V^A$  from Equation 66 therefore from IH2 we have  $(W, n, v_2, v'_2) \in [\tau'_2]_V^A$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A \subseteq \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

IH1:  $\lceil(\tau_1)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau'_1)\rceil_{\mathcal{V}}^A$  (Statement (1))

IH2:  $\lceil(\tau_2)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau'_2)\rceil_{\mathcal{V}}^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in \lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in \lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s1} = \text{inl } v_{i2}$ :

From Definition 1.32 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau_1\rceil_{\mathcal{V}}^A \tag{67}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_1\rceil_{\mathcal{V}}^A$$

From Equation 67 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_1\rceil_{\mathcal{V}}^A$$

(b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 1.32 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau_2\rceil_{\mathcal{V}}^A \tag{68}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_2\rceil_{\mathcal{V}}^A$$

From Equation 68 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_2\rceil_{\mathcal{V}}^A$$

4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A \subseteq \lceil((\text{Labeled } \ell' \tau'))\rceil_{\mathcal{V}}^A$

IH:  $\lceil(\tau)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau')\rceil_{\mathcal{V}}^A$

It suffices to prove:  $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell' \tau'))\rceil_{\mathcal{V}}^A$

This means we are given  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A$

From Definition 1.32 it means we have  $\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)$  (Sub-L0)

and it suffices to prove  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

Again from Definition 1.32 it means we need to prove that

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$$

Since we have (Sub-L0) and  $\ell \sqsubseteq \ell'$  therefore from Lemma 1.53 we have

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau)$$

2 cases arise:

(a)  $\ell' \sqsubseteq \mathcal{A}$ :

In this case from Definition 1.31 we know that  $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

From IH we also know that  $(W, n, v_1, v_2) \in [\tau']_{\mathcal{V}}^{\mathcal{A}}$

And from Definition 1.32 we get  $\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 1.31 we know that  $\forall j. (W.\theta_1, j, v_1) \in [\tau]_{\mathcal{V}}$  and  $(W.\theta_2, j, v_2) \in [\tau]_{\mathcal{V}}$

Since  $\tau <: \tau'$  therefore from Lemma 1.50 we get  $(W.\theta_1, j, v_1) \in [\tau']_{\mathcal{V}}$  and  $(W.\theta_2, j, v_2) \in [\tau']_{\mathcal{V}}$

And from Definition 1.32 we get  $\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $[((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}} \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

IH:  $[(\tau)]_{\mathcal{V}}^{\mathcal{A}} \subseteq [(\tau')]_{\mathcal{V}}^{\mathcal{A}}$

It suffices to prove:  $\forall (W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}}. (W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

This means we are given  $(W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 it means we have

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_{\mathcal{V}} \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) \right) \quad (\text{Sub-CG0}) \end{aligned}$$

And we need to prove

$$(W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$$

Again from Definition 1.32 it means we need to prove

$$\begin{aligned}
& \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\
& (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau') \right) \wedge \\
& \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) \right)
\end{aligned}$$

It means we need to prove:

$$\begin{aligned}
& \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\
& (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau') :
\end{aligned}$$

This means we are given  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t  
 $(k, H_1, H_2) \triangleright W_e, (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2)$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau')$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1)  $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau)$

Therefore from Lemma 1.53 we get  $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau)$

$$\begin{aligned}
& \text{(b) } \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) \right) :
\end{aligned}$$

Case  $l = 1$

Here we are given  $k, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

$$\text{i. } \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau' \rfloor_V :$$

Instantiating the second conjunct of (Sub-CG0) with the given  $k, \theta_e, H, j$  to get  
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau \rfloor_V$

Since  $\tau <: \tau'$  therefore from Lemma 1.50 we get  $(\theta', k - j, v'_1) \in \lfloor \tau' \rfloor_V$

$$\text{ii. } (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') :$$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')$$

Since  $\ell'_i \sqsubseteq \ell_i$  therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')$$

$$\text{iii. } (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) :$$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

Since  $\ell'_i \sqsubseteq \ell_i$  therefore we also get  
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$

Case  $l = 2$

Symmetric reasoning as in the previous  $l = 1$  case

6. CGsub-base:

Trivial

Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [(\tau)]_E^A. (W, n, e_1, e_2) \in [(\tau')]_E^A$$

This means given  $(W, n, e_1, e_2) \in [(\tau)]_E^A$

From Definition 1.33 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A \quad (\text{Sub-E0})$$

And it suffices to prove  $(W, n, e_1, e_2) \in [(\tau')]_E^A$

Again from Definition 1.33 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [(\tau')]_V^A$$

This means that given  $i < n$  s.t  $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$  we need to prove  $(W, n - i, v_1, v_2) \in [(\tau')]_V^A$

Instantiating (Sub-E0) with the given  $i$  we get  $(W, n - i, v_1, v_2) \in [\tau]_V^A$

From Statement (1) we get  $(W, n - i, v_1, v_2) \in [(\tau')]_V^A$  □

**Theorem 1.55** (NI for CG). *Say  $\text{bool} = (\text{unit} + \text{unit})$*

$\forall v_1, v_2, e, n'$ .

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2) \implies v'_1 = v'_2$$

*Proof.* Given some

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2)$$

And we need to prove

$$v'_1 = v'_2$$

From Theorem 1.52 we know that

$$\forall n. (\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \text{ bool}]_E^\perp$$

Similarly from Theorem 1.52 and Definition 1.41 we also get

$$\forall n. (\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\perp$$

From Definition 1.33 we get

$$\forall n. \forall i < n. e_1[v_1/x] \Downarrow_i v_{11} \wedge e_2 \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\perp$$

Instantiating it with  $n' + 1$  and then with  $0$ , from CG-val we have  $v_{11} = e[v_1/x]$  and  $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n' + 1, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathbb{V}}^{\perp}$$

From Definition 1.34 we have

$$\begin{aligned} & \left( \forall k \leq n' + 1, W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow_j^f (H_2', v_2'') \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\mathbf{b}]_{\mathbb{V}} \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp) \right) \end{aligned}$$

Instantiating the first conjunct with  $n' + 1, \emptyset, \emptyset, \emptyset$ . And then with  $v_1', v_2', n'$  we get  $\exists W' \sqsupseteq \emptyset. (1, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', 1, \perp, v_1', v_2', \text{bool})$

From Definition 1.31 and Definition 1.34 we get  $v_1' = v_2'$

□



### 1.3 CG to FG translation

#### 1.3.1 Type directed translation from CG to FG

CG types are translated into FG types by the following definition of  $\llbracket \cdot \rrbracket$

$$\begin{aligned} \llbracket \mathbf{b} \rrbracket &= \mathbf{b}^\perp \\ \llbracket \tau_1 \rightarrow \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \text{ref } \ell \tau \rrbracket &= (\text{ref } (\llbracket \tau \rrbracket) + \text{unit})^\ell^\perp \\ \llbracket \tau_1 \times \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \\ \llbracket \tau_1 + \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp & \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket &= (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket) + \text{unit})^{\ell_2}{}^\perp \\ \llbracket \text{Labeled } \ell \tau \rrbracket &= (\llbracket \tau \rrbracket + \text{unit})^\ell \end{aligned}$$

The translation judgment for expressions is of the form  $\boxed{\Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$ .

The translation for the pure calculus is omitted as it is straightforward.

$$\begin{array}{c} \frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{ label} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{\cdot}.e_F} \text{ unlabel} \\ \\ \frac{\Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F ())} \text{ toLabeled} \\ \\ \frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F)} \text{ ret} \\ \\ \frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_{\cdot}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{ bind} \\ \\ \frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(\text{new } (e_F))} \text{ ref} \\ \\ \frac{\Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F)} \text{ deref} \\ \\ \frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_{F1} := e_{F2})} \text{ assign} \end{array}$$

Figure 7: Expression translation from CG to FG

#### 1.3.2 Type preservation for CG to FG translation

**Theorem 1.56** (Type preservation, CG  $\rightsquigarrow$  FG).  $\forall \Gamma, e_C, \tau.$

$\Gamma \vdash e_C : \tau$  is a valid typing derivation in CG  $\implies$   
 $\exists e_F.$

$\Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$

$\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket$  is a valid typing derivation in FG

*Proof.* Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{ label}$$

$$\frac{\frac{\overline{\llbracket \Gamma \rrbracket \vdash_\top e_F : \llbracket \tau \rrbracket}}{\llbracket \Gamma \rrbracket \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\perp} \text{ IH}}{\llbracket \Gamma \rrbracket \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{ FG-inl}}{\llbracket \Gamma \rrbracket \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{ FG-sub}$$

2. unlabel:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{\cdot}.e_F} \text{ unlabel}$$

Main derivation:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_\top e_F : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{ IH}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_\top \lambda_{\cdot}.e_F : (\text{unit} \xrightarrow{\top} (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp} \text{ FG-lam}$$

3. toLabeled:

$$\frac{\Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F ())} \text{ toLabeled}$$

P2:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_\top e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{ IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{ FG-sub}$$

P1:

$$\frac{P2 \quad \overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} () : \text{unit}} \quad \mathcal{L} \vdash \ell_1 \sqcup \perp \sqsubseteq \ell_1 \quad \mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-app}$$

Main derivation:

$$\frac{P1 \quad \overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F()) : ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp} \text{ FG-inl}}{\llbracket \Gamma \rrbracket \vdash_\top \lambda_{\cdot}.\text{inl}(e_F()) : (\text{unit} \xrightarrow{\ell_1} ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp)^\perp} \text{ FG-lam}$$

4. ret:

$$\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{-}.\text{inl}(e_F)} \text{ ret}$$

$$\frac{\frac{\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : \llbracket \tau \rrbracket}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : \llbracket \tau \rrbracket} \text{ IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-sub} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-sub, FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(e_F) : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ FG-lam}$$

5. bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{ bind}$$

P1.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ FG-sub}$$

P1:

$$P1.1 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash (\ell \sqcup \perp) \sqsubseteq \ell_1 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_{F1}() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-app}$$

P2.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^{\perp}} \text{ IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqcup \ell_2 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^{\perp}} \text{ FG-sub}$$

P2:

$$P2.1 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\mathcal{L} \vdash (\ell \sqcup \ell_2 \sqcup \perp) \sqsubseteq \ell_3 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2}() : (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4}} \text{ FG-app}$$

P3:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, y : \text{unit} \vdash_{\ell \sqcup \ell_2} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\llbracket \Gamma \rrbracket, - : \text{unit}, y : \text{unit} \vdash_{\ell \sqcup \ell_2} \text{inr}() : (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4}} \text{ FG-sub, FG-inr}$$

Main derivation:

$$\begin{array}{c}
P1 \quad P2 \quad P3 \quad \frac{\overline{\mathcal{L} \vdash \ell_2 \sqsubseteq \ell_4} \text{ Given}}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4} \searrow \ell_2} \quad \frac{\overline{\ell_4 \sqsubseteq \ell'}}{\text{Given}} \\
\hline
\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\llbracket \tau' \rrbracket + \text{unit})^{\ell'}}{\text{FG-case, FG-sub}} \\
\hline
\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\text{unit} \xrightarrow{\ell} (\llbracket \tau' \rrbracket + \text{unit})^{\ell'})^{\perp} \text{ FG-lam}
\end{array}$$

6. ref:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{-}.\text{inl}(\text{new } (e_F))} \text{ ref}$$

P1:

$$\begin{array}{c}
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ IH, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ FG-sub} \\
\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell'} \searrow \ell} \\
\hline
\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{new } e_F : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp} \text{ FG-ref}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P1 \\
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(\text{new } e_F) : ((\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\perp}} \text{ FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(\text{new } e_F) : (\text{unit} \xrightarrow{\ell} ((\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\perp})^{\perp}} \text{ FG-lam}
\end{array}$$

7. deref:

$$\frac{\Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{-}.\text{inl}(e_F)} \text{ deref}$$

P2:

$$\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell})^{\perp}} \text{ IH}$$

P1:

$$\frac{P2 \quad \frac{\overline{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell} <: (\llbracket \tau \rrbracket + \text{unit})^{\ell}} \text{ Lemma 1.1} \quad \overline{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} !e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell}} \text{ FG-deref}}$$

Main derivation:

$$\begin{array}{c}
P1 \\
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \text{inl}(!e_F) : ((\llbracket \tau \rrbracket + \text{unit})^{\ell} + \text{unit})^{\perp}} \text{ FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\top} ((\llbracket \tau \rrbracket + \text{unit})^{\ell} + \text{unit})^{\perp})^{\perp}} \text{ FG-lam}
\end{array}$$

8. assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{\cdot} \text{inl}(e_{F1} := e_{F2})} \text{ assign}$$

P3:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F2} : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F2} : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ FG-sub}$$

P2:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{ IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{ FG-sub}$$

P1:

$$\frac{\begin{array}{c} P2 \quad P3 \quad \frac{\overline{\mathcal{L} \vdash \ell \sqsubseteq \ell'} \text{ Given}}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp)} \\ \hline \llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit} \end{array}}{\text{FG-assign}}$$

Main derivation:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^{\perp}} \text{ P1}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{\cdot} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^{\perp})^{\perp}} \text{ FG-lam}$$

9. sub:

$$\frac{\overline{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau' \rrbracket} \text{ IH} \quad \mathcal{L} \vdash \top \sqsubseteq \top \quad \frac{\mathcal{L} \vdash \tau' <: \tau}{\mathcal{L} \vdash \llbracket \tau' \rrbracket <: \llbracket \tau \rrbracket} \text{ Lemma 1.57}}{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket} \text{ FG-sub}$$

□

**Lemma 1.57** (Subtyping type preservation: CG to FG). *For any CG types  $\tau$  and  $\tau'$ ,  $\Sigma$ , and  $\Psi$ , if  $\mathcal{L} \vdash \tau <: \tau'$ , then  $\mathcal{L} \vdash \llbracket \tau \rrbracket <: \llbracket \tau' \rrbracket$ .*

*Proof.* Proof by induction on CG's subtyping relation

1. CGsub-base:

$$\overline{\mathcal{L} \vdash \llbracket \tau \rrbracket <: \llbracket \tau \rrbracket} \text{ Lemma 1.1}$$

2. CGsub-arrow:

$$\frac{\overline{\mathcal{L} \vdash \llbracket \tau'_1 \rrbracket <: \llbracket \tau_1 \rrbracket} \text{ IH1} \quad \overline{\mathcal{L} \vdash \llbracket \tau'_2 \rrbracket <: \llbracket \tau_2 \rrbracket} \text{ IH2} \quad \mathcal{L} \vdash \top \sqsubseteq \top}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^{\perp} <: (\llbracket \tau'_1 \rrbracket \xrightarrow{\top} \llbracket \tau'_2 \rrbracket)^{\perp}} \text{ FGsub-arrow}$$

$$\frac{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^{\perp} <: (\llbracket \tau'_1 \rrbracket \xrightarrow{\top} \llbracket \tau'_2 \rrbracket)^{\perp}}{\mathcal{L} \vdash \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2) \rrbracket <: \llbracket (\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \rrbracket} \text{ Definition of } \llbracket \cdot \rrbracket$$

3. CGsub-prod:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \llbracket \tau_2 \rrbracket <: \llbracket \tau'_2 \rrbracket}}{\text{IH2}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp <: (\llbracket \tau'_1 \rrbracket \times \llbracket \tau'_2 \rrbracket)^\perp} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket (\tau_1 \times \tau_2) \rrbracket <: \llbracket (\tau'_1 \times \tau'_2) \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

4. CGsub-sum:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \llbracket \tau_2 \rrbracket <: \llbracket \tau'_2 \rrbracket}}{\text{IH2}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp <: (\llbracket \tau'_1 \rrbracket + \llbracket \tau'_2 \rrbracket)^\perp} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket (\tau_1 + \tau_2) \rrbracket <: \llbracket (\tau'_1 + \tau'_2) \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

5. CGsub-labeled:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum} \\ \frac{\frac{\overline{\text{Labeled } \ell_1 \tau_1 <: \text{Labeled } \ell'_1 \tau'_1}}{\text{Given}}}{\ell_1 \sqsubseteq \ell'_1} \text{By inversion} \\ \hline \mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_1} \quad \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

6. CGsub-monad:

P3:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH}} \quad \frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum}$$

P2:

$$P3 \quad \frac{\frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o} \text{By inversion} \\ \hline \mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o} \quad \text{FGsub-label}$$

P1:

$$\frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \quad P2 \quad \frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o}) <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})$$

Main derivation:

$$P1 \quad \frac{\overline{\mathcal{L} \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}} \\ \hline \mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})^\perp \\ \hline \mathcal{L} \vdash \llbracket \mathbb{C} \ell_i \ell_o \tau_1 \rrbracket <: \llbracket \mathbb{C} \ell'_i \ell'_o \tau'_1 \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

□

### 1.3.3 Logical relation for CG to FG translation

**Definition 1.58** ( ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$

$$\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$$

**Definition 1.59** ( $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$

$$\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$$

**Definition 1.60** (Unary value relation).

$$\begin{aligned} \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\ \llbracket \text{unit} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket\} \\ \llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\ &\quad ({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\ &\quad \{({}^s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \\ &\quad \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}'}\} \\ \llbracket \text{ref } \ell \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s a, {}^t a) \mid {}^s\theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\ \llbracket \text{Labeled } \ell \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \\ &\quad \exists {}^s v', {}^t v'. {}^s v = \text{Lb}({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in \llbracket \tau \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ &\quad (k, H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H_s, {}^s v') \wedge i < k \implies \\ &\quad \exists H'_t, {}^t v'. (H_t, {}^t v) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ &\quad \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \rrbracket_V^{\hat{\beta}''}\} \end{aligned}$$

**Definition 1.61** (Unary expression relation).

$$\begin{aligned} \llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\quad \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies \\ &\quad \exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta\} \end{aligned}$$

**Definition 1.62** (Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a) \rrbracket_V^{\hat{\beta}} \end{aligned}$$

**Definition 1.63** (Value substitution).  $\delta^s : \text{Var} \mapsto \text{Val}$ ,  $\delta^t : \text{Var} \mapsto \text{Val}$

**Definition 1.64** (Unary interpretation of  $\Gamma$ ).

$$\begin{aligned} \llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ &\quad \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\} \end{aligned}$$

### 1.3.4 Soundness proof for CG to FG translation

**Lemma 1.65** (Monotonicity).  $\forall^s \theta, {}^s \theta', n, {}^s v, {}^t v, n', \beta, \beta'$ .

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**:

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}'}$$

Since  $({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}}$  therefore from Definition 1.60 we know that  ${}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 1.60  ${}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket$  we get the desired

2. Case **unit**:

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}'}$$

Since  $({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}}$  therefore from Definition 1.60 we know that  ${}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 1.60  ${}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket$  we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that  ${}^s v = ({}^s v_1, {}^s v_2)$  and  ${}^t v = ({}^t v_1, {}^t v_2)$ .

We also know that  $({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s \theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\underline{\text{IH1:}} ({}^s \theta', n', {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\underline{\text{IH2:}} ({}^s \theta', n', {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 1.60, IH1 and IH2 we get

$$({}^s \theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$



4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

$$\text{IH: } ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 1.60 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \rightarrow \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that

$$\forall {}^s\theta'' \sqsubseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 1.60 we are required to prove

$$\forall {}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some  ${}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$  s.t  $({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'}$  and we are required to prove

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with  ${}^s\theta'_1, {}^sv_2, {}^tv_2, j, \hat{\beta}''$  since  ${}^s\theta'_1 \sqsubseteq {}^s\theta' \sqsubseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case  $\text{ref } \ell \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{ref } \ell \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that  ${}^s v = {}^s a$  and  ${}^t v = {}^t a$ . We also know that

$${}^s\theta({}^s a) = \text{Labeled } \ell \ \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}$$

From Definition 1.60, Definition 1.58 and Definition 1.59 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

## 7. Case Labeled $\ell \ \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 1.60 it means

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, n, {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}$$

$$\text{IH: } ({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}'}$$

Similarly from Definition 1.60 we need to prove that

$$\exists {}^s v'', {}^t v''. {}^s v = \text{Lb}_\ell({}^s v'') \wedge {}^t v = \text{inl } {}^t v'' \wedge ({}^s\theta', n', {}^s v'', {}^t v'') \in [\tau]_V^{\hat{\beta}'}$$

We choose  ${}^s v''$  as  ${}^s v'$  and  ${}^t v''$  as  ${}^t v'$  and since from IH we know that  $({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}'}$

Therefore from Definition 1.60 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

## 8. Case $\mathbb{C} \ \ell_1 \ \ell_2 \ \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}'}$$

This means from Definition 1.60 we know that

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1.$$

$$(k, H_s, H_t) \triangleright_{\hat{\beta}_1} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies$$

$$\exists {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H'_s, H'_t) \triangleright_{\hat{\beta}_2} {}^s\theta' \wedge$$

$$\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', {}^t\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}_2} \wedge$$

$$(\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \quad (\text{CG0})$$

Similarly from Definition 1.60 we need to prove

$$\begin{aligned}
& \forall^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1. \\
& (k', H'_s, H'_t) \triangleright^{\hat{\beta}'_1} ({}^s \theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge i' < k' \implies \\
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s \theta') / \text{dom}({}^s \theta_e). {}^s \theta'(a) \searrow \ell_1)
\end{aligned}$$

This means we are given some  ${}^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  s.t.  $(k', H'_s, H'_t) \triangleright ({}^s \theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge i' < k'$

And we need to prove

$$\begin{aligned}
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta'', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s \theta') / \text{dom}({}^s \theta_e). {}^s \theta'(a) \searrow \ell_1)
\end{aligned}$$

Instantiating (CG0) with  ${}^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  we get the desired

□

**Lemma 1.66** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'.$

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $(\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 1.64 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 1.64 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 1.65 we get

$$\forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

□

**Lemma 1.67** (Unary monotonicity for  $H$ ).  $\forall {}^s \theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'.$

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s \theta$$

*Proof.* Given:  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta$

From Definition 1.62 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 1.62 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 1.65 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

□

**Theorem 1.68** (Fundamental theorem).  $\forall \Gamma, \tau, e, \delta^s, \delta^t, {}^s\theta, n$ .

$$\Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge$$

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$$

$\implies$

$$({}^s\theta, n, e_s, \delta^s, e_t, \delta^t) \in [\tau]_E^{\hat{\beta}}$$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. CF-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\}]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x, \delta^s, x, \delta^t) \in [\tau]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove that

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $x \delta^s \Downarrow_i {}^s v$

From cg-val we know that  $i = 0$ ,  ${}^s v = x \delta^s$ .

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-V0})$$

From fg-val we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$ . So we are left with proving

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we are given  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\}]_{\hat{V}}^{\hat{\beta}}$ , therefore from Definition 1.64 we get

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau]_{\hat{V}}^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \text{ in the context. So we are done.}$$

## 2. CF-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{ lam}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [\tau]_{\hat{E}}^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $(\lambda x. e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  ${}^s v = (\lambda x. e_s) \delta^s$ ,  ${}^t v = (\lambda x. e_t) \delta^t$ ,  $H'_t = H_t$  and  $i = 0$

It suffices to prove that

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}}$$

From Definition 1.60 it suffices to prove

$$\begin{aligned} & \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_{\hat{V}}^{\hat{\beta}'} \\ & \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_{\hat{E}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given  ${}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_{\hat{V}}^{\hat{\beta}'}$

And we need to prove

$$({}^s \theta', j, e_s[{}^s v/x] \delta^s, e_t[{}^t v/x] \delta^t) \in [\tau_2]_{\hat{E}}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{V}}^{\hat{\beta}}$  therefore from Lemma 1.66 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_1\}, e_t \cup \{x \mapsto {}^t v_1\}) \in [\tau_2]_E^{\hat{\beta}'}$$
 s.t

$$({}^s\theta', j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$$

We get (F-L0) directly from IH

3. CF-app:

$$\frac{\Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

$$\text{Also given is: } ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2]_E^{\hat{\beta}}$$

This means from Definition 1.61 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_E^{\hat{\beta}}$$

This means from Definition 1.61 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

This means from Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall k < n - j, {}^s v_2. e_{s2} \Downarrow_i {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta' \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

And we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-A2})$$

Since from (F-A1) we know that  $({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}}$  where  ${}^s v_1 = \lambda x. e'_s$  and  ${}^t v_1 = \lambda x. e'_t$

From Definition 1.60 we have

$$\begin{aligned} & \forall {}^s\theta'_3 \supseteq {}^s\theta, {}^s v, {}^t v, l < n - j, \hat{\beta}_3 \supseteq \hat{\beta}. ({}^s\theta'_3, l, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}_3} \\ & \implies ({}^s\theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}_3} \end{aligned}$$

Instantiating with  ${}^s\theta, {}^s v_2, {}^t v_2, n - j - k, \hat{\beta}$  we get

$$({}^s\theta, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (n - j - k, H_{s4}, H_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall k' < n - j - k, {}^s v_4. e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ & \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v_4, {}^t v_4) \in [\tau_2]_V^{\hat{\beta}} \wedge \\ & (n - j - k - k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t2}$ , from (F-A2) we know that  $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta$ . Instantiating  ${}^s v_4$  with  ${}^s v$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k' < i - j - k < n - j - k$  s.t  $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$ . therefore we have

$$\exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v, {}^t v_4) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k - k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-A3})$$

Since from cg-app we know that  $i = j + k + k'$  and  $H'_t = H'_{t4}$ ,  ${}^t v = {}^t v_4$  therefore we get (F-A0) from (F-A3) and Lemma 1.65 and Lemma 1.67

#### 4. CF-prod:

$$\frac{\Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{ prod}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t, \hat{\beta}.(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v.(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v.(H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n$  s.t  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v.(H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \quad (\text{F-P0})$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}.(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1.(H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1.(H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-P1})$$

IH2:

$$({}^s \theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}.(n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta \wedge \forall k < n - j. e_{s2} \delta^s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}, \hat{\beta}'_1$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-P2})$$

From cg-prod we know that  $i = j + k + 1$ ,  $H'_t = H'_{t2}$  and  ${}^t v = ({}^t v_1, {}^t v_2)$  therefore from Definition 1.60 and Lemma 1.65 we get  $({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}}$

And since we have  $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} s\theta$  therefore from Lemma 1.67 we also get

$$(n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta$$



5. CF-fst:

$$\frac{\Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1]_E^{\hat{\beta}}$  (F-F0)

This means from Definition 1.61 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j ({}^s v_1, -) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge \\ & (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and  ${}^s v_1$  with  ${}^s v$  since we know that  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j ({}^s v, -)$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge \\ & (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F1}) \end{aligned}$$

From cg-fst we know that  $i = j + 1$ ,  $H'_t = H'_{t1}$  and  ${}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}}$  therefore from Definition 1.60 and Lemma 1.65 we get

$$({}^s\theta, n - i, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$$

And since from (F-F1) we have  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 1.67 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

6. CF-snd:

Symmetric reasoning as in the CF-fst case

7. CF-inl:

$$\frac{\Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{CF-inl}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\tau_1 + \tau_2)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge \\ ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL1})$$

From cg-inl we know that  $i = j + 1$  and  $H'_t = H'_{t1}, {}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$  therefore from Definition 1.60 and Lemma 1.65 we get

$$({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}$$

And since from (F-IL1) we have  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 1.67 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Gamma \vdash \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{case}(e_t, x.e_{t1}, y.e_{t2})} \text{CF-case}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau]_E^{\hat{\beta}}$

This means from Definition 1.61 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C0)}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t.  $e_s \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C1)}$$

Two cases arise:

(a)  ${}^s v_1 = \text{inl}({}^s v'_1)$  and  ${}^t v_1 = \text{inl}({}^t v'_1)$ :

IH2:

$$({}^s\theta, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j, {}^s v_2. e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}} \wedge (n - \\ & j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i^s v$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k^s v$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-case1 we know that  $i = j + k + 1$  and  $H'_t = H'_{t2}, {}^t v = {}^t v_2$ . Since we know  $({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}}$  therefore from Definition 1.60 and Lemma 1.65 we get  $({}^s \theta, n - i, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}}$

And since from (F-C2) we have  $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$  therefore from Lemma 1.67 we get  $(n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$

(b)  ${}^s v_1 = \text{inr}({}^s v'_1)$  and  ${}^t v_1 = \text{inr}({}^t v'_1)$ :

Symmetric reasoning as in the previous case

10. CF-ret:

$$\frac{\Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Gamma \vdash \text{ret}(e_s) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda\_.\text{inl}(e_t)} \text{ret}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda\_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{ret}(e_s) \Downarrow_i^s v \implies \exists H'_t, {}^t v. (H_t, \lambda\_.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t  $\text{ret}(e_s) \delta^s \Downarrow_i^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda\_.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From CG-ret and FG-lam we know that  $i = 0, {}^s v = \text{ret}(e_s) \delta^s, {}^t v = \lambda\_.\text{inl}(e_t) \delta^t$  and  $H'_t = H_t$ .

So we need to prove

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda\_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda\_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ & (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \exists H'_t, {}^t v'. (H_t, (\lambda_{-} \text{inl}(e_t) ())) \delta^t \Downarrow \\ & (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k. \text{ Also from cg-ret we know that } H'_s = H_s$$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, (\lambda_{-} \text{inl}(e_t) ())) \delta^t \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies \\ & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s1}$  with  $H_s$  and  $H_{t1}$  with  $H_t$ . And since we know that  $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ . Therefore we have

$$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'} \wedge (k - f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-R1})$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_{t1}$ ,  ${}^t v'$  as  $\text{inl}({}^t v)$ ,  ${}^s\theta'$  as  ${}^s\theta_e$ ,  $\hat{\beta}''$  as  $\hat{\beta}'$ . Since from cg-ret we know that  $i = f + 1$  therefore from (F-R1) and Lemma 1.67 we know that

$$(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Next we choose  ${}^t v''$  as  ${}^t v$  (from F-R1) and from Lemma 1.65 we get  $({}^s\theta_e, k - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'}$  (we know from cg-ret that  ${}^s v' = {}^s v$ )

11. CF-bind:

$$\frac{\Gamma \vdash e_{s1} : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{t1} \quad \Gamma, x : \tau \vdash e_{s2} : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{t2} \quad \ell_i \sqsubseteq \ell_1 \quad \ell_i \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell_o}{\Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_{-} \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{ bind}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V'}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_{-} \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s,$   
 ${}^t v = \lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \\ & \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'''. (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'''} {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau']_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'''. (k - \\ & i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'''} {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau']_{V}^{\hat{\beta}''} \quad (\text{F-B0}) \end{aligned}$$

IH1:

$$({}^s \theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E^{\hat{\beta}}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k-j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V^{\hat{\beta}} \wedge (k-j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-B1.1})$$

From Definition 1.60 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k-j, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k-j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i-j < k-j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k-j-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k-j-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau]_V^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s \theta'', k-j-b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E^{\hat{\beta}''}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta \wedge \forall c < (k-j-b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k-j-b-c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V^{\hat{\beta}''} \wedge (k-j-b-c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \end{aligned}$$

Instantiating  $H_{s4}$  with  $H'_{s3}$  and  $H_{t4}$  with  $H'_{t3}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists c < i-j-b < k-j-b$  s.t  $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k-j-b-c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V^{\hat{\beta}''} \wedge (k-j-b-c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \quad (\text{F-B2.1}) \end{aligned}$$

From Definition 1.60 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k-j-b-c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1. \\ (m, H_{s5}, H_{t5}) \triangleright^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsupseteq {}^s \theta_e, \hat{\beta}''_1 \sqsubseteq \hat{\beta}''_2 . (m-d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s \theta''' \wedge \\ \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m-d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_V^{\hat{\beta}''_2} \end{aligned}$$

Instantiating  ${}^s\theta_e$  with  ${}^s\theta''$ ,  $H_{s5}$  with  $H'_{s3}$ ,  $H_{t5}$  with  $H'_{t3}$ ,  $m$  with  $k - j - b - c$  and  $\hat{\beta}_1''$  with  $\hat{\beta}''$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists d < i - j - b - c < k - j - b - c$  s.t.  $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s\theta''' \sqsupseteq {}^s\theta_e, \hat{\beta}_1'' \sqsubseteq \hat{\beta}_2''. (k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}_2''} {}^s\theta''' \wedge \\ & \exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}_2''} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose  $H'_{t1}$  as  $H'_{t5}$  and  ${}^t v'$  as  ${}^t v'_{h2}$ . Next we choose  ${}^s\theta'$  as  ${}^s\theta'''$  and  $\hat{\beta}''$  as  $\hat{\beta}_2''$  (both chosen from (F-B2)). Also from cg-bind we know that in (F-B0)  $H'_{s1}$  will be  $H'_{s5}$ .

Since  $(k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}_2''} {}^s\theta'''$  therefore Lemma 1.65 we get  $(k - i, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}_2''} {}^s\theta'''$

Also since from (F-B2) we have  $\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}_2''}$

Since  $i = j + b + c + d + 1$  therefore from Lemma 1.65 we get

$$\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}_2''}$$

12. CF-label:

$$\frac{\Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{ label}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ .

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}}$$

From Definition 1.61 we have



$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau)]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove  $({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}}$ . Since from cg-label we know that  $i = j + 1$ ,  ${}^s v = {}^s v_1$  and  ${}^t v = {}^t v_1$ . Therefore we get this from Definition 1.60, (F-LB1) and Lemma 1.65.

From Lemma 1.65 we get  $(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$

### 13. CF-toLabeled:

$$\frac{\Gamma \vdash e_s : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_t}{\Gamma \vdash \text{toLabeled}(e_s) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_. \text{inl}(e_t ())} \text{toLabeled}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda_. \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n$  s.t  $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda_. \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = \text{toLabeled}(e_s) \delta^s$ ,  ${}^t v = (\lambda_. \text{inl } e_t()) \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\forall^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl } e_t()) () \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell_2 \tau) \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl } e_t()) () \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell_2 \tau) \rrbracket_V^{\hat{\beta}''} \quad (\text{F-TL0}) \end{aligned}$$

IH:

$$({}^s \theta, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_E^{\hat{\beta}}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_V^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \quad (\text{F-TL1.1}) \end{aligned}$$

From Definition 1.60 we know have

$$\forall^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} & (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in \llbracket \tau \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in \llbracket \tau \rrbracket_V^{\hat{\beta}''} \quad (\text{F-TL1}) \end{aligned}$$

In order to prove (F-TL0) we choose  ${}^s\theta'$  as  ${}^s\theta''$  and  $\hat{\beta}'$  as  $\hat{\beta}''$  (both chosen from (F-TL2))  
Also from cg-toLabeled and fg-inl, fg-app we know that  $H'_s = H'_{s3}$  and  $H'_t = H'_{t3}$ , and  
 ${}^s v' = {}^s v'_{h1}$ ,  ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 1.65

14. CF-unlabel:

$$\frac{\Gamma \vdash e_s : \text{Labeled } \ell \tau \rightsquigarrow e_t}{\Gamma \vdash \text{unlabel}(e_s) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{-} e_t} \text{ unlabel}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-} e_t \delta^t) \in [\mathbb{C} \top (\ell) \tau]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \lambda_{-} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  
 $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \lambda_{-} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \\ \text{From cg-val and fg-val we know that } i = 0, {}^s v = \text{unlabel}(e_s) \delta^s, {}^t v = \lambda_{-} e_t \delta^t, H'_t = H_t \end{aligned}$$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-} e_t \delta^t) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''} \quad (\text{F-U0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h \cdot e_s \delta^s \Downarrow_f {}^s v_h &\implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} &{}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t.  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-U1})$$

In order to prove (F-U0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'$  as  ${}^t v_h$ ,  ${}^s\theta'$  as  ${}^s\theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From cg-unlabel and fg-app we also know that  $H'_{s1} = H_{s1}$  and  $H'_{t1} = H'_{t2}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

Since from (F-U1) we know that  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

Therefore from Lemma 1.67 we also get  $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'}:$$

Since from (F-U1) we have

$$({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

This means from Definition 1.60 we know that

$$\exists {}^s v_i, {}^t v_i. {}^s v_h = \text{Lb}_\ell({}^s v_i) \wedge {}^t v_h = \text{inl } {}^t v_i \wedge ({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-U2})$$

Since we know that  ${}^t v' = {}^t v_h$  and since from (F-U2) we have  ${}^t v_h = \text{inl } {}^t v_i$ . Therefore from we choose  ${}^t v''$  as  ${}^t v_i$  to get the first conjunct

From cg-unlabel we know that  ${}^s v = {}^s v_i$  and since we know that  $({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$

Therefore from Lemma 1.65 we also get  $({}^s\theta_e, k - i, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$

15. CF-ref:

$$\frac{\Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e_s : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_. \text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{new } e_s \delta^s, \lambda_. \text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_.\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_.\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}''} \quad (\text{F-N0}) \end{aligned}$$

From cg-ref we know that  ${}^s v' = a_s$  and from fg-ref, fg-inl we know that  ${}^t v' = \text{inl } a_t$ .

IH:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h. e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} \\ & {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-N1})$$

In order to prove (F-N0) we choose  $H'_{t1}$  as  $H'_{t2} \cup \{a_t \mapsto {}^t v_h\}$ ,  ${}^t v$  as  $a_t$ ,  ${}^s \theta'$  as  ${}^s \theta_n$  where  ${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}$

And we choose  $\hat{\beta}''$  as  $\hat{\beta}_n$  where  $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From cg-ref and fg-ref we also know that  $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

We need to prove

$$(a) (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}_n} {}^s \theta_n:$$

From Definition 1.62 it suffices to prove that

- $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$ :

Since  $\text{dom}({}^s \theta_e) \subseteq \text{dom}(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\} \text{ and } H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$$

Therefore we get  $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (\text{dom}({}^s \theta_e) \times \text{dom}(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}, H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\} \text{ and } \hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$$

Therefore we get  $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s \theta_n, k-i-1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s \theta_n(a)]_{V}^{\hat{\beta}_n}$ :

$$\forall (a_1, a_2) \in \hat{\beta}_n$$

- $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-N1) we know that  $({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}'}$

From Lemma 1.65 we get  $({}^s \theta_n, k-i-1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}_n}$

- $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$  therefore from Definition 1.62 we get

$$({}^s \theta_e, k-1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_{V}^{\hat{\beta}'}$$

From Lemma 1.65 we get

$$({}^s \theta_n, k-i-1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_n(a_1)]_{V}^{\hat{\beta}'}$$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_n, k-i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_{V}^{\hat{\beta}_n}:$$

We choose  ${}^t v''$  as  ${}^t v_h$  from (F-N1), fg-inl and fg-ref we know that  ${}^t v' = \text{inl } {}^t v_h$

In order to prove  $({}^s \theta_n, k-i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_{V}^{\hat{\beta}_n}$ , from Definition 1.60 it suffices to prove that

$${}^s \theta_n(a_s) = (\text{Labeled } \ell' \tau) \wedge (a_s, a_t) \in \hat{\beta}_n$$

We get this by construction of  ${}^s \theta_n$  and  $\hat{\beta}_n$

16. CF-deref:

$$\frac{\Gamma \vdash e_s : \text{ref } \ell \ \tau \rightsquigarrow e_t}{\Gamma \vdash !e_s : \mathbb{C} \top \perp (\text{Labeled } \ell \ \tau) \rightsquigarrow \lambda_{-}.\text{inl}(e_t)} \text{deref}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e_s \ \delta^s, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{E}}^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \ \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  and given some  $i < n$  s.t  $!e_s \ \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s\theta$$

From cg-val and fg-val we know that  $i = 0, {}^s v = !e_s \ \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_t) \ \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, !e_s \ \delta^s, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright} ({}^s\theta_e) \wedge (H_{s1}, !e_s \ \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta_e \wedge (H_{s1}, !e_s \ \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}''} \quad (\text{F-D0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \ell \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h \cdot e_s \delta^s \Downarrow_f {}^s v_h \implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{ref } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{ref } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'_1$  as  $H'_{t2}(a)$  (where  ${}^t v_h = a_t$  from fg-deref),  ${}^s\theta'$  as  ${}^s\theta_e$  and we choose  $\hat{\beta}''$  as  $\hat{\beta}'$ .

From cg-deref we also know that  $H'_{s1} = H_{s1}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

Since from (F-D1) we have  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$  and since  $f < i$  thfore from Lemma 1.67 we get  $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

Since from cg-deref and fg-deref we know that  ${}^s v_h = a_s$  and  ${}^t v_h = a_t$ .

Therefore from (F-D1) and from Definition 1.60 we know that

$${}^s\theta_e(a_s) = (\text{Labeled } \ell \tau) \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$  which means from Definition 1.62 we know that

$$({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \quad (\text{F-D2})$$

This means from Definition 1.60 we know that

$$\exists {}^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_\ell({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$$

We choose  ${}^t v''$  as  ${}^t v_i$  and we know that  ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$ . This proves the first conjunct.

Since from (F-D2) we have  $({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$  therefore from Lemma 1.65 we get

$$({}^s\theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

This proves the second conjunct.

17. CF-assign:



$$\frac{\Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_{s1} := e_{s2} : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s\theta$$

From cg-val and fg-val we know that  $i = 0, {}^s v = (e_{s1} := e_{s2}) \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright} ({}^s\theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} \\ & {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ & (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s\theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} \implies$$

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S1})$$

IH2:

$$({}^s\theta_e, k-f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k-f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} \implies$$

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Instantiating  $H_{s3}$  with  $H_{s1}$  and  $H_{t3}$  with  $H'_{t2}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists l < i-f < k-f$  s.t  $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$ .

Therefore we have

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s1}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S2})$$

In order to prove (F-S0) we choose  $H'_{t1}$  as  $H'_{t3}[a_t \mapsto {}^t v_{h3}]$ ,  ${}^t v'$  as  $(\ )$ ,  ${}^s\theta'$  as  ${}^s\theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From cg-assign and fg-assign we also know that  ${}^s v_{h2} = a_s$ ,  ${}^t v_{h2} = a_t$ ,  $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$  and  $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

$$(a) (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

From Definition 1.62 it suffices to prove that

- $dom({}^s\theta_e) \subseteq dom(H'_{s1})$ :

Since  $dom({}^s\theta_e) \subseteq dom(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ )

And since  $dom(H_{s1}) = dom(H'_{s1})$  therefore we also get

$$dom({}^s\theta_e) \subseteq dom(H'_{s1})$$

- $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ )

And since  $dom(H_{t1}) \subseteq dom(H'_{t1})$  therefore we also have  $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}' . ({}^s\theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$   
 $\forall (a_1, a_2) \in \hat{\beta}_n$

$$- (a_1, a_2) = (a_s, a_t):$$

Since from (F-S2) we know that  $({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

From Lemma 1.65 we get  $({}^s\theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

$$- (a_1, a_2) \neq (a_s, a_t):$$

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$  therefore  
from Definition 1.62 we get

$$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$$

From Lemma 1.65 we get

$$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$$

- (b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}_n}$ :

We choose  ${}^t v''$  as () from (F-S1), fg-inl and fg-assign we know that  ${}^t v' = \text{inl } ()$

To prove:  $({}^s\theta_n, k - i, (), ()) \in [\text{unit}]_V^{\hat{\beta}_n}$ ,

We get this directly from Definition 1.60

□

**Lemma 1.69** (Subtyping). *The following holds:*

$\forall, \tau, \tau'$ .

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V^{\hat{\beta}} \subseteq [(\tau')]_V^{\hat{\beta}}$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^{\hat{\beta}} \subseteq [(\tau')]_E^{\hat{\beta}}$$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}} . ({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

This means that given some  ${}^s\theta, n$  and  $\lambda x.e_i$  s.t  $({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}' .$$

$$({}^s\theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{S-A0})$$

And it suffices to prove:  $({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

Again from Definition 1.60 it suffices to prove:

$$\forall^s \theta'_1 \sqsupseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_1} \implies ({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_1}$$

This means that given some  ${}^s \theta'_1 \sqsubseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t.  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_1}$

And we are required to prove:  $({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_1}$

$$\text{IH: } \llbracket (\tau'_1) \rrbracket_V^{\hat{\beta}'_1} \subseteq \llbracket (\tau_1) \rrbracket_V^{\hat{\beta}'_1} \text{ (Statement (1))}$$

$$\llbracket (\tau_2) \rrbracket_E^{\hat{\beta}'_1} \subseteq \llbracket (\tau'_2) \rrbracket_E^{\hat{\beta}'_1} \text{ (Sub-A0, From Statement (2))}$$

Instantiating (S-A0) with  ${}^s \theta'_1, {}^s v_1, {}^t v_1, k, \hat{\beta}'_1$

Since  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_1}$  therefore from IH1 we know that  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'_1}$

As a result we get

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'_1}$$

From (Sub-A0), we know that

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_1}$$

## 2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2)) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V^{\hat{\beta}}$

$$\text{IH1: } \llbracket (\tau_1) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1) \rrbracket_V^{\hat{\beta}} \text{ (Statement (1))}$$

$$\text{IH2: } \llbracket (\tau_2) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2) \rrbracket_V^{\hat{\beta}} \text{ (Statement (1))}$$

It suffices to prove:

$$\forall ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2)) \rrbracket_V^{\hat{\beta}}. ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V^{\hat{\beta}}$$

This means that given  $({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2)) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s \theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s \theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V^{\hat{\beta}}$

Again from Definition 1.60, it suffices to prove:

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s \theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}$$

Since from (S-P0) we know that  $({}^s \theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}$  therefore from IH1 we have

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}}$$

Similarly since from (S-P0) we have  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$  therefore from IH2 we get  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2']_V^{\hat{\beta}}$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau_1' \quad \mathcal{L} \vdash \tau_2 <: \tau_2'}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau_1' + \tau_2'}$$

To prove:  $[\tau_1 + \tau_2]_V^{\hat{\beta}} \subseteq [(\tau_1' + \tau_2')]_V^{\hat{\beta}}$

IH1:  $[\tau_1]_V^{\hat{\beta}} \subseteq [\tau_1']_V^{\hat{\beta}}$  (Statement (1))

IH2:  $[\tau_2]_V^{\hat{\beta}} \subseteq [\tau_2']_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:  $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1' + \tau_2')]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1' + \tau_2')]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^sv = \text{inl } {}^sv_i$  and  ${}^tv = \text{inl } {}^tv_i$ :

From Definition 1.60 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1']_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1']_V^{\hat{\beta}}$$

(b)  ${}^sv = \text{inr } {}^sv_i$  and  ${}^tv = \text{inr } {}^tv_i$ :

Symmetric reasoning

4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $[\text{Labeled } \ell \tau]_V^{\hat{\beta}} \subseteq [\text{Labeled } \ell' \tau']_V^{\hat{\beta}}$

IH:  $[\tau]_V^{\hat{\beta}} \subseteq [\tau']_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^sv, {}^tv) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [(\text{Labeled } \ell' \tau')]_V^{\hat{\beta}}$$

This means that given some  $({}^s\theta, n, {}^sv, {}^tv) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$\exists s v', t v'. s v = \text{Lb}_\ell(s v') \wedge t v = \text{inl } t v' \wedge (s \theta, m, s v', t v') \in [\tau]_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$(s \theta, n, s v, t v) \in [((\text{Labeled } \ell' \tau'))]_V^{\hat{\beta}}$$

From Definition 1.60 it suffices to prove

$$\exists s v', t v'. s v = \text{Lb}_\ell(s v') \wedge t v = \text{inl } t v' \wedge (s \theta, m, s v', t v') \in [\tau']_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_1 \sqsubseteq \ell_1 \quad \mathcal{L} \vdash \ell_2 \sqsubseteq \ell'_2}{\mathcal{L} \vdash \mathbb{C} \ell_1 \ell_2 \tau <: \mathbb{C} \ell'_1 \ell'_2 \tau'}$$

$$\text{To prove: } [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}} \subseteq [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

It suffices to prove:

$$\forall (s \theta, n, s v, t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}}. (s \theta, n, s v, t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

$$\text{This means that given } (s \theta, n, s v, t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}}$$

Therefore from Definition 1.60 we are given:

$$\forall s \theta_e \sqsupseteq s \theta, H_s, H_t, i, s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} (s \theta_e) \wedge (H_s, s v) \Downarrow_i^f (H'_s, s v') \wedge i < k \implies$$

$$\exists H'_t, t v'. (H_t, t v()) \Downarrow (H'_t, t v') \wedge \exists s \theta' \sqsupseteq s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} s \theta' \wedge$$

$$\exists t v'' . t v' = \text{inl } t v'' \wedge (s \theta', k - i, s v', t v'') \in [\tau']_V^{\hat{\beta}''} \quad (\text{S-M0})$$

And we are required to prove

$$(s \theta, n, s v, t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

So again from Definition 1.60 we need to prove

$$\forall s \theta_{e1} \sqsupseteq s \theta, H_{s1}, H_{t1}, i_1, s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} (s \theta_{e1}) \wedge (H_{s1}, s v) \Downarrow_{i_1}^f (H'_{s1}, s v'_1) \wedge i_1 < k_1 \implies$$

$$\exists H'_{t1}, t v'_1. (H_{t1}, t v()) \Downarrow (H'_{t1}, t v'_1) \wedge \exists s \theta' \sqsupseteq s \theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} s \theta' \wedge$$

$$\exists t v''_1 . t v'_1 = \text{inl } t v''_1 \wedge (s \theta', k_1 - i_1, s v'_1, t v''_1) \in [\tau']_V^{\hat{\beta}''_1}$$

This means we are given some  $s \theta_{e1} \sqsupseteq s \theta, H_{s1}, H_{t1}, i_1, s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t.  $(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} (s \theta_{e1}) \wedge (H_{s1}, s v) \Downarrow_{i_1}^f (H'_{s1}, s v'_1) \wedge i_1 < k_1$

And we need to prove

$$\exists H'_{t1}, t v'_1. (H_{t1}, t v()) \Downarrow (H'_{t1}, t v'_1) \wedge \exists s \theta' \sqsupseteq s \theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} s \theta' \wedge$$

$$\exists t v''_1 . t v'_1 = \text{inl } t v''_1 \wedge (s \theta', k_1 - i_1, s v'_1, t v''_1) \in [\tau']_V^{\hat{\beta}''_1}$$

We instantiate (S-M0) with  ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$  we get

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_{e1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \end{aligned}$$

$$\text{IH: } [(\tau)]_{V'}^{\hat{\beta}''} \subseteq [(\tau')]_{V'}^{\hat{\beta}} \hat{\beta}'' \text{ (Statement (1))}$$

Since we have  $({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''}$  therefore from IH we get  $({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\tau')]_{V'}^{\hat{\beta}''}$

6. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s \theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}} . ({}^s \theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

This means that we are given  $({}^s \theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}$

From Definition 1.61 it means we have

$$\forall H_s, H_t . (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v . e_s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v . (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s \theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

From Definition 1.61 we need to prove

$$\forall H_{s1}, H_{t1} . (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1 . e_s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1 . (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau']_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

This further means that given  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$ . Also given some  $j < n, {}^s v_1$  s.t  $e_s \Downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t1}, {}^t v_1 . (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau']_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

Instantiating (Sub-E0) with the given  $H_{s1}, H_{t1}$  and  $j < n, {}^s v_1$ . We get

$$\exists H'_t, {}^t v . (H_{t1}, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau]_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we have  $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau]_{V'}^{\hat{\beta}}$  therefore from Statement(1) we get  $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [(\tau')]_{V'}^{\hat{\beta}}$

□

**Theorem 1.70** (Deriving CG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s1}, H'_{s2}$ .

let  $\text{bool} = (\text{unit} + \text{unit})$ .

$x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \wedge$

$\emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s1}, {}^s v'_1) \wedge$

$$\begin{aligned}
& (\emptyset, e_s[sv_2/x]) \Downarrow_{n_2}^f (H'_{s_2}, sv'_2) \\
& \implies \\
& sv'_1 = sv'_2
\end{aligned}$$

*Proof.* From the CG to FG translation we know that  $\exists e_t$  s.t

$$x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$$

Similarly we also know that  $\exists t v_1, t v_2$  s.t

$$\emptyset \vdash sv_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_1 \text{ and } \emptyset \vdash sv_2 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_2 \quad (\text{NI-0})$$

From type preservation theorem we know that

$$\begin{aligned}
& x : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \vdash_\top e_t : (\text{unit} \xrightarrow{\perp} ((\text{unit} + \text{unit})^\perp + \text{unit})^\perp)^\perp \\
& \emptyset \vdash_\top t v_1 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \\
& \emptyset \vdash_\top t v_2 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \quad (\text{NI-1})
\end{aligned}$$

Since we have  $\emptyset \vdash sv_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_1$

And since  $sv_1$  and  $t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 1.68 we have (we choose  $n$  s.t  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, sv_1, t v_1) \in [\text{Labeled } \top \text{ bool}]_E^\emptyset \quad (\text{NI-2})$$

And therefore from Definition 1.64 and (NI-2) we have

$$(\emptyset, n, (x \mapsto sv_1), (x \mapsto t v_1)) \in [x \mapsto \text{Labeled } \top \text{ bool}]_V^\emptyset$$

From (NI-0) we know that  $x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 1.68 to get

$$(\emptyset, n, e_s[sv_1/x], e_t[t v_1/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\emptyset \quad (\text{NI-3.1})$$

Applying Definition 1.61 on (NI-3.1) we get

$$\forall H_{s_2}, H_{t_2}. (n, H_{s_2}, H_{t_2}) \triangleright \hat{\beta} \emptyset \wedge \forall i < n. e_s[sv_1/x] \Downarrow_i sv \implies$$

$$\exists H'_{t_2}, t v. (H_{t_2}, e_t[t v_1/x]) \Downarrow (H'_{t_2}, t v) \wedge (\emptyset, n - i, sv, t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s_2}, H'_{t_2}) \triangleright \hat{\beta} \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  $sv = e_s[sv_1/x]$ .

Therefore we have

$$\exists H'_{t_2}, t v. (H_{t_2}, e_t[t v_1/x]) \Downarrow (H'_{t_2}, t v) \wedge (\emptyset, n, sv, t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n, H_{s_2}, H'_{t_2}) \triangleright \hat{\beta} \emptyset$$

From translation and from (NI-1) we know that  $t v = e_t[t v_1/x] = \lambda_. e_{b_1}$  and therefore from fg-val we have  $H'_{t_2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[sv_1/x], \lambda_. e_{b_1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$$

Expanding  $(\emptyset, n, e_s[sv_1/x], \lambda_. e_{b_1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$  using Definition 1.60 we get

$$\forall^s \theta_e \sqsupseteq \emptyset, H_{s_3}, H_{t_3}, i, sv'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s_3}, H_{t_3}) \triangleright \hat{\beta}' (s\theta_e) \wedge (H_{s_3}, e_s[sv_1/x]) \Downarrow_i^f (H'_{s_1}, sv'_1) \wedge i < k \implies$$

$$\begin{aligned}
& \exists H''_{t_1}, t v'', (H_{t_3}, (\lambda_. e_{b_1})()) \Downarrow (H''_{t_1}, t v'') \wedge \exists^s \theta' \sqsupseteq s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s_1}, H''_{t_1}) \triangleright \hat{\beta}'' s\theta' \wedge \exists t v_1'''. t v_1'' = \\
& \text{inl } t v_1''' \wedge (s\theta', k - i, sv_1'', t v_1''') \in [\text{bool}]_V^{\hat{\beta}''}
\end{aligned}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_1, sv'_1, n, \emptyset$  we get

$$\begin{aligned}
& \exists H''_{t_1}, t v''. (\emptyset, (\lambda_. e_{b_1})()) \Downarrow (H''_{t_1}, t v'') \wedge \exists^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s_1}, H''_{t_1}) \triangleright \hat{\beta}'' s\theta' \wedge \exists t v_1'''. t v_1'' = \\
& \text{inl } t v_1''' \wedge (s\theta', n - n_1, sv_1'', t v_1''') \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2})
\end{aligned}$$

Since we have  $\exists t v_1'''. t v_1'' = \text{inl } t v_1''' \wedge (s\theta', n - n_1, sv_1'', t v_1''') \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$ , therefore from Definition 1.60 we know that 2 cases arise



- ${}^s v'_1 = \text{inl}^s v'_{i1}$  and  ${}^t v'''_1 = \text{inl}^t v'_{i1}$ :

And from Definition 1.60 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means  ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr}^s v'_{i1}$  and  ${}^t v'''_1 = \text{inr}^t v'_{i1}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v'_1 = {}^t v'''_1$  (NI-3.3)

Similarly we can apply Theorem 1.68 with the other substitution to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_E^{\emptyset} \quad (\text{NI-4.1})$$

Applying Definition 1.61 on (NI-4.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} \emptyset \wedge \forall i < n, {}^s v_s.e_s[{}^s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  ${}^s v_s = e_s[{}^s v_2/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Also from (NI-1) and from translation we know that  ${}^t v = e_t[{}^t v_2/x] = \lambda_.e_{b2}$  and therefore from fg-val we know that  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\emptyset}$$

Expanding  $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\emptyset}$  using Definition 1.60 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s2}, {}^s v''_2) \wedge i < k \implies$$

$$\exists H''_{t2}, {}^t v'', (H_{t3}, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', k - i, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$  we get

$$\exists H''_{t2}, {}^t v'' . (\emptyset, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-4.2})$$

Since we have  $\exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''}$ , therefore from Definition 1.60 2 cases arise

- ${}^s v'_2 = \text{inl}^s v'_{i2}$  and  ${}^t v'''_2 = \text{inl}^t v'_{i2}$ :

And from Definition 1.60 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i2}, {}^t v'_{i2}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means  ${}^s v'_{i2} = {}^t v'_{i2} = ()$

- ${}^s v'_2 = \text{inr}^s v'_{i2}$  and  ${}^t v'''_2 = \text{inr}^t v'_{i2}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v_2' = {}^t v_2'''$  (NI-4.3)

From CG to FG translation we know that  $\exists^t v_{i1}. {}^t v_1 = \text{inl } {}^t v_{i1}$  and similarly  $\exists^t v_{i2}. {}^t v_2 = \text{inl } {}^t v_{i2}$

From (NI-1) since  $\emptyset \vdash_{\top} {}^t v_1 : (\text{bool}^{\perp} + \text{unit})^{\top}$  therefore from CG-inl we know that  $\emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\perp}$

And from CGsub-sum we know that  $\emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\top}$   
Therefore we also have  $\emptyset \vdash_{\perp} {}^t v_{i1} : \text{bool}^{\top}$  (NI-5.1)

Similarly we also have  $\emptyset \vdash_{\perp} {}^t v_{i2} : \text{bool}^{\top}$  (NI-5.2)

Next, let  $e_T = (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top}. \text{case}(e_t(), y.y, z. {}^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^{\perp}$

where  $\text{true} = \text{inl } ()$  and  $\text{false} = \text{inr } ()$

We claim  $u : \text{bool}^{\top} \vdash_{\perp} e_T : \text{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{true} : \text{bool}^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.1:

$$\frac{}{u : \text{bool}^{\top} \vdash_{\perp} u : \text{bool}^{\top}}$$

P2:

$$\frac{\text{P2.1} \quad \text{P2.2} \quad \text{P2.3} \quad \frac{}{\mathcal{L} \models (\text{bool}^{\perp} + \text{unit})^{\top} \searrow \perp}}{u : \text{bool}^{\top} \vdash_{\perp} (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : (\text{bool}^{\perp} + \text{unit})^{\top}}$$

P1.2:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^{\perp} + \text{unit})^{\perp})^{\perp}}{\text{NI-1}}}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} () : \text{unit}} \text{FG-unit}}{\frac{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp \quad \mathcal{L} \models (\text{bool}^{\perp} + \text{unit})^{\perp} \searrow \perp}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t() : (\text{bool}^{\perp} + \text{unit})^{\perp}} \text{FG-app}}$$

P1.1:

$$\frac{\text{P1.2} \quad \frac{}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, y : \text{bool}^{\perp} \vdash_{\perp} y : \text{bool}^{\perp}}{\text{FG-var}}}{\frac{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, z : \text{unit} \vdash_{\perp} \text{false} : \text{bool}^{\perp} \quad \mathcal{L} \models \text{bool}^{\perp} \searrow \perp}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z. {}^t v_b) : \text{bool}^{\perp}} \text{FG-case}}$$

P1:

$$\frac{\frac{P1.1}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top . \text{case}(e_t(), y.y, z.^t v_b)) : ((\text{bool}^\perp + \text{unit})^\top \xrightarrow{\perp} \text{bool}^\perp)^\perp}}$$

Main derivation:

$$\frac{\frac{P1 \quad P2 \quad \frac{}{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\mathcal{L} \models \text{bool}^\perp \searrow \perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top . \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^\perp} \text{FG-app}}$$

Assuming  $e_{b1}()$  reduces in  $n_{t1}$  steps in (NI-3.2) and  $e_{b2}()$  reduces in  $n_{t2}$  steps in (NI-4.2).

We instantiate Theorem 1.87 with  $e_T, ^t v_{i1}, ^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$  and  $\perp$  and therefore from (NI-3.3) and (NI-4.3) we get  $^t v_1''' = ^t v_2'''$  and thus  $^s v_1' = ^s v_2'$

□

## 1.4 FG to CG translation

### 1.4.1 Type directed (direct) translation from FG to CG

**Definition 1.71.**

$$\begin{aligned}
(\mathbf{b}) &= \mathbf{b} \\
(\mathbf{unit}) &= \mathbf{unit} \\
(\tau_1 \xrightarrow{\ell_e} \tau_2) &= (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2) \\
(\tau_1 \times \tau_2) &= (\tau_1) \times (\tau_2) \\
(\tau_1 + \tau_2) &= (\tau_1) + (\tau_2) \\
(\mathbf{ref} \ A^\ell) &= \mathbf{ref} \ \ell \ (A) \\
(A^\ell) &= \mathbf{Labeled} \ (\ell) \ (A)
\end{aligned}$$

For  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$ , define  $(\Gamma) = x_1 : (\tau_1), \dots, x_n : (\tau_n)$ .

We use a coercion function defined as follows:

$ \begin{aligned} \mathbf{coerce\_taint} &: \mathbb{C} \ pc \ \ell_c \ \tau' \rightarrow \mathbb{C} \ pc \ \perp \ \tau' \quad \text{when } \tau' = \mathbf{Labeled} \ \ell'_c \ \tau \text{ and } \ell_c \sqsubseteq \ell'_c \\ \mathbf{coerce\_taint} &\triangleq \lambda x. \mathbf{toLabeled}(\mathbf{bind}(x, y. \mathbf{unlabel}(y))) \end{aligned} $
---

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \mathbf{ret} \ x} \text{FC-var}$$

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \mathbf{ret}(\mathbf{Lb}(\lambda x. e_{c1}))} \text{FC-lam}$$

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} e_1 \ e_2 : \tau_2 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{bind}(\mathbf{unlabel} \ a, c. (c \ b)))))} \text{FC-app}$$

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{ret}(\mathbf{Lb}(a, b))))} \text{FC-prod}$$

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \mathbf{fst}(e) : \tau_1 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{fst}(b)))))} \text{FC-fst}$$

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} \mathbf{snd}(e) : \tau_2 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{snd}(b)))))} \text{FC-snd}$$

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinl}(a)))} \text{FC-inl}$$

$$\frac{\Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}$$

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \mathbf{case}(e, x. e_1, y. e_2) : \tau \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ a, b. \mathbf{case}(b, x. e_{c1}, y. e_{c2})))} \text{FC-case}$$

$$\frac{\Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \mathbf{new} \ (e) : (\mathbf{ref} \ \tau)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{new} \ (a), b. \mathbf{ret}(\mathbf{Lb}b)))} \text{FC-ref}$$

$$\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()))} \text{FC-assign}$$

### 1.4.2 Type preservation for FG to CG translation

**Theorem 1.72** (Type preservation: FG to CG). *If  $\Gamma \vdash_{pc} e : \tau$  in FG then there exists  $e'$  such that  $\Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$  such that there is a derivation of  $(\Gamma) \vdash e' : \mathbb{C} pc \perp (\tau)$  in CG.*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

$$\frac{}{(\Gamma), x : (\tau) \vdash x : (\tau)} \text{CG-var}$$

$$\frac{}{(\Gamma), x : (\tau) \vdash \text{ret } x : \mathbb{C} pc \perp (\tau)} \text{CG-ret}$$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\lambda x.e_{c1}))} \text{FC-lam}$$

$$T_0 = \mathbb{C} pc \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) = \mathbb{C} pc \perp \text{Labeled} \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2))$$

$$T_1 = \mathbb{C} pc \perp \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.0} = \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \ell_e \perp (\tau_2)$$

P1:

$$\frac{\frac{P2}{(\Gamma), x : (\tau_1) \vdash e_{c1} : T_{1.2}} \text{IH}}{(\Gamma) \vdash \lambda x.e_{c1} : T_{1.1}} \text{CG-lam}$$

Main derivation:

$$\frac{\frac{P1}{(\Gamma) \vdash (\text{Lb}(\lambda x.e_{c1})) : T_{1.0}} \text{CG-label}}{(\Gamma) \vdash \text{ret}(\text{Lb}(\lambda x.e_{c1})) : T_1} \text{CG-ret}$$

3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))))} \text{FC-app}$$

$$T_0 = \mathbb{C} \text{ pc } \perp ((\tau_1 \xrightarrow{\ell_i} \tau_2)^\ell) = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell ((\tau_1 \xrightarrow{\ell_i} \tau_2))$$

$$T_1 = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = \text{ Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \top \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.3} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.4} = \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.5} = \mathbb{C} \ell_e \ell (\tau_2)$$

$$T_{1.6} = \mathbb{C} \text{ pc } \ell (\mathbf{A}^{\ell_i})$$

$$T_{1.7} = \mathbb{C} \text{ pc } \ell \text{ Labeled } (\ell_i) (\mathbf{A})$$

$$T_{1.9} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i (\mathbf{A})$$

$$T_{1.10} = \mathbb{C} \text{ pc } \perp (\tau_2)$$

$$T_2 = \mathbb{C} \text{ pc } \perp (\tau_1)$$

$$T_{c4} = \text{ Labeled } \ell_i (\mathbf{A})$$

$$T_{c3} = \mathbb{C} \top \ell_i (\mathbf{A})$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i (\mathbf{A})$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i (\mathbf{A})$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i (\mathbf{A})$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}}$$

Pc1:

$$\frac{}{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{\frac{Pc1 \quad Pc2 \quad \frac{P0}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\text{CG-bind}}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-tolabeled}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}}$$

Pc:

$$\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}}$$

P6:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash b : (\tau_1)} \text{CG-var}$$

P5:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash c : T_{1.3}} \text{CG-var}$$

P4:

$$\frac{\frac{P5 \quad P6}{(\Gamma), a : T_{1.1}, b : \langle \tau_2 \rangle, c : T_{1.3} \vdash c b : T_{1.4}} \text{CG-app}}{(\Gamma), a : T_{1.1}, b : \langle \tau_2 \rangle, c : T_{1.3} \vdash c b : T_{1.5}} \text{CGSub-monad}$$

P3:

$$\frac{}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash a : T_{1.1}} \text{CG-var}$$

P2:

$$\frac{\frac{P3}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash \text{unlabel } a : T_{1.2}} \text{CG-unlabel} \quad P4}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash \text{bind}(\text{unlabel } a, c.(c b)) : T_{1.6}} \text{CG-bind}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{1.1} \vdash e_{c2} : T_2} \text{IH2, Weakening} \quad P2}{(\Gamma), a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b))) : T_{1.6}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Main derivation:

$$\frac{Pc \quad \frac{\frac{}{(\Gamma) \vdash e_{c1} : T_1} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.7}} \text{CG-bind}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.9}} \text{CG-app}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.10}} \text{Definition 1.71}$$

4. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\perp \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \text{Labeled } \perp \langle (\tau_1 \times \tau_2) \rangle$$

$$T_3 = \mathbb{C} \text{ pc } \perp \text{Labeled } \perp \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{3.1} = \text{Labeled } \perp \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_4 = \mathbb{C} \text{ pc } \perp \langle \tau_1 \rangle$$

$$T_5 = \mathbb{C} \text{ pc } \perp \langle \tau_2 \rangle$$

P4:

$$\frac{}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash a : \langle \tau_1 \rangle} \text{CG-var}$$

P3:

$$\frac{}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash b : \langle \tau_2 \rangle} \text{CG-var}$$

P2:

$$\frac{\frac{\frac{P3 \quad P4}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash (a, b) : \langle \tau_1 \rangle \times \langle \tau_2 \rangle} \text{CG-prod}}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_2 \rangle \vdash \text{Lb}(a, b) : T_{3.1}} \text{CG-label}}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_2 \rangle \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{\text{IH2} \quad P2}{(\Gamma), a : \langle \tau_1 \rangle \vdash e_{c2} : T_5} \text{CG-bind}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3}$$

Main derivation:

$$\frac{\frac{\frac{\text{IH1} \quad P1}{(\Gamma) \vdash e_{c1} : T_4} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{Definition 1.71}}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1}$$

5. FC-fst:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))} \text{FC-fst}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau_1 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\tau_1 \times \tau_2) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.3} = \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.4} = \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} \top \ell \langle \tau_1 \rangle$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell \langle \tau_1 \rangle$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell \langle \mathbb{A}^{\ell_i} \rangle$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp \langle \mathbb{A}^{\ell_i} \rangle$$

$$T_{c4} = \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle \mathbb{A} \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$



Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau_1 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{ CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{ From Definition of coerce\_taint}$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{ CG-var}}{\frac{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : \langle \tau_1 \rangle}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{ CG-ret}} \text{ CG-fst}$$

P1:

$$\frac{\overline{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{ CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1}} \text{ CG-bind}$$

P0:

$$\frac{\overline{(\Gamma) \vdash e_c : T_{2.2}} \text{ IH} \quad P1}{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1}} \text{ CG-bind}}{\frac{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2}}{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3}} \text{ Definition 1.71}}$$

Main derivation:

$$\frac{\overline{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.5}} \text{ CG-app}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.6}} \text{ Definition 1.71}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_1}$$

6. FC-snd:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))} \text{FC-snd}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau_2 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \langle (\tau_1 \times \tau_2) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.3} = \text{ Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.4} = \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} \top \ell \langle \tau_2 \rangle$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell \langle \tau_2 \rangle$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell \langle A^{\ell_i} \rangle$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp \langle A^{\ell_i} \rangle$$

$$T_{c4} = \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle A \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle A \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \langle A \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\frac{}{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\frac{}{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x. \text{toLabeled}(\text{bind}(x, y. \text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P2:

$$\frac{\frac{\frac{}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : \langle \tau_2 \rangle} \text{CG-snd}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b))) : T_{3.1}} \text{CG-bind}$$

P0:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_{2.2}} \text{IH} \quad P1}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.1}} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.2}}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 1.71}$$

Main derivation:

$$\frac{\frac{\frac{Pc \quad P0}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.5}} \text{CG-app}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.6}} \text{Definition 1.71}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_1}$$

7. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a. \text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \mathbb{C} \text{ pc} \perp ((\tau_1 + \tau_2)^\perp)$$

$$T_{1.1} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp ((\tau_1 + \tau_2))$$

$$T_{1.2} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp (\langle \tau_1 \rangle + \langle \tau_2 \rangle)$$

$$T_{1.3} = \text{Labeled} \perp (\langle \tau_1 \rangle + \langle \tau_2 \rangle)$$

$$T_2 = \mathbb{C} \text{ pc} \perp (\langle \tau_1 \rangle)$$

P1:

$$\frac{\frac{\frac{}{(\Gamma), a : \langle \tau_1 \rangle \vdash a : \langle \tau_1 \rangle} \text{CG-var}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{inl}(a) : \langle \tau_1 \rangle + \langle \tau_2 \rangle} \text{CG-inl}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{Lbinl}(a) : T_{1.3}} \text{CG-label}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{CG-ret}$$

Main derivation:

$$\frac{\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_2} \text{IH} \quad P1}{\langle \Gamma \rangle \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_{1.2}} \text{CG-bind}}{\langle \Gamma \rangle \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_1} \text{Definition 1.71}}$$

8. FC-inr:

$$\frac{\Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 + \tau_2)^\perp \rangle$$

$$T_{1.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp \langle (\tau_1 + \tau_2) \rangle$$

$$T_{1.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp \langle \tau_1 \rangle + \langle \tau_2 \rangle$$

$$T_{1.3} = \mathbf{Labeled} \perp \langle \tau_1 \rangle + \langle \tau_2 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle \tau_2 \rangle$$

P1:

$$\frac{\frac{\frac{\frac{}{\langle \Gamma \rangle, a : \langle \tau_2 \rangle \vdash a : \langle \tau_2 \rangle} \text{CG-var}}{\langle \Gamma \rangle, a : \langle \tau_2 \rangle \vdash \mathbf{inr}(a) : \langle \tau_1 \rangle + \langle \tau_2 \rangle} \text{CG-inr}}{\langle \Gamma \rangle, a : \langle \tau_2 \rangle \vdash \mathbf{Lbinr}(a) : T_{1.3}} \text{CG-label}}{\langle \Gamma \rangle, a : \langle \tau_2 \rangle \vdash \mathbf{ret}(\mathbf{Lbinr}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_2} \text{IH} \quad P1}{\langle \Gamma \rangle \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_{1.2}} \text{CG-bind}}{\langle \Gamma \rangle \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_1} \text{Definition 1.71}}$$

9. FC-case:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \mathbf{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.\mathbf{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case}}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 + \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ \langle \tau_1 + \tau_2 \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ (\langle \tau_1 \rangle + \langle \tau_2 \rangle)$$

$$T_{2.3} = \mathbf{Labeled} \ \ell \ (\langle \tau_1 \rangle + \langle \tau_2 \rangle)$$

$$T_{2.4} = \mathbb{C} \top \ \ell \ (\langle \tau_1 \rangle + \langle \tau_2 \rangle)$$

$$T_{2.5} = \langle \tau_1 \rangle + \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} (\text{pc } \sqcup \ell) \perp \langle \tau \rangle$$

$$T_4 = \mathbb{C} (pc \sqcup \ell) \ell (\tau)$$

$$T_5 = \mathbb{C} (pc) \ell (\mathbb{A}^{\ell_i})$$

$$T_{5.1} = \mathbb{C} (pc) \ell \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{5.3} = \mathbb{C} (pc) (\perp) \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{5.4} = \mathbb{C} (pc) (\perp) (\mathbb{A}^{\ell_i})$$

$$T_{c4} = \text{Labeled } \ell_i (\mathbb{A})$$

$$T_{c3} = \mathbb{C} \top \ell_i (\mathbb{A})$$

$$T_{c2} = \mathbb{C} pc \ell_i (\mathbb{A})$$

$$T_{c1} = \mathbb{C} pc \perp \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{c0} = \mathbb{C} pc \ell \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \overline{\mathcal{L} \models \ell \sqsubseteq \ell_i} \text{ Pg}}{\overline{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{ CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{ From Definition of coerce\_taint}$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}} \text{ CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.5}, x : (\tau_1) \vdash e_{c1} : T_3} \text{ IH2, Weakening}$$

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.5}, y : (\tau_2) \vdash e_{c2} : T_3} \text{ IH3, Weakening}}{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash \text{case}(b, x.e_{c1}, y.e_{c2}) : T_3} \text{ CG-case}$$

$$\begin{array}{c}
\text{P1:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_3} \text{CG-bind}}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_4} \text{CG-sub}
\end{array}$$

$$\begin{array}{c}
\text{P0:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.2}}{\text{IH1}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_5} \text{CG-bind}
\end{array}$$

$$\begin{array}{c}
\text{P0.2:} \\
\frac{P0}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_{5.1}} \text{Definition 1.71}
\end{array}$$

$$\begin{array}{c}
\text{P0.1:} \\
\frac{\frac{\frac{P_c \quad P0.2}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.3}}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.4}} \text{Definition 1.71}
\end{array}$$

Main derivation:

$$\frac{P0.1}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_1}$$

10. FC-ref:

$$\frac{\Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow_{pc}}{\Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})))} \text{FC-ref}$$

$$\begin{aligned}
T_1 &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } \tau)^\perp \rangle \\
T_{1.1} &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } A^{\ell_i})^\perp \rangle \\
T_{1.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \langle (\text{ref } A^{\ell_i}) \rangle \\
T_{1.3} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle \\
T_2 &= \mathbb{C} \text{ } pc \perp \langle \tau \rangle \\
T_{2.1} &= \mathbb{C} \text{ } pc \perp \langle A^{\ell_i} \rangle \\
T_{2.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled } \ell_i \langle A \rangle \\
T_{2.3} &= \text{Labeled } \ell_i \langle A \rangle \\
T_{2.4} &= \mathbb{C} \text{ } pc \perp \text{ref } \ell_i \langle A \rangle \\
T_{2.5} &= \text{ref } \ell_i \langle A \rangle \\
T_{2.51} &= \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle
\end{aligned}$$

$$\begin{array}{c}
\text{P2:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}} \text{CG-label}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}} \text{CG-ret}
\end{array}$$

P1:

$$\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{new } (a) : T_{2.4}} \text{ CG-new} \quad P2}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{ CG-bind}}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})) : T_{1.3}}$$

Main derivation:

$$\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{ IH} \quad P1}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{ CG-bind}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_{1.3}} \text{ Definition 1.71}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_1}$$

11. FC-deref:

$$\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{ FC-deref}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau' \rangle$$

$$T_{1.1} = \mathbb{C} \text{ pc } \perp \langle A'^{\ell'_i} \rangle$$

$$T_{1.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell'_i \langle A' \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\text{ref } \tau)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } \tau) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } A^{\ell_i}) \rangle$$

$$T_{2.3} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.4} = \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.6} = \text{ref } \ell_i \langle A \rangle$$

$$T_{2.7} = \mathbb{C} \top \perp \langle \text{Labeled } \ell_i \langle A \rangle \rangle$$

$$T_{2.8} = \mathbb{C} \top \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle$$

$$T_{2.9} = \mathbb{C} \text{ pc } \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle$$

$$T_{c4} = \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle A \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle A \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle A \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau' = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\frac{}{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\langle \Gamma \rangle \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\langle \Gamma \rangle \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P2:

$$\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}} \text{CG-var}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{CG-deref}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.8}} \text{CG-sub, Lemma 1.73}}$$

P1:

$$\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}} \text{CG-unlabel} \quad P2}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.3}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{Pc \quad P0}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.2}} \text{CG-app}}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{Definition 1.71}$$

12. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$T_1 = \mathbb{C} \text{ } pc \perp (\text{unit})$$

$$T_{1.1} = \mathbb{C} \text{ } pc \perp \text{unit}$$

$$T_2 = \mathbb{C} \text{ } pc \perp ((\text{ref } \tau)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } \tau))$$

$$T_{2.2} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } A^{\ell_i}))$$



$T_{2.3} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.4} = \text{Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.5} = \mathbb{C} \top (\ell) \text{ ref } \ell_i \text{ (A)}$

$T_{2.6} = \text{ref } \ell_i \text{ (A)}$

$T_{2.7} = \mathbb{C} (\text{pc } \sqcup \ell) \perp \text{unit}$

$T_{2.71} = \mathbb{C} (\text{pc } \sqcup \ell) \ell \text{unit}$

$T_{2.8} = \mathbb{C} \text{ pc } (\ell) \text{unit}$

$T_{2.9} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \text{unit}$

$T_3 = \mathbb{C} \text{ pc } \perp (\tau)$

$T_{3.1} = \mathbb{C} \text{ pc } \perp (\mathbb{A}^{\ell_i})$

$T_{3.2} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i \text{ (A)}$

$T_{3.3} = \text{Labeled } \ell_i \text{ (A)}$

P4:

$$\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{CG-var}$$

P5:

$$\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}} \text{CG-var}$$

P3:

$$\frac{\begin{array}{c} \text{P4} \quad \text{P5} \quad \frac{\frac{}{\mathcal{L} \vdash \tau \searrow (\text{pc } \sqcup \ell)}{\mathcal{L} \vdash (\text{pc } \sqcup \ell) \sqsubseteq \ell_i} \text{Given}}{\text{By inversion}} \\ \frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}} \text{CG-assign} \end{array}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.71}} \text{CGsub-monad}$$

P2:

$$\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{CG-bind} \quad \text{P3}}{\text{CG-unlabel}}$$

P1:

$$\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4} \vdash e_{c2} : T_{3.2}}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{CG-bind} \quad \text{IH2} \quad \text{P2}}{\text{CG-bind}}$$

P0:

$$\frac{\frac{}{\langle \Gamma \rangle \vdash e_{c1} : T_{2.3}}{\langle \Gamma \rangle \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.8}} \text{CG-bind} \quad \text{IH1} \quad \text{P1}}{\text{CG-bind}}$$

P0.1:

$$\frac{\text{P0}}{\langle \Gamma \rangle \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}} \text{CG-toLabeled}$$

Main derivation:

$$\frac{\text{P0.1} \quad \frac{}{\langle \Gamma \rangle, d : \text{Labeled } \ell \text{ unit} \vdash \text{ret}() : T_{1.1}}{\langle \Gamma \rangle \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) , d.\text{ret}()) : T_{1.1}}{\text{Main derivation}}$$

□

**Lemma 1.73** (Subtyping - Type preservation). *The following holds:*

1.  $\forall \tau, \tau'$ .

$$\mathcal{L} \vdash \tau <: \tau' \implies \langle \tau \rangle <: \langle \tau' \rangle$$

2.  $\forall A, A'$ .

$$\mathcal{L} \vdash A <: A' \implies \mathcal{L} \vdash \langle A \rangle <: \langle A' \rangle$$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\mathcal{L} \vdash A_1 <: A_2} \text{ By inversion} \quad P1}{\mathcal{L} \vdash \langle A_1 \rangle <: \langle A_2 \rangle} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\mathcal{L} \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \text{Labeled } \ell_1 \langle A_1 \rangle <: \text{Labeled } \ell_2 \langle A_2 \rangle} \text{ CGsub-labeled}}{\mathcal{L} \vdash \langle A_1^{\ell_1} \rangle <: \langle A_2^{\ell_2} \rangle}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\frac{\overline{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{ CG-refl}}{\mathcal{L} \vdash \langle \mathbf{b} \rangle <: \langle \mathbf{b} \rangle} \text{ Definition 1.71}$$

2. FGsub-ref:

$$\frac{\overline{\mathcal{L} \vdash \text{ref } \ell_i \langle A \rangle <: \text{ref } \ell_i \langle A \rangle} \text{ CG-refl}}{\mathcal{L} \vdash \langle \text{ref } A^{\ell_i} \rangle <: \langle \text{ref } A^{\ell_i} \rangle} \text{ Definition 1.71}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_1 \rangle <: \langle \tau'_1 \rangle} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_2 \rangle <: \langle \tau'_2 \rangle} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \langle \tau_1 \rangle \times \langle \tau_2 \rangle <: \langle \tau'_1 \rangle \times \langle \tau'_2 \rangle} \text{ CGsub-prod}}{\mathcal{L} \vdash \langle \tau_1 \times \tau_2 \rangle <: \langle \tau'_1 \times \tau'_2 \rangle} \text{ Definition 1.71}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_1 \rangle <: \langle \tau'_1 \rangle} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_2 \rangle <: \langle \tau'_2 \rangle} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \langle \tau_1 \rangle + \langle \tau_2 \rangle <: \langle \tau'_1 \rangle + \langle \tau'_2 \rangle} \text{ CGsub-prod}}{\mathcal{L} \vdash \langle \tau_1 + \tau_2 \rangle <: \langle \tau'_1 + \tau'_2 \rangle} \text{ Definition 1.71}$$

5. FGsub-arrow:

$$T_1 = \langle \tau_1 \rangle \rightarrow \mathbb{C} \ell_e \perp \langle \tau_2 \rangle$$

$$T_2 = \langle \tau'_1 \rangle \rightarrow \mathbb{C} \ell'_e \perp \langle \tau'_2 \rangle$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion, Weakening}}{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e} \text{ By inversion, Weakening}}{\mathcal{L} \vdash \mathbb{C} \ell_e \perp \langle \tau_2 \rangle <: \mathbb{C} \ell'_e \perp \langle \tau'_2 \rangle} \text{ IH(1), CGsub-monad}$$

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau'_1 <: \tau_1} \text{ By inversion, Weakening}}{\mathcal{L} \vdash \langle \tau'_1 \rangle <: \langle \tau_1 \rangle} \text{ IH(1)}$$

Main derivation:

$$\frac{P1 \quad P2}{\mathcal{L} \vdash (\tau_1 \xrightarrow{\ell_s} \tau_2) <: (\tau'_1 \xrightarrow{\ell'_s} \tau'_2)} \text{ Definition 1.71}$$

6. FGsub-unit:

$$\frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{ CGsub-unit}}{\mathcal{L} \vdash (\text{unit}) <: (\text{unit})} \text{ Definition 1.71}$$

□

### 1.4.3 Logical relation for FG to CG translation

**Definition 1.74** ( ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$

$$\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$$

**Definition 1.75** ( $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$

$$\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$$

**Definition 1.76** (Unary value relation).

$$\begin{aligned} [\mathbf{b}]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket \wedge {}^sv = {}^tv\} \\ [\text{unit}]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \text{unit} \rrbracket \wedge {}^tv \in \llbracket \text{unit} \rrbracket\} \\ [\tau_1 \times \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \mid \\ &\quad ({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}\} \\ [\tau_1 + \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^sv, \text{inl } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\tau_1]_V^{\hat{\beta}}\} \cup \\ &\quad \{({}^s\theta, m, \text{inr } {}^sv, \text{inr } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\tau_2]_V^{\hat{\beta}}\} \\ [\tau_1 \xrightarrow{\ell_s} \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \\ &\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in [\tau_1]_V^{\hat{\beta}'} \implies \\ &\quad ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'}\} \\ [\text{ref } \tau]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, a_s, a_t) \mid {}^s\theta(a_s) = \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}\} \\ [\mathbf{A}']_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, \text{Lb}({}^tv)) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\mathbf{A}]_V^{\hat{\beta}}\} \end{aligned}$$

**Definition 1.77** (Unary expression relation).

$$\begin{aligned} [\tau]_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\ &\quad \exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \\ &\quad \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [\tau]_V^{\hat{\beta}'}\} \end{aligned}$$

**Definition 1.78** (Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a_1)]_V^{\hat{\beta}} \end{aligned}$$

**Definition 1.79** (Value substitution).  $\delta^s : Var \mapsto Val, \delta^t : Var \mapsto Val$

**Definition 1.80** (Unary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V^{\hat{\beta}} \triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\}$$

#### 1.4.4 Soundness proof for FG to CG translation

**Lemma 1.81** (Monotonicity).  $\forall {}^s\theta, {}^s\theta', n, {}^s v, {}^t v, n', \beta, \beta'$ .

1.  $\forall A. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket A \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in \llbracket A \rrbracket_V^{\hat{\beta}'}$
2.  $\forall \tau. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}$

*Proof.* Proof by simultaneous induction on **A** and  $\tau$

Proof of statement (1)

We case analyze **A** in the last step

1. Case **b**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}}$  therefore from Definition 1.76 we know that  ${}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket$  and  ${}^s v = {}^t v$

Therefore from Definition 1.76 we get the desired

2. Case **unit**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in \llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}}$  therefore from Definition 1.76 we know that  ${}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 1.76 we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 we know that  ${}^sv = ({}^sv_1, {}^sv_2)$  and  ${}^tv = ({}^tv_1, {}^tv_2)$ .

We also know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1:  $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

IH2:  $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 1.76, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

IH:  $({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 1.76 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 we know that

${}^sv$  is of the form  $\lambda x.e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\lambda x.e_t$  (for some  $e_t$ ) s.t

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, {}^sv_1, {}^tv_1) \in [\tau_1]_{V}^{\hat{\beta}'_1} \implies \\ ({}^s\theta', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_{E}^{\hat{\beta}'_1} \quad (\text{A0}) \end{aligned}$$

Similarly from Definition 1.76 we are required to prove

$$\begin{aligned} \forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{V}^{\hat{\beta}''_1} \implies \\ ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''_1} \end{aligned}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t. } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{V}^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''}$$

Instantiating (A0) with  ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}''$  since

$${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n \text{ and } \hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ therefore we get}$$

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''}$$

## 6. Case ref $\tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}'}$$

From Definition 1.76 we know that  ${}^sv = a_s$  and  ${}^tv = a_t$ . We also know that

$${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$$

From Definition 1.76, Definition 1.74 and Definition 1.75 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}'}$$

### Proof of Statement (2)

Let  $\tau = \mathbf{A}^{\ell''}$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_{V}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 1.76 we know that

$$\exists {}^tv_i. {}^tv = \text{Lb}({}^tv_i) \text{ and } ({}^s\theta, n, {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_{V}^{\hat{\beta}'}$$

This means from Definition 1.76 we need to prove

$$({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

**Lemma 1.82** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}'}^{\hat{\beta}'}$$

*Proof.* Given:  $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $(\theta', n', \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}'}^{\hat{\beta}'}$

From Definition 1.80 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}}^{\hat{\beta}}$$

And again from Definition 1.80 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$ :

Since we know that  $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}}^{\hat{\beta}}$  (given)

Therefore from Lemma 1.81 we get

$$\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$$

□

**Lemma 1.83** (Unary monotonicity for  $H$ ).  $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$

From Definition 1.78 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$

And again from Definition 1.78 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$  (given)

Therefore from Lemma 1.81 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$



□

**Lemma 1.84** (Coercion lemma).  $\forall H, e, v.$

$$(H, e) \Downarrow_{-}^f (H', \text{Lb } v) \implies (H, \text{coerce\_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$$

*Proof.* Given:  $(H, e) \Downarrow_{-}^f (H', \text{Lb } v)$

To prove:  $(H, \text{coerce\_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and `cg-app` it suffices to prove that  $(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_{-}^f (H', \text{Lb } v)$

From `cg-tolabeled` it suffices to prove that  $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From `cg-bind` it suffices to prove that

1.  $(H, e) \Downarrow_{-}^f (H'_1, v_1)$ :

We are given that  $(H, e) \Downarrow_{-}^f (H', v)$  therefore we have  $H'_1 = H'$  and  $v'_1 = \text{Lb } v$

2.  $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$ :

It suffices to prove that

$$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v)$$

We get this directly from `cg-unlabel`

□

**Theorem 1.85** (Fundamental theorem).  $\forall \Gamma, \tau, e_s, e_t, pc, \delta^s, \delta^t, {}^s\theta, n, \hat{\beta}.$

$$\begin{aligned} & \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge \\ & ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}} \\ & \implies \\ & ({}^s\theta, n, e_s, \delta^s, e_t, \delta^t) \in [\tau]_{E}^{\hat{\beta}} \end{aligned}$$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\})]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x, \delta^s, \text{ret}(x), \delta^t) \in [\tau]_{E}^{\hat{\beta}}$

From Definition 1.77 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x, \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(x), \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge \\ & ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  $i = 0, {}^s v = x \delta^s$ . Also from cg-ret we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$

And we are required to prove

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ : Given

(b)  $({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}}$ :

Since we are given  $({}^s \theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\})]_V^{\hat{\beta}}$ , therefore from Definition 1.80 we get  $({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}}$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb} \lambda x. e_t)} \text{FC-lam}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (\lambda x. e_s) \delta^s, \text{ret}(\text{Lb} \lambda x. e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_E^{\hat{\beta}}$

From Definition 1.77 it suffices to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies$$

$$\exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\lambda x. e_t))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^{\hat{\beta}'}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  ${}^s v = (\lambda x. e_s) \delta^s, H'_s = H_s$  and  $i = 0$ . Also from cg-ret, cg-label and cg-FI we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\lambda x. e_t)) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ : Given

(b)  $({}^s\theta, n, \lambda x.e_s \delta^s, \text{Lb}(\lambda x.e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\hat{\beta}}^V$ :

From Definition 1.76 it suffices to prove that

$$({}^s\theta, n, \lambda x.e_s \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)]_{\hat{\beta}}^V$$

Again from Definition 1.76 it suffices to prove that

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V \implies \\ ({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2]_{\hat{\beta}'}^E \end{aligned}$$

This further means that given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t.  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V$

And we are required to prove

$$({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2]_{\hat{\beta}'}^E \quad (\text{F-L0})$$

Since we are given  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V$ , therefore from Definition 1.80 and Lemma 1.82 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\})]_{\hat{\beta}'}^V.$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2]_{\hat{\beta}'}^E$$

We get (F-L0) directly from IH

### 3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))))} \text{FC-app}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^V$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \in [\tau]_{\hat{\beta}}^E$$

This means from Definition 1.77 it suffices to prove

$$\begin{aligned} \forall H_s, H_t, (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v.(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v.(H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{\hat{\beta}'}^V \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright {}^s\theta$  and given some  $i < n, {}^s v$  s.t.  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v.(H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{\hat{\beta}'}^V \quad (\text{F-A0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{E}^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - \\ j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.0}) \end{aligned}$$

Since we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1}$  therefore from Definition 1.76 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 1.76 we know that  ${}^s v_1 = \lambda x. e'_s$  and  ${}^t v_i = \lambda x. e'_t$  s.t

$$\begin{aligned} \forall {}^s\theta''_1 \sqsupseteq {}^s\theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}''_1 \sqsubseteq \hat{\beta}'_1. \\ ({}^s\theta''_1, l, {}^s v', {}^t v') \in [\tau_1]_{V}^{\hat{\beta}''_1} \implies ({}^s\theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2]_{E}^{\hat{\beta}''_1} \quad (\text{F-A1}) \end{aligned}$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1]_{E}^{\hat{\beta}'_1}$$

This means from Definition 1.77 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2 \delta^t) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - \\ j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_{V}^{\hat{\beta}'_2} \end{aligned}$$

We instantiate with  $H'_{s1}, H'_{t1}$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s1}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$ .

This means we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - \\ j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_{V}^{\hat{\beta}'_2} \quad (\text{F-A2}) \end{aligned}$$

We instantiate (F-A1) with  $\theta''_1$  as  $\theta'_2$ ,  ${}^s v'$  as  ${}^s v_2$ ,  ${}^t v'$  as  ${}^t v_2$ ,  $l$  as  $n - j - k$  and  $\hat{\beta}''_1$  as  $\hat{\beta}'_2$ . Therefore we get

$$({}^s\theta'_2, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2]_{E}^{\hat{\beta}'_2}$$

From Definition 1.77 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s3}, {}^s v_3) \implies \\ \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s\theta'_3 \sqsupseteq {}^s\theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s\theta'_3 \wedge ({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3} \end{aligned}$$

Instantiating with  $H'_{s2}, H'_{t2}$ . since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists a < i - j - k < n - j - k$  s.t  $(H'_{s2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s3}, {}^s v_3)$

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s\theta'_3 \sqsupseteq {}^s\theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s\theta'_3 \wedge ({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3}) \end{aligned}$$

Let  $\tau_2 = A_2^{\ell_i}$ , since  $\tau_2 \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3}$$

Therefore from Definition 1.76 we know that

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, \text{Lb}^t v_{3i}) \in [\tau_2]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose  $H'_t$  as  $H'_{t3}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{3i})$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i}))$$

From cg-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.c \ b)[{}^t v_1/a][{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind again it suffices to prove

- \*  $(H'_{t2}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel and (F-A1) we know that  $H'_{t31} = H'_{t2}$  and  ${}^t v_{t2} = {}^t v_i = \lambda x. e'_t$

- \*  $((c \ b)[{}^t v_2/b][{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}:$

It suffices to prove that

$$((\lambda x. e'_t) {}^t v_2 \delta^t) \Downarrow {}^t v_{t21}$$

From cg-app we know that

$${}^t v_{t21} = e'_t[{}^t v_2/x] \delta^t$$

\*  $(H'_{t2}, {}^t v_{21}) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i}))$ :  
 From (F-A3) and (F-A3.1) we get the desired

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{V}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_3$  and  $\hat{\beta}'$  as  $\hat{\beta}'_3$ . From fg-app we know that  $i = j + k + a + 1$ ,  ${}^s v = {}^s v_3$  and  $H'_s = H'_{s3}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t3}$  and  ${}^t v = \text{Lb}({}^t v_3)$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta'$  from (F-A3) and Lemma 1.83

Since  ${}^t v = \text{Lb}({}^t v_3)$  therefore from Definition 1.76 it suffices to prove that

$$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3}$$

We get this directly from (F-A3) and Lemma 1.81

#### 4. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \in [(\tau_1 \times \tau_2)^\perp]_{E}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [(\tau_1 \times \tau_2)^\perp]_{V}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v_1, {}^s v_2$  s.t  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V}^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1]_{E}^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_{V}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_{V'}^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_{V'}^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists k < i - j < n - j$  s.t  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_{V'}^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose  $H_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_1, {}^t v_2)$

(a)  $(H_t, (\text{bind}(e_{t1}, a. \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :

From cg-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$ :  
From (F-P1) we know that  $H'_{tb1} = H'_{t1}$  and  ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From cg-bind it suffices to prove that
  - $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$ :  
From (F-P2) we know that  $H'_{tb2} = H'_{t2}$  and  ${}^t v_{tb2} = {}^t v_2$
  - $(H'_{t2}, \text{ret}(\text{Lb}(a, b))) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
We get this from cg-ret, (F-P1) and (F-P2)

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V'}^{\hat{\beta}'_1}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$  and since from fg-prod  $i = j + k + 1$  and  $H'_s = H'_{s2}$ . Therefore from (F-P2) and Lemma 1.83 we get

$$(n - i, H'_s, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'$$

In order to prove  $({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V'}^{\hat{\beta}'_2}$

From Definition 1.76 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2)]_{V'}^{\hat{\beta}'_2}$$

Since  ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$  therefore we get the desired from (F-P1), (F-P2), Definition 1.76 and Lemma 1.81

5. FC-fst:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t) \in [\tau_1]_{\hat{\beta}}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{\hat{\beta}'}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}}^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1}$  therefore from Definition 1.76 we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 1.76 we know that  ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$  and  ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$  s.t

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$



Let  $\tau_1 = \mathbf{A}_1^{\ell_i}$ , since  $\tau_1 \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\mathbf{A}_1^{\ell_i}]_V^{\hat{\beta}}$$

Therefore from Definition 1.76 we know that

$$({}^s\theta'_1, n - j, {}^s v_{i1}, \mathbf{Lb}^t v_{i11}) \in [\mathbf{A}_1]_V^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  ${}^t v_{i1} (= \mathbf{Lb}^t v_{i11})$  as we need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \mathbf{Lb}^t v_{i11}):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, \text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{i11})$ :

From (F-F1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{i11} = {}^t v_{i1} = \mathbf{Lb}({}^t v_{i1})$

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))[{}^t v_{i1}/a] \delta^t) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$ :

Again from cg-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_{i1}/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{i21})$ :

Since  ${}^t v_{i1} = \mathbf{Lb}({}^t v_{i1}, {}^t v_{i2})$  from (F-F1.1) and (F-F1.2) therefore we get the desired from cg-unlabel

So,  $H'_{t21} = H'_{t1}$  and  ${}^t v_{i21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$ :

We get the desired from cg-fst and cg-ret and (F-F1.3)

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v_{i1}) \in [\tau_1]_V^{\hat{\beta}'_1}$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . And from fg-fst we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-F1) and Lemma 1.83 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$$

Since from fg-fst we know that  ${}^s v = {}^s v_{i1}$  therefore from (F-F1.2) and Lemma 1.81 we get

$$({}^s\theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1]_V^{\hat{\beta}'_1}$$

## 6. FC-snd:

Symmetric reasoning as in the FC-fst case

## 7. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\mathbf{Lbinl}(a)))} \text{inl}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\mathbf{Lbinl}(a))) \delta^t) \in [(\tau_1 + \tau_2)^\perp]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\begin{aligned}
& \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}
\end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned}
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'} \quad (\text{F-IL0})
\end{aligned}$$

III:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned}
& \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\
& \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'_1}
\end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned}
& \exists H'_{t1}, {}^t v_1. (H_t, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'_1} \quad (\text{F-IL1})
\end{aligned}$$

In order to prove (F-IL0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $(\text{Lb inl}({}^t v_1))$  and we need to prove:

$$(a) (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-bind it suffices to prove that

$$i. (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$$

From (F-IL1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

$$ii. (H'_{t1}, \text{ret}(\text{Lbinl}(a)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-ret and (F-IL1)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'}:$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . Since from fg-inl we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-IL1) and Lemma 1.83 we get

$$(n - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta'_1$$

Now we need to prove  $({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'}$

Since  ${}^s v = \text{inl } {}^s v_1$  and  ${}^t v = \text{Lb}(\text{inl}({}^t v_1))$  therefore from Definition 1.76 it suffices to prove that

$$({}^s\theta', n - i, \text{inl } {}^s v_1, \text{inl } {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}'}$$

Since from (F-IL1) we know that  $({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$

Therefore from Lemma 1.81 and Definition 1.76 we get

$$({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}'}$$

8. FC-inr:

Symmetric reasoning as in the FC-inl case

9. FC-case:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau]_E^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

This means we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-C0})$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell]_E^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau]_V^{\hat{\beta}'_1}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell]_{V'}^{\hat{\beta}'_1}$  therefore from Definition 1.76 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2)]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

2 cases arise

(a)  ${}^s v_1 = \text{inl}({}^s v_{i1})$  and  ${}^t v_i = \text{inl}({}^t v_{i1})$ :

Also from Lemma 1.82 and Definition 1.80 we know that

$$({}^s \theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\})]_{V'}^{\hat{\beta}'_1}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t.  $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2}) \end{aligned}$$

Let  $\tau = A^{\ell_i}$  and since we know that  $\tau \searrow \ell$  therefore we have  $\ell \sqsubseteq \ell_i$

Since we have  $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$

Therefore from Definition 1.76 we have

$$({}^s \theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [A^{\ell_i}]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  ${}^t v_2 = \text{Lb}({}^t v_{2i})$

And we need to prove:

i.  $(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :

From Lemma 1.84 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :  
From (F-C1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1 / a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :  
From cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :  
Since from (F-C1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  therefore from cg-unlabel we know that  
 $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$
- $(\text{case}(b, x.e_{t1}, y.e_{t2})[{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$ :  
Since we know that in this case  ${}^t v_i = \text{inl}({}^t v_{i1})$   
Therefore from cg-case we know that  ${}^t v_{t22} = e_{t1}[{}^t v_{i1}/x] \delta^t$
- $(H'_{t1}, e_{t1}[{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \text{Lb}{}^t v_{2i})$ :  
From (F-C2) and (F-C2.1) we get the desired

ii.  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . Since from fg-case we know that  $i = j + k + 1$  and  $H'_s = H'_{s2}$  therefore from (F-C2) and Lemma 1.83 we get

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$$

Now we need to prove  $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'_2}$   
Since  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  and since from (F-C2) we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2}$$

Therefore from Lemma 1.81 and Definition 1.76 we get

$$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2}$$

(b)  ${}^s v_1 = \text{inr}({}^s v_{i1})$  and  ${}^t v_1 = \text{inr}({}^t v_{i1})$ :

Symmetric reasoning as in the previous case

10. FC-ref:

$$\frac{\Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } (e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)))} \text{ref}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{new } (e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b))) \delta^t) \delta^t \in [(\text{ref } \tau)^\perp]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ .

And we are required to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp]_V^{\hat{\beta}'} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \rfloor_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore we know that  $\exists j < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \rfloor_V^{\hat{\beta}'_1} \quad (\text{F-R1}) \end{aligned}$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_1 \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^t v = \text{Lb}(a_t)$ ,  ${}^s \theta'$  as  ${}^s \theta'_1 \cup \{a_s \mapsto \tau\}$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1})$ :  
From (F-R1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v)$ :

From cg-bind it suffices to prove that

- i.  $(H'_{t1}, \text{new } (a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v_{t2})$ :  
From cg-new we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v = a_t$
- ii.  $(H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_{t2})$ :  
From cg-ret we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_{t2} = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}:$$

From (F-R1) we know that  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$  and since  $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$ ,  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau\}$

Therefore from Definition 1.78 and Lemma 1.83 we get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'$

To prove:  $({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}$

Since we know that  ${}^s v = a_s$  and  ${}^t v = \text{Lb } a_t$  therefore we need to prove

$$({}^s \theta', n - i, a_s, \text{Lb}(a_t)) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}$$

From Definition 1.76 it suffices to prove that

$$({}^s \theta', n - i, a_s, a_t) \in \lfloor (\text{ref } \tau) \rfloor_V^{\hat{\beta}'_1}$$

Again from Definition 1.76 it suffices to prove that

$${}^s \theta'(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1$$

We get this by construction

11. FC-deref:

$$\frac{\Gamma \vdash_{pc} e_s : (\text{ref } \tau)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{deref}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))\delta^t) \in [\tau']_{\hat{\beta}}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, !e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau']_{\hat{\beta}'}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau']_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

IIH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}}^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1}$

From Definition 1.76 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 1.76 we know that  ${}^s v_1 = a_s$  and  ${}^t v_i = a_t$

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Since we are given that  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$  therefore from Definition 1.78 we know that

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [{}^s\theta(a_s)]_V^{\hat{\beta}}$$

which means we have

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [\tau]_V^{\hat{\beta}}$$

From Lemma 1.86 we know that

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [\tau']_V^{\hat{\beta}}$$

Let  $\tau' = A^{\ell_i}$  since  $\tau' \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$

Let  $v_g = H_t(a_t)$  therefore from Definition 1.76 we have

$$({}^s\theta, n-1, H_s(a_s), \text{Lb } v_{gi}) \in [\tau']_V^{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $H'_{t1}(a_t) = v_g = \text{Lb } v_{gi}$

$$(a) \ (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi})$$

From cg-bind it suffices to prove

$$i. \ (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$

$$ii. \ (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From cg-bind it suffices to prove that

$$A. \ (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$

$$B. \ (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

We get the desired from CG-deref, (F-DR1.2) and (F-DR1.3)

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n-i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s\theta', n-i, {}^s v, \text{Lb } v_{gi}) \in [\tau']_V^{\hat{\beta}'}:$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$

Therefore from (F-DR1) we get  $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$  and since  $i = j+1$  therefore from Lemma 1.83 we get  $(n-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$

Since from (F-DR1.2) we know that  $(a_s, a_t) \in \hat{\beta}'_1$  and  ${}^s\theta'_1(a_s) = \tau$ . Also from (F-

DR1) we have  $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$ . Therefore from Definition 1.77 we have  $(n-j-1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s\theta'_1(a_s)]_V^{\hat{\beta}'_1}$

Since  $i = j+1$ ,  ${}^s\theta'_1(a_s) = \tau$ ,  $H'_{s1}(a_s) = {}^s v$  and  $H'_{t1}(a_t) = {}^t v_g = \text{Lb } v_{gi}$

Therefore we get  $({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau']_V^{\hat{\beta}'}$

from (F-DR1.3) and Lemma 1.81



12. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_{s1} := e_{s2} : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{assign}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \in [\text{unit}]_E^{\hat{\beta}}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'} \quad (\text{F-AN0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \tau)^\ell]_E^{\hat{\beta}}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1}$  therefore from Definition 1.76 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 1.76 this further means that

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < n - j$  s.t.  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2} \quad (\text{F-AN2}) \end{aligned}$$

In order to prove (F-AN0) we choose  $H'_t$  as  $H'_{t2}[a_t \mapsto {}^s v_2]$ ,  ${}^t v$  as ()

We need to prove

$$(a) \ (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

$$- (H_t, \text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From cg-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti}):$$

where  ${}^t v_T = \text{Lb} {}^t v_{Ti}$

From cg-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$

From (F-AN1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12}):$

From cg-bind it suffices to prove

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13}):$

From (F-AN2) we know that  $H'_{t13} = H'_{t2}$  and  ${}^t v_{t13} = {}^t v_2$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v_{t12}):$

From cg-bind it suffices to prove that

- \*  $(H'_{t1}, \text{unlabel } a [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb} ({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)]_V^{\hat{\beta}'_1}$$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i = a_t$

- \*  $(H'_{t1}, (c := b) [{}^t v_1/a] [{}^t v_2/b] [{}^t v_i/c] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From cg-assign we know that  $H'_t = H'_{t1}[a_t \mapsto {}^t v_2]$  and  ${}^t v_{t12} = ()$

Since  ${}^t v_{t12} = {}^t v_{T_i} = ()$  therefore  ${}^t v_T = \text{Lb}()$

-  $(H'_T, \text{ret}()[{}^t v_T/d]) \delta^t \Downarrow^f (H'_t, ())$ :

From cg-ret and cg-val

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_2} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'_2}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$

In order to prove  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  it suffices to prove

- $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 1.78 we get  $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 1.78 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s \theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$ :

$\forall (a_1, a_2) \in \hat{\beta}'_2$ .

- $a_1 = a_s$  and  $a_1 = a_t$ :

Since from (F-AN2) we know that  $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 1.74 we know that  ${}^s \theta'_2(a_1) = \tau$

Therefore from Lemma 1.81 we get

$$({}^s \theta'_2, n - i - 1, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$  and  $a_1 \neq a_t$ :

From (F-AN2) since we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 1.78 we get

$$({}^s \theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

Since  $i = j + k + 1$  therefore from Lemma 1.81 we get

$$({}^s \theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

- $a_1 = a_s$  and  $a_1 \neq a_t$ :

This case cannot arise

- $a_1 \neq a_s$  and  $a_1 = a_t$ :

This case cannot arise

And in order to prove  $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'_2}$

Since we know that  ${}^s v = ()$  and  ${}^t v = ()$  therefore from Definition 1.76 we get  $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'_2}$

□

**Lemma 1.86** (Subtyping lemma). *The following holds:*

$\forall \mathcal{L}, \hat{\beta}$ .

1.  $\forall A, A'$ .

$$(a) \mathcal{L} \vdash A <: A' \implies [(A)]_{V'}^{\hat{\beta}} \subseteq [(A')]_{V'}^{\hat{\beta}}$$

2.  $\forall \tau, \tau'$ .

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau') \rrbracket_V^{\hat{\beta}}$$

$$(b) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_E^{\hat{\beta}} \subseteq \llbracket (\tau') \rrbracket_E^{\hat{\beta}}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau'_1) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau_1) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:  $\forall ({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}}$ .

$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

This means that given some  ${}^s\theta, m$  and  $\lambda x.e_s, (\lambda x.e_t)$  s.t

$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 1.76 we are given:

$$\begin{aligned} \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-L0}) \end{aligned}$$

And it suffices to prove:  $({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

Again from Definition 1.76, it suffices to prove:

$$\begin{aligned} \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L1}) \end{aligned}$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$  s.t  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2}$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with  ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$ . Since we have  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2}$  therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2)]_E^{\hat{\beta}} \subseteq [(\tau'_2)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2]_E^{\hat{\beta}'_2}$$

## 2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

$$\text{To prove: } [((\tau_1 \times \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau_1)]_V^{\hat{\beta}} \subseteq [(\tau'_1)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{IH2: } [(\tau_2)]_V^{\hat{\beta}} \subseteq [(\tau'_2)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}. ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, n$  and  ${}^s v_1, {}^s v_2, {}^t v_1, {}^t v_2$  s.t

$$({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}$$

Therefore from Definition 1.76 we are given:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \quad (\text{S-P0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

Again from Definition 1.76, it suffices to prove:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2]_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that  $({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$  therefore from IH1 we have

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}}$$

Similarly since we have  $({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$  from (S-P0) therefore from IH2 we have

$$({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2]_V^{\hat{\beta}}$$

## 3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

$$\text{To prove: } [((\tau_1 + \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau_1)]_V^{\hat{\beta}} \subseteq [(\tau'_1)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

IH2:  $[(\tau_2)]_V^{\hat{\beta}} \subseteq [(\tau'_2)]_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:  $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2))]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2))]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^sv = \text{inl } {}^sv_i$  and  ${}^tv = \text{inl } {}^tv_i$ :

From Definition 1.76 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b)  ${}^sv = \text{inr } {}^sv_i$  and  ${}^tv = \text{inr } {}^tv_i$ :

Symmetric reasoning as in the previous case

4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $[((\text{ref } \tau))]_V^{\hat{\beta}} \subseteq [((\text{ref } \tau))]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, a_s, a_t) \in [((\text{ref } \tau))]_V^{\hat{\beta}}. ({}^s\theta, n, a_s, a_t) \in [((\text{ref } \tau))]_V^{\hat{\beta}}$

We get this directly from Definition 1.76

5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $[((\mathbf{b}))]_V^{\hat{\beta}} \subseteq [((\mathbf{b}))]_V^{\hat{\beta}}$

Directly from Definition 1.76

6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $[((\text{unit}))]_V^{\hat{\beta}} \subseteq [((\text{unit}))]_V^{\hat{\beta}}$

Directly from Definition 1.76

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell' \sqsubseteq \ell'' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^{\ell'} <: A^{\ell''}} \text{FGsub-label}$$

To prove:  $[(A^{\ell'})]_V^{\hat{\beta}} \subseteq [(A^{\ell''})]_V^{\hat{\beta}}$

This means from Definition 1.76 we need to prove

$$\forall ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell'}]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell''}]_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell'}]_V^{\hat{\beta}}$

From Definition 1.76 it further means that we are given

$$({}^s\theta, n, {}^s v, {}^t v_i) \in [A]_V^{\hat{\beta}} \quad (\text{S-LB0})$$

And we need to prove

$$({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell''}]_V^{\hat{\beta}}$$

Again from Definition 1.76 it suffices to prove that

$$({}^s\theta, n, {}^s v, {}^t v_i) \in [A']_V^{\hat{\beta}}$$

Since  $\ell' \sqsubseteq \ell''$  and  $A' <: A''$  therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given:  $\mathcal{L} \vdash \tau <: \tau'$

To prove:  $[(\tau)]_E^{\hat{\beta}} \subseteq [(\tau')]_E^{\hat{\beta}}$

This means we need to prove that

$$\forall (\theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}. (\theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

This means given  $(\theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'_1} \quad (\text{S-E0})$$

And it suffices to prove that  $({}^s\theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$

Again from Definition 1.77 it means we need to prove

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1}$$

This means that given some  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} {}^s\theta$ . Also given some  $j < n, {}^s v_1$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

And we need to prove

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1} \quad (\text{S-E1})$$

Instantiating (S-E0) with  $H_{s1}, H_{t1}$  and with  $j, {}^s v_1$ . Then we get

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_t) & \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau]_V^{\hat{\beta}'_1} \end{aligned}$$

Since we have  $\tau <: \tau'$ . Therefore from IH (Statement 2(a)) we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) & \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1} \end{aligned}$$

□

**Theorem 1.87** (Deriving FG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, n_1, n_2, H'_{s1}, H'_{s2}, \perp$ .

$$\begin{aligned} & \text{Let } \mathbf{bool} = (\mathbf{unit} + \mathbf{unit}) \\ & x : \mathbf{bool}^\top \vdash_\perp e_s : \mathbf{bool}^\perp \wedge \\ & \emptyset \vdash_\perp {}^s v_1 : \mathbf{bool}^\top \wedge \emptyset \vdash_\perp {}^s v_2 : \mathbf{bool}^\top \wedge \\ & (\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^s v'_1) \wedge \\ & (\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^s v'_2) \wedge \\ & \implies \\ & {}^s v'_1 = {}^s v'_2 \end{aligned}$$

*Proof.* From the FG to CG translation we know that  $\exists e_t$  s.t

$$\begin{aligned} & x : \mathbf{bool}^\top \vdash e_s : \mathbf{bool}^\perp \rightsquigarrow e_t \\ & \text{Similarly we also know that } \exists {}^t v_1, {}^t v_2 \text{ s.t} \\ & \emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1 \text{ and } \emptyset \vdash {}^s v_2 : \mathbf{bool}^\top \rightsquigarrow {}^t v_2 \quad (\text{NI-0}) \end{aligned}$$

From type preservation theorem (choosing  $\alpha = \bar{\beta} = \perp$ ) we know that

$$\begin{aligned} & x : \text{Labeled } \top \mathbf{bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \mathbf{bool} \\ & \emptyset \vdash {}^t v_1 : \mathbb{C} \perp \perp \text{Labeled } \top \mathbf{bool} \\ & \emptyset \vdash {}^t v_2 : \mathbb{C} \perp \perp \text{Labeled } \top \mathbf{bool} \quad (\text{NI-1}) \end{aligned}$$

Since we have  $\emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 1.85 we have (we choose  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\mathbf{bool}^\top]_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 1.77 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \stackrel{\emptyset}{\triangleright} \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n - i, H'_s, H'_t) & \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{11}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\begin{aligned} & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n, H'_s, H'_t) & \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{11}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'_1} \quad (\text{NI-2.1}) \end{aligned}$$

From Definition 1.76 we know that

$${}^t v_{11} = \mathbf{Lb}({}^t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i11}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}'_1}$$

Again from Definition 1.76 we know that

Either a)  ${}^s v_1 = \mathbf{inl}()$  and  ${}^t v_{i11} = \mathbf{inl}()$  or b)  ${}^s v_1 = \mathbf{inr}()$  and  ${}^t v_{i11} = \mathbf{inr}()$

But in either case we have that  $\emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit})$  (NI-2.2)

As a result we have  $\emptyset \vdash {}^t v_{11} : \text{Labeled } \top (\mathbf{unit} + \mathbf{unit})$  (NI-2.3)

We give it typing derivation



$$\frac{\overline{\emptyset \vdash {}^t v_{i11} : (\text{unit} + \text{unit})}}{\emptyset \vdash \text{Lb}({}^t v_{i11}) : \text{Labeled } \top (\text{unit} + \text{unit})} \quad (\text{NI-2.2})$$

From Definition 1.80 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_{11})) \in \llbracket x \mapsto \text{bool}^\top \rrbracket_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 1.85 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_{11}/x]) \in \llbracket \text{bool}^\perp \rrbracket_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 1.77 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v_1'' \cdot (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v_1'') \implies \exists H'_{t1}, {}^t v_1'' \cdot (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v_1'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_1'', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_1, {}^s v_1'$  we get

$$\exists H'_{t1}, {}^t v_1'' \cdot (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v_1'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - n_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''} \quad (\text{NI-2.5})$$

Since we have  $({}^s \theta', n - n_1, {}^s v_1', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''}$  therefore from Definition 1.76 we have

$$\exists {}^t v_{i1}. {}^t v'' = \text{Lb}({}^t v_{i1}) \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_{i1}) \in \llbracket \text{bool} \rrbracket_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_1, {}^s v_1', {}^t v_{i1}) \in \llbracket (\text{unit} + \text{unit}) \rrbracket_V^{\hat{\beta}''}$  therefore from Definition 1.76 two cases arise

- ${}^s v_1' = \text{inl } {}^s v_{i11}$  and  ${}^t v_{i1} = \text{inl } {}^t v_{i11}$ :

From Definition 1.76 we have

$$({}^s \theta', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v_1' = \text{inr } {}^s v_{i11}$  and  ${}^t v_{i1} = \text{inr } {}^t v_{i11}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v_1' = {}^t v_{i1}$

$$\text{Similarly with other substitution we have } (\emptyset, n, {}^s v_2, {}^t v_2) \in \llbracket \text{bool}^\top \rrbracket_E^\emptyset \quad (\text{NI-3})$$

Therefore from Definition 1.77 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^\emptyset \emptyset \wedge \forall i < n, {}^s v \cdot (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{22} \cdot (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in \llbracket \text{bool}^\top \rrbracket_V^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, {}^t v_{22} \cdot (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in \llbracket \text{bool}^\top \rrbracket_V^{\hat{\beta}'} \quad (\text{NI-3.1})$$

From Definition 1.76 we know that

$${}^t v_2 = \text{Lb}({}^t v_{22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in \llbracket (\text{unit} + \text{unit}) \rrbracket_V^{\hat{\beta}'}$$

Again from Definition 1.76 we know that

Either a)  ${}^s v_2 = \text{inl}()$  and  ${}^t v_{i22} = \text{inl}()$  or b)  ${}^s v_2 = \text{inr}()$  and  ${}^t v_{i22} = \text{inr}()$   
 But in either case we have that  $\emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})$  (NI-3.2)

As a result we have  $\emptyset \vdash {}^t v_{22} : \text{Labeled } \top (\text{unit} + \text{unit})$  (NI-3.3)

We give it typing derivation

$$\frac{\overline{\emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})} \text{ (NI-3.2)}}{\emptyset \vdash \text{Lb}({}^t v_{i22}) : \text{Labeled } \top (\text{unit} + \text{unit})}$$

From Definition 1.80 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \text{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 1.85 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\text{bool}^\perp]_E^{\hat{\beta}'} \text{ (NI-3.4)}$$

From Definition 1.77 we get

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v_2''.(H_s, e_s[{}^s v_2/x]) \Downarrow_i (H_{s2}', {}^s v_2'') \implies \\ & \exists H_{t2}', {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H_{t2}', {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - i, H_{s2}', H_{t2}') \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_2'', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset, n_2, {}^s v_2'$  we get

$$\begin{aligned} & \exists H_{t2}', {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H_{t2}', {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - n_1, H_s', H_{t2}') \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_2', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''} \text{ (NI-3.5)} \end{aligned}$$

Since we have  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''}$  therefore from Definition 1.76 we have

$$\exists {}^t v_{i2}. {}^t v_2'' = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [\text{bool}]_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$  therefore from Definition 1.76 two cases arise

- ${}^s v_2' = \text{inl } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inl } {}^t v_{i22}$ :

From Definition 1.76 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v_2' = \text{inr } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inr } {}^t v_{i22}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v_2' = {}^t v_{i2}$

We know that  $\emptyset \vdash {}^t v_{11} : \text{Labeled } \top \text{ bool}$  (NI-2.3)

Also we have  $\emptyset \vdash {}^t v_{22} : \text{Labeled } \top \text{ bool}$  (NI-3.3)

Let  $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that  $x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{C} \perp \perp \text{ bool}$  by giving a typing derivation

P2:

$$\frac{\frac{}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool}}{\text{CG-var}}}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{C} \perp \perp \text{ bool}} \text{CG-unlabel}$$

P1:

$$\frac{}{x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \text{ bool}} \text{From (NI-1)}$$

Main derivation:

$$\frac{P1 \quad P2}{x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{C} \perp \perp \text{ bool}}$$

Say  $e_t[tv_{11}/x]$  reduces in  $n_{t1}$  steps in (NI-2.5) and  $e_t[tv_{22}/x]$  reduces in  $n_{t2}$  steps in (NI-3.5)

We instantiate Theorem 1.70 with  $e_T, {}^t v_{11}, {}^t v_{22}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$  and from (NI-2.5) and (NI-3.5) we have  ${}^t v_{i1} = {}^t v_{i2}$  and thus  ${}^s v'_1 = {}^s v'_2$

□

## 2 Part II: Alternate development with original HLIO in place of CG

### 2.1 Fine-grained IFC enforcement (FG)

#### 2.1.1 FG type system

**Syntax, types, constraints:**

Expressions	$e$	$::=$	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, x.e) \mid \text{new } e \mid !e \mid e := e \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet$
Labels	$\ell, pc$	$::=$	$l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
(Labeled) Types	$\tau$	$::=$	$A^\ell$
Unlabeled types	$A$	$::=$	$b \mid \tau \xrightarrow{\ell_c} \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \tau \mid \text{unit} \mid \forall \alpha. (\ell_e, \tau) \mid c \xrightarrow{\ell_c} \tau$
Constraints	$c$	$::=$	$\ell \sqsubseteq \ell \mid (c, c)$

**Lemma 2.1** (FG: Reflexivity of subtyping). *The following hold:*

1. For all  $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all  $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$ .

Proof of statement (1)

Let  $\tau = A^\ell$ . Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on  $A$ .

1.  $A = b$ :

$$\frac{}{\Sigma; \Psi \vdash b <: b} \text{FGsub-base}$$

2.  $A = \text{ref } \tau$ :

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}$$

**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var} \qquad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp} \text{FG-lam} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG-app} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp} \text{FG-prod} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG-fst} \qquad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp} \text{FG-inl} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau} \text{FG-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG-sub} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp} \text{FG-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}} \text{FG-assign} \\
\\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^\perp} \text{FG-unitI} \qquad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^\perp} \text{FG-FI} \\
\\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp} \text{FG-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG-CE}
\end{array}$$

Figure 8: Type system for FG

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A'^{\ell'}} \text{FGsub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2} \text{FGsub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{FGsub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{FGsub-constraint}
\end{array}$$

Figure 9: FG subtyping

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash A \text{ WF} \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi \vdash A^\ell \text{ WF}} \text{FG-wff-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 \text{ WF}} \text{FG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 \text{ WF}} \text{FG-wff-prod} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 + \tau_2 \text{ WF}} \text{FG-wff-sum} \\
\\
\frac{\text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \tau) \text{ WF}} \text{FG-wff-ref} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash (\forall \alpha. (\ell_e, \tau)) \text{ WF}} \text{FG-wff-forall} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ WF} \quad \text{FV}(c) \in \Sigma \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_\xi} \tau) \text{ WF}} \text{FG-wff-constraint}
\end{array}$$

Figure 10: Well-formedness relation for FG

4.  $A = \tau_1 + \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6.  $A = \text{unit}$ :

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$

7.  $A = \forall \alpha. \tau_i$ :

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}}$$

8.  $A = c \Rightarrow \tau_i$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash c \Rightarrow c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}}$$

□

### 2.1.2 FG semantics

Judgement:  $(H, e) \Downarrow_i (H', v)$

The semantics are described in Figure 11

### 2.1.3 Logical relation for $FG$

$W : ((Loc \mapsto Type) \times (Loc \mapsto Type) \times (Loc \leftrightarrow Loc))$

**Definition 2.2** (FG:  $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 2.3** (FG:  $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

$$\begin{array}{c}
\frac{(H, e_1) \Downarrow_i (H', \lambda x. e_i) \quad (H', e_2) \Downarrow_j (H'', v_2) \quad (H'', e_i[v_2/x]) \Downarrow_k (H''', v_3)}{(H, e_1 e_2) \Downarrow_{i+j+k+1} (H''', v_3)} \text{fg-app} \\
\\
\frac{(H, e_1) \Downarrow_i (H', v_1) \quad (H', e_2) \Downarrow_j (H'', v_2)}{(H, (e_1, e_2)) \Downarrow_{i+j+1} (H'', (v_1, v_2))} \text{fg-prod} \quad \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{fst}(e)) \Downarrow_{i+1} (H', v_1)} \text{fg-fst} \\
\\
\frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{snd}(e)) \Downarrow_{i+1} (H', v_2)} \text{fg-snd} \quad \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inl}(e)) \Downarrow_{i+1} (H', \text{inl}(v))} \text{fg-inl} \\
\\
\frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inr}(e)) \Downarrow_{i+1} (H', \text{inr}(v))} \text{fg-inr} \quad \frac{(H, e) \Downarrow_i (H', \text{inl } v) \quad (H', e_1[v/x]) \Downarrow_j (H'', v_1)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_1)} \text{fg-case1} \\
\\
\frac{(H, e) \Downarrow_i (H', \text{inr } v) \quad (H', e_2[v/x]) \Downarrow_j (H'', v_2)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_2)} \text{fg-case2} \\
\\
\frac{(H, e) \Downarrow_i (H', \Lambda e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e[]) \Downarrow_{i+j+1} (H'', v)} \text{fg-FE} \\
\\
\frac{(H, e) \Downarrow_i (H', \nu e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e\bullet) \Downarrow_{i+j+1} (H'', v)} \text{fg-CE} \\
\\
\frac{(H, e) \Downarrow_i (H', v) \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1} (H'[a \mapsto v], a)} \text{fg-ref} \quad \frac{(H, e) \Downarrow_i (H', a)}{(H, !e) \Downarrow_{i+1} (H', H(a))} \text{fg-deref} \\
\\
\frac{(H, e_1) \Downarrow_i (H', a) \quad (H', e_2) \Downarrow_j (H'', v)}{(H, e_1 := e_2) \Downarrow_{i+j+1} (H''[a \mapsto v], ())} \text{fg-assign} \quad \frac{e \in \{x, \lambda y. -, \Lambda -, \nu -\}}{(H, e) \Downarrow_0 (H, e)} \text{fg-val}
\end{array}$$

Figure 11: FG semantics



**Definition 2.4** (FG: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{b} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})\} \\
[\forall \alpha. (\ell_e, \tau)]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. \\
&\quad ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})\} \\
[c \xrightarrow{\ell_e} \tau]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n. \\
&\quad \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V^A \triangleq \begin{cases} \{(W, n, v_1, v_2) \mid (W, n, v_1, v_2) \in [\mathbf{A}]_V^A\} & \ell' \sqsubseteq \mathbf{A} \\ \{(W, n, v_1, v_2) \mid \forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in [\mathbf{A}]_V\} & \ell' \not\sqsubseteq \mathbf{A} \end{cases}$$

**Definition 2.5** (FG: Binary expression relation).

$$\begin{aligned}
[\tau]_E^A &\triangleq \{(W, n, e_1, e_2) \mid \\
&\quad \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge \\
&\quad (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\
&\quad \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_V^A\}
\end{aligned}$$

**Definition 2.6** (FG: Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V\} \\
[\tau_1 + \tau_2]_V &\triangleq \{(\theta, m, \text{inl } v) \mid (\theta, m, v) \in [\tau_1]_V\} \cup \{(\theta, m, \text{inr } v) \mid (\theta, m, v) \in [\tau_2]_V\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\
&\quad (\theta', j, e[v/x]) \in [\tau_2]_E^{\ell_e}\} \\
[\forall \alpha. (\ell_e, \tau)]_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \forall \ell' \in \mathcal{L}. (\theta', m', e) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}\} \\
[c \xrightarrow{\ell_e} \tau]_V &\triangleq \{(\theta, m, \nu e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \tau\}
\end{aligned}$$

$$[A^{\ell'}]_V \triangleq [A]_V$$

**Definition 2.7** (FG: Unary expression relation).

$$\begin{aligned} [\tau]_E^{pc} \triangleq & \{(\theta, n, e) \mid \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)\} \end{aligned}$$

**Definition 2.8** (FG: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$$

**Definition 2.9** (FG: Binary heap well formedness).

$$\begin{aligned} (n, H_1, H_2) \hat{\triangleright}^A W \triangleq & \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ & (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ & \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ & (W, n - 1, H_1(a_1), H_2(a_2)) \in [\![W.\theta_1(a_1)]\!]_V^A \wedge \\ & \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in [W.\theta_i(a_i)]_V \end{aligned}$$

**Definition 2.10** (FG: Label substitution).  $\sigma : Lvar \mapsto Label$

**Definition 2.11** (FG: Value substitution to value pairs).  $\gamma : Var \mapsto (Val, Val)$

**Definition 2.12** (FG: Value substitution to values).  $\delta : Var \mapsto Val$

**Definition 2.13** (FG: Unary interpretation of  $\Gamma$ ).

$$[\Gamma]_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V\}$$

**Definition 2.14** (FG: Binary interpretation of  $\Gamma$ ).

$$[\Gamma]_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A\}$$

### 2.1.4 Soundness proof for FG

**Lemma 2.15** (FG: Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n.$

*The following holds:*

1.  $\forall \mathbf{A}.$

$$(W, n, v_1, v_2) \in [\mathbf{A}]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbf{A}]_V$$

2.  $\forall \tau.$

$$(W, n, v_1, v_2) \in [\tau]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$$

*Proof.* Proof by simultaneous induction on  $\mathbf{A}$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $\mathbf{A}$  in the last step:

1. Case **b**:

From Definition 2.6

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A$  (S0)

IH1:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH2:  $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some  $m$  and we are required to prove:

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

Also from (S02) we know that given some  $m$  and we are required to prove:

$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH1 with  $m$  from (S01) to get

$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 2.6, we get

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH2 with  $m$  from (S02) to get

$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 2.6, we get

$(W.\theta_2, m, \text{inl}(v_{j1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{j2})$

Symmetric case as (a)

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in \lceil \tau_1 \xrightarrow{\ell_e} \tau_2 \rceil_V^A$

This means from Definition 2.4 we know that

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \lceil \tau_1 \rceil_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in \lceil \tau_2 \rceil_E^A)$   
 $\wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, i, e_1[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e})$   
 $\wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, k, e_2[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}) \quad (\text{L0})$

To prove:

(a)  $\forall m. (W.\theta_1, m, \lambda x.e_1) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V$ :

This means from Definition 2.6 we need to prove:

$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$

This further means that we have some  $\theta', j$  and  $v$  s.t

$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

And we need to prove:  $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

Therefore we get  $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$

(b)  $\forall m. (W.\theta_2, m, \lambda x.e_2) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V$ :

Similar reasoning with  $e_2$

5. Case  $\forall \alpha. (\ell_e, \tau)$ :

Given:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil \forall \alpha. (\ell_e, \tau) \rceil_V^A$

This means from Definition 2.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \quad (\text{F0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

This further means that we are given some  $\theta'$ ,  $m'$  and  $\ell_u$  s.t  $W.\theta_1 \sqsubseteq \theta'$ ,  $m' < m$  and  $\ell_u \in \mathcal{L}$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

Instantiating  $\theta_l, i$  and  $\ell''$  in the second conjunct of F0 with  $\theta'$ ,  $m'$  and  $\ell_u$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\ell_u \in \mathcal{L}$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

$$(b) \forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

Symmetric reasoning for  $e_2$

6. Case  $c \xrightarrow{\ell_e} \tau$ :

$$\text{Given: } (W, n, \nu e_1, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V^A$$

This means from Definition 2.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{C0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \nu e_1) \in [c \xrightarrow{\ell_e} \tau]_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

This further means that we are given some  $\theta'$  and  $m'$  s.t  $W.\theta_1 \sqsubseteq \theta'$ ,  $m' < m$  and  $\mathcal{L} \models c$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

Instantiating  $\theta_l, j$  in the second conjunct of C0 with  $\theta'$ ,  $m'$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\mathcal{L} \models c$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

$$(b) \forall m. (W.\theta_2, m, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V:$$

Symmetric reasoning for  $e_2$

7. Case ref  $\tau$ :

From Definition 2.4 and 2.6

Proof of statement (2)

Let  $\tau = A^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement(1))

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 2.4

□

**Lemma 2.16** (FG: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m'$ .

1.  $\forall \mathbf{A}. (\theta, m, v) \in \llbracket \mathbf{A} \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \mathbf{A} \rrbracket_V$

2.  $\forall \tau. (\theta, m, v) \in \llbracket \tau \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \tau \rrbracket_V$

*Proof.* Proof by simultaneous induction on  $\mathbf{A}$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $\mathbf{A}$  in the last step:

1. case  $\mathbf{b}$ :

Directly from Definition 2.6

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

This means from Definition 2.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V$

IH1 :  $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

IH2 :  $(\theta', m', v_2) \in \llbracket \tau_2 \rrbracket_V$

We get the desired from IH1, IH2 and Definition 2.6

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

This means from Definition 2.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V$

IH :  $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

Therefore from IH and Definition 2.6 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(\theta, m, (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

To prove:  $(\theta', m', (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

This means from Definition 2.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e} \quad (69)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating Equation 143 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

5. case ref  $\tau$ :

From Definition 2.6 and Definition 2.2

6. case  $\forall \alpha. (\ell_e, \tau)$ :

Given:  $(\theta, m, (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

To prove:  $(\theta', m', (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

This means from Definition 2.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E^{\ell_e[\ell_i/\alpha]} \quad (70)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

Instantiating Equation 70 with  $\theta''', k$  and  $\ell_j$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\ell_j \in \mathcal{L}$

Therefore we get  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

7. case  $c \xrightarrow{\ell_e} \tau$ :

Given:  $(\theta, m, (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

To prove:  $(\theta', m', (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

This means from Definition 2.6 we know that

$$\forall \theta'' . \theta \sqsubseteq \theta'' \wedge \forall j < m . \mathcal{L} \models c \implies (\theta'' , j, e_1) \in [\tau]_E^{\ell_e} \quad (71)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta''' . \theta' \sqsubseteq \theta''' \wedge \forall k < m' . \mathcal{L} \models c \implies (\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$$

This means that given some  $\theta''' , k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating Equation 71 with  $\theta''' , k$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Proof of statement (2)

Let  $\tau = A^\ell$

Since  $[A^\ell]_V = [A]_V$ , therefore from IH (statement 1) □

**Lemma 2.17** (FG: Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, \mathcal{A}, n, n'$ .

$$1. \forall A. (W, n, v_1, v_2) \in [A]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [A]_V^A$$

$$2. \forall \tau. (W, n, v_1, v_2) \in [\tau]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_V^A$$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We analyze the different cases of  $A$  in the last step:

1. Case  $b$ :

From Definition 2.4

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 2.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$$

$$\text{IH1 : } (W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$$

$$\text{IH2 : } (W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$$

From IH1, IH2 and Definition 2.4 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:



(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 2.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

This means from Definition 2.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A2)

Similarly from Definition 2.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we a required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we a required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W.\theta_2, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we a required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

5. Case ref  $\tau$ :

From Definition 2.4 and Definition 2.3

6. Case  $\forall\alpha.(\ell_e, \tau)$ :

Given:  $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

To prove:  $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

This means from Definition 2.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}.((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}.((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}.((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F2})$$

Similarly from Definition 2.4 we know that we are required to prove

$$(a) \forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}.((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A):$$

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with  $W'', n''$  and  $\ell''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . Also since  $n'' < n'$  and  $n' < n$  therefore  $n'' < n$ . And finally since  $\ell'' \in \mathcal{L}$  therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

$$(b) \forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}.((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_l \sqsupseteq W.\theta_1$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

$$(c) \forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}.((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some  $\theta_l \sqsupseteq W.\theta_2, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F2 with  $\theta_l, k$  and  $\ell''$ . And since  $\theta_l \sqsupseteq W.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta_l \sqsupseteq W.\theta_2$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

7. Case  $c \stackrel{\ell_e}{\Rightarrow} \tau$ :

Given:  $(W, n, (\nu e_1), (\nu e_2)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

To prove:  $(\theta', n', (\nu e_1), (\nu e_1)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

This means from Definition 2.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \quad (\text{BM-C0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \quad (\text{BM-C1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{BM-C2})$$

Similarly from Definition 2.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in [\tau]_E^A$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\mathcal{L} \models c$

And we are required to prove:  $(W'', n'', e_1, e_2) \in [\tau]_E^A$

Instantiating BM-C0 with  $W'', n''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . And since  $\mathcal{L} \models c$  therefore we get

$(W'', n'', e_1, e_2) \in [\tau]_E^A$

(b)  $\forall \theta'_i \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_1, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with  $\theta'_i, k$ . And since  $\theta'_i \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_i \sqsupseteq W.\theta_1$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

(c)  $\forall \theta'_i \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_2, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with  $\theta'_i, k$ . And since  $\theta'_i \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_i \sqsupseteq W.\theta_2$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

#### Proof of statement (2)

Let  $\tau = A^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement 1)

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

From Lemma 2.16 and Definition 2.4

□

**Lemma 2.18** (FG: Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .

$(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

*Proof.* Given:  $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove:  $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 2.13 it is given that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$

And again from Definition 2.13 we are required to prove that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$

- $dom(\Gamma) \subseteq dom(\delta)$ :

Given

- $\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$ :  
 Since we know that  $\forall x \in \text{dom}(\Gamma).(\theta, n, \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$  (given)  
 Therefore from Lemma 2.16 we get  
 $\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$

□

**Lemma 2.19** (FG: Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .  
 $(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \lfloor \Gamma \rfloor_V$

*Proof.* Given:  $(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(W', n', \gamma) \in \lfloor \Gamma \rfloor_V$

From Definition 2.14 it is given that  
 $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

And again from Definition 2.13 we are required to prove that  
 $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma)$ :  
 Given
- $\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$ :  
 Since we know that  $\forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$  (given)  
 Therefore from Lemma 2.17 we get  
 $\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

□

**Lemma 2.20** (FG: Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .  
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$   
 To prove:  $(n', H) \triangleright \theta$

From Definition 2.8 it is given that  
 $\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$

And again from Definition 2.13 we are required to prove that  
 $\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$ :  
 Given
- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :  
 Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)  
 Therefore from Lemma 2.16 we get  
 $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$

□

**Lemma 2.21** (FG: Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .  
 $(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 2.9 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W.\hat{\beta}). &(W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n-1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 2.9 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :  
 Given
- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :  
 Given
- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
 $\forall (a_1, a_2) \in (W.\hat{\beta})$ .
  - $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given
  - $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
 Given and from Lemma 2.17
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :  
 Given

□

**Theorem 2.22** (FG: Fundamental theorem unary).  $\forall \Sigma, \Psi, \Gamma, pc, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n$ .

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \\ \mathcal{L} \models \Psi \sigma \wedge \\ (\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V \implies \\ (\theta, n, e \delta) \in \lfloor \tau \sigma \rfloor_E^{pc} \end{aligned}$$

*Proof.* Proof by induction on *FG* typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove:  $(\theta, n, x \delta) \in \lfloor \tau \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, x \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-V0})
\end{aligned}$$

In order to prove FU-V0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $v' = \delta(x)$  and  $j = 1$

We need to prove the following:

- (a)  $\theta \sqsubseteq \theta \wedge (n - 1, H) \triangleright \theta \wedge (\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$ :
- $\theta \sqsubseteq \theta$ : From Definition 2.2
  - $(n - 1, H) \triangleright \theta$ : From Lemma 2.20
  - $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$ :  
Since we are given that  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$  and  $v' = \delta(x)$   
Therefore  $(\theta, n, v') \in \lfloor \Gamma(x) \sigma \rfloor_V$ , where  $\Gamma(x) = \tau$   
And finally from Lemma 2.16 we get  $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$
- (b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :  
Since  $H' = H$ , so we are done
- (c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :  
Since  $\theta' = \theta$ , so we are done

## 2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(\theta, \lambda x. e_i \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, (\lambda x. e_i) \delta) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-L0})
\end{aligned}$$

IH1:

$\forall \theta_i, v_x, n. (\theta_i, n, e_i \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$ , s.t  $(\theta_i, n, v_x) \in [\tau_1 \sigma]_V$

In order to prove FU-L0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $j = 0$  and  $v' = \lambda x. e_i \delta$

(a)  $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$ :

- $\theta \sqsubseteq \theta$ : From Definition 2.2
- $(n, H) \triangleright \theta$ : Given
- $(\theta, n, (\lambda x. e_i) \delta) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$ :

From Definition 2.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \forall v. (\theta'', j, v) \in [\tau_1 \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

This means given some  $\theta''$ ,  $j$  and  $v$  such that  $\theta \sqsubseteq \theta''$ ,  $j < n$  and  $(\theta'', j, v) \in [\tau_1 \sigma]_V$   
It suffices to prove that  $(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  and  $j < n$  therefore from Lemma 2.18 we have  
 $(\theta, j, \delta) \in [\Gamma \sigma]_V$

So we can apply IH1 instantiated with  $\theta''$ ,  $v$  and  $j$  we get  
 $(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

(b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $H' = H$  so we are done

(c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :

Since  $\theta' = \theta$  so we are done

### 3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-P0)} \end{aligned}$$

IH1:

$$\begin{aligned} & \forall n_1, H_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH1 with  $n, H$  and since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-P1})
\end{aligned}$$

From evaluation rule we know that  $v'_1 = \lambda x. e_i$ . Since from FU-P1 we know that

$$(\theta'_1, n - i, \lambda x. e_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < (n - i). \forall v. (\theta'', j, v) \in [\tau_1 \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma \quad (72)$$

IH2:

$$\begin{aligned}
& \forall n_2, \forall H_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall k < n_2. (H_2, (e_2) \delta) \Downarrow_k (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - k, v'_2) \in [(\tau_1) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH2 with  $n - i, H'_1$  and since we know that  $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - k, v'_2) \in [(\tau_1) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \sigma) \quad (\text{FU-P2})
\end{aligned}$$

Instantiating  $\theta'', j$  and  $v$  in Equation 72 with  $\theta'_2, n - i - k$  and  $v'_2$  from FU-P2 respectively, we get

$$(\theta'_2, n - i - k, e_i[v'_2/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma$$

This means from Definition 2.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - k, H_3) \triangleright \theta'_2 \wedge \forall l < (n - i - k). (H_3, e_i[v'_2/x]) \Downarrow_l (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \sigma)
\end{aligned}$$

Instantiating  $H_3$  with  $H'_2$  from FU-P2 and since we know that  $((n - i - k), H'_2) \triangleright \theta'_2$  and since the reduction happens therefore we have

$$\begin{aligned}
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-P3})
\end{aligned}$$

In order to prove FU-P0 we choose  $\theta'$  as  $\theta'_3$  from FU-P3. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + l$ . Now we are required to show



- (a)  $\theta \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$ :
- $\theta \sqsubseteq \theta'_3$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-P1,  $\theta'_1 \sqsubseteq \theta'_2$  from FU-P2 and  $\theta'_2 \sqsubseteq \theta'_3$  from FU-P3 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_3$
  - $((n - i - k - l), H'_3) \triangleright \theta'_3$ :  
From FU-P3 we get  $((n - i - k - l), H'_3) \triangleright \theta'_3$
  - $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$ :  
From FU-P3 we get  $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-P1, FU-P2 and FU-P3
- (c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta). \theta'_3(a) \searrow pc \sigma)$   
Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-P1, FU-P2 and FU-P3

#### 4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  s.t  $H \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-PA0)} \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

We instantiate IH1 with  $H$  and  $n$ . And since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-PA1)} \end{aligned}$$

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau_2) \sigma \rrbracket_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma)
\end{aligned}$$

We instantiate IH2 with  $H'_1$  and  $n - i$ . And since we know that  $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1, e_2) \ \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma) \quad \text{(FU-PA2)}
\end{aligned}$$

In order to prove FU-PA0 we choose  $\theta'$  as  $\theta'_2$  from FU-PA2. Also we know from the evaluation rule, that let  $v' = (v'_1, v'_2)$ ,  $H' = H'_2$  and  $n' = i + j + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v') \in [(\tau_1 \times \tau_2)^\perp]_V$ :
- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-PA1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-PA2 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_2$
  - $(n - i - j - 1, H'_2) \triangleright \theta'_2$ :  
From FU-PA2 we get  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 2.20 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$
  - $(\theta'_2, n - i - j, v') \in [(\tau_1 \times \tau_2)^\perp]_V$ :  
From Definition 2.6 it suffices to show
    - i.  $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$ :  
Since from FU-PA1 we know that  $(\theta'_1, n - i, v'_1) \in [(\tau_1) \ \sigma]_V$  and since  $\theta'_1 \sqsubseteq \theta'_2$  (from FU-PA2) therefore from Lemma 2.16 we get  $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$
    - ii.  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$ :  
From FU-PA2 we know that  $(\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V$  therefore from Lemma 2.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$   
From FU-PA1 and FU-PA2
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma)$   
From FU-PA1 and FU-PA2

5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove:  $(\theta, n, \text{fst}(e_i) \ \delta) \in [\tau_1 \ \sigma]_E^{pc \ \sigma}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{fst}(e_i) \ \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \ \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma)
\end{aligned}$$

This means that given some heap  $H$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{fst}(e_i) \ \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-F0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-F1}) \end{aligned}$$

From evaluation rule we know that  $v'_1 = (v''_1, v''_2)$

In order to prove FU-F0 we choose  $\theta'$  as  $\theta'_1$  from FU-P1. Also we know that  $H' = H'_1$  and  $v' = v''_1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v'_1) \in [\tau_1 \sigma]_V$ :
  - $\theta \sqsubseteq \theta'_1$ :  
From FU-F1
  - $(n - i - 1, H'_1) \triangleright \theta'_1$ :  
From FU-F1 we know  $(n - i, H'_1) \triangleright \theta'_1$  therefore from Lemma 2.20 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$
  - $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$ :  
Since from FU-F1 we know that  $(\theta'_1, n - i, (v''_1, v''_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$   
Therefore from Definition 2.6 we know that  $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$   
From Lemma 2.16 we get  $(\theta'_1, n - i - 1, v''_1) \in [\tau_1 \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
From FU-F1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$   
From FU-F1

6. FG-snd:

Symmetric case to FG-fst

7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(\theta, n, \text{inl}(e_i) \delta) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^{pc \sigma}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-LE0)}
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-LE1)}
\end{aligned}$$

In order to prove FU-LE0 we choose  $\theta'$  as  $\theta'_1$  from FU-LE1. Also we know from the evaluation rule, that let  $v' = \text{inl}(v'_1)$ ,  $H' = H'_1$  and  $n' = i + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-LE1
  - $(n - i - 1, H') \triangleright \theta'_1$ :  
From FU-LE1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 2.20 we get  $(n - i - 1, H') \triangleright \theta'_1$
  - $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$ :  
Since  $v' = \text{inl}(v'_1)$  and from FU-LE1 we know that  $(\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V$   
Therefore from Definition 2.6 we get  $(\theta'_1, n - i, v') \in [(\tau_1 + \tau_2) \sigma]_V$   
From Lemma 2.16 we get  $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
From FU-LE1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$   
From FU-LE1

8. FG-inr:

Symmetric case to FG-inl

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau}$$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t.  $(n, H) \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-C0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_c) \delta) \Downarrow_i (H'_1, v'_c) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-C1}) \end{aligned}$$

2 cases arise:

(a)  $v'_c = \text{inl}(v_{ci})$ :

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_1) \delta \cup \{x \mapsto v_{ci}\}) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Instantiating IH2 with  $H'_1$  and  $n-i$  since we know that  $H'_1 \triangleright \theta'_1 \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{FU-C2}) \end{aligned}$$

In order to prove FU-C0 we choose  $\theta'$  as  $\theta'_2$  from FU-C2. Also we know that  $H' = H'_2$ ,  $v' = v'_2$  and  $n' = i + j + 1$ . Now we are required to show

- i.  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$ :
- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-C1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-C2 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_2$
  - $(n - i - j - 1, H'_2) \triangleright \theta'_2$ :  
From FU-C2 we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 2.20 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$
  - $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$ :  
From FU-C2 we know that  $(\theta'_2, n - i - j, v'_2) \in \lfloor \tau \sigma \rfloor_V$  therefore from Lemma 2.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$
- ii.  $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$ :  
Since from FU-C2 we know that  
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell')$   
therefore we also have  
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$   
and from FU-C1 we know that  
 $(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$   
Combining the two we get  
 $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
- iii.  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$ :  
Since from FU-C2 we know that  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$   
therefore we also have  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc) \sigma)$   
and from FU-C1 we know that  
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma)$   
Combining the two we get  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$
- (b)  $v'_c = \text{inr}(v_{ci})$ :  
Symmetric case as  $v'_c = \text{inl}(v_{ci})$

10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(\theta, n, \text{new } (e_i) \delta) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\text{ref } \tau)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\mathbf{ref} \ \tau)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma) \quad (\text{FU-R0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \ \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \mathbf{new} \ (e_i) \ \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma) \quad (\text{FU-R1})
\end{aligned}$$

From the evaluation rule we know that  $H' = H'_1[a \mapsto v'_1]$  where  $a \notin \text{dom}(H'_1)$ ,  $v' = a$  and  $n' = i + 1$ . In order to prove FU-R0 we choose  $\theta'$  as  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$ . Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_2 \wedge (n - i - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - 1, v') \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V:$$

- $\theta \sqsubseteq \theta'_2$ :

From FU-R1 we know that  $\theta \sqsubseteq \theta'_1$  therefore from Definition 2.2  $\theta \sqsubseteq \theta'_2$

- $(n - i - 1, H') \triangleright \theta'_2$ :

From FU-R1 we know that  $(n - i, H'_1) \triangleright \theta'_1$ . Therefore from Lemma 2.20 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$ .

We also know that  $(\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V$  (from FU-R1) therefore from Lemma 2.16 we get  $(\theta'_1, n - i - 1, v'_1) \in [\tau \ \sigma]_V$

Since  $H' = H'_1[a \mapsto v'_1]$  and  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$  therefore from Definition 2.8 we get  $(n - i - 1, H') \triangleright \theta'_2$

- $(\theta'_2, n - i - 1, a) \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V$ :

Since  $\theta'_2(a) = \tau \ \sigma$  therefore from Definition 2.6 we get  $(\theta'_2, n - i - 1, a) \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$$

From FU-R1

$$(c) \ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma):$$

We get this from FU-R1 and  $\tau \ \sigma \searrow pc \ \sigma$  (given)

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\mathbf{ref} \ \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(\theta, n, (!e_i) \ \delta) \in [\tau' \ \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (!e_i) \ \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \ \sigma]_V \wedge
\end{aligned}$$

$$\begin{aligned}
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t.  $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-D0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-D1})
\end{aligned}$$

In order to prove FU-D0 we choose  $\theta'$  as  $\theta'_1$  from FU-D1. Also we know from the evaluation rule, that  $H' = H'_1$ ,  $v' = H'_1(a)$ ,  $v'_1 = a$  and  $n' = i + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-D1
  - $(n - i - 1, H') \triangleright \theta'_1$ :  
From FU-D1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 2.20 we get  $(n - i - 1, H') \triangleright \theta'_1$
  - $(\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$ :  
Since from FU-D1 we know that  $(n - i, H'_1) \triangleright \theta'_1$  therefore from the Definition 2.8 we get  $(\theta'_1, n - i, H'_1(a)) \in [\tau \sigma]_V$   
From Lemma 2.16 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau \sigma]_V$   
Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.24 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
From FU-D1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$   
From FU-D1

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}}$$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [\text{unit } \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove



$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\mathbf{ref} \ \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\mathbf{ref} \ \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-A1})
\end{aligned}$$

IH2:

$$\begin{aligned}
& \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq (n_2 - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with  $H'_1$  and since we know that  $H'_1 \triangleright \theta'_1 \wedge (H, (e_1 := e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq (n - i - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-A2})
\end{aligned}$$

In order to prove FU-A0 we choose  $\theta'$  as  $\theta'_2$  from FU-A2. Also we know from the evaluation rule, assign, that let  $v'_1 = a_1$ ,  $H' = H'_2[a_1 \mapsto v'_2]$ ,  $v' = ()$  and  $n' = i + j + 1$ . Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, ()) \in [\mathbf{unit}]_V:$$

- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-A1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-A2 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_2$
- $(n - i - j - 1, H') \triangleright \theta'_2$ :  
From Definition 2.8 it suffices to prove that
  - i.  $\text{dom}(\theta'_2) \subseteq \text{dom}(H')$ : From FU-A2
  - ii.  $\forall a \in \text{dom}(\theta'_2). (\theta'_2, n - i - j - 1, H'(a)) \in [\theta'_2(a)]_V$ :  
 $\forall a \in \text{dom}(\theta'_2).$

- $a = a_1$ :  
From FU-A2 (since we know that  $(\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$ )  
Therefore from Lemma 2.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$
  - $a \neq a_1$ :  
From FU-A2 (since we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 2.20  
we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$ )
  - $(\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$ :  
From Definition 2.6
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
 $\forall a \in \text{dom}(H)$ .
- $a = a_1$ :  
Since we know that  $H(a_1) \neq H'(a_1)$  and  $\theta(a_1) = \tau = \mathbf{A}^{\ell_i}$  (given)  
It is given that  $\tau \sigma \searrow pc \sigma$  therefore  $pc \sigma \sqsubseteq \ell_i \sigma$
  - $a \neq a_1$ :  
From FU-A2
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$   
 From FU-A2

### 13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove:  $(\theta, n, (\Lambda e_i) \delta) \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\Lambda e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\Lambda e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-FI0}) \end{aligned}$$

IH1:

$$\forall n_1, \theta_i, \ell' \in \mathcal{L}. (\theta_i, n_1, e_i \delta) \in \llbracket \tau \sigma \cup \{\alpha \mapsto \ell''\} \rrbracket_E^{\ell_e \sigma \cup \{\alpha \mapsto \ell''\}}$$

In order to prove FU-FI0 we choose  $\theta'$  as  $\theta$ . Also we know from the evaluation rule, that  $H' = H$  and  $n' = 0$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V$ :
- $\theta \sqsubseteq \theta$ : From Definition 2.2
  - $(n, H) \triangleright \theta$ : Given

- $(\theta, n, (\Lambda e_i)\delta) \in [(\forall\alpha.(\ell_e, \tau))^\perp]_V \sigma$ :

From Definition 2.6 it suffices to prove that

$$\forall\theta''.\theta \sqsubseteq \theta'' \wedge \forall j < n.\forall\ell_d \in \mathcal{L} \implies (\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$$

This means given some  $\theta'', j$  and  $\ell_d$  such that  $\theta \sqsubseteq \theta'', j < n$  and  $\ell_d \in \mathcal{L}$

It suffices to prove that  $(\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$

Instantiating IH1 with  $j, \theta''$  and  $\ell_d$  we get  $(\theta_i, j, e_i \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell_d\}]_E^{\ell_e \sigma \cup \{\alpha \mapsto \ell_d\}}$

- (b)  $(\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $H' = H$  so we are done

- (c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta(a) \searrow pc)$ :

Since  $\theta' = \theta$  so we are done

#### 14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\forall\alpha.(\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i [] : \tau[\ell''/\alpha]}$$

To prove:  $(\theta, n, (e_i[]) \delta) \in [\tau[\ell''/\alpha] \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n.(n, H) \triangleright \theta \wedge \forall n' < n.(H, (e_i[]) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \quad (\text{FU-FE0}) \end{aligned}$$

IH:

$$\begin{aligned} & \forall H_1, n_1.(n_1, H_1) \triangleright \theta \wedge \forall i < n_1.(H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \quad (\text{FU-FE1}) \end{aligned}$$

From evaluation rule we know that  $v'_1 = \Lambda e_{i1}$ . Since from FU-FE1 we know that

$$(\theta'_1, n - i, \Lambda e_{i1}) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta'' . \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i . \forall \ell_g \in \mathcal{L} \implies (\theta'' , j , e_{i1}) \in [\tau[\ell_g/\alpha] \sigma]_E^{\ell_e[\ell_g/\alpha]} \sigma \quad (73)$$

Instantiating Equation 73 with  $\theta'_1$ ,  $n - i - 1$  and  $\ell''$  we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha]} \sigma$$

This means from Definition 2.7 we have

$$\begin{aligned} & \forall H_3 . (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1 . (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\ & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \end{aligned}$$

Instantiating  $H_3$  with  $H'_1$  from FU-FE1 and since we know that  $(n - i - 1, H'_1) \triangleright \theta'_1$  (Lemma 2.20) and since we know that  $e_i \Downarrow \gamma \downarrow_1$  reduces in  $n'$  steps where  $n' = i + k + 1$  and since  $n' < n$  therefore we have  $k < n - i - 1$  s.t.  $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$ . Therefore we get

$$\begin{aligned} & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-FE2}) \end{aligned}$$

In order to prove FU-FE0 we choose  $\theta'$  as  $\theta'_3$  from FU-FE2. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$ :
- $\theta \sqsubseteq \theta'_3$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-FE1 and  $\theta'_1 \sqsubseteq \theta'_3$  from FU-FE2 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_3$
  - $(n - i - k - 1, H'_3) \triangleright \theta'_3$ :  
From FU-FE2 we know that  $(n - i - k - 1, H'_3) \triangleright \theta'_3$
  - $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$ :  
From FU-FE2 we know that  $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H) . H(a) \neq H'_3(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
Since  $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  therefore we get the desired from FU-FE1 and FU-FE2
- (c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta) . \theta'_3(a) \searrow pc \sigma)$   
Since  $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  therefore we get the desired from FU-FE1 and FU-FE2

## 15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e_i : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove:  $(\theta, n, (\nu e_i) \delta) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n . (n, H) \triangleright \theta \wedge \forall n' < n . (H, (\nu e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta' . \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V \wedge \\ & (\forall a . H(a) \neq H'(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta) . \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\nu e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CI0}) \end{aligned}$$

IH1:

$$\forall \theta_i, n_1. (\theta_i, n_1, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma \text{ such that } \mathcal{L} \models c \sigma$$

In order to prove FU-FI0 we choose  $\theta'$  as  $\theta$ . Also we know from the evaluation rule, that  $H' = H$ ,  $v' = \nu e_i \delta$  and  $n' = 0$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in [(c \xrightarrow{\ell_e} \tau)^\perp]_V \sigma:$$

•  $\theta \sqsubseteq \theta$ : From Definition 2.2

•  $(n, H) \triangleright \theta$ : Given

•  $(\theta, n, (\nu e_i) \delta) \in [(c \xrightarrow{\ell_e} \tau)^\perp]_V \sigma$ :

From Definition 2.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \mathcal{L} \models c \sigma \implies (\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$$

This means given some  $\theta''$  such that  $\theta \sqsubseteq \theta''$ ,  $j < n$  and  $\mathcal{L} \models c$

It suffices to prove that  $(\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$

Instantiating IH1 with  $\theta''$  and  $j$  we get  $(\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$

$$(b) (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell'):$$

Since  $H' = H$  so we are done

$$(c) (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc):$$

Since  $\theta' = \theta$  so we are done

16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i \bullet : \tau}$$

To prove:  $(\theta, n, (e_i \bullet) \delta) \in [\tau \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v') \implies \\ \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CE0}) \end{aligned}$$

IH:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH with  $H$  and  $n$ . And since we know that  $(n, H) \triangleright \theta \wedge (H, (e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-CE1)}
\end{aligned}$$

From evaluation rule we know that  $v'_1 = \nu e_{i1}$ . Since from FU-CE1 we know that

$$(\theta'_1, n - i, \nu e_{i1}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i. \mathcal{L} \models c \sigma \implies (\theta'', j, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma} \quad (74)$$

Instantiating Equation 74 with  $\theta'_1$  and  $n - i - 1$  since we know that  $\mathcal{L} \models c \sigma$  therefore we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

This means from Definition 2.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1. (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma)
\end{aligned}$$

Instantiating  $H_3$  with  $H'_1$  from FU-CE1 and since we know that  $(n - i - 1, H'_1) \triangleright \theta'_1$  (Lemma 2.20) and since we know that  $e_i \bullet \gamma \Downarrow_1$  reduces in  $n'$  steps where  $n' = i + k + 1$  and since  $n' < n$  therefore we have  $k < n - i - 1$  s.t.  $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$ . Therefore we get

$$\begin{aligned}
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma) \quad \text{(FU-CE2)}
\end{aligned}$$

In order to prove FU-CE0 we choose  $\theta'$  as  $\theta'_3$  from FU-CE2. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + 1$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V:$$

$$\bullet \theta \sqsubseteq \theta'_3:$$

Since  $\theta \sqsubseteq \theta'_1$  from FU-CE1 and  $\theta'_1 \sqsubseteq \theta'_3$  from FU-CE2 therefore from Definition 2.2 we get  $\theta \sqsubseteq \theta'_3$

$$\bullet (n - i - k - 1, H'_3) \triangleright \theta'_3:$$

From FU-CE3 we know that  $(n - i - k - 1, H'_3) \triangleright \theta'_3$

$$\bullet (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V:$$

From FU-CE3 we know that  $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$

$$(b) (\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$$

Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-CE1 and FU-CE2

(c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta)).\theta'_3(a) \searrow pc \sigma$

Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-CE1 and FU-CE2

□

**Lemma 2.23** (FG: Expression subtyping with closed labels and types).  $\forall pc, pc', \tau$ .

$$\mathcal{L} \models pc \sqsubseteq pc' \implies \lfloor \tau \rfloor_E^{pc'} \subseteq \lfloor \tau \rfloor_E^{pc}$$

*Proof.* Given:  $\mathcal{L} \models pc \sqsubseteq pc'$

$$\text{To prove: } \lfloor (\tau) \rfloor_E^{pc'} \subseteq \lfloor (\tau) \rfloor_E^{pc}$$

This means we need to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc'} . (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc}$$

This means given  $\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc'}$

It suffices to prove that  $(\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc}$

From Definition 2.7 for the chosen  $\theta, n, e$  we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc' \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall k < n.(H_1, e) \Downarrow_k (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta)).\theta'_1(a) \searrow pc \end{aligned}$$

This means that we are given some  $H_1$  and  $k$  such that  $(n, H_1) \triangleright \theta$ ,  $k < n$  and  $(H_1, e) \Downarrow_k (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta)).\theta'_1(a) \searrow pc \end{aligned}$$

Instantiate  $H$  in (A) with  $H_1$  and then we choose  $\theta'_1$  as  $\theta'$

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \lfloor \tau \rfloor_V$ :

Given

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $pc \sqsubseteq pc'$  and we are given

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell')$$

Therefore

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc$ :

We are given

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc'$$

and since  $pc \sqsubseteq pc'$  Therefore

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc$$

□

**Lemma 2.24** (FG: Subtyping unary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1.  $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (A \sigma) \rrbracket_V \subseteq \llbracket (A' \sigma) \rrbracket_V$$

2.  $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_V \subseteq \llbracket (\tau' \sigma) \rrbracket_V$$

$$(b) \forall pc. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{pc}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

IH1:  $\llbracket (\tau'_1 \sigma) \rrbracket_V \subseteq \llbracket (\tau_1 \sigma) \rrbracket_V$  (Statement 2(a))

IH2:  $\forall pc. \llbracket (\tau_2 \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_E^{pc}$  (Statement 2(b))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e \sigma} \quad (75)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e \sigma}$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e \sigma}$

Since  $(\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in \llbracket \tau_1 \sigma \rrbracket_V$

As a result from Equation 75 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e \sigma}$$



From IH2, we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell_e \sigma}$$

Since  $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $[\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V \subseteq [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

IH1:  $[\lceil\tau_1 \sigma\rceil]_V \subseteq [\lceil\tau'_1 \sigma\rceil]_V$  (Statement 2(a))

IH2:  $[\lceil\tau_2 \sigma\rceil]_V \subseteq [\lceil\tau'_2 \sigma\rceil]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, (v_1, v_2)) \in [\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V. (\theta, n, (v_1, v_2)) \in [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V$

Therefore from Definition 2.6 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \quad (76)$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

Again from Definition 2.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 76 we know that  $(\theta, n, v_1) \in [\tau_1 \sigma]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2 \sigma]_V$  from Equation 76 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

## 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[\lceil(\tau_1 + \tau_2) \sigma\rceil]_V \subseteq [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

IH1:  $[\lceil\tau_1 \sigma\rceil]_V \subseteq [\lceil\tau'_1 \sigma\rceil]_V$  (Statement 2(a))

IH2:  $[\lceil\tau_2 \sigma\rceil]_V \subseteq [\lceil\tau'_2 \sigma\rceil]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, v_s) \in [\lceil(\tau_1 + \tau_2) \sigma\rceil]_V. (\theta, v_s) \in [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

This means that given:  $(\theta, n, v_s) \in [\lceil(\tau_1 + \tau_2) \sigma\rceil]_V$

And it suffices to prove:  $(\theta, n, v_s) \in [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_1 \sigma \rfloor_V \quad (77)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

From Equation 77 and IH1 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_2 \sigma \rfloor_V \quad (78)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

From Equation 78 and IH2 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

#### 4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $\lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V \subseteq \lfloor (\forall \alpha. (\ell'_e, \tau_2)) \sigma \rfloor_V$

IH1:  $\forall pc. \lfloor (\tau_1 \sigma) \rfloor_E^{pc} \subseteq \lfloor (\tau_2 \sigma) \rfloor_E^{pc}$  (Statement 2(b))

It suffices to prove:  $\forall (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V. (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

This means that given:  $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])} \quad (79)$$

And it suffices to prove:  $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

This means that given some  $\theta_2, j < n, \ell' \in \mathcal{L}$  s.t  $\theta \sqsubseteq \theta_2$

And we are required to prove:  $(\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$  therefore from Equation 79 we have

$$(\theta_2, j, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

Since  $\mathcal{L} \models \ell'_e (\sigma \cup [\alpha \mapsto \ell']) \sqsubseteq \ell_e (\sigma \cup [\alpha \mapsto \ell'])$  therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell'_e (\sigma \cup [\alpha \mapsto \ell'])}$$

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

$$\text{To prove: } [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V \subseteq [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{IH1: } \forall pc. [(\tau_1 \sigma)]_E^{pc} \subseteq [(\tau_2 \sigma)]_E^{pc} \text{ (Statement 2(b))}$$

$$\text{It suffices to prove: } \forall (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V. (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{This means that given: } (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V$$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma} \quad (80)$$

$$\text{And it suffices to prove: } (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

This means that given some  $\theta_2, j$  s.t  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

$$\text{And we are required to prove: } (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$  therefore from Equation 80 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell_e \sigma}$$

Since  $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

$$\text{To prove: } [((\text{ref } \tau) \sigma)]_V \subseteq [((\text{ref } \tau) \sigma)]_V$$

$$\text{It suffices to prove: } \forall (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V. (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V$$

Trivial

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{b}) \sigma) \rrbracket_V$

Directly from Definition 2.6

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit}) \sigma) \rrbracket_V \subseteq \llbracket ((\text{unit}) \sigma) \rrbracket_V$

Directly from Definition 2.6

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^\ell <: \mathbf{A}^{\ell'}} \text{FGsub-label}$$

To prove:  $\llbracket ((\mathbf{A}^\ell) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}^{\ell'}) \sigma) \rrbracket_V$

From Definition 2.6 it suffices to prove:  $\llbracket ((\mathbf{A}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}') \sigma) \rrbracket_V$

This we get directly from IH (Statement 1(a))

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\llbracket (\tau \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}. (\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$$

This means given  $(\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}$

It suffices to prove that  $(\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

From Definition 2.7 we know we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall i < n.(H, e) \Downarrow_i (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - i, H') \triangleright \theta' \wedge (\theta', n - i, v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall j < n.(H_1, e) \Downarrow_j (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some  $H_1$  and  $j < n$  s.t  $(n, H_1) \triangleright \theta \wedge (H_1, e) \Downarrow_j (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n-j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n-j, v') \in [\tau' \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate  $H$  in (A) with  $H_1$  and  $i$  with  $j$  then we choose  $\theta'_1$  as  $\theta'$   
Also we have IH1 as  $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$  (Statement 2(a))

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$ :  
We are given  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau \sigma]_V$   
From IH1 we know that  $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$   
Therefore,  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$
- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :  
Given
- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$ :  
Given

□

**Lemma 2.25** (FG: Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$   
 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 2.14 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case  $i = 1$

Given some  $m$  we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i)$ :  
 $\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i)$   
Therefore,  $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i))$  (Given)
- $\forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$ :  
We are given:  $\forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$   
Therefore from Lemma 2.15 we know that  
 $\forall m'. (W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$   
Instantiating  $m'$  with  $m$  we get  
 $(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 2$

Symmetric case as  $i = 1$

□

**Theorem 2.26** (FG: Fundamental theorem binary).  $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}} &\implies \\ (W, n, e (\gamma \downarrow_1), e (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}} & \end{aligned}$$

*Proof.* Proof by induction on the typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove:  $(W, n, x (\gamma \downarrow_1), x (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}}$

Say  $e_1 = x (\gamma \downarrow_1)$  and  $e_2 = x (\gamma \downarrow_2)$

From Definition of  $[\tau]_E^{\mathcal{A}}$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge \forall j < n. (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) &\implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} & \end{aligned}$$

This means given some  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove:  $\exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Here

- $H'_1 = H_1$  and  $H'_2 = H_2$
- $e_1 = v'_1 = \gamma(x) \downarrow_1$
- $e_2 = v'_2 = \gamma(x) \downarrow_2$
- $j = 1$

We choose  $W' = W$ .

- $W \sqsupseteq W$ : From Definition 2.3

- $(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$ :

Since we know that  $(n, H_1, H_2) \triangleright^{\mathcal{A}} W$  therefore from Lemma 2.21 we get

$$(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$$

- $(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$ :

We are given that  $(W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$  therefore from Lemma 2.19 we get

$$(W, n - 1, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$$

which means from Definition 2.14 we have

$$(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$$

2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_E^A$

Say  $e_1 = \lambda x. e (\gamma \downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \end{aligned}$$

This means that given  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

It suffices to prove:

$$\exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \quad (\text{FB-L0})$$

IH1:

$$\forall W, n. (W, n, e (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (v_1, v_2)) \in [\tau_1 \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \lambda x. e (\gamma \downarrow_1)$ ,  $v'_2 = e_2 = \lambda x. e (\gamma \downarrow_2)$  and  $j = 0$ . In order to prove FB-L0 we choose  $W' = W$  and we need to prove the following:

- $W \sqsubseteq W$ : From Definition 2.3
- $(n, H_1, H_2) \triangleright^A W$ : Given
- $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A$

From Definition 2.4 it suffices to prove that:

$$\forall W'' \sqsupseteq W, k < n, v_1, v_2.$$

$$((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, k, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$

This means that we need to prove the following:

$$\begin{aligned} - \forall W'' \sqsupseteq W, k < n, v_1, v_2. ((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A): \end{aligned}$$

This means given  $W'' \sqsupseteq W, k < n, v_1, v_2$  s.t  $((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove:  $(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

We instantiate IH1 with  $W''$  and  $k$   
 And since  $(W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$  therefore we get  
 $(W'', k, e(\gamma \downarrow_1)[v_1/x], e(\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

–  $\forall \theta_l \sqsupseteq W.\theta_1, k, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$   
 $(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$ :

This means that we are given  $\theta_l, k$  and  $v_c$  s.t

$\theta_l \sqsupseteq W.\theta_1$  and  $(\theta_l, k, v_c) \in [\tau_1 \sigma]_V$

And we are required to prove:

$(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

It is given to us that

$\forall v_1, v_2. (W, n, \gamma \in [\Gamma]_V^A$

Therefore from Lemma 2.25 we know that

$\forall m. (W.\theta_1, m, (\gamma \downarrow_1) \in [\Gamma]_V$

Therefore, we can apply Theorem 2.22 to obtain

$\forall m. (W.\theta_1, m, \lambda x.e \gamma \downarrow_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V$

From Definition 2.6 it means that we have

$\forall m. \forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1 \sigma]_V \implies$

$(\theta', j, e[v/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

We instantiate  $m$  with some  $l > k$ ,  $\theta'$  with  $\theta_l$ ,  $j$  with  $k$  and  $v$  with  $v_c$  to get

$W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since we show that  $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V$  therefore we get

$(\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

–  $\forall \theta_l \sqsupseteq W.\theta_2, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$

$(\theta_l, k, e(\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$ :

Symmetric case as above

### 3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2)(\gamma \downarrow_1), (e_1 e_2)(\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$\forall H_1, H_2, n' < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow$   
 $(H'_2, v'_2) \implies$

$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$

This further means that given  $H_1, H_2, n' < n$  s.t

$(n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$

It suffices to prove



$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_{\mathcal{V}}^A \quad (\text{FB-A0})$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}, i < n. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge (H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$ .  $(H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A \quad (81)$$

$$\underline{\text{IH2}}: (W'_1, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{j1}, H_{j2}, j < (n - i). (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge (H_{j1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{j1}$  with  $H'_1$  and  $H_{j2}$  with  $H'_2$  in IH2. Since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Also,  $e_1$  reduces to value  $\gamma \downarrow_1$  in  $i < n'$  steps. Therefore  $\exists j < n' - i < n - i$  s.t  $(H_{i1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1})$ .  $(H_{i2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A \quad (82)$$

We case analyze on  $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$  from Equation 81

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_{\mathcal{V}}^A$$

$$\text{Let } v'_1 = \lambda x.e_{h1} \text{ and } v'_2 = \lambda x.e_{h2}$$

Again from Definition 2.4 it means that

$$\forall W'_{h1} \sqsupseteq W'_1, j_1 < (n - i), v_1, v_2.$$

$$((W'_{h1}, j_1, v_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}}^A \implies (W'_{h1}, j_1, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_1, m_1, v_c.$$

$$\wedge ((\theta_{l1}, m_1, v_1) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_1, e_{h1}[v_1/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_2, m_1, v_c.$$

$$\wedge (\theta_{l1}, m_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_2, e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$

We instantiate  $W'_{h1}$  with  $W'_2$  obtained from Equation 82. Similarly we also instantiate  $v_1$  and  $v_2$  with  $v'_{j1}$  and  $v'_{j2}$  respectively from Equation 82, and  $j_1$  with  $n - i - j$ . And we get

$$(W'_2, n - i - j, e_{h1}[v'_{j1}/x], e_{h2}[v'_{j2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 2.5 we get

$$\begin{aligned} & \forall H_1, H_2, k_e < (n - i - j). (n - i - j, H_1, H_2) \triangleright^A W'_2 \wedge \\ & (H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1}) \wedge (H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2}) \implies \\ & \exists W' \sqsupseteq W'_2. (n - i - j - k_e, H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', n - i - j - k_e, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \end{aligned}$$

Instantiating  $H_1$  with  $H'_{j1}$  and  $H_2$  with  $H'_{j2}$  obtained from Equation 82. And since we know that  $e_1 e_2$  reduces with  $\gamma \downarrow_1$  in  $n' < n$  steps. And  $e_2$  reduces to value  $\gamma \downarrow_1$  in  $j < n' - 1 < n - i$  steps. Therefore  $\exists k_e = n' - i - j < n - i - j$  s.t  $(H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1})$ .  $(H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W' \sqsupseteq W'_2. ((n - i - j - k_e), H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', (n - i - j - k_e), v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \quad (83)$$

This concludes the proof in this case.

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From FB-A0 we know that we need to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau_2 \sigma = A^{\ell_i}$  and since  $\tau_2 \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V \quad (84)$$

In this case from Definition 2.6 we know that

$$\forall m. (W'_1.\theta_1, m, \lambda x. e_{h1}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (85)$$

$$\forall m. (W'_1.\theta_2, m, \lambda x. e_{h2}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (86)$$

Applying Definition 2.6 on Equation 85 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

where  $\theta = W'_1.\theta_1$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 1 + t_1). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC1})$$

Since from Equation 82 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 2.15 we get

$$\forall m. (W'_2.\theta_1, m, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(W'_2.\theta_1, m_1 + 1 + t_1, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_1$ ,  $j_1$  with  $m_1 + t_1$  and  $v$  with  $v'_{j_1}$  from Equation 82.

$$\text{Therefore we get } (W'_2.\theta_1, m_1 + 1 + t_1, e_{h1}[v'_{j_1}/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\forall H. (m_1 + 1 + t_1, H) \triangleright W'_2.\theta_1 \wedge \forall k_c < (m_1 + 1 + t_1). (H, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_1, v'_1) \implies$$

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1 + t_1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1 + t_1 - k_c), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

Since from Equation 82 we have  $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 2.27 we get  $\forall m. (m, H'_{j_1}) \triangleright W'_2.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get  $(m_1 + 1 + t_1, H'_{j_1}) \triangleright W'_2.\theta_1$

Now instantiating  $H$  with  $H'_{j_1}$  from Equation 82 and  $k_c$  with  $t_1$  we get

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{R1})$$

Similarly we can apply Definition 2.6 on Equation 86 to get

$$\forall m. \forall \theta'_2. (m, W'_1.\theta_2) \sqsubseteq \theta'_2 \wedge \forall j_2 < m. \forall v. (\theta'_2, j_2, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta'_2, j_2, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

We instantiate  $m$  with  $m_2 + 2 + t_2$  where  $t_2$  is the number of steps in which  $e_{h2}$  reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta', j_1, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC2})$$

Since from Equation 82 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 2.15 we get

$$\forall m. (W'_2.\theta_2, m, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_2$ ,  $j_1$  with  $m_2 + 1 + t_2$  and  $v$  with  $v'_{j_2}$  from Equation 82 in FB-AC2 we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, e_{h_2}[v'_{j_2}/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma$$

From Definition 2.7, we get

$$\begin{aligned} \forall H.(m_2 + 1 + t_2, H) \triangleright W'_2.\theta_2 \wedge \forall k_c < (m_2 + 1 + t_2).(H, e_{h_2}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_2, v'_2) \implies \\ \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1 + t_2 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1 + t_2 - k_c)v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 82 we have  $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 2.27 we get  $\forall m.(m, H'_{j_2}) \triangleright W'_2.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get  $(m_2 + 1 + t_2, H'_{j_2}) \triangleright W'_2.\theta_2$

Now Instantiating  $H$  with  $H'_{j_2}$  from Equation 82 and  $k_c$  with  $t_2$ .

$$\begin{aligned} \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{R2}) \end{aligned}$$

In order to prove FB-A0 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_2.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 2.9 it suffices to show that

–  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$ :

From R1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from R2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$ :

Since from Equation 82 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$  therefore from

Definition 2.9 we know that  $(W'_2.\hat{\beta}) \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  therefore

$(W'_2.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$

$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

4 cases arise for each  $(a_1, a_2) \in W'_2.\hat{\beta}$

i.  $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

We know from Equation 82 that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$

Therefore from Definition 2.9 we have

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

Since  $W'.\hat{\beta} = W'_2.\hat{\beta}$  by construction therefore

$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 2.2

$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From Equation 82 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$

This means from Definition 2.9 that

$$\forall (a_{i_1}, a_{i_2}) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_2 \sqsubseteq W'$  and  $n - n' - 1 < n - i - j - 1$  (since  $n' = i + j + t_1$  where  $t_1$  is the number of steps taken by  $e_{h_1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce and  $j$  is the number of steps taken by  $e_2 \gamma \downarrow_1$  to reduce) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii.  $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From R1 and R2 we know that

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_2.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_2.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from R1 and R2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ . Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$
 and

$$(\theta'_2, m_2, H'_2(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii.  $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From R2 we know that

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e \sigma$  in the world before the modification. Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 82 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$  that means from Definition 2.9 that  $(W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$ . Since  $(\ell_e \sigma) \sqsubseteq \ell'$  therefore from Definition 2.4 we know that  $H'_{j_1}(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_2.\theta_1, m, H'_{j_1}(a_1)) \in W'_2.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_2.\theta_2, m, H'_{j_2}(a_2)) \in W'_2.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{j_1}(a_1)) \in \theta'_1(a_1)$$

Since from R2 we know that  $(m_2+1, H'_2) \triangleright \theta'_2$  therefore from Definition 2.8 we know that  $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V:$$

$$\underline{i = 1}$$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we instantiate Equation 85 and Equation 86 with  $m + 2 + t_1$  and  $m + 2 + t_2$  respectively. This will give us

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

and

$$\exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma))$$

Since we have  $(m+1, H'_1) \triangleright \theta'_1$  and  $(m+1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 2.8

$$\underline{i = 2}$$

Symmetric to  $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau_2 \sigma]_V^A:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From R1 and R2 we and Definition 2.4 we get the desired.

#### 4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\perp \sigma]_E^A$

Say  $e_1 = (e_1, e_2) (\gamma \downarrow_1)$  and  $e_2 = (e_1, e_2) (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \times \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies$$

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$$

This means that given some  $H_1, H_2$  and  $n' < n$  s.t

$$(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{\mathcal{V}}^A \quad (87)$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{p11}, H_{p12}. (n, H_{p11}, H_{p12}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11}) \wedge (H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{p11}$  with  $H_1$  and  $H_{p22}$  with  $H_2$  in IH1 and since the  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12})$ . Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (88)$$

$$\underline{\text{IH2}} (W, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{p21}, H_{p22}. (n - i, H_{p21}, H_{p22}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22}) \implies$$

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{p21}$  with  $H'_{p11}$  and  $H_{p22}$  with  $H'_{p21}$  and in IH2. Since  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  has reduced with  $i < n'$  steps. Therefore we know that  $\exists j < n' - i < n - i$  s.t  $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$ . Hence we get

since the  $(e_1, e_2)$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$ . Hence we get

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (89)$$

In order to prove Equation 87 we instantiate  $W'$  in Equation 87 with  $W'_2$  we are required to show the following:

- $W \sqsubseteq W'_2$ :  
Since  $W \sqsubseteq W'_1$  from Equation 88 and  $W'_1 \sqsubseteq W'_2$  from Equation 89  
Therefore,  $W \sqsubseteq W'_2$  from Definition 2.3

- $(n - n', H'_1, H'_2) \overset{A}{\triangleright} W'$ :

Here  $n' = i + j + 1$

From evaluation rule of products we know that  $H'_1 = H'_{p21}$  and  $H'_2 = H'_{p22}$

From Equation 89 we know that  $(n - i - j, H'_{p21}, H'_{p22}) \overset{A}{\triangleright} W'_2$

Therefore from Lemma 2.21 we get  $(n - i - j - 1, H'_{p21}, H'_{p22}) \overset{A}{\triangleright} W'_2$

- $(W', n - i - j - 1, v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$ :

From evaluation rule of products we know that  $v'_1 = (v'_{p11}, v'_{p21})$  and  $v'_2 = (v'_{p12}, v'_{p22})$

We are required to show

$$- (W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A:$$

From Equation 88 and Equation 89 we know that

$$(W'_2, n - i - j, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

Therefore from Lemma 2.17 we get

$$(W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

## 5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove:  $(W, n, (\text{fst}(e_i)) (\gamma \downarrow_1), (\text{fst}(e_i)) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

Say  $e_1 = (\text{fst}(e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{fst}(e_i)) (\gamma \downarrow_2)$

From Definition 2.5 it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A \quad (90)$$

### IH1

$$(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \overset{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \overset{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{fst}(e_i)$  reduces to value reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{fst}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$ . Hence we get



$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A \quad (91)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$  from Equation 91

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V^A$$

Let  $v'_{i1} = (v_{i1}, v_{i2})$  and  $v'_{i2} = (v_{j1}, v_{j2})$

Again from Definition 2.4 it means that

$$(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A \wedge (W'_1, n - i, v_{i2}, v_{j2}) \in [\tau_2 \sigma]_V^A \quad (F1)$$

In order to prove Equation 90 we choose  $W'$  as  $W'_1$  and from the evaluation rule of fst we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . Also, from reduction rules we know that  $n' = i + 1$ . And then we need to show:

- $W \sqsubseteq W'_1$ :

Directly from Equation 91

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

Since from Equation 91 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 2.21 we get  $(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$

- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$ :

From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$

From F1 we know that  $(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 2.17 we get  $(W'_1, n - i - 1, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.6 we know that

$$(a) \forall m. (W'_1.\theta_1, m, v'_{i1}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V \text{ and}$$

$$(b) \forall m. (W'_1.\theta_2, m, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V$$

where

$$v'_{i1} = (v_{i1}, v_{i2}) \text{ and } v'_{i2} = (v_{j1}, v_{j2})$$

In order to prove Equation 90 we choose  $W'$  as  $W'_1$  and from the evaluation rule of fst we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . And then we need to show:

- $W \sqsubseteq W'_1$ :

Directly from Equation 91

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

From Equation 91 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 2.21 we get

$$(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$ :  
 From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$   
 Let  $\tau_1 = \mathbf{A}^{\ell_i}$  Since  $\tau_1 \sigma \searrow \ell$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \quad (92)$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \quad (93)$$

Instantiating (a) with  $m_1$

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_1, m_1, v_{i2}) \in [\tau_2 \sigma]_V \quad (94)$$

Instantiating (b) with  $m_2$

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_2, m_2, v_{j2}) \in [\tau_2 \sigma]_V \quad (95)$$

From Equation 94 and Equation 95 we get

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \text{ and } (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

## 6. FG-snd:

Symmetric case as FG-fst

## 7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(W, n, (\text{inl}(e_i))(\gamma \downarrow_1), (\text{inl}(e_i))(\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^A$

Say  $e_1 = (\text{inl}(e_i))(\gamma \downarrow_1)$  and  $e_2 = (\text{inl}(e_i))(\gamma \downarrow_2)$

From Definition of  $[(\tau_1 + \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A \quad (96)$$

IH1  $(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t.  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \quad (97)$$

Instantiating  $W'$  in Equation 96 with  $W'_1$ . Also from reduction relation we know that  $n' = i + 1$  we are required to show the following:

- $W \sqsubseteq W'_1$ :

Directly from Equation 97

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 97 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$$(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$ :

From evaluation rule of  $\text{inl}$  we know that  $v'_1 = \text{inl}(v'_{i1})$  and  $v'_2 = \text{inl}(v'_{i2})$

We are required to show

$$- (W'_1, n - n', v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A:$$

From Equation 97 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 2.17 we get

$$(W'_1, n - i - 1, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$$

8. FG-inr:

Symmetric case to FG-inl.

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{i1} : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_i, x.e_{i1}, y.e_{i2}) : \tau}$$

To prove:  $(W, (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1), (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

Say  $e_1 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1)$  and  $e_2 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ & \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A \quad (98)$$

$$\underline{\text{IH1}} \ (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_{\mathcal{V}}^A \quad (99)$$

IH2:

$$(W'_1, n - i, (e_{i1}) (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\}), (e_{i1}) (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{j1}, H_{j2}. (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_1, e_{i1} (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\})) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \Downarrow (H'_{j2}, v'_{j2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{j1}$  with  $H'_1$  and  $H_{j2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A \quad (100)$$

IH3:

$$(W'_1, n - i, (e_{i2}) (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\}), (e_{i2}) (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{k1}, H_{k2}. (n - i, H_{k1}, H_{k2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < n - i. (H_1, e_{i2} (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\})) \Downarrow_k (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \Downarrow (H'_{k2}, v'_{k2}) \implies$$

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{k1}$  with  $H'_1$  and  $H_{k2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i2}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{k2}, v'_{k2})$ . Hence we get

$$\exists W'_3 \sqsupseteq W'_1.(n - i - k, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau] \sigma \uparrow_V^A \quad (101)$$

We case analyze  $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 + \tau_2)^\ell \sigma] \uparrow_V^A$  from Equation 99

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 2.4 2 further cases arise:

- $v'_1 = \text{inl}(v_{i1})$  and  $v'_2 = \text{inl}(v_{i2})$ :

In this case from Definition 2.4 we know that  $(W, n - i, v_{i1}, v_{i2}) \in [\tau_1 \sigma] \uparrow_V^A$

In order to prove Equation 98 we choose  $W'$  as  $W'_2$  from Equation 100 and from the first evaluation rule of case we know that  $H'_1 = H'_{j1}$  and  $H'_2 = H'_{j2}$ . Also we know from the evaluation rule that  $n' = i + j + 1$ . And then we need to show:

- \*  $W \sqsubseteq W'_2$ :

Since  $W \sqsubseteq W'_1$  from Equation 99 and  $W'_1 \sqsubseteq W'_2$  from Equation 100

Therefore,  $W \sqsubseteq W'_2$  from Definition 2.3

- \*  $(n - n', H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$ :

From Equation 100 we know that  $(n - i - j, H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$

Therefore from Lemma 2.21 we get

$$(n - i - j - 1, H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$$

- \*  $(W'_2, n - n', v'_1, v'_2) \in [\tau \sigma] \uparrow_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{j1}$  and  $v'_2 = v'_{j2}$

From Equation 100 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [\tau \sigma] \uparrow_V^A$

Therefore from Lemma 2.17 we get

$$(W'_2, n - i - j - 1, v'_{j1}, v'_{j2}) \in [\tau \sigma] \uparrow_V^A$$

- $v'_1 = \text{inr}(v_{i1})$  and  $v'_2 = \text{inr}(v_{i2})$ :

In this case from Definition 2.4 we know that  $(W, v_{i1}, v_{i2}) \in [\tau_2 \sigma] \uparrow_V^A$

In order to prove Equation 98 we choose  $W'$  as  $W'_3$  from Equation 101 and from the second evaluation rule of case we know that  $H'_1 = H'_{k1}$  and  $H'_2 = H'_{k2}$ . Also we know from the evaluation rule that  $n' = i + k + 1$ . And then we need to show:

- \*  $W \sqsubseteq W'_3$ :

Since  $W \sqsubseteq W'_1$  from Equation 99 and  $W'_1 \sqsubseteq W'_3$  from Equation 101

Therefore,  $W \sqsubseteq W'_3$  from Definition 2.3

- \*  $(n - n', H'_1, H'_2) \overset{A}{\triangleright} W'_3$ :

From Equation 101 we know that  $(n - i - k, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3$

Therefore from Lemma 2.21 we get

$$(n - i - k - 1, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3$$

- \*  $(W'_3, n - n', v'_1, v'_2) \in [\tau \sigma] \uparrow_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{k1}$  and  $v'_2 = v'_{k2}$

From Equation 101 we know that  $(W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau \sigma] \uparrow_V^A$

Therefore from Lemma 2.17 we get

$$(W'_3, n - i - k - 1, v'_{k1}, v'_{k2}) \in [\tau \sigma] \uparrow_V^A$$

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

The following cases arise:

- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case1 :  
Exactly the same reasoning as in the  $v'_1 = \text{inl}(v_{i1})$  and  $v'_2 = \text{inl}(v_{i2})$  subcase of the  $\ell \sigma \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case2 :  
Exactly the same reasoning as in the  $v'_1 = \text{inr}(v_{i1})$  and  $v'_2 = \text{inr}(v_{i2})$  subcase of the  $\ell \sigma \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case2 :

From Equation 98 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau \sigma = A^{\ell_i}$  and since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (102)$$

Since we know that  $(W, n, \gamma) \in [\Gamma]_V^A$  (given) therefore from Lemma 2.25 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore by instantiating it at  $m_1 + 1 + j$  we know that

$$(W.\theta_1, m_1 + 1 + j, \gamma \downarrow_1) \in [\Gamma]_V \quad (103)$$

Next we apply Theorem 2.22 on  $e_{i1} \gamma \downarrow_1$ . Here  $j$  is the number of steps in which  $e_{i1} \gamma \downarrow_1$  reduces. We use  $\gamma \downarrow_1 \cup \{x \mapsto v'_{s1}\}$  as the unary substitution to get

$$(W.\theta_1, m_1 + 1 + j, e_{i1} \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$$

This means from Definition 2.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_1 + 1 + j, H_{c1}) \triangleright W_1.\theta_1 \wedge \forall l_c < (m_1 + 1 + j). (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1 + j - l_c, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1 + j - l_c, v'_c) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W_1.\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since from Equation 99 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 2.27 we get  $\forall m. (m, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + j$  we get  $(m_1 + 1 + j, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $H_{c2}$  with  $H'_1$  from Equation 99 and  $l_c$  with  $j$  we get  
 $\exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC1})$

Similarly we apply Theorem 2.22 on  $e_{i2} \gamma \downarrow_2$ . Here  $j_2$  is the number of steps in which  $e_{i2} \gamma \downarrow_2$  reduces. We use  $\gamma \downarrow_2 \cup \{y \mapsto v'_{s2}\}$  as the unary substitution to get  
 $(W_1.\theta_2, m_2 + 1 + j_2, e_{i2} \gamma \downarrow_1 \cup \{y \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$

This means from Definition 2.7 we get

$\forall H_{c2}. (m_2 + 1 + j_2, H_{c1}) \triangleright W_1.\theta_2 \wedge \forall l_c < m_2 + 1 + j_2. (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \Downarrow_{k_c}$   
 $(H'_{c2}, v'_c) \implies$   
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1 + j_2 - l_c, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1 + j_2 - l_c, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$

Since from Equation 99 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 2.27 we get  $\forall m. (m, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + j_2$  we get  $(m_2 + 1 + j_2, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $H_{c2}$  with  $H'_2$  (from Equation 99) and  $l_c$  with  $j_2$  to get  
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC2})$

We choose

$W_n.\theta_1 = \theta'_1$  (from CC1)  
 $W_n.\theta_2 = \theta'_2$  (from CC2)  
 $W_n.\hat{\beta} = W'_1.\hat{\beta}$  (from Equation 99)

In order to prove Equation 98 we choose  $W'$  as  $W_n$

i.  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 2.9 it suffices to show that

–  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$ :

From (CC1) we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from (CC2) we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$ :

Since from Equation 99 we have  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Definition 2.9 we get  $(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore  
 $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$   
 $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A$ :

4 cases arise for each  $a_1$  and  $a_2$

- A.  $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$ :  
 $\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$ :

We know from Equation 99 that  $(n - i, H'_1, H'_2) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall(a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 2.2

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From Equation 99 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$

This means from Definition 2.9 that

$$\forall(a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$  (since  $n' = i + t_1 + 1$  where  $t_1$  is the number of steps taken by  $e_{i1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

- B.  $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From (CC1) and (CC2) we know that

$$(\forall a. H'_1(a) \neq H'_{c1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from (CC1) and (CC2),  $(m_1 + 1, H'_{c1}) \triangleright \theta'_1$  and  $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$ .

Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_{c1}(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$

$$(\theta'_2, m_2, H'_{c2}(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 2.4 we get (here  $H'_1 = H'_{c1}$  and  $H'_2 = H'_{c2}$ )

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

- C.  $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$



From (CC2) we know that

$$(\forall a. H_2'(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1'.\theta_2(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $(pc \sqcup \ell) \sigma$  in the world before the modification. Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 99 we know that  $(n - i, H_1', H_2') \stackrel{\mathcal{A}}{\triangleright} W_1'$  that means from Definition 2.9 that  $(W_1', n - i - 1, H_1'(a_1), H_2'(a_2)) \in \lceil W_1'.\theta_1(a_1) \rceil_V^{\mathcal{A}}$ . Since  $((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$  therefore from Definition 2.4 we know that  $H_1'(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W_1'.\theta_1, m, H_1'(a_1)) \in W_1'.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W_1'.\theta_2, m, H_2'(a_2)) \in W_1'.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 2.16 we get

$$(\theta_1', m_1, H_1'(a_1)) \in \theta_1'(a_1)$$

Since from (CC2) we know that  $(m_2 + 1, H_{c2}') \triangleright \theta_2'$  therefore from Definition 2.8 we know that  $(\theta_2', m_2, H_{c2}'(a_2)) \in \theta_2'(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H_{c1}'(a_1), H_{c2}'(a_2)) \in \lceil \theta_1'(a_1) \rceil_V^{\mathcal{A}}$$

$$\text{D. } H_{j1}'(a_1) \neq H_1'(a_1) \vee H_{j2}'(a_2) = H_2'(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V:$$

$$\underline{i = 1}$$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V$$

Like before we apply Theorem 2.22 on  $e_{i1} \gamma_1$  and  $e_{i2} \gamma_2$  but this time using  $m + 1 + i$  and  $m + 1 + j$  where  $i$  and  $j$  are the number of steps in which  $e_{i1} \gamma_1$  and  $e_{i2} \gamma_2$  reduces respectively. This will give us

$$\begin{aligned} & \exists \theta_1'. W_1.\theta_1 \sqsubseteq \theta_1' \wedge (m + 1, H_{c2}') \triangleright \theta_1' \wedge (\theta_1', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_1') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta_2'. W_1.\theta_2 \sqsubseteq \theta_2' \wedge (m + 1, H_{c2}') \triangleright \theta_2' \wedge (\theta_2', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_2') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since we have  $(m + 1, H_{c1}') \triangleright \theta_1'$  and  $(m + 1, H_{c2}') \triangleright \theta_2'$  therefore we get the desired from Definition 2.8

$$\underline{i = 2}$$

Symmetric to  $i = 1$

$$\text{ii. } (W', n - n' - 1, v'_1, v'_2) \in \lceil \tau_2 \sigma \rceil_V^{\mathcal{A}}:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CC1 and CC2 we and Definition 2.4 we get the desired.

- (d) Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case1 :  
Symmetric case as before

10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow_{pc}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(W, (\text{new } (e_i)) (\gamma \downarrow_1), (\text{new } (e_i)) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\perp \sigma]_E^A$

Say  $e_1 = (\text{new } (e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{new } (e_i)) (\gamma \downarrow_2)$

From Definition of  $[(\text{ref } \tau)^\perp \sigma]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \quad (104)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow \\ (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{ref}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$ . s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since  $\text{ref}(e_i)$  reduces with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \quad (105)$$

From the evaluation rule of  $\text{ref}$  we know that  $H'_1 = H'_{i1} \cup \{a_{n1} \mapsto v_{i1}\}$  and  $H'_2 = H'_{i2} \cup \{a_{n2} \mapsto v_{i2}\}$

Inorder to prove Equation 104 we instantiate  $W'$  with  $W_n$  where  $W_n$  is

$$W_n.\theta_1 = W'_1.\theta_1 \cup \{a_{n1} \mapsto \tau\}$$

$$W_n.\theta_2 = W'_1.\theta_2 \cup \{a_{n2} \mapsto \tau\}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$$

Also we know that  $n' = i + 1$

We are now required to prove

- $W \sqsubseteq W_n$ :

From Equation 105 we know that  $W \sqsubseteq W'_1$  and  $W'_1 \sqsubseteq W_n$  by construction. Therefore from Definition 2.3,  $W \sqsubseteq W_n$

- $(n - n', H'_1, H'_2) \hat{\triangleright}^A W_n$ :

From Definition 2.9 it suffices to show that

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H'_2):$$

From Equation 105 and by construction of  $W_n$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_1)):$$

From Equation 105 and by construction of  $W_n$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, n - n', H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A:$$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

From Equation 105 and by construction of  $W_n$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A:$$

From Equation 105 since we know that  $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$  that means

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in [\![W'_1.\theta_1(a_1)]\!]_V^A$$

Therefore from Lemma 2.17 we get  $(n - i - 2 = n - n' - 1, \text{ since } n' = i + 1)$

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 2, H'_1(a_1), H'_2(a_2)) \in [\![W'_1.\theta_1(a_1)]\!]_V^A$$

Since  $W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$  and from Equation 105 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\![\tau \sigma]\!]_V^A$

Therefore combining the two we get

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A$$

$$- \forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V:$$

From Equation 105 we have  $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$  that means from Definition 2.9 we have

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W'_1.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V$$

Also from Equation 105 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\![\tau \sigma]\!]_V^A$

Therefore from Lemma 2.15 and Lemma 2.16 we get

$$\forall m. (W'_1.\theta_1, m, v'_{i1}) \in [\![\tau \sigma]\!]_V$$

and

$$\forall m. (W'_1.\theta_2, m, v'_{i2}) \in [\![\tau \sigma]\!]_V$$

Combining the two we get

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V$$

- $(W_n, n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A$ :

Here  $v'_1 = a_{n1}$  and  $v'_2 = a_{n2}$

Since  $(a_{n1}, a_{n2}) \in W_n$  and also  $W_n.\theta_1(a_{n1}) = W_n.\theta_1(a_{n1}) = \tau$

Therefore from Definition 2.4  $(W_n, v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A$

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(W, n, !(e_i)) (\gamma \downarrow_1), !(e_i) (\gamma \downarrow_2) \in [(\tau') \sigma]_E^A$

Say  $e_1 = !(e_i) (\gamma \downarrow_1)$  and  $e_2 = !(e_i) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \downarrow \\ (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A \quad (106)$$

$$\mathbf{IH1} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_2, v'_2) \implies \\ \exists W'_1 \sqsubseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $!(e_i)$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t.  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_1, v'_1)$ . Similarly since  $!e_i$  reduces to value with  $\gamma \downarrow_2$  therefore  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_2, v'_2)$ . Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A \quad (107)$$

We case analyze on  $(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A$  from Equation 107

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } (\tau \sigma))]_V^A$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 2.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{D1})$$

In order to prove Equation 106 we instantiate  $W'$  with  $W'_1$ . Also we know that  $n' = i + 1$

- $W'_1 \sqsubseteq W$ :

From Equation 107

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 107 we know that

$(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

–  $(W'_1, n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A$ :

From the evaluation rule of deref we know that  $v'_1 = H'_1(a_{i1})$  and  $v'_2 = H'_2(a_{i2})$

Since from Equation 107 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ , therefore from Definition 2.9 we know that

$(W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in [W'_1.\theta_1(a_{i1})]_V^A$

And from D1 we know that  $W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau$

Therefore  $(W'_1, v'_1, v'_2) \in [(\tau) \sigma]_V^A$

Since  $\tau \sigma <: \tau' \sigma$  Therefore from Lemma 2.28, we get

$(W'_1, n - i - 1, v'_1, v'_2) \in [(\tau') \sigma]_V^A$

• Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From the evaluation rule of deref we know that  $v'_{i1} = a_1$  and  $v'_{i2} = a_2$

In this case from Definition 2.4 we know that

$$\forall m_1. (W'_1.\theta_1, m_1, a_1) \in [(\text{ref } \tau) \sigma]_V \quad (108)$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, a_2) \in [(\text{ref } \tau) \sigma]_V \quad (109)$$

In order to prove Equation 106 we choose  $W'$  as  $W'_1$ . And then we need to show:

–  $W \sqsubseteq W'_1$ :

Directly from Equation 107

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 107 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

–  $(W'_1, n - n', v'_1, v'_2) \in [\tau' \sigma]_V^A$ :

Let  $\tau' = A^{\ell_i}$  Since  $\tau' \sigma \searrow_{\ell}$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it suffices to prove that

$\forall m_1. (W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V$

and

$\forall m_2. (W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V \quad (110)$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V \quad (111)$$

Since from Equation 107 we know that  $(n-i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 2.27 we get

$$\forall m_{h1}.(m_{h1}, H'_1) \triangleright W'_1.\theta_1 \quad (112)$$

$$\forall m_{h2}.(m_{h2}, H'_2) \triangleright W'_1.\theta_2 \quad (113)$$

Instantiating  $m_{h1}$  in Equation 112 with  $m_1 + 1$  we get  $(m_1, H'_1) \triangleright W'_1.\theta_1$

Therefore from Definition 2.8, we get

$$\forall a \in \text{dom}(W'_1.\theta_1).(W'_1.\theta_1, m_1, H'_1(a)) \in [W'_1.\theta_1(a)]_V$$

Instantiating  $a$  with  $a_1$  we get  $(W'_1.\theta_1, m_1, H'_1(a_1)) \in [W'_1.\theta_1(a)]_V$

Since  $W'_1.\theta_1(a_{i1}) = \tau$  therefore we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau \sigma]_V$$

and since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.24 we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V$$

Similarly we also get

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V$$

Finally from Definition 2.4 we get

$$(W'_1, v'_1, v'_2) \in [(\tau') \sigma]_V^A$$

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} := e_{i2} : \text{unit}}$$

To prove:  $(W, n, (e_{i1} := e_{i2}) (\gamma \downarrow_1), (e_{i1} := e_{i2}) (\gamma \downarrow_2)) \in [(\text{unit}) \sigma]_E^A$

Say  $e_1 = (e_{i1} := e_{i2}) (\gamma \downarrow_1)$  and  $e_2 = (e_{i1} := e_{i2}) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2.(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_V^A$$

This further means that given

$$\forall H_1, H_2.(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_V^A \quad (114)$$

IH1  $(W, n, (e_{i1}) (\gamma \downarrow_1), (e_{i1}) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\ell \sigma]_E^A$

This means from Definition 2.5 we get

$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_1) \wedge (H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_2) \implies$

$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_{i1} := e_{i2})$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{i1}, v'_{i1})$ . Similarly since  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_2$  therefore we also have  $(H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_V^A \quad (115)$$

IH2  $(W, n - i, (e_{i2}) (\gamma \downarrow_1), (e_{i2}) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

This means from Definition 2.5 we get

$\forall H_{j1}, H_{j2}.(n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies$

$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A$

Instantiating  $H_{j1}$  with  $H'_{i1}$  and  $H_{j2}$  with  $H'_{i2}$  in IH2 and since the  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  reduces  $\gamma \downarrow_1$  with  $i < n'$  steps therefore  $\exists j < (n' - i) < (n - i)$  s.t  $(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1})$ . Similarly we also have  $(H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A \quad (116)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_V^A$  from Equation 115

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau \sigma))]_V^A$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 2.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \sigma \quad (\text{A1})$$

In order to prove Equation 114 we instantiate  $W'$  with  $W'_2$

- $W'_2 \sqsupseteq W$ :

Since  $W'_1 \sqsupseteq W$  from Equation 115 and  $W'_2 \sqsupseteq W'_1$  from Equation 116

Therefore from Definition 2.3 we get  $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$ :

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

Inorder to prove  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$  we need to show:

- \*  $dom(W'_2.\theta_1) \subseteq dom(H'_1) \wedge dom(W'_2.\theta_2) \subseteq dom(H'_2)$ :  
Directly from Equation 116
- \*  $W'_2.\hat{\beta} \subseteq (dom(W'_2.\theta_1) \times dom(W'_2.\theta_1))$ :  
Directly from Equation 116
- \*  $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge$   
 $(W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$ :
- (a)  $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$ :  
 $\forall(a_1, a_2) \in (W'_2.\hat{\beta})$ .
  - i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :  
From A1 we know that  $W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$   
and since  $W'_1 \sqsubseteq W'_2$  therefore from Lemma 2.16 we get  $W'_2.\theta_1(a_1) =$   
 $W'_2.\theta_2(a_2) = \tau$
  - ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise
  - iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise
  - iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 116 and Lemma 2.17
- (b)  $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). (W'_2, n - n', H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$ :  
 $\forall(a_1, a_2) \in (W'_2.\hat{\beta})$ .
  - i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :  
Since  $H'_1(a_{i1}) = v'_{j1}$  and  $H'_1(a_{i2}) = v'_{j2}$   
From A1 we know that  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$   
And since from Equation 116 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in$   
 $\lceil (\tau) \sigma \rceil_V^A$   
Therefore from Lemma 2.17 we get  
 $(W'_2, n - j - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$
  - ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise
  - iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise
  - iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 116 and from Lemma 2.17
- \*  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in dom(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V$ :  
When  $i = 1$   
Given some  $m$   
 $\forall a_1 \in dom(W'_2.\theta_1)$ .
  - when  $a_1 = a_{i1}$ :  
From Equation 116 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \lceil (\tau) \sigma \rceil_V^A$  thus  
from Lemma 2.15 we know that  
 $\forall m_1. (W'_2.\theta_1, m_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$   
  
Instantiating with  $m$  we get  
 $(W'_2.\theta_1, m, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$
  - Otherwise:  
From Equation 116 and Lemma 2.27
- When  $i = 2$   
Similar reasoning as with  $i = 1$
- $(W'_1, n - n', val'_1, v'_2) \in \lceil (unit) \sigma \rceil_V^A$ :  
From evaluation rule assign we know that  $v'_1 = v'_2 = ()$   
Directly from Definition 2.4



- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Definition 2.4 we know that this would mean that

$$\forall m_1. (W'_1.\theta_1, m_1, a_{i1}) \in [(\text{ref } \tau) \sigma]_V \quad (117)$$

$$\forall m_2. (W'_1.\theta_2, m_2, a_{i2}) \in [(\text{ref } \tau) \sigma]_V \quad (118)$$

In order to prove Equation 114 we instantiate  $W'$  with  $W'_2$  and then we need to show that:

- $W'_2 \sqsupseteq W$ :

Since  $W'_1 \sqsupseteq W$  from Equation 115 and  $W'_2 \sqsupseteq W'_1$  from Equation 116

Therefore from Definition 2.3 we get  $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W'_2$ :

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

In order to prove  $(n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W'_2$  we need to show:

$$* \text{ dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2):$$

Directly from Equation 116

$$* W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2)):$$

Directly from Equation 116

$$* \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}:$$

- (a) When  $(a_{i1}, a_{i2}) \in W'_2.\hat{\beta}$ :

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :

Instantiating Equation 117 and Equation 118 with  $n - n' - 1$  we get

$$W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$$

and since  $W'_1 \sqsubseteq W'_2$  therefore from Definition 2.3 we get  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

From Equation 116 we know that  $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^{\mathcal{A}}$

Therefore  $(W'_2, H_1(a_{i1})', H_2(a_{i2})') \in [(\tau) \sigma]_V^{\mathcal{A}}$

- ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise

- iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise

- iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 116

- (b) When  $(a_{i1}, a_{i2}) \notin W'_2.\hat{\beta}$ :

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise

- ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ :

From Equation 116 we know that  $(n - i - j, H'_{j1}, H'_{j2}) \stackrel{\mathcal{A}}{\triangleright} W'_2$  and since

$(a_{i1}, a_2) \in W'_2.\hat{\beta}$  therefore from Definition 2.9 we know that

$$(W'_2.\theta_1(a_{i1}) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j1}(a_{i1}), H'_{j2}(a_2)) \in [W'_2.\theta_1(a_{i1})]_V^{\mathcal{A}}) \quad (119)$$

Instantiating Equation 117 and Equation 118 with  $n - i - j - 1$  we get  $W'_1.\theta_1(a_{i1}) = \tau \sigma$  therefore from monotonicity we also have  $W'_2.\theta_1(a_{i1}) = \tau \sigma$ .

As a result from Equation 119 we get  $W'_2.\theta_2(a_2) = \tau \sigma$

Also since from Equation 119  $(W'_2, n - i - j - 1, H'_{j1}(a_{i1}), H'_{j2}(a_2)) \in [\tau \sigma]_{\mathcal{V}}^A$  and  $\tau \sigma \searrow \ell, \ell \sigma \not\sqsubseteq \mathcal{A}$  therefore from Lemma 2.15 we know that

$$\forall m.(W'_2.\theta_1, m, H'_{j1}(a_{i1})) \in \lfloor \tau \sigma \rfloor_V \quad (120)$$

$$\forall m.(W'_2.\theta_2, m, H'_{j2}(a_2)) \in \lfloor \tau \sigma \rfloor_V \quad (121)$$

Instantiating  $m$  with  $n - i - j - 1$  in Equation 120 and Equation 121 to get

$$(W'_2.\theta_1, n - i - j - 1, H'_{j1}(a_{i1})) \in \lfloor \tau \sigma \rfloor_V$$

and

$$(W'_2.\theta_2, n - i - j - 1, H'_{j2}(a_2)) \in \lfloor \tau \sigma \rfloor_V$$

Since  $H'_1(a_{i1}) = v'_{j1}$  and  $H'_2(a_2) = H'_{j2}(a_2)$

Again from Equation 116 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$ . This means from Lemma 2.15 and instantiating it with  $n - i - j - 1$  we get

$$(W'_2.\theta_1, n - i - j - 1, v'_{j1}) \in \lfloor (\tau) \sigma \rfloor_V \quad (122)$$

Therefore from Equation 121 and Equation 122 we have

$$(W'_2, n - i - j - 1, H'_1(a_{i1}), H'_2(a_2)) \in [\tau \sigma]_{\mathcal{V}}^A$$

iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ :

Symmetric case as (ii)

iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ :

From Equation 116 and Definition 2.9

\*  $\forall i \in \{1, 2\}.\forall m.\forall a_i \in \text{dom}(W'_2.\theta_i).(W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V$ :

When  $i = 1$

Given some  $m$

$\forall a_1 \in \text{dom}(W'_2.\theta_1)$ .

· when  $a_1 = a_{i1}$ :

From Equation 116 we know that  $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$  thus from Lemma 2.15 we know that

$$(W'_2.\theta_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$$

· Otherwise:

From Equation 116 and Lemma 2.27

When  $i = 2$

Similar reasoning as with  $i = 1$

–  $(W'_1, n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_{\mathcal{V}}^A$ :

From evaluation rule assign we know that  $v'_1 = v'_2 = ()$

Directly from Definition 2.4

13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove:  $(W, n, \Lambda e_i (\gamma \downarrow_1), \Lambda e_i (\gamma \downarrow_2)) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_E^A$

Say  $e_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $e_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition of  $[(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$$

This means that given  $\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A \quad (123)$$

IH1  $(W, n, (e_i (\gamma \downarrow_1), (e_i (\gamma \downarrow_2))) \in [\tau \sigma]_E^A$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $v'_2 = e_2 = \Lambda e_i (\gamma \downarrow_2)$ . We choose  $W' = W$  and we know that  $n' = 0$  we need to show the following:

- $W \sqsubseteq W$ : From Definition 2.3
- $(n, H_1, H_2) \stackrel{A}{\triangleright} W$ : Given
- $(W, n, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$   
Here  $v'_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $v'_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition 2.4 it suffices to prove

$$\begin{aligned} & \forall W' \sqsupseteq W. \forall \ell' \in \mathcal{L}. \forall j < n. \\ & ((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A) \\ & \wedge \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma) \\ & \wedge \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma) \end{aligned}$$

This means given some  $W' \sqsupseteq W$ ,  $\ell' \in \mathcal{L}$  and  $j < n$  we need to show that

$$\begin{aligned} & - \forall W' \sqsupseteq W. \forall \ell' \in \mathcal{L}. \forall j < n. \\ & ((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A): \end{aligned}$$

This means that given some  $W' \sqsupseteq W$ ,  $\ell' \in \mathcal{L}$ ,  $j < n$  we need to prove

$$((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A)$$

From Definition 2.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \triangleright^A W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

This means for some  $H_{s1}$  and  $H_{s2}$  and some  $m < j$  we are given  $(j, H_{s1}, H_{s2}) \triangleright^A W \wedge m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$

And we need to show that

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

We instantiate IH1 with  $H_{s1}$ ,  $H_{s2}$ ,  $m$  and  $\sigma \cup \{\alpha \mapsto \ell'\}$  to obtain

$$\exists W'_1 \sqsupseteq W. (n - m, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - m, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \cup \{\alpha \mapsto \ell'\}$$

Since  $j < n$  therefore from Lemma 2.21 and Lemma 2.17 we get

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

$$- \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

From Lemma 2.25 we know that  $(W'.\theta_1, \gamma \downarrow_1) \in [\Gamma]_V$ . Therefore, we can apply Theorem 2.22 with  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$\forall k. (W'.\theta_1, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

From Lemma 2.16 we get

$$\forall \theta_l \sqsupseteq W'. \theta_1. \forall k. (\theta_l, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

$$- \forall \theta_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

Similar reasoning as in the previous case

#### 14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash_{pc} \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell''/\alpha]}$$

To prove:  $(W, n, (e[]) (\gamma \downarrow_1), (e[]) (\gamma \downarrow_2)) \in [(\tau[\ell''/\alpha]) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e[]) (\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[]) (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e[]) (\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[]) (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A \quad (124)$$

$$\underline{\text{IH}} (W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH and since the  $(e[])$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly  $(e[])$  also reduces to value with  $\gamma \downarrow_2$  therefore we also have  $(H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A \quad (125)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A$  from Equation 125

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.4 we know that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_V^A$$

Here  $v'_{i1} = \Lambda e_{i1}$  and  $v'_{i2} = \Lambda e_{i2}$

This further means that we have

$$\begin{aligned} \forall W'' \sqsupseteq W'_1. \forall \ell' \in \mathcal{L}. \forall j < n - i. ((W'', j, e_{i1}, e_{i2}) \in [\tau[\ell'/\alpha]]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_l, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_{i1}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_l, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_{i2}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \end{aligned} \quad (E1)$$

Instantiating the first conjunct of (E1) with  $W'_1$ ,  $\ell''$  and  $n - i - 1$  we get

$$((W'_1, n - i - 1, e_{i1}, e_{i2}) \in [\tau[\ell''/\alpha]]_E^A)$$

Therefore from Definition 2.5 we get

$$\forall H_1, H_2. (n - i - 1, H_1, H_2) \triangleright^A W'_1 \wedge \forall k < (n - i - 1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Instantiating  $H_1$  and  $H_2$  with  $H'_{i1}$  and  $H'_{i2}$  and since  $e[]$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e$  with  $\gamma \downarrow_1$  reduces in  $i < n' < n$  steps. Therefore  $\exists k < (n' - i - 1)$  steps in which  $e_{i1}$  reduces. Also since  $e[]$  reduces to value with  $\gamma \downarrow_2$  therefore  $e_{i2}$  must also reduce. As a result we get

$$\exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Since  $n' = i + k + 1$  therefore we are done

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Equation 124 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau[\ell''/\alpha] \sigma = \mathbf{A}^{\ell_i}$  and since  $\tau[\ell''/\alpha] \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V \quad (126)$$

In this case from Definition 2.6 we know that

$$\forall m. (W'_1.\theta_1, m, \Lambda e_{h1}) \in [\forall \alpha. (\ell_e, \tau) \sigma]_V \quad (127)$$

$$\forall m. (W'_1.\theta_2, m, \Lambda e_{h2}) \in [\forall \alpha. (\ell_e, \tau) \sigma]_V \quad (128)$$

Applying Definition 2.6 on Equation 127 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_1$$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-FE1})$$

Instantiating  $\theta'$  with  $W'_1.\theta_1$ ,  $j_1$  with  $m_1 + t_1 + 1$  and  $\ell'$  with  $\ell''$

$$\text{Therefore we get } (W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in [\tau[\ell''/\alpha]]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\begin{aligned} & \forall H. (m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \downarrow_{k_c} (H'_1, v'_1) \implies \\ & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in [\tau[\ell''/\alpha]]_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 125 we have

$$(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $H$  with  $H'_{j1}$  from Equation 125 and  $k_c$  with  $t_1$ , we get

$$\begin{aligned} & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau[\ell''/\alpha]]_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad (\text{CF1}) \end{aligned}$$

Similarly applying Definition 2.6 to Equation 128 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}[v/x]) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate  $m$  with  $m_2 + 1 + t_2$  where  $t_2$  is the number of steps in which  $e_{h_2}$  reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-FE2})$$

Instantiating  $\theta'$  with  $W'_1.\theta_2$ ,  $j_1$  with  $m_2 + t_2 + 1$  and  $\ell'$  with  $\ell''$

$$\text{Therefore we get } (W'_1.\theta_2, m_2 + t_2 + 1, e_{h_2}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma$$

From Definition 2.7, we get

$$\begin{aligned} \forall H. (m_2 + t_2 + 1, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1). (H, e_{h_2}) \Downarrow_{k_c} (H'_2, v'_1) \implies \\ \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in [\tau[\ell''/\alpha]]_V \sigma \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 125 we have

$$(n - i, H'_{i_1}, H'_{i_2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i_2}) \triangleright W'_1.\theta_2$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(m_2 + 1 + t_2, H'_{i_2}) \triangleright W'_1.\theta_2$$

Instantiating  $H$  with  $H'_{j_2}$  from Equation 125 and  $k_c$  with  $t_2$ , we get

$$\begin{aligned} \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau[\ell''/\alpha]]_V \sigma \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad (\text{CF2}) \end{aligned}$$

In order to prove Equation 124 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_1.\beta)$ . Now we need to show two things:

$$(a) (n - n', H'_1, H'_2) \triangleright W':$$

From Definition 2.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2):$$

From CF1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CF2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since  $(n - i, H'_{j_1}, H'_{j_2}) \triangleright W'_1$  therefore from Definition 2.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CF1 and CF2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$$

$$- \forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge \\ (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A:$$

4 cases arise for each  $a_1$  and  $a_2$

$$i. H'_{i_1}(a_1) = H'_1(a_1) \wedge H'_{i_2}(a_2) = H'_2(a_2):$$

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

We know from Equation 125 that  $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CF1 and CF2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 2.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From Equation 125 we know that  $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$

This means from Definition 2.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$  (since  $i < n'$ ) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

Same as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From CF1 and CF2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e[\ell''/\alpha] \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from CF1 and CF2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ . Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii.  $H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

Same as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From CF2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$



This means that  $a_2$  was protected at  $\ell_e[\ell''/\alpha]$   $\sigma$  in the world before the modification. Since  $pc \ \sigma \sqcup \ell \ \sigma \sqsubseteq \ell_e[\ell''/\alpha] \ \sigma$  (given) and  $\ell \ \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e[\ell''/\alpha] \ \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 125 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$  that means from Definition 2.9 that  $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$ . Since  $(\ell_e[\ell''/\alpha] \ \sigma) \sqsubseteq \ell'$  therefore from Definition 2.4 we know that  $H'_{i1}(a_1)$  must also have a label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$$

Since from CF2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$  therefore from Definition 2.8 we know that  $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W'.\theta_i(a_i)]_V$$

$i = 1$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W'.\theta_i(a_i)]_V$$

Like before we apply Theorem 2.22 on  $e_{h1}$  and  $e_{h2}$  but this time  $m + 2 + t_1$  and  $m + 2 + t_2$  where  $t_1$  and  $t_2$  are the number of steps in which  $e_{h1}$  and  $e_{h2}$  reduces respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau[\ell''/\alpha] \ \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \ \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \ \sigma)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau[\ell''/\alpha] \ \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \ \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \ \sigma)) \end{aligned}$$

Since we have  $(m + 1, H'_1) \triangleright \theta'_1$  and  $(m + 1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 2.8

$i = 2$

Symmetric to  $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau[\ell''/\alpha] \ \sigma]_V^A:$$

Let  $\tau[\ell''/\alpha] = \mathbf{A}^{\ell_i}$  Since  $\tau[\ell''/\alpha] \ \sigma \searrow \ell \ \sigma$  and since  $\ell \ \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \ \sigma \not\sqsubseteq \mathcal{A}$

From CF1 and CF2 we and Definition 2.4 we get the desired.

15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove:  $(W, n, \nu e (\gamma \downarrow_1), \nu e (\gamma \downarrow_2)) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_E^A$

Say  $e_1 = \nu e (\gamma \downarrow_1)$  and  $e_2 = \nu e (\gamma \downarrow_2)$

From Definition of  $[(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A$$

This means that given  $\forall H_1, H_2. (n', H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A \quad (129)$$

IH1  $(W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \nu e (\gamma \downarrow_1)$  and  $v'_2 = e_2 = \nu e (\gamma \downarrow_2)$ . We choose  $W' = W$  and we know that  $n' = 0$ . We need to show the following:

- $W \sqsubseteq W$ : From Definition 2.3

- $(n, H_1, H_2) \stackrel{A}{\triangleright} W$ : Given

- $(W, n, v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A$

Here  $v'_1 = \nu e (\gamma \downarrow_1)$  and  $v'_2 = \nu e (\gamma \downarrow_2)$

From Definition 2.4 it suffices to prove

$$\forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma} \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$$

We need to prove:

$$- \forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A:$$

This means given some  $W' \sqsupseteq W$ ,  $j < n$  and given that  $\mathcal{L} \models c \sigma$  we need to show that

$$(W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A$$

From Definition 2.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \stackrel{A}{\triangleright} W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

This means for some  $H_{s1}, H_{s2}, m < j$  s.t

$$(H_{s1}, H_{s2}) \triangleright^A W \wedge (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$$

And we need to show that

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

We instantiate IH1 with  $H_{s1}, H_{s2}$  and  $m$  to obtain

$$\exists W'_1 \sqsupseteq W.(n - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, n - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

Since  $j < n$  therefore from Lemma 2.21 and Lemma 2.17 we get

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

$$- \forall \theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}:$$

This means given  $\theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c$

We need to prove:  $(\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$

From Lemma 2.25 we know that  $\forall m_1. (W'.\theta_1, m_1, \gamma \downarrow_1) \in [\Gamma]_{\mathcal{V}}$ . Therefore by instantiating  $m_1$  at  $j$  we can apply Theorem 2.22 to get

$$(\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$$

$$- \forall \theta_l \sqsupseteq W.\theta_2, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}:$$

Symmetric reasoning as in the previous case

#### 16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau}$$

To prove:  $(W, n, (e \bullet) (\gamma \downarrow_1), (e \bullet) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e \bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e \bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A$$

This further means that given

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e \bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e \bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A \quad (130)$$

$$\underline{\text{IH}} (W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n.(H_{i1}, e (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH and since the  $(e \bullet)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since

(e•) reduces to value with  $\gamma \downarrow_2$  therefore also have  $(H_{i2}, e (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n-i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A \quad (131)$$

We case analyze on  $(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A$  from Equation 131

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.4 we know that

$$(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A$$

Here  $v'_{i1} = \nu e_{i1}$  and  $v'_{i2} = \nu e_{i2}$

This further means that we have

$$\begin{aligned} \forall W' \sqsupseteq W. \forall j < n-i. \mathcal{L} \models c \sigma &\implies ((W', j, e_{i1}, e_{i2}) \in [\tau \sigma]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i1}) \in [\tau \sigma]_E^{\ell_e} \sigma) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i2}) \in [\tau \sigma]_E^{\ell_e} \sigma) \end{aligned} \quad (CE1)$$

Instantiating the first conjunct of (CE1) with  $W'_1, \ell''$  and  $n-i-1$  we get

$$((W'_1, n-i-1, e_{i1}, e_{i2}) \in [\tau \sigma]_E^A)$$

Therefore from Definition 2.5 we get

$$\begin{aligned} \forall H_1, H_2. (n-i-1, H_1, H_2) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < (n-i-1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge \\ (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) &\implies \\ \exists W''' \sqsupseteq W'_1. ((n-i-1) - k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1) - k, v'_1, v'_2) &\in [(\tau) \sigma]_V^A \end{aligned}$$

Instantiating  $H_1$  and  $H_2$  with  $H'_{i1}$  and  $H'_{i2}$  and since  $e[]$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e$  with  $\gamma \downarrow_1$  reduces in  $i < n' < n$  steps. Therefore  $\exists k < (n' - i - 1)$  steps in which  $e_{i1}$  reduces. Also since  $e[]$  reduces to value with  $\gamma \downarrow_2$  therefore  $e_{i2}$  must also reduce. As a result we get

$$\exists W''' \sqsupseteq W'_1. ((n-i-1) - k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Since  $n' = i + k + 1$  therefore we are done

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Equation 130 we know that we need to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau \sigma = \mathbf{A}^{\ell_i}$  and since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (132)$$

In this case from Definition 2.6 we know that

$$\forall m.(W'_1.\theta_1, m, \nu e_{h1}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (133)$$

$$\forall m.(W'_1.\theta_2, m, \nu e_{h2}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (134)$$

Applying Definition 2.6 to Equation 133 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma} \text{ where } \theta = W'_1.\theta_1$$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-CE1})$$

Instantiating  $\theta'$  with  $W'_1.\theta_1$ ,  $j_1$  with  $m_1 + t_1 + 1$  and since we know that  $\mathcal{L} \models c \sigma$ . Therefore we get

$$(W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\begin{aligned} \forall H.(m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 131 we have

$$(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $H$  with  $H'_{i1}$  from Equation 131 and  $k_c$  with  $t_1$ , we get

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_{i1}) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_{i1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{CCE1}) \end{aligned}$$

Similarly applying Definition 2.6 to Equation 134 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}) \in [\tau \sigma]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate  $m$  with  $m_2 + 2 + t_2$  where  $t_2$  is the number of steps in which  $e_{h_2}$  reduces  $\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}) \in \lfloor \tau \rfloor_E^{\ell_e[\ell'/\alpha]}$  (FB-CE2)

Instantiating  $\theta'$  with  $W'_1.\theta_2$ ,  $j_1$  with  $m_2 + t_2 + 1$  and  $\ell'$  with  $\ell''$

Therefore we get  $(W'_1.\theta_2, m_2 + t_2 + 1, e_{h_2}) \in \lfloor \tau \sigma \rfloor_E^{\ell_e \sigma}$

From Definition 2.7, we get

$$\begin{aligned} \forall H.(m_2 + t_2, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1).(H, e_{h_2}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 131 we have

$$(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(m_2 + 1 + t_2, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $H$  with  $H'_{i2}$  from Equation 125 and  $k_c$  with  $t_2$ , we get

$$\begin{aligned} \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{CCE2}) \end{aligned}$$

In order to prove Equation 130 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_1.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 2.9 it suffices to show that

–  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$ :

From CCE1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CCE2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 2.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$ :

Since  $(n - i, H'_{j1}, H'_{j2}) \triangleright W'_1$  therefore from Definition 2.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CCE1 and CCE2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore  $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

4 cases arise for each  $a_1$  and  $a_2$

i.  $H'_{i1}(a_1) = H'_1(a_1) \wedge H'_{i2}(a_2) = H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

We know from Equation 125 that  $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CCE1 and CCE2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 2.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From Equation 131 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

This means from Definition 2.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$

(since  $i < n'$ ) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

$$\text{ii. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE1 and CCE2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from CCE1 and CCE2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ .

Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$\text{iii. } H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e \sigma$  in the world before the modification. Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ .

And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 131 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$  that means from Definition 2.9 that  $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$ . Since  $(\ell_e \sigma) \sqsubseteq \ell'$  therefore from Definition 2.4 we know that  $H'_{i1}(a_1)$  must have a label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$$

Since from CCE2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$  therefore from Definition 2.8 we know that  $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$\text{iv. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

$i = 1$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

Like before we apply Theorem 2.22 on  $e_{h1}$  and  $e_{h2}$  but this time  $m + 2 + t_1$  and  $m + 2 + t_2$  where  $t_1$  and  $t_2$  are the number of steps in which  $e_{h1}$  and  $e_{h2}$  reduces respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in \lceil \tau \sigma \rceil_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in \lceil \tau \sigma \rceil_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since we have  $(m + 1, H'_1) \triangleright \theta'_1$  and  $(m + 1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 2.8

$i = 2$

Symmetric to  $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A:$$

Let  $\tau = \mathbf{A}^{\ell_i}$  Since  $\tau \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CCE1 and CCE2 we and Definition 2.4 we get the desired.

□

**Lemma 2.27** (FG: Binary heap well formedness implies unary heap well formedness).  $\forall H_1, H_2, W. (n, H_1, H_2) \triangleright W \implies \forall i \in \{1, 2\}. \forall m. (m, H_i) \triangleright W.\theta_i$

*Proof.* Directly from Definition 2.9

□



**Lemma 2.28** (FG: Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1.  $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (A \sigma) \rrbracket_V^A \subseteq \llbracket (A' \sigma) \rrbracket_V^A$$

2.  $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_V^A \subseteq \llbracket (\tau' \sigma) \rrbracket_V^A$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_E^A \subseteq \llbracket (\tau' \sigma) \rrbracket_E^A$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$

IH1:  $\llbracket (\tau'_1 \sigma) \rrbracket_V^A \subseteq \llbracket (\tau_1 \sigma) \rrbracket_V^A$

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_E^A \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_E^A$

It suffices to prove:

$$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$

From Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \llbracket \tau_1 \sigma \rrbracket_V^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_l, j, e_1[v_1/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e} \sigma) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_l, j, e_2[v_c/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e} \sigma) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 2.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in \llbracket \tau'_1 \sigma \rrbracket_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ \llbracket \tau'_2 \sigma \rrbracket_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta'_l, k, e_1[v'_c/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e} \sigma) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta'_l, k, e_2[v'_c/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e} \sigma) \end{aligned}$$

This means given some  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$  we need to prove:

(a)  $\forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A)$  :

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (135)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 135 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using IH2 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$ :

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$  and since  $\tau'_1 \sigma <: \tau_1 \sigma$  therefore from Lemma 2.24 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (136)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \quad (137)$$

Therefore from Equation 136 and 137 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since  $\tau_2 \sigma <: \tau'_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 2.24 and 2.23 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$ :

Similar reasoning as in the previous case

## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $[((\tau_1 \times \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

IH1:  $[(\tau_1 \sigma)]_V^A \subseteq [(\tau'_1 \sigma)]_V^A$

IH2:  $[(\tau_2 \sigma)]_V^A \subseteq [(\tau'_2 \sigma)]_V^A$

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A$

Therefore from Definition 2.4 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (138)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_{\mathcal{V}}^A$

Again from Definition 2.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

Since from Equation 138 we know that  $(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$

Similarly since  $(W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A$  from Equation 138 therefore from IH2 we have  $(W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$

### 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A \subseteq [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

IH1:  $[(\tau_1 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_1 \sigma)]_{\mathcal{V}}^A$

IH2:  $[(\tau_2 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_2 \sigma)]_{\mathcal{V}}^A$

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s1} = \text{inl } v_{i2}$ :

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (139)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

From Equation 139 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

(b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (140)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

From Equation 140 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $\lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A \subseteq \lceil (\forall \alpha. (\ell'_e, \tau_2)) \sigma \rceil_V^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau_2 \sigma) \rceil_V^A$

IH2:  $\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$

It suffices to prove:  $\forall (W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A$ .

$(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

This means that given:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E^{\ell_e[\ell'/\alpha]})$  (Sub-F1)

And it suffices to prove:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:

$\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}) \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$

This means we are required to show:

(a)  $\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$ :

By instantiating the first conjunct of Sub-F1 with  $W'', n''$  and  $\ell''$  we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from IH1 instantiated at  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$ :

By instantiating the second conjunct of Sub-F1 with  $\theta'_l$  and  $\ell''$  we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha] \sigma})$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 2.24 and Lemma 2.23 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\ell'_e[\ell''/\alpha] \sigma})$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$ :

Similar reasoning as in the previous case

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \Longrightarrow c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove:  $\lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A \subseteq \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

IH:  $\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$

It suffices to prove:  $\forall (W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A. (W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma &\Longrightarrow (W', n', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 &\Longrightarrow (\theta_l, k, e_1) \in \lceil \tau_1 \sigma \rceil_E^{\ell'_e \sigma} \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 &\Longrightarrow (\theta_l, k, e_2) \in \lceil \tau_1 \sigma \rceil_E^{\ell_e \sigma} \quad (\text{Sub-C1}) \end{aligned}$$

And it suffices to prove:  $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma &\Longrightarrow (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 &\Longrightarrow (\theta'_l, j, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma} \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 &\Longrightarrow (\theta'_l, j, e_2) \in \lceil \tau_2 \sigma \rceil_E^{\ell_e \sigma} \end{aligned}$$

This means that we are required to show the following:

(a)  $\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \Longrightarrow (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$ :

We are given  $W'' \sqsupseteq W, n'' < n$  also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \Longrightarrow c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with  $W''$  and  $n''$  we know that the following holds

$$(W'', n'', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A$$

Therefore from IH we get  $(W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \Longrightarrow (\theta'_l, k, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma}$ :

We are given some  $\theta'_l \sqsupseteq W.\theta_1, k$ , also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \Longrightarrow c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with  $\theta'_l$  we know that the following holds

$$(\theta'_l, k, e_1) \in \lceil \tau_1 \sigma \rceil_E^{\ell_e \sigma}$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 2.23 and Lemma 2.24 we get

$$(\theta'_l, k, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma}$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \Longrightarrow (\theta'_l, j, e_2) \in \lceil \tau_2 \sigma \rceil_E^{\ell_e \sigma}$ :

Similar reasoning as in the previous case

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\lceil ((\text{ref } \tau) \sigma) \rceil_V^A \subseteq \lceil ((\text{ref } \tau) \sigma) \rceil_V^A$

Directly from Definition 2.4

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\lceil ((\mathbf{b}) \sigma) \rceil_V^A \subseteq \lceil ((\mathbf{b}) \sigma) \rceil_V^A$

Directly from Definition 2.4

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\lceil ((\text{unit}) \sigma) \rceil_V^A \subseteq \lceil ((\text{unit}) \sigma) \rceil_V^A$

Directly from Definition 2.4

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^\ell <: \mathbf{A}^{\ell'}} \text{FGsub-label}$$

To prove:  $\lceil ((\mathbf{A}^\ell) \sigma) \rceil_V^A \subseteq \lceil ((\mathbf{A}^{\ell'}) \sigma) \rceil_V^A$

2 cases arise

1.  $\ell \sigma \sqsubseteq \ell' \sigma$ :

From Definition 2.4 it suffices to prove:  $\lceil ((\mathbf{A}) \sigma) \rceil_V^A \subseteq \lceil ((\mathbf{A}')) \sigma \rceil_V^A$

This we get directly from IH (Statement (1))

2.  $\ell \sigma \not\sqsubseteq \ell' \sigma$ :

We need to prove that

$$\forall (W, n, v_1, v_2) \in \lceil \mathbf{A} \sigma \rceil_V^A. (W, n, v_1, v_2) \in \lceil \mathbf{A}' \sigma \rceil_V^A$$

From Definition 2.4 it suffices to prove:

$$\forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in \lceil \mathbf{A} \sigma \rceil_V. (W(n). \theta_i, m, v_i) \in \lceil \mathbf{A}' \sigma \rceil_V$$

Since  $\mathbf{A} \sigma <: \mathbf{A}' \sigma$  therefore from Lemma 2.24 we get the desired

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\lceil (\tau \sigma) \rceil_E^A \subseteq \lceil (\tau' \sigma) \rceil_E^A$

This means we need to prove that

$$\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A. (W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$$

This means given  $\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A$

It suffices to prove that  $(W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$

From Definition 2.5 we know we are given:

$$\begin{aligned} & \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. & (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A \quad (\text{Sub-exp1}) \end{aligned}$$

And we need prove that

$$\begin{aligned} & \forall H_{21}, H_{22}, k < n. (n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22}) \implies \\ \exists W'' \sqsupseteq W. & (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \end{aligned}$$

This means that we are given some  $H_{21}, H_{22}$  and  $k < n$  such that  $(n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22})$

It suffices to prove:

$$\exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \quad (141)$$

Instantiating (Sub-exp1) with  $H_{21}, H_{22}$  and  $k$  we get

$$\exists W' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W' \wedge (W', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \quad (142)$$

We choose  $W''$  in Equation 141 as  $W'$  from Equation 142 and we are done □

**Theorem 2.29** (FG: NI). *Say*  $\text{bool} = (\text{unit} + \text{unit})$

$\forall v_1, v_2, e, \tau, n_1.$

$\emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$\emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2) \implies$

$v'_1 = v'_2$

*Proof.* Given some

$\emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$\emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2)$

We need to prove

$v'_1 = v'_2$

From Theorem 2.26 we have

$\forall n. (\emptyset, n, v_1, v_2) \in \lceil \text{bool}^{\top} \rceil_E^{\perp}$

Therefore from Theorem 2.26 and from Definition 2.14 we have

$\forall n. (\emptyset, n, e[v_1/x], e[v_1/x]) \in \lceil \text{bool}^{\perp} \rceil_E^{\perp}$

Therefore from Definition 2.5 we know that

$$\forall n. (\forall H_1, H_2, j < n. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A)$$

Instantiating with  $n_1 + 1$  and then with  $\emptyset, \emptyset, n_1$  we get

$$\exists W' \sqsupseteq W. (1, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', 1, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A$$

Since we have  $(W', 1, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A$  therefore from Definition 2.4 we get  $v'_1 = v'_2$   
 $\square$



## 2.2 Coarse-grained IFC enforcement (CG)

### 2.2.1 CG type system

**Syntax, types, constraints:**

Expressions	$e ::= x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid \text{new } e \mid !e \mid e := e \mid () \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid \text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::= l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::= \mathbf{b} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{unit} \mid \forall \alpha. \tau \mid c \Rightarrow \tau \mid \text{Labeled } \ell \tau \mid \mathbb{C} \ell_i \ell_o \tau$
Constraints	$c ::= \ell \sqsubseteq \ell \mid (c, c)$

**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types  $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau, \text{unit}$  are included.)

$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{CG-label}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau} \text{CG-unlabel}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)} \text{CG-toLabeled}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_i \ell_i \tau} \text{CG-ret}$
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau'} \text{CG-bind}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau} \text{CG-sub}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)} \text{CG-ref}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)} \text{CG-deref}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit}} \text{CG-assign}$	
$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \Lambda e : \forall \alpha. \tau} \text{CG-FI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha]} \text{CG-FE}$
$\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \nu e : c \Rightarrow \tau} \text{CG-CI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau} \text{CG-CE}$

Figure 12: Type system for CG

### 2.2.2 CG semantics

Judgement:  $e \Downarrow_i v$  and  $(H, e) \Downarrow_i^f (H', v)$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \tau <: \tau} \text{CGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{CGsub-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{CGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{CGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{CGsub-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'} \text{CGsub-monad} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{CGsub-forall} \qquad \frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2} \text{CGsub-constraint}
\end{array}$$

Figure 13: CG subtyping

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ } WF} \text{CG-wff-base} \qquad \frac{}{\Sigma; \Psi \vdash \text{unit} \text{ } WF} \text{CG-wff-unit} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \xrightarrow{\ell_e} \tau_2) \text{ } WF} \text{CG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \times \tau_2) \text{ } WF} \text{CG-wff-times} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 + \tau_2) \text{ } WF} \text{CG-wff-sum} \qquad \frac{\text{FV}(\ell) = \emptyset \quad \text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \ell \tau) \text{ } WF} \text{CG-wff-ref} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\forall \alpha. (\tau)) \text{ } WF} \text{CG-wff-forall} \qquad \frac{\Sigma; \Psi, c \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_c} \tau) \text{ } WF} \text{CG-wff-constraint} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\text{Labeled } \ell \tau) \text{ } WF} \text{CG-wff-labeled} \qquad \frac{\Sigma; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\mathbb{C} \ell_i \ell_o \tau) \text{ } WF} \text{CG-wff-monad}
\end{array}$$

Figure 14: Well-formedness relation for CG

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x. e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 e_2 \Downarrow_{i+j+k+1} v_3} \text{cg-app} \qquad \frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{cg-prod} \\
\\
\frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{cg-fst} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{cg-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{cg-inl} \\
\\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{cg-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{cg-case1} \\
\\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{cg-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{cg-Lb} \\
\\
\frac{e \Downarrow_i \Lambda e_i \quad e_i \Downarrow_j v}{e[] \Downarrow_{i+j+1} v} \text{cg-FE} \qquad \frac{e \Downarrow_i \nu e_i \quad e_i \Downarrow_j v}{e\bullet \Downarrow_{i+j+1} v} \text{cg-CE} \qquad \frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-ret} \\
\\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{cg-bind} \\
\\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-unlabel} \qquad \frac{e \Downarrow_i v \quad (H, v) \Downarrow_j^f (H', v')}{(H, \text{toLabeled}(e)) \Downarrow_{i+j+1}^f (H', \text{Lb}(v'))} \text{cg-toLabeled} \\
\\
\frac{e \Downarrow_i \text{Lb}v \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1}^f (H[a \mapsto \text{Lb}v], a)} \text{cg-ref} \qquad \frac{e \Downarrow_i a}{(H, !e) \Downarrow_{i+1}^f (H, H(a))} \text{cg-deref} \\
\\
\frac{e_1 \Downarrow_i a \quad e_2 \Downarrow_j \text{Lb}v}{(H, e_1 := e_2) \Downarrow_{i+j+1}^f (H[a \mapsto \text{Lb}v], ())} \text{cg-assign} \\
\\
\frac{e \in \{x, \lambda y. -, \Lambda -, \nu -, \text{ret} -, \text{bind}(-, -, -), \text{unlabel}(-), \text{toLabeled}(-), \text{new }(-), !-, - := -\}}{e \Downarrow_0 e} \text{cg-val}
\end{array}$$

### 2.2.3 Logical relation for CG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \leftrightarrow \text{Loc}))$

**Definition 2.30** (CG:  $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 2.31** (CG:  $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

**Definition 2.32** (CG: Value Equivalence).

$$\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_V^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_V \wedge (W.\theta_2, j, v_2) \in [\tau]_V & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

**Definition 2.33** (CG: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \rightarrow \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\
[\forall \alpha. \tau]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, \ell' \in \mathcal{L}. \\
&\quad ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E\} \\
[ c \Rightarrow \tau ]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n. \\
&\quad \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E\} \\
[\text{ref } \ell \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \text{Labeled } \ell \tau\} \\
[\text{Labeled } \ell \tau]_V^A &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\
[\mathbb{C} \ell_1 \ell_2 \tau]_V^A &\triangleq \{(W, n, v_1, v_2) \mid \\
&\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\
&\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
&\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\
&\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\}
\end{aligned}$$

**Definition 2.34** (CG: Binary expression relation).

$$[\tau]_E^A \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A\}$$

**Definition 2.35** (CG: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x.e) \mid \forall \theta' \supseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \forall \alpha. \tau \rrbracket_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta', j < m. \forall \ell' \in \mathcal{L}. (\theta', j, e) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E\} \\
\llbracket c \Rightarrow \tau \rrbracket_V &\triangleq \{(\theta, m, \nu e) \mid \mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < m. (\theta', j, e) \in \llbracket \tau \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v)) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbb{C} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \supseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \supseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \mathit{dom}(\theta') \setminus \mathit{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

**Definition 2.36** (CG: Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

**Definition 2.37** (CG: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \mathit{dom}(\theta) \subseteq \mathit{dom}(H) \wedge \forall a \in \mathit{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

**Definition 2.38** (CG: Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \mathit{dom}(W.\theta_1) \subseteq \mathit{dom}(H_1) \wedge \mathit{dom}(W.\theta_2) \subseteq \mathit{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\mathit{dom}(W.\theta_1) \times \mathit{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A) \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \mathit{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

**Definition 2.39** (CG: Label substitution).  $\sigma : \mathit{Lvar} \mapsto \mathit{Label}$

**Definition 2.40** (CG: Value substitution to value pairs).  $\gamma : \mathit{Var} \mapsto (\mathit{Val}, \mathit{Val})$

**Definition 2.41** (CG: Value substitution to values).  $\delta : \mathit{Var} \mapsto \mathit{Val}$

**Definition 2.42** (CG: Unary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \mathit{dom}(\Gamma) \subseteq \mathit{dom}(\delta) \wedge \forall x \in \mathit{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

**Definition 2.43** (CG: Binary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \mathit{dom}(\Gamma) \subseteq \mathit{dom}(\gamma) \wedge \forall x \in \mathit{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

## 2.2.4 Soundness proof for CG

**Lemma 2.44** (CG: Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

*Proof.* Proof by induction on  $\tau$

1. Case b:

From Definition 2.35

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 2.35, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 2.35, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 2.33 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

$$\text{IH1: } \forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \text{ and}$$

$$\text{IH2: } \forall m_2. (W.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$$

From (S01) we know that given some  $m$  and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

Also from (S02) we know that given some  $m$  and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH1 with  $m$  from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$$

Therefore from Definition 2.35, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with  $m$  from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 2.35, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

$$(b) \ v_1 = \text{inr}(v_{i2}) \text{ and } v_2 = \text{inr}(v_{j2})$$

Symmetric reasoning as in the (a) case above

#### 4. Case $\tau_1 \rightarrow \tau_2$ :

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V^A$$

This means from Definition 2.33 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V \implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E) \quad (\text{L0}) \end{aligned}$$

To prove:

$$(a) \ \forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \rightarrow \tau_2]_V:$$

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

This further means that we have some  $\theta', j$  and  $v$  s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

$$(b) \ \forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V:$$

Similar reasoning with  $e_2$

5. Case  $\forall\alpha.\tau$ :

Given:  $(W, n, \Lambda e_1, \Lambda e_2) \in [\forall\alpha.\tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} &\forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E) \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E) \end{aligned} \quad (\text{F0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall\alpha.\tau]_V$ :

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$$

This further means that we are given some  $\theta', m'$  and  $\ell_u$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\ell_u \in \mathcal{L}$

And we need to prove:  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

Instantiating  $\theta_l, i$  and  $\ell''$  in the second conjunct of F0 with  $\theta', m'$  and  $\ell_u$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\ell_u \in \mathcal{L}$

Therefore we get  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

(b)  $\forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall\alpha.\tau]_V$ :

Symmetric reasoning for  $e_2$

6. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, \nu e_1, \nu e_2) \in [c \Rightarrow \tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} &\forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E \end{aligned} \quad (\text{C0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \nu e_1) \in [c \Rightarrow \tau]_V$ :

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E$$

This further means that we are given some  $\theta'$  and  $m'$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\mathcal{L} \models c$

And we need to prove:  $(\theta', m', e_1) \in [\tau]_E$

Instantiating  $\theta_l, j$  in the second conjunct of C0 with  $\theta', m'$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta', m', e_1) \in [\tau]_E$

(b)  $\forall m. (W.\theta_2, m, \nu e_2) \in [c \Rightarrow \tau]_V$ :

Symmetric reasoning for  $e_2$

7. Case ref  $\ell \tau$ :

From Definition 2.33 and 2.35



8. Case Labeled  $\ell \tau$ :

Given  $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 2.32 we know that

$$(W, n, v_1, v_2) \in [\tau]_V^A$$

Therefore from IH we get  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 2.32

9. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 2.35 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case  $l = 1$

And given some  $m$  and  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with  $l = 1$  and the given  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  we get the desired.

Case  $l = 2$

Symmetric reasoning as in the previous case above

□

**Lemma 2.45** (CG: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$(\theta, m, v) \in \lfloor \tau \rfloor_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \lfloor \tau \rfloor_V$

*Proof.* Proof by induction on  $\tau$

1. case b:

Directly from Definition 2.35

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V \wedge (\theta, m, v_2) \in \lfloor \tau_2 \rfloor_V$

IH1 :  $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

IH2 :  $(\theta', m', v_2) \in \lfloor \tau_2 \rfloor_V$

We get the desired from IH1, IH2 and Definition 2.35

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V$

IH :  $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

Therefore from IH and Definition 2.35 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(\theta, m, (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

To prove:  $(\theta', m', (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E \quad (143)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V \implies (\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$

Instantiating Equation 143 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

5. case ref  $\ell \tau$ :

From Definition 2.35 and Definition 2.30

6. case  $\forall\alpha.\tau$ :

Given:  $(\theta, m, (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

To prove:  $(\theta', m', (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

This means from Definition 2.35 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E \quad (144)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

Instantiating Equation 144 with  $\theta''', k$  and  $\ell_j$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\ell_j \in \mathcal{L}$

Therefore we get  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

7. case  $c \Rightarrow \tau$ :

Given:  $(\theta, m, (\nu e_1)) \in [c \Rightarrow \tau]_V$

To prove:  $(\theta', m', (\nu e_1)) \in [c \Rightarrow \tau]_V$

This means from Definition 2.35 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \mathcal{L} \models c \implies (\theta'', j, e_1) \in [\tau]_E \quad (145)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \mathcal{L} \models c \implies (\theta''', k, e_1) \in [\tau]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau]_E$

Instantiating Equation 145 with  $\theta''', k$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta''', k, e_1) \in [\tau]_E$

8. case Labeled  $\ell \tau$ :

Given:  $(\theta, m, (\text{Lb } v)) \in \lfloor \text{Labeled } \ell \tau \rfloor_V$

To prove:  $(\theta', m', (\text{Lb } v)) \in \lfloor \text{Labeled } \ell \tau \rfloor_V$

This means from Definition 2.35 we know that  $(\theta, m, v) \in \lfloor \tau \rfloor_V$

IH:  $(\theta', m', v) \in \lfloor \tau \rfloor_V$

Therefore from IH and Definition 2.35 we get the desired

9. case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(\theta, m, e) \in \lfloor \mathbb{C} \ell_1 \ell_2 \tau \rfloor_V$

To prove:  $(\theta', m', e) \in \lfloor \mathbb{C} \ell_1 \ell_2 \tau \rfloor_V$

This means from Definition 2.35 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 2.35 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t. } (k_1, H) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0),  $k$  with  $k_1$ ,  $\theta_e$  with  $\theta_{e1}$ ,  $H$  with  $H_1$  and  $j$  with  $j_1$ . We know that  $k_1 < m' < m$ ,  $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$ ,  $(k_1, H_1) \triangleright \theta_{e1}$ ,  $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$  and  $i_1 + j_1 < k_1$ . Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

**Lemma 2.46** (CG: Monotonicity binary). *The following holds:*

$$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$$

$$(W, n, v_1, v_2) \in \lfloor \tau \rfloor_V^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in \lfloor \tau \rfloor_V^{\mathcal{A}}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**, unit:

From Definition 2.33

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 :  $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 :  $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 2.33 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 2.33 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i1})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

This means from Definition 2.33 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E)$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)$  (BM-A2)

Similarly from Definition 2.33 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$$

#### 5. Case ref $\ell \tau$ :

From Definition 2.33 and Definition 2.31

#### 6. Case $\forall \alpha. \tau$ :

Given:  $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall \alpha. \tau]_V^A$

To prove:  $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall \alpha. \tau]_V^A$

This means from Definition 2.33 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F2})$$

Similarly from Definition 2.33 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with  $W'', n''$  and  $\ell''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . Also since  $n'' < n'$  and  $n' < n$  therefore  $n'' < n$ . And finally since  $\ell'' \in \mathcal{L}$  therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_l \sqsupseteq W.\theta_1$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

7. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

To prove:  $(\theta', n', (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

This means from Definition 2.33 we know that the following holds

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in \lceil \tau \rceil_E^A$  (BM-C0)

$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in \lfloor \tau \rfloor_E$  (BM-C1)

$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in \lfloor \tau \rfloor_E$  (BM-C2)

Similarly from Definition 2.33 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\mathcal{L} \models c$

And we are required to prove:  $(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

Instantiating BM-C0 with  $W'', n''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . And since  $\mathcal{L} \models c$  therefore we get

$(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_1 \sqsupseteq W.\theta_1$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta_l, k, e_2) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

8. Case Labeled  $\ell \tau$ :

Given:  $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

To prove:  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

From Definition 2.33 2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Therefore from IH we know that  $(W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Hence from Definition 2.33 we get  $(W', n', (\text{Lb}v_1), (\text{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case we know that  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $(W.\theta_2, m, v_2) \in [\tau]_V$

Since  $W.\theta_1 \sqsubseteq W'.\theta_1$  (from Definition 2.31). Therefore from Lemma 2.45 we know that  $\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since  $W.\theta_2 \sqsubseteq W'.\theta_2$  (from Definition 2.31). Therefore from Lemma 2.45 we know that

$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$

Finally from Definition 2.33 we get  $(W', n', (\text{Lb}v_1), (\text{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

9. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(W', n', v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 2.33 we are given that

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{BM-M0}) \end{aligned}$$

Similarly from Definition 2.33 it suffices to prove that

$$\begin{aligned} & \text{(a) } \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right): \\ & \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\ & (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given  $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  and since we know that  $n' \leq n$  and  $W \sqsubseteq W'$  we get the desired

$$\begin{aligned} & \text{(b) } \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right): \end{aligned}$$

Similar reasoning as in the previous case but using Lemma 2.45



□

**Lemma 2.47** (CG: Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .  
 $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$   
 To prove:  $(\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.42 it is given that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

And again from Definition 2.42 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

- $dom(\Gamma) \subseteq dom(\delta)$ :  
 Given
- $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$ :  
 Since we know that  $\forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$  (given)  
 Therefore from Lemma 2.45 we get  
 $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

□

**Lemma 2.48** (CG: Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .  
 $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.43 it is given that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

And again from Definition 2.42 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

- $dom(\Gamma) \subseteq dom(\gamma)$ :  
 Given
- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$ :  
 Since we know that  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$  (given)  
 Therefore from Lemma 2.46 we get  
 $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

□

**Lemma 2.49** (CG: Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .  
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 2.37 it is given that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

And again from Definition 2.42 we are required to prove that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$ :

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :

Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)

Therefore from Lemma 2.45 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

**Lemma 2.50** (CG: Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 2.38 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W.\hat{\beta}). & (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). & (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 2.38 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :

Given

- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

$$\forall (a_1, a_2) \in (W.\hat{\beta}).$$

–  $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given

–  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

Given and from Lemma 2.46

- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :

Given

□

**Theorem 2.51** (CG: Fundamental theorem unary).  $\forall \Sigma, \Psi, \Gamma, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash e : \tau \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge \\ & (\theta, n, \delta) \in [\Gamma \sigma]_V \implies \\ & (\theta, n, e \delta) \in [\tau \sigma]_E \end{aligned}$$

*Proof.* Proof by induction on CG typing derivation

1. CG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, x \delta) \in [\tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some  $i < n$  s.t  $x \delta \Downarrow_i v$

(from *cg - val* we know that  $v = x \delta$  and  $i = 0$ )

It suffices to prove  $(\theta, n, x \delta) \in [\tau \sigma]_V$  (FU-V0)

Since  $(\theta, n, \delta) \in [\Gamma' \sigma]_V$  where  $\Gamma' = \Gamma \cup \{x : \tau\}$ . Therefore from Definition 2.42 we know that  $(\theta, n, \delta(x)) \in [\Gamma'(x) \sigma]_V$

So we are done.

2. CG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e' : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from *cg - val* we know that  $v = \lambda x. e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-L0})$$

From Definition 2.35 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1 \sigma]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

This means given some  $\theta'', v', j$  s.t  $\theta'' \sqsupseteq \theta, j < n$  and  $(\theta'', j, v') \in [\tau_1 \sigma]_V$  (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.47 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E, \text{ s.t. } (\theta_i, j, v_x) \in [\tau_1 \sigma]_V$$

Instantiating IH with  $\theta''$  and  $v'$  from (FU-L1) we get  $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$

3. CG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2 \sigma]_V$$

This means that given some  $i < n$  s.t.  $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2 \sigma]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

From *cg - app* we know that  $v_1 = \lambda x. e'$ . Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-P1})$$

This means from Definition 2.35 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1 \sigma]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2 \sigma]_E \quad (146)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t.  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V \quad (\text{FU-P2})$$

Instantiating Equation 146 with  $\theta, (n - j - k), v_2$  and since we know that  $(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$  therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in [\tau_2 \sigma]_E$$

This means from Definition 2.36 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore we know that  $\exists J < i < n$  s.t  $i = j + k + J$  (since  $j + k + J < n$  therefore  $J < n - j - k$ ) and  $e'[v_2/x] \Downarrow_J v_f$

Therefore we have  $(\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$

Since we know that  $i = j + k + J$  and  $v = v_J$  therefore we get  $(\theta, n - i, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$  (so FU-P0 is proved)

4. CG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

Since we know that  $(e_1, e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from *cg - prod* we know that  $i = j + k + 1$  and  $v = (v_1, v_2)$  therefore from Definition 2.35 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \text{ and } (\theta, n - j - k - 1, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

We get this from (FU-PA1) and Lemma 2.45 and from (FU-PA2) and Lemma 2.45

5. CG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, \text{fst}(e') \delta) \in \lfloor \tau_1 \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau_1 \sigma \rfloor_V$

This means that given some  $i < n$  s.t  $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$(\theta, n - i, v) \in \lfloor \tau_1 \sigma \rfloor_V$  (FU-F0)

IH1:

$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$

Since we know that  $\text{fst}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j (v_1, v_2)$ . This means we have

$(\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$

From Definition 2.35 we know the following holds

$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$  and  $(\theta, n - j, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$  (FU-F1)

From *cg-fst* we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-F0), we are required to prove

$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$

We get this from (FU-F1) and Lemma 2.45

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, \text{inl}(e') \delta) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$

This means that given some  $i < n$  s.t  $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$(\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$  (FU-LE0)

IH1:

$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$

Since we know that  $\text{inl}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-LE1})$$

From *cg - inl* we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-LE0) we are required to prove

$$(\theta, n - j - 1, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

From Definition 2.35 it suffices to prove

$$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

We get this from (FU-LE1) and Lemma 2.45

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in \lfloor \tau \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$$

This means that given some  $i < n$  s.t.  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e_c \delta \Downarrow_j v_c$ . This means we have

$$(\theta, n - j, v_c) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V \quad (\text{FU-C1})$$

2 cases arise:

(a)  $v_c = \text{inl}(v_l)$ :

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t.  $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$ . This means we have

$$(\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C2})$$

From *cg - case1* we know that  $i = j + k + 1$  and  $v = v_1$ . Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau \sigma \rfloor_V$$

We get this from (FU-C2) and Lemma 2.45

(b)  $v_c = \text{inr}(v_r)$ :

Symmetric reasoning as in the previous case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \Lambda e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\forall \alpha. \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from *cg - val* we know that  $v = \Lambda e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. \tau) \sigma]_V \quad (\text{FU-FI0})$$

From Definition 2.35 it further suffices to prove

$$\forall \theta'. \theta \sqsubseteq \theta', j < n. \forall \ell' \in \mathcal{L}. (\theta', j, e' \delta) \in [\tau[\ell'/\alpha]]_E$$

This means given some  $\theta', j, \ell' \in \mathcal{L}$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-FI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau[\ell'/\alpha] \sigma]_E \quad (\text{FU-FI2})$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.47 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1)

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E$$

(FU-FI2) is obtained directly from IH

11. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \nu e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(c \Rightarrow \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\nu e' \delta \Downarrow_i v$

(from *cg - val* we know that  $v = \nu e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_V \quad (\text{FU-CI0})$$

From Definition 2.35 it further suffices to prove



$$\mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < n. (\theta', j, e' \delta) \in [\tau]_E$$

This means given  $\mathcal{L} \models c$  and some  $\theta', j$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-CI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau \sigma]_E \quad (\text{FU-CI2})$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.47 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1). Also we know that  $\mathcal{L} \models c \sigma$  therefore  $\mathcal{L} \models (\Sigma \cup \{c\}) \sigma$

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma]_E$$

(FU-CI2) is obtained directly from IH

12. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' [] \delta) \in [\tau[\ell/\alpha] \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. e' [] \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V$$

This means that given some  $i < n$  s.t  $e' [] \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V \quad (\text{FU-FE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [\forall \alpha. \tau]_E$$

From Definition 2.36 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \Lambda e_{h_1} \implies (\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$$

Since  $e' [] \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \Lambda e_i$

Therefore we know that  $(\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$

From Definition 2.35 we know that

$$\forall \theta'' \sqsupseteq \theta, x < (n - h_1), \ell_h \in \mathcal{L}. (\theta'', x, e_{h_1}) \in [(\tau[\ell_h/\alpha]) \sigma]_E$$

Instantiating  $\theta''$  with  $\theta, x$  with  $n - h_1 - 1$  and  $\ell_h$  with  $\ell$ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [(\tau[\ell/\alpha]) \sigma]_E$$

From Definition 2.36 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $e' [] \delta$  reduces in  $i$  steps therefore from cg-FE we know that  $(i = h_1 + h_2 + 1)$  and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ . Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' \bullet \delta) \in [\tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. e' \bullet \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some  $i < n$  s.t.  $e' \bullet \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau \sigma]_V \quad (\text{FU-CE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [c \Rightarrow \tau \sigma]_E$$

From Definition 2.36 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \nu e_{h_1} \implies (\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$$

Since  $e' \bullet \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \nu e_{h_1}$

Therefore we know that  $(\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$

From Definition 2.35 we know that

$$\mathcal{L} \models c \sigma \implies \forall \theta'' \sqsupseteq \theta, x < (n - h_1). (\theta'', x, e_{h_1}) \in [\tau \sigma]_E$$

Since we know that  $\mathcal{L} \models c \sigma$  and then we instantiate  $\theta''$  with  $\theta$ ,  $x$  with  $n - h_1 - 1$ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [\tau \sigma]_E$$

From Definition 2.36 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since  $e' \bullet \delta$  reduces in  $i$  steps therefore from cg-CE we know that  $(i = h_1 + h_2 + 1)$  and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ . Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since we know that  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in [\tau \sigma]_V$$

14. CG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } (e') : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $\text{new } (e') \delta \Downarrow_i v$

(from  $cg - val$  we know that  $v = \text{new } (e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{new } (e') \delta) \in \lfloor \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma \rfloor_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor (\text{ref } \ell' \tau) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
Also from  $cg - ref$  we know that  $v' = a$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in \lfloor (\text{ref } \ell' \tau) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_E$$

From Definition 2.36 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$$

Since we know that  $(H, \text{new } (e') \delta) \Downarrow_j^f (H', a)$  therefore from  $cg - ref$  we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h$$

Therefore we have

$$(\theta_e, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V \quad (\text{FU-R2})$$

In order to prove (FU-R0) we choose  $\theta'$  as  $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_n$ :

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \lfloor \theta_n(a) \rfloor_V$$

- $\text{dom}(\theta_n) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H) \cup \{a\}$

We know that  $\text{dom}(\theta_n) = \text{dom}(\theta_e) \cup \{a\}$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \lfloor \theta_n(a) \rfloor_V$ :

Since from (FU-R2) we know that  $(\theta_h, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$

Since  $\theta_h \sqsubseteq \theta_n$  and  $k - j - 1 < n - l$  (since  $k < n$  and  $l < j$ ) therefore from

Lemma 2.45 we know that  $(\theta_n, k - j - 1, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$

- (b)  $(\theta_n, k - j - 1, a) \in [(\text{ref } \ell' \tau)]_V$ :  
 From Definition 2.35 it suffices to prove that  $\theta_n(a) = \text{Labeled } \ell' \tau$   
 We get this by construction of  $\theta_n$
- (c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :  
 Holds vacuously
- (d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :  
 From *CG-ref* we know that  $\ell \sqsubseteq \ell'$

15. CG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (!e') \delta) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. !e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

(From *cg-val* we know that  $v = !e' \delta$  and  $i = 0$ )

This means that given some  $i < n$  s.t.  $!e' \delta \Downarrow_i !e' \delta$

It suffices to prove

$(\theta, n, !e' \delta) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

From Definition 2.35 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell') \quad (\text{FU-D0})$

IH:

$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau) \sigma]_E$

From Definition 2.36 this means we have

$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau) \sigma]_V$

Since we know that  $(H, !e') \Downarrow_j^f (H', a)$  therefore from *cg-deref* we know that

$\exists l < j < k$  s.t.  $e' \delta \Downarrow_l v_h, v_h = a$

Therefore we have

$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau) \sigma]_V \quad (\text{FU-D1})$

In order to prove (FU-D0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$   
And since  $H' = H$  (from *cg - deref*) so we are done

- $\forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that

$$\forall a \in \text{dom}(\theta_e).(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since  $H' = H$  and from Lemma 2.45 we get

$$\forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b)  $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau)]_V$ :

From *cg - deref* we know that  $H = H'$  and  $v' = H(a)$

From (FU-D1) and Definition 2.35 we know that  $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that

$$\forall a \in \text{dom}(\theta_e).(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from *cg - deref* we know that  $j \geq 1$ . Therefore from Lemma 2.45 we get

$$(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau)]_V$$

(c)  $(\forall a. H(a) \neq H'(a)) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell'$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

16. CG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit}}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_V$$

This means that given some  $i < n$  s.t  $(e_1 := e_2) \delta \Downarrow_i v$ .

It suffices to prove

$$(\theta, n - i, ()) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from *cg - assign* we know that  $v' = ()$

It suffices to prove

$$\begin{aligned}
& \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, ()) \in [\text{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists l < j < k$  s.t  $e_1 \delta \Downarrow_l a$ . This means we have

$$(\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k-l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k-l-m, v_2) \in [\text{Labeled } \ell' \tau \sigma]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists m < j-l$  (since  $j < k$  therefore  $j-l < k-l$ ) s.t  $e_2 \delta \Downarrow_m v_2$ . This means we have

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k-j, H') \triangleright \theta_e$ :

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$$

•  $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H)$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

•  $\forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

$\forall a \in \text{dom}(\theta_e)$ .

i.  $H(a) = H'(a)$ :

Since  $(k, H) \triangleright \theta_e$  therefore from Definition 2.37 we know that

$$(\theta_e, k-1, H(a)) \in [\theta_e(a)]_V$$

Therefore from Lemma 2.45 we get

$$(\theta_e, k-1-j, H(a)) \in [\theta_e(a)]_V$$

ii.  $H(a) \neq H'(a)$ :

From *cg-assign* we know that  $H'(a) = v_2$

From (FU-A1) we know that  $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that  $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore we get

$$(\theta, k-j+1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore from Lemma 2.45 we get

$$(\theta, k-j-1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

(b)  $(\theta_e, k-j-1, ()) \in [\text{unit}]_V$ :

From Definition 2.35

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$ :

From *CG - assign* we know that  $\ell \sqsubseteq \ell'$

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

17. CG-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{Lb}(e') \delta) \in [\text{Labeled } \ell \tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

This means we are given some  $i < n$  s.t  $\text{Lb}(e') \delta \Downarrow_i v$  and we are required to prove

$(\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

Let  $v = \text{Lb}(v_i)$ . This means from Definition 2.35 we are required to prove

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

IH:  $(\theta, n, e' \delta) \in [\tau \sigma]_E$

This means from Definition 2.36 we have

$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau \sigma]_V$

Since we know that  $\text{Lb}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_i$

Therefore we have  $(\theta, n - j, v_i) \in [\tau \sigma]_V$

From *cg - label* we know that  $i = j + 1$  therefore from Lemma 2.45 we have

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

18. CG-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \delta \Downarrow_i v$

(from *cg - val* we know that  $v = \text{unlabel}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from *cg - unlabel* we know that  $H' = H$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-U0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell \tau) \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from *cg - unlabel* we know that

$$\exists h_1 < j < k \text{ s.t. } e' \delta \Downarrow_{h_1} \text{Lb } v'$$

This means we have

$$(\theta_e, k - h_1, \text{Lb } v') \in [(\text{Labeled } \ell \tau) \sigma]_V$$

This means from Definition 2.35 we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-U1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.49 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-U1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 2.45 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :

Holds vacuously

19. CG-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_i \ell_i \tau}$$



Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

This means we are given some  $i < n$  s.t  $\text{ret}(e') \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

(from *cg - val* we know that  $v = \text{ret}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
Also from *cg - ret* we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [\tau \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [\tau \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from *cg - ret* we know that

$$\exists h_1 < j < k \text{ s.t } e' \delta \Downarrow_{h_1} v'$$

This means we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.49 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-R1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 2.45 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :

Holds vacuously

20. CG-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau'}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

This means we are given some  $i < n$  s.t  $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

(from *cg - val* we know that  $v = \text{bind}(e_1, x.e_2) \delta$  and  $i = 0$ )

Therefore we need to prove

$$(\theta, n, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore from *cg - bind* we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V$$

From Definition 2.35 we know that

$$\begin{aligned}
& \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h_1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\
& \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell)
\end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t.  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.49 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned}
& \exists \theta'' \sqsupseteq \theta_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell) \quad (\text{FU-B1})
\end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in [(\mathbb{C} \ell_i \ell \tau') \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_2 < k - h_1 - J. e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$  therefore from *cg - bind* we know that  $\exists h_2 < j - h_1 - J < k - h_1 - J$  s.t.  $e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

From Definition 2.35 we know that

$$\begin{aligned}
& \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v'') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell)
\end{aligned}$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists v_{h_2}, i$  s.t.  $(v'' \Downarrow_i v_{h_2})$ . From *cg - val* we know that  $v_{h_2} = v''$  and  $i = 0$ . Instantiating  $k_{h_2}$  with  $k - h_1 - J - h_2$ ,  $\theta'_e$  with  $\theta''$ ,  $H$  with  $H'$  (from FU-B1) and  $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$  s.t.  $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$ . And since we already know that  $(k - h_1, H') \triangleright \theta''$  therefore from Lemma 2.49 we get  $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned}
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v'') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell) \quad (\text{FU-B2})
\end{aligned}$$

We get (FU-B0) by choosing  $\theta'$  as  $\theta''$  (from FU-B2)

21. CG-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma]_V$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \delta \Downarrow_i v$

(from *cg - val* we know that  $v = \text{toLabeled}(e') \delta$  and  $i = 0$ )

It suffices to prove

$(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma]_V$

From Definition 2.35 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled} \ell_o \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

And given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
 Also from *cg - tolabeled* we know that  $H' = H$

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled} \ell_o \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-TL0})$

IIH:

$(\theta_e, k, e' \delta) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_V$

Since  $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$  therefore from *cg - tolabeled* we know that  $\exists h_1 < j < k$   
 s.t  $e' \delta \Downarrow_{h_1} v_1$

Therefore we get  $(\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_V$

From Definition 2.35 we know that

$\forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J. (k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies$   
 $\exists \theta'' \sqsupseteq \theta'_e. (k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in [\tau \sigma]_V \wedge$   
 $(\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell)$

Instantiating  $k_{h_1}$  with  $k - h_1$ ,  $H_h$  with  $H$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{toLabeled}(e')) \Downarrow_j^f$   
 $(H', v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already  
 know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.49 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$\exists \theta'' \sqsupseteq \theta'_e. (k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in [\tau \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell) \quad (\text{FU-TL1})$

In order to prove (FU-TL0) we choose  $\theta'$  as  $\theta''$ . Now we need to prove the following

- (a)  $(k - j, H') \triangleright \theta''$ :  
 Since  $(k - h_1 - J, H') \triangleright \theta''$  and  $j = h_1 + J + 1$  therefore from Lemma 2.49 we get  
 $(k - j, H') \triangleright \theta''$
- (b)  $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau \sigma)]_V$ :  
 From *cg - tolabeled* we know that  $v' = \text{toLabeled}(v_1)$   
 From Definition 2.33 it suffices to prove that  $(\theta'', k - j - 1, v_1) \in [\tau \sigma]_V$   
 We get this from (FU-TL1) and Lemma 2.45
- (c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :  
 Directly from (FU-TL1)
- (d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :  
 Directly from (FU-TL1)

□

**Lemma 2.52** (CG: Subtyping unary). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

1.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$
2.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E \subseteq [(\tau' \sigma)]_E$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

IH1:  $[(\tau'_1 \sigma)]_V \subseteq [(\tau_1 \sigma)]_V$  (Statement (1))

$[(\tau_2 \sigma)]_E \subseteq [(\tau'_2 \sigma)]_E$  (Sub-A0, From Statement (2))

It suffices to prove:  $\forall (\theta, n, \lambda x. e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V. (\theta, n, \lambda x. e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

This means that given some  $\theta, n$  and  $\lambda x. e_i$  s.t  $(\theta, n, \lambda x. e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in [\tau_1 \sigma]_V \implies (\theta_1, i, e_i[v/x]) \in [\tau_2 \sigma]_E \quad (147)$$

And it suffices to prove:  $(\theta, n, \lambda x. e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in [\tau'_1 \sigma]_V \implies (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$

Since  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in [\tau_1 \sigma]_V$

As a result from Equation 147 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2 \sigma]_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

## 2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[\tau_1 \times \tau_2 \sigma]_V \subseteq [(\tau'_1 \times \tau'_2) \sigma]_V$

IH1:  $[\tau_1 \sigma]_V \subseteq [\tau'_1 \sigma]_V$  (Statement (1))

IH2:  $[\tau_2 \sigma]_V \subseteq [\tau'_2 \sigma]_V$  (Statement (1))

It suffices to prove:  $\forall(\theta, n, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V. (\theta, n, (v_1, v_2)) \in [(\tau'_1 \times \tau'_2) \sigma]_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$

Therefore from Definition 2.35 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \tag{148}$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [(\tau'_1 \times \tau'_2) \sigma]_V$

Again from Definition 2.35, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 148 we know that  $(\theta, n, v_1) \in [\tau_1 \sigma]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2 \sigma]_V$  from Equation 148 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

## 3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $[(\tau_1 + \tau_2) \sigma]_V \subseteq [(\tau'_1 + \tau'_2) \sigma]_V$

IH1:  $[\tau_1 \sigma]_V \subseteq [\tau'_1 \sigma]_V$  (Statement (1))

IH2:  $[\tau_2 \sigma]_V \subseteq [\tau'_2 \sigma]_V$  (Statement (1))

It suffices to prove:  $\forall(\theta, n, v_s) \in [(\tau_1 + \tau_2) \sigma]_V. (\theta, v_s) \in [(\tau'_1 + \tau'_2) \sigma]_V$

This means that given:  $(\theta, n, v_s) \in [((\tau_1 + \tau_2) \sigma)]_V$

And it suffices to prove:  $(\theta, n, v_s) \in [((\tau'_1 + \tau'_2) \sigma)]_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 2.35 we are given:

$$(\theta, n, v_i) \in [\tau_1 \sigma]_V \quad (149)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_1 \sigma]_V$$

From Equation 149 and IH1 we know that

$$(\theta, n, v_i) \in [\tau'_1 \sigma]_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 2.35 we are given:

$$(\theta, n, v_i) \in [\tau_2 \sigma]_V \quad (150)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_2 \sigma]_V$$

From Equation 150 and IH2 we know that

$$(\theta, n, v_i) \in [\tau'_2 \sigma]_V$$

#### 4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $[((\forall \alpha. \tau_1) \sigma)]_V \subseteq [(\forall \alpha. \tau_2) \sigma]_V$

It suffices to prove:  $\forall (\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_1) \sigma)]_V. (\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_2) \sigma)]_V$

This means that given:  $(\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_1) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E \quad (151)$$

And it suffices to prove:  $(\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_2) \sigma)]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

This means that given some  $\theta_2, j < n, \ell' \in \mathcal{L}$  s.t  $\theta \sqsubseteq \theta_2$

And we are required to prove:  $(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$  therefore from Equation 151 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell']))]_E \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell']))]_E \text{ (Sub-F0, Statement (2))}$$

From (Sub-F0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $[\!(c_1 \Rightarrow \tau_1) \sigma\!]_V \subseteq [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

It suffices to prove:  $\forall (\theta, n, \nu e_i) \in [\!(c_1 \Rightarrow \tau_1) \sigma\!]_V. (\theta, n, \nu e_i) \in [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

This means that given:  $(\theta, n, \nu e_i) \in [\!(c_1 \Rightarrow \tau_1) \sigma\!]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E \quad (152)$$

And it suffices to prove:  $(\theta, n, \nu e_i) \in [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$

This means that given some  $\theta_2, j$  s.t  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

And we are required to prove:  $(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$  and  $\mathcal{L} \models c_2 \sigma \implies c_1 \sigma$  therefore from Equation 152 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E$$

$$[(\tau_1 \sigma)]_E \subseteq [(\tau_2 \sigma)]_E \text{ (Sub-C0, Statement (2))}$$

From (Sub-C0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $[\!(\text{Labeled } \ell \tau) \sigma\!]_V \subseteq [\!(\text{Labeled } \ell' \tau') \sigma\!]_V$

IH:  $[(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$  (Statement (1))



It suffices to prove:

$$\forall(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell \tau) \sigma)]_V. (\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V$$

This means that given some  $\theta, n$  and  $\text{Lb}(e_i)$  s.t  $(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell \tau) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$(\theta, n, v_i) \in [(\tau \sigma)]_V \quad (\text{SL})$$

And we are required to prove that

$$(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V$$

From Definition 2.35 it suffices to prove

$$(\theta, n, v_i) \in [(\tau' \sigma)]_V$$

We get this directly from (SL) and IH

## 7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $[((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$

IH:  $[(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$  (Statement (1))

It suffices to prove:

$$\forall(\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V. (\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$$

This means that given some  $\theta, n$  and  $e$  s.t  $(\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) & \quad (\text{SC0}) \end{aligned}$$

And we are required to prove

$$(\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$$

So again from Definition 2.35 we need to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau' \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) & \end{aligned}$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$   
(SC1)

And we need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau' \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) & \end{aligned}$$

We instantiate (SC0) with  $k, \theta_e, H, j$  from (SC1) and we get

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Since  $\tau \sigma <: \tau' \sigma$  therefore from IH we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \sigma \rfloor_V$$

And since  $\ell'_i \sqsubseteq \ell_i$  therefore we also have

$$\begin{aligned} & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) \end{aligned}$$

#### 8. CGsub-base:

Trivial

#### Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E. (\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

This means that we are given  $(\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E$

From Definition 2.36 it means we have

$$\forall i < n. e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

From Definition 2.36 we need to prove

$$\forall i < n. e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

This further means that given some  $i < n$  s.t  $e \downarrow_i v$ , it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

Instantiating (Sub-E0) with the given  $i$  we get  $(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$

Finally from Statement(1) we get  $(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$

□

**Lemma 2.53** (CG: Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 2.43 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case  $i = 1$

Given some  $m$  we need to show:

- $dom(\Gamma) \subseteq dom(\gamma \downarrow_i)$ :  
 $dom(\gamma) = dom(\gamma \downarrow_i)$   
Therefore,  $dom(\Gamma) \subseteq (dom(\gamma) = dom(\gamma \downarrow_i))$  (Given)
- $\forall x \in dom(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$ :  
We are given:  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$   
Therefore from Lemma 2.44 we know that  
 $\forall m'. (W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$   
Instantiating  $m'$  with  $m$  we get  
 $(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 2$

Symmetric reasoning as in the  $i = 1$  case above

□

**Theorem 2.54** (CG: Fundamental theorem binary).  $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$   
 $\Sigma; \Psi; \Gamma \vdash e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_V^A \implies$   
 $(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

*Proof.* Proof by induction on the typing derivation

1. CG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

To prove:  $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

Say  $e_1 = x(\gamma \downarrow_1)$  and  $e_2 = x(\gamma \downarrow_2)$

From Definition 2.34 it suffices to prove that

$$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some  $i < n$  s.t  $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2$

We are required to prove:  $(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

From *cg-val* we know that  $x(\gamma \downarrow_1) \downarrow x(\gamma \downarrow_1)$  and  $x(\gamma \downarrow_2) \downarrow x(\gamma \downarrow_2)$

This means  $v'_1 = x(\gamma \downarrow_1)$  and  $v'_2 = x(\gamma \downarrow_2)$

Since  $(W, n, \gamma) \in [\tau]_V^A$ . Therefore from Definition 2.43 we know that

$$(W, n, v'_1, v'_2) \in [\tau]_V^A$$

From Lemma 2.46 we get

$$(W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

2. CG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove:  $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

Say  $e_1 = \lambda x. e (\gamma \downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \rightarrow \tau_2) \sigma]_E^A$  it suffices to prove that

$$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some  $i < n$  s.t  $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From *cg - val* we know that  $v'_1 = (\lambda x. e_i) \gamma \downarrow_1$  and  $v'_2 = (\lambda x. e_i) \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\lambda x. e_i) \gamma \downarrow_1, (\lambda x. e_i) \gamma \downarrow_2) \in [\tau]_V^A$$

From Definition 2.33 it suffices to prove

$\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E^A) \wedge$$

$\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E) \wedge$$

$\forall \theta_l \sqsupseteq W. \theta_2, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E) \quad (\text{FB-L0})$$

IH:

$$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in [\Gamma]_V^A$$

In order to prove (FB-L0) we need to prove the following:

(a)  $\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A):$$

This means given some  $W' \sqsupseteq W, j < n, v_1, v_2$  s.t.  $(W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove  $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A$

We get this by instantiating IH with  $W'$  and  $j$

(b)  $\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E):$$

This means given some  $\theta_l \sqsupseteq W. \theta_1, v_c, j$  s.t  $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

We need to prove:  $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E$

It is given to us that

$$(W, n, \gamma) \in [\Gamma]_V^A$$

Therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$

Instantiating  $m$  with  $j$  we get

$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 2.48 we know that

$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$

Since we know that  $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

Therefore we also have

$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1 \sigma\}]_V$

Therefore, we can apply Theorem 2.51 to obtain

$(\theta_l, j, e[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_V$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E):$

Similar reasoning as in the previous case

### 3. CG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$

This further means that given some  $i < n$  s.t  $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$

IH1:  $(W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

This means from Definition 2.34 we know that

$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e_1 \gamma \downarrow_2 \downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$

From *cg - app* we know that  $val_{h1} = \lambda x. e_{h1}$  and  $val_{h2} = \lambda x. e_{h2}$

From Definition 2.33 this further means

$\forall W' \sqsupseteq W, J < (n - j), v_1, v_2.$

$((W', J, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', J, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2 \sigma]_E) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \quad \text{(FB-A1)}$

IH2:  $(W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_j v_{h1'} \wedge e_2 \gamma \downarrow_2 \Downarrow v_{h2'} \implies (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \Downarrow_k v_{h1'}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_2 \gamma \downarrow_2 \Downarrow v_{h2'}$

$$\text{This means we have } (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-A2})$$

Instantiating the first conjunct of (FB-A1) as follows  $W'$  with  $W$ ,  $J$  with  $n - j - k$ ,  $v_1$  and  $v_2$  with  $v'_{h1}$  and  $v'_{h2}$  respectively, we obtain

$$(W, n - j - k, e_{h1}[v'_{h1}/x], e_{h2}[v'_{h2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 2.34

$$\forall l < n - j - k. (e_{h1}[v'_{h1}/x]) \gamma \downarrow_l \Downarrow v_{f1} \wedge e_{h2}[v'_{h2}/x] \Downarrow v_{f2} \implies (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists l < i - j - k < n - j - k$  s.t  $e_{h1}[v'_{h1}/x] \Downarrow_l v_{f1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_{h2}[v'_{h2}/x] \Downarrow v_{f2}$

$$\text{Therefore we have } (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since  $i = j + k + l$  threfore we are done

#### 4. CG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

To prove:  $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. (e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2}) \implies (W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2})$

We are required to prove

$$(W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A \quad (\text{FB-P0})$$

IH1:  $(W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \Downarrow_j v_{f1} \wedge e_1 \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2})$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A \quad (\text{FB-P1})$$

IH2:  $(W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \Downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \Downarrow_j v_{f_2}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \Downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 2.33 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathcal{V}}^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since  $i = j + k + 1$  therefore from (FB-P1) and (FB-P2) and from Lemma 2.46 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

5. CG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove:  $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

Since we know that  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$  therefore  $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

From Definition 2.33 we know that

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Since from  $cg - fst$   $i = j + 1$  therefore from Lemma 2.46 we get

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove:  $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{inl}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1}) \implies \\ (W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \downarrow_1 \downarrow_i v_{f1} \wedge e' \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

Since we know that  $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \downarrow_j v_{f1}$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$  therefore  $e' \gamma \downarrow_2 \downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 2.33 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

From  $cg - \text{inl}$  since  $i = j + 1$  therefore from (FB-IL1) and Lemma 2.46 we get (FB-IL0)

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove:  $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \downarrow v_{f2} \implies \\ (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \downarrow v_{f2}$



We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_i v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t.  $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$  therefore  $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-C1})$$

2 cases arise

- (a)  $v_{h1} = \text{inl}(v_1)$  and  $v'_{h1} = \text{inl}(v'_1)$ :

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_i v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies \\ (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t.  $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

From *cg - case1* we know that  $i = j + k + 1$  therefore from Lemma 2.46 we get (FB-C0)

- (b)  $v_{h1} = \text{inr}(v_1)$  and  $v'_{h1} = \text{inr}(v'_1)$ :  
Symmetric case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

To prove:  $(W, n, \Lambda e' (\gamma \downarrow_1), \Lambda e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (\Lambda e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

This means given some  $i < n$  s.t.  $(\Lambda e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \Downarrow v_{f2}$

From *cg - val* we know that  $v_{f1} = (\Lambda e') \gamma \downarrow_1$  and  $v_{f2} = (\Lambda e') \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\Lambda e')\gamma \downarrow_1, (\Lambda e')\gamma \downarrow_2) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Let  $e_1 = (\Lambda e')\gamma \downarrow_1$  and  $e_2 = (\Lambda e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} & \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ & \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ & \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FI0}) \end{aligned}$$

$$\underline{\text{IH}}: \forall W, n. (W, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E^A$$

In order to prove (FB-FI0) we need to prove the following

$$(a) \quad \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A):$$

This means given  $W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}$  and we are required to prove

$$(W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

$$(b) \quad \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E:$$

This means given  $\theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j$  and we are required to prove

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

Since from Lemma 2.53

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 2.46 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 2.51 to get

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

$$(c) \quad \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E:$$

Symmetric reasoning as before

11. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

$$\text{To prove: } (W, n, e'[](\gamma \downarrow_1), e'[](\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (e'[]) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e'[]) \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

This means given some  $i < n$  s.t  $(e'[]) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e'[]) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE0})$$

$$\underline{\text{IH}}: (W, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (e')\gamma \downarrow_1 \downarrow_i v_{h1} \wedge (e')\gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Since we know that  $(e')\gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e')\gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e' \gamma \downarrow_2 \downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$

From  $cg - FE$  we know that  $v_{h1} = \Lambda e_{h1}$  and  $v_{h2} = \Lambda e_{h2}$

From Definition 2.33 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, k < (n - j), \ell' \in \mathcal{L}. ((W', k, e_{h1}, e_{h2}) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h1}) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h2}) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FE1}) \end{aligned}$$

Instantiating the first conjunct of (FB-FE1) with  $W, n - j - 1$  and  $\ell$  we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau[\ell/\alpha] \sigma]_E^A$$

This means from Definition 2.34 we know that

$$\forall l < n - j - 1. (e_{h1}) \downarrow_l v_{f1} \wedge e_{h2} \downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

Since we know that  $(e')\gamma \downarrow_1 \downarrow_i v_{f1}$  therefore from  $cg - FE$  we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t  $e_{h1} \gamma \downarrow_1 \downarrow_j v_{f1}$ . Similarly since  $(e')\gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e_{h2} \gamma \downarrow_2 \downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-FE2) we get (FB-FE0)

## 12. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

To prove:  $(W, n, \nu e' (\gamma \downarrow_1), \nu e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (\nu e')\gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(c \Rightarrow \tau) \sigma]_V^A$$

This means given some  $i < n$  s.t  $(\nu e')\gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \downarrow v_{f2}$

From  $cg - val$  we know that  $v_{f1} = (\nu e')\gamma \downarrow_1$  and  $v_{f2} = (\nu e')\gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\nu e')\gamma \downarrow_1, (\nu e')\gamma \downarrow_2) \in [(c \Rightarrow \tau) \sigma]_V^A$$

Let  $e_1 = (\nu e')\gamma \downarrow_1$  and  $e_2 = (\nu e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau \sigma]_E \quad (\text{FB-CI0}) \end{aligned}$$

IH:  $\forall W, n. (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

In order to prove (FB-CI0) we need to prove the following

(a)  $\forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \sigma \implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A$ :

This means given  $W' \sqsupseteq W, j < n, \mathcal{L} \models c \sigma$  and we are required to prove

$$(W', j, e_1, e_2) \in [\tau \sigma]_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

(b)  $\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma \implies (\theta_l, j, e_1) \in [\tau \sigma]_E$ :

This means given  $\theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma$  and we are required to prove

$$(\theta_l, j, e_1) \in [\tau \sigma]_E$$

Since from Lemma 2.53  $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 2.46 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 2.51 to get

$$(\theta_l, j, e_1) \in [\tau \sigma]_E$$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau \sigma]_E$ :

Symmetric reasoning as before

13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

To prove:  $(W, n, e' \bullet (\gamma \downarrow_1), e' \bullet (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

This means given some  $i < n$  s.t.  $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-CE0})$$

IH:  $(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. e' \gamma \downarrow_1 \downarrow_i v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(c \Rightarrow \tau) \sigma]_V^A$$

Since we know that  $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t.  $e' \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e' \gamma \downarrow_2 \downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in [(c \Rightarrow \tau) \sigma]_V^A$

From  $cg - CE$  we know that  $v_{h1} = \nu e_{h1}$  and  $v_{h2} = \nu e_{h2}$

From Definition 2.33 this further means

$$\forall W' \sqsupseteq W, k < n - j. \mathcal{L} \models c \sigma \implies (W', k, e_1, e_2) \in [\tau \sigma]_E^A \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_1) \in [\tau \sigma]_E \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_2) \in [\tau \sigma]_E \quad (\text{FB-CE1})$$

Instantiating the first conjunct of (FB-CE1) with  $W$ ,  $n - j - 1$  and since we know that  $\mathcal{L} \models c \sigma$  therefore we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau \sigma]_E^A$$

This means from Definition 2.34 we know that

$$\forall l < n - j - 1. (e_{h1}) \Downarrow_l v_{f1} \wedge e_{h2} \Downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

Since we know that  $(e' \bullet) \gamma \Downarrow_1 \Downarrow_i v_{f1}$  therefore from cg-CE we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t  $e_{h1} \gamma \Downarrow_1 \Downarrow_l v_{f1}$ . Similarly since  $(e' \bullet) \gamma \Downarrow_2 \Downarrow v_{f2}$  therefore  $e_{h2} \gamma \Downarrow_2 \Downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-CE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-CE2) we get (FB-CE0)

14. CG-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove:  $(W, n, \text{Lb}(e') (\gamma \Downarrow_1), \text{Lb}(e') (\gamma \Downarrow_2)) \in [\text{Labeled } \ell \tau \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1}) \implies (W, n - i, \text{Lb}(v_{f1}), \text{Lb}(v'_{f1})) \in [\text{Labeled } \ell \tau \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$

We are required to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\text{Labeled } \ell \tau \sigma]_V^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \Downarrow_1), e' (\gamma \Downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \Downarrow_1 \Downarrow_i v_{f1} \wedge e' \gamma \Downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A$$

Since we know that  $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \Downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $\text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$  therefore  $e' \gamma \Downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 2.33 it suffices to prove that

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A$$

From *cg-label* we know that  $i = j + 1$ . Therefore we get the desired from (FB-LB1) and Lemma 2.46

15. CG-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau}$$

To prove:  $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From *cg - val* we know that  $v_{f1} = \text{unlabel}(e') \gamma \downarrow_1$  and  $v'_{f1} = \text{unlabel}(e') \gamma \downarrow_2$ . Also  $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A$$

This means from Definition 2.33 we need to prove

Let  $e_1 = \text{unlabel}(e') \gamma \downarrow_1$  and  $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau']_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

We need to show

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma): \end{aligned}$$

Also given is some  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{FB-U0})$$

$$\underline{\text{IH}}: (W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{Labeled } \ell \tau) \sigma]_E^A$$

This means from Definition 2.34 we are given

$$\begin{aligned} \forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) &\implies \\ (W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau) \sigma]_V^A & \end{aligned}$$

Since we know that

$$\begin{aligned} (H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2) \wedge j < k &\text{ therefore} \\ \exists I < j < k \text{ s.t } e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) & \end{aligned}$$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h_1}), \text{Lb}(v'_{h_1})) \in [(\text{Labeled } \ell \ \tau) \ \sigma]_V^A$$

This means from Definition 2.33 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \ \sigma, v_{h_1}, v'_{h_1}, \tau \ \sigma) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose  $W'$  as  $W_e$  and from *cg - unlabel* we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . And we already know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore from Lemma 2.50 we get  $(k - j, H_1, H_2) \triangleright W_e$

From *cg - unlabel* we know that  $v'_1, v'_2$  in (FB-U0) is  $v_{h_1}, v'_{h_1}$  respectively. And since from (FB-U1) we know that  $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \ \sigma, v_{h_1}, v'_{h_1}, \tau \ \sigma)$ . Therefore from Lemma 2.55 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, (\ell_i \sqcup \ell) \ \sigma, v_{h_1}, v'_{h_1}, \tau \ \sigma)$$

$$(b) \ \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \ \sigma]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \ \tau \wedge \ell_i \ \sigma \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \ \sigma) \right):$$

Case  $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t } (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k$$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \ \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \ \tau \wedge \ell_i \ \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \ \sigma)$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

$$\text{Instantiating } m \text{ with } k \text{ we get } (W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in [(\mathbb{C} \ \ell_i \ \ell_i \sqcup \ell \ \tau) \ \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c < k. (\text{unlabel } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ \ell_i \ \ell_i \sqcup \ell \ \tau) \ \sigma]_V$$

This further means that given some  $c < k$  s.t  $(\text{unlabel } e')\gamma \downarrow_1 \Downarrow_c v$ . From *cg - val* we know that  $c = 0$  and  $v = (\text{unlabel } e')\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in [(\mathbb{C} \ \ell_i \ \ell_i \sqcup \ell \ \tau) \ \sigma]_V$$

From Definition 2.35 we have

$$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1)$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

16. CG-tolabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

To prove:  $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$   
From *cg - val* we know that  $v_{f1} = \text{toLabeled}(e') \gamma \downarrow_1$ ,  $v_{f2} = \text{toLabeled}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Let  $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$  and  $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 2.33 we are required to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma)) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \end{aligned}$$

We need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma): \end{aligned}$$

This means that we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t

$$(k, H_1, H_2) \triangleright W_e \text{ and } (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma) \quad (\text{FB-TL0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, n - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_V^A$$

Since we know that  $(H_1, \text{toLabeled}(e') \gamma \downarrow_1) \Downarrow_j (H'_1, v'_1)$  and  $(H_2, \text{toLabeled}(e') \gamma \downarrow_1) \Downarrow_j (H'_2, v'_2)$ . Therefore from *cg - val* we know that  $\exists J < j < k \leq n$  s.t  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly we also know that  $e' \gamma \downarrow_2 \Downarrow v'_{h1}$



This means we have

$$(W_e, k - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_V^A$$

From Definition 2.33 we know that

$$\begin{aligned} & (\forall k_1 \leq (k - J), W_e'' \sqsupseteq W_e. \forall H_1'', H_2''. (k_1, H_1'', H_2'') \triangleright W_e'' \wedge \forall v_1'', v_2'', m. \\ & (H_1'', v_{h1}) \Downarrow_m^f (H_1', v_1'') \wedge (H_2'', v'_{h1}) \Downarrow^f (H_2', v_2'') \wedge m < k_1 \implies \\ & \exists W' \sqsupseteq W_e''. (k_1 - m, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1 - m, \ell_o, v_1'', v_2'', \tau \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \quad (\text{FB-TL1}) \end{aligned}$$

We instantiate  $W_e''$  with  $W_e$ ,  $H_1''$  with  $H_1$ ,  $H_2''$  with  $H_2$  and  $k_1$  with  $k$  in (FB-TL1).

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1', v_1') \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \Downarrow^f (H_2', v_2')$ , therefore  $\exists m < j < k \leq n$  s.t  $(H_1, v_{h1}) \Downarrow_m^f (H_1', v_1') \wedge (H_2, v'_{h1}) \Downarrow^f (H_2', v_2')$

This means we have

$$\begin{aligned} & \exists W' \sqsupseteq W_e. (k - m, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - m, \ell_o, v_1'', v_2'', \tau \sigma) \\ & (\text{FB-TL2}) \end{aligned}$$

In order to prove (FB-TL0) we choose  $W'$  as  $W'$  from (FB-TL2). Since from *cg - tolabeled* we know that  $v_1' = \text{Lb}_{\ell_o}(v_1'')$ ,  $v_2' = \text{Lb}_{\ell_o}(v_2'')$  and  $j = m + 1$ , therefore from Lemma 2.50 we get  $(k - j, H_1', H_2') \triangleright W'$ .

Since we have by assumption that  $\ell_i \sqsubseteq \ell_o$  therefore the following cases arise

i.  $\ell_i \sqsubseteq \ell_o \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$(W', k - j, v_1', v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since  $v_1' = \text{Lb}_{\ell_o}(v_1'')$  and  $v_2' = \text{Lb}_{\ell_o}(v_2'')$ . Therefore from Definition 2.33 it suffices to prove that

$$\text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v_1'', v_2'', \tau \sigma)$$

We get this from (FB-TL2) and Lemma 2.55

ii.  $(\ell_i \sqsubseteq \ell_o) \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W', m, v_1') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \text{ and } \forall m. (W', m, v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V$$

Since  $\ell_o \not\sqsubseteq \mathcal{A}$  therefore we get this from (FB-TL2), Definition 2.32 and Definition 2.35

iii.  $(\ell_i \sqsubseteq \mathcal{A} \sqsubseteq \ell_o)$ :

In this case from Definition 2.32 it suffices to prove that

$$(W', k - j, v_1', v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since  $v_1' = \text{Lb}_{\ell_o}(v_1'')$  and  $v_2' = \text{Lb}_{\ell_o}(v_2'')$ . Therefore from Definition 2.33 it suffices to prove that

$$\forall m. (W', m, v_1'') \in [\tau \sigma]_V \text{ and } \forall m. (W', m, v_2'') \in [\tau \sigma]_V$$

We obtain this directly from (FB-TL2) and Definition 2.32

$$\begin{aligned} & (\text{b}) \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \end{aligned}$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{Labeled } \ell_o \tau \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_E$

This means from Definition 2.36 we get

$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$

Instantiating  $c$  with 0 and from  $cg - val$  we know  $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$

$\exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\text{Labeled } \ell_o \tau \sigma]_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

17. CG-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_i \ell_i \tau}$$

To prove:  $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. \text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies$   
 $(W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^A$

This means that given some  $i < n$  s.t  $\text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From  $cg - val$  we know that  $v_{f1} = \text{ret}(e')\gamma \downarrow_1$ ,  $v_{f2} = \text{ret}(e')\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$(W, n, \text{ret}(e')\gamma \downarrow_1, \text{ret}(e')\gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^A$

Let  $v_1 = \text{ret}(e')\gamma \downarrow_1$  and  $v_2 = \text{ret}(e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall v, i. (e_l \Downarrow_i v_l) \implies \right. \\ & \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \end{aligned}$$

It suffices to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau): \end{aligned}$$

We are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

From  $cg - \text{ret}$  we know that  $H'_1 = H_1$  and  $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in \lceil \tau \sigma \rceil_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in \lceil \tau \sigma \rceil_V^A$$

Since we know that  $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow_j^f (H_2, v'_2)$ , therefore  $\exists J < j < k$  s.t  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly  $e' \gamma \downarrow_2 \Downarrow v'_{h1}$ .

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in \lceil \tau \sigma \rceil_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_e$  and from  $cg - \text{ret}$  we know that  $v'_1 = v_{h1}$  and  $v'_2 = v'_{h1}$ . We need to prove the following:

i.  $(k - j, H_1, H_2) \triangleright W_e$ :

Since we have  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.50 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_i, v'_1, v'_2, \tau)$ :

2 cases arise:

A.  $\ell_i \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A$$

Since  $j = J + 1$  therefore we get this from (FB-R1) and Lemma 2.46

B.  $\ell_i \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in \lfloor \tau \sigma \rfloor_V \text{ and } \forall m. (W_e, m, v'_2) \in \lfloor \tau \sigma \rfloor_V$$

We get this From (FB-R1) and Lemma 2.44

(b)  $\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_E$

This means from Definition 2.36 we get

$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_V$

Instantiating  $c$  with 0 and from  $cg - val$  we know that  $v = (\text{ret } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$   
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \lfloor \tau \sigma \rfloor_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

18. CG-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_b : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_l, x.e_b) : \mathbb{C} \ell_i \ell_o \tau'}$$

To prove:  $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in \lfloor \mathbb{C} \ell_i \ell_o \tau' \sigma \rfloor_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} \implies$   
 $(W, n - i, v_{f1}, v'_{f1}) \in \lfloor \mathbb{C} \ell_i \ell_o \tau' \sigma \rfloor_V^A$

This means that given some  $i < n$  s.t  $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From  $cg - val$  we know that  $v_{f1} = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$ ,  $v_{f2} = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b)\gamma \downarrow_1, \text{bind}(e_l, x.e_b)\gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V^A$$

Let  $v_1 = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$  and  $v_2 = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$

This means from Definition 2.33 we need to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma): \end{aligned}$$

This means we are given some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t.  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t.  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} & \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ & (W_e, k - f, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell \tau \sigma]_V^A \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t.  $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell \tau \sigma]_V^A$$

This means from Definition 2.33 we have

$$\begin{aligned} & \left( \forall K \leq (k - f), W_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W_e \wedge \forall v''_1, v''_2, J. \right. \\ & (H''_1, v_{h1}) \Downarrow_J^f (H''_1, v''_1) \wedge (H''_2, v'_{h1}) \Downarrow^f (H''_2, v''_2) \wedge J < K \implies \\ & \left. \exists W'' \sqsupseteq W_e. (K - J, H''_1, H''_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell \sigma, v''_1, v''_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

Instantiating  $K$  with  $(k - f)$ ,  $W'_e$  with  $W_e$ ,  $H'_1$  with  $H_1$  and  $H'_2$  with  $H_2$  in the first conjunct of the above equation. Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.50 we also have  $(k - f, H_1, H_2) \triangleright W_e$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J < j - f < k - f$  s.t  $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v'_{h1}) \wedge (H_2, v'_{h1}) \Downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e.(k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{FB-B1})$$

From Definition 2.32 two cases arise:

i.  $\ell \sigma \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W'', k - f - J, v'_1, v'_2) \in [\tau \sigma]_V^A$

IH2:

$$(W'', k - f - J, e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}), e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\})) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall s < k - f - J. e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}) \Downarrow_s v_{h2} \wedge e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\}) \Downarrow v'_{h2} \implies (W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_V^A$$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f - J < k - f - J$  s.t  $e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}) \Downarrow_s v_{h2} \wedge e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_V^A$$

This means from Definition 2.33 we know that

$$\left( \forall K_s \leq (k - f - J - s), W_s \sqsupseteq W''. \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right.$$

$$(H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies$$

$$\left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_i, v'_1, v'_2, \tau' \sigma) \right) \wedge$$

$$\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \Big)$$

Instantiating  $K_s$  with  $(k - f - J - s)$ ,  $W_s$  with  $W''$ ,  $H_1$  with  $H'_1$  and  $H_2$  with  $H_2$ . Since we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Lemma 2.50 we also have  $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J_s < j - f - J - s < k - f - J - s$  s.t  $(H'_1, v'_1) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_2) \Downarrow^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s. (k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma) \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose  $W'$  as  $W'_s$ . From  $cg - bind$  we know that  $H'_1 = H'_{s1}$ ,  $H'_2 = H'_{s2}$ ,  $v'_1 = v'_{s1}$ ,  $v'_2 = v'_{s2}$  and  $j = f + J + s + J_s + 1$ . And we need to prove:

- A.  $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$ :  
 Since from (FB-B2) we know that  $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$  therefore from Lemma 2.50 we get  
 $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$
- B.  $ValEq(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$ :  
 Since from (FB-B2) we know that  $ValEq(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$  therefore from Lemma 2.55 we get  
 $ValEq(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$

ii.  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma)$$

Since  $\ell_i \sigma \sqsubseteq \ell \sigma \sqsubseteq \ell_o \sigma$  (by assumption) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore we have  $\ell_o \sigma \not\sqsubseteq \mathcal{A}$

This means that from Definition 2.32 it suffices to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'.\theta_1, m_{u1}, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge \forall m_{u2}. (W'.\theta_2, m_{u2}, v'_2) \in \llbracket \tau' \sigma \rrbracket_V$$

This means given some  $m_{u1}, m_{u2}$  and we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in \llbracket \tau' \sigma \rrbracket_V \quad (\text{FB-B01})$$

In this case we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in \llbracket \tau \sigma \rrbracket_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in \llbracket \tau \sigma \rrbracket_V \quad (\text{FB-B3})$$

Since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_1 \Downarrow_j v'_1$  therefore  $\exists J_1 < j - f - J < k - f - J$  s.t  $(e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \Downarrow_{J_1} v'_1$ . Similarly,  $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$  s.t  $(H'_1, v'_1) \Downarrow_{J'_1}^f -$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  in the first conjunct of (FB-B3)

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in \llbracket \tau \sigma \rrbracket_V$$

Since  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V^A$  therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  we get  $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$

From Lemma 2.47 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V \quad (\text{FB-B4})$$

Now we can apply Theorem 2.51 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_E$$

This means from Definition 2.36 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \Downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_V \quad (\text{FB-B5})$$

Instantiating  $c_1$  with  $J_1$  in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_V$$

From Definition 2.35 we have

$$\forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \Downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies$$

$$\exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge$$

$$\begin{aligned} (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1)/\text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Instantiating  $K$  with  $m_{u1} + 1 + J'_1$ ,  $\theta'_e$  with  $W''.\theta_1$ ,  $H_1$  with  $H'_1$  (from FB-B1) and  $J_2$  with  $J'_1$  we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1.(m_{u1} + 1, H_1'') \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1)/\text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow v'_2$ . Say this reduction happens in  $t$  steps. Therefore  $\exists t_1 < t < k \leq n$  s.t  $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v'_2\} \downarrow_{t_1} v_{l2}$  and similarly  $\exists t_2 < t - t_1 < k - t_1$  s.t  $(H, v_{l2})\gamma \downarrow_2 \downarrow_{t_2}^f (H_2'', v_2'')$

Again since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow_t v'_2$  therefore  $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$  s.t  $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v_2''\} \downarrow_{J_2} v'_2$ . Similarly  $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$  s.t  $(H_2', v_2') \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with  $m_{u2} + 1 + J_2 + J'_2$  we get  $(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v_2'') \in [\tau \sigma]_V$

Again since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.53 we know that  $\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $m_{u2} + 1 + J_2 + J'_2$  we get  $(W.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 2.47 we know that

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 2.51 to get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v_2''\}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c_2 < (m_{u2} + 1 + J_2 + J'_2). (e_b)\gamma \downarrow_2 \cup \{x \mapsto v_2''\} \downarrow_{c_2} v_{o2} \implies (W''.\theta_2, m_{u2} + 1 + J_2 - c_2, v_{o2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V \quad (\text{FB-B8})$$

Instantiating  $c_2$  with  $J_2$  in (FB-B8) we get

$$(W''.\theta_2, m_{u2} + 1 + J'_2, v_{o2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

From Definition 2.35 we have

$$\forall K \leq (m_{u2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3. (K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o2}) \downarrow_{J_3}^f (H_2'', v_2'') \wedge J_3 < K \implies$$

$$\begin{aligned} \exists \theta'_2 \sqsupseteq \theta'_e.(K - J_3, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v_2') \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2)/\text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \end{aligned}$$

Instantiating  $K$  with  $m_{u2} + 1 + J'_2$ ,  $\theta'_e$  with  $W''.\theta_2$ ,  $H_2$  with  $H_2'$  (from FB-B1) and  $J_3$  with  $J'_2$ , we get

$$\begin{aligned} \exists \theta'_2 \sqsupseteq W''.\theta_2.(m_{u2} + 1, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, m_{u2} + 1, v_2') \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2)/\text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose  $W'$  as  $W_n$  where  $W_n$  is defined as follows:  $W_n.\theta_1 = \theta'_1$  (From (FB-B6))



$$W_n.\theta_2 = \theta'_2 \text{ (From (FB-B9))}$$

$$W_n.\hat{\beta} = W''.\hat{\beta} \text{ (From (FB-B1))}$$

It suffices to prove

- $(k - j, H'_1, H''_2) \triangleright W_n$ :

From Definition 2.38 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H''_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H''_2):$$

From (FB-B6) we know that  $(m_{u1}+1, H''_1) \triangleright \theta'_1$  therefore from Definition 2.37 we know that  $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H''_1)$

Similarly from (FB-B9) we know that  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$  therefore from Definition 2.37 we know that  $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H''_2)$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Definition 2.38 we know that  $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq W_n.\theta_1$  and  $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k-j-1, H'_1(a_1), H''_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A):$$

4 cases arise for each  $(a_1, a_2) \in W_n.\hat{\beta}$

$$A. H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

We know from that  $(k - f - J, H'_1, H'_2) \triangleright W''$

Therefore from Definition 2.38 we have

$$\forall (a'_1, a'_2) \in (W''.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

Since  $W_n.\hat{\beta} = W''.\hat{\beta}$  by construction therefore

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

From (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq \theta'_1$  and  $W''.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 2.30

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

To prove:

$$(W_n, k - j - 1, H'_1(a_1), H''_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

From (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$

This means from Definition 2.38 we know that

$$\forall (a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) = W''.\theta_2(a_{i2}) \wedge (W'', k - f - J - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in \lceil W''.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'' \sqsubseteq W_n$  and  $k - j - 1 < k - f - J - 1$  (since  $j = f + J + J_1 + 1$  therefore from Lemma 2.46 we get

$$(W_n, k - j - 1, H_1'(a_1), H_2'(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

B.  $H_1'(a_1) \neq H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$ :

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B6) and (FB-B9) we know that

$$(\forall a. H_1'(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

$$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell' \not\sqsubseteq \mathcal{A}$ .

Also from (FB-B6) and (FB-B9),  $(m_{u1}+1, H_1'') \triangleright \theta_1'$  and  $(m_{u2}+1, H_2'') \triangleright \theta_2'$ .

Therefore from Definition 2.37 we have

$$(\theta_1', m_{u1}, H_1''(a_1)) \in [\theta_1'(a_1)]_V \text{ and}$$

$$(\theta_2', m_{u2}, H_2''(a_1)) \in [\theta_2'(a_2)]_V$$

Since  $m_{u1}$  and  $m_{u2}$  are arbitrary indices therefore from Definition 2.33 we get

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in [\theta_1'(a_1)]_V^A$$

C.  $H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$ :

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B9) we know that

$$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell' \not\sqsubseteq \mathcal{A}$ .

Since from (FB-B1) we know that  $(k - f - J, H_1', H_2') \stackrel{A}{\triangleright} W''$  that means from Definition 2.38 that  $(W'', k - f - J - 1, H_1'(a_1), H_2'(a_2)) \in [W''.\theta_1(a_1)]_V^A$ . Since  $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$  and since  $\ell' \not\sqsubseteq \mathcal{A}$  therefore from Definition 2.33 and Definition 2.32 we know that

Therefore

$$\forall m. (W''.\theta_1, m, H_1'(a_1)) \in W''.\theta_1(a_1) \quad (\text{F})$$

Instantiating the (F) with  $m_{u1}$  and using Lemma 2.45 we get

$$(\theta'_1, m_{u1}, H'_1(a_1)) \in \theta'_1(a_1)$$

Since from (FB-B9) we know that  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$  therefore from Definition 2.37 we know that  $(\theta'_2, m_{u2}, H''_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.33 we get

$$(W', k - j - 1, H'_1(a_1), H''_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$D. H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

Symmetric reasoning as in the previous case

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in [W_n.\theta_i(a_i)]_V:$$

Case  $i = 1$

Given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in [W_n.\theta_i(a_i)]_V$$

This further means that given some  $a_1 \in \text{dom}(W_n.\theta_1)$  we need to show

$$(W_n.\theta_1, m, H''_1(a_1)) \in [W_n.\theta_1(a_1)]_V$$

Since  $W_n.\theta_1 = \theta'_1$ , it suffices to prove

$$(\theta'_1, m, H''_1(a_1)) \in [\theta'_1(a_1)]_V$$

Like before we apply Theorem 2.51 on  $e_b \gamma \downarrow_1 \cup \{x \mapsto v'_1\}$  but this time at  $m + 1 + J_1 + J'_1$  to get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1.(m + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) / \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Since we have  $(m + 1, H''_1) \triangleright \theta'_1$  therefore from Definition 2.37 we get the desired.

Case  $i = 2$

Similar reasoning as in the  $i = 1$  case

$$\bullet (W'.\theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V:$$

We get this from (FB-B6), (FB-B9) and Lemma 2.45 we get the desired

19. CG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } (e') : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)}$$

To prove:  $(W, n, \text{new } (e') (\gamma \downarrow_1), \text{new } (e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{new } (e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{new } (e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \downarrow v'_{f1}$

From *cg-val* we know that  $v_{f1} = \text{new } (e') \gamma \downarrow_1$ ,  $v_{f2} = \text{new } (e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{new } (e') \gamma \downarrow_1, \text{new } (e') \gamma \downarrow_2) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V^A$$

Let  $v_1 = \text{new } (e')\gamma \downarrow_1$  and  $v_2 = \text{new } (e')\gamma \downarrow_2$

From Definition 2.33 we are required to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_l) \in [(\text{ref } \ell' \tau)]_V \sigma \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove the following:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma): \end{aligned}$$

This means we are given some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also we are given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma) \quad \text{(FB-R0)}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} & \forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ & (W_e, k-f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e' \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k-f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \quad \text{(FB-R1)}$$

In order to prove (FB-R0) we choose  $W'$  as  $W_n$  where

$$W_n. \theta_1 = W_e. \theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n. \theta_2 = W_e. \theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n. \hat{\beta} = W_e. \hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

$$\text{i. } (k-j, H'_1, H'_2) \triangleright W_n:$$

From Definition 2.38 it suffices to prove:

$$\text{dom}(W_n. \theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n. \theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_n. \hat{\beta}) \subseteq (\text{dom}(W_n. \theta_1) \times \text{dom}(W_n. \theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_n. \hat{\beta}). (W_n. \theta_1(a_1) = W_n. \theta_2(a_2) \wedge$$

$$(W_n, (k-j) - 1, H'_1(a_1), H'_2(a_2)) \in [W_n. \theta_1(a_1)]_V^A) \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n. \theta_i). (W_n. \theta_i, m, H_i(a_i)) \in [W_n. \theta_i(a_i)]_V$$

This means we need to prove

- $dom(W_n.\theta_1) \subseteq dom(H'_1) \wedge dom(W_n.\theta_2) \subseteq dom(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (dom(W_n.\theta_1) \times dom(W_n.\theta_2))$ :

We know that  $dom(W_n.\theta_1) = dom(W_e.\theta_1) \cup \{a_1\}$  and  $dom(W_n.\theta_2) = dom(W_e.\theta_2) \cup \{a_2\}$

Also  $dom(H'_1) = dom(H_1) \cup \{a_1\}$  and  $dom(H'_2) = dom(H_2) \cup \{a_2\}$

Therefore from  $(k, H_1, H_2) \triangleright W_e$  and from construction of  $W_n$  we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2) \wedge (W_n, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_n.\theta_1(a'_1) \rceil_V^A)$ :

$\forall (a'_1, a'_2) \in (W_n.\hat{\beta})$ .

- A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

Since from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \sigma \rceil_V^A$

And since from  $cg - ref$  we know that  $H'_1(a_1) = v_{h1}$ ,  $H'_2(a_2) = v'_{h1}$  and  $j = f + 1$  therefore from Lemma 2.46 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

- B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise

- C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise

- D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.38

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in dom(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in \lfloor W_n.\theta_i(a'_i) \rfloor_V$ :

When  $i = 1$

Given some  $m$

$\forall a'_1 \in dom(W_n.\theta_1)$ .

- when  $a'_1 = a_1$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

And from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \sigma \rceil_V^A$

Therefore from Lemma 2.44 get the desired

- Otherwise:

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.38

When  $i = 2$

Similar reasoning as with  $i = 1$

- ii.  $ValEq(\mathcal{A}, W_n, k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma)$ :

From  $cg - ref$  we know that  $v'_1 = a_1$  and  $v'_2 = a_2$

2 cases arise:

- A.  $\ell \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$(W_n, k - j, a_1, a_2) \in (\text{ref } \ell' \tau) \sigma$$

From Definition 2.33 it suffices to prove

$$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

This holds from construction of  $W_n$

B.  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$\forall m. (W_n.\theta_1, m, a_1) \in (\text{ref } \ell' \tau) \sigma$  and  $(W_n.\theta_2, m, a_2) \in (\text{ref } \ell' \tau) \sigma$

From Definition 2.35 this means for any given  $m$  we need to prove that

$W_n.\theta_1(a_1) \in (\text{Labeled } \ell' \tau) \sigma$  and  $W_n.\theta_2(a_2) \in (\text{Labeled } \ell' \tau) \sigma$

This holds from construction of  $W_n$

- (b)  $\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^{\mathcal{A}}$  therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_E$

This means from Definition 2.36 we get

$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$

This further means that given some  $c < k$  s.t.  $\text{ref } (e')\gamma \downarrow_1 \Downarrow_c v$ . From  $cg - val$  we know that  $c = 0$  and  $v = \text{ref } (e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$   
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

20. CG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)}$$

To prove:  $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. !e' \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f_1} &\implies \\ (W, n - i, v_{f_1}, v'_{f_1}) &\in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $!e' \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f_1}$

From  $cg - val$  we know that  $v_{f_1} = !e' \gamma \downarrow_1$ ,  $v_{f_2} = !e' \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, !e' \gamma \downarrow_1, !e' \gamma \downarrow_2) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A$$

Let  $v_1 = !e' \gamma \downarrow_1$  and  $v_2 = !e' \gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma): \end{aligned}$$

This means we are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\begin{aligned} &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma) \\ &(\text{FB-D0}) \end{aligned}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{ref } \ell \tau) \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} &\implies \\ (W_e, k - f, v_{h_1}, v'_{h_1}) &\in [(\text{ref } \ell \tau) \sigma]_V^A \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau) \sigma]_V^A \quad (\text{FB-D1})$$

In order to prove (FB-D0) we choose  $W'$  as  $W_e$ . Also from  $cg - deref$  we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . Also we know that  $v_{h_1} = a_1$  and  $v'_{h_1} = a_2$ .

- $(k - j, H_1, H_2) \triangleright W_e$ :

Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.50 we get  
 $(k - j, H_1, H_2) \triangleright W_e$

- $ValEq(\mathcal{A}, W_e, k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)$ :

From *cg - ref* we know that  $v'_1 = H_1(a_1)$  and  $v'_2 = H_2(a_2)$

2 cases arise:

- $\ell' \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$(W_e, k - j, v'_1, v'_2) \in (\text{Labeled } \ell \tau) \sigma$$

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau \sigma]_V^{\mathcal{A}}$

Therefore from Definition 2.33 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that  $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau \sigma]_V^{\mathcal{A}}$

From Lemma 2.46 we get  $(W_e, k - j, H_1(a_1), H_2(a_2)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\mathcal{A}}$

- $\ell' \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W_e.\theta_1, m, H_1(a_1)) \in (\text{Labeled } \ell \tau) \sigma \text{ and } (W_e.\theta_2, m, H_2(a_2)) \in (\text{Labeled } \ell \tau) \sigma$$

(FB-B2)

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau \sigma]_V^{\mathcal{A}}$

Therefore from Definition 2.33 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that  $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau \sigma]_V^{\mathcal{A}}$

Finally from Lemma 2.44 we get (FB-B2)

$$(b) \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma):$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma)$$

Since  $(W, n, \gamma) \in [\Gamma]_V^{\mathcal{A}}$  therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, (!e' \gamma \downarrow_1) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_E$$

This means from Definition 2.36 we get



$$\forall c < k. !e'\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$$

Instantiating  $c$  with 0 and from  $cg - val$  we know that  $v = !e'\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, !e'\gamma \downarrow_1) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$$

From Definition 2.35 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

21. CG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_l := e_r : \mathbb{C} \ell \ell \text{ unit}}$$

To prove:  $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. (e_l := e_r) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $(e_l := e_r) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \downarrow v'_{f1}$

From  $cg - val$  we know that  $v_{f1} = (e_l := e_r) \gamma \downarrow_1$ ,  $v_{f2} = (e_l := e_r) \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_V^A$$

Let  $e_1 = (e_l := e_r) \gamma \downarrow_1$  and  $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ \left. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \right. \\ \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}) \right) \wedge \\ \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{unit}]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

- (a)  $\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2.$   
 $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$   
 $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}):$

This means we are given some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

And finally given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit})$$

(FB-A0)

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\text{ref } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies$$

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_V^A \quad (\text{FB-A1})$$

IH2:

$$(W_e, k - f, e_r (\gamma \downarrow_1), e_r (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall s < k - f. e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2} \implies$$

$$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists s < j - f < k - f$  s.t  $e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \quad (\text{FB-A2})$$

In order to prove (FB-A0) we choose  $W'$  as  $W_e$ . Also from *cg - assign* we know that  $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$  and  $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$ , and  $j = f + s + 1$

We need to prove the following:

- i.  $(k - j, H'_1, H'_2) \triangleright W_e$ :

Say  $v_{h1} = a_1$  and  $v'_{h1} = a_2$

From Definition 2.38 it suffices to prove:

$$\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a_1) = W_e.\theta_2(a_2) \wedge$$

$$(W_e, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [W_e.\theta_1(a_1)]_V^A) \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [W_e.\theta_i(a_i)]_V$$

This means we need to prove

- $dom(W_e.\theta_1) \subseteq dom(H'_1) \wedge dom(W_e.\theta_2) \subseteq dom(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (dom(W_e.\theta_1) \times dom(W_e.\theta_2))$ :  
 Since  $dom(H_1) = dom(H'_1)$  and  $dom(H_2) = dom(H'_2)$ , and also we know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore we obtain the desired directly from Definition 2.38
  - $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \llbracket W_e.\theta_1(a'_1) \rrbracket_V^A)$ :  
 $\forall (a'_1, a'_2) \in (W_e.\hat{\beta})$ .
    - A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :  
 From (FB-A1) and from Definition 2.33 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$   
 Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \llbracket \text{Labeled } \ell' \tau \sigma \rrbracket_V^A$   
 And since from *cg - assign* we know that  $H'_1(a_1) = v_{h2}, H'_2(a_2) = v'_{h2}$  and  $j = f + s + 1$  thfore from Lemma 2.46 we get  
 $(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \llbracket W_e.\theta_1(a_1) \rrbracket_V^A$
    - B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise
    - C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise
    - D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :  
 Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.38
  - $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in dom(W_e.\theta_i). (W_e.\theta_i, m, H_i(a'_i)) \in \llbracket W_e.\theta_i(a'_i) \rrbracket_V$ :  
When  $i = 1$   
 Given some  $m$   
 $\forall a'_1 \in dom(W_e.\theta_1)$ .
    - when  $a'_1 = a_1$ :  
 From (FB-A1) and from Definition 2.33 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$   
 Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \llbracket \text{Labeled } \ell' \tau \sigma \rrbracket_V^A$   
 Therefore from Lemma 2.44 get the desired
    - Otherwise:  
 Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.38When  $i = 2$   
 Similar reasoning as with  $i = 1$
  - ii.  $ValEq(\mathcal{A}, W_e, k - j, \ell, (), (), \text{unit})$ :  
 Holds directly from Definition 2.32 and Definition 2.33
- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \text{unit} \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \ell \sigma)$
- Case  $l = 1$   
 Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$
- We need to prove  
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket (\text{unit}) \sigma \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sigma \sqsubseteq \ell'') \wedge$   
 $(\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \ell \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.53 we know that  
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c < k. (e_l := e_r)\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_V$$

Instantiating  $c$  with 0 and from  $cg - val$  we know that  $v = (e_l := e_r)\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_V$

From Definition 2.35 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

□

**Lemma 2.55** (CG: Equivalence of values).  $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies$$

$$\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$$

*Proof.* Given that  $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$ . From Definition 2.32 two cases arise

1.  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, i, v_1, v_2) \in [\tau]_V^A$

2 cases arise

(a)  $\ell' \sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^A$  therefore from Lemma 2.46 we know that  $(W', j, v_1, v_2) \in [\tau]_V^A$

And thus from Definition 2.32 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^A$  therefore from Lemma 2.44 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.45 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.32 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Given is  $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.45 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.32 we know that  $ValEq(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

**Lemma 2.56** (CG: Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

1.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$
2.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^A \subseteq [(\tau' \sigma)]_E^A$

*Proof.* Proof of statement (1)

Proof by induction on the  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

IH1:  $[(\tau'_1 \sigma)]_V^A \subseteq [(\tau_1 \sigma)]_V^A$  (Statement 1)

$[(\tau_2 \sigma)]_E^A \subseteq [(\tau'_2 \sigma)]_E^A$  (Sub-A0 From Statement 2)

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

From Definition 2.33 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2 \sigma]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 2.33 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2 \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_l, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_l, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E) \end{aligned}$$

This means need to prove:

(a)  $\forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A)$  :

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (153)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 153 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using (Sub-A0) we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E)$ :

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$  and since  $\tau'_1 \sigma <: \tau_1 \sigma$  therefore from Lemma 2.52 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (154)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E) \quad (155)$$

Therefore from Equation 154 and 155 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E$

Since  $\tau_2 \sigma <: \tau'_2 \sigma$  therefore from Lemma 2.52 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E)$ :

Similar reasoning as in the previous case

## 2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[((\tau_1 \times \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

IH1:  $[(\tau_1 \sigma)]_V^A \subseteq [(\tau'_1 \sigma)]_V^A$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_V^A \subseteq [(\tau'_2 \sigma)]_V^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A$

Therefore from Definition 2.33 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (156)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_{\mathcal{V}}^A$

Again from Definition 2.33, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

Since from Equation 156 we know that  $(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$

Similarly since  $(W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A$  from Equation 156 therefore from IH2 we have  $(W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$

### 3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $[((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A \subseteq [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

IH1:  $[(\tau_1 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_1 \sigma)]_{\mathcal{V}}^A$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_2 \sigma)]_{\mathcal{V}}^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s1} = \text{inl } v_{i2}$ :

From Definition 2.33 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (157)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

From Equation 157 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

(b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 2.33 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (158)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

From Equation 158 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $[((\forall \alpha. \tau_1) \sigma)]_V^A \subseteq [(\forall \alpha. \tau_2) \sigma]_V^A$

$\forall \sigma. [(\tau_1 \sigma)]_E^A \subseteq [(\tau_2 \sigma)]_E^A$  (Sub-F2, From Statement (2))

It suffices to prove:  $\forall (W, n, \Lambda e_1, \Lambda e_2) \in [((\forall \alpha. \tau_1) \sigma)]_V^A$ .

$(W, n, \Lambda e_1, \Lambda e_2) \in [((\forall \alpha. \tau_2) \sigma)]_V^A$

This means that given:  $(W, n, \Lambda e_1, \Lambda e_2) \in [((\forall \alpha. (\tau_1)) \sigma)]_V^A$

Therefore from Definition 2.33 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E) \quad (\text{Sub-F1}) \end{aligned}$$

And it suffices to prove:  $(W, n, \Lambda e_1, \Lambda e_2) \in [((\forall \alpha. \tau_2) \sigma)]_V^A$

Again from Definition 2.33, it suffices to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W. \theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E) \wedge \\ \forall \theta'_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E) \end{aligned}$$

This means we are required to show:

(a)  $\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$ :

By instantiating the first conjunct of Sub-F1 with  $W'', n''$  and  $\ell''$  we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from Sub-F2 instantiated at  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

(b)  $\forall \theta'_l \sqsupseteq W. \theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E)$ :

By instantiating the second conjunct of Sub-F1 with  $\theta'_l$  and  $\ell''$  we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E)$$

Since  $\tau_1 \sigma \cup \{\alpha \mapsto \ell''\} <: \tau_2 \sigma \cup \{\alpha \mapsto \ell''\}$  therefore from Lemma 2.52 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E)$$

(c)  $\forall \theta'_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E)$ :

Similar reasoning as in the previous case

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2}$$



To prove:  $\lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A \subseteq \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

$\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$  (Sub-C0, From Statement (2))

It suffices to prove:  $\forall (W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A. (W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A$

Therefore from Definition 2.33 we are given:

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma \implies (W', n', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_1) \in \lfloor \tau_1 \sigma \rfloor_E \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_2) \in \lfloor \tau_1 \sigma \rfloor_E$  (Sub-C1)

And it suffices to prove:  $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

Again from Definition 2.33, it suffices to prove:

$\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_1) \in \lfloor \tau_2 \sigma \rfloor_E \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E$

This means that we are required to show the following:

(a)  $\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$ :

We are given  $W'' \sqsupseteq W, n'' < n$  also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with  $W''$  and  $n''$  we know that the following holds

$(W'', n'', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A$

Therefore from (Sub-C0) we get  $(W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$

(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \implies (\theta'_l, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E$ :

We are given some  $\theta'_l \sqsupseteq W.\theta_1, k$ , also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with  $\theta'_l$  we know that the following holds

$(\theta'_l, k, e_1) \in \lfloor \tau_1 \sigma \rfloor_E$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  therefore from Lemma 2.52 we get

$(\theta'_l, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E$ :

Similar reasoning as in the previous case

## 6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\lceil ((\text{Labeled } \ell \tau) \sigma) \rceil_V^A \subseteq \lceil ((\text{Labeled } \ell' \tau') \sigma) \rceil_V^A$

IH:  $\lceil (\tau \sigma) \rceil_V^A \subseteq \lceil (\tau' \sigma) \rceil_V^A$

It suffices to prove:  $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_V^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^A$

This means we are given  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_V^A$

From Definition 2.33 it means we have  $\text{ValEq}(\mathcal{A}, W, \ell \sigma, n, v_1, v_2, \tau \sigma)$  (Sub-L0)

and it suffices to prove  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^A$

Again from Definition 2.33 it means we need to prove that

$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

Since we have (Sub-L0) and  $\ell \sigma \sqsubseteq \ell' \sigma$  therefore from Lemma 2.55 we have

$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau \sigma)$

2 cases arise:

(a)  $\ell' \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 we know that  $(W, n, v_1, v_2) \in [\tau \sigma]_V^A$

From IH we also know that  $(W, n, v_1, v_2) \in [\tau' \sigma]_V^A$

And from Definition 2.33 we get  $\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

(b)  $\ell' \sigma \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.32 we know that  $\forall j. (W.\theta_1, j, v_1) \in [\tau \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau \sigma]_V$

Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.52 we get  $(W.\theta_1, j, v_1) \in [\tau' \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau' \sigma]_V$

And from Definition 2.33 we get  $\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $[((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$

IH:  $[(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$

It suffices to prove:  $\forall (W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A. (W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$

This means we are given  $(W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A$

From Definition 2.33 it means we have

$(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j.$

$(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$

$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma)) \wedge$

$\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge$

$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \Downarrow \ell_i \sigma))$  (Sub-CG0)

And we need to prove

$$(W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$$

Again from Definition 2.33 it means we need to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) \right) \end{aligned}$$

It means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau' \sigma): \end{aligned}$$

This means we are given  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t

$$(k, H_1, H_2) \triangleright W_e, (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma)$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1)  $\text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma)$

Therefore from Lemma 2.55 we get  $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau \sigma)$

$$\begin{aligned} \text{(b)} \quad & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right): \end{aligned}$$

Case  $l = 1$

Here we are given  $k, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

$$\text{i. } \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau' \sigma]_V:$$

Instantiating the second conjunct of (Sub-CG0) with the given  $k, \theta_e, H, j$  to get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau \sigma]_V$$

Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.52 we get  $(\theta', k - j, v'_1) \in [\tau' \sigma]_V$

$$\text{ii. } (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell'):$$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell')$$

Since  $\ell'_i \sigma \sqsubseteq \ell_i \sigma$  therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell')$$

- iii.  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$ :  
 Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get  
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$   
 Since  $\ell'_i \sigma \sqsubseteq \ell_i \sigma$  therefore we also get  
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$

Case  $l = 2$

Symmetric reasoning as in the previous  $l = 1$  case

8. CGsub-base:

Trivial

Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [\tau \sigma]_E^A. (W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$$

This means given  $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

From Definition 2.34 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A \quad (\text{Sub-E0})$$

And it suffices to prove  $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

Again from Definition 2.34 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$$

This means that given  $i < n$  s.t  $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$  we need to prove  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$

Instantiating (Sub-E0) with the given  $i$  we get  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$

From Statement (1) we get  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$  □

**Theorem 2.57** (CG: NI).  $\forall v_1, v_2, e, \tau, n.$

$$(\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \mathbf{b}]_V^\perp \wedge$$

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v'_2) \implies v'_1 = v'_2$$

*Proof.* Given some  $v_1, v_2, e, \tau, H_1, H_2, W, n, j < n$  s.t

$$(\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \mathbf{b}]_V^\perp \wedge$$

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v'_2)$$

We need to prove

$$v'_1 = v'_2$$

Since we are given  $(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp$

Therefore from Definition 2.34 we know that

$$\forall i < n. e_1[v_1/x] \Downarrow_i v_{11} \wedge e_2 \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\mathbb{C} \perp \perp \mathbf{b}]_V^\perp$$

From CG-val we know that  $i = 0$ ,  $v_{11} = e[v_1/x]$  and  $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_V^\perp$$

From Definition 2.35 we have

$$\begin{aligned}
& \left( \forall k \leq n, W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right. \\
& \forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow^f (H_2', v_2'') \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \right) \wedge \\
& \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\mathbf{b}]_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp) \right)
\end{aligned}$$

Instantiating the first conjunct with  $n, \emptyset, \emptyset, \emptyset$ .

Since we know that

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v_1') \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v_2')$$

Therefore we instantiate  $v_1''$  with  $v_1'$ ,  $v_2''$  with  $v_2'$ ,  $j$  with  $n'$  to get

$$\exists W' \sqsupseteq \emptyset. (n - n', H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b})$$

From Definition 2.32 and Definition 2.35 we get  $v_1' = v_2'$

□

## 2.3 CG to FG translation

### 2.3.1 Type directed translation from CG to FG

CG types are translated into FG types by the following definition of  $\llbracket \cdot \rrbracket$

$$\begin{array}{ll}
\llbracket \mathbf{b} \rrbracket = \mathbf{b}^\perp & \llbracket \text{ref } \ell \ \tau \rrbracket = (\text{ref } (\llbracket \tau \rrbracket) + \text{unit})^\ell{}^\perp \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \mathbb{C} \ \ell_i \ \ell_o \ \tau \rrbracket = (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket) + \text{unit})^{\ell_o}{}^\perp \\
\llbracket \tau_1 \times \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp & \llbracket c \Rightarrow \tau \rrbracket = (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \\
\llbracket \tau_1 + \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp & \llbracket \forall \alpha. \tau \rrbracket = (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \\
\llbracket \text{Labeled } \ell \ \tau \rrbracket = (\llbracket \tau \rrbracket)^\ell
\end{array}$$

The translation judgment for expressions is of the form  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$ . Its rules are shown below.

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{var} \\
\frac{\Sigma; \Psi; \Gamma, x : \tau \vdash e : \tau' \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \lambda x. e : \tau \rightarrow \tau' \rightsquigarrow \lambda x. e_F} \text{lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau \rightarrow \tau' \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash e_1 \ e_2 : \tau' \rightsquigarrow e_{F1} \ e_{F2}} \text{app} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2 \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{F1}, e_{F2})} \text{prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e) : \tau_1 \rightsquigarrow \text{fst}(e_F)} \text{fst} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{snd}(e) : \tau_1 \rightsquigarrow \text{snd}(e_F)} \text{snd} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inl}(e_F)} \text{inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inr}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inr}(e_F)} \text{inr} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 + \tau_2 \rightsquigarrow e_F \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{case}(e_F, x.e_{F1}, y.e_{F2})} \text{case} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \ \tau) \rightsquigarrow \text{inl}(e_F)} \text{label} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \ \ell_i \ (\ell_i \sqcup \ell) \ \tau \rightsquigarrow \lambda_. e_F} \text{unlabel} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ \ell_i \ \ell_o \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ \ell_i \ \ell_i \ (\text{Labeled } \ell_o \ \tau) \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{toLabeled} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ \ell_i \ \ell_i \ \tau \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ret}
\end{array}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau' \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{bind} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda_.\text{inl}(\text{new } (e_F))} \text{ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_.\text{inl}(e_F)} \text{deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit} \rightsquigarrow \lambda_.\text{inl}(e_{F1} := e_{F2})} \text{assign} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F} \text{sub} \\
\\
\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \Lambda e : \forall \alpha. \tau \rightsquigarrow \Lambda e_F} \text{FI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \rightsquigarrow e_F \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha] \rightsquigarrow e_F[]} \text{FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \nu e : c \Rightarrow \tau \rightsquigarrow \nu e_F} \text{CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau \rightsquigarrow e_F \bullet} \text{CE}
\end{array}$$

### 2.3.2 Type preservation for CG to FG translation

**Assumption 2.58.**  $\forall e, \tau, \Sigma, \Psi, \Gamma, \ell_i, \ell_o.$

$$\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell_i \ell_o \tau \implies \ell_i \sqsubseteq \ell_o$$

**Theorem 2.59** (CG  $\rightsquigarrow$  FG: Type preservation).  $\forall \Sigma, \Psi, \Gamma, e_C, \tau.$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \text{ is a valid typing derivation in CG} \implies$$

$$\exists e_F.$$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$$

$$\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau] \text{ is a valid typing derivation in FG}$$

*Proof.* Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\perp}} \text{FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\ell}} \text{FG-sub}$$

2. unlabel:

P1:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i \sqcup \ell \quad \frac{\overline{\Sigma; \Psi \vdash ([\tau] + \text{unit}) <: ([\tau] + \text{unit})}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell} <: ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{Lemma 2.1}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell} <: ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{FGsub-label}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : ([\tau] + \text{unit})^\ell} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top \quad P1}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} \lambda_{-}.e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_i \sqcup \ell})^\perp} \text{FG-lam}}$$

3. toLabeled:

P2:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_o})^\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_o})^\perp} \text{FG-sub}}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} () : \text{unit}} \text{P2} \quad \Sigma; \Psi \vdash \ell_i \sqcup \perp \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F() : ([\tau] + \text{unit})^{\ell_o}} \text{FG-app}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{inl}(e_F()) : (([\tau] + \text{unit})^{\ell_o} + \text{unit})^{\ell_i}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_F()) : (\text{unit} \xrightarrow{\ell_i} (([\tau] + \text{unit})^{\ell_o} + \text{unit})^{\ell_i})^\perp} \text{FG-lam}}$$

4. ret:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : [\tau]} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : [\tau]} \text{FG-sub} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{inl}(e_F) : ([\tau] + \text{unit})^{\ell_i}} \text{FG-sub, FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_F) : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_i})^\perp}$$

5. bind:

P1.1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F1} : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^\ell)^\perp} \text{IH1, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_{F1} : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^\ell)^\perp} \text{FG-sub}}$$

P1:

$$\frac{P1.1 \quad \frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} () : \text{unit}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi \vdash (\ell_i \sqcup \perp) \sqsubseteq \ell_i \quad \frac{}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell \searrow \perp} \text{FG-app}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_{F1}() : ([\tau] + \text{unit})^\ell} \text{FG-app}$$



P2.1:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}}{\Sigma; \Psi \vdash \ell \sqsubseteq \top} \text{IH2, Weakening}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}} \text{FG-sub}$$

P2:

$$\frac{\frac{\text{P2.1} \quad \frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi \vdash \perp \sqsubseteq \ell_o} \text{FG-var}}{\Sigma; \Psi \vdash (\ell \sqcup \perp) \sqsubseteq \ell} \quad \frac{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell \sqcup \perp} e_{F2}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app}$$

P3:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_o}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-sub, FG-inr}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; \Gamma \vdash e_2 : \mathbb{C} \ell \ell_o \tau}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_o} \text{Assumption 2.58}}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \ell} \text{FG-case}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-lam, weak}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\text{unit} \xrightarrow{\ell_i} ([\tau'] + \text{unit})^{\ell_o})^{\perp}} \text{FG-lam, weak}$$

6. ref:

P1:

$$\frac{\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : ([\tau] + \text{unit})^{\ell'} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_F : ([\tau] + \text{unit})^{\ell'}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{FG-sub}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow \ell} \text{FG-ref}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{new } e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp}} \text{FG-ref}$$

Main derivation:

$$\frac{\frac{\text{P1} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(\text{new } e_F) : ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(\text{new } e_F) : (\text{unit} \xrightarrow{\ell} ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell})^{\perp}} \text{FG-lam}$$

7. deref:

P2:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : (\text{ref}([\tau] + \text{unit})^{\ell})^{\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell' \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} e_F : (\text{ref}([\tau] + \text{unit})^{\ell})^{\perp}} \text{FG-sub}}$$

P1:

$$P2 \quad \frac{\frac{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell \text{ Lemma 2.1} \quad \Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell \searrow \perp}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} !e_F : ([\tau] + \text{unit})^\ell} \text{FG-deref}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} !e_F : ([\tau] + \text{unit})^\ell}$$

Main derivation:

$$P1 \quad \frac{\frac{\Sigma; \Psi \vdash \perp \sqsubseteq \ell' \quad \frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'}} \text{FG-lam}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\ell'} (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'})^\perp}$$

8. assign:

P3:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F2} : ([\tau] + \text{unit})^{\ell'} \text{ IH2, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}}$$

P2:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp \text{ IH1, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp}$$

P1:

$$P2 \quad P3 \quad \frac{\frac{\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp)} \text{FG-assign} \quad \text{Given}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}} \text{FG-assign}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\ell} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^\ell)^\perp} \text{FG-lam}}$$

9. sub:

$$\frac{\frac{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau'] \text{ IH} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top \quad \frac{\Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi \vdash [\tau'] <: [\tau]} \text{Lemma 2.60}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{FG-sub}}$$

10. FI:

$$\frac{\frac{\Sigma, \alpha; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau] \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp} \text{FG-FI}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp}$$

11. FE:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (\forall \alpha. (\top, [\tau]))^{\perp}} \text{ IH}}{\text{FV}(\ell) \in \Sigma \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau[\ell/\alpha]] \searrow \perp} \text{ FG-FE}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F [] : [\tau][\ell/\alpha]} \text{ FG-FE}$$

12. CI:

$$\frac{\overline{\Sigma; \Psi, c; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \nu e_F : (c \overset{\top}{\Rightarrow} [\tau])^{\perp}} \text{ FG-CI}$$

13. CE:

$$\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (c \overset{\top}{\Rightarrow} [\tau])^{\perp}} \text{ IH} \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau] \searrow \perp}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F \bullet : [\tau]} \text{ FG-CE}$$

□

**Lemma 2.60** (CG  $\rightsquigarrow$  FG: Subtyping). *For any CG types  $\tau$  and  $\tau'$ ,  $\Sigma$ , and  $\Psi$ , if  $\Sigma; \Psi \vdash \tau <: \tau'$ , then  $\Sigma; \Psi \vdash [\tau] <: [\tau']$ .*

*Proof.* Proof by induction on CG's subtyping relation

1. CGsub-base:

$$\overline{\Sigma; \Psi \vdash [\tau] <: [\tau]} \text{ Lemma 2.1}$$

2. CGsub-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau'_1] <: [\tau_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau'_2] <: [\tau'_2]} \text{ IH2} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top}{\Sigma; \Psi \vdash ([\tau_1] \overset{\top}{\rightarrow} [\tau_2])^{\perp} <: ([\tau'_1] \overset{\top}{\rightarrow} [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \overset{\ell_s}{\rightarrow} \tau_2)) <: ((\tau'_1 \overset{\ell'_s}{\rightarrow} \tau'_2))} \text{ Definition of } [\cdot]$$

3. CGsub-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\Sigma; \Psi \vdash ([\tau_1] \times [\tau_2])^{\perp} <: ([\tau'_1] \times [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \times \tau_2)) <: ((\tau'_1 \times \tau'_2))} \text{ Definition of } [\cdot]$$

4. CGsub-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\Sigma; \Psi \vdash ([\tau_1] + [\tau_2])^{\perp} <: ([\tau'_1] + [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 + \tau_2)) <: ((\tau'_1 + \tau'_2))} \text{ Definition of } [\cdot]$$

5. CGsub-labeled:

$$\begin{array}{c}
\frac{\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})}}{\text{FGsub-sum}}} \quad \text{FGsub-sum} \\
\frac{\frac{\overline{\text{Labeled } \ell_1 \tau_1 <: \text{Labeled } \ell'_1 \tau'_1}}{\text{Given}}}{\frac{\ell_1 \sqsubseteq \ell'_1}{\text{By inversion}}} \quad \text{By inversion} \\
\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_1}}}{\text{FGsub-arrow}} \quad \text{FGsub-arrow} \\
\hline
\overline{\Sigma; \Psi \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket} \quad \text{Definition of } \llbracket \cdot \rrbracket
\end{array}$$

6. CGsub-monad:

P3:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} <: \llbracket \tau'_1 \rrbracket}}{\text{IH}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})}}{\text{FGsub-sum}}} \quad \text{FGsub-sum}$$

P2:

$$\frac{P3 \quad \frac{\overline{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\frac{\overline{\Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}}{\text{By inversion}}} \quad \text{By inversion} \quad \text{FGsub-label} \\
\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o}}}{\text{FGsub-label}}$$

P1:

$$\frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \quad P2 \quad \frac{\overline{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\frac{\overline{\Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i}}{\text{FGsub-arrow}}} \quad \text{FGsub-arrow} \\
\overline{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o}) <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma; \Psi \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}}}{\frac{\overline{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})^\perp}}{\text{Definition of } \llbracket \cdot \rrbracket}} \quad \text{Definition of } \llbracket \cdot \rrbracket \\
\overline{\Sigma; \Psi \vdash \llbracket \mathbb{C} \ell_i \ell_o \tau_1 \rrbracket <: \llbracket \mathbb{C} \ell'_i \ell'_o \tau'_1 \rrbracket}$$

7. CGsub-forall:

P1:

$$\frac{\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket} <: \llbracket \tau' \rrbracket}}{\text{IH, Weakening}} \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \top \sqsubseteq \top}}{\text{FGsub-forall}}}{\frac{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))}}{\text{FGsub-forall}}} \quad \text{FGsub-forall}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}}}{\frac{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))^\perp}}{\text{FGsub-label}}} \quad \text{FGsub-label} \\
\overline{\Sigma; \Psi \vdash \llbracket \forall \alpha. \tau \rrbracket <: \llbracket \forall \alpha. \tau' \rrbracket}$$

8. CGsub-constraint:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket} <: \llbracket \tau' \rrbracket} \text{IH} \quad \frac{\overline{\Sigma; \Psi \vdash \top \sqsubseteq \top} \quad \frac{\overline{\Sigma; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given}}{\overline{\Sigma; \Psi \vdash c' \Longrightarrow c}} \text{By inversion}}{\overline{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket) <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)}} \text{FGsub-constra}}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\overline{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket)^\perp <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)^\perp}} \text{FGsub-label}}{\overline{\Sigma; \Psi \vdash \llbracket c \Rightarrow \tau \rrbracket <: \llbracket c' \Rightarrow \tau' \rrbracket}}}$$

□

**Lemma 2.61** (CG  $\rightsquigarrow$  FG: Preservation of well-formedness).  $\forall \Sigma, \Psi, \tau$ .

$$\Sigma; \Psi \vdash \tau \text{ WF} \Longrightarrow \Sigma; \Psi \vdash \llbracket \tau \rrbracket \text{ WF}$$

*Proof.* Proof by induction on the  $\tau$  WF relation.

1. CG-wff-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base}}{\overline{\Sigma; \Psi \vdash \mathbf{b}^\perp \text{ WF}} \text{FG-wff-label}}$$

2. CG-wff-unit:

$$\overline{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit}$$

3. CG-wff-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \text{IH2}}{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket) \text{ WF}} \text{FG-wff-arrow}}}{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \text{ WF}} \text{FG-wff-label}}$$

4. CG-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \text{IH2}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket \text{ WF}} \text{FG-wff-prod}}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \rrbracket \text{ WF}} \text{FG-wff-label}}$$

5. CG-wff-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \text{IH2}}{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket \rrbracket \text{ WF}} \text{FG-wff-prod}}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \rrbracket \text{ WF}} \text{FG-wff-label}}$$

6. CG-wff-ref:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ Given} \\
\frac{}{\text{FV}(\tau) = \emptyset} \text{ By inversion} \\
\hline
\text{FV}(\llbracket \tau \rrbracket) = \emptyset \quad \text{Lemma 2.62} \\
\hline
\frac{}{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ Given} \\
\frac{}{\text{FV}(\text{unit}) = \emptyset} \text{ By inversion} \\
\hline
\frac{}{\Sigma; \Psi \vdash \text{FV}((\llbracket \tau \rrbracket + \text{unit})^\ell) = \emptyset} \text{ FG-wff-ref} \\
\frac{}{\Sigma; \Psi \vdash \text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF} \text{ FG-wff-label} \\
\hline
\Sigma; \Psi \vdash (\text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp \ WF
\end{array}$$

7. CG-wff-forall:

$$\begin{array}{c}
\frac{}{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \\
\frac{}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) \ WF} \text{ FG-wff-forall} \\
\hline
\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

8. CG-wff-constraint:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi, c \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \\
\frac{}{\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket) \ WF} \text{ FG-wff-constraint} \\
\hline
\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

9. CG-wff-labeled:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \\
\frac{}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF} \text{ FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF \quad \text{CG-wff-label}
\end{array}$$

10. CG-wff-monad:

P1:

$$\frac{}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF \quad \text{FG-wff-sum}$$

Main derivation:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \quad \frac{}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_o} \ WF} \text{ FG-wff-label} \\
\frac{}{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \text{unit})^{\ell_o}) \ WF} \text{ FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \text{unit})^{\ell_o})^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

□

**Lemma 2.62** (CG  $\rightsquigarrow$  FG: Free variable lemma).  $\forall \tau. FV(\llbracket \tau \rrbracket) \subseteq FV(\tau)$

*Proof.* Proof by induction on the CG types,  $\tau$

1.  $\tau = \mathbf{b}$ :

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ = & FV(\mathbf{b}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & \emptyset \\ = & FV(\mathbf{b}) \end{aligned}$$

2.  $\tau = \mathbf{unit}$ :

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ = & FV(\mathbf{unit}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & \emptyset \\ = & FV(\mathbf{unit}) \end{aligned}$$

3.  $\tau = \tau_1 \rightarrow \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 \rightarrow \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \rightarrow \tau_2) \end{aligned}$$

4.  $\tau = \tau_1 \times \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 \times \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \times \tau_2) \end{aligned}$$

5.  $\tau = \tau_1 + \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 + \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 + \tau_2) \end{aligned}$$

6.  $\tau = \mathbf{ref} \ell_i \tau_i$ :

$$\begin{aligned} & FV(\llbracket \mathbf{ref} \ell_i \tau_i \rrbracket) \\ = & FV(\mathbf{ref} (\llbracket \tau_i \rrbracket + \mathbf{unit})^{\ell_i})^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_i \rrbracket) \cup FV(\ell_i) \\ \subseteq & FV(\tau_i) \cup FV(\ell_i) \quad \text{IH} \\ = & FV(\mathbf{ref} \ell_i \tau_i) \end{aligned}$$

7.  $\tau = \forall \alpha. \tau_i$ :

$$\begin{aligned} & FV(\llbracket \forall \alpha. \tau_i \rrbracket) \\ = & FV(\forall \alpha. (\top, \llbracket \tau_i \rrbracket))^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_i \rrbracket) - \{\alpha\} \\ \subseteq & FV(\tau_i) - \{\alpha\} \quad \text{IH} \\ = & FV(\forall \alpha. \tau_i) \end{aligned}$$

8.  $\tau = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket c \Rightarrow \tau_i \rrbracket) \\
= & \text{FV}(c \xrightarrow{\top} \llbracket \tau_i \rrbracket)^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\llbracket \tau_i \rrbracket) \\
\subseteq & \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\tau_i) && \text{IH} \\
= & \text{FV}(c \Rightarrow \tau_i)
\end{aligned}$$

9.  $\tau = \text{Labeled } \ell_i \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \text{Labeled } \ell_i \tau_i \rrbracket) \\
= & \text{FV}(\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i} && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \\
\subseteq & \text{FV}(\tau_i) \cup \text{FV}(\ell_i) && \text{IH} \\
= & \text{FV}(\text{Labeled } \ell_i \tau_i)
\end{aligned}$$

10.  $\tau = \mathbb{C} \ell_i \ell_o \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \mathbb{C} \ell_i \ell_o \tau_i \rrbracket) \\
= & \text{FV}(\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_o})^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) \\
\subseteq & \text{FV}(\tau_i) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) && \text{IH} \\
= & \text{FV}(\mathbb{C} \ell_i \ell_o \tau_i)
\end{aligned}$$

□

### 2.3.3 Logical relation for CG to FG translation

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type})) \times (\text{Loc} \leftrightarrow \text{Loc})$

**Definition 2.63** (CG  $\rightsquigarrow$  FG:  ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq \forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

**Definition 2.64** (CG  $\rightsquigarrow$  FG:  $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq \forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$



**Definition 2.65** (CG  $\rightsquigarrow$  FG: Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\
[\mathbf{unit}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\
&\quad (s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 + \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}}\} \cup \\
&\quad \{(s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \\
&\quad \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_V^{\hat{\beta}'}\} \\
[\forall \alpha. \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \Lambda e_s, \Lambda e_t) \mid \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_V^{\hat{\beta}'}\} \\
[c \Rightarrow \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \nu e_s, \nu e_t) \mid \mathcal{L} \models c \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau]_V^{\hat{\beta}'}\} \\
[\text{ref } \ell \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s a, {}^t a) \mid {}^s \theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\
[\text{Labeled } \ell \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid \\
&\quad \exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge (s\theta, m, {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}\} \\
[\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
&\quad (k, H_s, H_t) \hat{\triangleright} ({}^s \theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \\
&\quad \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''. (k - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge \\
&\quad \exists {}^t v''. {}^t v = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''}\}
\end{aligned}$$

**Definition 2.66** (CG  $\rightsquigarrow$  FG: Unary expression relation).

$$\begin{aligned}
[\tau]_E^{\hat{\beta}} &\triangleq \{(s\theta, n, e_s, e_t) \mid \\
&\quad \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s \theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies \\
&\quad \exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge (s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta\}
\end{aligned}$$

**Definition 2.67** (CG  $\rightsquigarrow$  FG: Unary heap well formedness).

$$\begin{aligned}
(n, H_s, H_t) \hat{\triangleright} {}^s \theta &\triangleq \text{dom}({}^s \theta) \subseteq \text{dom}(H_s) \wedge \\
&\quad \hat{\beta} \subseteq (\text{dom}({}^s \theta) \times \text{dom}(H_t)) \wedge \\
&\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s \theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s \theta(a)]_V^{\hat{\beta}}
\end{aligned}$$

**Definition 2.68** (CG  $\rightsquigarrow$  FG: Label substitution).  $\sigma : \text{Lvar} \mapsto \text{Label}$

**Definition 2.69** (CG  $\rightsquigarrow$  FG: Value substitution to values).  $\delta^s : \text{Var} \mapsto \text{Val}, \delta^t : \text{Var} \mapsto \text{Val}$

**Definition 2.70** (CG  $\rightsquigarrow$  FG: Unary interpretation of  $\Gamma$ ).

$$\begin{aligned}
[\Gamma]_V^{\hat{\beta}} &\triangleq \{(s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\
&\quad \forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}\}
\end{aligned}$$

### 2.3.4 Soundness proof for CG to FG translation

**Lemma 2.71** (CG  $\rightsquigarrow$  FG: Monotonicity).  $\forall {}^s \theta, {}^s \theta', n, {}^s v, {}^t v, n', \beta, \beta'$ .

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}}$  therefore from Definition 2.65 we know that  ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 2.65  ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$  we get the desired

2. Case **unit**:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}}$  therefore from Definition 2.65 we know that  ${}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 2.65  ${}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket$  we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that  ${}^sv = ({}^sv_1, {}^sv_2)$  and  ${}^tv = ({}^tv_1, {}^tv_2)$ .

We also know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\text{IH1: } ({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\text{IH2: } ({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 2.65, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

$$\text{IH: } ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 2.65 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \rightarrow \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that

$$\forall {}^s\theta'' \sqsubseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 2.65 we are required to prove

$$\forall {}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some  ${}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$  s.t  $({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'}$  and we are required to prove

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with  ${}^s\theta'_1, {}^sv_2, {}^tv_2, j, \hat{\beta}''$  since  ${}^s\theta'_1 \sqsubseteq {}^s\theta' \sqsubseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case  $\forall\alpha.\tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\forall\alpha.\tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\forall \alpha. \tau]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 2.65 we know that  ${}^s v = \Lambda e'_s$  and  ${}^t v = \Lambda e'_t$ . And

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'' . ({}^s\theta'', j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''} \quad (\text{F0})$$

Similarly from Definition 2.65 we are required to prove

$$\forall {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

This means we are given some  ${}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

Instantiating (F0) with  ${}^s\theta''_1, j, \hat{\beta}''_1$  since  ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$  therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

7. Case  $c \Rightarrow \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [c \Rightarrow \tau]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \Rightarrow \tau]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 2.65 we know that  ${}^s v = \nu(e'_s)$  and  ${}^t v = \nu(e'_t)$ . And

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta'', j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 2.65 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

This means we are given some  $\mathcal{L} \models c, {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

Since  $\mathcal{L} \models c$  and instantiating (C0) with  ${}^s\theta''_1, j, \hat{\beta}''_1$  since  ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$  therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

8. Case ref  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \ell \tau]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \ell \tau]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that  ${}^sv = {}^sa$  and  ${}^tv = {}^ta$ . We also know that  ${}^s\theta({}^sa) = \text{Labeled } \ell \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}$

From Definition 2.65, Definition 2.63 and Definition 2.64 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \ell \tau]_V^{\hat{\beta}'}$$

9. Case Labeled  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}'}$$

From Definition 2.65 it means

$$\exists {}^sv', {}^tv'. {}^sv = \text{Lb}_\ell({}^sv') \wedge {}^tv = \text{inl } {}^tv' \wedge ({}^s\theta, n, {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}}$$

$$\underline{\text{IH:}} ({}^s\theta', n', {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}}$$

Similarly from Definition 2.65 we need to prove that

$$\exists {}^sv'', {}^tv''. {}^sv = \text{Lb}_\ell({}^sv'') \wedge {}^tv = \text{inl } {}^tv'' \wedge ({}^s\theta', n', {}^sv'', {}^tv'') \in [\tau]_V^{\hat{\beta}'}$$

We choose  ${}^sv''$  as  ${}^sv'$  and  ${}^tv''$  as  ${}^tv'$  and since from IH we know that  $({}^s\theta', n', {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}'}$

Therefore from Definition 2.65 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}'}$$

10. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}'}$$

This means from Definition 2.65 we know that

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^sv', {}^tv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1.$$

$$(k, H_s, H_t) \hat{\beta}_1 \triangleright ({}^s\theta_e) \wedge (H_s, {}^sv) \Downarrow_i^f (H'_s, {}^sv') \wedge i < k \implies$$

$$\exists {}^tv'. (H_t, {}^tv()) \Downarrow (H'_t, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H'_s, H'_t) \hat{\beta}_2 \triangleright {}^s\theta' \wedge$$

$$\exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', {}^t\theta', k - i, {}^sv', {}^tv'') \in [\tau]_V^{\hat{\beta}_2} \wedge$$

$$(\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \quad (\text{CG0})$$

Similarly from Definition 2.65 we need to prove

$$\begin{aligned}
& \forall {}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1. \\
& (k', H'_s, H'_t) \triangleright^{\hat{\beta}'_1} ({}^s\theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge i' < k' \implies \\
& \exists {}^t v'' . (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2 . (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'' \wedge \\
& \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a . H_s(a) \neq H'_s(a) \implies \exists \ell' . {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e) . {}^s\theta'(a) \searrow \ell_1)
\end{aligned}$$

This means we are given some  ${}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  s.t.  $(k', H'_s, H'_t) \triangleright ({}^s\theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge i' < k'$

And we need to prove

$$\begin{aligned}
& \exists {}^t v'' . (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2 . (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'' \wedge \\
& \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta'', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a . H_s(a) \neq H'_s(a) \implies \exists \ell' . {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e) . {}^s\theta'(a) \searrow \ell_1)
\end{aligned}$$

Instantiating (CG0) with  ${}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  we get the desired

□

**Lemma 2.72** (CG  $\rightsquigarrow$  FG: Unary monotonicity for  $\Gamma$ ).  $\forall {}^s\theta, {}^s\theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies ({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 2.70 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma) . ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 2.70 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma) . ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x \in \text{dom}(\Gamma) . ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x \in \text{dom}(\Gamma) . ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 2.71 we get

$$\forall x \in \text{dom}(\Gamma) . ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

□

**Lemma 2.73** (CG  $\rightsquigarrow$  FG: Unary monotonicity for  $H$ ).  $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$

From Definition 2.67 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 2.67 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 2.71 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

□

**Theorem 2.74** (CG  $\rightsquigarrow$  FG: Fundamental theorem).  $\forall \Gamma, \tau, e, \delta^s, \delta^t, \sigma, {}^s\theta, n.$

$$\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge$$

$$\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$$

$\implies$

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. CF-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

$$\text{Also given is: } \mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\} \sigma]_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 it suffices to prove that

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $x \delta^s \downarrow_i {}^s v$

From cg-val we know that  $i = 0$ ,  ${}^s v = x \delta^s$ .

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-V0})$$

From fg-val we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$ . So we are left with proving

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we are given  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau \sigma\} \sigma]_{V}^{\hat{\beta}}$ , therefore from Definition 2.70 we get

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_{V}^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \text{ in the context. So we are done.}$$

2. CF-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{ lam}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $(\lambda x. e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  ${}^s v = (\lambda x. e_s) \delta^s$ ,  ${}^t v = (\lambda x. e_t) \delta^t$ ,  $H'_t = H_t$  and  $i = 0$

It suffices to prove that

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\begin{aligned} & \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'} \\ & \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \end{aligned}$$

This means that we are given  ${}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

And we need to prove

$$({}^s \theta', j, e_s[{}^s v/x] \delta^s, e_t[{}^t v/x] \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \quad (\text{F-L0})$$



Since  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$  therefore from Lemma 2.72 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_1\}, e_t \cup \{x \mapsto {}^t v_1\}) \in [\tau_2 \sigma]_E^{\hat{\beta}'}$$
 s.t

$$({}^s\theta', j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$$

We get (F-L0) directly from IH

3. CF-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 it suffices to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

(F-A0)

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.66 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

(F-A1)

IH2:

$$({}^s\theta, n-j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.66 it suffices to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n-j, {}^s v_2.e_{s2} \Downarrow_i {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n-j-k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n-j-k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta'_2 \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i-j < n-j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

And we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n-j-k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n-j-k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A2})$$

Since from (F-A1) we know that  $({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}}$  where  ${}^s v_1 = \lambda x.e'_s$  and  ${}^t v_1 = \lambda x.e'_t$

From Definition 2.65 we have

$$\begin{aligned} \forall {}^s\theta'_3 \sqsupseteq {}^s\theta, {}^s v, {}^t v, l < n-j, \hat{\beta}_3 \sqsupseteq \hat{\beta}. ({}^s\theta'_3, l, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}_3} \\ \implies ({}^s\theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}_3} \end{aligned}$$

Instantiating with  ${}^s\theta, {}^s v_2, {}^t v_2, n-j-k, \hat{\beta}$  we get

$$({}^s\theta, n-j-k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} \forall H_{s4}, H_{t4}. (n-j-k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k' < n-j-k, {}^s v_4.e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n-j-k-k', {}^s v_4, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge \\ (n-j-k-k', H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t2}$ , from (F-A2) we know that  $(n-j-k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$ . Instantiating  ${}^s v_4$  with  ${}^s v$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k' < i-j-k < n-j-k$  s.t  $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$ . therefore we have

$$\exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n-j-k-k', {}^s v, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n-j-k-k', H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A3})$$

Since from cg-app we know that  $i = j + k + k'$  and  $H'_t = H'_{t4}$ ,  ${}^t v = {}^t v_4$  therefore we get (F-A0) from (F-A3) and Lemma 2.71 and Lemma 2.73

#### 4. CF-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{ prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t, \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \quad (\text{F-P0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-P1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j. e_{s2} \delta^s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}, \hat{\beta}'_1$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-P2})$$

From cg-prod we know that  $i = j + k + 1$ ,  $H'_t = H'_{t2}$  and  ${}^t v = ({}^t v_1, {}^t v_2)$  therefore from Definition 2.65 and Lemma 2.71 we get  $({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$

And since we have  $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^s \theta$  therefore from Lemma 2.73 we also get

$$(n - i, H_s, H'_{t2}) \hat{\triangleright}^s \theta$$

5. CF-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$  (F-F0)

This means from Definition 2.66 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v &\implies \\ \exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \quad (\text{F-F0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j ({}^s v_1, -) &\implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge \\ (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and  ${}^s v_1$  with  ${}^s v$  since we know that  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j ({}^s v, -)$ .

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge \\ (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-F1}) \end{aligned}$$

From cg-fst we know that  $i = j + 1$ ,  $H'_t = H'_{t1}$  and  ${}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$  therefore from Definition 2.65 and Lemma 2.71 we get  $({}^s\theta, n - i, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$

And since from (F-F1) we have  $(n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta$  therefore from Lemma 2.73 we get

$$(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$$

6. CF-snd:

Symmetric reasoning as in the CF-fst case

7. CF-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s), \delta^s, \text{inl}(e_t), \delta^t) \in [(\tau_1 + \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t.  $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge \\ ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t.  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL1})$$

From cg-inl we know that  $i = j + 1$  and  $H'_t = H'_{t1}, {}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from Definition 2.65 and Lemma 2.71 we get

$$({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}$$

And since from (F-IL1) we have  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 2.73 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \mathbf{case}(e_t, x.e_{t1}, y.e_{t2})} \text{ case}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  $\mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C0)}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C1)}$$

Two cases arise:

(a)  ${}^s v_1 = \mathbf{inl}({}^s v'_1)$  and  ${}^t v_1 = \mathbf{inl}({}^t v'_1)$ :

IH2:

$$({}^s\theta, n-j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright} {}^s\theta \wedge \forall k < n - j, {}^s v_2.e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - \\ & j - k, H_{s2}, H'_{t2}) \hat{\triangleright} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - \\ & j - k, H_s, H'_{t2}) \hat{\triangleright} {}^s\theta \end{aligned}$$

From cg-case1 we know that  $i = j + k + 1$  and  $H'_t = H'_{t2}, {}^t v = {}^t v_2$ . Since we know  $({}^s\theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$  therefore from Definition 2.65 and Lemma 2.71 we get  $({}^s\theta, n - i, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$

And since from (F-C2) we have  $(n - j - k, H_s, H'_{t2}) \hat{\triangleright} {}^s\theta$  therefore from Lemma 2.73 we get  $(n - i, H_s, H'_{t2}) \hat{\triangleright} {}^s\theta$

(b)  ${}^s v_1 = \text{inr}({}^s v'_1)$  and  ${}^t v_1 = \text{inr}({}^t v'_1)$ :

Symmetric reasoning as in the previous case

10. CF-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \Lambda e_s : \forall \alpha. \tau \rightsquigarrow \Lambda e_t} \text{FI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \Lambda e_s \delta^s, \Lambda e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 we know that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. \Lambda e_s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \Lambda e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  and given some  $i < n$  s.t  $(\Lambda e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s, {}^t v = (\Lambda e_t) \delta^t, i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s\theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s\theta$$

We know  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  from the context. So, we are only left to prove

$$({}^s\theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{\hat{\beta}'}^E$$

This means that we are given  ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{\hat{\beta}'}^E \quad (\text{F-FI0})$$

Since  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^V$  therefore from Lemma 2.72 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}'}^V$$

IH:

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_{\hat{\beta}'}^E$$

We get (F-FI0) directly from IH

#### 11. CF-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \forall \alpha. \tau \rightsquigarrow e_t \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e_s [] : \tau[\ell/\alpha] \rightsquigarrow e_t []} \text{FE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^V$

To prove:  $({}^s\theta, n, e_s [] \delta^s, e_t [] \delta^t) \in [\tau[\ell/\alpha] \sigma]_{\hat{\beta}}^E$

From Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v. e_s [] \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{\hat{\beta}}^V \wedge (n - i, H_s, H'_t) \triangleright {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $e_s [] \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{\hat{\beta}}^V \wedge (n - i, H_s, H'_t) \triangleright {}^s\theta \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_{\hat{\beta}}^E$$

This means from Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_{\hat{\beta}}^V \wedge (n - j, H_{s1}, H'_{t1}) \triangleright {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s []) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n, {}^s v_1$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have



$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$   
(F-FE1)

From cg-FE we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_1 = \Lambda e'_t$

Therefore we have

$({}^s \theta, n - j, \Lambda e'_s, \Lambda e'_t) \in [(\forall \alpha. \tau) \sigma]_{V}^{\hat{\beta}}$

This means from Definition 2.65 we have

$\forall {}^s \theta' \sqsubseteq {}^s \theta, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s \theta', k, e'_s, e'_t) \in [\tau[\ell'/\alpha] \sigma]_{E}^{\hat{\beta}_2}$

Instantiating  ${}^s \theta'$  with  ${}^s \theta$ ,  $k$  with  $n - j - 1$ ,  $\ell'$  with  $\ell$   $\sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$({}^s \theta, n - j - 1, e'_s, e'_t) \in [\tau[\ell/\alpha] \sigma]_{E}^{\hat{\beta}}$

From Definition 2.66 we get

$\forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s \theta'_1 \wedge \forall k < n - j - 1, {}^s v_2. e'_s \Downarrow_k {}^s v_2 \implies$   
 $\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - j - 1 -$   
 $k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-FE1) we know that  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$  therefore from Lemma 2.73 we get  $(n - j - 1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$

Since we know that  $e_s \Downarrow_i \delta^s$  and from cg-FE we know that  $i = j + k + 1$  (for some  $k$ ) and  $i < n$  therefore we have  $k < n - j - 1$  s.t  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - j - 1 -$   
 $k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$  (F-FE2)

Since  $H'_t = H_{t2'}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-FE0) directly from (F-FE2)

12. CF-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \nu e_s : c \Rightarrow \tau \rightsquigarrow \nu e_t} \text{ CI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \nu e_s \delta^s, \nu e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_{E}^{\hat{\beta}}$

This means from Definition 2.66 we know that

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n. \nu e_s \Downarrow_i {}^s v \implies$   
 $\exists H'_t, {}^t v. (H_t, \nu e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(c \Rightarrow \tau) \hat{\beta} \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t

$(\nu e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  ${}^t v = (\nu e_t) \delta^t$ ,  $i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\mathcal{L} \models c \sigma \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

This means that we are given  $\mathcal{L} \models c \sigma$  and  ${}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-CI0})$$

Since  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^{\hat{\beta}}$  therefore from Lemma 2.72 we also have

$$({}^s \theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

And since we know that  $\mathcal{L} \models c \sigma$  therefore

$$\underline{\text{IH:}} ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

### 13. CF-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : c \Rightarrow \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e_s \bullet : \tau \rightsquigarrow e_t \bullet} \text{CE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, e_s \bullet \delta^s, e_t \bullet \delta^t) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}}$

From Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v.e_s \bullet \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This further means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t  $e_s \bullet \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}_E}^{\hat{\beta}}$$

This means from Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}_V}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s \bullet) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}_V}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE1})$$

From cg-CE we know that  ${}^s v_1 = \nu e'_s$  and  ${}^t v_1 = \nu e'_t$

Therefore we have

$$({}^s\theta, n - j, \nu e'_s, \nu e'_t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}_V}^{\hat{\beta}}$$

This means from Definition 2.65 we have

$$\forall {}^s\theta' \sqsupseteq {}^s\theta'_1, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s\theta', k, e'_s, e'_t) \in [\tau \sigma]_{\hat{\beta}_E}^{\hat{\beta}_2}$$

Instantiating  ${}^s\theta'$  with  ${}^s\theta$ ,  $k$  with  $n - j - 1$ ,  $\ell'$  with  $\ell \sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$$({}^s\theta, n - j - 1, e'_s, e'_t) \in [\tau \sigma]_{\hat{\beta}_E}^{\hat{\beta}}$$

From Definition 2.66 we get

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s\theta'_1 \wedge \forall k < n - j - 1. e'_s \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{\beta}_V}^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-CE1) we know that  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 2.73 we get  $(n - j - 1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$

Since we know that  $e_s \bullet \delta^s \Downarrow_i {}^s v$  and from cg-CE we know that  $i = j + k + 1$  (for some  $k$ ) and  $i < n$  therefore we have  $k < n - j - 1$  s.t  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{\beta}_V}^{\hat{\beta}} \wedge (n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE2})$$

Since  $H'_t = H_{t2'}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-CE0) directly from (F-CE2)

14. CF-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e_s) : \mathbb{C} \ell_i \ell_i \tau \rightsquigarrow \lambda_{\cdot} \text{.inl}(e_t)} \text{ret}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}_V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{ret}(e_s) \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot}\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  $\text{ret}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot}\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

From CG-ret and FG-lam we know that  $i = 0$ ,  ${}^s v = \text{ret}(e_s) \delta^s$ ,  ${}^t v = \lambda_{\cdot}\text{inl}(e_t) \delta^t$  and  $H'_t = H_t$ .

So we need to prove

$$({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ & (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \exists H'_t, {}^t v'. (H_t, (\lambda_{\cdot}\text{inl}(e_t) ()) \delta^t) \Downarrow \\ & (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k. \text{ Also from cg-ret we know that } H'_s = H_s$$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, (\lambda_{\cdot}\text{inl}(e_t) ()) \delta^t) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-R0}) \end{aligned}$$

IIH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies \\ & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s1}$  with  $H_s$  and  $H_{t1}$  with  $H_t$ . And since we know that  $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ . Therefore we have

$$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s \theta_e, k-f, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \wedge (k-f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-R1})$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_{t1}$ ,  ${}^t v'$  as  $\text{inl}({}^t v)$ ,  ${}^s \theta'$  as  ${}^s \theta_e$ ,  $\hat{\beta}''$  as  $\hat{\beta}'$ . Since from cg-ret we know that  $i = f + 1$  therefore from (F-R1) and Lemma 2.73 we know that  $(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

Next we choose  ${}^t v''$  as  ${}^t v$  (from F-R1) and from Lemma 2.71 we get  $({}^s \theta_e, k - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'}$  (we know from cg-ret that  ${}^s v' = {}^s v$ )

15. CF-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \mathbb{C} \ell_i \ell \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_{s2} : \mathbb{C} \ell \ell_o \tau' \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{bind}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s,$

$${}^t v = \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$$

And we need to prove

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda\_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsupseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \dots \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \\ & \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \dots \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - \\ & i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \quad (\text{F-B0}) \end{aligned}$$

IH1:

$$({}^s \theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_E^{\hat{\beta}}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \quad (\text{F-B1.1}) \end{aligned}$$

From Definition 2.65 we know have

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}' . \\ & (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta, H_{s3}$  with  $H_{s1}, H_{t3}$  with  $H'_{t2}, m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s \theta'', k - j - b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_E^{\hat{\beta}''}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta \wedge \forall c < (k - j - b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ & \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_{\hat{V}}^{\hat{\beta}''} \wedge (k - \\ & j - b - c, H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta'' \end{aligned}$$

Instantiating  $H_{s4}$  with  $H'_{s3}$  and  $H_{t4}$  with  $H'_{t3}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists c < i - j - b < k - j - b$  s.t.  $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$ .

Therefore we have

$$\exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_{\hat{V}}^{\hat{\beta}''} \wedge (k - j - b - c, H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta'' \quad (\text{F-B2.1})$$

From Definition 2.65 we know have

$$\begin{aligned} & \forall {}^s \theta_e \sqsubseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k - j - b - c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1. \\ & (m, H_{s5}, H_{t5}) \hat{\triangleright}^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsubseteq {}^s \theta_e, \hat{\beta}''_1 \sqsubseteq \hat{\beta}''_2. (m - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta''' \wedge \\ & \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta''$ ,  $H_{s5}$  with  $H'_{s3}$ ,  $H_{t5}$  with  $H'_{t3}$ ,  $m$  with  $k - j - b - c$  and  $\hat{\beta}''$  with  $\hat{\beta}''$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists d < i - j - b - c < k - j - b - c$  s.t.  $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsubseteq {}^s \theta_e, \hat{\beta}''_1 \sqsubseteq \hat{\beta}''_2. (k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta''' \wedge \\ & \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose  $H'_{t1}$  as  $H'_{t5}$  and  ${}^t v'$  as  ${}^t v'_{h2}$ . Next we choose  ${}^s \theta'$  as  ${}^s \theta'''$  and  $\hat{\beta}''$  as  $\hat{\beta}''_2$  (both chosen from (F-B2)). Also from cg-bind we know that in (F-B0)  $H'_{s1}$  will be  $H'_{s5}$ .

Since  $(k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta'''$  therefore Lemma 2.71 we get  $(k - i, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta'''$

Also since from (F-B2) we have  $\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$

Since  $i = j + b + c + d + 1$  therefore from Lemma 2.71 we get

$$\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$$

16. CF-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{label}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n$  s.t.  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ .

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$  therefore  $\exists j < i < n$  s.t.  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove  $({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$ . Since from cg-label we know that  $i = j + 1$ ,  ${}^s v = {}^s v_1$  and  ${}^t v = {}^t v_1$ . Therefore we get this from Definition 2.65, (F-LB1) and Lemma 2.71.

From Lemma 2.71 we get  $(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$

17. CF-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \mathbb{C} \ell_i \ell_o \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e_s) : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \rightsquigarrow \lambda_. \text{inl}(e_t ())} \text{toLabeled}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{E}}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove



$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda\_ \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n$  s.t  $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda\_ \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{toLabeled}(e_s) \delta^s,$   
 ${}^t v = (\lambda\_ \text{inl } e_t()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda\_ \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda\_ \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_ \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_ \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}''} \quad (\text{F-TL0}) \end{aligned}$$

IH:

$$({}^s \theta, k, e_s \delta^s, e_t \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$$

It means from Definition 2.66 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies$$

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k-j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k-j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-TL1.1})$$

From Definition 2.65 we know have

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k-j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies$$

$$\exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k-j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i-j < k-j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k-j-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k-j-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-TL1})$$

In order to prove (F-TL0) we choose  ${}^s \theta'$  as  ${}^s \theta''$  and  $\hat{\beta}'$  as  $\hat{\beta}''$  (both chosen from (F-TL2))

Also from cg-toLabeled and fg-inl, fg-app we know that  $H'_s = H'_{s3}$  and  $H'_t = H'_{t3}$ , and  ${}^s v' = {}^s v'_{h1}$ ,  ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 2.71

18. CF-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e_s) : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \rightsquigarrow \lambda_{\cdot} e_t} \text{ unlabel}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{E}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = \text{unlabel}(e_s) \delta^s$ ,  ${}^t v = \lambda_{\cdot} e_t \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} e_t)()) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} e_t)()) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \quad (\text{F-U0}) \end{aligned}$$

III:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{E}}^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}'} \wedge (k - \\ f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-U1})$$

In order to prove (F-U0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'$  as  ${}^t v_h$ ,  ${}^s \theta'$  as  ${}^s \theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From cg-unlabel and fg-app we also know that  $H'_{s1} = H_{s1}$  and  $H'_{t1} = H'_{t2}$

We need to prove

(a)  $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$ :

Since from (F-U1) we know that  $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

Therefore from Lemma 2.73 we also get  $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

(b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}'}$ :

Since from (F-U1) we have

$({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}'}$

This means from Definition 2.65 we know that

$\exists {}^s v_i, {}^t v_i . {}^s v_h = \text{Lb}_\ell({}^s v_i) \wedge {}^t v_h = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$  (F-U2)

Since we know that  ${}^t v' = {}^t v_h$  and since from (F-U2) we have  ${}^t v_h = \text{inl } {}^t v_i$ . Therefore from we choose  ${}^t v''$  as  ${}^t v_i$  to get the first conjunct

From cg-unlabel we know that  ${}^s v = {}^s v_i$  and since we know that  $({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$

Therefore from Lemma 2.71 we also get  $({}^s \theta_e, k - i, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$

19. CF-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e_s : \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda\_.\text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{new } e_s \delta^s, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^s v$  s.t  $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{s\theta}$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  from the context so we are left with proving

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda\_.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ .

$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$   
 $\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl}(\text{new } e_t))(\_)) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge$   
 $\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}''}$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k$ .

And we need to prove

$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl}(\text{new } e_t))(\_)) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge$   
 $\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-N0})$

From cg-ref we know that  ${}^s v' = a_s$  and from fg-ref, fg-inl we know that  ${}^t v' = \text{inl } a_t$ .

IH:

$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell' \tau) \sigma]_E^{\hat{\beta}'}$

It means from Definition 2.66 that we need to prove

$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies$   
 $\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k -$   
 $f, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k -$   
 $f, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-N1})$

In order to prove (F-N0) we choose  $H'_{t1}$  as  $H'_{t2} \cup \{a_t \mapsto {}^t v_h\}$ ,  ${}^t v$  as  $a_t$ ,  ${}^s\theta'$  as  ${}^s\theta_n$  where  ${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$

And we choose  $\hat{\beta}''$  as  $\hat{\beta}_n$  where  $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From cg-ref and fg-ref we also know that  $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

We need to prove

(a)  $(k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}_n} {}^s\theta_n$ :

From Definition 2.67 it suffices to prove that

- $\text{dom}({}^s\theta_n) \subseteq \text{dom}(H'_{s1})$ :

Since  $\text{dom}({}^s\theta_e) \subseteq \text{dom}(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$ )

And since we know that

${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$  and  $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

Therefore we get  $\text{dom}({}^s\theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s\theta_n) \times \text{dom}(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ )

And since we know that

${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$ ,  $H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\}$  and  $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

Therefore we get  $\hat{\beta}_n \subseteq (\text{dom}({}^s\theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s\theta_n, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_n(a)]_V^{\hat{\beta}_n}$ :

$\forall (a_1, a_2) \in \hat{\beta}_n$

–  $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-N1) we know that  $({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'}$

From Lemma 2.71 we get  $({}^s\theta_n, k - i - 1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$

–  $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$  therefore  
from Definition 2.67 we get

$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$

From Lemma 2.71 we get

$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_n(a_1)]_V^{\hat{\beta}'}$

- (b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$ :

We choose  ${}^t v''$  as  ${}^t v_h$  from (F-N1), fg-inl and fg-ref we know that  ${}^t v' = \text{inl } {}^t v_h$

In order to prove  $({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$ , from Definition 2.65 it suffices to prove that

${}^s\theta_n(a_s) = (\text{Labeled } \ell' \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}_n$

We get this by construction of  ${}^s\theta_n$  and  $\hat{\beta}_n$

20. CF-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{ref } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash !e_s : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda \_ . \text{inl}(e_t)} \text{deref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e_s \delta^s, \lambda \_ . \text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \delta^s \Downarrow_i {}^s v \implies$

$\exists H'_t. {}^t v. (H_t, \lambda \_ . \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t

$!e_s \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-} \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = !e_s \delta^s$ ,  ${}^t v = \lambda_{-} \text{inl}(e_t) \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, !e_s \delta^s, \lambda_{-} \text{inl}(e_t) \delta^t) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-D0})$$

IH:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies$$

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'_1$  as  $H'_{t2}(a)$  (where  ${}^t v_h = a_t$  from fg-deref),  ${}^s \theta'$  as  ${}^s \theta_e$  and we choose  $\hat{\beta}''$  as  $\hat{\beta}'$ .

From cg-deref we also know that  $H'_{s1} = H_{s1}$

We need to prove

(a)  $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$ :

Since from (F-D1) we have  $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$  and since  $f < i$  therefore from Lemma 2.73 we get  $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

(b)  $\exists^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$ :

Since from cg-deref and fg-deref we know that  ${}^s v_h = a_s$  and  ${}^t v_h = a_t$ .

Therefore from (F-D1) and from Definition 2.65 we know that

$${}^s \theta_e(a_s) = (\text{Labeled } \ell \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that  $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$  which means from Definition 2.67 we know that

$$({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-D2})$$

This means from Definition 2.65 we know that

$$\exists^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_{\ell}({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_{V}^{\hat{\beta}'}$$

We choose  ${}^t v''$  as  ${}^t v_i$  and we know that  ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$ . This proves the first conjunct.

Since from (F-D2) we have  $({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$  therefore from Lemma 2.71 we get

$$({}^s \theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$$

This proves the second conjunct.

21. CF-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_{s1} := e_{s2} : \mathbb{C} \ell \ell \text{ unit} \rightsquigarrow \lambda_{-} \text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$



From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = (e_{s1} := e_{s2}) \delta^s$ ,  ${}^t v = \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} \\ {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s \theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} \implies \\ \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - \\ f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-S1})$$

IH2:

$$({}^s \theta_e, k - f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k - f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} \implies \\ & \exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - \\ & f - l, H_{s3}, H'_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s3}$  with  $H_{s1}$  and  $H_{t3}$  with  $H'_{t2}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists l < i - f < k - f$  s.t.  $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$ .

Therefore we have

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - f - l, H_{s1}, H'_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S2})$$

In order to prove (F-S0) we choose  $H'_{t1}$  as  $H'_{t3}[a_t \mapsto {}^t v_{h3}]$ ,  ${}^t v'$  as  $()$ ,  ${}^s\theta'$  as  ${}^s\theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$ . From cg-assign and fg-assign we also know that  ${}^s v_{h2} = a_s$ ,  ${}^t v_{h2} = a_t$ ,  $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$  and  $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

$$(a) (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e:$$

From Definition 2.67 it suffices to prove that

- $dom({}^s\theta_e) \subseteq dom(H'_{s1})$ :

Since  $dom({}^s\theta_e) \subseteq dom(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$ )

And since  $dom(H_{s1}) = dom(H'_{s1})$  therefore we also get

$$dom({}^s\theta_e) \subseteq dom(H'_{s1})$$

- $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$ )

And since  $dom(H_{t1}) \subseteq dom(H'_{t1})$  therefore we also have  $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$ :

$$\forall (a_1, a_2) \in \hat{\beta}'_n$$

- $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-S2) we know that  $({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'}$

From Lemma 2.71 we get  $({}^s\theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'}$

- $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$  therefore

from Definition 2.67 we get

$$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$$

From Lemma 2.71 we get

$$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$$

$$(b) \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}'_n}:$$

We choose  ${}^t v''$  as  $()$  from (F-S1), fg-inl and fg-assign we know that  ${}^t v' = \text{inl } ()$

To prove:  $({}^s\theta_n, k - i, (), ()) \in [\text{unit}]_{V}^{\hat{\beta}'_n}$ ,

We get this directly from Definition 2.65

□

**Lemma 2.75** (CG  $\rightsquigarrow$  FG: Subtyping). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

1.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$
2.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

This means that given some  ${}^s\theta, n$  and  $\lambda x.e_i$  s.t  $({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$ .

$({}^s\theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}'}$  (S-A0)

And it suffices to prove:  $({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 2.65 it suffices to prove:

$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$ .

$({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1} \implies ({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

This means that given some  ${}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t  $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1}$

And we are required to prove:  $({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

IH:  $[(\tau'_1 \sigma)]_V^{\hat{\beta}'_1} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}'_1}$  (Statement (1))

$[(\tau_2 \sigma)]_E^{\hat{\beta}'_1} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}'_1}$  (Sub-A0, From Statement (2))

Instantiating (S-A0) with  ${}^s\theta'_1, {}^s v_1, {}^t v_1, k, \hat{\beta}'_1$

Since  $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_1}$  therefore from IH1 we know that  $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1}$

As a result we get

$({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$

From (Sub-A0), we know that

$({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_1}$

2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$\forall ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given  $({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s\theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$$

Since from (S-P0) we know that  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}}$  therefore from IH1 we have  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}}$

Similarly since from (S-P0) we have  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$  therefore from IH2 we get  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}}$

3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:  $\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^s v = \text{inl } {}^s v_i$  and  ${}^t v = \text{inl } {}^t v_i$ :

From Definition 2.65 we are given:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

(b)  ${}^s v = \text{inr } {}^s v_i$  and  ${}^t v = \text{inr } {}^t v_i$ :

Symmetric reasoning

#### 4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $[(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'} \quad (\text{S-F0})$$

And it suffices to prove:  $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$\forall {}^s \theta'_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

This means that given  ${}^s \theta_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:  $({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$

Instantiating (S-F0) with  ${}^s \theta_1, k, \ell'_1, \hat{\beta}'_1$  we get

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_1[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \quad (\text{Sub-F0, Statement (2)})$$

From (Sub-F0), we know that

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

#### 5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $\llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\mathcal{L} \models c_1 \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'} \quad (\text{S-C0})$$

And it suffices to prove:  $({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$\mathcal{L} \models c_2 \sigma \implies \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means that given  $\mathcal{L} \models c_2, {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

since we know that  $c_2 \implies c_1$  and since  $\mathcal{L} \models c_2 \sigma$  therefore  $\mathcal{L} \models c_1 \sigma$ . Next we instantiate (S-C0) with  ${}^s\theta'_1, k, \hat{\beta}'_1$  to get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

$$\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_1} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_1} \quad (\text{Sub-C0, Statement (2)})$$

Therefore from (Sub-C0), we get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$

IH:  $\llbracket (\tau \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in [\tau' \sigma]_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

## 7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_1 \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell'_2}{\Sigma; \Psi \vdash \mathbb{C} \ell_1 \ell_2 \tau <: \mathbb{C} \ell'_1 \ell'_2 \tau'}$$

$$\text{To prove: } [((\mathbb{C} \ell_i \ell_2 \tau) \sigma)]_V^{\hat{\beta}} \subseteq [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

It suffices to prove:

$$\forall ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

$$\text{This means that given } ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 2.65 we are given:

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies$$

$$\exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge$$

$$\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{S-M0})$$

And we are required to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

So again from Definition 2.65 we need to prove

$$\forall {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1 \implies$$

$$\exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} {}^s\theta' \wedge$$

$$\exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}$$

$$\text{This means we are given some } {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1 \text{ s.t. } (k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1}$$

$$({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1$$

And we need to prove

$$\exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} {}^s\theta' \wedge$$

$$\exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}$$

We instantiate (S-M0) with  ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$  we get

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_{e_1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

$$\text{IH: } [(\tau \sigma)]_V^{\hat{\beta}''} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}} \hat{\beta}'' \text{ (Statement (1))}$$

Since we have  $({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''}$  therefore from IH we get  $({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\tau' \sigma)]_V^{\hat{\beta}}$

8. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s \theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}} . ({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means that we are given  $({}^s \theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

From Definition 2.66 it means we have

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s \theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

From Definition 2.66 we need to prove

$$\forall H_{s_1}, H_{t_1}. (n, H_{s_1}, H_{t_1}) \hat{\triangleright} {}^s \theta \wedge \forall j < n, {}^s v_1. e_s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t_1}, {}^t v_1. (H_{t_1}, e_t) \Downarrow (H'_{t_1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s_1}, H'_{t_1}) \hat{\triangleright} {}^s \theta$$

This further means that given  $H_{s_1}, H_{t_1}$  s.t  $(n, H_{s_1}, H_{t_1}) \hat{\triangleright} {}^s \theta$ . Also given some  $j < n, {}^s v_1$  s.t  $e_s \Downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t_1}, {}^t v_1. (H_{t_1}, e_t) \Downarrow (H'_{t_1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s_1}, H'_{t_1}) \hat{\triangleright} {}^s \theta$$

Instantiating (Sub-E0) with the given  $H_{s_1}, H_{t_1}$  and  $j < n, {}^s v_1$ . We get

$$\exists H'_t, {}^t v. (H_{t_1}, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s_1}, H'_t) \hat{\triangleright} {}^s \theta$$

Since we have  $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}}$  therefore from Statement(1) we get  $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [(\tau' \sigma)]_V^{\hat{\beta}}$

□

**Theorem 2.76** (CG  $\rightsquigarrow$  FG: Deriving CG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s_1}, H'_{s_2}$ .

let  $\text{bool} = (\text{unit} + \text{unit})$ .

$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \wedge$

$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s_1}, {}^s v'_1) \wedge$

$(\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2}^f (H'_{s_2}, {}^s v'_2)$

$\implies$

${}^s v'_1 = {}^s v'_2$



*Proof.* From the CG to FG translation we know that  $\exists e_t$  s.t

$$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$$

Similarly we also know that  $\exists {}^t v_1, {}^t v_2$  s.t

$$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1 \text{ and } \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$$

From type preservation theorem we know that

$$\begin{aligned} \emptyset, \emptyset, x : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \vdash_\top e_t : (\text{unit} \xrightarrow{\perp} ((\text{unit} + \text{unit})^\perp + \text{unit})^\perp)^\perp \\ \emptyset, \emptyset, \emptyset \vdash_\top {}^t v_1 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \\ \emptyset, \emptyset, \emptyset \vdash_\top {}^t v_2 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \quad (\text{NI-1}) \end{aligned}$$

Since we have  $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 2.74 we have (we choose  $n$  s.t  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\text{Labeled } \top \text{ bool}]_E^\emptyset \quad (\text{NI-2})$$

And therefore from Definition 2.70 and (NI-2) we have

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_1)) \in [x \mapsto \text{Labeled } \top \text{ bool}]_V^\emptyset$$

From (NI-0) we know that  $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 2.74 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\emptyset \quad (\text{NI-3.1})$$

Applying Definition 2.66 on (NI-3.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^\beta \emptyset \wedge \forall i < n. e_s[{}^s v_1/x] \Downarrow_i {}^s v \implies$$

$$\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n - i, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright}^\beta \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  ${}^s v = e_s[{}^s v_1/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \hat{\triangleright}^\beta \emptyset$$

From translation and from (NI-1) we know that  ${}^t v = e_t[{}^t v_1/x] = \lambda_{\cdot}. e_{b1}$  and therefore from fg-val we have  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_1/x], \lambda_{\cdot}. e_{b1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$$

Expanding  $(\emptyset, n, e_s[{}^s v_1/x], \lambda_{\cdot}. e_{b1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$  using Definition 2.65 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \hat{\triangleright}^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_1/x]) \Downarrow_i^f (H'_{s1}, {}^s v''_1) \wedge i < k \implies$$

$$\begin{aligned} \exists H''_{t1}, {}^t v'', (H_{t3}, (\lambda_{\cdot}. e_{b1})()) \Downarrow (H''_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H''_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_1. {}^t v''_1 = \\ \text{inl } {}^t v'''_1 \wedge ({}^s \theta', k - i, {}^s v''_1, {}^t v'''_1) \in [\text{bool}]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_1, {}^s v'_1, n, \emptyset$  we get

$$\begin{aligned} \exists H''_{t1}, {}^t v'' . (\emptyset, (\lambda_{\cdot}. e_{b1})()) \Downarrow (H''_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s1}, H''_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_1. {}^t v''_1 = \\ \text{inl } {}^t v'''_1 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_1) \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2}) \end{aligned}$$

Since we have  $\exists {}^t v'''_1. {}^t v''_1 = \text{inl } {}^t v'''_1 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_1) \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$ , therefore from Definition 2.65 we know that 2 cases arise

- ${}^s v'_1 = \text{inl}^s v'_{i1}$  and  ${}^t v'''_1 = \text{inl}^t v'_{i1}$ :

And from Definition 2.65 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\text{unit}]_{\mathcal{V}}^{\hat{\beta}''}$$

which means  ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr}^s v'_{i1}$  and  ${}^t v'''_1 = \text{inr}^t v'_{i1}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v'_1 = {}^t v'''_1$  (NI-3.3)

Similarly we can apply Theorem 2.74 with the other substitution to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_E^{\emptyset} \quad (\text{NI-4.1})$$

Applying Definition 2.66 on (NI-4.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} \emptyset \wedge \forall i < n, {}^s v_s.e_s[{}^s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  ${}^s v_s = e_s[{}^s v_2/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Also from (NI-1) and from translation we know that  ${}^t v = e_t[{}^t v_2/x] = \lambda_.e_{b2}$  and therefore from fg-val we know that  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\emptyset}$$

Expanding  $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\emptyset}$  using Definition 2.65 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s2}, {}^s v''_2) \wedge i < k \implies$$

$$\exists H''_{t2}, {}^t v'', (H_{t3}, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', k - i, {}^s v''_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$  we get

$$\exists H''_{t2}, {}^t v'' . (\emptyset, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''} \quad (\text{NI-4.2})$$

Since we have  $\exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''}$ , therefore from Definition 2.65 2 cases arise

- ${}^s v'_2 = \text{inl}^s v'_{i2}$  and  ${}^t v'''_2 = \text{inl}^t v'_{i2}$ :

And from Definition 2.65 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i2}, {}^t v'_{i2}) \in [\text{unit}]_{\mathcal{V}}^{\hat{\beta}''}$$

which means  ${}^s v'_{i2} = {}^t v'_{i2} = ()$

- ${}^s v'_2 = \text{inr}^s v'_{i2}$  and  ${}^t v'''_2 = \text{inr}^t v'_{i2}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v_2' = {}^t v_2'''$  (NI-4.3)

From CG to FG translation we know that  $\exists {}^t v_{i1}. {}^t v_{i1} = \text{inl } {}^t v_{i1}$  and similarly  $\exists {}^t v_{i2}. {}^t v_{i2} = \text{inl } {}^t v_{i2}$

From (NI-1) since  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : (\text{bool}^{\perp} + \text{unit})^{\top}$  therefore from CG-inl we know that  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\perp}$

And from CGsub-sum we know that  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\top}$

Therefore we also have  $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i1} : \text{bool}^{\top}$  (NI-5.1)

Similarly we also have  $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i2} : \text{bool}^{\top}$  (NI-5.2)

Next, let  $e_T = (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top}. \text{case}(e_t(), y.y, z. {}^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^{\perp}$

where  $\text{true} = \text{inl } ()$  and  $\text{false} = \text{inr } ()$

We claim  $\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} e_T : \text{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{true} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.1:

$$\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} u : \text{bool}^{\top}}$$

P2:

$$\frac{\frac{P2.1 \quad P2.2 \quad P2.3 \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\top} \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : (\text{bool}^{\perp} + \text{unit})^{\top}}}$$

P1.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^{\perp} + \text{unit})^{\perp})^{\perp}}{\text{NI-1}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} () : \text{unit}}{\text{FG-unit}}}{\frac{\frac{}{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\perp} \searrow \perp}}{\text{FG-app}}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t() : (\text{bool}^{\perp} + \text{unit})^{\perp}}$$

P1.1:

$$\frac{\frac{P1.2 \quad \frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, y : \text{bool}^{\perp} \vdash_{\perp} y : \text{bool}^{\perp}}{\text{FG-var}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, z : \text{unit} \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-var}} \quad \frac{}{\emptyset, \emptyset \models \text{bool}^{\perp} \searrow \perp}}{\text{FG-case}}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z. {}^t v_b) : \text{bool}^{\perp}}$$

P1:

$$\frac{\frac{P1.1}{\emptyset, \emptyset, u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) : ((\text{bool}^\perp + \text{unit})^\top \xrightarrow{\perp} \text{bool}^\perp)^\perp}}$$

Main derivation:

$$\frac{\frac{P1 \quad P2 \quad \frac{}{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\emptyset, \emptyset \models \text{bool}^\perp \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^\perp} \text{FG-app}}$$

Assuming  $e_{b1}()$  reduces in  $n_{t1}$  steps in (NI-3.2) and  $e_{b2}()$  reduces in  $n_{t2}$  steps in (NI-4.2).

We instantiate Theorem 2.29 with  $e_T, ^t v_{i1}, ^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$  and  $\perp$  and therefore from (NI-3.3) and (NI-4.3) we get  $^t v_1''' = ^t v_2'''$  and thus  $^s v_1' = ^s v_2'$

□

## 2.4 Translation from FG to FG<sup>-</sup>

### 2.4.1 FG<sup>-</sup> typesystem

**Lemma 2.77** (FG<sup>-</sup>: Reflexivity of subtyping). *The following hold:*

1. For all  $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all  $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$ .

Proof of statement (1)

Let  $\tau = A^\ell$ . Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on  $A$ .

1.  $A = \mathbf{b}$ :

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2.  $A = \text{ref } \tau$ :

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}}$$

4.  $A = \tau_1 + \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6.  $A = \text{unit}$ :

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$

**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow pc}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau'} \text{FG}^- \text{-var} \quad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^{pc}} \text{FG}^- \text{-lam} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG}^- \text{-app} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^{pc}} \text{FG}^- \text{-prod} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG}^- \text{-fst} \quad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^{pc}} \text{FG}^- \text{-inl} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x. e_1, y. e_2) : \tau} \text{FG}^- \text{-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG}^- \text{-sub} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^{pc}} \text{FG}^- \text{-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG}^- \text{-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^{pc}} \text{FG}^- \text{-assign} \\
\\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^{pc}} \text{FG}^- \text{-unitI} \quad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^{pc}} \text{FG}^- \text{-FI} \\
\\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG}^- \text{-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^{pc}} \text{FG}^- \text{-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG}^- \text{-CE}
\end{array}$$

Figure 15: Type system for  $\text{FG}^-$

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A'^{\ell'}} \text{FG}^- \text{-sub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FG}^- \text{-sub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FG}^- \text{-sub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FG}^- \text{-sub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FG}^- \text{-sub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FG}^- \text{-sub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FG}^- \text{-sub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{FG}^- \text{-sub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{FG}^- \text{-sub-constraint}
\end{array}$$

Figure 16:  $\text{FG}^-$  subtyping

7.  $A = \forall \alpha. \tau_i$ :

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}$$

8.  $A = c \Rightarrow \tau_i$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash c \implies c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}$$

□

### 2.4.2 Type translation

We define a translation of types, indexed by a label  $\ell$  (which represents a  $pc$  joined with all outer labels) below. This is written  $\llbracket \tau \rrbracket_\ell$ .

**Definition 2.78** ( $FG \rightsquigarrow FG^-$ : Type translation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_\ell &= \mathbf{b} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_\ell &= \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha \\
\llbracket \tau_1 \times \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell \times \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \tau_1 + \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \text{ref } \tau \rrbracket_\ell &= \text{ref } \llbracket \tau \rrbracket_\perp \\
\llbracket \text{unit} \rrbracket_\ell &= \text{unit} \\
\llbracket \forall \gamma. (\ell_e, \tau) \rrbracket_\ell &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_\ell &= \forall \alpha. \alpha, (((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket A' \rrbracket_\ell &= (\llbracket A \rrbracket_{\ell \sqcup \ell'})^{\ell \sqcup \ell'}
\end{aligned}$$

Translation judgement:

$$\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau \rrbracket_{pc'}} \text{ where}$$

$pc' \sqsubseteq pc$  and  $\forall i \in 1 \dots n. \ell_i \sqsubseteq pc'$   
 $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$   
 $\Gamma' = x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$

### 2.4.3 Type preservation: $FG$ to $FG^-$

**Theorem 2.79** ( $FG \rightsquigarrow FG^-$ : Type preservation). *Suppose (1)  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$  and (2)  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$  in  $FG$ . Suppose  $\ell_1, \dots, \ell_n$  and  $pc'$  are arbitrary labels with free variables in  $\Sigma$  such that (3)  $\Sigma; \Psi \vdash pc' \sqsubseteq pc$  and (4) For each  $i \in [1, n]$ ,  $\Sigma; \Psi \vdash \ell_i \sqsubseteq pc'$ .*

*Let  $\Gamma'$  be the  $FG^-$  context  $x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$ . Then,  $\Sigma; \Psi; \Gamma' \vdash_{pc'} e : \llbracket \tau \rrbracket_{pc'}$  in  $FG^-$ .*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. var:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} x : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}} \text{var}$$

$$\frac{\llbracket \tau \rrbracket_{\ell_n} <: \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}}$$

2. lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma, x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{\ell'_e} e_m : \llbracket \tau_2 \rrbracket_{\ell'_e} \quad \ell_{n+1} \sqsubseteq \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : T_1} \text{lam}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{1.2} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$



P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Given}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Weakening}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma', x : \llbracket \tau_1 \rrbracket_\beta \vdash_\alpha e_m : \llbracket \tau_2 \rrbracket_\alpha} \text{IH}}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{P1}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma' \vdash_\alpha \lambda x. e_m : T_{1.3}} \text{FG}^- \text{-lam}}{\Sigma, \alpha, \beta; \Psi; \Gamma' \vdash_\alpha \nu(\lambda x. e_m) : T_{1.2}} \text{FG}^- \text{-CI}}{\Sigma, \alpha; \Psi; \Gamma' \vdash_\alpha \Lambda(\nu(\lambda x. e_m)) : T_{1.1}} \text{FG}^- \text{-FI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\Lambda(\nu(\lambda x. e_m))) : T_1} \text{FG}^- \text{-FI}$$

3. app:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau_2 \rrbracket_{pc'}} \text{app}}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = (\forall \beta. (pc' \sqcup \ell), (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge \beta \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell))$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.4} = (\llbracket \tau_1 \rrbracket_{(pc')} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

P7:

$$\overline{pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P6:

$$\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}} \text{IH2}$$

P5:

$$\overline{\Sigma; \Psi \vdash T_{1.3} \searrow pc' \sqcup \ell} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\overline{\Sigma; \Psi \vdash T_{1.2} \searrow pc' \sqcup \ell} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\overline{\Sigma; \Psi \vdash T_{1.1} \searrow pc' \sqcup \ell} \text{Definition of } \llbracket \cdot \rrbracket$$

P2:

$$\overline{pc' \sqcup pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P1:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1} \text{ IH1} \quad P2 \quad P3}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square : T_{1.1}} \text{ FG}^- \text{-FE} \quad P2 \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square \square : T_{1.2}} \text{ FG}^- \text{-FE}}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad \overline{\Sigma; \Psi \vdash c_1} \quad P2 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.3}} \text{ FG}^- \text{-CE} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.4}} \text{ FG}^- \text{-sub} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell}} \text{ FG}^- \text{-app}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 2.82}$$

4. prod:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ prod}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'}} \text{ IH1} \quad \overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ FG}^- \text{-prod}}$$

5. fst:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ fst}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-fst}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ Lemma 2.82}}$$

6. snd:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ snd}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-snd}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 2.82}}$$

7. inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{inl}$$

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{FG}^{-}\text{-inl}$$

8. inr:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{inr}$$

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inr}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inr}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{FG}^{-}\text{-inr}$$

9. case:

$$\frac{\begin{array}{l} \Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell} \\ \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell} \\ \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{\ell_{n+2}} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell} \end{array}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{case}$$

P2:

$$\frac{}{\Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH3 @ } pc' \sqcup \ell$$

P1:

$$\frac{}{\Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH2 @ } pc' \sqcup \ell$$

Main derivation:

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH1} \quad \text{P1} \quad \text{P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{FG}^{-}\text{-case} \quad \text{Lemma 2.82}$$

10. sub:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc''} e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'} \quad \Sigma; \Psi \vdash_{pc} pc \sqsubseteq pc'' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{sub}$$

$$\frac{\frac{\frac{pc' \sqsubseteq pc \sqsubseteq pc''}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'}} \text{IH} \quad \frac{\tau' <: \tau}{\llbracket \tau' \rrbracket_{pc'} <: \llbracket \tau \rrbracket_{pc'}} \text{Lemma 2.80}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{IH}$$

11. ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{ref}$$

P1:

$$\frac{\frac{\text{Given} \quad \Sigma; \Psi \vdash pc' \sqsubseteq pc}{\Sigma; \Psi \vdash \tau \searrow pc}}{\Sigma; \Psi \vdash \tau \searrow pc'} \text{Lemma 2.85}}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp \searrow pc'}$$

Main derivation:

$$\frac{\frac{\frac{\text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_\perp} \text{Lemma 2.82} \quad P1}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{FG}^- \text{-new}}$$

12. deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{deref}$$

$$\frac{\frac{\frac{\tau <: \tau'}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp <: \llbracket \tau' \rrbracket_{pc' \sqcup \ell}} \text{Lemma 2.80} \quad \frac{\text{IH1}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc'} \text{Definition of } \searrow}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc' \sqcup \ell}}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{Lemma 2.82}$$

13. assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{assign}$$

P1:

$$\frac{\frac{\text{IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp} \text{Lemma 2.82} \quad \frac{\text{Given} \quad \tau \searrow pc}{\tau \searrow pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp}$$

Main derivation:

$$\frac{\frac{\text{IH1} \quad P1 \quad \frac{\tau \searrow (\ell \sqcup pc)}{\llbracket \tau \rrbracket_\perp \searrow \ell \sqcup pc'} \text{Lemma 2.85}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{FG}^- \text{-assign}$$

14. unitI:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \mathbf{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} () : \mathbf{unit}^{pc'}} \text{unitI}$$

$$\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} () : \mathbf{unit}^{pc'}} \text{FG}^- \text{-unitI}$$

15. FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma, \alpha; \Psi; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_{\alpha})^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_{\alpha})^\alpha)^\alpha$$

$$T_{1.2} = (\forall \gamma. \alpha, \llbracket \tau \rrbracket_{\alpha})^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_{\alpha}$$

P1:

$$\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_{\alpha} e_m : T_{1.3}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_{\alpha} \Lambda e_m : T_{1.2}} \text{FG}^- \text{-FI}$$

Main derivation:

$$\frac{\frac{P1}{\Sigma, \alpha; \Psi; \Gamma' \vdash_{\alpha} \nu(\Lambda e_m) : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FG}^- \text{-FI}$$

16. CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma; \Psi, c; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \overset{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{CI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_{\alpha})^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_{\alpha})^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_{\alpha}$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi, c_1; \Gamma' \vdash_{\alpha} e_m : T_{1.2}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha; \Psi; \Gamma' \vdash_{\alpha} \nu e_m : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{FG}^- \text{-FI}$$

17. FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \gamma. (\ell_e, \tau))^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \text{FV}(\ell') \in \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\gamma] \quad \Sigma; \Psi \vdash \tau[\ell'/\gamma] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau[\ell'/\gamma] \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}} \text{FE}$$

$$T_1 = (\forall \alpha. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha) \alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}[\ell'/\gamma]$$

$$T_{1.31} = \llbracket \tau[\ell'/\gamma] \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.4} = \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}$$

P5:

$$\frac{}{T_{1.2} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\frac{}{T_{1.1} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\frac{}{(pc' \sqcup \ell) \sqsubseteq (pc \sqcup \ell) \sqsubseteq \ell_e} \text{Given}$$

P2:

$$\frac{\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{IH} \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : T_{1.1}} \text{FG}^- \text{-FE}$$

P1:

$$\frac{P2 \quad P3 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : T_{1.2}} \text{FG}^- \text{-CE}$$

P0:

$$\frac{P1 \quad \frac{\frac{}{\Sigma; \Psi \vdash T_{1.3} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket \quad P2}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : T_{1.3}} \text{FG}^- \text{-FE}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : T_{1.31}} \text{Lemma 2.84}$$

Main derivation:

$$\frac{P0 \quad \frac{}{\Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell} \text{Lemma 2.82}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet \square : T_{1.4}} \text{Lemma 2.82}$$

18. CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\rightsquigarrow} \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : \llbracket \tau \rrbracket_{pc'}} \text{CE}$$

$$T_1 = (\forall \alpha. \alpha, ((c \wedge (pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = ((c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} \llbracket \tau \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.3} = \llbracket \tau \rrbracket_{pc'}$$

$$c_1 = (c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

P3:

$$\frac{\overline{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_e} \text{ Given}}{\Sigma; \Psi \vdash (pc' \sqcup \ell) \sqsubseteq \ell_e}$$

P2:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)}$$

P1:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)}$$

P0:

$$\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{ IH} \quad P1}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.1}} \text{ FG}^- \text{-FE} \quad \frac{\overline{\Sigma; \Psi \vdash c} \text{ Given, Weakening} \quad P3}{\Sigma; \Psi \vdash c_1} \text{ P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.2}} \text{ FG}^- \text{-CE}$$

Main derivation:

$$\frac{P0.1 \quad \frac{\overline{\tau \searrow \ell} \text{ Given}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}} \text{ Lemma 2.82}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}}$$

□

**Lemma 2.80** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping).  $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$  and the following holds:

1.  $\forall \tau, \tau'.$

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_\ell <: \llbracket \tau' \rrbracket_{\ell'}$$

2.  $\forall A, A'.$

$$\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \llbracket A \rrbracket_\ell <: \llbracket A' \rrbracket_{\ell'}$$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash A_1 <: A_2} \text{ By inversion} \quad P1}{\Sigma; \Psi \vdash ([A_1]_{\ell \sqcup \ell_1}) <: ([A_2]_{\ell' \sqcup \ell_2})} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion} \quad \frac{\overline{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{ Given}}{\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2}}{\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash ([A_1]_{\ell \sqcup \ell_1})^{\ell \sqcup \ell_1} <: ([A_2]_{\ell' \sqcup \ell_2})^{\ell' \sqcup \ell_2}}{\Sigma; \Psi \vdash [A_1^{\ell_1}]_{\ell} <: [A_2^{\ell_2}]_{\ell'}}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{ FG}^- \text{ sub-base}}{\Sigma; \Psi \vdash [\mathbf{b}]_{\ell} <: [\mathbf{b}]_{\ell'}} \text{ Definition of } [\cdot]$$

2. FGsub-ref:

$$\frac{\overline{\Sigma; \Psi \vdash \text{ref } [\tau_i]_{\perp} <: \text{ref } [\tau_i]_{\perp}} \text{ FG}^- \text{ sub-ref}}{\Sigma; \Psi \vdash [\text{ref } \tau_i]_{\ell} <: [\text{ref } \tau_i]_{\ell'}} \text{ Definition of } [\cdot]$$

3. FGsub-prod:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash [\tau_1]_{\ell} <: [\tau'_1]_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash [\tau_2]_{\ell} <: [\tau'_2]_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash [\tau_1]_{\ell} \times [\tau_2]_{\ell} <: [\tau'_1]_{\ell} \times [\tau'_2]_{\ell'}}{\Sigma; \Psi \vdash [\tau_1 \times \tau_2]_{\ell} <: [\tau'_1 \times \tau'_2]_{\ell'}} \text{ FG}^- \text{ sub-prod}}{\Sigma; \Psi \vdash [\tau_1 \times \tau_2]_{\ell} <: [\tau'_1 \times \tau'_2]_{\ell'}} \text{ Definition of } [\cdot]$$



4. FGsub-sum:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'} + \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ FG}^- \text{ sub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

5. FGsub-arrow:

$$T_1 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)$$

$$T_{1.0} = \forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)$$

$$T_{1.1} = ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)$$

$$T_{1.2} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$

$$T_2 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha)$$

$$T_{2.0} = \forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)$$

$$T_{2.1} = ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)$$

$$T_{2.2} = (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha$$

$$c_2 = (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha)$$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\alpha <: \llbracket \tau'_2 \rrbracket_\alpha} \text{ IH(1) with } \ell = \ell' = \alpha$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau'_1 <: \tau_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{ IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.3} <: T_{2.3}} \text{ FG}^- \text{ sub-arrow}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell'_e \sqsubseteq \ell_e}}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e}}{\Sigma, \alpha, \beta; \Psi \vdash c_2 \implies c_1} \quad \text{P1}}{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.2} <: T_{2.2}}}{\text{Weakening, FG}^- \text{-sub-label}}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}}} \text{FG}^- \text{-sub-constraint}$$

P0.1:

$$\frac{\text{P0}}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \text{FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{\frac{\frac{\text{P0.1}}{\Sigma; \Psi \vdash T_1 <: T_2}}{\text{FG}^- \text{-sub-label}}}{\Sigma; \Psi \vdash \left[ \left[ \tau_1 \xrightarrow{\ell_\xi} \tau_2 \right]_\ell <: \left[ \left[ \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2 \right]_{\ell'} \right]_{\ell'}} \text{Definition of } \llbracket \cdot \rrbracket$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_\ell <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{FG}^- \text{-sub-unit} \text{ Definition of } \llbracket \cdot \rrbracket$$

7. FGsub-forall:

$$\begin{aligned}
T_1 &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
T_{1.0} &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
T_{1.1} &= (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
c_1 &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \\
T_{1.2} &= \llbracket \tau \rrbracket_\alpha \\
T_2 &= \forall \alpha. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha)^\alpha \\
T_{2.0} &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
T_{2.1} &= (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
c_2 &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \\
T_{2.2} &= \llbracket \tau' \rrbracket_\alpha
\end{aligned}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e}}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

P1:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}}{\text{IH}} \quad \text{FG}^- \text{-sub-forall} \quad \frac{}{\Sigma; \Psi \vdash c_2 \implies c_1} \text{P0}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \quad \text{FG}^- \text{-sub-constraint}}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \quad \text{FG}^- \text{-sub-forall}}{\Sigma; \Psi \vdash T_1 <: T_2} \quad \text{FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{}{\Sigma; \Psi \vdash \llbracket \forall \gamma. \tau_1 \rrbracket_\ell <: \llbracket \forall \gamma. \tau_2 \rrbracket_{\ell'}} \text{P0.1} \quad \text{Definition of } \llbracket \cdot \rrbracket$$

8. FGsub-constraint:

$$T_1 = \forall \alpha. \alpha, ((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.1} = ((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_\alpha$$

$$c_1 = (c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_2 = \forall \alpha. \alpha, ((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau' \rrbracket_\alpha)^\alpha$$

$$T_{2.1} = ((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau' \rrbracket_\alpha)^\alpha$$

$$T_{2.2} = \llbracket \tau' \rrbracket_\alpha$$

$$c_2 = (c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e)$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau_2$$

P1:

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash c' \implies c} \text{By inversion}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e} \text{P1}}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket_\alpha <: \llbracket \tau' \rrbracket_\alpha} \text{IH}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \text{FG}^- \text{-sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{Definition of } \llbracket \cdot \rrbracket_\ell$$

□

**Lemma 2.81** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping with label). *If  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell'$ , then  $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau \rrbracket_{\ell'}$  in FG<sup>-</sup>.*

*Proof.* From Lemma 2.80 with  $\tau = \tau'$  and from Lemma 2.77 □

**Lemma 2.82** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping for  $\tau \searrow \ell$ ). *If  $\Sigma; \Psi \vdash \tau \searrow \ell$ , then  $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'}$  in FG<sup>-</sup>.*

*Proof.* Since  $\Sigma; \Psi \vdash \tau \searrow \ell$ , there exists  $\ell''$  such that  $\tau = \mathbf{A}^{\ell''}$  and  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$ . Now we have:

$$\begin{aligned} & \Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'} \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell \sqcup \ell'} <: \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell'} && (\tau = \mathbf{A}^{\ell''}) \\ = & \Sigma; \Psi \vdash (\llbracket \mathbf{A} \rrbracket_{\ell \sqcup \ell' \sqcup \ell''})^{\ell \sqcup \ell' \sqcup \ell''} <: (\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell''})^{\ell' \sqcup \ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell \sqcup \ell''} <: \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \end{aligned}$$

The last statement holds by Lemma 2.81, since  $\Sigma; \Psi \vdash \ell \sqcup \ell'' \sqsubseteq \ell''$  follows from our earlier assertion that  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$ . □

**Lemma 2.83** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Lemma for protection relation).  $\forall \Sigma, \Psi, \alpha, \tau, \ell, \ell'$ .

$$\Sigma, \alpha; \Psi \vdash \tau \searrow \ell \implies \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell[\ell'/\alpha], \text{ where } FV(\ell') \in \Sigma$$

*Proof.* Say  $\tau = \mathbf{A}^{\ell_g}$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell_g} \text{By inversion on } \Sigma, \alpha; \Psi \vdash \tau \searrow \ell}}{\Sigma; \Psi \vdash \ell[\ell'/\alpha] \sqsubseteq \ell_g[\ell'/\alpha]} \text{Substitution over constraints}}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_g}[\ell'/\alpha] \searrow \ell[\ell'/\alpha]} \text{Definition of } \searrow$$

□

**Lemma 2.84** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Substitution lemma). *For all  $\ell, \ell'$  the following hold:*

1.  $\forall \tau. \llbracket \tau \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \tau[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$
2.  $\forall \mathbf{A}. \llbracket \mathbf{A} \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $\mathbf{A}$

Proof of statement (1)

$$\begin{aligned} & \text{Let } \tau = \mathbf{A}^{\ell_i} \\ & \llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell}[\ell'/\alpha] \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell})^{\ell_i \sqcup \ell}[\ell'/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell}[\ell'/\alpha])^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} \\ = & (\llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]})^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} && \text{IH(2) on } \mathbf{A} \\ = & \llbracket (\mathbf{A}[\ell'/\alpha])^{\ell_i[\ell'/\alpha]} \rrbracket_{\ell[\ell'/\alpha]} \\ = & \llbracket \mathbf{A}^{\ell_i}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \end{aligned}$$

Proof of statement (2)

We consider cases of  $\mathbf{A}$

1.  $\mathbf{A} = \mathbf{b}$ :

$$\begin{aligned} & \llbracket \mathbf{b} \rrbracket_{\ell}[\ell'/\alpha] \\ = & \mathbf{b}[\ell'/\alpha] && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \mathbf{b} && \alpha \notin FV(\mathbf{b}) \\ = & \llbracket \mathbf{b} \rrbracket_{\ell} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \llbracket \mathbf{b}[\ell'/\alpha] \rrbracket_{\ell} \end{aligned}$$

2.  $A = \text{ref } \tau_i$ :

$$\begin{aligned}
& \llbracket \text{ref } \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \text{ref } \llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \text{ref } (\llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha]) \\
&= \text{ref } \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\perp} && \text{IH(1) on } \tau_i \\
&= \llbracket \text{ref } \tau_i[\ell'/\alpha] \rrbracket_{\ell}
\end{aligned}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] \times \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \times \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] \times \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 \times \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

4.  $A = \tau_1 + \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 + \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} + \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] + \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} + \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] + \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 + \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell \sqsubseteq \beta_1 \sqsubseteq \ell_e \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta}[\ell'/\alpha] \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1}[\ell'/\alpha])^{\beta_1})^{\beta_1})^{\beta_1} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1[\ell'/\alpha] \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1} \\
& \quad \text{(IH1 on } \tau_1 \text{ and } \tau_2 \text{)} \\
&= \llbracket (\tau_1[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} \\
&= \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

6.  $A = \forall \gamma. \tau_i$ :

$$\begin{aligned}
& \llbracket \forall \gamma. \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta. \beta, ((\ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \quad \text{IH1 on } \tau_i \\
&= \llbracket \forall \beta. \ell_e[\ell'/\alpha], \tau_i[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

7.  $A = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& \llbracket c \Rightarrow \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
= & \forall \beta. \beta, (((c \wedge \ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \text{IH1 on } \tau_i \\
= & \left[ \left[ (c[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_i[\ell'/\alpha]) \right] \right]_{\ell[\ell'/\alpha]} \\
= & \left[ (c \xrightarrow{\ell_e} \tau_i)[\ell'/\alpha] \right]_{\ell[\ell'/\alpha]}
\end{aligned}$$

□

**Lemma 2.85** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Preservation of protection relation).  $\forall \tau, \ell, \ell'$ .

$$\tau \searrow \ell \implies \llbracket \tau \rrbracket_{\ell'} \searrow \ell$$

*Proof.* Let  $\tau = \mathbf{A}^{\ell_i}$

$$\begin{array}{c}
\frac{}{\tau \searrow \ell} \text{ Given} \\
\frac{}{\ell \sqsubseteq \ell_i} \text{ Given} \\
\frac{}{\ell \sqsubseteq (\ell' \sqcup \ell_i)} \text{ By inversion} \\
\frac{}{(\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell_i})^{\ell' \sqcup \ell_i} \searrow \ell} \text{ Definition of } \llbracket \cdot \rrbracket \\
\frac{}{\llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell'} \searrow \ell} \\
\hline
\llbracket \tau \rrbracket_{\ell'} \searrow \ell
\end{array}$$

□

## 2.5 FG to CG translation

### 2.5.1 Type directed (direct) translation from FG to CG

**Definition 2.86** (FG  $\rightsquigarrow$  CG: Type translation).

$$\begin{aligned}
\langle \mathbf{b} \rangle_\ell &= \mathbf{b} \\
\langle \mathbf{unit} \rangle_\ell &= \mathbf{unit} \\
\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle_\ell &= \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha \\
\langle \tau_1 \times \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell \\
\langle \tau_1 + \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell \\
\langle \mathbf{ref} \ A^{\ell'} \rangle_\ell &= \mathbf{ref} \ \ell' \ \langle A \rangle_{\ell'} \\
\langle \forall \alpha. (\ell_e, \tau) \rangle_\ell &= \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} \\
\langle c \xrightarrow{\ell_e} \tau \rangle_\ell &= \forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_\alpha \\
\langle A^{\ell'} \rangle_\ell &= \mathbf{Labeled} \ (\ell \sqcup \ell') \ \langle A \rangle_{\ell \sqcup \ell'}
\end{aligned}$$

For  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$  and  $\bar{\ell} = \ell_1, \dots, \ell_n$ , define  $\langle \Gamma \rangle_{\bar{\ell}} = x_1 : \langle \tau_1 \rangle_{\ell_1}, \dots, x_n : \langle \tau_n \rangle_{\ell_n}$ .

We use a coercion function defined as follows:

$$\begin{aligned}
\mathbf{coerce\_taint} &: \mathbb{C} \gamma \alpha_c \tau' \rightarrow \mathbb{C} \gamma \gamma \tau' \quad \text{when } \tau' = \mathbf{Labeled} \ \alpha'_c \ \tau \text{ and } \Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \\
\mathbf{coerce\_taint} &\triangleq \lambda x. \mathbf{toLabeled}(\mathbf{bind}(x, y. \mathbf{unlabel}(y)))
\end{aligned}$$

$$\begin{aligned}
&\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \mathbf{ret} \ x} \text{FC-var} \\
&\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \mathbf{ret}(\mathbf{Lb}(\mathbf{\Lambda}\mathbf{\Lambda}\mathbf{\Lambda}(\nu(\lambda x. e_{c1}))))} \text{FC-lam} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 \ e_2 : \tau_2 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{bind}(\mathbf{unlabel} \ a, c. (c \bullet b)))))} \text{FC-app} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{ret}(\mathbf{Lb}(a, b))))} \text{FC-prod} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{fst}(e) : \tau_1 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{fst}(b)))))} \text{FC-fst} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{snd}(e) : \tau_2 \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{snd}(b)))))} \text{FC-snd} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinl}(a)))} \text{FC-inl} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}
\end{aligned}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)))} \text{FC-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()))} \text{FC-assign} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda (\nu(e_c))))} \text{FC-FI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda (\nu(e_c))))} \text{FC-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-CE}
\end{array}$$

## 2.5.2 Type preservation for FG to CG translation

**Lemma 2.87** (Coercion lemma - typing).  $\forall \Sigma, \Psi, \Gamma, \alpha_c, \alpha'_c, \tau.$

$$\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \implies$$

$$\Sigma; \Psi; \Gamma \vdash \text{coerce\_taint} : \mathbb{C} \gamma \alpha_c \text{ Labeled } \alpha'_c \tau \rightarrow \mathbb{C} \gamma \gamma \text{ Labeled } \alpha'_c \tau$$

*Proof.*  $T_{c4} = \text{Labeled } \alpha'_c \tau$

$$T_{c3} = \mathbb{C} \alpha_c \alpha'_c \tau$$

$$T_{c2} = \mathbb{C} \gamma \alpha'_c \tau$$

$$T_{c1} = \mathbb{C} \gamma \gamma \text{ Labeled } \alpha'_c \tau$$

$$T_{c0} = \mathbb{C} \gamma \alpha_c \text{ Labeled } \alpha'_c \tau$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}}{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel} \quad \frac{}{\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c} \text{Given}}{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-var}$$



Pc1:

$$\frac{}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\Sigma; \Psi; \Gamma \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\Sigma; \Psi; \Gamma \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

□

**Theorem 2.88** (FG  $\rightsquigarrow$  CG: Type preservation). *Suppose  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$  in FG. Then, there exists  $e'$  such that  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$  and for any  $\alpha', \beta', \gamma'$  with  $\beta' \sqcup \gamma' \sqsubseteq pc \sqcap \alpha'$ , there is a derivation of  $\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e' : \mathbb{C} \gamma' \gamma' (\tau)_{\alpha'}$  in CG.*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{\frac{\frac{}{(\Gamma)_{\beta'_o}(x) = (\tau)_{\beta'_o}} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash x : (\tau)_{\beta'_o}} \text{CG-var} \quad \frac{\frac{}{\Sigma; \Psi \vdash \beta'_o \sqcup \gamma'_o \sqsubseteq \alpha'_o \sqcap pc} \text{Given}}{\Sigma; \Psi \vdash \beta'_o \sqsubseteq \alpha'_o} \text{Lemma 2.89, CG-sub}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash x : (\tau)_{\alpha'_o}} \text{CG-ret}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash \text{ret } x : \mathbb{C} \gamma'_o \gamma'_o (\tau)_{\alpha'_o}} \text{CG-ret}$$

2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1}))))} \text{FC-lam}$$

$$\begin{aligned} T_0 &= \mathbb{C} \gamma'_j \gamma'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp)_{\alpha'_j} = \mathbb{C} \gamma'_j \gamma'_j \text{Labeled } \alpha'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2))_{\alpha'_j} \\ T_1 &= \mathbb{C} \gamma'_j \gamma'_j \text{Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.0} &= \text{Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.1} &= \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.2} &= (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.3} &= (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.4} &= \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \end{aligned}$$

P3:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \gamma_j \sqsubseteq \alpha'_j \sqcap pc \quad \text{Given, Weakening}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqsubseteq \alpha'_j}$$

P2:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e \quad P3}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e}$$

P1:

$$\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}}, x : \langle \tau_1 \rangle_{\beta_t} \vdash e_{c1} : T_{1.4} \quad \text{IH}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x. e_{c1} : T_{1.3}} \quad P2}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x. e_{c1} : T_{1.3}} \quad \text{CG-lam}$$

P0:

$$\frac{\Sigma; \Psi \vdash \overline{\beta'_j} \sqcup \gamma'_j \sqsubseteq \alpha'_j \quad \text{Given}}{\Sigma; \Psi \vdash \gamma_j \sqsubseteq \alpha_j}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \nu(\lambda x. e_{c1}) : T_{1.2} \quad P1}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1})) : T_{1.1}} \quad \text{CG-CI}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1}))) : T_{1.0}} \quad \text{3 applications CG-FI} \quad P0}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1})))) : T_1} \quad \text{CG-label} \quad \text{CG-ret}$$

3. FC-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c \square \square \square \bullet) b))))} \quad \text{FC-app}$$

$$\beta' = \bigcup_{\beta_i \in \overline{\beta'}} \beta_i$$

$$T_0 = \mathbb{C} \gamma' \gamma' \langle (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rangle_{\beta' \sqcup \gamma'} = \mathbb{C} \gamma' \gamma' \text{ Labeled } (\beta' \sqcup \gamma' \sqcup \ell) \langle (\tau_1 \xrightarrow{\ell_e} \tau_2) \rangle_{\beta' \sqcup \gamma' \sqcup \ell}$$

$$T_1 = \mathbb{C} \gamma' \gamma' \text{ Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha$$

$$T_{1.1} = \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha$$

$$T_{1.2} = \mathbb{C} \gamma' (\gamma' \sqcup (\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha$$

$$T_{1.3} = \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha$$

$$T_{1.4} = \forall \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.5} = \forall \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.6} = (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow T_{1.7}$$

$$T_{1.7} = (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.8} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.10} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i})_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.11} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.12} = \mathbb{C} (\gamma') (\gamma') \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.13} = \mathbb{C} (\gamma') (\gamma') \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma') (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma')}$$

$$T_2 = \mathbb{C} (\gamma') (\gamma') (\tau_2)_{(\beta' \sqcup \gamma')}$$

P8:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash b : (\tau_2)_{(\beta' \sqcup \gamma')}} \text{CG-var}$$

P7:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \alpha' \sqcap pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha' \sqcap pc \sqsubseteq \ell_e}$$

P6:

$$\frac{P7 \quad \frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash (\ell \sqcup \beta' \sqcup \gamma') \sqsubseteq \ell_e}$$

P5:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c : T_{1.3}} \text{CG-var}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[] : T_{1.4}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][] : T_{1.5}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][][] : T_{1.6}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][][] \bullet : T_{1.7}} \text{CG-CE} \quad P6$$

P4:

$$\frac{\frac{P5 \quad P8}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash (c[][][] \bullet) b : T_{1.8}} \text{CG-app}}$$

P3:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash a : T_{1.1}} \text{CG-var}$$

P2:

$$\frac{\frac{P3}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash \text{unlabel } a : T_{1.2}} \text{CG-unlabel} \quad P4}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash \text{bind}(\text{unlabel } a, c.(c[][][] \bullet) b) : T_{1.9}} \text{CG-bind}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash e_{c2} : T_2} \text{IH2, Weakening} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.9}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, e_{c1} : T_1} \text{IH1 with } (\beta' \sqcup \gamma'), \bar{\beta}', \gamma' \quad P1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, e_{c1} : T_1} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.10}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.11}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_{1.12}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_{1.13}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_2}$$

4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 \times \tau_2)^\perp)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 \times \tau_2))_{\alpha'}$$

$$T_3 = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_{3.1} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_4 = \mathbb{C} \gamma' \gamma' (\tau_1)_{\alpha'}$$

$$T_5 = \mathbb{C} \gamma' \gamma' (\tau_2)_{\alpha'}$$

P4:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{CG-var}$$

P3:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash b : (\tau_2)_{\alpha'}} \text{CG-var}$$

P2:

$$\frac{\frac{\frac{P3 \quad P4}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash (a, b) : (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}} \text{CG-prod}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{Lb}(a, b) : T_{3.1}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'} \vdash e_{c2} : T_5} \text{IH2} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_{c1} : T_4}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1} \text{CG-bind}}{\text{Definition 2.86}}$$

5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{FC-fst}$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau_1 \rangle_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (\tau_1 \times \tau_2)^\ell \rangle_{\alpha'}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup \alpha' \langle (\tau_1 \times \tau_2) \rangle_{\alpha' \sqcup \ell}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup \alpha' \langle \tau_1 \rangle_{\alpha' \sqcup \ell} \times \langle \tau_2 \rangle_{\alpha' \sqcup \ell}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup \alpha' \langle \tau_1 \rangle_{\alpha' \sqcup \ell} \times \langle \tau_2 \rangle_{\alpha' \sqcup \ell}$$

$$T_{2.4} = \langle \tau_1 \rangle_{\alpha' \sqcup \ell} \times \langle \tau_2 \rangle_{\alpha' \sqcup \ell}$$

$$T_{2.5} = \mathbb{C} (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) \langle \tau_1 \rangle_{\alpha' \sqcup \ell} \times \langle \tau_2 \rangle_{\alpha' \sqcup \ell}$$

$$T_3 = \mathbb{C} (\gamma' \sqcup \alpha' \sqcup \ell) (\gamma' \sqcup \alpha' \sqcup \ell) \langle \tau_1 \rangle_{\alpha' \sqcup \ell}$$

$$T_{3.1} = \mathbb{C} (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) \langle \tau_1 \rangle_{\alpha' \sqcup \ell}$$

$$T_{3.2} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) \langle \tau_1 \rangle_{\alpha' \sqcup \ell}$$

$$T_{3.3} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) \langle \mathbf{A}^{\ell_i} \rangle_{\alpha' \sqcup \ell}$$

$$T_{3.4} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \langle \mathbf{A} \rangle_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.5} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \langle \mathbf{A} \rangle_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.6} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup \alpha' \langle \mathbf{A} \rangle_{\alpha' \sqcup \ell_i}$$

$$T_{3.7} = \mathbb{C} (\gamma') (\gamma') \langle \mathbf{A}^{\ell_i} \rangle_{\alpha'}$$

P2:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}}{\text{CG-var}} \quad \text{CG-fst}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : \langle \tau_1 \rangle_{\alpha' \sqcup \ell}} \quad \text{CG-ret}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1}} \text{CG-bind}$$

P0:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e_c : T_{2.2}} \text{IH} \quad P1 \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1} \quad \text{CG-bind} \\
\frac{}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.4} \quad \text{Definition 2.86} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.5} \quad \text{Lemma 2.87}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P0 \quad \frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{By inversion}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.6} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.7} \quad \text{Definition 2.86} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_1
\end{array}$$

6. FC-snd:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd}$$

$$\begin{aligned}
T_1 &= \mathbb{C} \ \gamma' \ \gamma' \ (\tau_2)_{\alpha'} \\
T_2 &= \mathbb{C} \ \gamma' \ \gamma' \ ((\tau_1 \times \tau_2)^\ell)_{\alpha'} \\
T_{2.1} &= \mathbb{C} \ \gamma' \ \gamma' \ \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1 \times \tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.2} &= \mathbb{C} \ \gamma' \ \gamma' \ \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.3} &= \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.4} &= ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.5} &= \mathbb{C} \ (\gamma') \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_3 &= \mathbb{C} \ (\gamma' \sqcup \alpha' \sqcup \ell) \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.1} &= \mathbb{C} \ (\gamma') \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.2} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.3} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ (\mathbf{A}^{\ell_i})_{\alpha' \sqcup \ell} \\
T_{3.4} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i} \\
T_{3.5} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i} \\
T_{3.6} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell_i} \\
T_{3.7} &= \mathbb{C} \ (\gamma') \ (\gamma') \ (\mathbf{A}^{\ell_i})_{\alpha'}
\end{aligned}$$

P2:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : (\tau_2)_{\alpha' \sqcup \ell}} \text{CG-snd}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))) : T_{3.1}} \text{CG-bind}$$

P0:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.1}} \text{CG-bind}}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.2}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.4}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.5}}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell}{\text{By inversion}}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.6}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.7}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_1}}$$

7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_{1.3} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau_1)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{inl}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{CG-inl}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{Lbinl}(a) : T_{1.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{CG-ret}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_{1.2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_1} \text{Definition 2.86}}$$

8. FC-inr:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_{1.3} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau_2)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{CG-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{inr}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{CG-inr}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{Lbinr}(a) : T_{1.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{ret}(\text{Lbinr}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_{1.2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_1} \text{Definition 2.86}}$$

9. FC-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case}}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}' } \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' (\tau)_{(\alpha')}$$

$$T_2 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) (\tau_1 + \tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$





10. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow_{pc}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new}(e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb})))} \text{FC-ref}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\text{ref } \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' ((\text{ref } A^{\ell_i})^\perp)_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\text{ref } A^{\ell_i}))_{\alpha'}$$

$$T_{1.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i (A)_{\ell_i}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' (A^{\ell_i})_{(\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.4} = \mathbb{C} \gamma' \gamma' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.5} = \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.51} = \text{Labeled } \alpha' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.6} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i (A)_{\ell_i}$$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash A^{\ell_i} \searrow_{pc}}{\text{Given}} \quad \text{By inversion} \quad \frac{}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq pc} \text{Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_i}}{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq \ell_i}$$

P2:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}} \quad \text{CG-label} \quad \frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}}{\text{CG-ret}} \quad P3}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{2.6}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{new}(a) : T_{2.4}}{\text{CG-new}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{CG-bind}}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb}))) : T_{1.3}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb}))) : T_1} \text{Definition 2.86}$$

11. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau' \rangle_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \langle \mathbf{A}^{\ell'_i} \rangle_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell'_i \sqcup \alpha' \langle \mathbf{A}' \rangle_{\ell'_i \sqcup \alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (\text{ref } \tau)^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) \langle (\text{ref } \tau) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) \langle (\text{ref } \mathbf{A}^{\ell_i}) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.4} = \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.5} = \mathbb{C} \gamma' \beta' \sqcup \gamma' \sqcup \ell (\text{ref } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.6} = (\text{ref } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.7} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.8} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i \langle \mathbf{A} \rangle_{\ell_i})$$

$$T_{2.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell'_i \langle \mathbf{A}' \rangle_{\ell'_i})$$

$$T_{2.10} = \mathbb{C} (\gamma') (\gamma') (\text{Labeled } \beta' \sqcup \gamma' \sqcup \ell \sqcup \ell'_i \langle \mathbf{A}' \rangle_{\ell'_i})$$

$$T_{2.11} = \mathbb{C} (\gamma') (\gamma') (\text{Labeled } \alpha \sqcup \ell'_i \langle \mathbf{A}' \rangle_{\ell'_i})$$

P2:

$$\frac{\frac{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}}{\text{CG-var}}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{CG-deref}$$

P1:

$$\frac{\frac{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash e_c : T_{2.3}}{P1}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.8}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}}{P0}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.10}} \text{Lemma 2.87}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{2.10}} \text{Lemma 2.87}}{\frac{\frac{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell}{\text{Given}}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion} \quad \frac{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha'}{\text{Given}}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{CG-sub}$$

12. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow_{\Delta} (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \text{unit} \rangle_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{unit}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (\text{ref } \tau)^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (\text{ref } \tau) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (\text{ref } A^{\ell_i}) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i \langle A \rangle_{\ell_i}$$

$$T_{2.4} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i \langle A \rangle_{\ell_i}$$

$$T_{2.5} = \mathbb{C} \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i \langle A \rangle_{\ell_i}$$

$$T_{2.6} = \text{ref } \ell_i \langle A \rangle_{\ell_i}$$

$$T_{2.7} = \mathbb{C} \ell \sqcup (\beta' \sqcup \gamma') \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.8} = \mathbb{C} \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.9} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_3 = \mathbb{C} \gamma' \gamma' \langle \tau \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{3.1} = \mathbb{C} \gamma' \gamma' \langle A^{\ell_i} \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{3.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') \langle A \rangle_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') \langle A \rangle_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.4} = \text{Labeled } \ell_i \langle A \rangle_{\ell_i}$$

P4:

$$\frac{}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{CG-var}$$

P5:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}}{\text{CG-var}}}{\Sigma; \Psi \vdash \tau = A^{\ell_i} \searrow_{\Delta} (pc \sqcup \ell)} \text{Given}}{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq pc} \text{Given}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.4}} \text{CG-var}$$

P3:

$$\frac{P4 \quad P5}{\Sigma; \Psi; \langle \Gamma \rangle_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}} \text{CG-assign}$$

P2:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}}{\text{CG-unlabel}} \quad P3}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{CG-bind}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash e_{c2} : T_{3.2}}{\text{IH2}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_{c1} : T_{2.3}}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.8}} \text{CG-bind}$$

P0.1:

$$\frac{P0}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}} \text{CG-toLabeled}$$

Main derivation:

$$\frac{P0.1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()) : T_{1.1}} \text{CG-bind}$$

13. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c))))} \text{FC-FI}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\forall \alpha. (\ell_e, \tau))^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\forall \alpha. (\ell_e, \tau))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \forall \alpha. \forall \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.1} = (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.2} = \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \alpha' \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi, (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2}{\text{IH, Weakening}}}{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \nu(e_c) : T_{2.1}} \text{CG-CI}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \Lambda \Lambda \Lambda(\nu(e_c)) : T_{2.2}} \text{CG-FI}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c))) : T_{2.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c)))) : T_{1.2}} \text{CG-ret}$$

14. FC-FE:

$$\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.\text{b}[]\bullet)))} \text{FC-FE}$$

$$\beta' = \bigcup_{\beta_i \in \overline{\beta'}} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau[\ell''/\alpha] \rangle_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (\forall \alpha. (\ell_e, \tau))^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (\forall \alpha. (\ell_e, \tau)) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.4} = \mathbb{C} \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.5} = \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.6} = \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}[\ell''/\alpha]$$

$$T_{2.7} = \forall \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.8} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.81} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.9} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.10} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.11} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau[\ell''/\alpha] \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.12} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \mathbf{A}^{\ell_i}[\ell''/\alpha] \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.13} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled } \ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.14} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled } \ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma' \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.15} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma' \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.16} = \mathbb{C} (\gamma') (\gamma') \langle \mathbf{A}[\ell''/\alpha]^{\ell_i[\ell''/\alpha]} \rangle_{(\beta' \sqcup \gamma')}$$

P3:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e[\ell''/\alpha]}{\text{Given}} \quad \frac{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq pc \sqsubseteq \ell_e[\ell''/\alpha]}{\text{Given}}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq \ell_e[\ell''/\alpha])}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha])}$$

P2:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[] : T_{2.6}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][] : T_{2.7}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.81}} \text{CG-FE} \quad P3}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.9}} \text{CG-CE}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.10}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.10}} \text{CG-bind}$$

P0.1:

$$\frac{\frac{\Sigma; \Psi \vdash A[\ell''/\alpha]^{l_i[\ell''/\alpha]} \searrow \ell}{\text{Given}}}{\Sigma; \Psi \vdash \ell \sqsubseteq l_i[\ell''/\alpha]} \text{By inversion}$$

P0.2:

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.11}}{\text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.12}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.13}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.14}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.15}} \text{Lemma 2.87}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.11}} \text{P0}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.12}} \text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.13}} \text{P0.1}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.14}} \text{Lemma 2.87}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.15}} \text{Lemma 2.87}}$$

Main derivation:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_1}{\text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_1} \text{P0.2}}$$

15. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{FC-CI}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((c \xrightarrow{\ell_e} \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((c \xrightarrow{\ell_e} \tau))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.3} = \text{Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.4} = \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.5} = (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\Sigma, \alpha_i, \gamma_i; \Psi, (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma) \vdash e_c : T_2}{\text{IH, Weakening}}}{\Sigma; \Psi; \Gamma \vdash \nu(e_c) : T_{1.5}} \text{CG-CI}}{\Sigma; \Psi; \Gamma \vdash \Lambda\Lambda(\nu(e_c)) : T_{1.4}} \text{CG-FI}}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(\Lambda\Lambda(\nu(e_c))) : T_{1.3}} \text{CG-label}}{\Sigma; \Psi; \Gamma \vdash \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) : T_{1.2}} \text{CG-ret}}$$

16. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.\text{b}[\square]\square\bullet)))} \text{FC-CE}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau \rangle_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (c \xrightarrow{\ell_e} \tau)^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (c \xrightarrow{\ell_e} \tau) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.4} = \mathbb{C} \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.5} = \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.6} = \forall \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.71} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.8} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.10} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \mathbf{A}^{\ell_i} \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.11} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.12} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.13} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.14} = \mathbb{C} (\gamma') (\gamma') \langle \mathbf{A}^{\ell_i} \rangle_{(\beta' \sqcup \gamma')}$$

P2:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[\square] : T_{2.6}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[\square]\square : T_{2.71}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[\square]\square\bullet : T_{2.8}} \text{CG-CE}} \text{CG-var}$$

P1:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.\text{b}[\square]\square\bullet) : T_{2.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}} \text{CG-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.\text{b}[\square]\square\bullet) : T_{2.9}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.\text{b}[\square]\square\bullet)) : T_{2.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}} \text{IH} \quad P1$$



Main derivation:

$$\begin{array}{c}
P0 \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])) : T_{2.10} \quad \text{CG-bind} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])) : T_{2.11} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.12} \quad \text{Lemma 2.87} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.13} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.14} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_1
\end{array}$$

□

**Lemma 2.89** (FG  $\rightsquigarrow$  CG: Subtyping preservation).  $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$  and the following holds:

1.  $\forall \tau, \tau'.$   
 $\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau' \rrbracket_{\ell'}$
2.  $\forall A, A'.$   
 $\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \llbracket A \rrbracket_{\ell} <: \llbracket A' \rrbracket_{\ell'}$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\begin{array}{c}
\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash A_1 <: A_2 \quad \text{By inversion} \quad P1 \\
\hline
\Sigma; \Psi \vdash (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2}) \quad \text{IH(2) on } A_1 <: A_2
\end{array}$$

P1:

$$\begin{array}{c}
\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2 \quad \text{By inversion} \quad \overline{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2
\end{array}$$

Main derivation:

$$\begin{array}{c}
P1 \quad P2 \\
\hline
\Sigma; \Psi \vdash \text{Labeled } \ell \sqcup \ell_1 (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: \text{Labeled } \ell' \sqcup \ell_2 (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2}) \quad \text{CGsub-labeled} \\
\hline
\Sigma; \Psi \vdash \llbracket A_1^{\ell_1} \rrbracket_{\ell} <: \llbracket A_2^{\ell_2} \rrbracket_{\ell'}
\end{array}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\begin{array}{c}
\overline{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \quad \text{CG-refl} \\
\hline
\Sigma; \Psi \vdash \llbracket \mathbf{b} \rrbracket_{\ell} <: \llbracket \mathbf{b} \rrbracket_{\ell'} \quad \text{Definition 2.86}
\end{array}$$

2. FGsub-ref:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{ref } \ell_i \llbracket A \rrbracket_{\ell_i} <: \text{ref } \ell_i \llbracket A \rrbracket_{\ell_i}} \text{CG-ref}}{\Sigma; \Psi \vdash \llbracket \text{ref } A^{\ell_i} \rrbracket_{\ell} <: \llbracket \text{ref } A^{\ell_i} \rrbracket_{\ell'}} \text{Definition 2.86}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell} \times \llbracket \tau'_2 \rrbracket_{\ell'}} \text{CGsub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \times \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} + \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell} + \llbracket \tau'_2 \rrbracket_{\ell'}} \text{CGsub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

5. FGsub-arrow:

$$T_1 = \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.0} = \forall \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.2} = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.3} = (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$c_1 = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.0} = \forall \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.2} = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.3} = (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$c_2 = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\frac{}{\text{Given}}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_2)_\alpha <: \mathbb{C} \gamma \gamma (\tau'_2)_\alpha} \text{IH(1) with } \ell = \ell' = \alpha, \text{ CGsub-monad}$$

P2:

$$\frac{\frac{\frac{}{\text{Given}}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau'_1 <: \tau_1} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma; \Psi \vdash T_{1.3} <: T_{2.3}} \text{CGsub-arrow}$$

P0.1:

$$\frac{\frac{\frac{\frac{}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell \sqsubseteq \ell'}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha) \Rightarrow (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha)}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \ell'_e) \Rightarrow (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \ell_e)} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash c_2 \Rightarrow c_1}$$

P0:

$$\frac{P0.1 \quad \frac{\frac{}{P1}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.3} <: T_{2.3}} \text{CGsub-arrow}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}$$

Main derivation:

$$\frac{P0}{\Sigma; \Psi \vdash \llbracket \tau_1 \xrightarrow{\ell_\xi} \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{CGsub-unit}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_{\ell} <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{Definition 2.86}$$

7. FGsub-forall:

$$T_1 = \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.0} = \forall \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.2} = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.3} = \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$c_1 = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.0} = \forall \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.2} = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.3} = \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$c_2 = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \tau_1 <: \tau_2} \text{Given, Weakening}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\tau_1)_{\alpha'} <: \tau_{2\alpha'}} \text{IH(1) with } \ell = \ell' = \alpha'}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'} <: \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}} \text{IH(1) with } \ell = \ell' = \alpha'$$

P2:

$$\frac{\frac{}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell'_e \sqsubseteq \ell_e)} \text{Given}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \ell'_e) \Rightarrow (\ell \sqcup \gamma \sqsubseteq \ell_e)}$$

P1:

$$\frac{\frac{}{(\ell \sqsubseteq \ell')} \text{Given}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \alpha') \Rightarrow (\ell \sqcup \gamma \sqsubseteq \alpha')}$$

P0:

$$\frac{\frac{P1 \quad P2}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash c_2 \Rightarrow c_1}}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad P3}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash T_{1.2} <: T_{2.2}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}}{\Sigma; \Psi \vdash \llbracket \forall \alpha. \tau_1 \rrbracket_{\ell} <: \llbracket \forall \alpha. \tau_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

8. FGsub-constraint:

$$T_1 = \forall \alpha, \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.0} = \forall \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.1} = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.2} = \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$C_1 = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.0} = \forall \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.1} = (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.2} = \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$C_2 = (c_2 \wedge \ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$

P1:

$$\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash \tau_1 <: \tau_2} \text{Given, Weakening}}{\Sigma, \alpha, \gamma; \Psi \vdash (\tau_1)_\alpha <: \tau_{2\alpha}} \text{IH(1) with } \ell = \ell' = \alpha}{\Sigma, \alpha, \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_1)_\alpha <: \mathbb{C} \gamma \gamma (\tau_2)_\alpha}$$

P0:

$$\frac{\frac{}{\Sigma; \Psi \vdash c_2 \Rightarrow c_1} \text{Given}}{\Sigma, \alpha, \gamma; \Psi \vdash c_2 \wedge (\ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow c_1 \wedge (\ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)} \text{Weakening, } \ell \sqsubseteq \ell', \ell'_e \sqsubseteq \ell_e}{\Sigma, \alpha, \gamma; \Psi \vdash C_2 \Rightarrow C_1}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad P1}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.1} <: T_{2.1}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}}{\Sigma; \Psi \vdash \left[ \left[ c_1 \stackrel{\ell_e}{\Rightarrow} \tau_1 \right]_\ell <: \left[ \left[ c_2 \stackrel{\ell'_e}{\Rightarrow} \tau_2 \right]_{\ell'} \right]} \text{Definition 2.86}$$

□

**Lemma 2.90** (FG  $\rightsquigarrow$  CG: Preservation of well-formedness). *For all  $\Sigma, \Psi$  and  $\ell$  s.t  $FV(\ell) \in \Sigma$  the following hold:*

$$1. \forall \tau. \Sigma; \Psi \vdash \tau \text{ WF} \Rightarrow \Sigma; \Psi \vdash (\tau)_\ell \text{ WF}$$

$$2. \forall \mathbf{A}. \Sigma; \Psi \vdash \mathbf{A} \text{ WF} \Rightarrow \Sigma; \Psi \vdash (\mathbf{A})_\ell \text{ WF}$$

*Proof.* Proof by simultaneous induction on the WF relation of FG

Proof of statement (1)

Let  $\tau = \mathbf{A}^{\ell'}$

$$\frac{\frac{\overline{\text{FV}(\ell') \in \Sigma} \text{ By inversion}}{\text{FV}(\ell' \sqcup \ell) \in \Sigma} \text{ IH(2) on A}}{\Sigma; \Psi \vdash \langle \mathbf{A} \rangle_{\ell' \sqcup \ell} WF} \text{ CG-wff-labeled}}{\Sigma; \Psi \vdash \text{Labeled } \ell' \sqcup \ell \langle \mathbf{A} \rangle_{\ell' \sqcup \ell} WF} \text{ CG-wff-labeled}$$

Proof of statement (2)

We proceed by case analyzing the last rule of given  $WF$  judgment.

1. FG-wff-base:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} WF} \text{ CG-wff-base}$$

2. FG-wff-unit:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} WF} \text{ CG-wff-unit}$$

3. FG-wff-arrow:

P1:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau_2 \rangle_\alpha WF} \text{ IH(1) on } \tau_2}}{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha WF} \text{ CG-wff-monad}}$$

P0:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau_1 \rangle_\beta WF} \text{ IH(1) on } \tau_1 \quad P1}}{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash (\langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-arrow}}$$

Main derivation:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi \vdash ((\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-constraint}}{\Sigma; \Psi \vdash (\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-forall}} \text{ P0}$$

4. FG-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell WF} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash \langle \tau_2 \rangle_\ell WF} \text{ IH(1) on } \tau_2}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell WF} \text{ CG-wff-prod}}$$

5. FG-wff-sum:

$$\frac{\overline{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell WF} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash \langle \tau_2 \rangle_\ell WF} \text{ IH(1) on } \tau_2}{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell WF} \text{ CG-wff-prod}}$$

6. FG-wff-ref:

Let  $\tau = A^{\ell'}$

$$\frac{\frac{\overline{FV(A) = \emptyset} \text{ By inversion} \quad \overline{FV(\ell') = \emptyset} \text{ By inversion}}{\overline{FV(\langle A \rangle_{\ell'}) = \emptyset}} \text{ Lemma 2.91}}{\Sigma; \Psi \vdash \text{ref } \ell' \langle A \rangle_{\ell'} WF} \text{ CG-wff-ref}$$

7. FG-wff-forall:

$$\frac{\frac{\frac{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash \langle \tau \rangle_{\alpha'} WF} \text{ IH(1) on } \tau}{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} WF}} \text{ CG-wff-monad}}{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} WF}} \text{ CG-wff-constraint}}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'}) WF}} \text{ CG-wff-forall}$$

8. FG-wff-constraint:

$$\frac{\frac{\frac{\overline{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau \rangle_{\alpha} WF} \text{ IH(1) on } \tau}{\overline{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha} WF}} \text{ CG-wff-monad}}{\overline{\Sigma, \alpha, \gamma; \Psi \vdash (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha} WF}} \text{ CG-wff-constraint}}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha}) WF}} \text{ CG-wff-forall}$$

□

**Lemma 2.91** (FG  $\rightsquigarrow$  CG: Free variable lemma).  $\forall \Sigma, \ell. FV(\ell) \in \Sigma$ , the following hold

1.  $\forall \tau. FV(\langle \tau \rangle_{\ell}) \subseteq FV(\tau) \cup FV(\ell)$
2.  $\forall A. FV(\langle A \rangle_{\ell}) \subseteq FV(A) \cup FV(\ell)$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$

Proof for (1)

Let  $\tau = A^{\ell_i}$

$$\begin{aligned} & FV(\langle A^{\ell_i} \rangle) \\ &= FV(\text{Labeled } \ell_i \sqcup \ell \langle A \rangle_{\ell_i \sqcup \ell}) && \text{Definition 2.86} \\ &= FV(\ell_i) \cup FV(\ell) \cup FV(\langle A \rangle_{\ell_i \sqcup \ell}) \\ &\subseteq FV(\ell_i) \cup FV(\ell) \cup FV(A) && \text{IH(2) on } A \\ &= FV(A^{\ell_i}) \cup FV(\ell) \end{aligned}$$

Proof for (2)

1.  $A = \mathbf{b}$ :

$$\begin{aligned} & FV(\langle \mathbf{b} \rangle_{\ell}) \\ &= FV(\mathbf{b}) && \text{Definition 2.86} \\ &\subseteq FV(\mathbf{b}) \cup FV(\ell) \end{aligned}$$

2.  $A = \text{unit}$ :

$$\begin{aligned} & FV(\langle \text{unit} \rangle_{\ell}) \\ &= FV(\text{unit}) && \text{Definition 2.86} \\ &\subseteq FV(\text{unit}) \cup FV(\ell) \end{aligned}$$

3.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle_\ell) \\
&= \text{FV}(\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) && \text{Definition 2.86} \\
&= \text{FV}(\ell) \cup \text{FV}(\langle \tau_1 \rangle_\beta) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_2 \rangle_\alpha) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 \xrightarrow{\ell_e} \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
4.  $A = \tau_1 \times \tau_2$ :
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 \times \tau_2 \rangle_\ell) \\
&= \text{FV}(\langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell) && \text{Definition 2.86} \\
&= \text{FV}(\langle \tau_1 \rangle_\ell) \cup \text{FV}(\langle \tau_2 \rangle_\ell) \cup \text{FV}(\ell) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 \times \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
5.  $A = \tau_1 + \tau_2$ :
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 + \tau_2 \rangle_\ell) \\
&= \text{FV}(\langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell) && \text{Definition 2.86} \\
&= \text{FV}(\langle \tau_1 \rangle_\ell) \cup \text{FV}(\langle \tau_2 \rangle_\ell) \cup \text{FV}(\ell) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 + \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
6.  $A = \text{ref } \tau_i$ :
- Let  $\tau_i = A_i^{\ell_i}$
- $$\begin{aligned}
& \text{FV}(\langle \text{ref } \tau_i \rangle_\ell) \\
&= \text{FV}(\text{ref } \ell_i \langle A_i \rangle) && \text{Definition 2.86} \\
&= \text{FV}(\ell_i) \cup \text{FV}(\langle A_i \rangle) \\
&\subseteq \text{FV}(\ell_i) \cup \text{FV}(A_i) \cup \text{FV}(\ell) && \text{IH(2) on } A_i \\
&= \text{FV}(\text{ref } A_i^{\ell_i}) \cup \text{FV}(\ell) \\
&= \text{FV}(\text{ref } \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$
7.  $A = \forall \alpha. (\ell_e, \tau_i)$ :
- $$\begin{aligned}
& \text{FV}(\langle \forall \alpha. (\ell_e, \tau_i) \rangle) \\
&= \text{FV}(\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau_i \rangle_{\alpha'}) && \text{Definition 2.86} \\
&= \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_i \rangle) \\
&\subseteq \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_i) && \text{IH(1) on } \tau_i \\
&= \text{FV}(\ell) \cup \text{FV}(\forall \alpha. (\ell_e, \tau_i))
\end{aligned}$$
8.  $A = c \xrightarrow{\ell_e} \tau_i$ :
- $$\begin{aligned}
& \text{FV}(\langle c \xrightarrow{\ell_e} \tau_i \rangle) \\
&= \text{FV}(\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau_i \rangle_\alpha) && \text{Definition 2.86} \\
&= \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\langle \tau_i \rangle) \cup \text{FV}(\ell) \\
&\subseteq \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\tau_i) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_i \\
&= \text{FV}(c \xrightarrow{\ell_e} \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$

□



### 2.5.3 Logical relation for FG to CG translation

**Definition 2.92** (FG  $\rightsquigarrow$  CG:  ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$   
 $\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

**Definition 2.93** (FG  $\rightsquigarrow$  CG:  $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$   
 $\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

**Definition 2.94** (FG  $\rightsquigarrow$  CG: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket \wedge {}^sv = {}^tv\} \\
\llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \mid \\
&\quad ({}^s\theta, m, {}^sv_1, {}^tv_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^sv, \text{inl } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\
&\quad \{({}^s\theta, m, \text{inr } {}^sv, \text{inr } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \Lambda \Lambda(\nu(\lambda x. e_t))) \mid \\
&\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \implies \\
&\quad ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \forall \alpha. (\ell_e, \tau) \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \Lambda e_s, \Lambda \Lambda(\nu(e_t))) \mid \\
&\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \nu e_s, \Lambda(\nu(e_t))) \mid \\
&\quad \mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \text{ref } \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, a_s, a_t) \mid {}^s\theta(a_s) = \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}\} \\
\llbracket \mathbf{A}' \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, \text{Lb}({}^tv)) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \mathbf{A} \rrbracket_V^{\hat{\beta}}\}
\end{aligned}$$

**Definition 2.95** (FG  $\rightsquigarrow$  CG: Unary expression relation).

$$\begin{aligned}
\llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\
&\quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\
&\quad \exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \\
&\quad \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}\}
\end{aligned}$$

**Definition 2.96** (FG  $\rightsquigarrow$  CG: Unary heap well formedness).

$$\begin{aligned}
(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\
&\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\
&\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a_1) \rrbracket_V^{\hat{\beta}}
\end{aligned}$$

**Definition 2.97** (FG  $\rightsquigarrow$  CG: Label substitution).  $\sigma : \text{Lvar} \mapsto \text{Label}$

**Definition 2.98** (FG  $\rightsquigarrow$  CG: Value substitution to values).  $\delta^s : \text{Var} \mapsto \text{Val}, \delta^t : \text{Var} \mapsto \text{Val}$

**Definition 2.99** (FG  $\rightsquigarrow$  CG: Unary interpretation of  $\Gamma$ ).

$$\begin{aligned}
\llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\
&\quad \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\}
\end{aligned}$$

#### 2.5.4 Soundness proof for FG to CG translation

**Lemma 2.100** (FG  $\rightsquigarrow$  CG: Monotonicity).  $\forall {}^s\theta, {}^s\theta', n, {}^s v, {}^t v, n', \beta, \beta'$ .

1.  $\forall A. ({}^s\theta, n, {}^s v, {}^t v) \in [A]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in [A]_V^{\hat{\beta}'}$
2.  $\forall \tau. ({}^s\theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$

*Proof.* Proof by simultaneous induction on A and  $\tau$

Proof of statement (1)

We case analyze A in the last step

1. Case **b**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [b]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [b]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^s v, {}^t v) \in [b]_V^{\hat{\beta}}$  therefore from Definition 2.94 we know that  ${}^s v \in \llbracket b \rrbracket \wedge {}^t v \in \llbracket b \rrbracket$  and  ${}^s v = {}^t v$

Therefore from Definition 2.94 we get the desired

2. Case **unit**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}}$  therefore from Definition 2.94 we know that  ${}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket$

Therefore from Definition 2.94 we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that  ${}^s v = ({}^s v_1, {}^s v_2)$  and  ${}^t v = ({}^t v_1, {}^t v_2)$ .

We also know that  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1:  $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

IH2:  $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 2.94, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

IH:  $({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 2.94 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that

${}^sv$  is of the form  $\lambda x.e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\Lambda\Lambda\Lambda(\nu(\lambda x.e_t))$  (for some  $e_t$ ) s.t

$$({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'_1} \quad (\text{A0})$$

Similarly from Definition 2.94 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''} \implies ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

Instantiating (A0) with  ${}^s\theta'', {}^s v_2, {}^t v_2, k, \hat{\beta}''$  since  
 ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get  
 $({}^s\theta'', k, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$

6. Case  $\forall\alpha.\tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that  ${}^s v = \Lambda e'_s$  and  ${}^t v = \Lambda\Lambda(\nu(e_t))$  s.t

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'_1} \quad (\text{F0})$$

Similarly from Definition 2.94 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''_1}$$

This means we are given  ${}^s\theta''_1 \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

Instantiating (F0) with  ${}^s\theta''_1, k, \hat{\beta}''$  since  ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$   
therefore we get

$$({}^s\theta''_1, k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

7. Case  $c \stackrel{\ell_e}{\Rightarrow} \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that  ${}^s v = \nu(e'_s)$  and  ${}^t v = \Lambda\Lambda(\nu(e_t))$ . And

$$\mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, e_s, e_t) \in [\tau]_E^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 2.94 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''_1}$$

This means we are given  $\mathcal{L} \models c, {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

Since  $\mathcal{L} \models c$  and instantiating (C0) with  ${}^s\theta''_1, k, \hat{\beta}''$  since  ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta$ ,  $k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

8. Case ref  $\tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that  ${}^s v = a_s$  and  ${}^t v = a_t$ . We also know that  ${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$

From Definition 2.94, Definition 2.92 and Definition 2.93 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

Proof of Statement (2)

Let  $\tau = \mathbf{A}^{\ell''}$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \text{ and } ({}^s\theta, n, {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}'}$$

This means from Definition 2.94 we need to prove

$$({}^s\theta', n', {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

**Lemma 2.101** (FG  $\rightsquigarrow$  CG: Unary monotonicity for  $\Gamma$ ).  $\forall {}^s\theta, {}^s\theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies ({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 2.99 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$$

And again from Definition 2.99 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

- $dom(\Gamma) \subseteq dom(\delta^s) \wedge dom(\Gamma) \subseteq dom(\delta^t)$ :

Given

- $\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x_i \in dom(\Gamma). (^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 2.100 we get

$$\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

□

**Lemma 2.102** (FG  $\rightsquigarrow$  CG: Unary monotonicity for  $H$ ).  $\forall ^s\theta, H_s, H_t, n, n', \hat{\beta}$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}'} ^s\theta$

From Definition 2.96 it is given that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 2.96 we are required to prove that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

- $dom(^s\theta) \subseteq dom(H_S)$ :

Given

- $\hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 2.100 we get

$$\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

□

**Lemma 2.103** (Coercion lemma).  $\forall H, e, v$ .

$$(H, e) \Downarrow_-^f (H', \text{Lb } v) \implies (H, \text{coerce\_taint } e) \Downarrow_-^f (H', \text{Lb } v)$$

*Proof.* Given:  $(H, e) \Downarrow_-^f (H', \text{Lb } v)$

To prove:  $(H, \text{coerce\_taint } e) \Downarrow_-^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and cg-app it suffices to prove that

$$(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_-^f (H', \text{Lb } v)$$

From cg-tolabeled it suffices to prove that  
 $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From cg-bind it suffices to prove that

1.  $(H, e) \Downarrow_{-}^f (H'_1, v_1)$ :

We are given that  $(H, e) \Downarrow_{-}^f (H', v)$  therefore we have  $H'_1 = H'$  and  $v'_1 = \text{Lb } v$

2.  $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$ :

It suffices to prove that

$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v)$ :

We get this directly from cg-unlabel

□

**Theorem 2.104** (FG  $\rightsquigarrow$  CG: Fundamental theorem).  $\forall \Sigma, \Psi, \Gamma, \tau, e_s, e_t, pc, \mathcal{L}, \delta^s, \delta^t, \sigma, {}^s\theta, n, \hat{\beta}$ .

$\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge$

$\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

$\implies$

$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x \delta^s, \text{ret}(x) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

From Definition 2.95 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(x) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge \\ & ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  $i = 0, {}^s v = x \delta^s$ . Also from cg-ret we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$

And we are required to prove

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ : Given

(b)  $({}^s\theta, n, {}^sv, {}^tv) \in [\tau]_V^{\hat{\beta}}$ :

Since we are given  $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_V^{\hat{\beta}}$ , therefore from Definition 2.99 we get  $({}^s\theta, n, {}^sv, {}^tv) \in [\tau \sigma]_V^{\hat{\beta}}$

2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t))))} \text{FC-lam}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\lambda x. e_s) \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^{\hat{\beta}}$

From Definition 2.95 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^sv. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv) \implies \\ & \exists H'_t, {}^tv. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \Downarrow^f (H'_t, {}^tv) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^sv$  s.t  $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv)$

From fg-val we know that  ${}^sv = (\lambda x. e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret, cg-label and cg-FI we know that  $H'_t = H_t$  and  ${}^tv = (\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ : Given

(b)  $({}^s\theta, n, \lambda x. e_s \delta^s, (\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}}$ :

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \lambda x. e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(\lambda x. e_t))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\begin{aligned} & \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'} \implies \\ & ({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \end{aligned}$$

This further means that given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$

And we are required to prove

$$({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \quad (\text{F-L0})$$



Since we are given  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'}$ , therefore from Definition 2.99 and Lemma 2.101 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\}) \sigma]_{V'}^{\hat{\beta}'}$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2 \sigma]_E^{\hat{\beta}'}$$

We get (F-L0) directly from IH

### 3. FC-app:

$$\frac{\begin{array}{c} \Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \rightsquigarrow e_{t1} \\ \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell \end{array}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square\square]\bullet) b))))} \text{FC-app}}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square\square]\bullet) b)))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 it suffices to prove

$$\begin{array}{l} \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \end{array}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{array}{l} \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-A0}) \end{array}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{array}{l} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \end{array}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2)^\ell \sigma]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.0})$$

Since we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2)^\ell \sigma]_{V}^{\hat{\beta}'_1}$  therefore from Definition 2.94 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2) \sigma]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 2.94 we know that  ${}^s v_1 = \lambda x. e'_s$  and  ${}^t v_i = \Lambda \Lambda \Lambda (\nu (\lambda x. e'_t))$  s.t

$$\forall {}^s \theta''_1 \sqsupseteq {}^s \theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1.$$

$$({}^s \theta''_1, l, {}^s v', {}^t v') \in [\tau_1 \sigma]_{V}^{\hat{\beta}''_1} \implies ({}^s \theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}''_1} \quad (\text{F-A1})$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_{E}^{\hat{\beta}'_1}$$

This means from Definition 2.95 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2 \delta^t) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'_2} \end{aligned}$$

We instantiate with  $H'_{s1}, H'_{t1}$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s1}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$ .

This means we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'_2} \quad (\text{F-A2})$$

We instantiate (F-A1) with  $\theta''_1$  as  $\theta'_2$ ,  ${}^s v'$  as  ${}^s v_2$ ,  ${}^t v'$  as  ${}^t v_2$ ,  $l$  as  $n - j - k$  and  $\hat{\beta}''_1$  as  $\hat{\beta}'_2$ . Therefore we get

$$({}^s \theta'_2, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'_2}$$

From Definition 2.95 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s3}, {}^s v_3) \implies \\ \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_{V}^{\hat{\beta}'_3} \end{aligned}$$

Instantiating with  $H'_{s2}, H'_{t2}$ . since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists a < i - j - k < n - j - k$  s.t  $(H'_{s2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s3}, {}^s v_3)$

Therefore we have

$$\exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2.$$

$$(n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3})$$

Let  $\tau_2 \sigma = A_2^{\ell_i}$ , since  $\tau_2 \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}$$

Therefore from Definition 2.94 we know that

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, \text{Lb}^t v_{3i}) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose  $H'_t$  as  $H'_{t3}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{3i})$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i})$$

From cg-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

From cg-bind again it suffices to prove

- \*  $(H'_{t2}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel and (F-A1) we know that  $H'_{t31} = H'_{t2}$  and  ${}^t v_{t2} = {}^t v_i = \Lambda\Lambda\Lambda(\nu(\lambda x.e'_t))$

- \*  $((c[] [] \bullet) b) [{}^t v_2/b] [{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}:$

It suffices to prove that

$$(((\Lambda\Lambda\Lambda(\nu(\lambda x.e'_t))) [] [] \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-FE it suffices to prove that

$$(((\Lambda\Lambda(\nu(\lambda x.e'_t))) [] \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

Again from cg-FE applied two times it suffices to prove that

$$((\nu(\lambda x.e'_t) \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-CE it suffices to prove that

$$(((\lambda x.e'_t) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-app we know that

$${}^t v_{t21} = e'_t [{}^t v_2/x] \delta^t$$

- \*  $(H'_{t2}, {}^t v_{t21}) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

We get this from (F-A3) and (F-A3.1)

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}:$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_3$  and  $\hat{\beta}'$  as  $\hat{\beta}'_3$ . From fg-app we know that  $i = j + k + a + 1$ ,  ${}^s v = {}^s v_3$  and  $H'_s = H'_{s3}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t3}$  and  ${}^t v = \text{Lb}({}^t v_3)$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_3} {}^s\theta'$  from (F-A3) and Lemma 2.102

Since  ${}^t v = \text{Lb}({}^t v_3)$  therefore from Definition 2.94 it suffices to prove that

$$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}'_3}$$

We get this directly from (F-A3) and Lemma 2.100

4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v_1, {}^s v_2$  s.t  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2 \wedge ({}^s \theta'_1, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists k < i - j < n - j$  s.t  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose  $H_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_1, {}^t v_2)$

$$(a) (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2)):$$

From cg-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$ :  
From (F-P1) we know that  $H'_{tb1} = H'_{t1}$  and  ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))[{}^t v_1/a] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From cg-bind it suffices to prove that
  - $(H_t, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$ :  
From (F-P2) we know that  $H'_{tb2} = H'_{t2}$  and  ${}^t v_{tb2} = {}^t v_2$
  - $(H'_{t2}, \text{ret}(\text{Lb}(a, b)))[{}^t v_1/a][{}^t v_2/b] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From cg-ret, (F-P1) and (F-P2)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$  and since from fg-prod  $i = j + k + 1$  and  $H'_s = H'_{s2}$ . Therefore from (F-P2) and Lemma 2.102 we get

$$(n - i, H'_s, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2$$

In order to prove  $({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^{\hat{\beta}'_2}$

From Definition 2.94 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}'_2}$$

Since  ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$  therefore we get the desired from (F-P1), (F-P2), Definition 2.94 and Lemma 2.100

## 5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t \in [\tau_1 \sigma]_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 2.94 we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 2.94 we know that  ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$  and  ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$  s.t

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$

Let  $\tau_1 \sigma = A_1^{\ell_i}$ , since  $\tau_1 \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$\text{Since } ({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1}$$

Therefore from Definition 2.94 we know that

$$({}^s\theta'_1, n - j, {}^s v_{i1}, \text{Lb}({}^t v_{i1})) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{i1})$  as we need to prove

(a)  $(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$ :

From Lemma 2.103 it suffices to prove that

$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1)$ :

We get this from (F-F1)

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$ :

Again from cg-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :

Since  ${}^t v_1 = \text{Lb}({}^t v_{i1}, {}^t v_{i2})$  from (F-F1.1) and (F-F1.2) therefore we get the desired from cg-unlabel

So,  $H_{t21} = H'_{t1}$  and  ${}^t v_{t21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$ :

We get this from cg-fst, cg-ret and (F-F1.2) and (F-F1.3)

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_{\mathcal{V}}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . And from fg-fst we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-F1) and Lemma 2.102 we get

$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$

Since from fg-fst we know that  ${}^s v = {}^s v_{i1}$  therefore from (F-F1.2) and Lemma 2.100 we get

$({}^s \theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{\mathcal{V}}^{\hat{\beta}'_1}$

## 6. FC-snd:

Symmetric reasoning as in the FC-fst case

## 7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))} \text{inl}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\mathcal{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\mathcal{E}}^{\hat{\beta}}$

This means from Definition 2.95 we have

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies$   
 $\exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$

$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'}$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_t, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-IL1}) \end{aligned}$$

In order to prove (F-IL0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $(\text{Lb inl}({}^t v_1))$  and we need to prove:

$$(a) (H'_{t1}, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-bind it suffices to prove that

$$i. (H'_{t1}, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$$

From (F-IL1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

$$ii. (H'_{t1}, \text{ret}(\text{Lbinl}(a))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

We get this from cg-ret, (F-IL1)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . Since from fg-inl we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-IL1) and Lemma 2.102 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$$

Now we need to prove  $({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'}$

Since  ${}^s v = \text{inl } {}^s v_1$  and  ${}^t v = \text{Lb}(\text{inl}({}^t v_1))$  therefore from Definition 2.94 it suffices to prove that

$$({}^s \theta', n - i, \text{inl } {}^s v_1, \text{inl } {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_{V}^{\hat{\beta}'}$$

Since from (F-IL1) we know that  $({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

Therefore from Lemma 2.100 and Definition 2.94 we get

$$({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2) \sigma]_{V}^{\hat{\beta}'}$$

## 8. FC-inr:

Symmetric reasoning as in the FC-inl case



9. FC-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-C0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 2.94 we know that

$$\exists^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

2 cases arise

(a)  ${}^s v_1 = \text{inl}({}^s v_{i1})$  and  ${}^t v_i = \text{inl}({}^t v_{i1})$ :

Also from Lemma 2.101 and Definition 2.99 we know that

$$({}^s \theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\}) \sigma]_{V'}^{\hat{\beta}'_1}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \end{aligned}$$

$$(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t.  $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1.$$

$$(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2})$$

Let  $\tau \sigma = \mathbf{A}_2^{\ell_i}$ , since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

Therefore from Definition 2.94 we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{2i})$

And we need to prove:

i.  $(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

•  $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :

From (F-C1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

•  $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{2i}))$ :

From cg-bind it suffices to prove that

–  $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :

Since from (F-C1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  therefore from cg-unlabel we know that

$$H'_{t21} = H'_{t1} \text{ and } {}^t v_{t21} = {}^t v_i$$

–  $(\text{case}(b, x.e_{t1}, y.e_{t2}) [{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$ :

Since we know that in this case  ${}^t v_i = \text{inl}({}^t v_{i1})$

Therefore from cg-case we know that  ${}^t v_{t22} = e_{t1} [{}^t v_{i1}/x] \delta^t$

–  $(H'_{t1}, e_{t1}[{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \text{Lb}{}^t v_{2i})$ :  
 We get this from (F-C2) and (F-C2.1)

ii.  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'}$ :  
 We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . Since from fg-case we know that  $i = j + k + 1$  and  $H'_s = H'_{s2}$  therefore from (F-C2) and Lemma 2.102 we get

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$$

Now we need to prove  $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

Since  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  and since from (F-C2) we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

Therefore from Lemma 2.100 and Definition 2.94 we get

$$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

(b)  ${}^s v_1 = \text{inr}({}^s v_{i1})$  and  ${}^t v_1 = \text{inr}({}^t v_{i1})$ :

Symmetric reasoning as in the previous case

## 10. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_s : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t))))} \text{FI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \lfloor \Gamma \sigma \rfloor_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \Lambda e_s \delta^s, \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \in \lfloor \tau \sigma \rfloor_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ : Given

(b)  $({}^s\theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_t)))) \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau))^\perp \sigma]_{\hat{\beta}}^{\hat{\beta}}$ :

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(e_t))) \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau)) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{E}}^{\hat{\beta}'_1}$$

This further means that given  ${}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we need to prove

$$({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g]]_{\hat{E}}^{\hat{\beta}'_1} \quad (\text{F-FI0})$$

$$\underline{\text{IH}}: ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha_g \mapsto \ell'\}]_{\hat{E}}^{\hat{\beta}'_1}$$

We get (F-FI0) directly from IH

## 11. FC-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\forall\alpha_g.(\ell_e, \tau))^\ell \rightsquigarrow e_t \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash_{pc} \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s [] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))} \text{FE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s [] \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet))) \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s []) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, e_s []) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau))^\ell \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall\alpha_g.(\ell_e, \tau))^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t.  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-FE1}) \end{aligned}$$

Since from (F-FE1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V}^{\hat{\beta}'}$  therefore from Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\forall \alpha_g. (\ell_e, \tau)) \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-FE1.1})$$

Therefore from Definition 2.94 we have

$$\begin{aligned} & {}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \Lambda \nu e'_t \\ & \forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \ell'' \in \mathcal{L}, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in [\tau[\ell''/\alpha_g] \sigma]_{E}^{\hat{\beta}'_1} \quad (\text{F-FE1.2}) \end{aligned}$$

We instantiate with  ${}^s \theta'_1, \ell', n - j - 1, \hat{\beta}'$  we get  $({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in [\tau[\ell'/\alpha_g] \sigma]_{E}^{\hat{\beta}'}$

From Definition 2.95 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  and from fg-FE we know that  $i = j + k + 1 < n$  therefore we know that  $k < n - j - 1$  s.t.  $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-FE1.3}) \end{aligned}$$

Let  $\tau[\ell'/\alpha] \sigma = \mathbf{A}^{\ell_i}$ , since  $\tau[\ell'/\alpha] \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$\begin{aligned} & ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \\ & \text{Therefore from Definition 2.94 we know that} \\ & ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-FE1.4}) \end{aligned}$$

In order to prove (F-FE0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{2i})$ . We need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i})):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :  
From (F-FE1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.b[][\bullet]))[{}^t v_1/a] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :  
Again from cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a)[^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12})$ :  
From (F-FE1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$   
Therefore from cg-unlabel we have  $H'_{t12} = H'_{t1}$  and  ${}^t v_{t12} = {}^t v_i$
- $(b[\square\square\square\bullet][^t v_i/b] \delta^t) \Downarrow {}^t v_{t13}$ :  
From (F-FE1.2) we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_i = \Lambda\Lambda\Lambda\nu e'_t$   
Therefore from cg-FE and cg-CE we know that  ${}^t v_{t13} = e'_t$
- $(H'_{t1}, e'_t \Downarrow^f (H'_{t2}, \text{Lb}{}^t v_{2i}))$   
From (F-FE1.3) and (F-FE1.4) we get the desired.

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}''$ . From fg-FE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v'_2$ ,  ${}^t v = {}^t v'_2$ ,  $H'_s = H'_{s2}$  and  $H'_t = H'_{t2}$ .

Therefore from (F-FE1.3) we get the  $(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2$

To prove:  $({}^s \theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''}$

We get this directly from (F-FE1.3)

## 12. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{ CI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \nu e \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ : Given

(b)  $({}^s\theta, n, \nu e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\perp \sigma \rrbracket_V^{\hat{\beta}}$ :

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\mathcal{L} \models c \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'}$$

This further means that given  $\mathcal{L} \models c \sigma$  and  ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'} \quad (\text{F-CI0})$$

$$\underline{\text{IH}}: ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

### 13. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (c \xrightarrow{\ell_\xi} \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b[][\bullet])))} \text{CE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s \bullet \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b[][\bullet]))) \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v.(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b[][\bullet]))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b[][\bullet]))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, e_s \llbracket \cdot \rrbracket) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-CE1}) \end{aligned}$$

Since from (F-CE1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$  therefore from Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in \llbracket (c \xrightarrow{\ell_\xi} \tau) \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-CE1.1})$$

Therefore from Definition 2.94 we have

$$\begin{aligned} & {}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \nu e'_t \\ & \mathcal{L} \models c \sigma \implies \forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{F-CE1.2}) \end{aligned}$$

Since we know that  $\mathcal{L} \models c \sigma$ , we instantiate with  ${}^s \theta'_1, n - j - 1, \hat{\beta}'$  to get

$$({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1}$$

From Definition 2.95 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_1} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, e_s \llbracket \cdot \rrbracket) \Downarrow_i (H'_s, {}^s v)$  and since from fg-CE we know that  $i = j + k + 1 < n$  therefore we know that  $k < n - j - 1$  s.t  $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_1} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \quad (\text{F-CE1.3}) \end{aligned}$$

Let  $\tau \sigma = \mathbf{A}^{\ell_i}$ , since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1}$$

Therefore from Definition 2.94 we know that

$$({}^s \theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}({}^t v_2)) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \quad (\text{F-CE1.4})$$

In order to prove (F-CE0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_2)$ . We need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b \llbracket \cdot \rrbracket \bullet)))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_2)):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b \llbracket \cdot \rrbracket \bullet)))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_2))$$

From cg-bind it suffices to prove that



- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :  
From (F-CE1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.b[][\bullet]) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb} {}^t v_{2i})$ :  
Again from cg-bind it suffices to prove that
  - $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12})$ :  
From (F-CE1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$   
Therefore from cg-unlabel we have  $H'_{t12} = H'_{t1}$  and  ${}^t v_{t12} = {}^t v_i$
  - $(b[][\bullet]) [{}^t v_i/b] \delta^t \Downarrow {}^t v_{t13}$ :  
From (F-CE1.2) we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_i = \Lambda \Lambda \nu e'_t$   
Therefore from cg-FE and cg-CE we know that  ${}^t v_{t13} = e'_t$
  - $(H'_{t1}, e'_t \Downarrow^f (H'_{t2}, \text{Lb} {}^t v_{2i}))$   
We get the desired from From (F-CE1.3) and (F-CE1.4)

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}''$ . From fg-CE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v'_2$ ,  ${}^t v = {}^t v'_2$ ,  $H'_s = H'_{s2}$  and  $H'_t = H'_{t2}$ .

Therefore from (F-CE1.3) we get the  $(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2$

To prove:  $({}^s \theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau \sigma]_{V}^{\hat{\beta}''}$

From (F-CE1.3) we know that  $({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V}^{\hat{\beta}''}$

14. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b)))} \text{ref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{new } (e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b)) \delta^t) \delta^t) \in [(\text{ref } \tau)^\perp \sigma]_{E}^{\hat{\beta}}$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ .

And we are required to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore we know that  $\exists j < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_1} \quad (\text{F-R1}) \end{aligned}$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_1 \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^t v = \text{Lb}(a_t)$ ,  ${}^s \theta'$  as  ${}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1})$ :  
From (F-R1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$
- $(H'_1, \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2})$ :

From cg-bind it suffices to prove that

- i.  $(H'_1, \text{new } (a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2})$ :  
From cg-new we know that  $H'_{t2} = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_{t2} = a_t$
- ii.  $(H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_t)$ :  
From cg-ret we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_t = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}:$$

From (F-R1) we know that  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$  and since  $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$ ,  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$

Therefore from Definition 2.96 and Lemma 2.102 we get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'$

To prove:  $({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}$

Since we know that  ${}^s v = a_s$  and  ${}^t v = \text{Lb } a_t$  therefore we need to prove

$$({}^s \theta', n - i, a_s, \text{Lb}(a_t)) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}$$

From Definition 2.94 it suffices to prove that

$$({}^s \theta', n - i, a_s, a_t) \in \lfloor (\text{ref } \tau) \sigma \rfloor_V^{\hat{\beta}'_1}$$

Again from Definition 2.94 it suffices to prove that

$${}^s \theta'(a_s) = \tau \sigma \wedge (a_s, a_t) \in \hat{\beta}'_1$$

We get this by construction

15. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\text{ref } \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{deref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \in [\tau' \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, !e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$

From Definition 2.94 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 2.94 we know that  ${}^s v_1 = a_s$  and  ${}^t v_i = a_t$

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Let  $\tau' \sigma = A^{\ell_i}$ , since  $\tau' \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

Let  $v_g = H_t(a_t)$  therefore from Definition 1.76 we have

$$({}^s\theta, n-1, H_s(a_s), \text{Lb}v_{gi}) \in [\tau']_{\hat{V}}^{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $H'_{t1}(a_t) = v_g = \text{Lb}v_{gi}$

$$(a) \ (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \delta^t \Downarrow^f (H'_{t1}, \text{Lb}v_{gi}):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \delta^t \Downarrow^f (H'_{t1}, \text{Lb}v_{gi})$$

From cg-bind it suffices to prove

$$i. \ (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$

$$ii. \ (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t2}):$$

From cg-bind it suffices to prove that

$$A. \ (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$

$$B. \ (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_t, \text{Lb}v_{gi}):$$

Since from (F-DR1.2) we know that  ${}^t v_i = a_t$  therefore from cg-deref we know that  $H'_t = H'_{t1}$  and  ${}^t v = H'_{t1}(a_t) = v_g = \text{Lb}v_{gi}$

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n-i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$

Therefore from (F-DR1) we get  $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$  and since  $i = j+1$  therefore from Lemma 2.102 we get  $(n-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$

Since from (F-DR1.2) we know that  $(a_s, a_t) \in \hat{\beta}'_1$  and  ${}^s\theta'_1(a_s) = \tau$ . Also from (F-DR1) we have  $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$ . Therefore from Definition 2.95 we have  $(n-j-1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s\theta'_1(a_s)]_{\hat{V}}^{\hat{\beta}'_1}$

Since  $i = j+1$ ,  ${}^s\theta'_1(a_s) = \tau$ ,  $H'_{s1}(a_s) = {}^s v$  and  $H'_{t1}(a_t) = {}^t v$

Therefore we get

$$({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

Finally from Lemma 2.105 we get

$$({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'}$$

## 16. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} := e_{s2} : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{assign}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \in \llbracket \text{unit } \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsubseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsubseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'} \quad (\text{F-AN0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsubseteq {}^s\theta, \hat{\beta}'_1 \sqsubseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t.  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsubseteq {}^s\theta, \hat{\beta}'_1 \sqsubseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$  therefore from Definition 2.94 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in \llbracket (\text{ref } \tau) \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 2.94 this further means that

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \end{aligned}$$

$$(n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < n - j$  s.t.  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1.$$

$$(n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2} \wedge$$

(F-AN2)

In order to prove (F-AN0) we choose  $H'_t$  as  $H'_{t2}[a_t \mapsto {}^s v_2]$ ,  ${}^t v$  as ()

We need to prove

$$(a) (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

$$- (H_t, (\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From cg-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$$

where  ${}^t v_T = \text{Lb}^t v_{Ti}$

From cg-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$

From (F-AN1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12}):$

From cg-bind it suffices to prove

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13}):$

From (F-AN2) we know that  $H'_{t13} = H'_{t2}$  and  ${}^t v_{t13} = {}^t v_2$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From cg-bind it suffices to prove that

- \*  $(H'_{t1}, \text{unlabel } a [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\mathcal{V}}^{\hat{\beta}'_1}$$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i = a_t$

- \*  $(H'_{t1}, (c := b) [{}^t v_1/a] [{}^t v_2/b] [{}^t v_i/c] \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti}):$

From cg-assign we know that  $H'_T = H'_{t1}[a_t \mapsto {}^t v_2]$  and  ${}^t v_{Ti} = ()$

Since  ${}^t v_{t12} = {}^t v_{Ti} = ()$  therefore  ${}^t v_T = \text{Lb}()$

$$- (H'_T, \text{ret}() [{}^t v_T/d]) \delta^t) \Downarrow^f (H'_t, ()):$$

From cg-ret and cg-val

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2}:$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$

In order to prove  $(n - i, H'_s, H'_t) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2$  it suffices to prove

- $\text{dom}({}^s\theta'_2) \subseteq \text{dom}(H'_s)$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$  therefore from Definition 2.96 we get  $\text{dom}({}^s\theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s\theta'_2) \times \text{dom}(H'_t))$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$  therefore from Definition 2.96 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s\theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s\theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s\theta'_2(a_1)]_V^{\hat{\beta}'_2}$ :

$$\forall (a_1, a_2) \in \hat{\beta}'_2.$$

- $a_1 = a_s$  and  $a_1 = a_t$ :

Since from (F-AN2) we know that  $({}^s\theta'_2, n - j - k, {}^sv_2, {}^tv_2) \in [\tau \sigma]_V^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 2.92 we know that  ${}^s\theta'_2(a_1) = \tau \sigma$

Therefore from Lemma 2.100 we get

$$({}^s\theta'_2, n - i - 1, {}^sv_2, {}^tv_2) \in [\tau \sigma]_V^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$  and  $a_1 \neq a_t$ :

From (F-AN2) since we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$  therefore from Definition 2.96 we get

$$({}^s\theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s\theta'_2(a_1) \sigma]_V^{\hat{\beta}'_2}$$

Since  $i = j + k + 1$  therefore from Lemma 2.100 we get

$$({}^s\theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s\theta'_2(a_1) \sigma]_V^{\hat{\beta}'_2}$$

- $a_1 = a_s$  and  $a_1 \neq a_t$ :

This case cannot arise

- $a_1 \neq a_s$  and  $a_1 = a_t$ :

This case cannot arise

And in order to prove  $({}^s\theta', n - i, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'_2}$

Since we know that  ${}^sv = ()$  and  ${}^tv = ()$  therefore from Definition 2.94 we get  $({}^s\theta', n - i, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'_2}$

□

**Lemma 2.105** (FG  $\rightsquigarrow$  CG: Semantic Subtyping lemma). *The following holds:*

$$\forall \Sigma, \Psi, \sigma, \mathcal{L}, \hat{\beta}.$$

1.  $\forall A, A'$ .

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies [(A \sigma)]_V^{\hat{\beta}} \subseteq [(A' \sigma)]_V^{\hat{\beta}}$$

2.  $\forall \tau, \tau'$ .

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau'_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{It suffices to prove: } \forall ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}. \\ ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, m$  and  $\lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))$  s.t

$$({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 2.94 we are given:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-L0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 2.94, it suffices to prove:

$$\forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L1})$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$  s.t  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with  ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$ . Since we have  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$  therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2 \sigma)]_E^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2}$$



## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, n$  and  ${}^s v_1, {}^s v_2, {}^t v_1, {}^t v_2$  s.t

$$({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Therefore from Definition 2.94 we are given:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.94, it suffices to prove:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that  $({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}}$  therefore from IH1 we have  $({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}}$

Similarly since we have  $({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$  from (S-P0) therefore from IH2 we have  $({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}}$

## 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:  $\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^s v = \text{inl } {}^s v_i$  and  ${}^t v = \text{inl } {}^t v_i$ :

From Definition 2.94 we are given:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1' \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b)  ${}^s v = \text{inr } {}^s v_i$  and  ${}^t v = \text{inr } {}^t v_i$ :

Symmetric reasoning as in the previous case

#### 4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $[(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

This means that given  $({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 2.94 we have:

$$\forall {}^s \theta'_1 \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_1, j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-F0})$$

And we need to prove

$$({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

Again from Definition 2.94 it means we need to prove

$$\forall {}^s \theta'_2 \sqsupseteq {}^s \theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s \theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2}$$

This means that given  ${}^s \theta'_2 \sqsupseteq {}^s \theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s \theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-F1})$$

Instantiating (S-F0) with  ${}^s \theta'_2, k, \ell'', \hat{\beta}'_2$  and we get

$$({}^s \theta'_2, k, e_s, e_t) \in [\tau_1[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2}$$

IH:  $[(\tau_1 \sigma \cup \{\alpha \mapsto \ell''\})]_E^{\hat{\beta}'_2} \subseteq [(\tau_2 \sigma \cup \{\alpha \mapsto \ell''\})]_E^{\hat{\beta}'_2}$  (Statement 2(b))

Therefore from IH we get the desired

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \Longrightarrow c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell'_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove:  $\llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given:  $({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.94 we are given:

$$\mathcal{L} \models c_1 \sigma \Longrightarrow \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-C0})$$

And it suffices to prove:

$$({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Again from Definition 2.94 it means that we need to prove:

$$\mathcal{L} \models c_2 \sigma \Longrightarrow \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2}$$

This means that given that  $\mathcal{L} \models c_2 \sigma$  and  ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-C1})$$

Instantiating (S-C0) with  ${}^s\theta'_2, k, \hat{\beta}'_2$  we get  $({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_2}$

IH:  $\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_2} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_2}$  (Statement 2(b))

Finally from IH we get  $({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2}$

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

We get this directly from Definition 2.94

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{b}) \sigma) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 2.94

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{unit}) \sigma) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 2.94

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell' \sqsubseteq \ell'' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^{\ell'} <: \mathbf{A}^{\ell''}} \text{FGsub-label}$$

To prove:  $\llbracket ((\mathbf{A}^{\ell'}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{A}^{\ell''}) \sigma) \rrbracket_V^{\hat{\beta}}$

This means from Definition 2.94 we need to prove

$$\forall ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell'} \sigma \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell''} \sigma \rrbracket_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell'} \sigma \rrbracket_V^{\hat{\beta}}$

From Definition 2.94 it further means that we are given

$$({}^s\theta, n, {}^s v, {}^t v_i) \in \llbracket \mathbf{A} \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-LB0})$$

And we need to prove

$$({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell''} \sigma \rrbracket_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$({}^s\theta, n, {}^s v, {}^t v_i) \in \llbracket \mathbf{A}' \sigma \rrbracket_V^{\hat{\beta}}$$

Since  $\ell' \sqsubseteq \ell''$  and  $\mathbf{A}' <: \mathbf{A}''$  therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$

This means we need to prove that

$$\forall ({}^s\theta, n, e_s, e_t) \in \llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}}. ({}^s\theta, n, e_s, e_t) \in \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$$

This means given  $({}^s\theta, n, e_s, e_t) \in \llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}}$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{S-E0})$$

And it suffices to prove that  $({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$

Again from Definition 2.95 it means we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1}$$

This means that given some  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} s\theta$ . Also given some  $j < n, {}^s v_1$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1} \quad (\text{S-E1}) \end{aligned}$$

Instantiating (S-E0) with  $H_{s1}, H_{t1}$  and with  $j, {}^s v_1$ . Then we get

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Since we have  $\tau <: \tau'$ . Therefore from IH (Statement 2(a)) we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

□

**Theorem 2.106** (FG  $\rightsquigarrow$  CG: Deriving FG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, n_1, n_2, H'_{s1}, H'_{s2}, pc$ .

Let  $\text{bool} = (\text{unit} + \text{unit})$

$$\begin{aligned} & \emptyset, \emptyset, x : \text{bool}^\top \vdash_{pc} e_s : \text{bool}^\perp \wedge \\ & \emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_1 : \text{bool}^\top \wedge \emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_2 : \text{bool}^\top \wedge \\ & (\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^s v'_1) \wedge \\ & (\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^s v'_2) \wedge \\ & \implies \\ & {}^s v'_1 = {}^s v'_2 \end{aligned}$$

*Proof.* From the FG to CG translation we know that  $\exists e_t$  s.t

$$\emptyset, \emptyset, x : \text{bool}^\top \vdash e_s : \text{bool}^\perp \rightsquigarrow e_t$$

Similarly we also know that  $\exists {}^t v_1, {}^t v_2$  s.t

$$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow {}^t v_1 \text{ and } \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{bool}^\top \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$$

From type preservation theorem (choosing  $\alpha = \gamma = \bar{\beta} = \perp$ ) we know that

$$\emptyset, \emptyset, x : \text{Labeled} \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{ Labeled} \perp \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash {}^t v_1 : \mathbb{C} \perp \perp \text{ Labeled} \top \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash {}^t v_2 : \mathbb{C} \perp \perp \text{ Labeled} \top \text{ bool} \quad (\text{NI-1})$$

Since we have  $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 2.104 we have (we choose  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\text{bool}^\top]_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{11}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\begin{aligned} & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ & (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{11}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'} \quad (\text{NI-2.1}) \end{aligned}$$

From Definition 2.94 we know that

$${}^t v_{11} = \mathbf{Lb}({}^t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i11}) \in [(\mathbf{unit} + \mathbf{unit}) \sigma]_V^{\hat{\beta}'}$$

Again from Definition 2.94 we know that

$$\begin{aligned} & \text{Either a) } {}^s v_1 = \mathbf{inl}() \text{ and } {}^t v_{i11} = \mathbf{inl}() \text{ or b) } {}^s v_1 = \mathbf{inr}() \text{ and } {}^t v_{i11} = \mathbf{inr}() \\ & \text{But in either case we have that } \emptyset, \emptyset, \emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.2}) \end{aligned}$$

$$\text{As a result we have } \emptyset, \emptyset, \emptyset \vdash {}^t v_{11} : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.3})$$

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit})} \quad (\text{NI-2.2})}{\emptyset, \emptyset, \emptyset \vdash \mathbf{Lb}({}^t v_{i11}) : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit})}$$

From Definition 2.99 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_{11})) \in [x \mapsto \mathbf{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 2.104 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_{11}/x]) \in [\mathbf{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 2.95 we get

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v''_1. (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v''_1) \implies \\ & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta'' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge ({}^s \theta'', n - i, {}^s v''_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset, n_1, {}^s v'_1$  we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - n_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge ({}^s \theta'', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \quad (\text{NI-2.5}) \end{aligned}$$

Since we have  $({}^s \theta'', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''}$  therefore from Definition 2.94 we have

$$\exists {}^t v_{i1}. {}^t v'' = \mathbf{Lb}({}^t v_{i1}) \wedge ({}^s \theta'', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [\mathbf{bool} \sigma]_V^{\hat{\beta}''}$$

Since  $({}^s \theta'', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$  therefore from Definition 2.94 two cases arise

- ${}^s v'_1 = \mathbf{inl} {}^s v_{i11}$  and  ${}^t v_{i1} = \mathbf{inl} {}^t v_{i11}$ :

From Definition 2.94 we have

$$({}^s \theta'', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v'_1 = \text{inr } {}^s v_{i11}$  and  ${}^t v_{i1} = \text{inr } {}^t v_{i11}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v'_1 = {}^t v_{i1}$

Similarly with other substitution we have  $(\emptyset, n, {}^s v_2, {}^t v_2) \in [\text{bool}^\top]_E^\emptyset$  (NI-3)

Therefore from Definition 2.95 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in [\text{bool}^\top \sigma]_{\hat{V}}^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [\text{bool}^\top \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{NI-3.1})$$

From Definition 2.94 we know that

$${}^t v_2 = \text{Lb}({}^t v_{i22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i22}) \in [(\text{unit} + \text{unit}) \sigma]_{\hat{V}}^{\hat{\beta}'}$$

Again from Definition 2.94 we know that

Either a)  ${}^s v_2 = \text{inl}()$  and  ${}^t v_{i22} = \text{inl}()$  or b)  ${}^s v_2 = \text{inr}()$  and  ${}^t v_{i22} = \text{inr}()$   
But in either case we have that  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})$  (NI-3.2)

As a result we have  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : \text{Labeled } \top (\text{unit} + \text{unit})$  (NI-3.3)

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})} \quad (\text{NI-3.2})}{\emptyset, \emptyset, \emptyset \vdash \text{Lb}({}^t v_{i22}) : \text{Labeled } \top (\text{unit} + \text{unit})}$$

From Definition 2.99 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \text{bool}^\top]_{\hat{V}}^{\hat{\beta}'}$$

Therefore we can apply Theorem 2.104 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\text{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-3.4})$$

From Definition 2.95 we get

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} \emptyset \wedge \forall i < n, {}^s v''_2. (H_s, e_s[{}^s v_2/x]) \Downarrow_i (H'_{s2}, {}^s v''_2) \implies \\ \exists H'_{t2}, {}^t v''_2. (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s2}, H'_{t2}) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v''_2, {}^t v''_2) \in [\text{bool}^\perp \sigma]_{\hat{V}}^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_2, {}^s v''_2$  we get

$$\exists H'_{t2}, {}^t v''_2. (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ (n - n_1, H'_s, H'_{t2}) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v''_2, {}^t v''_2) \in [\text{bool}^\perp \sigma]_{\hat{V}}^{\hat{\beta}''} \quad (\text{NI-3.5})$$

Since we have  $({}^s \theta', n - n_2, {}^s v''_2, {}^t v''_2) \in [\text{bool}^\perp \sigma]_{\hat{V}}^{\hat{\beta}''}$  therefore from Definition 2.94 we have

$$\exists {}^t v_{i2}. {}^t v''_2 = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v''_2, {}^t v_{i2}) \in [\text{bool} \sigma]_{\hat{V}}^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_2, {}^s v''_2, {}^t v_{i2}) \in [(\text{unit} + \text{unit})]_{\hat{V}}^{\hat{\beta}''}$  therefore from Definition 2.94 two cases arise

- ${}^s v'_2 = \text{inl } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inl } {}^t v_{i22}$ :

From Definition 2.94 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v'_1 = \text{inr } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inr } {}^t v_{i22}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v'_2 = {}^t v_{i2}$

We know that  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i1} : \text{Labeled } \top \text{ bool}$  (NI-2.3)

Also we have  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i2} : \text{Labeled } \top \text{ bool}$  (NI-3.3)

Let  $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that  $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{C} \perp \perp \text{ bool}$  by giving a typing derivation P2:

$$\frac{\frac{}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool}}{\text{CG-var}}}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{C} \perp \perp \text{ bool}} \text{CG-unlabel}$$

P1:

$$\frac{}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \text{ bool}} \text{From (NI-1)}$$

Main derivation:

$$\frac{P1 \quad P2}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{C} \perp \perp \text{ bool}}$$

Say  $e_t[{}^t v_{i1}/x]$  reduces in  $n_{t1}$  steps in (NI-2.5) and  $e_t[{}^t v_{i2}/x]$  reduces in  $n_{t2}$  steps in (NI-3.5)

We instantiate Theorem 2.57 with  $e_T, {}^t v_{i1}, {}^t v_{i2}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$  and from (NI-2.5) and (NI-3.5) we have  ${}^t v_{i1} = {}^t v_{i2}$  and thus  ${}^s v'_1 = {}^s v'_2$

□