

# On the Expressiveness and Semantics of Information Flow Types (Technical appendix)

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# 1 Fine-grained IFC enforcement (FG)

## 1.1 FG type system

**Syntax, types, constraints:**

Expressions	$e$	$::=$	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, x.e) \mid \text{new } e \mid !e \mid e := e \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet$
Labels	$\ell, pc$	$::=$	$l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
(Labeled) Types	$\tau$	$::=$	$A^\ell$
Unlabeled types	$A$	$::=$	$\mathbf{b} \mid \tau \xrightarrow{\ell_e} \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \tau \mid \text{unit} \mid \forall \alpha. (\ell_e, \tau) \mid c \xrightarrow{\ell_e} \tau$
Constraints	$c$	$::=$	$\ell \sqsubseteq \ell \mid (c, c)$

**Lemma 1.1** (FG: Reflexivity of subtyping). *The following hold:*

1. For all  $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all  $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$ .

Proof of statement (1)

Let  $\tau = A^\ell$ . Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on  $A$ .

1.  $A = \mathbf{b}$ :

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2.  $A = \text{ref } \tau$ :

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}$$

4.  $A = \tau_1 + \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}$$

**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var} \qquad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp} \text{FG-lam} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG-app} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp} \text{FG-prod} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG-fst} \quad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp} \text{FG-inl} \\
\\
\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau} \text{FG-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG-sub} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp} \text{FG-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}} \text{FG-assign} \\
\\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^\perp} \text{FG-unitI} \qquad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^\perp} \text{FG-FI} \\
\\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp} \text{FG-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG-CE}
\end{array}$$

Figure 1: Type system for FG

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A^{\ell'}} \text{FGsub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2} \text{FGsub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_\xi} \tau_1 <: c_2 \xrightarrow{\ell'_\xi} \tau_2} \text{FGsub-constraint}
\end{array}$$

Figure 2: FG subtyping

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash A \text{ WF} \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi \vdash A^\ell \text{ WF}} \text{FG-wff-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 \text{ WF}} \text{FG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 \text{ WF}} \text{FG-wff-prod} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 + \tau_2 \text{ WF}} \text{FG-wff-sum} \\
\\
\frac{\text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \tau) \text{ WF}} \text{FG-wff-ref} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma \cup \{\alpha\}}{\Sigma; \Psi \vdash (\forall \alpha. (\ell_e, \tau)) \text{ WF}} \text{FG-wff-forall} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ WF} \quad \text{FV}(c) \in \Sigma \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_\xi} \tau) \text{ WF}} \text{FG-wff-constraint}
\end{array}$$

Figure 3: Well-formedness relation for FG

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6.  $A = \text{unit}$ :

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$

7.  $A = \forall \alpha. \tau_i$ :

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}}$$

8.  $A = c \Rightarrow \tau_i$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash c \Rightarrow c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}}$$

□

## 1.2 FG semantics

Judgement:  $(H, e) \Downarrow_i (H', v)$

The semantics are described in Figure 4

## 1.3 Model for FG

$W : ((Loc \mapsto Type) \times (Loc \mapsto Type) \times (Loc \leftrightarrow Loc))$

**Definition 1.2** (FG:  $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 1.3** (FG:  $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

$$\begin{array}{c}
\frac{(H, e_1) \Downarrow_i (H', \lambda x. e_i) \quad (H', e_2) \Downarrow_j (H'', v_2) \quad (H'', e_i[v_2/x]) \Downarrow_k (H''', v_3)}{(H, e_1 e_2) \Downarrow_{i+j+k+1} (H''', v_3)} \text{fg-app} \\
\\
\frac{(H, e_1) \Downarrow_i (H', v_1) \quad (H', e_2) \Downarrow_j (H'', v_2)}{(H, (e_1, e_2)) \Downarrow_{i+j+1} (H'', (v_1, v_2))} \text{fg-prod} \quad \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{fst}(e)) \Downarrow_{i+1} (H', v_1)} \text{fg-fst} \\
\\
\frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{snd}(e)) \Downarrow_{i+1} (H', v_2)} \text{fg-snd} \quad \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inl}(e)) \Downarrow_{i+1} (H', \text{inl}(v))} \text{fg-inl} \\
\\
\frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inr}(e)) \Downarrow_{i+1} (H', \text{inr}(v))} \text{fg-inr} \quad \frac{(H, e) \Downarrow_i (H', \text{inl } v) \quad (H', e_1[v/x]) \Downarrow_j (H'', v_1)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_1)} \text{fg-case1} \\
\\
\frac{(H, e) \Downarrow_i (H', \text{inr } v) \quad (H', e_2[v/x]) \Downarrow_j (H'', v_2)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_2)} \text{fg-case2} \\
\\
\frac{(H, e) \Downarrow_i (H', \Lambda e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e[]) \Downarrow_{i+j+1} (H'', v)} \text{fg-FE} \\
\\
\frac{(H, e) \Downarrow_i (H', \nu e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e\bullet) \Downarrow_{i+j+1} (H'', v)} \text{fg-CE} \\
\\
\frac{(H, e) \Downarrow_i (H', v) \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1} (H'[a \mapsto v], a)} \text{fg-ref} \quad \frac{(H, e) \Downarrow_i (H', a)}{(H, !e) \Downarrow_{i+1} (H', H(a))} \text{fg-deref} \\
\\
\frac{(H, e_1) \Downarrow_i (H', a) \quad (H', e_2) \Downarrow_j (H'', v)}{(H, e_1 := e_2) \Downarrow_{i+j+1} (H''[a \mapsto v], ())} \text{fg-assign} \quad \frac{e \in \{x, \lambda y. -, \Lambda -, \nu -\}}{(H, e) \Downarrow_0 (H, e)} \text{fg-val}
\end{array}$$

Figure 4: FG semantics

**Definition 1.4** (FG: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{b} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})\} \\
[\forall \alpha. (\ell_e, \tau)]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. \\
&\quad ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})\} \\
[c \xrightarrow{\ell_e} \tau]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n. \\
&\quad \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V^A \triangleq \begin{cases} \{(W, n, v_1, v_2) \mid (W, n, v_1, v_2) \in [\mathbf{A}]_V^A\} & \ell' \sqsubseteq \mathbf{A} \\ \{(W, n, v_1, v_2) \mid \forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in [\mathbf{A}]_V\} & \ell' \not\sqsubseteq \mathbf{A} \end{cases}$$

**Definition 1.5** (FG: Binary expression relation).

$$\begin{aligned}
[\tau]_E^A &\triangleq \{(W, n, e_1, e_2) \mid \\
&\quad \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge \\
&\quad (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\
&\quad \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_V^A\}
\end{aligned}$$

**Definition 1.6** (FG: Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V\} \\
[\tau_1 + \tau_2]_V &\triangleq \{(\theta, m, \text{inl } v) \mid (\theta, m, v) \in [\tau_1]_V\} \cup \{(\theta, m, \text{inr } v) \mid (\theta, m, v) \in [\tau_2]_V\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\
&\quad (\theta', j, e[v/x]) \in [\tau_2]_E^{\ell_e}\} \\
[\forall \alpha. (\ell_e, \tau)]_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \forall \ell' \in \mathcal{L}. (\theta', m', e) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}\} \\
[c \xrightarrow{\ell_e} \tau]_V &\triangleq \{(\theta, m, \nu e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \tau\}
\end{aligned}$$



$$[A^{\ell'}]_V \triangleq [A]_V$$

**Definition 1.7** (FG: Unary expression relation).

$$\begin{aligned} [\tau]_E^{pc} \triangleq & \{(\theta, n, e) \mid \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)\} \end{aligned}$$

**Definition 1.8** (FG: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$$

**Definition 1.9** (FG: Binary heap well formedness).

$$\begin{aligned} (n, H_1, H_2) \hat{\triangleright}^A W \triangleq & \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ & (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ & \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ & (W, n - 1, H_1(a_1), H_2(a_2)) \in [W.\theta_1(a_1)]_V^A \wedge \\ & \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in [W.\theta_i(a_i)]_V \end{aligned}$$

**Definition 1.10** (FG: Label substitution).  $\sigma : Lvar \mapsto Label$

**Definition 1.11** (FG: Value substitution to value pairs).  $\gamma : Var \mapsto (Val, Val)$

**Definition 1.12** (FG: Value substitution to values).  $\delta : Var \mapsto Val$

**Definition 1.13** (FG: Unary interpretation of  $\Gamma$ ).

$$[\Gamma]_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V\}$$

**Definition 1.14** (FG: Binary interpretation of  $\Gamma$ ).

$$[\Gamma]_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A\}$$

## 1.4 Soundness proof for FG

**Lemma 1.15** (FG: Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n.$

*The following holds:*

1.  $\forall \mathcal{A}.$

$$(W, n, v_1, v_2) \in [A]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [A]_V$$

2.  $\forall \tau.$

$$(W, n, v_1, v_2) \in [\tau]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $A$  in the last step:

1. Case **b**:

From Definition 1.6

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 1.6, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 1.6, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A$  (S0)

IH1:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH2:  $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some  $m$  and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

Also from (S02) we know that given some  $m$  and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

We instantiate IH1 with  $m$  from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

We instantiate IH2 with  $m$  from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{j2})$

Symmetric case as (a)

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in \lceil \tau_1 \xrightarrow{\ell_e} \tau_2 \rceil_V^A$$

This means from Definition 1.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \lceil \tau_1 \rceil_V^A &\implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in \lceil \tau_2 \rceil_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in \lfloor \tau_1 \rfloor_V &\implies (\theta_l, i, e_1[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in \lfloor \tau_1 \rfloor_V &\implies (\theta_l, k, e_2[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}) \end{aligned} \quad (\text{L0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \lambda x.e_1) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V$ :

This means from Definition 1.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

This further means that we have some  $\theta', j$  and  $v$  s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in \lfloor \tau_1 \rfloor_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

(b)  $\forall m. (W.\theta_2, m, \lambda x.e_2) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V$ :

Similar reasoning with  $e_2$

5. Case  $\forall \alpha. (\ell_e, \tau)$ :

$$\text{Given: } (W, n, \Lambda e_1, \Lambda e_2) \in \lceil \forall \alpha. (\ell_e, \tau) \rceil_V^A$$

This means from Definition 1.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \quad (\text{F0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

This means from Definition 1.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

This further means that we are given some  $\theta'$ ,  $m'$  and  $\ell_u$  s.t  $W.\theta_1 \sqsubseteq \theta'$ ,  $m' < m$  and  $\ell_u \in \mathcal{L}$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

Instantiating  $\theta_l, i$  and  $\ell''$  in the second conjunct of F0 with  $\theta', m'$  and  $\ell_u$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\ell_u \in \mathcal{L}$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

$$(b) \forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

Symmetric reasoning for  $e_2$

6. Case  $c \xrightarrow{\ell_e} \tau$ :

$$\text{Given: } (W, n, \nu e_1, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V^A$$

This means from Definition 1.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{C0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \nu e_1) \in [c \xrightarrow{\ell_e} \tau]_V:$$

This means from Definition 1.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

This further means that we are given some  $\theta'$  and  $m'$  s.t  $W.\theta_1 \sqsubseteq \theta'$ ,  $m' < m$  and  $\mathcal{L} \models c$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

Instantiating  $\theta_l, j$  in the second conjunct of C0 with  $\theta', m'$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\mathcal{L} \models c$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

$$(b) \forall m. (W.\theta_2, m, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V:$$

Symmetric reasoning for  $e_2$

7. Case ref  $\tau$ :

From Definition 1.4 and 1.6

Proof of statement (2)

Let  $\tau = A^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement(1))

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 1.4

□

**Lemma 1.16** (FG: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m'$ .

1.  $\forall \mathbf{A}. (\theta, m, v) \in \llbracket \mathbf{A} \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \mathbf{A} \rrbracket_V$

2.  $\forall \tau. (\theta, m, v) \in \llbracket \tau \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \tau \rrbracket_V$

*Proof.* Proof by simultaneous induction on  $\mathbf{A}$  and  $\tau$

Proof of statement (1)

We analyze the various cases of  $\mathbf{A}$  in the last step:

1. case  $\mathbf{b}$ :

Directly from Definition 1.6

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V$

IH1 :  $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

IH2 :  $(\theta', m', v_2) \in \llbracket \tau_2 \rrbracket_V$

We get the desired from IH1, IH2 and Definition 1.6

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V$

IH :  $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

Therefore from IH and Definition 1.6 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(\theta, m, (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

To prove:  $(\theta', m', (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

This means from Definition 1.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e} \quad (1)$$

Similarly from Definition 1.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating Equation 75 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

5. case ref  $\tau$ :

From Definition 1.6 and Definition 1.2

6. case  $\forall \alpha. (\ell_e, \tau)$ :

Given:  $(\theta, m, (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

To prove:  $(\theta', m', (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

This means from Definition 1.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E^{\ell_e[\ell_i/\alpha]} \quad (2)$$

Similarly from Definition 1.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

Instantiating Equation 2 with  $\theta''', k$  and  $\ell_j$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\ell_j \in \mathcal{L}$

Therefore we get  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

7. case  $c \xrightarrow{\ell_e} \tau$ :

Given:  $(\theta, m, (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

To prove:  $(\theta', m', (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

This means from Definition 1.6 we know that

$$\forall \theta'' . \theta \sqsubseteq \theta'' \wedge \forall j < m . \mathcal{L} \models c \implies (\theta'' , j, e_1) \in [\tau]_E^{\ell_e} \quad (3)$$

Similarly from Definition 1.6 we know that we are required to prove

$$\forall \theta''' . \theta' \sqsubseteq \theta''' \wedge \forall k < m' . \mathcal{L} \models c \implies (\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$$

This means that given some  $\theta''' , k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating Equation 3 with  $\theta''' , k$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Proof of statement (2)

Let  $\tau = A^\ell$

Since  $[A^\ell]_V = [A]_V$ , therefore from IH (statement 1) □

**Lemma 1.17** (FG: Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, \mathcal{A}, n, n'$ .

1.  $\forall A. (W, n, v_1, v_2) \in [A]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [A]_V^A$
2.  $\forall \tau. (W, n, v_1, v_2) \in [\tau]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_V^A$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We analyze the different cases of  $A$  in the last step:

1. Case  $b$ :

From Definition 1.4

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 :  $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 :  $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 1.4 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 1.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i2})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

This means from Definition 1.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})$  (BM-A2)

Similarly from Definition 1.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e})$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A2 with  $\theta'_l, k$  and  $v'_c$  we get

$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$



5. Case ref  $\tau$ :

From Definition 1.4 and Definition 1.3

6. Case  $\forall\alpha.(\ell_e, \tau)$ :

Given:  $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

To prove:  $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

This means from Definition 1.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F2})$$

Similarly from Definition 1.4 we know that we are required to prove

$$(a) \quad \forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A):$$

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with  $W'', n''$  and  $\ell''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . Also since  $n'' < n'$  and  $n' < n$  therefore  $n'' < n$ . And finally since  $\ell'' \in \mathcal{L}$  therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

$$(b) \quad \forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_l \sqsupseteq W.\theta_1$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

$$(c) \quad \forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some  $\theta_l \sqsupseteq W.\theta_2, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F2 with  $\theta_l, k$  and  $\ell''$ . And since  $\theta_l \sqsupseteq W.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta_l \sqsupseteq W.\theta_2$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

7. Case  $c \stackrel{\ell_e}{\Rightarrow} \tau$ :

Given:  $(W, n, (\nu e_1), (\nu e_2)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

To prove:  $(\theta', n', (\nu e_1), (\nu e_1)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

This means from Definition 1.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \quad (\text{BM-C0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \quad (\text{BM-C1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{BM-C2})$$

Similarly from Definition 1.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in [\tau]_E^A$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\mathcal{L} \models c$

And we are required to prove:  $(W'', n'', e_1, e_2) \in [\tau]_E^A$

Instantiating BM-C0 with  $W'', n''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . And since  $\mathcal{L} \models c$  therefore we get

$(W'', n'', e_1, e_2) \in [\tau]_E^A$

(b)  $\forall \theta'_i \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_1, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with  $\theta'_i, k$ . And since  $\theta'_i \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_i \sqsupseteq W.\theta_1$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

(c)  $\forall \theta'_i \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$ :

This means that we are given some  $\theta'_i \sqsupseteq W'.\theta_2, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with  $\theta'_i, k$ . And since  $\theta'_i \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_i \sqsupseteq W.\theta_2$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

#### Proof of statement (2)

Let  $\tau = A^\ell$

2 cases arise:

1.  $\ell \sqsubseteq \mathcal{A}$ :

From IH (statement 1)

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

From Lemma 1.16 and Definition 1.4

□

**Lemma 1.18** (FG: Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .

$(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

*Proof.* Given:  $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove:  $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 1.13 it is given that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$

And again from Definition 1.13 we are required to prove that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$

•  $dom(\Gamma) \subseteq dom(\delta)$ :

Given

- $\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$ :

Since we know that  $\forall x \in \text{dom}(\Gamma).(\theta, n, \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$  (given)

Therefore from Lemma 1.16 we get

$$\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$$

□

**Lemma 1.19** (FG: Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .

$$(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \lfloor \Gamma \rfloor_V$$

*Proof.* Given:  $(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(W', n', \gamma) \in \lfloor \Gamma \rfloor_V$

From Definition 1.14 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$$

And again from Definition 1.13 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma)$ :

Given

- $\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$ :

Since we know that  $\forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$  (given)

Therefore from Lemma 1.17 we get

$$\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$$

□

**Lemma 1.20** (FG: Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .

$$(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 1.8 it is given that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

And again from Definition 1.13 we are required to prove that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$ :

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :

Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)

Therefore from Lemma 1.16 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

**Lemma 1.21** (FG: Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .  
 $(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 1.9 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall(a_1, a_2) \in (W.\hat{\beta}). &(W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n-1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 1.9 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :  
 Given
- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :  
 Given
- $\forall(a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
 $\forall(a_1, a_2) \in (W.\hat{\beta})$ .
  - $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given
  - $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :  
 Given and from Lemma 1.17
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :  
 Given

□

**Theorem 1.22** (FG: Fundamental theorem unary).  $\forall \Sigma, \Psi, \Gamma, pc, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n$ .

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \\ \mathcal{L} \models \Psi \sigma \wedge \\ (\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V \implies \\ (\theta, n, e \delta) \in \lfloor \tau \sigma \rfloor_E^{pc} \end{aligned}$$

*Proof.* Proof by induction on *FG* typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove:  $(\theta, n, x \delta) \in \lfloor \tau \sigma \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, x \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-V0})
\end{aligned}$$

In order to prove FU-V0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $v' = \delta(x)$  and  $j = 1$

We need to prove the following:

- (a)  $\theta \sqsubseteq \theta \wedge (n - 1, H) \triangleright \theta \wedge (\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$ :
- $\theta \sqsubseteq \theta$ : From Definition 1.2
  - $(n - 1, H) \triangleright \theta$ : From Lemma 1.20
  - $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$ :  
Since we are given that  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$  and  $v' = \delta(x)$   
Therefore  $(\theta, n, v') \in \lfloor \Gamma(x) \sigma \rfloor_V$ , where  $\Gamma(x) = \tau$   
And finally from Lemma 1.16 we get  $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$
- (b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :  
Since  $H' = H$ , so we are done
- (c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :  
Since  $\theta' = \theta$ , so we are done

## 2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(\theta, \lambda x. e_i \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, (\lambda x. e_i) \delta) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $j < n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-L0})
\end{aligned}$$

IH1:

$\forall \theta_i, v_x, n. (\theta_i, n, e_i \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$ , s.t  $(\theta_i, n, v_x) \in [\tau_1 \sigma]_V$

In order to prove FU-L0 we instantiate  $\theta'$  with  $\theta$ . From reduction relation we know that  $H' = H$ ,  $j = 0$  and  $v' = \lambda x. e_i \delta$

(a)  $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$ :

- $\theta \sqsubseteq \theta$ : From Definition 1.2
- $(n, H) \triangleright \theta$ : Given
- $(\theta, n, (\lambda x. e_i) \delta) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$ :

From Definition 1.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \forall v. (\theta'', j, v) \in [\tau_1 \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

This means given some  $\theta''$ ,  $j$  and  $v$  such that  $\theta \sqsubseteq \theta''$ ,  $j < n$  and  $(\theta'', j, v) \in [\tau_1 \sigma]_V$

It suffices to prove that  $(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  and  $j < n$  therefore from Lemma 1.18 we have

$$(\theta, j, \delta) \in [\Gamma \sigma]_V$$

So we can apply IH1 instantiated with  $\theta''$ ,  $v$  and  $j$  we get

$$(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

(b)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $H' = H$  so we are done

(c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$ :

Since  $\theta' = \theta$  so we are done

### 3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-P0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall n_1, H_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma)
\end{aligned}$$

Instantiating IH1 with  $n, H$  and since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 \ e_2) \ \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma) \quad (\text{FU-P1})
\end{aligned}$$

From evaluation rule we know that  $v'_1 = \lambda x. e_i$ . Since from FU-P1 we know that

$$(\theta'_1, n - i, \lambda x. e_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \ \sigma]_V$$

This means from Definition 1.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < (n - i). \forall v. (\theta'', j, v) \in [\tau_1 \ \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \ \sigma]_E^{\ell_e \ \sigma} \quad (4)$$

IH2:

$$\begin{aligned}
& \forall n_2, \forall H_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall k < n_2. (H_2, (e_2) \ \delta) \Downarrow_k (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - k, v'_2) \in [(\tau_1) \ \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma)
\end{aligned}$$

Instantiating IH2 with  $n - i, H'_1$  and since we know that  $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1 \ e_2) \ \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - k, v'_2) \in [(\tau_1) \ \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma) \quad (\text{FU-P2})
\end{aligned}$$

Instantiating  $\theta'', j$  and  $v$  in Equation 4 with  $\theta'_2, n - i - k$  and  $v'_2$  from FU-P2 respectively, we get

$$(\theta'_2, n - i - k, e_i[v'_2/x]) \in [\tau_2 \ \sigma]_E^{\ell_e \ \sigma}$$

This means from Definition 1.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - k, H_3) \triangleright \theta'_2 \wedge \forall l < (n - i - k). (H_3, e_i[v'_2/x]) \Downarrow_l (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \ \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \ \sigma)
\end{aligned}$$

Instantiating  $H_3$  with  $H'_2$  from FU-P2 and since we know that  $((n - i - k), H'_2) \triangleright \theta'_2$  and since the reduction happens therefore we have

$$\begin{aligned}
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \ \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \ \sigma) \quad (\text{FU-P3})
\end{aligned}$$

In order to prove FU-P0 we choose  $\theta'$  as  $\theta'_3$  from FU-P3. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + l$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$ :
- $\theta \sqsubseteq \theta'_3$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-P1,  $\theta'_1 \sqsubseteq \theta'_2$  from FU-P2 and  $\theta'_2 \sqsubseteq \theta'_3$  from FU-P3 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_3$
  - $((n - i - k - l), H'_3) \triangleright \theta'_3$ :  
From FU-P3 we get  $((n - i - k - l), H'_3) \triangleright \theta'_3$
  - $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$ :  
From FU-P3 we get  $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-P1, FU-P2 and FU-P3
- (c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta). \theta'_3(a) \searrow pc \sigma)$   
Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-P1, FU-P2 and FU-P3

#### 4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  s.t  $H \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-PA0)} \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

We instantiate IH1 with  $H$  and  $n$ . And since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-PA1)} \end{aligned}$$

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau_2) \sigma \rrbracket_V \wedge \end{aligned}$$



$$\begin{aligned}
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma)
\end{aligned}$$

We instantiate IH2 with  $H'_1$  and  $n - i$ . And since we know that  $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma) \quad \text{(FU-PA2)}
\end{aligned}$$

In order to prove FU-PA0 we choose  $\theta'$  as  $\theta'_2$  from FU-PA2. Also we know from the evaluation rule, that let  $v' = (v'_1, v'_2)$ ,  $H' = H'_2$  and  $n' = i + j + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v') \in [(\tau_1 \times \tau_2)^\perp]_V$ :
- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-PA1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-PA2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$
  - $(n - i - j - 1, H'_2) \triangleright \theta'_2$ :  
From FU-PA2 we get  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.20 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$
  - $(\theta'_2, n - i - j, v') \in [(\tau_1 \times \tau_2)^\perp]_V$ :  
From Definition 1.6 it suffices to show
    - i.  $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$ :  
Since from FU-PA1 we know that  $(\theta'_1, n - i, v'_1) \in [(\tau_1) \ \sigma]_V$  and since  $\theta'_1 \sqsubseteq \theta'_2$  (from FU-PA2) therefore from Lemma 1.16 we get  $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$
    - ii.  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$ :  
From FU-PA2 we know that  $(\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V$  therefore from Lemma 1.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$   
From FU-PA1 and FU-PA2
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma)$   
From FU-PA1 and FU-PA2

5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove:  $(\theta, n, \text{fst}(e_i) \delta) \in [\tau_1 \ \sigma]_E^{pc \ \sigma}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \ \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma)
\end{aligned}$$

This means that given some heap  $H$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-F0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-F1}) \end{aligned}$$

From evaluation rule we know that  $v'_1 = (v''_1, v''_2)$

In order to prove FU-F0 we choose  $\theta'$  as  $\theta'_1$  from FU-F1. Also we know that  $H' = H'_1$  and  $v' = v''_1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v'_1) \in [\tau_1 \sigma]_V$ :
  - $\theta \sqsubseteq \theta'_1$ :  
From FU-F1
  - $(n - i - 1, H'_1) \triangleright \theta'_1$ :  
From FU-F1 we know  $(n - i, H'_1) \triangleright \theta'_1$  therefore from Lemma 1.20 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$
  - $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$ :  
Since from FU-F1 we know that  $(\theta'_1, n - i, (v''_1, v''_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$   
Therefore from Definition 1.6 we know that  $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$   
From Lemma 1.16 we get  $(\theta'_1, n - i - 1, v''_1) \in [\tau_1 \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
From FU-F1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$   
From FU-F1

6. FG-snd:

Symmetric case to FG-fst

7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(\theta, n, \text{inl}(e_i) \delta) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^{pc \sigma}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-LE0)}
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-LE1)}
\end{aligned}$$

In order to prove FU-LE0 we choose  $\theta'$  as  $\theta'_1$  from FU-LE1. Also we know from the evaluation rule, that let  $v' = \text{inl}(v'_1)$ ,  $H' = H'_1$  and  $n' = i + 1$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V:$$

- $\theta \sqsubseteq \theta'_1$ :

From FU-LE1

- $(n - i - 1, H') \triangleright \theta'_1$ :

From FU-LE1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 1.20 we get  $(n - i - 1, H') \triangleright \theta'_1$

- $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$ :

Since  $v' = \text{inl}(v'_1)$  and from FU-LE1 we know that  $(\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V$

Therefore from Definition 1.6 we get  $(\theta'_1, n - i, v') \in [(\tau_1 + \tau_2) \sigma]_V$

From Lemma 1.16 we get  $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$

$$(b) (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$$

From FU-LE1

$$(c) (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$$

From FU-LE1

8. FG-inr:

Symmetric case to FG-inl

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau}$$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t.  $(n, H) \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-C0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_c) \delta) \Downarrow_i (H'_1, v'_c) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $H \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-C1}) \end{aligned}$$

2 cases arise:

(a)  $v'_c = \text{inl}(v_{ci})$ :

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_1) \delta \cup \{x \mapsto v_{ci}\}) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Instantiating IH2 with  $H'_1$  and  $n-i$  since we know that  $H'_1 \triangleright \theta'_1 \wedge (H'_1, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{FU-C2}) \end{aligned}$$

In order to prove FU-C0 we choose  $\theta'$  as  $\theta'_2$  from FU-C2. Also we know that  $H' = H'_2$ ,  $v' = v'_2$  and  $n' = i + j + 1$ . Now we are required to show

- i.  $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$ :
- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-C1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-C2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$
  - $(n - i - j - 1, H'_2) \triangleright \theta'_2$ :  
From FU-C2 we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.20 we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$
  - $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$ :  
From FU-C2 we know that  $(\theta'_2, n - i - j, v'_2) \in \lfloor \tau \sigma \rfloor_V$  therefore from Lemma 1.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$
- ii.  $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$ :  
Since from FU-C2 we know that  
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell')$   
therefore we also have  
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$   
and from FU-C1 we know that  
 $(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$   
Combining the two we get  
 $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
- iii.  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$ :  
Since from FU-C2 we know that  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$   
therefore we also have  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc) \sigma)$   
and from FU-C1 we know that  
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma)$   
Combining the two we get  
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$
- (b)  $v'_c = \text{inr}(v_{ci})$ :  
Symmetric case as  $v'_c = \text{inl}(v_{ci})$

10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(\theta, n, \text{new } (e_i) \delta) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\text{ref } \tau)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\text{ref } \tau)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma) \quad (\text{FU-R0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \ \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \ \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma) \quad (\text{FU-R1})
\end{aligned}$$

From the evaluation rule we know that  $H' = H'_1[a \mapsto v'_1]$  where  $a \notin \text{dom}(H'_1)$ ,  $v' = a$  and  $n' = i + 1$ . In order to prove FU-R0 we choose  $\theta'$  as  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$ . Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_2 \wedge (n - i - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - 1, v') \in [(\text{ref } \tau)^\perp \ \sigma]_V:$$

- $\theta \sqsubseteq \theta'_2$ :

From FU-R1 we know that  $\theta \sqsubseteq \theta'_1$  therefore from Definition 1.2  $\theta \sqsubseteq \theta'_2$

- $(n - i - 1, H') \triangleright \theta'_2$ :

From FU-R1 we know that  $(n - i, H'_1) \triangleright \theta'_1$ . Therefore from Lemma 1.20 we get  $(n - i - 1, H'_1) \triangleright \theta'_1$ .

We also know that  $(\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V$  (from FU-R1) therefore from Lemma 1.16 we get  $(\theta'_1, n - i - 1, v'_1) \in [\tau \ \sigma]_V$

Since  $H' = H'_1[a \mapsto v'_1]$  and  $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$  therefore from Definition 1.8 we get  $(n - i - 1, H') \triangleright \theta'_2$

- $(\theta'_2, n - i - 1, a) \in [(\text{ref } \tau)^\perp \ \sigma]_V$ :

Since  $\theta'_2(a) = \tau \ \sigma$  therefore from Definition 1.6 we get  $(\theta'_2, n - i - 1, a) \in [(\text{ref } \tau)^\perp \ \sigma]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$$

From FU-R1

$$(c) \ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma):$$

We get this from FU-R1 and  $\tau \ \sigma \searrow pc \ \sigma$  (given)

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(\theta, n, (!e_i) \ \delta) \in [\tau' \ \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (!e_i) \ \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \ \sigma]_V \wedge
\end{aligned}$$

$$\begin{aligned}
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t.  $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-D0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-D1})
\end{aligned}$$

In order to prove FU-D0 we choose  $\theta'$  as  $\theta'_1$  from FU-D1. Also we know from the evaluation rule, that  $H' = H'_1$ ,  $v' = H'_1(a)$ ,  $v'_1 = a$  and  $n' = i + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$ :
- $\theta \sqsubseteq \theta'_1$ :  
From FU-D1
  - $(n - i - 1, H') \triangleright \theta'_1$ :  
From FU-D1 we know that  $(n - i, H') \triangleright \theta'_1$  therefore from Lemma 1.20 we get  $(n - i - 1, H') \triangleright \theta'_1$
  - $(\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$ :  
Since from FU-D1 we know that  $(n - i, H'_1) \triangleright \theta'_1$  therefore from the Definition 1.8 we get  $(\theta'_1, n - i, H'_1(a)) \in [\tau \sigma]_V$   
From Lemma 1.16 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau \sigma]_V$   
Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 1.24 we get  $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
From FU-D1
- (c)  $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$   
From FU-D1

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}}$$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [\text{unit } \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\mathbf{ref} \ \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\mathbf{ref} \ \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-A1})
\end{aligned}$$

IH2:

$$\begin{aligned}
& \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq (n_2 - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with  $H'_1$  and since we know that  $H'_1 \triangleright \theta'_1 \wedge (H, (e_1 := e_2) \delta) \Downarrow (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq (n - i - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-A2})
\end{aligned}$$

In order to prove FU-A0 we choose  $\theta'$  as  $\theta'_2$  from FU-A2. Also we know from the evaluation rule, assign, that let  $v'_1 = a_1$ ,  $H' = H'_2[a_1 \mapsto v'_2]$ ,  $v' = ()$  and  $n' = i + j + 1$ . Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, ()) \in [\mathbf{unit}]_V:$$

- $\theta \sqsubseteq \theta'_2$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-A1 and  $\theta'_1 \sqsubseteq \theta'_2$  from FU-A2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_2$
- $(n - i - j - 1, H') \triangleright \theta'_2$ :  
From Definition 1.8 it suffices to prove that
  - i.  $\text{dom}(\theta'_2) \subseteq \text{dom}(H')$ : From FU-A2
  - ii.  $\forall a \in \text{dom}(\theta'_2). (\theta'_2, n - i - j - 1, H'(a)) \in [\theta'_2(a)]_V$ :  
 $\forall a \in \text{dom}(\theta'_2).$



- $a = a_1$ :  
From FU-A2 (since we know that  $(\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$ )  
Therefore from Lemma 1.16 we get  $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$
  - $a \neq a_1$ :  
From FU-A2 (since we know that  $(n - i - j, H'_2) \triangleright \theta'_2$  therefore from Lemma 1.20  
we get  $(n - i - j - 1, H'_2) \triangleright \theta'_2$ )
  - $(\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$ :  
From Definition 1.6
- (b)  $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$   
 $\forall a \in \text{dom}(H)$ .
- $a = a_1$ :  
Since we know that  $H(a_1) \neq H'(a_1)$  and  $\theta(a_1) = \tau = \mathbf{A}^{\ell_i}$  (given)  
It is given that  $\tau \sigma \searrow pc \sigma$  therefore  $pc \sigma \sqsubseteq \ell_i \sigma$
  - $a \neq a_1$ :  
From FU-A2
- (c)  $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$   
From FU-A2

### 13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove:  $(\theta, n, (\Lambda e_i) \delta) \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\Lambda e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\Lambda e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-FI0)} \end{aligned}$$

IH1:

$$\forall n_1, \theta_i, \ell' \in \mathcal{L}. (\theta_i, n_1, e_i \delta) \in \llbracket \tau \sigma \cup \{\alpha \mapsto \ell''\} \rrbracket_E^{\ell_e \sigma \cup \{\alpha \mapsto \ell''\}}$$

In order to prove FU-FI0 we choose  $\theta'$  as  $\theta$ . Also we know from the evaluation rule, that  $H' = H$  and  $n' = 0$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V$ :
- $\theta \sqsubseteq \theta$ : From Definition 1.2
  - $(n, H) \triangleright \theta$ : Given

- $(\theta, n, (\Lambda e_i)\delta) \in [(\forall\alpha.(\ell_e, \tau))^\perp]_V \sigma$ :

From Definition 1.6 it suffices to prove that

$$\forall\theta''.\theta \sqsubseteq \theta'' \wedge \forall j < n.\forall\ell_d \in \mathcal{L} \implies (\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$$

This means given some  $\theta'', j$  and  $\ell_d$  such that  $\theta \sqsubseteq \theta'', j < n$  and  $\ell_d \in \mathcal{L}$

It suffices to prove that  $(\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$

Instantiating IH1 with  $j, \theta''$  and  $\ell_d$  we get  $(\theta_i, j, e_i \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell_d\}]_E^{\ell_e} \sigma \cup \{\alpha \mapsto \ell_d\}$

- (b)  $(\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $H' = H$  so we are done

- (c)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta(a) \searrow pc)$ :

Since  $\theta' = \theta$  so we are done

#### 14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\forall\alpha.(\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i [] : \tau[\ell''/\alpha]}$$

To prove:  $(\theta, n, (e_i[]) \delta) \in [\tau[\ell''/\alpha] \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n.(n, H) \triangleright \theta \wedge \forall n' < n.(H, (e_i[]) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \quad (\text{FU-FE0}) \end{aligned}$$

IH:

$$\begin{aligned} & \forall H_1, n_1.(n_1, H_1) \triangleright \theta \wedge \forall i < n_1.(H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH with  $H$  and  $n$ . Since we know that  $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned} & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \quad (\text{FU-FE1}) \end{aligned}$$

From evaluation rule we know that  $v'_1 = \Lambda e_{i1}$ . Since from FU-FE1 we know that

$$(\theta'_1, n - i, \Lambda e_{i1}) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V$$

This means from Definition 1.6 we have

$$\forall \theta'' . \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i . \forall \ell_g \in \mathcal{L} \implies (\theta'' , j , e_{i1}) \in [\tau[\ell_g/\alpha] \sigma]_E^{\ell_e[\ell_g/\alpha]} \sigma \quad (5)$$

Instantiating Equation 5 with  $\theta'_1$ ,  $n - i - 1$  and  $\ell''$  we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha]} \sigma$$

This means from Definition 1.7 we have

$$\begin{aligned} & \forall H_3 . (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1 . (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\ & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \end{aligned}$$

Instantiating  $H_3$  with  $H'_1$  from FU-FE1 and since we know that  $(n - i - 1, H'_1) \triangleright \theta'_1$  (Lemma 1.20) and since we know that  $e_i \Downarrow \gamma \downarrow_1$  reduces in  $n'$  steps where  $n' = i + k + 1$  and since  $n' < n$  therefore we have  $k < n - i - 1$  s.t.  $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$ . Therefore we get

$$\begin{aligned} & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-FE2}) \end{aligned}$$

In order to prove FU-FE0 we choose  $\theta'$  as  $\theta'_3$  from FU-FE2. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + 1$ . Now we are required to show

- (a)  $\theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$ :
- $\theta \sqsubseteq \theta'_3$ :  
Since  $\theta \sqsubseteq \theta'_1$  from FU-FE1 and  $\theta'_1 \sqsubseteq \theta'_3$  from FU-FE2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_3$
  - $(n - i - k - 1, H'_3) \triangleright \theta'_3$ :  
From FU-FE2 we know that  $(n - i - k - 1, H'_3) \triangleright \theta'_3$
  - $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$ :  
From FU-FE2 we know that  $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$
- (b)  $(\forall a \in \text{dom}(H) . H(a) \neq H'_3(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$   
Since  $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  therefore we get the desired from FU-FE1 and FU-FE2
- (c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta) . \theta'_3(a) \searrow pc \sigma)$   
Since  $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  therefore we get the desired from FU-FE1 and FU-FE2

## 15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e_i : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove:  $(\theta, n, (\nu e_i) \delta) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n . (n, H) \triangleright \theta \wedge \forall n' < n . (H, (\nu e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta' . \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V \wedge \\ & (\forall a . H(a) \neq H'(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta) . \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (\nu e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (c \xrightarrow{\ell_e} \tau)^\perp \sigma \rrbracket_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CI0}) \end{aligned}$$

IH1:

$$\forall \theta_i, n_1. (\theta_i, n_1, e_i \delta) \in \llbracket \tau \sigma \rrbracket_E^{\ell_e} \sigma \text{ such that } \mathcal{L} \models c \sigma$$

In order to prove FU-FI0 we choose  $\theta'$  as  $\theta$ . Also we know from the evaluation rule, that  $H' = H$ ,  $v' = \nu e_i \delta$  and  $n' = 0$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \llbracket (c \xrightarrow{\ell_e} \tau)^\perp \rrbracket_V \sigma:$$

•  $\theta \sqsubseteq \theta$ : From Definition 1.2

•  $(n, H) \triangleright \theta$ : Given

•  $(\theta, n, (\nu e_i) \delta) \in \llbracket (c \xrightarrow{\ell_e} \tau)^\perp \rrbracket_V \sigma$ :

From Definition 1.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \mathcal{L} \models c \sigma \implies (\theta'', j, e_i \delta) \in \llbracket \tau \sigma \rrbracket_E^{\ell_e} \sigma$$

This means given some  $\theta''$  such that  $\theta \sqsubseteq \theta''$ ,  $j < n$  and  $\mathcal{L} \models c$

It suffices to prove that  $(\theta'', j, e_i \delta) \in \llbracket \tau \sigma \rrbracket_E^{\ell_e} \sigma$

Instantiating IH1 with  $\theta''$  and  $j$  we get  $(\theta'', j, e_i \delta) \in \llbracket \tau \sigma \rrbracket_E^{\ell_e} \sigma$

$$(b) (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell'):$$

Since  $H' = H$  so we are done

$$(c) (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc):$$

Since  $\theta' = \theta$  so we are done

16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i \bullet : \tau}$$

To prove:  $(\theta, n, (e_i \bullet) \delta) \in \llbracket \tau \sigma \rrbracket_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v') \implies \\ \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap  $H$  and  $n$  s.t  $(n, H) \triangleright \theta \wedge (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CE0}) \end{aligned}$$

IH:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH with  $H$  and  $n$ . And since we know that  $(n, H) \triangleright \theta \wedge (H, (e_i) \delta) \Downarrow_{n'} (H', v')$  therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-CE1})
\end{aligned}$$

From evaluation rule we know that  $v'_1 = \nu e_{i1}$ . Since from FU-CE1 we know that

$$(\theta'_1, n - i, \nu e_{i1}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V$$

This means from Definition 1.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i. \mathcal{L} \models c \sigma \implies (\theta'', j, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma} \quad (6)$$

Instantiating Equation 6 with  $\theta'_1$  and  $n - i - 1$  since we know that  $\mathcal{L} \models c \sigma$  therefore we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

This means from Definition 1.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1. (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma)
\end{aligned}$$

Instantiating  $H_3$  with  $H'_1$  from FU-CE1 and since we know that  $(n - i - 1, H'_1) \triangleright \theta'_1$  (Lemma 1.20) and since we know that  $e_i \bullet \gamma \Downarrow_1$  reduces in  $n'$  steps where  $n' = i + k + 1$  and since  $n' < n$  therefore we have  $k < n - i - 1$  s.t.  $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$ . Therefore we get

$$\begin{aligned}
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-CE2})
\end{aligned}$$

In order to prove FU-CE0 we choose  $\theta'$  as  $\theta'_3$  from FU-CE2. Also we know that  $H' = H'_3$ ,  $v' = v'_3$  and  $n' = i + k + 1$ . Now we are required to show

$$(a) \theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V:$$

- $\theta \sqsubseteq \theta'_3$ :

Since  $\theta \sqsubseteq \theta'_1$  from FU-CE1 and  $\theta'_1 \sqsubseteq \theta'_3$  from FU-CE2 therefore from Definition 1.2 we get  $\theta \sqsubseteq \theta'_3$

- $(n - i - k - 1, H'_3) \triangleright \theta'_3$ :

From FU-CE3 we know that  $(n - i - k - 1, H'_3) \triangleright \theta'_3$

- $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$ :

From FU-CE3 we know that  $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$

$$(b) (\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$$

Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-CE1 and FU-CE2

(c)  $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta)).\theta'_3(a) \searrow pc \sigma$

Since  $pc \sigma \sqsubseteq \ell_e \sigma$  therefore we get the desired from FU-CE1 and FU-CE2

□

**Lemma 1.23** (FG: Expression subtyping with closed labels and types).  $\forall pc, pc', \tau$ .

$$\mathcal{L} \models pc \sqsubseteq pc' \implies \llbracket \tau \rrbracket_E^{pc'} \subseteq \llbracket \tau \rrbracket_E^{pc}$$

*Proof.* Given:  $\mathcal{L} \models pc \sqsubseteq pc'$

$$\text{To prove: } \llbracket (\tau) \rrbracket_E^{pc'} \subseteq \llbracket (\tau) \rrbracket_E^{pc}$$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc'} . (\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}$$

This means given  $\forall (\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc'}$

It suffices to prove that  $(\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}$

From Definition 1.7 for the chosen  $\theta, n, e$  we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc') \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall k < n.(H_1, e) \Downarrow_k (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket \tau \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some  $H_1$  and  $k$  such that  $(n, H_1) \triangleright \theta$ ,  $k < n$  and  $(H_1, e) \Downarrow_k (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket \tau \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate  $H$  in (A) with  $H_1$  and then we choose  $\theta'_1$  as  $\theta'$

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \llbracket \tau \rrbracket_V$ :

Given

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :

Since  $pc \sqsubseteq pc'$  and we are given

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell')$$

Therefore

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$ :

We are given

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc')$$

and since  $pc \sqsubseteq pc'$  Therefore

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$$

□

**Lemma 1.24** (FG: Subtyping unary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1.  $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (A \sigma) \rrbracket_V \subseteq \llbracket (A' \sigma) \rrbracket_V$$

2.  $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_V \subseteq \llbracket (\tau' \sigma) \rrbracket_V$$

$$(b) \forall pc. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{pc}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove:  $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

IH1:  $\llbracket (\tau'_1 \sigma) \rrbracket_V \subseteq \llbracket (\tau_1 \sigma) \rrbracket_V$  (Statement 2(a))

IH2:  $\forall pc. \llbracket (\tau_2 \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_E^{pc}$  (Statement 2(b))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V$

Therefore from Definition 1.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e \sigma} \quad (7)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V$

Again from Definition 1.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e \sigma}$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e \sigma}$

Since  $(\theta_2, j, v) \in \llbracket \tau'_1 \sigma \rrbracket_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in \llbracket \tau_1 \sigma \rrbracket_V$

As a result from Equation 7 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e \sigma}$$

From IH2, we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell_e \sigma}$$

Since  $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 1.23 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $[\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V \subseteq [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

IH1:  $[\lceil\tau_1 \sigma\rceil]_V \subseteq [\lceil\tau'_1 \sigma\rceil]_V$  (Statement 2(a))

IH2:  $[\lceil\tau_2 \sigma\rceil]_V \subseteq [\lceil\tau'_2 \sigma\rceil]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, (v_1, v_2)) \in [\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V. (\theta, n, (v_1, v_2)) \in [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [\lceil(\tau_1 \times \tau_2) \sigma\rceil]_V$

Therefore from Definition 1.6 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \quad (8)$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [\lceil(\tau'_1 \times \tau'_2) \sigma\rceil]_V$

Again from Definition 1.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 8 we know that  $(\theta, n, v_1) \in [\tau_1 \sigma]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2 \sigma]_V$  from Equation 8 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

## 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[\lceil(\tau_1 + \tau_2) \sigma\rceil]_V \subseteq [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

IH1:  $[\lceil\tau_1 \sigma\rceil]_V \subseteq [\lceil\tau'_1 \sigma\rceil]_V$  (Statement 2(a))

IH2:  $[\lceil\tau_2 \sigma\rceil]_V \subseteq [\lceil\tau'_2 \sigma\rceil]_V$  (Statement 2(a))

It suffices to prove:  $\forall(\theta, n, v_s) \in [\lceil(\tau_1 + \tau_2) \sigma\rceil]_V. (\theta, v_s) \in [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

This means that given:  $(\theta, n, v_s) \in [\lceil(\tau_1 + \tau_2) \sigma\rceil]_V$

And it suffices to prove:  $(\theta, n, v_s) \in [\lceil(\tau'_1 + \tau'_2) \sigma\rceil]_V$

2 cases arise



(a)  $v_s = \text{inl } v_i$ :

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_1 \sigma \rfloor_V \quad (9)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

From Equation 9 and IH1 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_2 \sigma \rfloor_V \quad (10)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

From Equation 10 and IH2 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

#### 4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $\lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V \subseteq \lfloor (\forall \alpha. (\ell'_e, \tau_2)) \sigma \rfloor_V$

IH1:  $\forall pc. \lfloor (\tau_1 \sigma) \rfloor_E^{pc} \subseteq \lfloor (\tau_2 \sigma) \rfloor_E^{pc}$  (Statement 2(b))

It suffices to prove:  $\forall (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V. (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

This means that given:  $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V$

Therefore from Definition 1.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])} \quad (11)$$

And it suffices to prove:  $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

Again from Definition 1.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

This means that given some  $\theta_2, j < n, \ell' \in \mathcal{L}$  s.t  $\theta \sqsubseteq \theta_2$

And we are required to prove:  $(\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$  therefore from Equation 11 we have

$$(\theta_2, j, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

Since  $\mathcal{L} \models \ell'_e (\sigma \cup [\alpha \mapsto \ell']) \sqsubseteq \ell_e (\sigma \cup [\alpha \mapsto \ell'])$  therefore from Lemma 1.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell'_e (\sigma \cup [\alpha \mapsto \ell'])}$$

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

$$\text{To prove: } [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V \subseteq [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{IH1: } \forall pc. [(\tau_1 \sigma)]_E^{pc} \subseteq [(\tau_2 \sigma)]_E^{pc} \text{ (Statement 2(b))}$$

$$\text{It suffices to prove: } \forall (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V. (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{This means that given: } (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V$$

Therefore from Definition 1.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma} \quad (12)$$

$$\text{And it suffices to prove: } (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

Again from Definition 1.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

This means that given some  $\theta_2, j$  s.t  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

$$\text{And we are required to prove: } (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$  therefore from Equation 12 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell_e \sigma}$$

Since  $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 1.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

$$\text{To prove: } [((\text{ref } \tau) \sigma)]_V \subseteq [((\text{ref } \tau) \sigma)]_V$$

$$\text{It suffices to prove: } \forall (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V. (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V$$

Trivial

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{b}) \sigma) \rrbracket_V$

Directly from Definition 1.6

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit}) \sigma) \rrbracket_V \subseteq \llbracket ((\text{unit}) \sigma) \rrbracket_V$

Directly from Definition 1.6

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^\ell <: \mathbf{A}^{\ell'}} \text{FGsub-label}$$

To prove:  $\llbracket ((\mathbf{A}^\ell) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}^{\ell'}) \sigma) \rrbracket_V$

From Definition 1.6 it suffices to prove:  $\llbracket ((\mathbf{A}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}') \sigma) \rrbracket_V$

This we get directly from IH (Statement 1(a))

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\llbracket (\tau \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}. (\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$$

This means given  $(\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}$

It suffices to prove that  $(\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

From Definition 1.7 we know we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall i < n.(H, e) \Downarrow_i (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - i, H') \triangleright \theta' \wedge (\theta', n - i, v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall j < n.(H_1, e) \Downarrow_j (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some  $H_1$  and  $j < n$  s.t  $(n, H_1) \triangleright \theta \wedge (H_1, e) \Downarrow_j (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n-j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n-j, v') \in [\tau' \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate  $H$  in (A) with  $H_1$  and  $i$  with  $j$  then we choose  $\theta'_1$  as  $\theta'$   
Also we have IH1 as  $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$  (Statement 2(a))

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$ :  
We are given  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau \sigma]_V$   
From IH1 we know that  $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$   
Therefore,  $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$
- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$ :  
Given
- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$ :  
Given

□

**Lemma 1.25** (FG: Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$   
 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 1.14 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case  $i = 1$

Given some  $m$  we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i)$ :  
 $\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i)$   
Therefore,  $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i))$  (Given)
- $\forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$ :  
We are given:  $\forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$   
Therefore from Lemma 1.15 we know that  
 $\forall m'. (W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$   
Instantiating  $m'$  with  $m$  we get  
 $(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 2$

Symmetric case as  $i = 1$

□

**Theorem 1.26** (FG: Fundamental theorem binary).  $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}} &\implies \\ (W, n, e (\gamma \downarrow_1), e (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}} & \end{aligned}$$

*Proof.* Proof by induction on the typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove:  $(W, n, x (\gamma \downarrow_1), x (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}}$

Say  $e_1 = x (\gamma \downarrow_1)$  and  $e_2 = x (\gamma \downarrow_2)$

From Definition of  $[\tau]_E^{\mathcal{A}}$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge \forall j < n. (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) &\implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} & \end{aligned}$$

This means given some  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove:  $\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Here

- $H'_1 = H_1$  and  $H'_2 = H_2$
- $e_1 = v'_1 = \gamma(x) \downarrow_1$
- $e_2 = v'_2 = \gamma(x) \downarrow_2$
- $j = 1$

We choose  $W' = W$ .

- $W \sqsubseteq W$ : From Definition 1.3

- $(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$ :

Since we know that  $(n, H_1, H_2) \triangleright^{\mathcal{A}} W$  therefore from Lemma 1.21 we get

$$(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$$

- $(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$ :

We are given that  $(W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$  therefore from Lemma 1.19 we get

$$(W, n - 1, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$$

which means from Definition 1.14 we have

$$(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$$

2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove:  $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_E^A$

Say  $e_1 = \lambda x. e (\gamma \downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \end{aligned}$$

This means that given  $H_1, H_2$  and  $j$  s.t  $(n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

It suffices to prove:

$$\exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \quad (\text{FB-L0})$$

IH1:

$$\forall W, n. (W, n, e (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (v_1, v_2)) \in [\tau_1 \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \lambda x. e (\gamma \downarrow_1)$ ,  $v'_2 = e_2 = \lambda x. e (\gamma \downarrow_2)$  and  $j = 0$ . In order to prove FB-L0 we choose  $W' = W$  and we need to prove the following:

- $W \sqsubseteq W$ : From Definition 1.3
- $(n, H_1, H_2) \triangleright^A W$ : Given
- $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A$

From Definition 1.4 it suffices to prove that:

$$\forall W'' \sqsupseteq W, k < n, v_1, v_2.$$

$$((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, k, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$

This means that we need to prove the following:

$$\begin{aligned} - \forall W'' \sqsupseteq W, k < n, v_1, v_2. ((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A): \end{aligned}$$

This means given  $W'' \sqsupseteq W, k < n, v_1, v_2$  s.t  $((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove:  $(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

We instantiate IH1 with  $W''$  and  $k$   
 And since  $(W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$  therefore we get  
 $(W'', k, e(\gamma \downarrow_1)[v_1/x], e(\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

–  $\forall \theta_l \sqsupseteq W.\theta_1, k, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$   
 $(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$ :

This means that we are given  $\theta_l, k$  and  $v_c$  s.t

$\theta_l \sqsupseteq W.\theta_1$  and  $(\theta_l, k, v_c) \in [\tau_1 \sigma]_V$

And we are required to prove:

$(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

It is given to us that

$\forall v_1, v_2. (W, n, \gamma \in [\Gamma]_V^A$

Therefore from Lemma 1.25 we know that

$\forall m. (W.\theta_1, m, (\gamma \downarrow_1) \in [\Gamma]_V$

Therefore, we can apply Theorem 1.22 to obtain

$\forall m. (W.\theta_1, m, \lambda x. e \gamma \downarrow_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V$

From Definition 1.6 it means that we have

$\forall m. \forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1 \sigma]_V \implies$

$(\theta', j, e[v/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

We instantiate  $m$  with some  $l > k$ ,  $\theta'$  with  $\theta_l$ ,  $j$  with  $k$  and  $v$  with  $v_c$  to get

$W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since we show that  $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V$  therefore we get

$(\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

–  $\forall \theta_l \sqsupseteq W.\theta_2, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$

$(\theta_l, k, e(\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$ :

Symmetric case as above

### 3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2)(\gamma \downarrow_1), (e_1 e_2)(\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 1.5 we need to prove:

$\forall H_1, H_2, n' < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow$   
 $(H'_2, v'_2) \implies$

$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$

This further means that given  $H_1, H_2, n' < n$  s.t

$(n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_{\mathcal{V}}^A \quad (\text{FB-A0})$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}, i < n. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge (H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$ .  $(H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A \quad (13)$$

$$\underline{\text{IH2}}: (W'_1, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{j1}, H_{j2}, j < (n - i). (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge (H_{j1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{j1}$  with  $H'_1$  and  $H_{j2}$  with  $H'_2$  in IH2. Since the  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps. Also,  $e_1$  reduces to value  $\gamma \downarrow_1$  in  $i < n'$  steps. Therefore  $\exists j < n' - i < n - i$  s.t  $(H_{i1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1})$ .  $(H_{i2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A \quad (14)$$

We case analyze on  $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$  from Equation 13

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_{\mathcal{V}}^A$$

$$\text{Let } v'_1 = \lambda x.e_{h1} \text{ and } v'_2 = \lambda x.e_{h2}$$

Again from Definition 1.4 it means that

$$\forall W'_{h1} \sqsupseteq W'_1, j_1 < (n - i), v_1, v_2.$$

$$((W'_{h1}, j_1, v_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}}^A \implies (W'_{h1}, j_1, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_1, m_1, v_c.$$

$$\wedge ((\theta_{l1}, m_1, v_1) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_1, e_{h1}[v_1/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_2, m_1, v_c.$$

$$\wedge (\theta_{l1}, m_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_2, e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$



We instantiate  $W'_{h1}$  with  $W'_2$  obtained from Equation 14. Similarly we also instantiate  $v_1$  and  $v_2$  with  $v'_{j1}$  and  $v'_{j2}$  respectively from Equation 14, and  $j_1$  with  $n - i - j$ . And we get

$$(W'_2, n - i - j, e_{h1}[v'_{j1}/x], e_{h2}[v'_{j2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 1.5 we get

$$\begin{aligned} & \forall H_1, H_2, k_e < (n - i - j). (n - i - j, H_1, H_2) \triangleright^A W'_2 \wedge \\ & (H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1}) \wedge (H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2}) \implies \\ & \exists W' \sqsupseteq W'_2. (n - i - j - k_e, H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', n - i - j - k_e, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \end{aligned}$$

Instantiating  $H_1$  with  $H'_{j1}$  and  $H_2$  with  $H'_{j2}$  obtained from Equation 14. And since we know that  $e_1 e_2$  reduces with  $\gamma \downarrow_1$  in  $n' < n$  steps. And  $e_2$  reduces to value  $\gamma \downarrow_1$  in  $j < n' - 1 < n - i$  steps. Therefore  $\exists k_e = n' - i - j < n - i - j$  s.t  $(H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1})$ .  $(H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2})$  is known because  $(e_1 e_2)$  reduces to value with  $\gamma \downarrow_2$ . Hence we get

$$\exists W' \sqsupseteq W'_2. ((n - i - j - k_e), H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', (n - i - j - k_e), v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \quad (15)$$

This concludes the proof in this case.

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From FB-A0 we know that we need to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau_2 \sigma = A^{\ell_i}$  and since  $\tau_2 \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V \quad (16)$$

In this case from Definition 1.6 we know that

$$\forall m. (W'_1.\theta_1, m, \lambda x. e_{h1}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (17)$$

$$\forall m. (W'_1.\theta_2, m, \lambda x. e_{h2}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (18)$$

Applying Definition 1.6 on Equation 17 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

where  $\theta = W'_1.\theta_1$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 1 + t_1). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC1})$$

Since from Equation 14 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 1.15 we get

$$\forall m. (W'_2.\theta_1, m, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(W'_2.\theta_1, m_1 + 1 + t_1, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_1$ ,  $j_1$  with  $m_1 + t_1$  and  $v$  with  $v'_{j_1}$  from Equation 14.

$$\text{Therefore we get } (W'_2.\theta_1, m_1 + 1 + t_1, e_{h1}[v'_{j_1}/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

From Definition 1.7, we get

$$\begin{aligned} \forall H. (m_1 + 1 + t_1, H) \triangleright W'_2.\theta_1 \wedge \forall k_c < (m_1 + 1 + t_1). (H, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1 + t_1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1 + t_1 - k_c), v'_1) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 14 we have  $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 1.27 we get  $\forall m. (m, H'_{j_1}) \triangleright W'_2.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get  $(m_1 + 1 + t_1, H'_{j_1}) \triangleright W'_2.\theta_1$

Now instantiating  $H$  with  $H'_{j_1}$  from Equation 14 and  $k_c$  with  $t_1$  we get

$$\begin{aligned} \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{R1}) \end{aligned}$$

Similarly we can apply Definition 1.6 on Equation 18 to get

$$\forall m. \forall \theta'_2. (m, W'_1.\theta_2) \sqsubseteq \theta'_2 \wedge \forall j_2 < m. \forall v. (\theta'_2, j_2, v) \in [\tau_1 \sigma]_V \implies (\theta'_2, j_2, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

We instantiate  $m$  with  $m_2 + 2 + t_2$  where  $t_2$  is the number of steps in which  $e_{h2}$  reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies (\theta', j_1, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC2})$$

Since from Equation 14 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 1.15 we get

$$\forall m. (W'_2.\theta_2, m, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating  $\theta'$  with  $W'_2.\theta_2$ ,  $j_1$  with  $m_2 + 1 + t_2$  and  $v$  with  $v'_{j_2}$  from Equation 14 in FB-AC2 we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, e_{h_2}[v'_{j_2}/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

From Definition 1.7, we get

$$\begin{aligned} \forall H.(m_2 + 1 + t_2, H) \triangleright W'_2.\theta_2 \wedge \forall k_c < (m_2 + 1 + t_2).(H, e_{h_2}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_2, v'_2) \implies \\ \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1 + t_2 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1 + t_2 - k_c)v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 14 we have  $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 1.27 we get  $\forall m.(m, H'_{j_2}) \triangleright W'_2.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get  $(m_2 + 1 + t_2, H'_{j_2}) \triangleright W'_2.\theta_2$

Now Instantiating  $H$  with  $H'_{j_2}$  from Equation 14 and and  $k_c$  with  $t_2$ .

$$\begin{aligned} \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{R2}) \end{aligned}$$

In order to prove FB-A0 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_2.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

–  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$ :

From R1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from R2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$ :

Since from Equation 14 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$  therefore from

Definition 1.9 we know that  $(W'_2.\hat{\beta}) \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  therefore

$(W'_2.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$

$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

4 cases arise for each  $(a_1, a_2) \in W'_2.\hat{\beta}$

i.  $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

We know from Equation 14 that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$

Therefore from Definition 1.9 we have

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

Since  $W'.\hat{\beta} = W'_2.\hat{\beta}$  by construction therefore

$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

From R1 and R2 we know that  $W'_2.\theta_1 \sqsubseteq \theta'_1$  and  $W'_2.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From Equation 14 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$

This means from Definition 1.9 that

$$\forall (a_{i_1}, a_{i_2}) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_2 \sqsubseteq W'$  and  $n - n' - 1 < n - i - j - 1$  (since  $n' = i + j + t_1$  where  $t_1$  is the number of steps taken by  $e_{h_1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce and  $j$  is the number of steps taken by  $e_2 \gamma \downarrow_1$  to reduce) therefore from Lemma 1.17 we get

$$(W', n - n' - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii.  $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From R1 and R2 we know that

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_2.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_2.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from R1 and R2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ . Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V \text{ and}$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii.  $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From R2 we know that

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e \sigma$  in the world before the modification. Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 14 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$  that means from Definition 1.9 that  $(W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$ . Since  $(\ell_e \sigma) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H'_{j_1}(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_2.\theta_1, m, H'_{j_1}(a_1)) \in W'_2.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_2.\theta_2, m, H'_{j_2}(a_2)) \in W'_2.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.16 we get

$$(\theta'_1, m_1, H'_{j_1}(a_1)) \in \theta'_1(a_1)$$

Since from R2 we know that  $(m_2+1, H'_2) \triangleright \theta'_2$  therefore from Definition 1.8

$$\text{we know that } (\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V:$$

$$\underline{i = 1}$$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we instantiate Equation 17 and Equation 18 with  $m + 2 + t_1$  and  $m + 2 + t_2$  respectively. This will give us

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

and

$$\exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma))$$

Since we have  $(m+1, H'_1) \triangleright \theta'_1$  and  $(m+1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 1.8

$$\underline{i = 2}$$

Symmetric to  $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau_2 \sigma]_V^A:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From R1 and R2 we and Definition 1.4 we get the desired.

#### 4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove:  $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\perp \sigma]_E^A$

Say  $e_1 = (e_1, e_2) (\gamma \downarrow_1)$  and  $e_2 = (e_1, e_2) (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \times \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies$$

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$$

This means that given some  $H_1, H_2$  and  $n' < n$  s.t

$$(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{\mathcal{V}}^A \quad (19)$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{p11}, H_{p12}. (n, H_{p11}, H_{p12}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11}) \wedge (H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{p11}$  with  $H_1$  and  $H_{p22}$  with  $H_2$  in IH1 and since the  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12})$ . Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (20)$$

$$\underline{\text{IH2}} (W, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{p21}, H_{p22}. (n - i, H_{p21}, H_{p22}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22}) \implies$$

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{p21}$  with  $H'_{p11}$  and  $H_{p22}$  with  $H'_{p21}$  and in IH2. Since  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  has reduced with  $i < n'$  steps. Therefore we know that  $\exists j < n' - i < n - i$  s.t  $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21})$ . Similarly since we know that  $(e_1, e_2)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$ . Hence we get

since the  $(e_1, e_2)$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$ . Hence we get

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (21)$$

In order to prove Equation 19 we instantiate  $W'$  in Equation 19 with  $W'_2$  we are required to show the following:

- $W \sqsubseteq W'_2$ :  
Since  $W \sqsubseteq W'_1$  from Equation 20 and  $W'_1 \sqsubseteq W'_2$  from Equation 21  
Therefore,  $W \sqsubseteq W'_2$  from Definition 1.3

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'$ :

Here  $n' = i + j + 1$

From evaluation rule of products we know that  $H'_1 = H'_{p21}$  and  $H'_2 = H'_{p22}$

From Equation 21 we know that  $(n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.21 we get  $(n - i - j - 1, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

- $(W', n - i - j - 1, v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$ :

From evaluation rule of products we know that  $v'_1 = (v'_{p11}, v'_{p21})$  and  $v'_2 = (v'_{p12}, v'_{p22})$

We are required to show

$$- (W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

From Equation 20 and Equation 21 we know that

$$(W'_2, n - i - j, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

Therefore from Lemma 1.17 we get

$$(W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

## 5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove:  $(W, n, (\text{fst}(e_i)) (\gamma \downarrow_1), (\text{fst}(e_i)) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

Say  $e_1 = (\text{fst}(e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{fst}(e_i)) (\gamma \downarrow_2)$

From Definition 1.5 it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A \quad (22)$$

### IH1

$$(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{fst}(e_i)$  reduces to value reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{fst}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A \quad (23)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$  from Equation 23

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V^A$$

Let  $v'_{i1} = (v_{i1}, v_{i2})$  and  $v'_{i2} = (v_{j1}, v_{j2})$

Again from Definition 1.4 it means that

$$(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A \wedge (W'_1, n - i, v_{i2}, v_{j2}) \in [\tau_2 \sigma]_V^A \quad (\text{F1})$$

In order to prove Equation 22 we choose  $W'$  as  $W'_1$  and from the evaluation rule of fst we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . Also, from reduction rules we know that  $n' = i + 1$ . And then we need to show:

- $W \sqsubseteq W'_1$ :

Directly from Equation 23

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

Since from Equation 23 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 1.21 we get  $(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$

- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$ :

From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$

From F1 we know that  $(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 1.17 we get  $(W'_1, n - i - 1, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 1.6 we know that

$$(a) \forall m. (W'_1.\theta_1, m, v'_{i1}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V \text{ and}$$

$$(b) \forall m. (W'_1.\theta_2, m, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V$$

where

$$v'_{i1} = (v_{i1}, v_{i2}) \text{ and } v'_{i2} = (v_{j1}, v_{j2})$$

In order to prove Equation 22 we choose  $W'$  as  $W'_1$  and from the evaluation rule of fst we know that  $H'_1 = H'_{i1}$  and  $H'_2 = H'_{i2}$ . And then we need to show:

- $W \sqsubseteq W'_1$ :

Directly from Equation 23

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

From Equation 23 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 1.21 we get

$$(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$$



- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$ :  
 From the evaluation rule we know that  $v'_1 = v_{i1}$  and  $v'_2 = v_{j1}$   
 Let  $\tau_1 = \mathbf{A}^{\ell_i}$  Since  $\tau_1 \sigma \searrow \ell$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \quad (24)$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \quad (25)$$

Instantiating (a) with  $m_1$

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_1, m_1, v_{i2}) \in [\tau_2 \sigma]_V \quad (26)$$

Instantiating (b) with  $m_2$

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_2, m_2, v_{j2}) \in [\tau_2 \sigma]_V \quad (27)$$

From Equation 26 and Equation 27 we get

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \text{ and } (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

## 6. FG-snd:

Symmetric case as FG-fst

## 7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove:  $(W, n, (\text{inl}(e_i))(\gamma \downarrow_1), (\text{inl}(e_i))(\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^A$

Say  $e_1 = (\text{inl}(e_i))(\gamma \downarrow_1)$  and  $e_2 = (\text{inl}(e_i))(\gamma \downarrow_2)$

From Definition of  $[(\tau_1 + \tau_2)^\perp \sigma]_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A \quad (28)$$

IH1  $(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore we know that  $\exists i < n' < n$  s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since we know that  $\text{inl}(e_i)$  reduces to value with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \quad (29)$$

Instantiating  $W'$  in Equation 28 with  $W'_1$ . Also from reduction relation we know that  $n' = i + 1$  we are required to show the following:

- $W \sqsubseteq W'_1$ :

Directly from Equation 29

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ :

From Equation 29 we know that  $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 1.21 we get

$$(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$ :

From evaluation rule of  $\text{inl}$  we know that  $v'_1 = \text{inl}(v'_{i1})$  and  $v'_2 = \text{inl}(v'_{i2})$

We are required to show

$$- (W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A:$$

From Equation 29 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 1.17 we get

$$(W'_1, n - i - 1, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$$

8. FG-inr:

Symmetric case to FG-inl.

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{i1} : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_i, x.e_{i1}, y.e_{i2}) : \tau}$$

To prove:  $(W, (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1), (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

Say  $e_1 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1)$  and  $e_2 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ & \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A \quad (30)$$

$$\underline{\text{IH1}} \ (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value with both  $\gamma \downarrow_1$  and  $\gamma \downarrow_2$  therefore we know that  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^A \quad (31)$$

IH2:

$$(W'_1, n - i, (e_{i1}) (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\}), (e_{i1}) (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{j1}, H_{j2}. (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_1, e_{i1} (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\})) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \Downarrow (H'_{j2}, v'_{j2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A$$

Instantiating  $H_{j1}$  with  $H'_1$  and  $H_{j2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A \quad (32)$$

IH3:

$$(W'_1, n - i, (e_{i2}) (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\}), (e_{i2}) (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{k1}, H_{k2}. (n - i, H_{k1}, H_{k2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < n - i. (H_1, e_{i2} (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\})) \Downarrow_k (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \Downarrow (H'_{k2}, v'_{k2}) \implies$$

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [(\tau) \sigma]_V^A$$

Instantiating  $H_{k1}$  with  $H'_1$  and  $H_{k2}$  with  $H'_2$  in IH2. Also instantiating  $W$  with  $W'_1$ . Since the  $(\text{case}(e_i, x.e_{i2}, y.e_{i2}))$  reduces to value in both runs therefore we know that  $(H_1, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{k2}, v'_{k2})$ . Hence we get

$$\exists W'_3 \sqsubseteq W'_1.(n - i - k, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau \sigma]_V^A \quad (33)$$

We case analyze  $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^A$  from Equation 31

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 1.4 2 further cases arise:

- $v'_1 = \text{inl}(v_{i1})$  and  $v'_2 = \text{inl}(v_{i2})$ :

In this case from Definition 1.4 we know that  $(W, n - i, v_{i1}, v_{i2}) \in [\tau_1 \sigma]_V^A$

In order to prove Equation 30 we choose  $W'$  as  $W'_2$  from Equation 32 and from the first evaluation rule of case we know that  $H'_1 = H'_{j1}$  and  $H'_2 = H'_{j2}$ . Also we know from the evaluation rule that  $n' = i + j + 1$ . And then we need to show:

- \*  $W \sqsubseteq W'_2$ :

Since  $W \sqsubseteq W'_1$  from Equation 31 and  $W'_1 \sqsubseteq W'_2$  from Equation 32

Therefore,  $W \sqsubseteq W'_2$  from Definition 1.3

- \*  $(n - n', H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2$ :

From Equation 32 we know that  $(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.21 we get

$$(n - i - j - 1, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2$$

- \*  $(W'_2, n - n', v'_1, v'_2) \in [\tau \sigma]_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{j1}$  and  $v'_2 = v'_{j2}$

From Equation 32 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [\tau \sigma]_V^A$

Therefore from Lemma 1.17 we get

$$(W'_2, n - i - j - 1, v'_{j1}, v'_{j2}) \in [\tau \sigma]_V^A$$

- $v'_1 = \text{inr}(v_{i1})$  and  $v'_2 = \text{inr}(v_{i2})$ :

In this case from Definition 1.4 we know that  $(W, v_{i1}, v_{i2}) \in [\tau_2 \sigma]_V^A$

In order to prove Equation 30 we choose  $W'$  as  $W'_3$  from Equation 33 and from the second evaluation rule of case we know that  $H'_1 = H'_{k1}$  and  $H'_2 = H'_{k2}$ . Also we know from the evaluation rule that  $n' = i + k + 1$ . And then we need to show:

- \*  $W \sqsubseteq W'_3$ :

Since  $W \sqsubseteq W'_1$  from Equation 31 and  $W'_1 \sqsubseteq W'_3$  from Equation 33

Therefore,  $W \sqsubseteq W'_3$  from Definition 1.3

- \*  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_3$ :

From Equation 33 we know that  $(n - i - k, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3$

Therefore from Lemma 1.21 we get

$$(n - i - k - 1, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3$$

- \*  $(W'_3, n - n', v'_1, v'_2) \in [\tau \sigma]_V^A$ :

From the evaluation rule we know that  $v'_1 = v'_{k1}$  and  $v'_2 = v'_{k2}$

From Equation 33 we know that  $(W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau \sigma]_V^A$

Therefore from Lemma 1.17 we get

$$(W'_3, n - i - k - 1, v'_{k1}, v'_{k2}) \in [\tau \sigma]_V^A$$

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

The following cases arise:

- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case1 :  
Exactly the same reasoning as in the  $v'_1 = \text{inl}(v_{i1})$  and  $v'_2 = \text{inl}(v_{i2})$  subcase of the  $\ell \sigma \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case2 :  
Exactly the same reasoning as in the  $v'_1 = \text{inr}(v_{i1})$  and  $v'_2 = \text{inr}(v_{i2})$  subcase of the  $\ell \sigma \not\sqsubseteq \mathcal{A}$  case before.
- Reduction of  $e_1$  happens via Case1 and Reduction of  $e_2$  happens via Case2 :

From Equation 30 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau \sigma = A^{\ell_i}$  and since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$

From Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (34)$$

Since we know that  $(W, n, \gamma) \in [\Gamma]_V^A$  (given) therefore from Lemma 1.25 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore by instantiating it at  $m_1 + 1 + j$  we know that

$$(W.\theta_1, m_1 + 1 + j, \gamma \downarrow_1) \in [\Gamma]_V \quad (35)$$

Next we apply Theorem 1.22 on  $e_{i1} \gamma \downarrow_1$ . Here  $j$  is the number of steps in which  $e_{i1} \gamma \downarrow_1$  reduces. We use  $\gamma \downarrow_1 \cup \{x \mapsto v'_{s1}\}$  as the unary substitution to get

$$(W.\theta_1, m_1 + 1 + j, e_{i1} \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$$

This means from Definition 1.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_1 + 1 + j, H_{c1}) \triangleright W_1.\theta_1 \wedge \forall l_c < (m_1 + 1 + j). (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1 + j - l_c, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1 + j - l_c, v'_c) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W_1.\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since from Equation 31 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.27 we get  $\forall m. (m, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $m$  with  $m_1 + 1 + j$  we get  $(m_1 + 1 + j, H'_1) \triangleright W'_1.\theta_1$

Instantiating  $H_{c2}$  with  $H'_1$  from Equation 31 and  $l_c$  with  $j$  we get  
 $\exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC1})$

Similarly we apply Theorem 1.22 on  $e_{i2} \gamma \downarrow_2$ . Here  $j_2$  is the number of steps in which  $e_{i2} \gamma \downarrow_2$  reduces. We use  $\gamma \downarrow_2 \cup \{y \mapsto v'_{s2}\}$  as the unary substitution to get  
 $(W_1.\theta_2, m_2 + 1 + j_2, e_{i2} \gamma \downarrow_1 \cup \{y \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$

This means from Definition 1.7 we get

$\forall H_{c2}. (m_2 + 1 + j_2, H_{c1}) \triangleright W_1.\theta_2 \wedge \forall l_c < m_2 + 1 + j_2. (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \Downarrow_{k_c}$   
 $(H'_{c2}, v'_c) \implies$   
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1 + j_2 - l_c, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1 + j_2 - l_c, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_2). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$

Since from Equation 31 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.27 we get  $\forall m. (m, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $m$  with  $m_2 + 1 + j_2$  we get  $(m_2 + 1 + j_2, H'_2) \triangleright W'_1.\theta_2$

Instantiating  $H_{c2}$  with  $H'_2$  (from Equation 31) and  $l_c$  with  $j_2$  to get  
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$   
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_2). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC2})$

We choose

$W_n.\theta_1 = \theta'_1$  (from CC1)  
 $W_n.\theta_2 = \theta'_2$  (from CC2)  
 $W_n.\hat{\beta} = W'_1.\hat{\beta}$  (from Equation 31)

In order to prove Equation 30 we choose  $W'$  as  $W_n$

i.  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

–  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$ :

From (CC1) we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from (CC2) we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

–  $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$ :

Since from Equation 31 we have  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Definition 1.9 we get  $(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore  
 $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

–  $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$   
 $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A$ :

4 cases arise for each  $a_1$  and  $a_2$

- A.  $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$ :  
 $\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$ :

We know from Equation 31 that  $(n - i, H'_1, H'_2) \triangleright W'_1$

Therefore from Definition 1.9 we have

$$\forall(a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From (CC1) and (CC2) we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From Equation 31 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$

This means from Definition 1.9 that

$$\forall(a_{i_1}, a_{i_2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$  (since  $n' = i + t_1 + 1$  where  $t_1$  is the number of steps taken by  $e_{i_1}$ ,  $i$  is the number of steps taken by  $e_1 \gamma \downarrow_1$  to reduce) therefore from Lemma 1.17 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

- B.  $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From (CC1) and (CC2) we know that

$$(\forall a. H'_1(a) \neq H'_{c_1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H'_{c_2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from (CC1) and (CC2),  $(m_1 + 1, H'_{c_1}) \triangleright \theta'_1$  and  $(m_2 + 1, H'_{c_2}) \triangleright \theta'_2$ .

Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_{c_1}(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$

$$(\theta'_2, m_2, H'_{c_2}(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we get (here  $H'_1 = H'_{c_1}$  and  $H'_2 = H'_{c_2}$ )

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

- C.  $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$ :

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From (CC2) we know that

$$(\forall a. H_2'(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1'.\theta_2(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $(pc \sqcup \ell) \sigma$  in the world before the modification. Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 31 we know that  $(n - i, H_1', H_2') \stackrel{\mathcal{A}}{\triangleright} W_1'$  that means from Definition 1.9 that  $(W_1', n - i - 1, H_1'(a_1), H_2'(a_2)) \in \lceil W_1'.\theta_1(a_1) \rceil_V^{\mathcal{A}}$ . Since  $((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H_1'(a_1)$  must also be protected at some label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W_1'.\theta_1, m, H_1'(a_1)) \in W_1'.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W_1'.\theta_2, m, H_2'(a_2)) \in W_1'.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.16 we get

$$(\theta_1', m_1, H_1'(a_1)) \in \theta_1'(a_1)$$

Since from (CC2) we know that  $(m_2 + 1, H_{c2}') \triangleright \theta_2'$  therefore from Definition 1.8 we know that  $(\theta_2', m_2, H_{c2}'(a_2)) \in \theta_2'(a_2)$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H_{c1}'(a_1), H_{c2}'(a_2)) \in \lceil \theta_1'(a_1) \rceil_V^{\mathcal{A}}$$

$$\text{D. } H_{j1}'(a_1) \neq H_1'(a_1) \vee H_{j2}'(a_2) = H_2'(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V:$$

$$\underline{i = 1}$$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V$$

Like before we apply Theorem 1.22 on  $e_{i1} \gamma_1$  and  $e_{i2} \gamma_2$  but this time using  $m + 1 + i$  and  $m + 1 + j$  where  $i$  and  $j$  are the number of steps in which  $e_{i1} \gamma_1$  and  $e_{i2} \gamma_2$  reduces respectively. This will give us

$$\begin{aligned} & \exists \theta_1'. W_1.\theta_1 \sqsubseteq \theta_1' \wedge (m + 1, H_{c2}') \triangleright \theta_1' \wedge (\theta_1', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_1') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta_2'. W_1.\theta_2 \sqsubseteq \theta_2' \wedge (m + 1, H_{c2}') \triangleright \theta_2' \wedge (\theta_2', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_2') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since we have  $(m + 1, H_{c1}') \triangleright \theta_1'$  and  $(m + 1, H_{c2}') \triangleright \theta_2'$  therefore we get the desired from Definition 1.8

$$\underline{i = 2}$$

Symmetric to  $i = 1$

$$\text{ii. } (W', n - n' - 1, v'_1, v'_2) \in \lceil \tau_2 \sigma \rceil_V^{\mathcal{A}}:$$

Let  $\tau_2 = \mathbf{A}^{\ell_i}$  Since  $\tau_2 \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CC1 and CC2 we and Definition 1.4 we get the desired.

- (d) Reduction of  $e_1$  happens via Case2 and Reduction of  $e_2$  happens via Case1 :  
Symmetric case as before



10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow_{pc}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove:  $(W, (\text{new } (e_i)) (\gamma \downarrow_1), (\text{new } (e_i)) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\perp \sigma]_E^A$

Say  $e_1 = (\text{new } (e_i)) (\gamma \downarrow_1)$  and  $e_2 = (\text{new } (e_i)) (\gamma \downarrow_2)$

From Definition of  $[(\text{ref } \tau)^\perp \sigma]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \quad (36)$$

IH1  $(W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow \\ (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $\text{ref}(e_i)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$ . s.t  $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since  $\text{ref}(e_i)$  reduces with  $\gamma \downarrow_2$  therefore we know that  $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \quad (37)$$

From the evaluation rule of  $\text{ref}$  we know that  $H'_1 = H'_{i1} \cup \{a_{n1} \mapsto v_{i1}\}$  and  $H'_2 = H'_{i2} \cup \{a_{n2} \mapsto v_{i2}\}$

Inorder to prove Equation 36 we instantiate  $W'$  with  $W_n$  where  $W_n$  is

$$W_n.\theta_1 = W'_1.\theta_1 \cup \{a_{n1} \mapsto \tau\}$$

$$W_n.\theta_2 = W'_1.\theta_2 \cup \{a_{n2} \mapsto \tau\}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$$

Also we know that  $n' = i + 1$

We are now required to prove

- $W \sqsubseteq W_n$ :

From Equation 37 we know that  $W \sqsubseteq W'_1$  and  $W'_1 \sqsubseteq W_n$  by construction. Therefore from Definition 1.3,  $W \sqsubseteq W_n$

- $(n - n', H'_1, H'_2) \hat{\triangleright}^A W_n$ :

From Definition 1.9 it suffices to show that

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H'_2):$$

From Equation 37 and by construction of  $W_n$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_1)):$$

From Equation 37 and by construction of  $W_n$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, n - n', H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

From Equation 37 and by construction of  $W_n$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

From Equation 37 since we know that  $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$  that means

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Therefore from Lemma 1.17 we get  $(n - i - 2 = n - n' - 1, \text{ since } n' = i + 1)$

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 2, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Since  $W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$  and from Equation 37 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in \lceil \tau \sigma \rceil_V^A$

Therefore combining the two we get

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

$$- \forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V:$$

From Equation 37 we have  $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$  that means from Definition 1.9 we have

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W'_1.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

Also from Equation 37 we know that  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in \lceil \tau \sigma \rceil_V^A$

Therefore from Lemma 1.15 and Lemma 1.16 we get

$$\forall m. (W'_1.\theta_1, m, v'_{i1}) \in \lfloor \tau \sigma \rfloor_V$$

and

$$\forall m. (W'_1.\theta_2, m, v'_{i2}) \in \lfloor \tau \sigma \rfloor_V$$

Combining the two we get

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

- $(W_n, n - n', v'_1, v'_2) \in \lceil (\text{ref } \tau)^\perp \sigma \rceil_V^A$ :

Here  $v'_1 = a_{n1}$  and  $v'_2 = a_{n2}$

Since  $(a_{n1}, a_{n2}) \in W_n$  and also  $W_n.\theta_1(a_{n1}) = W_n.\theta_1(a_{n1}) = \tau$

Therefore from Definition 1.4  $(W_n, v'_1, v'_2) \in \lceil (\text{ref } \tau)^\perp \sigma \rceil_V^A$

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove:  $(W, n, !(e_i)) (\gamma \downarrow_1), !(e_i)) (\gamma \downarrow_2) \in \lceil (\tau') \sigma \rceil_E^A$

Say  $e_1 = (!e_i) (\gamma \downarrow_1)$  and  $e_2 = (!e_i) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !e_i(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !e_i(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_{\mathcal{V}}^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !e_i(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !e_i(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_{\mathcal{V}}^A \quad (38)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\ell \sigma]_E^A$$

This means from Definition 1.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A \end{aligned}$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $!e_i$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t.  $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$ . Similarly since  $!e_i$  reduces to value with  $\gamma \downarrow_2$  therefore  $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A \quad (39)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A$  from Equation 39

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau \sigma))]_{\mathcal{V}}^A$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1 \cdot \hat{\beta} \wedge W'_1 \cdot \theta_1(a_{i1}) = W'_1 \cdot \theta_2(a_{i2}) = \tau \quad (\text{D1})$$

In order to prove Equation 38 we instantiate  $W'$  with  $W'_1$ . Also we know that  $n' = i + 1$

$$- W'_1 \sqsupseteq W:$$

From Equation 39

$$- (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1:$$

From Equation 39 we know that

$$(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 1.21 we get

$$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

–  $(W'_1, n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A$ :

From the evaluation rule of deref we know that  $v'_1 = H'_1(a_{i1})$  and  $v'_2 = H'_2(a_{i2})$

Since from Equation 39 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$ , therefore from Definition 1.9 we know that

$$(W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in [W'_1.\theta_1(a_{i1})]_V^A$$

And from D1 we know that  $W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau$

Therefore  $(W'_1, v'_1, v'_2) \in [(\tau) \sigma]_V^A$

Since  $\tau \sigma <: \tau' \sigma$  Therefore from Lemma 1.28, we get

$$(W'_1, n - i - 1, v'_1, v'_2) \in [(\tau') \sigma]_V^A$$

• Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From the evaluation rule of deref we know that  $v'_{i1} = a_1$  and  $v'_{i2} = a_2$

In this case from Definition 1.4 we know that

$$\forall m_1. (W'_1.\theta_1, m_1, a_1) \in [(\text{ref } \tau) \sigma]_V \quad (40)$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, a_2) \in [(\text{ref } \tau) \sigma]_V \quad (41)$$

In order to prove Equation 38 we choose  $W'$  as  $W'_1$ . And then we need to show:

–  $W \sqsubseteq W'_1$ :

Directly from Equation 39

–  $(n - n', H'_1, H'_2) \triangleright^A W'_1$ :

From Equation 39 we know that  $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 1.21 we get

$$(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$$

–  $(W'_1, n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A$ :

Let  $\tau' = A^{\ell_i}$  Since  $\tau' \sigma \not\sqsubseteq \ell$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v'_1) \in [(\tau') \sigma]_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v'_2) \in [(\tau') \sigma]_V$$

This means given  $m_1$  and it suffices to prove:

$$(W'_1.\theta_1, m_1, v'_1) \in [(\tau') \sigma]_V \quad (42)$$

Similarly given  $m_2$ , it suffices to prove:

$$(W'_1.\theta_2, m_2, v'_2) \in [(\tau') \sigma]_V \quad (43)$$

Since from Equation 39 we know that  $(n - i, H'_1, H'_2) \triangleright W'_1$  therefore from Lemma 1.27 we get

$$\forall m_{h1}. (m_{h1}, H'_1) \triangleright W'_1.\theta_1 \quad (44)$$

$$\forall m_{h2}.(m_{h2}, H_2') \triangleright W_1'.\theta_2 \quad (45)$$

Instantiating  $m_{h1}$  in Equation 44 with  $m_1 + 1$  we get  $(m_1, H_1') \triangleright W_1'.\theta_1$

Therefore from Definition 1.8, we get

$$\forall a \in \text{dom}(W_1'.\theta_1).(W_1'.\theta_1, m_1, H_1'(a)) \in \llbracket W_1'.\theta_1(a) \rrbracket_V$$

Instantiating  $a$  with  $a_1$  we get  $(W_1'.\theta_1, m_1, H_1'(a_1)) \in \llbracket W_1'.\theta_1(a) \rrbracket_V$

Since  $W_1'.\theta_1(a_{i1}) = \tau$  therefore we get

$$(W_1'.\theta_1, m_1, v_1') \in \llbracket \tau \sigma \rrbracket_V$$

and since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 1.24 we get

$$(W_1'.\theta_1, m_1, v_1') \in \llbracket \tau' \sigma \rrbracket_V$$

Similarly we also get

$$(W_1'.\theta_2, m_2, v_2') \in \llbracket \tau' \sigma \rrbracket_V$$

Finally from Definition 1.4 we get

$$(W_1', v_1', v_2') \in \llbracket (\tau') \sigma \rrbracket_V^A$$

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} := e_{i2} : \text{unit}}$$

To prove:  $(W, n, (e_{i1} := e_{i2}) (\gamma \downarrow_1), (e_{i1} := e_{i2}) (\gamma \downarrow_2)) \in \llbracket (\text{unit}) \sigma \rrbracket_E^A$

Say  $e_1 = (e_{i1} := e_{i2}) (\gamma \downarrow_1)$  and  $e_2 = (e_{i1} := e_{i2}) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H_1', v_1') \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H_2', v_2') \implies$$

$$\exists W' \sqsupseteq W.(n - n', H_1', H_2') \triangleright^A W' \wedge (W', n - n', v_1', v_2') \in \llbracket (\text{unit}) \sigma \rrbracket_V^A$$

This further means that given

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H_1', v_1') \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H_2', v_2')$$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H_1', H_2') \triangleright^A W' \wedge (W', n - n', v_1', v_2') \in \llbracket (\text{unit}) \sigma \rrbracket_V^A \quad (46)$$

$$\underline{\text{IH1}} (W, n, (e_{i1}) (\gamma \downarrow_1), (e_{i1}) (\gamma \downarrow_2)) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n.(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow_i (H_{i1}', v_1') \wedge (H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H_{i2}', v_2') \implies$$

$$\exists W_1' \sqsupseteq W.(n - i, H_{i1}', H_{i2}') \triangleright^A W_1' \wedge (W_1', n - i, v_1', v_2') \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH1 and since the  $(e_{i1} := e_{i2})$  reduces to value with both  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow (H_{i1}', v_{i1}')$ .

Similarly since  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_2$  therefore we also have  $(H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A \quad (47)$$

$$\underline{\text{IH2}} (W, n - i, (e_{i2}) (\gamma \downarrow_1), (e_{i2}) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{j1}, H_{j2}.(n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i.(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{j1}$  with  $H'_{i1}$  and  $H_{j2}$  with  $H'_{i2}$  in IH2 and since the  $(e_{i1} := e_{i2})$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e_1$  reduces  $\gamma \downarrow_1$  with  $i < n'$  steps therefore  $\exists j < (n' - i) < (n - i)$  s.t  $(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1})$ . Similarly we also have  $(H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ . Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A \quad (48)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A$  from Equation 47

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau \sigma))]_{\mathcal{V}}^A$$

Let  $v'_{i1} = a_{i1}$  and  $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \sigma \quad (\text{A1})$$

In order to prove Equation 46 we instantiate  $W'$  with  $W'_2$

$$- W'_2 \sqsupseteq W:$$

Since  $W'_1 \sqsupseteq W$  from Equation 47 and  $W'_2 \sqsupseteq W'_1$  from Equation 48

Therefore from Definition 1.3 we get  $W'_2 \sqsupseteq W$

$$- (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2:$$

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

In order to prove  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$  we need to show:

$$* \text{dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2):$$

Directly from Equation 48

$$* W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2)):$$

Directly from Equation 48

$$* \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_{\mathcal{V}}^A:$$

- (a)  $\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2):$   
 $\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$
- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :  
 From A1 we know that  $W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$   
 and since  $W'_1 \sqsubseteq W'_2$  therefore from Lemma 1.16 we get  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$
  - ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise
  - iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise
  - iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 48 and Lemma 1.17
- (b)  $\forall (a_1, a_2) \in (W'_2.\hat{\beta}). (W'_2, n - n', H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A:$   
 $\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$
- i. When  $a_1 = a_{i1}$  and  $a_2 = a_{i2}$ :  
 Since  $H'_1(a_{i1}) = v'_{j1}$  and  $H'_1(a_{i2}) = v'_{j2}$   
 From A1 we know that  $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$   
 And since from Equation 48 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \lceil (\tau) \sigma \rceil_V^A$   
 Therefore from Lemma 1.17 we get  
 $(W'_2, n - j - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$
  - ii. When  $a_1 = a_{i1}$  and  $a_2 \neq a_{i2}$ : This case cannot arise
  - iii. When  $a_1 \neq a_{i1}$  and  $a_2 = a_{i2}$ : This case cannot arise
  - iv. When  $a_1 \neq a_{i1}$  and  $a_2 \neq a_{i2}$ : From Equation 48 and from Lemma 1.17
- \*  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V:$   
When  $i = 1$   
 Given some  $m$   
 $\forall a_1 \in \text{dom}(W'_2.\theta_1).$
- when  $a_1 = a_{i1}$ :  
 From Equation 48 we know that  $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \lceil (\tau) \sigma \rceil_V^A$  thus  
 from Lemma 1.15 we know that  
 $\forall m_1. (W'_2.\theta_1, m_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$   
 Instantiating with  $m$  we get  
 $(W'_2.\theta_1, m, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$
  - Otherwise:  
 From Equation 48 and Lemma 1.27
- When  $i = 2$   
 Similar reasoning as with  $i = 1$
- $(W'_1, n - n', \text{val}'_1, v'_2) \in \lceil (\text{unit}) \sigma \rceil_V^A:$   
 From evaluation rule assign we know that  $v'_1 = v'_2 = ()$   
 Directly from Definition 1.4

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Definition 1.4 we know that this would mean that

$$\forall m_1. (W'_1.\theta_1, m_1, a_{i1}) \in \lfloor (\text{ref } \tau) \sigma \rfloor_V \quad (49)$$

$$\forall m_2. (W'_1.\theta_2, m_2, a_{i2}) \in \lfloor (\text{ref } \tau) \sigma \rfloor_V \quad (50)$$

In order to prove Equation 46 we instantiate  $W'$  with  $W'_2$  and then we need to show that:

–  $W'_2 \sqsupseteq W$ :

Since  $W'_1 \sqsupseteq W$  from Equation 47 and  $W'_2 \sqsupseteq W'_1$  from Equation 48

Therefore from Definition 1.3 we get  $W'_2 \sqsupseteq W$

–  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$ :

From the evaluation rule assign we know that

$$H'_1 = H'_{j_1}[a_{i_1} \mapsto v'_{j_1}] \text{ and } H'_2 = H'_{j_2}[a_{i_2} \mapsto v'_{j_2}]$$

In order to prove  $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$  we need to show:

\*  $\text{dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2)$ :

Directly from Equation 48

\*  $W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$ :

Directly from Equation 48

\*  $\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \llbracket W_2.\theta_1(a_1) \rrbracket_V^A$ :

(a) When  $(a_{i_1}, a_{i_2}) \in W'_2.\hat{\beta}$ :

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$ .

i. When  $a_1 = a_{i_1}$  and  $a_2 = a_{i_2}$ :

Instantiating Equation 49 and Equation 50 with  $n - n' - 1$  we get

$$W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$$

and since  $W'_1 \sqsubseteq W'_2$  therefore from Definition 1.3 we get  $W'_2.\theta_1(a_1) =$

$$W'_2.\theta_2(a_2) = \tau$$

From Equation 48 we know that  $(W'_2, v'_{j_1}, v'_{j_2}) \in \llbracket (\tau) \sigma \rrbracket_V^A$

Therefore  $(W'_2, H_1(a_{i_1})', H_2(a_{i_2})') \in \llbracket (\tau) \sigma \rrbracket_V^A$

ii. When  $a_1 = a_{i_1}$  and  $a_2 \neq a_{i_2}$ : This case cannot arise

iii. When  $a_1 \neq a_{i_1}$  and  $a_2 = a_{i_2}$ : This case cannot arise

iv. When  $a_1 \neq a_{i_1}$  and  $a_2 \neq a_{i_2}$ : From Equation 48

(b) When  $(a_{i_1}, a_{i_2}) \notin W'_2.\hat{\beta}$ :

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$ .

i. When  $a_1 = a_{i_1}$  and  $a_2 = a_{i_2}$ : This case cannot arise

ii. When  $a_1 = a_{i_1}$  and  $a_2 \neq a_{i_2}$ :

From Equation 48 we know that  $(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$  and since

$(a_{i_1}, a_2) \in W'_2.\hat{\beta}$  therefore from Definition 1.9 we know that

$$(W'_2.\theta_1(a_{i_1}) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in \llbracket W'_2.\theta_1(a_{i_1}) \rrbracket_V^A) \quad (51)$$

Instantiating Equation 49 and Equation 50 with  $n - i - j - 1$  we get

$$W'_1.\theta_1(a_{i_1}) = \tau \sigma \text{ therefore from monotonicity we also have } W'_2.\theta_1(a_{i_1}) = \tau \sigma.$$

As a result from Equation 51 we get  $W'_2.\theta_2(a_2) = \tau \sigma$

Also since from Equation 51  $(W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in \llbracket \tau \sigma \rrbracket_V^A$  and  $\tau \sigma \searrow \ell, \ell \sigma \not\sqsubseteq \mathcal{A}$  therefore from Lemma 1.15 we know that



$$\forall m.(W'_2.\theta_1, m, H'_{j_1}(a_{i_1})) \in \lfloor \tau \sigma \rfloor_V \quad (52)$$

$$\forall m.(W'_2.\theta_2, m, H'_{j_2}(a_2)) \in \lfloor \tau \sigma \rfloor_V \quad (53)$$

Instantiating  $m$  with  $n - i - j - 1$  in Equation 52 and Equation 53 to get

$$(W'_2.\theta_1, n - i - j - 1, H'_{j_1}(a_{i_1})) \in \lfloor \tau \sigma \rfloor_V$$

and

$$(W'_2.\theta_2, n - i - j - 1, H'_{j_2}(a_2)) \in \lfloor \tau \sigma \rfloor_V$$

Since  $H'_1(a_{i_1}) = v'_{j_1}$  and  $H'_2(a_2) = H'_{j_2}(a_2)$

Again from Equation 48 we know that  $(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in \lceil (\tau) \sigma \rceil_V^A$ .

This means from Lemma 1.15 and instantiating it with  $n - i - j - 1$  we get

$$(W'_2.\theta_1, n - i - j - 1, v'_{j_1}) \in \lfloor (\tau) \sigma \rfloor_V \quad (54)$$

Therefore from Equation 53 and Equation 54 we have

$$(W'_2, n - i - j - 1, H'_1(a_{i_1}), H'_2(a_2)) \in \lceil \tau \sigma \rceil_V^A$$

iii. When  $a_1 \neq a_{i_1}$  and  $a_2 = a_{i_2}$ :

Symmetric case as (ii)

iv. When  $a_1 \neq a_{i_1}$  and  $a_2 \neq a_{i_2}$ :

From Equation 48 and Definition 1.9

\*  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V$ :

When  $i = 1$

Given some  $m$

$\forall a_1 \in \text{dom}(W'_2.\theta_1)$ .

· when  $a_1 = a_{i_1}$ :

From Equation 48 we know that  $(W'_2, v'_{j_1}, v'_{j_2}) \in \lceil (\tau) \sigma \rceil_V^A$  thus from Lemma 1.15 we know that

$$(W'_2.\theta_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$$

· Otherwise:

From Equation 48 and Lemma 1.27

When  $i = 2$

Similar reasoning as with  $i = 1$

–  $(W'_1, n - n', v'_1, v'_2) \in \lceil (\text{unit}) \sigma \rceil_V^A$ :

From evaluation rule assign we know that  $v'_1 = v'_2 = ()$

Directly from Definition 1.4

13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove:  $(W, n, \Lambda e_i (\gamma \downarrow_1), \Lambda e_i (\gamma \downarrow_2)) \in \lceil (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rceil_E^A$

Say  $e_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $e_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition of  $\lceil (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rceil_E^A$  it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$$

This means that given  $\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A \quad (55)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $v'_2 = e_2 = \Lambda e_i (\gamma \downarrow_2)$ . We choose  $W' = W$  and we know that  $n' = 0$  we need to show the following:

- $W \sqsubseteq W$ : From Definition 1.3
- $(n, H_1, H_2) \stackrel{A}{\triangleright} W$ : Given
- $(W, n, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$   
Here  $v'_1 = \Lambda e_i (\gamma \downarrow_1)$  and  $v'_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition 1.4 it suffices to prove

$$\forall W' \sqsubseteq W. \forall \ell' \in \mathcal{L}. \forall j < n.$$

$$((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A) \\ \wedge \forall \theta_l \sqsubseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e \sigma}) \\ \wedge \forall \theta_l \sqsubseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e \sigma})$$

This means given some  $W' \sqsubseteq W$ ,  $\ell' \in \mathcal{L}$  and  $j < n$  we need to show that

$$- \forall W' \sqsubseteq W. \forall \ell' \in \mathcal{L}. \forall j < n. \\ ((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A):$$

This means that given some  $W' \sqsubseteq W$ ,  $\ell' \in \mathcal{L}$ ,  $j < n$  we need to prove

$$((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A)$$

From Definition 1.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \stackrel{A}{\triangleright} W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsubseteq W. (j - m, H'_{s1}, H'_{s2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

This means for some  $H_{s1}$  and  $H_{s2}$  and some  $m < j$  we are given  $(j, H_{s1}, H_{s2}) \stackrel{A}{\triangleright} W \wedge m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$

And we need to show that

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_{\mathcal{V}}^A$$

We instantiate IH1 with  $H_{s1}$ ,  $H_{s2}$ ,  $m$  and  $\sigma \cup \{\alpha \mapsto \ell'\}$  to obtain

$$\exists W'_1 \sqsupseteq W.(n - m, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - m, v'_{i1}, v'_{i2}) \in [\tau \sigma]_{\mathcal{V}}^A \cup \{\alpha \mapsto \ell'\}$$

Since  $j < n$  therefore from Lemma 1.21 and Lemma 1.17 we get

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_{\mathcal{V}}^A$$

$$- \forall \theta_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}.((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

From Lemma 1.25 we know that  $(W'.\theta_1, \gamma \downarrow_1) \in [\Gamma]_{\mathcal{V}}$ . Therefore, we can apply Theorem 1.22 with  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$\forall k. (W'.\theta_1, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

From Lemma 1.16 we get

$$\forall \theta_l \sqsupseteq W'.\theta_1. \forall k. (\theta_l, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

$$- \forall \theta_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}.((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

Similar reasoning as in the previous case

#### 14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell''/\alpha]}$$

To prove:  $(W, n, (e[])(\gamma \downarrow_1), (e[])(\gamma \downarrow_2)) \in [(\tau[\ell''/\alpha]) \sigma]_{\mathcal{E}}^A$

This means from Definition 1.5 we need to prove:

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, (e[])(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[])(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_{\mathcal{V}}^A$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, (e[])(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[])(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_{\mathcal{V}}^A \quad (56)$$

$$\underline{\text{IH}} (W, n, (e)(\gamma \downarrow_1), (e)(\gamma \downarrow_2)) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_{\mathcal{E}}^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH and since the  $(e[])$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t.  $(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly  $(e[])$  also reduces to value with  $\gamma \downarrow_2$  therefore we also have  $(H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W.(n-i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n-i, v'_{i1}, v'_{i2}) \in [(\forall \alpha.(\ell_e, \tau))^\ell \sigma]_V^A \quad (57)$$

We case analyze on  $(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(\forall \alpha.(\ell_e, \tau))^\ell \sigma]_V^A$  from Equation 57

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 1.4 we know that

$$(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(\forall \alpha.(\ell_e, \tau)) \sigma]_V^A$$

Here  $v'_{i1} = \Lambda e_{i1}$  and  $v'_{i2} = \Lambda e_{i2}$

This further means that we have

$$\begin{aligned} \forall W'' \sqsupseteq W'_1. \forall \ell' \in \mathcal{L}. \forall j < n-i. ((W'', j, e_{i1}, e_{i2}) \in [\tau[\ell'/\alpha]]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_1, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_{i1}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_2, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_{i2}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \end{aligned} \quad (E1)$$

Instantiating the first conjunct of (E1) with  $W'_1, \ell''$  and  $n-i-1$  we get

$$((W'_1, n-i-1, e_{i1}, e_{i2}) \in [\tau[\ell''/\alpha]]_E^A)$$

Therefore from Definition 1.5 we get

$$\forall H_1, H_2. (n-i-1, H_1, H_2) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < (n-i-1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W''' \sqsupseteq W'_1. ((n-i-1)-k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1)-k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Instantiating  $H_1$  and  $H_2$  with  $H'_{i1}$  and  $H'_{i2}$  and since  $e[]$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e$  with  $\gamma \downarrow_1$  reduces in  $i < n' < n$  steps. Therefore  $\exists k < (n-i-1)$  steps in which  $e_{i1}$  reduces. Also since  $e[]$  reduces to value with  $\gamma \downarrow_2$  therefore  $e_{i2}$  must also reduce. As a result we get

$$\exists W''' \sqsupseteq W'_1. ((n-i-1)-k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1)-k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Since  $n' = i + k + 1$  therefore we are done

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Equation 56 we know that we need to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau[\ell''/\alpha] \sigma = \mathbf{A}^{\ell_i}$  and since  $\tau[\ell''/\alpha] \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

$$\text{This means in order to prove } \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

From Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V) \wedge (W'.\theta_1, m_2, v'_2) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V \quad (58)$$

In this case from Definition 1.6 we know that

$$\forall m.(W'_1.\theta_1, m, \Lambda e_{h1}) \in \lfloor \forall \alpha.(\ell_e, \tau) \sigma \rfloor_V \quad (59)$$

$$\forall m.(W'_1.\theta_2, m, \Lambda e_{h2}) \in \lfloor \forall \alpha.(\ell_e, \tau) \sigma \rfloor_V \quad (60)$$

Applying Definition 1.6 on Equation 59 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in \lfloor \tau[\ell'/\alpha] \rfloor_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_1$$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces  $\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in \lfloor \tau[\ell'/\alpha] \rfloor_E^{\ell_e[\ell'/\alpha]}$  (FB-FE1)

Instantiating  $\theta'$  with  $W'_1.\theta_1$ ,  $j_1$  with  $m_1 + t_1 + 1$  and  $\ell'$  with  $\ell''$

Therefore we get  $(W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_E^{\ell_e}$

From Definition 1.7, we get

$$\begin{aligned} \forall H.(m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 57 we have

$$(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$$

Therefore from Lemma 1.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $H$  with  $H'_{j1}$  from Equation 57 and  $k_c$  with  $t_1$ , we get

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad (\text{CF1}) \end{aligned}$$

Similarly applying Definition 1.6 to Equation 60 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}[v/x]) \in \lfloor \tau[\ell'/\alpha] \rfloor_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate  $m$  with  $m_2 + 1 + t_2$  where  $t_2$  is the number of steps in which  $e_{h2}$  reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}) \in \lfloor \tau[\ell'/\alpha] \rfloor_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-FE2})$$

Instantiating  $\theta'$  with  $W'_1.\theta_2$ ,  $j_1$  with  $m_2 + t_2 + 1$  and  $\ell'$  with  $\ell''$

Therefore we get  $(W'_1.\theta_2, m_2 + t_2 + 1, e_{h2}) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_E^{\ell_e[\ell''/\alpha]}$

From Definition 1.7, we get

$$\begin{aligned} & \forall H.(m_2 + t_2 + 1, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1).(H, e_{h2}) \Downarrow_{k_c} (H'_2, v'_1) \implies \\ & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 57 we have

$$(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 1.27 we get

$$\forall m. (m, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(m_2 + 1 + t_2, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $H$  with  $H'_{j2}$  from Equation 57 and  $k_c$  with  $t_2$ , we get

$$\begin{aligned} & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in \lfloor \tau[\ell''/\alpha] \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad \text{(CF2)} \end{aligned}$$

In order to prove Equation 56 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_1.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2):$$

From CF1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CF2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since  $(n - i, H'_{j1}, H'_{j2}) \triangleright W'_1$  therefore from Definition 1.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CF1 and CF2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$$

$$- \forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

4 cases arise for each  $a_1$  and  $a_2$

i.  $H'_{i1}(a_1) = H'_1(a_1) \wedge H'_{i2}(a_2) = H'_2(a_2)$ :

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2):$$

We know from Equation 57 that  $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 1.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CF1 and CF2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From Equation 57 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

This means from Definition 1.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$

(since  $i < n'$ ) therefore from Lemma 1.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

Same as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From CF1 and CF2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e[\ell''/\alpha] \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from CF1 and CF2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ . Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V \text{ and}$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii.  $H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2)$ :

\*  $W'.\theta_1(a_1) = W'.\theta_2(a_2)$ :

Same as in the previous case

\*  $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$ :

From CF2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e[\ell''/\alpha] \sigma$  in the world before the modification. Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e[\ell''/\alpha] \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 57 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$  that means from Definition 1.9 that  $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$ . Since  $(\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H'_{i1}(a_1)$  must also have a label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_2) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.16 we get  
 $(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$

Since from CF2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$  therefore from Definition 1.8 we know that  $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

iv.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2)$ :

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

$i = 1$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we apply Theorem 1.22 on  $e_{h1}$  and  $e_{h2}$  but this time  $m + 2 + t_1$  and  $m + 2 + t_2$  where  $t_1$  and  $t_2$  are the number of steps in which  $e_{h1}$  and  $e_{h2}$  reduces respectively. This will give us

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

and

$$\begin{aligned} \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since we have  $(m + 1, H'_1) \triangleright \theta'_1$  and  $(m + 1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 1.8

$i = 2$

Symmetric to  $i = 1$

(b)  $(W', n - n' - 1, v'_1, v'_2) \in [\tau[\ell''/\alpha] \sigma]_V^A$ :

Let  $\tau[\ell''/\alpha] = \mathbf{A}^{\ell_i}$  Since  $\tau[\ell''/\alpha] \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CF1 and CF2 we and Definition 1.4 we get the desired.

## 15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove:  $(W, n, \nu e (\gamma \downarrow_1), \nu e (\gamma \downarrow_2)) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_E^A$

Say  $e_1 = \nu e (\gamma \downarrow_1)$  and  $e_2 = \nu e (\gamma \downarrow_2)$

From Definition of  $[(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_E^A$  it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^A \end{aligned}$$

This means that given  $\forall H_1, H_2. (n', H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$



We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^A \quad (61)$$

$$\underline{\text{IH1}} (W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that  $H'_1 = H_1$ ,  $H'_2 = H_2$ ,  $v'_1 = e_1 = \nu e (\gamma \downarrow_1)$  and  $v'_2 = e_2 = \nu e (\gamma \downarrow_2)$ . We choose  $W' = W$  and we know that  $n' = 0$ . We need to show the following:

- $W \sqsubseteq W$ : From Definition 1.3

- $(n, H_1, H_2) \triangleright^A W$ : Given

- $(W, n, v'_1, v'_2) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^A$

Here  $v'_1 = \nu e (\gamma \downarrow_1)$  and  $v'_2 = \nu e (\gamma \downarrow_2)$

From Definition 1.4 it suffices to prove

$$\forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e} \sigma \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e} \sigma$$

We need to prove:

$$- \forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A:$$

This means given some  $W' \sqsupseteq W$ ,  $j < n$  and given that  $\mathcal{L} \models c \sigma$  we need to show that

$$(W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A$$

From Definition 1.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \triangleright^A W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

This means for some  $H_{s1}, H_{s2}, m < j$  s.t

$$(H_{s1}, H_{s2}) \triangleright^A W \wedge (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$$

And we need to show that

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

We instantiate IH1 with  $H_{s1}, H_{s2}$  and  $m$  to obtain

$$\exists W'_1 \sqsupseteq W. (n - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, n - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

Since  $j < n$  therefore from Lemma 1.21 and Lemma 1.17 we get

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

- $\forall \theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$ :  
This means given  $\theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c$   
We need to prove:  $(\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$   
From Lemma 1.25 we know that  $\forall m_1. (W'.\theta_1, m_1, \gamma \downarrow_1) \in [\Gamma]_V$ . Therefore by  
instantiating  $m_1$  at  $j$  we can apply Theorem 1.22 to get  
 $(\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$
- $\forall \theta_l \sqsupseteq W.\theta_2, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$ :  
Symmetric reasoning as in the previous case

16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau}$$

To prove:  $(W, n, (e\bullet)(\gamma \downarrow_1), (e\bullet)(\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

This means from Definition 1.5 we need to prove:

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e\bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e\bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e\bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e\bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A \quad (62)$$

$$\underline{\text{IH}} (W, n, (e)(\gamma \downarrow_1), (e)(\gamma \downarrow_2)) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^A$$

Instantiating  $H_{i1}$  with  $H_1$  and  $H_{i2}$  with  $H_2$  in IH and since the  $(e\bullet)$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps therefore  $\exists i < n' < n$  s.t  $(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$ . Similarly since  $(e\bullet)$  reduces to value with  $\gamma \downarrow_2$  therefore also have  $(H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$ . Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^A \quad (63)$$

We case analyze on  $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^A$  from Equation 63

- Case  $\ell \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 1.4 we know that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^{\mathcal{A}}$$

Here  $v'_{i1} = \nu e_{i1}$  and  $v'_{i2} = \nu e_{i2}$

This further means that we have

$$\begin{aligned} \forall W' \sqsupseteq W. \forall j < n - i. \mathcal{L} \models c \sigma &\implies ((W', j, e_{i1}, e_{i2}) \in [\tau \sigma]_E^{\mathcal{A}}) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma}) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i2}) \in [\tau \sigma]_E^{\ell_e \sigma}) \end{aligned} \quad (\text{CE1})$$

Instantiating the first conjunct of (CE1) with  $W'_1, \ell''$  and  $n - i - 1$  we get

$$((W'_1, n - i - 1, e_{i1}, e_{i2}) \in [\tau \sigma]_E^{\mathcal{A}})$$

Therefore from Definition 1.5 we get

$$\begin{aligned} \forall H_1, H_2. (n - i - 1, H_1, H_2) \triangleright^{\mathcal{A}} W'_1 \wedge \forall k < (n - i - 1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge \\ (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) &\implies \\ \exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^{\mathcal{A}} W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) &\in [(\tau) \sigma]_V^{\mathcal{A}} \end{aligned}$$

Instantiating  $H_1$  and  $H_2$  with  $H'_{i1}$  and  $H'_{i2}$  and since  $e[]$  reduces to value with  $\gamma \downarrow_1$  in  $n' < n$  steps and  $e$  with  $\gamma \downarrow_1$  reduces in  $i < n' < n$  steps. Therefore  $\exists k < (n' - i - 1)$  steps in which  $e_{i1}$  reduces. Also since  $e[]$  reduces to value with  $\gamma \downarrow_2$  therefore  $e_{i2}$  must also reduce. As a result we get

$$\begin{aligned} \exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^{\mathcal{A}} W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) &\in [(\tau[\ell''/\alpha]) \sigma]_V^{\mathcal{A}} \\ \text{Since } n' = i + k + 1 \text{ therefore we are done} \end{aligned}$$

- Case  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From Equation 62 we know that we need to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^{\mathcal{A}}$$

In this case since we know that  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Let  $\tau \sigma = \mathbf{A}^{\ell_i}$  and since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove  $\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^{\mathcal{A}}$

From Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given  $m_1$  and  $m_2$  it suffices to prove:

$$(\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (64)$$

In this case from Definition 1.6 we know that

$$\forall m. (W'_1.\theta_1, m, \nu e_{h1}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (65)$$

$$\forall m. (W'_1.\theta_2, m, \nu e_{h2}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (66)$$

Applying Definition 1.6 to Equation 65 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma} \text{ where } \theta = W'_1.\theta_1$$

We instantiate  $m$  with  $m_1 + 2 + t_1$  where  $t_1$  is the number of steps in which  $e_{h1}$  reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-CE1})$$

Instantiating  $\theta'$  with  $W'_1.\theta_1$ ,  $j_1$  with  $m_1 + t_1 + 1$  and since we know that  $\mathcal{L} \models c \sigma$ . Therefore we get

$$(W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

From Definition 1.7, we get

$$\begin{aligned} \forall H. (m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 63 we have

$$(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 1.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $m$  with  $m_1 + 1 + t_1$  we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating  $H$  with  $H'_{i1}$  from Equation 63 and  $k_c$  with  $t_1$ , we get

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{CCE1}) \end{aligned}$$

Similarly applying Definition 1.6 to Equation 66 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}) \in [\tau \sigma]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate  $m$  with  $m_2 + 2 + t_2$  where  $t_2$  is the number of steps in which  $e_{h2}$  reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}) \in [\tau]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-CE2})$$

Instantiating  $\theta'$  with  $W'_1.\theta_2$ ,  $j_1$  with  $m_2 + t_2 + 1$  and  $\ell'$  with  $\ell''$

$$\text{Therefore we get } (W'_1.\theta_2, m_2 + t_2 + 1, e_{h2}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

From Definition 1.7, we get

$$\begin{aligned} \forall H. (m_2 + t_2, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1). (H, e_{h2}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in [\tau \sigma]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma))
\end{aligned}$$

Since from Equation 63 we have

$$(n - i, H'_{i1}, H'_{i2}) \overset{A}{\triangleright} W'_1$$

Therefore from Lemma 1.27 we get

$$\forall m. (m, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $m$  with  $m_2 + 1 + t_2$  we get

$$(m_2 + 1 + t_2, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating  $H$  with  $H'_{i2}$  from Equation 57 and  $k_c$  with  $t_2$ , we get

$$\begin{aligned}
& \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{CCE2})
\end{aligned}$$

In order to prove Equation 62 we choose  $W'$  to be  $(\theta'_1, \theta'_2, W'_1.\beta)$ . Now we need to show two things:

(a)  $(n - n', H'_1, H'_2) \triangleright W'$ :

From Definition 1.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2):$$

From CCE1 we know that  $(m_1 + 1, H'_1) \triangleright \theta'_1$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CCE2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$ , therefore from Definition 1.8 we get  $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since  $(n - i, H'_{j1}, H'_{j2}) \triangleright W'_1$  therefore from Definition 1.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CCE1 and CCE2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  therefore

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$$

$$\begin{aligned}
- \forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge \\
(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:
\end{aligned}$$

4 cases arise for each  $a_1$  and  $a_2$

i.  $H'_{i1}(a_1) = H'_1(a_1) \wedge H'_{i2}(a_2) = H'_2(a_2)$ :

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

We know from Equation 57 that  $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 1.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since  $W'.\hat{\beta} = W'_1.\hat{\beta}$  by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CCE1 and CCE2 we know that  $W'_1.\theta_1 \sqsubseteq \theta'_1$  and  $W'_1.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 1.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From Equation 63 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

This means from Definition 1.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'_1 \sqsubseteq W'$  and  $n - n' - 1 < n - i - 1$  (since  $i < n'$ ) therefore from Lemma 1.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

$$\text{ii. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE1 and CCE2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Also from CCE1 and CCE2,  $(m_1 + 1, H'_1) \triangleright \theta'_1$  and  $(m_2 + 1, H'_2) \triangleright \theta'_2$ .

Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V \text{ and}$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since  $m_1$  and  $m_2$  are arbitrary indices therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$\text{iii. } H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that  $a_2$  was protected at  $\ell_e \sigma$  in the world before the modification. Since  $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$  (given) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_e \sigma \not\sqsubseteq \mathcal{A}$ . And thus,  $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 63 we know that  $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$  that means from Definition 1.9 that  $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$ . Since  $(\ell_e \sigma) \sqsubseteq \ell'$  therefore from Definition 1.4 we know that  $H'_{i1}(a_1)$  must have a label  $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_2) \quad (\text{S})$$

Instantiating the (F) with  $m_1$  and using Lemma 1.16 we get  
 $(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$

Since from CCE2 we know that  $(m_2 + 1, H'_2) \triangleright \theta'_2$  therefore from Definition 1.8 we know that  $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

iv.  $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2)$ :

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

$i = 1$

This means that given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we apply Theorem 1.22 on  $e_{h1}$  and  $e_{h2}$  but this time  $m + 2 + t_1$  and  $m + 2 + t_2$  where  $t_1$  and  $t_2$  are the number of steps in which  $e_{h1}$  and  $e_{h2}$  reduces respectively. This will give us

$$\exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

and

$$\exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma))$$

Since we have  $(m + 1, H'_1) \triangleright \theta'_1$  and  $(m + 1, H'_2) \triangleright \theta'_2$  therefore we get the desired from Definition 1.8

$i = 2$

Symmetric to  $i = 1$

(b)  $(W', n - n' - 1, v'_1, v'_2) \in [\tau \sigma]_V^A$ :

Let  $\tau = A^{\ell_i}$  Since  $\tau \sigma \searrow \ell \sigma$  and since  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore  $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CCE1 and CCE2 we and Definition 1.4 we get the desired.

□

**Lemma 1.27** (FG: Binary heap well formedness implies unary heap well formedness).  $\forall H_1, H_2, W. (n, H_1, H_2) \triangleright W \implies \forall i \in \{1, 2\}. \forall m. (m, H_i) \triangleright W.\theta_i$

*Proof.* Directly from Definition 1.9

□

**Lemma 1.28** (FG: Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1.  $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies [(A \sigma)]_V^A \subseteq [(A' \sigma)]_V^A$$

2.  $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^A \subseteq [(\tau' \sigma)]_E^A$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^A$$

$$\text{IH1: } [(\tau'_1 \sigma)]_V^A \subseteq [(\tau_1 \sigma)]_V^A$$

$$\text{IH2: } [(\tau_2 \sigma)]_E^A \subseteq [(\tau'_2 \sigma)]_E^A$$

It suffices to prove:

$$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^A$$

$$\text{This means that given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^A$$

$$\text{And it suffices to prove: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^A$$

From Definition 1.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A &\implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V &\implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V &\implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 1.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A &\implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2 \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V &\implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e} \sigma) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V &\implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e} \sigma) \end{aligned}$$

This means given some  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$  we need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A) :$$

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (67)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 67 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using IH2 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$



(b)  $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$ :

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$  and since  $\tau'_1 \sigma <: \tau_1 \sigma$  therefore from Lemma 1.24 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (68)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \quad (69)$$

Therefore from Equation 68 and 69 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since  $\tau_2 \sigma <: \tau'_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 1.24 and 1.23 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

(c)  $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$ :

Similar reasoning as in the previous case

## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $\lceil [(\tau_1 \times \tau_2) \sigma] \rceil_V^A \subseteq \lceil [(\tau'_1 \times \tau'_2) \sigma] \rceil_V^A$

IH1:  $\lceil [\tau_1 \sigma] \rceil_V^A \subseteq \lceil [\tau'_1 \sigma] \rceil_V^A$

IH2:  $\lceil [\tau_2 \sigma] \rceil_V^A \subseteq \lceil [\tau'_2 \sigma] \rceil_V^A$

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil [(\tau_1 \times \tau_2) \sigma] \rceil_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil [(\tau'_1 \times \tau'_2) \sigma] \rceil_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil [(\tau_1 \times \tau_2) \sigma] \rceil_V^A$

Therefore from Definition 1.4 we are given:

$$(W, n, v_1, v'_1) \in \lceil [\tau_1 \sigma] \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil [\tau_2 \sigma] \rceil_V^A \quad (70)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil [(\tau'_1 \times \tau'_2) \sigma] \rceil_V^A$

Again from Definition 1.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in \lceil [\tau'_1 \sigma] \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil [\tau'_2 \sigma] \rceil_V^A$$

Since from Equation 70 we know that  $(W, n, v_1, v'_1) \in \lceil [\tau_1 \sigma] \rceil_V^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in \lceil [\tau'_1 \sigma] \rceil_V^A$

Similarly since  $(W, n, v_2, v'_2) \in \lceil [\tau_2 \sigma] \rceil_V^A$  from Equation 70 therefore from IH2 we have  $(W, n, v_2, v'_2) \in \lceil [\tau'_2 \sigma] \rceil_V^A$

### 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $\lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A \subseteq \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau'_1 \sigma) \rceil_V^A$

IH2:  $\lceil (\tau_2 \sigma) \rceil_V^A \subseteq \lceil (\tau'_2 \sigma) \rceil_V^A$

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A. (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s2} = \text{inl } v_{i2}$ :

From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1 \sigma \rceil_V^A \tag{71}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_V^A$$

From Equation 71 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_V^A$$

(b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2 \sigma \rceil_V^A \tag{72}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_V^A$$

From Equation 72 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_V^A$$

### 4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $\lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A \subseteq \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau_2 \sigma) \rceil_V^A$

IH2:  $\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$

It suffices to prove:  $\forall (W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A.$

$(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

This means that given:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A$

Therefore from Definition 1.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \quad (\text{Sub-F1}) \end{aligned}$$

And it suffices to prove:  $(W, n, \Lambda_{e_1}, \Lambda_{e_2}) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^A$

Again from Definition 1.4, it suffices to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}) \end{aligned}$$

This means we are required to show:

$$(a) \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A):$$

By instantiating the first conjunct of Sub-F1 with  $W'', n''$  and  $\ell''$  we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from IH1 instantiated at  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

$$(b) \forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}):$$

By instantiating the second conjunct of Sub-F1 with  $\theta'_l$  and  $\ell''$  we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha]})$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 1.24 and Lemma 1.23 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\ell'_e[\ell''/\alpha]})$$

$$(c) \forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}):$$

Similar reasoning as in the previous case

## 5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell'_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove:  $[((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^A \subseteq [((c_2 \xrightarrow{\ell'_e} \tau_2)) \sigma]_V^A$

IH:  $[(\tau_1 \sigma)]_E^A \subseteq [(\tau_2 \sigma)]_E^A$

It suffices to prove:  $\forall (W, n, \nu_{e_1}, \nu_{e_2}) \in [((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^A. (W, n, \nu_{e_1}, \nu_{e_2}) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^A$

This means that given:  $(W, n, \nu_{e_1}, \nu_{e_2}) \in [((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^A$

Therefore from Definition 1.4 we are given:

$$\begin{aligned}
\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma &\implies (W', n', e_1, e_2) \in [\tau_1 \sigma]_E^A \wedge \\
\forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 &\implies (\theta_l, k, e_1) \in [\tau_1 \sigma]_E^{\ell_e \sigma} \wedge \\
\forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 &\implies (\theta_l, k, e_2) \in [\tau_1 \sigma]_E^{\ell_e \sigma} \quad (\text{Sub-C1})
\end{aligned}$$

And it suffices to prove:  $(W, n, \nu e_1, \nu e_2) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^A$

Again from Definition 1.4, it suffices to prove:

$$\begin{aligned}
\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma &\implies (W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A \wedge \\
\forall \theta'_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 &\implies (\theta'_l, j, e_1) \in [\tau_2 \sigma]_E^{\ell'_e \sigma} \wedge \\
\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 &\implies (\theta'_l, j, e_2) \in [\tau_2 \sigma]_E^{\ell'_e \sigma}
\end{aligned}$$

This means that we are required to show the following:

$$(a) \quad \forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A:$$

We are given  $W'' \sqsupseteq W, n'' < n$  also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with  $W''$  and  $n''$  we know that the following holds

$$(W'', n'', e_1, e_2) \in [\tau_1 \sigma]_E^A$$

Therefore from IH we get  $(W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A$

$$(b) \quad \forall \theta'_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \implies (\theta'_l, k, e_1) \in [\tau_2 \sigma]_E^{\ell'_e \sigma}:$$

We are given some  $\theta'_l \sqsupseteq W.\theta_1, k$ , also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with  $\theta'_l$  we know that the following holds

$$(\theta'_l, k, e_1) \in [\tau_1 \sigma]_E^{\ell_e \sigma}$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  and  $\ell'_e \sigma \sqsubseteq \ell_e \sigma$  therefore from Lemma 1.23 and Lemma 1.24 we get

$$(\theta'_l, k, e_1) \in [\tau_2 \sigma]_E^{\ell'_e \sigma}$$

$$(c) \quad \forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_2) \in [\tau_2 \sigma]_E^{\ell'_e \sigma}:$$

Similar reasoning as in the previous case

## 6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $[((\text{ref } \tau) \sigma)]_V^A \subseteq [((\text{ref } \tau) \sigma)]_V^A$

Directly from Definition 1.4

## 7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $[((\mathbf{b}) \sigma)]_V^A \subseteq [((\mathbf{b}) \sigma)]_V^A$

Directly from Definition 1.4

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\lceil ((\text{unit}) \sigma) \rceil_V^A \subseteq \lceil ((\text{unit}) \sigma) \rceil_V^A$

Directly from Definition 1.4

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A^{\ell'}} \text{FGsub-label}$$

To prove:  $\lceil ((A^\ell) \sigma) \rceil_V^A \subseteq \lceil ((A^{\ell'}) \sigma) \rceil_V^A$

2 cases arise

1.  $\ell \sigma \sqsubseteq \ell' \sigma$ :

From Definition 1.4 it suffices to prove:  $\lceil ((A) \sigma) \rceil_V^A \subseteq \lceil ((A') \sigma) \rceil_V^A$

This we get directly from IH (Statement (1))

2.  $\ell \sigma \not\sqsubseteq \ell' \sigma$ :

We need to prove that

$$\forall (W, n, v_1, v_2) \in \lceil A \sigma \rceil_V^A. (W, n, v_1, v_2) \in \lceil A' \sigma \rceil_V^A$$

From Definition 1.4 it suffices to prove:

$$\forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in \lceil A \sigma \rceil_V. (W(n). \theta_i, m, v_i) \in \lceil A' \sigma \rceil_V$$

Since  $A \sigma <: A' \sigma$  therefore from Lemma 1.24 we get the desired

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\lceil (\tau \sigma) \rceil_E^A \subseteq \lceil (\tau' \sigma) \rceil_E^A$

This means we need to prove that

$$\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A. (W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$$

This means given  $\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A$

It suffices to prove that  $(W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$

From Definition 1.5 we know we are given:

$$\begin{aligned} & \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A \quad (\text{Sub-exp1}) \end{aligned}$$

And we need prove that

$$\begin{aligned} & \forall H_{21}, H_{22}, k < n. (n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22}) \implies \\ \exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \end{aligned}$$

This means that we are given some  $H_{21}, H_{22}$  and  $k < n$  such that  $(n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22})$

It suffices to prove:

$$\exists W'' \sqsupseteq W.(n - k, H'_{21}, H'_{22}) \stackrel{A}{\triangleright} W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (73)$$

Instantiating (Sub-exp1) with  $H_{21}$ ,  $H_{22}$  and  $k$  we get

$$\exists W' \sqsupseteq W.(n - k, H'_{21}, H'_{22}) \stackrel{A}{\triangleright} W' \wedge (W', n - k, v'_{21}, v'_{22}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (74)$$

We choose  $W''$  in Equation 73 as  $W'$  from Equation 74 and we are done  $\square$

**Theorem 1.29** (FG: NI). *Say*  $\text{bool} = (\text{unit} + \text{unit})$

$$\begin{aligned} & \forall v_1, v_2, e, \tau, n_1. \\ & \emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top} \\ & \emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge \\ & (\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_{-} (-, v'_2) \implies \\ & v'_1 = v'_2 \end{aligned}$$

*Proof.* Given some

$$\begin{aligned} & \emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top} \\ & \emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge \\ & (\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_{-} (-, v'_2) \end{aligned}$$

We need to prove

$$\frac{v'_1 = v'_2}{v'_1 = v'_2}$$

From Theorem 1.26 we have

$$\forall n. (\emptyset, n, v_1, v_2) \in [\text{bool}^{\top}]_{\mathcal{E}}^{\perp}$$

Therefore from Theorem 1.26 and from Definition 1.14 we have

$$\forall n. (\emptyset, n, e[v_1/x], e[v_1/x]) \in [\text{bool}^{\perp}]_{\mathcal{E}}^{\perp}$$

Therefore from Definition 1.5 we know that

$$\forall n. (\forall H_1, H_2, j < n. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W' \sqsupseteq W.(n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [(\text{unit} + \text{unit})^{\perp}]_{\mathcal{V}}^A)$$

Instantiating with  $n_1 + 1$  and then with  $\emptyset, \emptyset, n_1$  we get

$$\exists W' \sqsupseteq W.(1, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', 1, v'_1, v'_2) \in [(\text{unit} + \text{unit})^{\perp}]_{\mathcal{V}}^A$$

Since we have  $(W', 1, v'_1, v'_2) \in [(\text{unit} + \text{unit})^{\perp}]_{\mathcal{V}}^A$  therefore from Definition 1.4 we get  $v'_1 = v'_2$   $\square$

## 2 Coarse-grained IFC enforcement (SLIO\*)

### 2.1 SLIO\* type system

### 2.2 SLIO\* semantics

Judgement:  $e \Downarrow_i v$  and  $(H, e) \Downarrow_i^f (H', v)$

**Syntax, types, constraints:**

Expressions	$e ::= x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid \text{new } e \mid !e \mid e := e \mid () \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid \text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::= l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::= \mathbf{b} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{unit} \mid \forall \alpha. \tau \mid c \Rightarrow \tau \mid \text{Labeled } \ell \tau \mid \text{SLIO } \ell_i \ell_o \tau$
Constraints	$c ::= \ell \sqsubseteq \ell \mid (c, c)$

**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types  $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau$ , **unit** are included.)

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{SLIO}^*\text{-label} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau} \text{SLIO}^*\text{-unlabel} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{SLIO } \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)} \text{SLIO}^*\text{-toLabeled} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \text{SLIO } \ell_i \ell_i \tau} \text{SLIO}^*\text{-ret} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{SLIO } \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \text{SLIO } \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \text{SLIO } \ell_i \ell_o \tau'} \text{SLIO}^*\text{-bind} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau} \text{SLIO}^*\text{-sub} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \text{SLIO } \ell \ell (\text{ref } \ell' \tau)} \text{SLIO}^*\text{-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e : \text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)} \text{SLIO}^*\text{-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \text{SLIO } \ell \ell \text{ unit}} \text{SLIO}^*\text{-assign} \\
\\
\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \Lambda e : \forall \alpha. \tau} \text{SLIO}^*\text{-FI} \qquad \frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha]} \text{SLIO}^*\text{-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \nu e : c \Rightarrow \tau} \text{SLIO}^*\text{-CI} \qquad \frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau} \text{SLIO}^*\text{-CE}
\end{array}$$

Figure 5: Type system for SLIO\*



$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \tau <: \tau} \text{SLIO}^* \text{-sub-refl} \qquad \frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{SLIO}^* \text{-sub-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{SLIO}^* \text{-sub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{SLIO}^* \text{-sub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{SLIO}^* \text{-sub-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \text{SLIO } \ell_i \ell_o \tau <: \text{SLIO } \ell'_i \ell'_o \tau'} \text{SLIO}^* \text{-sub-monad} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{SLIO}^* \text{-sub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2} \text{SLIO}^* \text{-sub-constraint}
\end{array}$$

Figure 6: SLIO\* subtyping

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ } WF} \text{SLIO}^*\text{-wff-base} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{unit} \text{ } WF} \text{SLIO}^*\text{-wff-unit} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \rightarrow \tau_2) \text{ } WF} \text{SLIO}^*\text{-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \times \tau_2) \text{ } WF} \text{SLIO}^*\text{-wff-times} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 + \tau_2) \text{ } WF} \text{SLIO}^*\text{-wff-sum} \qquad \frac{\text{FV}(\ell) = \emptyset \quad \text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \ell \ \tau) \text{ } WF} \text{SLIO}^*\text{-wff-ref} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\forall \alpha. \ \tau) \text{ } WF} \text{SLIO}^*\text{-wff-forall} \qquad \frac{\Sigma; \Psi, c \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (c \Rightarrow \tau) \text{ } WF} \text{SLIO}^*\text{-wff-constraint} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ } WF \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi \vdash (\text{Labeled } \ell \ \tau) \text{ } WF} \text{SLIO}^*\text{-wff-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ } WF \quad \text{FV}(\ell_i) \in \Sigma \quad \text{FV}(\ell_o) \in \Sigma}{\Sigma; \Psi \vdash (\text{SLIO } \ell_i \ \ell_o \ \tau) \text{ } WF} \text{SLIO}^*\text{-wff-monad}
\end{array}$$

Figure 7: Well-formedness relation for SLIO\*

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x. e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 \ e_2 \Downarrow_{i+j+k+1} v_3} \text{SLIO}^*\text{-Sem-app} \\
\\
\frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{SLIO}^*\text{-Sem-prod} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{SLIO}^*\text{-Sem-fst} \\
\\
\frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{SLIO}^*\text{-Sem-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{SLIO}^*\text{-Sem-inl} \\
\\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{SLIO}^*\text{-Sem-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{SLIO}^*\text{-Sem-case1} \\
\\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{SLIO}^*\text{-Sem-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{SLIO}^*\text{-Sem-Lb} \\
\\
\frac{e \Downarrow_i \Lambda e_i \quad e_i \Downarrow_j v}{e[] \Downarrow_{i+j+1} v} \text{SLIO}^*\text{-Sem-FE} \qquad \frac{e \Downarrow_i \nu e_i \quad e_i \Downarrow_j v}{e \bullet \Downarrow_{i+j+1} v} \text{SLIO}^*\text{-Sem-CE} \\
\\
\frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{SLIO}^*\text{-Sem-ret} \\
\\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{SLIO}^*\text{-Sem-bind} \\
\\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{SLIO}^*\text{-Sem-unlabel}
\end{array}$$

### 2.3 Model for SLIO\*

$W : ((Loc \mapsto Type) \times (Loc \mapsto Type) \times (Loc \leftrightarrow Loc))$

**Definition 2.1** (SLIO\*:  $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 2.2** (SLIO\*:  $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

**Definition 2.3** (SLIO\*: Value Equivalence).

$$ValEq(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_{\mathcal{V}} \wedge (W.\theta_2, j, v_2) \in [\tau]_{\mathcal{V}} & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

**Definition 2.4** (SLIO\*: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \rightarrow \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\
[\forall \alpha. \tau]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, \ell' \in \mathcal{L}. \\
&\quad ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E\} \\
[ c \Rightarrow \tau ]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n. \\
&\quad \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E\} \\
[\text{ref } \ell \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \text{Labeled } \ell \tau\} \\
[\text{Labeled } \ell \tau]_V^A &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\
[\text{SLIO } \ell_1 \ell_2 \tau]_V^A &\triangleq \{(W, n, v_1, v_2) \mid \\
&\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\
&\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
&\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\
&\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\}
\end{aligned}$$

**Definition 2.5** (SLIO\*: Binary expression relation).

$$[\tau]_E^A \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A\}$$

**Definition 2.6** (SLIO\*: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta' \sqsupseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \forall \alpha. \tau \rrbracket_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta', j < m. \forall \ell' \in \mathcal{L}. (\theta', j, e) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E\} \\
\llbracket c \Rightarrow \tau \rrbracket_V &\triangleq \{(\theta, m, \nu e) \mid \mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < m. (\theta', j, e) \in \llbracket \tau \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v)) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbf{SLIO} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

**Definition 2.7** (SLIO\*: Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

**Definition 2.8** (SLIO\*: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

**Definition 2.9** (SLIO\*: Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

**Definition 2.10** (SLIO\*: Label substitution).  $\sigma : \text{Lvar} \mapsto \text{Label}$

**Definition 2.11** (SLIO\*: Value substitution to value pairs).  $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

**Definition 2.12** (SLIO\*: Value substitution to values).  $\delta : \text{Var} \mapsto \text{Val}$

**Definition 2.13** (SLIO\*: Unary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

**Definition 2.14** (SLIO\*: Binary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

## 2.4 Soundness proof for SLIO\*

**Lemma 2.15** (SLIO\*: Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

*Proof.* Proof by induction on  $\tau$

1. Case b:

From Definition 2.6

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 2.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

$$\text{IH1: } \forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \text{ and}$$

$$\text{IH2: } \forall m_2. (W.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$$

From (S01) we know that given some  $m$  and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

Also from (S02) we know that given some  $m$  and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH1 with  $m$  from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$$

Therefore from Definition 2.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with  $m$  from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 2.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

$$(b) \ v_1 = \text{inr}(v_{i2}) \text{ and } v_2 = \text{inr}(v_{j2})$$

Symmetric reasoning as in the (a) case above

#### 4. Case $\tau_1 \rightarrow \tau_2$ :

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V^A$$

This means from Definition 2.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V \implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E) \quad (\text{L0}) \end{aligned}$$

To prove:

$$(a) \ \forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \rightarrow \tau_2]_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

This further means that we have some  $\theta', j$  and  $v$  s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

$$(b) \ \forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V:$$

Similar reasoning with  $e_2$

5. Case  $\forall\alpha.\tau$ :

Given:  $(W, n, \Lambda e_1, \Lambda e_2) \in [\forall\alpha.\tau]_V^A$

This means from Definition 2.4 we know that

$$\begin{aligned} & \forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E) \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E) \end{aligned} \quad (\text{F0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall\alpha.\tau]_V$ :

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$$

This further means that we are given some  $\theta', m'$  and  $\ell_u$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\ell_u \in \mathcal{L}$

And we need to prove:  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

Instantiating  $\theta_l, i$  and  $\ell''$  in the second conjunct of F0 with  $\theta', m'$  and  $\ell_u$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\ell_u \in \mathcal{L}$

Therefore we get  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

(b)  $\forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall\alpha.\tau]_V$ :

Symmetric reasoning for  $e_2$

6. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, \nu e_1, \nu e_2) \in [c \Rightarrow \tau]_V^A$

This means from Definition 2.4 we know that

$$\begin{aligned} & \forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E \end{aligned} \quad (\text{C0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \nu e_1) \in [c \Rightarrow \tau]_V$ :

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E$$

This further means that we are given some  $\theta'$  and  $m'$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\mathcal{L} \models c$

And we need to prove:  $(\theta', m', e_1) \in [\tau]_E$

Instantiating  $\theta_l, j$  in the second conjunct of C0 with  $\theta', m'$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta', m', e_1) \in [\tau]_E$

(b)  $\forall m. (W.\theta_2, m, \nu e_2) \in [c \Rightarrow \tau]_V$ :

Symmetric reasoning for  $e_2$

7. Case ref  $\ell \tau$ :

From Definition 2.4 and 2.6



8. Case Labeled  $\ell \tau$ :

Given  $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 2.3 we know that

$$(W, n, v_1, v_2) \in [\tau]_V^A$$

Therefore from IH we get  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 2.3

9. Case SLIO  $\ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\text{SLIO } \ell_1 \ell_2 \tau]_V^A$

This means from Definition 2.4 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\text{SLIO } \ell_1 \ell_2 \tau]_V$

This means from Definition 2.6 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case  $l = 1$

And given some  $m$  and  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with  $l = 1$  and the given  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  we get the desired.

Case  $l = 2$

Symmetric reasoning as in the previous case above

□

**Lemma 2.16** (SLIO\*: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$$(\theta, m, v) \in \lfloor \tau \rfloor_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \lfloor \tau \rfloor_V$$

*Proof.* Proof by induction on  $\tau$

1. case b:

Directly from Definition 2.6

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

This means from Definition 2.6 we know that

$$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V \wedge (\theta, m, v_2) \in \lfloor \tau_2 \rfloor_V$$

$$\text{IH1} : (\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$$

$$\text{IH2} : (\theta', m', v_2) \in \lfloor \tau_2 \rfloor_V$$

We get the desired from IH1, IH2 and Definition 2.6

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

This means from Definition 2.6 we know that

$$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V$$

$$\text{IH} : (\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$$

Therefore from IH and Definition 2.6 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(\theta, m, (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

To prove:  $(\theta', m', (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

This means from Definition 2.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E \quad (75)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V \implies (\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$

Instantiating Equation 75 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

5. case ref  $\ell \tau$ :

From Definition 2.6 and Definition 2.1

6. case  $\forall\alpha.\tau$ :

Given:  $(\theta, m, (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

To prove:  $(\theta', m', (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

This means from Definition 2.6 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E \quad (76)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

Instantiating Equation 76 with  $\theta''', k$  and  $\ell_j$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\ell_j \in \mathcal{L}$

Therefore we get  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

7. case  $c \Rightarrow \tau$ :

Given:  $(\theta, m, (\nu e_1)) \in [c \Rightarrow \tau]_V$

To prove:  $(\theta', m', (\nu e_1)) \in [c \Rightarrow \tau]_V$

This means from Definition 2.6 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \mathcal{L} \models c \implies (\theta'', j, e_1) \in [\tau]_E \quad (77)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \mathcal{L} \models c \implies (\theta''', k, e_1) \in [\tau]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau]_E$

Instantiating Equation 77 with  $\theta''', k$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta''', k, e_1) \in [\tau]_E$

8. case Labeled  $\ell \tau$ :

Given:  $(\theta, m, (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

To prove:  $(\theta', m', (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

This means from Definition 2.6 we know that  $(\theta, m, v) \in \llbracket \tau \rrbracket_V$

IH:  $(\theta', m', v) \in \llbracket \tau \rrbracket_V$

Therefore from IH and Definition 2.6 we get the desired

9. case SLIO  $\ell_1 \ell_2 \tau$ :

Given:  $(\theta, m, e) \in \llbracket \text{SLIO } \ell_1 \ell_2 \tau \rrbracket_V$

To prove:  $(\theta', m', e) \in \llbracket \text{SLIO } \ell_1 \ell_2 \tau \rrbracket_V$

This means from Definition 2.6 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 2.6 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t. } (k_1, H) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0),  $k$  with  $k_1$ ,  $\theta_e$  with  $\theta_{e1}$ ,  $H$  with  $H_1$  and  $j$  with  $j_1$ . We know that  $k_1 < m' < m$ ,  $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$ ,  $(k_1, H_1) \triangleright \theta_{e1}$ ,  $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$  and  $i_1 + j_1 < k_1$ . Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

**Lemma 2.17** (SLIO\*: Monotonicity binary). *The following holds:*

$$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**, unit:

From Definition 2.4

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 :  $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 :  $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 2.4 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 2.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i1})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

This means from Definition 2.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E)$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)$  (BM-A2)

Similarly from Definition 2.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$$

#### 5. Case ref $\ell \tau$ :

From Definition 2.4 and Definition 2.2

#### 6. Case $\forall \alpha. \tau$ :

Given:  $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall \alpha. \tau]_V^A$

To prove:  $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall \alpha. \tau]_V^A$

This means from Definition 2.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F2})$$

Similarly from Definition 2.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with  $W'', n''$  and  $\ell''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . Also since  $n'' < n'$  and  $n' < n$  therefore  $n'' < n$ . And finally since  $\ell'' \in \mathcal{L}$  therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_l \sqsupseteq W.\theta_1$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

7. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

To prove:  $(\theta', n', (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

This means from Definition 2.4 we know that the following holds

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in \lceil \tau \rceil_E^A$  (BM-C0)

$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in \lfloor \tau \rfloor_E$  (BM-C1)

$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in \lfloor \tau \rfloor_E$  (BM-C2)

Similarly from Definition 2.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\mathcal{L} \models c$

And we are required to prove:  $(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

Instantiating BM-C0 with  $W'', n''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . And since  $\mathcal{L} \models c$  therefore we get

$(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_1 \sqsupseteq W.\theta_1$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta_l, k, e_2) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

8. Case Labeled  $\ell \tau$ :

Given:  $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

To prove:  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

From Definition 2.4 2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Therefore from IH we know that  $(W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Hence from Definition 2.4 we get  $(W', n', (\mathbf{Lb}v_1), (\mathbf{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case we know that  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $(W.\theta_2, m, v_2) \in [\tau]_V$

Since  $W.\theta_1 \sqsubseteq W'.\theta_1$  (from Definition 2.2). Therefore from Lemma 2.16 we know that

$\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since  $W.\theta_2 \sqsubseteq W'.\theta_2$  (from Definition 2.2). Therefore from Lemma 2.16 we know that

$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$

Finally from Definition 2.4 we get  $(W', n', (\mathbf{Lb}v_1), (\mathbf{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

9. Case  $\text{SLIO } \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\text{SLIO } \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(W', n', v_1, v_2) \in [\text{SLIO } \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 2.4 we are given that

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{BM-M0}) \end{aligned}$$

Similarly from Definition 2.4 it suffices to prove that

$$\begin{aligned} & \text{(a) } \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right): \\ & \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\ & (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given  $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  and since we know that  $n' \leq n$  and  $W \sqsubseteq W'$  we get the desired

$$\begin{aligned} & \text{(b) } \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right): \end{aligned}$$

Similar reasoning as in the previous case but using Lemma 2.16



□

**Lemma 2.18** (SLIO\*: Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .  
 $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$   
 To prove:  $(\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.13 it is given that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

And again from Definition 2.13 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

- $dom(\Gamma) \subseteq dom(\delta)$ :  
 Given
- $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$ :  
 Since we know that  $\forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$  (given)  
 Therefore from Lemma 2.16 we get  
 $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

□

**Lemma 2.19** (SLIO\*: Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .  
 $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.14 it is given that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

And again from Definition 2.13 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

- $dom(\Gamma) \subseteq dom(\gamma)$ :  
 Given
- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$ :  
 Since we know that  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$  (given)  
 Therefore from Lemma 2.17 we get  
 $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

□

**Lemma 2.20** (SLIO\*: Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .  
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 2.8 it is given that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

And again from Definition 2.13 we are required to prove that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$ :

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :

Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)

Therefore from Lemma 2.16 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

**Lemma 2.21** (SLIO\*: Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 2.9 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) &\in (W.\hat{\beta}).(W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 2.9 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :

Given

- $\forall (a_1, a_2) \in (W.\hat{\beta}).(W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

$$\forall (a_1, a_2) \in (W.\hat{\beta}).$$

–  $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given

–  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

Given and from Lemma 2.17

- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :

Given

□

**Theorem 2.22** (SLIO\*: Fundamental theorem unary).  $\forall \Sigma, \Psi, \Gamma, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash e : \tau \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge \\ & (\theta, n, \delta) \in [\Gamma \sigma]_V \implies \\ & (\theta, n, e \delta) \in [\tau \sigma]_E \end{aligned}$$

*Proof.* Proof by induction on SLIO\* typing derivation

1. SLIO\*-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{SLIO*}-\text{var}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, x \delta) \in [\tau \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some  $i < n$  s.t  $x \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = x \delta$  and  $i = 0$ )

It suffices to prove  $(\theta, n, x \delta) \in [\tau \sigma]_V$  (FU-V0)

Since  $(\theta, n, \delta) \in [\Gamma' \sigma]_V$  where  $\Gamma' = \Gamma \cup \{x : \tau\}$ . Therefore from Definition 2.13 we know that  $(\theta, n, \delta(x)) \in [\Gamma'(x) \sigma]_V$

So we are done.

2. SLIO\*-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e' : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \lambda x. e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-L0})$$

From Definition 2.6 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1 \sigma]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

This means given some  $\theta'', v', j$  s.t  $\theta'' \sqsupseteq \theta, j < n$  and  $(\theta'', j, v') \in [\tau_1 \sigma]_V$  (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.18 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E, \text{ s.t. } (\theta_i, j, v_x) \in [\tau_1 \sigma]_V$$

Instantiating IH with  $\theta''$  and  $v'$  from (FU-L1) we get  $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$

### 3. SLIO\*-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2 \sigma]_V$$

This means that given some  $i < n$  s.t  $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2 \sigma]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

From SLIO\*-Sem-app we know that  $v_1 = \lambda x. e'$ . Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-P1})$$

This means from Definition 2.6 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1 \sigma]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2 \sigma]_E \quad (78)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V \quad (\text{FU-P2})$$

Instantiating Equation 78 with  $\theta, (n - j - k), v_2$  and since we know that  $(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$  therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in [\tau_2 \sigma]_E$$

This means from Definition 2.7 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore we know that  $\exists J < i < n$  s.t  $i = j + k + J$  (since  $j + k + J < n$  therefore  $J < n - j - k$ ) and  $e'[v_2/x] \Downarrow_J v_f$

Therefore we have  $(\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$

Since we know that  $i = j + k + J$  and  $v = v_J$  therefore we get  $(\theta, n - i, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$  (so FU-P0 is proved)

#### 4. SLIO\*-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

Since we know that  $(e_1, e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from SLIO\*-Sem-prod we know that  $i = j + k + 1$  and  $v = (v_1, v_2)$  therefore from Definition 2.6 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \text{ and } (\theta, n - j - k - 1, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

We get this from (FU-PA1) and Lemma 2.16 and from (FU-PA2) and Lemma 2.16

#### 5. SLIO\*-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{fst}(e') \delta) \in [\tau_1 \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_1 \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_1 \sigma]_V \quad (\text{FU-F0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$$

Since we know that  $\text{fst}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j (v_1, v_2)$ . This means we have

$$(\theta, n - j, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$$

From Definition 2.6 we know the following holds

$$(\theta, n - j, v_1) \in [\tau_1 \sigma]_V \text{ and } (\theta, n - j, v_2) \in [\tau_2 \sigma]_V \quad (\text{FU-F1})$$

From SLIO\*-Sem-fst we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-F0), we are required to prove

$$(\theta, n - j - 1, v_1) \in [\tau_1 \sigma]_V$$

We get this from (FU-F1) and Lemma 2.16

6. SLIO\*-snd:

Symmetric reasoning as in the SLIO\*-fst case above

7. SLIO\*-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{inl}(e') \delta) \in [(\tau_1 + \tau_2) \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 + \tau_2) \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [(\tau_1 + \tau_2) \sigma]_V \quad (\text{FU-LE0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [\tau_1 \sigma]_V$$

Since we know that  $\text{inl}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-LE1})$$

From SLIO\*-Sem-inl we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-LE0) we are required to prove

$$(\theta, n - j - 1, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

From Definition 2.6 it suffices to prove

$$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

We get this from (FU-LE1) and Lemma 2.16

8. SLIO\*-inr:

Symmetric reasoning as in the SLIO\*-inl case above

9. SLIO\*-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in \lfloor \tau \sigma \rfloor_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$$

This means that given some  $i < n$  s.t.  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e_c \delta \Downarrow_j v_c$ . This means we have

$$(\theta, n - j, v_c) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V \quad (\text{FU-C1})$$

2 cases arise:

(a)  $v_c = \text{inl}(v_l)$ :

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t.  $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$ . This means we have

$$(\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C2})$$

From SLIO\*-Sem-case1 we know that  $i = j + k + 1$  and  $v = v_1$ . Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau \sigma \rfloor_V$$

We get this from (FU-C2) and Lemma 2.16

(b)  $v_c = \text{inr}(v_r)$ :

Symmetric reasoning as in the previous case

10. SLIO\*-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \Lambda e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\forall \alpha. \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \Lambda e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. \tau) \sigma]_V \quad (\text{FU-FI0})$$

From Definition 2.6 it further suffices to prove

$$\forall \theta'. \theta \sqsubseteq \theta', j < n. \forall \ell' \in \mathcal{L}. (\theta', j, e' \delta) \in [\tau[\ell'/\alpha]]_E$$

This means given some  $\theta', j, \ell' \in \mathcal{L}$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-FI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau[\ell'/\alpha] \sigma]_E \quad (\text{FU-FI2})$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.18 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1)

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E$$

(FU-FI2) is obtained directly from IH

11. SLIO\*-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \nu e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(c \Rightarrow \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\nu e' \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \nu e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_V \quad (\text{FU-CI0})$$

From Definition 2.6 it further suffices to prove



$$\mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < n. (\theta', j, e' \delta) \in [\tau]_E$$

This means given  $\mathcal{L} \models c$  and some  $\theta', j$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-CI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau \sigma]_E \quad (\text{FU-CI2})$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 2.18 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1). Also we know that  $\mathcal{L} \models c \sigma$  therefore  $\mathcal{L} \models (\Sigma \cup \{c\}) \sigma$

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma]_E$$

(FU-CI2) is obtained directly from IH

12. SLIO\*-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' [] \delta) \in [\tau[\ell/\alpha] \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. e' [] \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V$$

This means that given some  $i < n$  s.t  $e' [] \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V \quad (\text{FU-FE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [\forall \alpha. \tau]_E$$

From Definition 2.7 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \Lambda e_{h_1} \implies (\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$$

Since  $e' [] \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \Lambda e_i$

Therefore we know that  $(\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$

From Definition 2.6 we know that

$$\forall \theta'' \sqsupseteq \theta, x < (n - h_1), \ell_h \in \mathcal{L}. (\theta'', x, e_{h_1}) \in [(\tau[\ell_h/\alpha]) \sigma]_E$$

Instantiating  $\theta''$  with  $\theta, x$  with  $n - h_1 - 1$  and  $\ell_h$  with  $\ell$ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [(\tau[\ell/\alpha]) \sigma]_E$$

From Definition 2.7 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $e' [] \delta$  reduces in  $i$  steps therefore from SLIO\*-Sem-FE we know that  $(i = h_1 + h_2 + 1)$  and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ .

Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

13. SLIO\*-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' \bullet \delta) \in [\tau \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. e' \bullet \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some  $i < n$  s.t.  $e' \bullet \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau \sigma]_V \quad (\text{FU-CE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [c \Rightarrow \tau \sigma]_E$$

From Definition 2.7 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \nu e_{h_1} \implies (\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$$

Since  $e' \bullet \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \nu e_{h_1}$

Therefore we know that  $(\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$

From Definition 2.6 we know that

$$\mathcal{L} \models c \sigma \implies \forall \theta'' \sqsupseteq \theta, x < (n - h_1). (\theta'', x, e_{h_1}) \in [\tau \sigma]_E$$

Since we know that  $\mathcal{L} \models c \sigma$  and then we instantiate  $\theta''$  with  $\theta, x$  with  $n - h_1 - 1$ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [\tau \sigma]_E$$

From Definition 2.7 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since  $e' \bullet \delta$  reduces in  $i$  steps therefore from SLIO\*-Sem-CE we know that  $(i = h_1 + h_2 + 1)$  and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ . Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since we know that  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in [\tau \sigma]_V$$

14. SLIO\*-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } (e') : \text{SLIO } \ell \ell (\text{ref } \ell' \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{new } (e') \delta) \in [\text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma]_E$

This means that from Definition 2.7 we need to prove

$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_V$

This means that given some  $i < n$  s.t  $\text{new } (e') \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \text{new } (e') \delta$  and  $i = 0$ )

It suffices to prove

$(\theta, n, \text{new } (e') \delta) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_V$

From Definition 2.6 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket (\text{ref } \ell' \tau) \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
 Also from SLIO\*-Sem-ref we know that  $v' = a$

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in \llbracket (\text{ref } \ell' \tau) \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0})$

IH:

$(\theta_e, k, e' \delta) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_E$

From Definition 2.7 this means we have

$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V$

Since we know that  $(H, \text{new } (e')) \Downarrow_j^f (H', a)$  therefore from SLIO\*-Sem-ref we know that  
 $\exists l < j < k$  s.t  $e' \delta \Downarrow_l v_h$

Therefore we have

$(\theta_e, n - l, v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V \quad (\text{FU-R2})$

In order to prove (FU-R0) we choose  $\theta'$  as  $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_n$ :

From Definition 2.8 it suffices to prove that

$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \llbracket \theta_n(a) \rrbracket_V$

- $\text{dom}(\theta_n) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H) \cup \{a\}$

We know that  $\text{dom}(\theta_n) = \text{dom}(\theta_e) \cup \{a\}$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \llbracket \theta_n(a) \rrbracket_V$ :

Since from (FU-R2) we know that  $(\theta_h, n - l, v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V$

Since  $\theta_h \sqsubseteq \theta_n$  and  $k - j - 1 < n - l$  (since  $k < n$  and  $l < j$ ) therefore from Lemma 2.16 we know that  $(\theta_n, k - j - 1, v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V$

- (b)  $(\theta_n, k - j - 1, a) \in [(\text{ref } \ell' \tau)]_V$ :  
 From Definition 2.6 it suffices to prove that  $\theta_n(a) = \text{Labeled } \ell' \tau$   
 We get this by construction of  $\theta_n$
- (c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :  
 Holds vacuously
- (d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :  
 From SLIO\*-ref we know that  $\ell \sqsubseteq \ell'$

15. SLIO\*-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (!e') \delta) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E$

This means that from Definition 2.7 we need to prove

$\forall i < n. !e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

(From SLIO\*-Sem-val we know that  $v = !e' \delta$  and  $i = 0$ )

This means that given some  $i < n$  s.t  $!e' \delta \Downarrow_i !e' \delta$

It suffices to prove

$(\theta, n, !e' \delta) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

From Definition 2.6 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell') \quad (\text{FU-D0})$

IH:

$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau) \sigma]_E$

From Definition 2.7 this means we have

$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau) \sigma]_V$

Since we know that  $(H, !e') \Downarrow_j^f (H', a)$  therefore from SLIO\*-Sem-deref we know that

$\exists l < j < k$  s.t  $e' \delta \Downarrow_l v_h, v_h = a$

Therefore we have

$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau) \sigma]_V \quad (\text{FU-D1})$

In order to prove (FU-D0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 2.8 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

And since  $H' = H$  (from SLIO\*-Sem-deref) so we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since  $H' = H$  and from Lemma 2.16 we get

$$\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b)  $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau)]_V$ :

From SLIO\*-Sem-deref we know that  $H = H'$  and  $v' = H(a)$

From (FU-D1) and Definition 2.6 we know that  $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from SLIO\*-Sem-deref we know that  $j \geq 1$ . Therefore from Lemma 2.16 we get

$$(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau)]_V$$

(c)  $(\forall a. H(a) \neq H'(a)) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell'$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

16. SLIO\*-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \text{SLIO } \ell \ell \text{ unit}}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [(\text{SLIO } \ell \ell \text{ unit}) \sigma]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\text{SLIO } \ell \ell \text{ unit}) \sigma]_V$$

This means that given some  $i < n$  s.t  $(e_1 := e_2) \delta \Downarrow_i v$ .

It suffices to prove

$$(\theta, n - i, ()) \in [(\text{SLIO } \ell \ell \text{ unit}) \sigma]_V$$

From Definition 2.6 it suffices to prove

$$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from SLIO\*-Sem-assign we know that  $v' = ()$

It suffices to prove

$$\begin{aligned}
& \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, ()) \in [\text{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists l < j < k$  s.t  $e_1 \delta \Downarrow_l a$ . This means we have

$$(\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k-l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k-l-m, v_2) \in [\text{Labeled } \ell' \tau \sigma]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists m < j-l$  (since  $j < k$  therefore  $j-l < k-l$ ) s.t  $e_2 \delta \Downarrow_m v_2$ . This means we have

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k-j, H') \triangleright \theta_e$ :

From Definition 2.8 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H)$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

$\forall a \in \text{dom}(\theta_e)$ .

i.  $H(a) = H'(a)$ :

Since  $(k, H) \triangleright \theta_e$  therefore from Definition 2.8 we know that

$$(\theta_e, k-1, H(a)) \in [\theta_e(a)]_V$$

Therefore from Lemma 2.16 we get

$$(\theta_e, k-1-j, H(a)) \in [\theta_e(a)]_V$$

ii.  $H(a) \neq H'(a)$ :

From SLIO\*-Sem-assign we know that  $H'(a) = v_2$

From (FU-A1) we know that  $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that  $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore we get

$$(\theta, k-j+1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore from Lemma 2.16 we get

$$(\theta, k-j-1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

(b)  $(\theta_e, k-j-1, ()) \in [\text{unit}]_V$ :

From Definition 2.6

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$ :

From SLIO\*-assign we know that  $\ell \sqsubseteq \ell'$

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

17. SLIO\*-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{Lb}(e') \delta) \in [\text{Labeled } \ell \tau \sigma]_E$

This means that from Definition 2.7 we need to prove

$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

This means we are given some  $i < n$  s.t  $\text{Lb}(e') \delta \Downarrow_i v$  and we are required to prove

$(\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

Let  $v = \text{Lb}(v_i)$ . This means from Definition 2.6 we are required to prove

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

IH:  $(\theta, n, e' \delta) \in [\tau \sigma]_E$

This means from Definition 2.7 we have

$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau \sigma]_V$

Since we know that  $\text{Lb}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_i$

Therefore we have  $(\theta, n - j, v_i) \in [\tau \sigma]_V$

From SLIO\*-Sem-label we know that  $i = j + 1$  therefore from Lemma 2.16 we have

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

18. SLIO\*-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{unlabel}(e') \delta) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E$

This means that from Definition 2.7 we need to prove

$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \text{unlabel}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$$

From Definition 2.6 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from SLIO\*-Sem-unlabel we know that  $H' = H$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-U0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell \tau) \sigma]_E$$

This means that from Definition 2.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from SLIO\*-Sem-unlabel we know that

$$\exists h_1 < j < k \text{ s.t. } e' \delta \Downarrow_{h_1} \text{Lb } v'$$

This means we have

$$(\theta_e, k - h_1, \text{Lb } v') \in [(\text{Labeled } \ell \tau) \sigma]_V$$

This means from Definition 2.6 we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-U1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.20 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-U1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 2.16 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :

Holds vacuously

19. SLIO\*-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \text{SLIO } \ell_i \ell_i \tau}$$



Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{ret}(e') \delta) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_V$$

This means we are given some  $i < n$  s.t  $\text{ret}(e') \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_V$$

(from SLIO\*-Sem-val we know that  $v = \text{ret}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{ret}(e') \delta) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_V$$

From Definition 2.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
Also from SLIO\*-Sem-ret we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [\tau \sigma]_E$$

This means that from Definition 2.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [\tau \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from SLIO\*-Sem-ret we know that

$$\exists h_1 < j < k \text{ s.t } e' \delta \Downarrow_{h_1} v'$$

This means we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.20 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-R1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 2.16 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Holds vacuously

- (d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :  
Holds vacuously

20. SLIO\*-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{SLIO } \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \text{SLIO } \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \text{SLIO } \ell_i \ell_o \tau'}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_V$$

This means we are given some  $i < n$  s.t  $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_V$$

(from SLIO\*-Sem-val we know that  $v = \text{bind}(e_1, x.e_2) \delta$  and  $i = 0$ )

Therefore we need to prove

$$(\theta, n, v) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_V$$

From Definition 2.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_E$$

This means that from Definition 2.7 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore from SLIO\*-Sem-bind we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_V$$

From Definition 2.6 we know that

$$\begin{aligned}
& \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h_1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\
& \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell)
\end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.20 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned}
& \exists \theta'' \sqsupseteq \theta_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell) \quad (\text{FU-B1})
\end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in \llbracket (\text{SLIO } \ell_i \ell \tau') \sigma \rrbracket_E$$

This means that from Definition 2.7 we need to prove

$$\forall h_2 < k - h_1 - J. e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in \llbracket (\text{SLIO } \ell \ell_o \tau') \sigma \rrbracket_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$  therefore from SLIO\*-Sem-bind we know that

$$\exists h_2 < j - h_1 - J < k - h_1 - J \text{ s.t } e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in \llbracket (\text{SLIO } \ell \ell_o \tau') \sigma \rrbracket_V$$

From Definition 2.6 we know that

$$\begin{aligned}
& \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell)
\end{aligned}$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists v_{h_2}, i$  s.t  $(v'' \Downarrow_i v_{h_2})$ . From SLIO\*-Sem-val we know that  $v_{h_2} = v''$  and  $i = 0$ . Instantiating  $k_{h_2}$  with  $k - h_1 - J - h_2$ ,  $\theta'_e$  with  $\theta''$ ,  $H$  with  $H'$  (from FU-B1) and  $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$  s.t  $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$ . And since we already know that  $(k - h_1, H') \triangleright \theta''$  therefore from Lemma 2.20 we get  $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned}
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell) \quad (\text{FU-B2})
\end{aligned}$$

We get (FU-B0) by choosing  $\theta'$  as  $\theta''$  (from FU-B2)

21. SLIO\*-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{SLIO } \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{toLabeled}(e') \delta) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_E$

This means that from Definition 2.7 we need to prove

$$\forall i < n. \text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \delta \Downarrow_i v$

(from SLIO\*-Sem-val we know that  $v = \text{toLabeled}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{toLabeled}(e') \delta) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$$

From Definition 2.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

And given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

Also from SLIO\*-Sem-tolabeled we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-TL0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_E$$

This means that from Definition 2.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_V$$

Since  $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$  therefore from SLIO\*-Sem-tolabeled we know that  $\exists h_1 < j < k$  s.t  $e' \delta \Downarrow_{h_1} v_1$

Therefore we get  $(\theta, k - h_1, v_1) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_V$

From Definition 2.6 we know that

$$\begin{aligned} \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J. (k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ \exists \theta'' \sqsupseteq \theta'_e. (k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell) \end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1$ ,  $H_h$  with  $H$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{toLabeled}(e')) \Downarrow_j^f (H', v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 2.20 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} & \exists \theta'' \sqsupseteq \theta'_e.(k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell) \quad (\text{FU-TL1}) \end{aligned}$$

In order to prove (FU-TL0) we choose  $\theta'$  as  $\theta''$ . Now we need to prove the following

(a)  $(k - j, H') \triangleright \theta''$ :

Since  $(k - h_1 - J, H') \triangleright \theta''$  and  $j = h_1 + J + 1$  therefore from Lemma 2.20 we get  $(k - j, H') \triangleright \theta''$

(b)  $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau \sigma)]_V$ :

From SLIO\*-Sem-tolabeled we know that  $v' = \text{toLabeled}(v_1)$

From Definition 2.4 it suffices to prove that  $(\theta'', k - j - 1, v_1) \in [\tau \sigma]_V$

We get this from (FU-TL1) and Lemma 2.16

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Directly from (FU-TL1)

(d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :

Directly from (FU-TL1)

□

**Lemma 2.23** (SLIO\*: Subtyping unary). *The following holds:*

$$\forall \Sigma, \Psi, \sigma, \tau, \tau'.$$

$$1. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$$

$$2. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E \subseteq [(\tau' \sigma)]_E$$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. SLIO\*sub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

IH1:  $[(\tau'_1 \sigma)]_V \subseteq [(\tau_1 \sigma)]_V$  (Statement (1))

$[(\tau_2 \sigma)]_E \subseteq [(\tau'_2 \sigma)]_E$  (Sub-A0, From Statement (2))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V. (\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V$

Therefore from Definition 2.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in [\tau_1 \sigma]_V \implies (\theta_1, i, e_i[v/x]) \in [\tau_2 \sigma]_E \quad (79)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

Again from Definition 2.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in [\tau'_1 \sigma]_V \implies (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$

Since  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in [\tau_1 \sigma]_V$

As a result from Equation 79 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2 \sigma]_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

## 2. SLIO\*sub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[((\tau_1 \times \tau_2) \sigma)]_V \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V$

IH1:  $[(\tau_1 \sigma)]_V \subseteq [(\tau'_1 \sigma)]_V$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_V \subseteq [(\tau'_2 \sigma)]_V$  (Statement (1))

It suffices to prove:  $\forall (\theta, n, (v_1, v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V. (\theta, n, (v_1, v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V$

Therefore from Definition 2.6 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \tag{80}$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V$

Again from Definition 2.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 80 we know that  $(\theta, n, v_1) \in [\tau_1 \sigma]_V$  therefore from IH1 we have  $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since  $(\theta, n, v_2) \in [\tau_2 \sigma]_V$  from Equation 80 therefore from IH2 we have  $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

3. SLIO\*sub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V$  (Statement (1))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:  $\forall (\theta, n, v_s) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V. (\theta, v_s) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

This means that given:  $(\theta, n, v_s) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V$

And it suffices to prove:  $(\theta, n, v_s) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_1 \sigma \rrbracket_V \tag{81}$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_1 \sigma \rrbracket_V$$

From Equation 81 and IH1 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_1 \sigma \rrbracket_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_2 \sigma \rrbracket_V \tag{82}$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_2 \sigma \rrbracket_V$$

From Equation 82 and IH2 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_2 \sigma \rrbracket_V$$

4. SLIO\*sub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $\llbracket ((\forall \alpha. \tau_1) \sigma) \rrbracket_V \subseteq \llbracket ((\forall \alpha. \tau_2) \sigma) \rrbracket_V$

It suffices to prove:  $\forall (\theta, n, \Lambda e_i) \in \llbracket ((\forall \alpha. \tau_1) \sigma) \rrbracket_V. (\theta, n, \Lambda e_i) \in \llbracket ((\forall \alpha. \tau_2) \sigma) \rrbracket_V$

This means that given:  $(\theta, n, \Lambda e_i) \in \llbracket ((\forall \alpha. \tau_1) \sigma) \rrbracket_V$

Therefore from Definition 2.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E \quad (83)$$

And it suffices to prove:  $(\theta, n, \Lambda e_i) \in [(\forall \alpha. \tau_2) \sigma]_V$

Again from Definition 2.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

This means that given some  $\theta_2, j < n, \ell' \in \mathcal{L}$  s.t  $\theta \sqsubseteq \theta_2$

And we are required to prove:  $(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$  therefore from Equation 83 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell']))]_E \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell']))]_E \text{ (Sub-F0, Statement (2))}$$

From (Sub-F0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

#### 5. SLIO\*sub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $[((c_1 \Rightarrow \tau_1) \sigma)]_V \subseteq [((c_2 \Rightarrow \tau_2) \sigma)]_V$

It suffices to prove:  $\forall (\theta, n, \nu e_i) \in [((c_1 \Rightarrow \tau_1) \sigma)]_V. (\theta, n, \nu e_i) \in [((c_2 \Rightarrow \tau_2) \sigma)]_V$

This means that given:  $(\theta, n, \nu e_i) \in [((c_1 \Rightarrow \tau_1) \sigma)]_V$

Therefore from Definition 2.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E \quad (84)$$

And it suffices to prove:  $(\theta, n, \nu e_i) \in [((c_2 \Rightarrow \tau_2) \sigma)]_V$

Again from Definition 2.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$

This means that given some  $\theta_2, j$  s.t  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

And we are required to prove:  $(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$  and  $\mathcal{L} \models c_2 \sigma \implies c_1 \sigma$  therefore from Equation 84 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E$$

$$[(\tau_1 \sigma)]_E \subseteq [(\tau_2 \sigma)]_E \text{ (Sub-C0, Statement (2))}$$

From (Sub-C0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$



6. SLIO\*sub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V \subseteq \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V$

IH:  $\llbracket (\tau \sigma) \rrbracket_V \subseteq \llbracket (\tau' \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:

$\forall (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V. (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V$

This means that given some  $\theta, n$  and  $\text{Lb}(e_i)$  s.t  $(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V$

Therefore from Definition 2.6 we are given:

$(\theta, n, v_i) \in \llbracket (\tau \sigma) \rrbracket_V$  (SL)

And we are required to prove that

$(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V$

From Definition 2.6 it suffices to prove

$(\theta, n, v_i) \in \llbracket (\tau' \sigma) \rrbracket_V$

We get this directly from (SL) and IH

7. SLIO\*sub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \text{SLIO } \ell_i \ell_o \tau <: \text{SLIO } \ell'_i \ell'_o \tau'}$$

To prove:  $\llbracket ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rrbracket_V \subseteq \llbracket ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rrbracket_V$

IH:  $\llbracket (\tau \sigma) \rrbracket_V \subseteq \llbracket (\tau' \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:

$\forall (\theta, n, e) \in \llbracket ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rrbracket_V. (\theta, n, e) \in \llbracket ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rrbracket_V$

This means that given some  $\theta, n$  and  $e$  s.t  $(\theta, n, e) \in \llbracket ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rrbracket_V$

Therefore from Definition 2.6 we are given:

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \sigma \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \quad (\text{SC0})$

And we are required to prove

$(\theta, n, e) \in \llbracket ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rrbracket_V$

So again from Definition 2.6 we need to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau' \sigma \rrbracket_V \wedge$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$   
(SC1)

And we need to prove

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$$

We instantiate (SC0) with  $k, \theta_e, H, j$  from (SC1) and we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$$

Since  $\tau \sigma <: \tau' \sigma$  therefore from IH we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \sigma \rfloor_V$$

And since  $\ell'_i \sqsubseteq \ell_i$  therefore we also have

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$$

## 8. SLIO\* sub-base:

Trivial

### Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E. (\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

This means that we are given  $(\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E$

From Definition 2.7 it means we have

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

From Definition 2.7 we need to prove

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

This further means that given some  $i < n$  s.t.  $e \Downarrow_i v$ , it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

Instantiating (Sub-E0) with the given  $i$  we get  $(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$

Finally from Statement(1) we get  $(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$

□

**Lemma 2.24** (SLIO\*: Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in \lfloor \Gamma \rfloor_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \Downarrow_i) \in \lfloor \Gamma \rfloor_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 2.14 we know that we are given:

$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

And we are required to prove:

$\forall i \in \{1, 2\}. \forall m.$

$dom(\Gamma) \subseteq dom(\gamma \downarrow_i) \wedge \forall x \in dom(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 1$

Given some  $m$  we need to show:

- $dom(\Gamma) \subseteq dom(\gamma \downarrow_1)$ :

$$dom(\gamma) = dom(\gamma \downarrow_1)$$

Therefore,  $dom(\Gamma) \subseteq (dom(\gamma) = dom(\gamma \downarrow_1))$  (Given)

- $\forall x \in dom(\Gamma). (W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$ :

We are given:  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

Therefore from Lemma 2.15 we know that

$$\forall m'. (W.\theta_1, m', \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Instantiating  $m'$  with  $m$  we get

$$(W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Case  $i = 2$

Symmetric reasoning as in the  $i = 1$  case above

□

**Theorem 2.25** (SLIO\*: Fundamental theorem binary).  $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$

$$\Sigma; \Psi; \Gamma \vdash e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_V^A \implies$$

$$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

*Proof.* Proof by induction on the typing derivation

1. SLIO\*-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{SLIO}^*\text{-var}$$

To prove:  $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

Say  $e_1 = x(\gamma \downarrow_1)$  and  $e_2 = x(\gamma \downarrow_2)$

From Definition 2.5 it suffices to prove that

$$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some  $i < n$  s.t  $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2$

We are required to prove:  $(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

From SLIO\*-Sem-val we know that  $x (\gamma \downarrow_1) \Downarrow x (\gamma \downarrow_1)$  and  $x (\gamma \downarrow_2) \Downarrow x (\gamma \downarrow_2)$

This means  $v'_1 = x (\gamma \downarrow_1)$  and  $v'_2 = x (\gamma \downarrow_2)$

Since  $(W, n, \gamma) \in [\tau]_V^A$ . Therefore from Definition 2.14 we know that

$$(W, n, v'_1, v'_2) \in [\tau]_V^A$$

From Lemma 2.17 we get

$$(W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

2. SLIO\*-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove:  $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

Say  $e_1 = \lambda x. e (\gamma \downarrow_1)$  and  $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of  $[(\tau_1 \rightarrow \tau_2) \sigma]_E^A$  it suffices to prove that

$$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some  $i < n$  s.t  $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From SLIO\*-Sem-val we know that  $v'_1 = (\lambda x. e_i) \gamma \downarrow_1$  and  $v'_2 = (\lambda x. e_i) \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\lambda x. e_i) \gamma \downarrow_1, (\lambda x. e_i) \gamma \downarrow_2) \in [\tau]_V^A$$

From Definition 2.4 it suffices to prove

$$\forall W' \sqsupseteq W, j < n, v_1, v_2.$$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, v_c, j.$$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, v_c, j.$$

$$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E) \quad (\text{FB-L0})$$

IH:

$$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in [\Gamma]_V^A$$

In order to prove (FB-L0) we need to prove the following:

$$(a) \forall W' \sqsupseteq W, j < n, v_1, v_2.$$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A):$$

This means given some  $W' \sqsupseteq W, j < n, v_1, v_2$  s.t.  $(W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove  $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A$

We get this by instantiating IH with  $W'$  and  $j$

- (b)  $\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$   
 $((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E):$   
This means given some  $\theta_l \sqsupseteq W.\theta_1, v_c, j$  s.t  $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$   
We need to prove:  $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E$

It is given to us that

$$(W, n, \gamma) \in [\Gamma]_V^A$$

Therefore from Lemma 2.24 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating  $m$  with  $j$  we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

From Lemma 2.19 we know that

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Since we know that  $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

Therefore we also have

$$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1 \sigma\}]_V$$

Therefore, we can apply Theorem 2.22 to obtain

$$(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_V$$

- (c)  $\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$   
 $((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E):$   
Similar reasoning as in the previous case

### 3. SLIO\*-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

This further means that given some  $i < n$  s.t  $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

$$\underline{\text{IH1:}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e_1 \gamma \downarrow_2 \downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$

From SLIO\*-Sem-app we know that  $val_{h_1} = \lambda x.e_{h_1}$  and  $val_{h_2} = \lambda x.e_{h_2}$

From Definition 2.4 this further means

$$\begin{aligned}
& \forall W' \sqsupseteq W, J < (n - j), v_1, v_2. \\
& ((W', J, v_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}}^A \implies (W', J, e_{h_1}[v_1/x], e_{h_2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_1, v_c, j. \\
& ((\theta_l, j, v_c) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2 \sigma]_E) \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_2, v_c, j. \\
& ((\theta_l, j, v_c) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \quad (\text{FB-A1})
\end{aligned}$$

$$\underline{\text{IH2}}: (W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall k < n - j.e_2 \gamma \downarrow_1 \downarrow_j v_{h_1'} \wedge e_2 \gamma \downarrow_2 \downarrow v_{h_2'} \implies (W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \downarrow_k v_{h_1'}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \downarrow v_{h_2'}$

$$\text{This means we have } (W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (\text{FB-A2})$$

Instantiating the first conjunct of (FB-A1) as follows  $W'$  with  $W$ ,  $J$  with  $n - j - k$ ,  $v_1$  and  $v_2$  with  $v'_{h_1}$  and  $v'_{h_2}$  respectively, we obtain

$$(W, n - j - k, e_{h_1}[v'_{h_1}/x], e_{h_2}[v'_{h_2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 2.5

$$\forall l < n - j - k.(e_{h_1}[v'_{h_1}/x]) \gamma \downarrow_l v_{f_1} \wedge e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2} \implies (W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$ . Therefore  $\exists l < i - j - k < n - j - k$  s.t  $e_{h_1}[v'_{h_1}/x] \downarrow_l v_{f_1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f_2}$  therefore  $e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2}$

$$\text{Therefore we have } (W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since  $i = j + k + l$  therefore we are done

#### 4. SLIO\*-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

$$\text{To prove: } (W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall i < n.(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2}) \implies (W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2})$

We are required to prove

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A \quad (\text{FB-P0})$$

$$\underline{\text{IH1}}: (W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge e_1 \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathbb{V}}^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \Downarrow_j v_{f_1}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{f_1}$

This means we have

$$(W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathbb{V}}^A \quad (\text{FB-P1})$$

$$\underline{\text{IH2}}: (W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \Downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathbb{V}}^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \Downarrow_j v_{f_2}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \Downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathbb{V}}^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 2.4 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathbb{V}}^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathbb{V}}^A$$

Since  $i = j + k + 1$  therefore from (FB-P1) and (FB-P2) and from Lemma 2.17 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathbb{V}}^A$$

5. SLIO\*-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove:  $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathbb{V}}^A$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1 \sigma]_{\mathbb{V}}^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathbb{V}}^A$$

Since we know that  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$  therefore  $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f1}, v_{f2}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

From Definition 2.4 we know that

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

Since from SLIO\*-Sem-fst  $i = j + 1$  therefore from Lemma 2.17 we get

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

6. SLIO\*-snd:

Symmetric reasoning as in the SLIO\*-fst case above

7. SLIO\*-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove:  $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{inl}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1}) \implies \\ (W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we have:

$$\forall j < n. e' \gamma \downarrow_1 \downarrow_i v_{f1} \wedge e' \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

Since we know that  $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \downarrow_j v_{f1}$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$  therefore  $e' \gamma \downarrow_2 \downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 2.4 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

From SLIO\*-Sem-inl since  $i = j + 1$  therefore from (FB-IL1) and Lemma 2.17 we get (FB-IL0)

8. SLIO\*-inr:

Symmetric reasoning as in the SLIO\*-inl case above



9. SLIO\*-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove:  $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2}$

We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_j v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$  therefore  $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-C1})$$

2 cases arise

- (a)  $v_{h1} = \text{inl}(v_1)$  and  $v'_{h1} = \text{inl}(v'_1)$ :

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.5 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

From SLIO\*-Sem-case1 we know that  $i = j + k + 1$  therefore from Lemma 2.17 we get (FB-C0)

- (b)  $v_{h1} = \text{inr}(v_1)$  and  $v'_{h1} = \text{inr}(v'_1)$ :

Symmetric case

10. SLIO\*-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

To prove:  $(W, n, \Lambda e' (\gamma \downarrow_1), \Lambda e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 2.5 it suffices to prove that

$$\forall i < n. (\Lambda e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

This means given some  $i < n$  s.t.  $(\Lambda e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \downarrow v_{f2}$

From SLIO\*-Sem-val we know that  $v_{f1} = (\Lambda e') \gamma \downarrow_1$  and  $v_{f2} = (\Lambda e') \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\Lambda e') \gamma \downarrow_1, (\Lambda e') \gamma \downarrow_2) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Let  $e_1 = (\Lambda e') \gamma \downarrow_1$  and  $e_2 = (\Lambda e') \gamma \downarrow_2$

From Definition 2.4 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FI0}) \end{aligned}$$

IH:  $\forall W, n. (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E^A$

In order to prove (FB-FI0) we need to prove the following

(a)  $\forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A)$ :

This means given  $W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}$  and we are required to prove

$$(W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

(b)  $\forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$ :

This means given  $\theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j$  and we are required to prove

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

Since from Lemma 2.24

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 2.17 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 2.22 to get

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E$ :

Symmetric reasoning as before

11. SLIO\*-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

To prove:  $(W, n, e' [] (\gamma \downarrow_1), e' [] (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 2.5 it suffices to prove that

$$\forall i < n. (e' []) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e' []) \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

This means given some  $i < n$  s.t.  $(e' []) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e' []) \gamma \downarrow_2 \Downarrow v_{f2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE0})$$

IH:  $(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 2.5 it suffices to prove that

$$\forall i < n. (e' \gamma \downarrow_1 \Downarrow_i v_{h1} \wedge (e' \gamma \downarrow_2 \Downarrow v_{h2}) \implies (W, n - i, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Since we know that  $(e' []) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t.  $e' \gamma \downarrow_1 \Downarrow_j v_{h1}$ . Similarly since  $(e' []) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e' \gamma \downarrow_2 \Downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$

From SLIO\*-Sem-FE we know that  $v_{h1} = \Lambda e_{h1}$  and  $v_{h2} = \Lambda e_{h2}$

From Definition 2.4 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, k < (n - j), \ell' \in \mathcal{L}. ((W', k, e_{h1}, e_{h2}) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h1}) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h2}) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FE1}) \end{aligned}$$

Instantiating the first conjunct of (FB-FE1) with  $W, n - j - 1$  and  $\ell$  we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau[\ell/\alpha] \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall l < n - j - 1. (e_{h1}) \Downarrow_l v_{f1} \wedge e_{h2} \Downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

Since we know that  $(e' []) \gamma \downarrow_1 \Downarrow_i v_{f1}$  therefore from SLIO\*-Sem-FE we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t.  $e_{h1} \gamma \downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $(e' []) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_{h2} \gamma \downarrow_2 \Downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-FE2) we get (FB-FE0)

12. SLIO\*-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

To prove:  $(W, n, \nu e' (\gamma \downarrow_1), \nu e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 2.5 it suffices to prove that

$$\forall i < n. (\nu e')\gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in \llbracket (c \Rightarrow \tau) \sigma \rrbracket_V^A$$

This means given some  $i < n$  s.t  $(\nu e')\gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \Downarrow v_{f2}$

From SLIO\*-Sem-val we know that  $v_{f1} = (\nu e')\gamma \downarrow_1$  and  $v_{f2} = (\nu e')\gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\nu e')\gamma \downarrow_1, (\nu e')\gamma \downarrow_2) \in \llbracket (c \Rightarrow \tau) \sigma \rrbracket_V^A$$

Let  $e_1 = (\nu e')\gamma \downarrow_1$  and  $e_2 = (\nu e')\gamma \downarrow_2$

From Definition 2.4 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c &\implies (W', j, e_1, e_2) \in \llbracket \tau \sigma \rrbracket_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c &\implies (\theta_l, j, e_1) \in \llbracket \tau \sigma \rrbracket_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c &\implies (\theta_l, j, e_2) \in \llbracket \tau \sigma \rrbracket_E \quad (\text{FB-CI0}) \end{aligned}$$

$$\underline{\text{IH}}: \forall W, n. (W, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in \llbracket \tau \sigma \rrbracket_E^A$$

In order to prove (FB-CI0) we need to prove the following

$$(a) \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \sigma \implies (W', j, e_1, e_2) \in \llbracket \tau \sigma \rrbracket_E^A:$$

This means given  $W' \sqsupseteq W, j < n, \mathcal{L} \models c \sigma$  and we are required to prove

$$(W', j, e_1, e_2) \in \llbracket \tau \sigma \rrbracket_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

$$(b) \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma \implies (\theta_l, j, e_1) \in \llbracket \tau \sigma \rrbracket_E:$$

This means given  $\theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma$  and we are required to prove

$$(\theta_l, j, e_1) \in \llbracket \tau \sigma \rrbracket_E$$

Since from Lemma 2.24  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in \llbracket \Gamma \rrbracket_V$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$$

And from Lemma 2.17 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$$

Therefore we can apply Theorem 2.22 to get

$$(\theta_l, j, e_1) \in \llbracket \tau \sigma \rrbracket_E$$

$$(c) \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in \llbracket \tau \sigma \rrbracket_E:$$

Symmetric reasoning as before

### 13. SLIO\*-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

To prove:  $(W, n, e' \bullet (\gamma \downarrow_1), e' \bullet (\gamma \downarrow_2)) \in \llbracket \tau \sigma \rrbracket_E^A$

From Definition 2.5 it suffices to prove that

$$\forall i < n. (e' \bullet)\gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e' \bullet)\gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in \llbracket \tau \sigma \rrbracket_V^A$$

This means given some  $i < n$  s.t  $(e' \bullet) \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge (e' \bullet) \gamma \downarrow_2 \Downarrow v_{f_2}$

We are required to prove:

$$(W, n - i, v_{f_1}, v_{f_2}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-CE0})$$

$$\underline{\text{IH}}: (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$$

From Definition 2.5 it suffices to prove that

$$\forall i < n. e' \gamma \downarrow_1 \Downarrow_i v_{h_1} \wedge e' \gamma \downarrow_2 \Downarrow v_{h_2} \implies (W, n - i, v_{h_1}, v_{h_2}) \in [(c \Rightarrow \tau) \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e' \bullet) \gamma \downarrow_1 \Downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j v_{h_1}$ . Similarly since  $(e' \bullet) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e' \gamma \downarrow_2 \Downarrow v_{h_2}$

$$\text{This means we have } (W, n - j, v_{h_1}, v_{h_2}) \in [(c \Rightarrow \tau) \sigma]_{\mathcal{V}}^A$$

From SLIO\*-Sem-CE we know that  $v_{h_1} = \nu e_{h_1}$  and  $v_{h_2} = \nu e_{h_2}$

From Definition 2.4 this further means

$$\forall W' \sqsupseteq W, k < n - j. \mathcal{L} \models c \sigma \implies (W', k, e_1, e_2) \in [\tau \sigma]_E^A \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_1) \in [\tau \sigma]_E \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_2) \in [\tau \sigma]_E \quad (\text{FB-CE1})$$

Instantiating the first conjunct of (FB-CE1) with  $W, n - j - 1$  and since we know that  $\mathcal{L} \models c \sigma$  therefore we get

$$(W, n - j - 1, e_{h_1}, e_{h_2}) \in [\tau \sigma]_E^A$$

This means from Definition 2.5 we know that

$$\forall l < n - j - 1. (e_{h_1}) \Downarrow_l v_{f_1} \wedge e_{h_2} \Downarrow v_{f_2} \implies (W, n - j - 1 - l, v_{f_1}, v_{f_2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e' \bullet) \gamma \downarrow_1 \Downarrow_i v_{f_1}$  therefore from SLIO\*-Sem-CE we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t  $e_{h_1} \gamma \downarrow_1 \Downarrow_l v_{f_1}$ . Similarly since  $(e' \bullet) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e_{h_2} \gamma \downarrow_2 \Downarrow v_{f_2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f_1}, v_{f_2}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-CE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-CE2) we get (FB-CE0)

14. SLIO\*-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove:  $(W, n, \text{Lb}(e') (\gamma \downarrow_1), \text{Lb}(e') (\gamma \downarrow_2)) \in [\text{Labeled } \ell \tau \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. \text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f_1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f_1}) \implies$$

$$(W, n - i, \text{Lb}(v_{f_1}), \text{Lb}(v'_{f_1})) \in [\text{Labeled } \ell \tau \sigma]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f_1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f_1})$

We are required to prove

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\text{Labeled } \ell \tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.5 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - j, v_{f_1}, v'_{f_1}) \in [\tau \sigma]_V^A$$

Since we know that  $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f_1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j v_{f_1}$ . Similarly since  $\text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f_1})$  therefore  $e' \gamma \downarrow_2 \Downarrow v'_{f_1}$

This means we have

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau \sigma]_V^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 2.4 it suffices to prove that

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau \sigma]_V^A$$

From SLIO\*-Sem-label we know that  $i = j+1$ . Therefore we get the desired from (FB-LB1) and Lemma 2.17

15. SLIO\*-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau}$$

To prove:  $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

From SLIO\*-Sem-val we know that  $v_{f_1} = \text{unlabel}(e') \gamma \downarrow_1$  and  $v'_{f_1} = \text{unlabel}(e') \gamma \downarrow_2$ . Also  $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\text{SLIO } \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A$$

This means from Definition 2.4 we need to prove

Let  $e_1 = \text{unlabel}(e') \gamma \downarrow_1$  and  $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau']_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

We need to show

- (a)  $\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2.$   
 $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$   
 $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma):$

Also given is some  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{FB-U0})$$

$$\underline{\text{IH:}} (W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{Labeled } \ell \tau) \sigma]_E^A$$

This means from Definition 2.5 we are given

$$\forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) \implies$$

$$(W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau) \sigma]_V^A$$

Since we know that

$$(H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \text{ therefore}$$

$$\exists I < j < k \text{ s.t } e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1})$$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau) \sigma]_V^A$$

This means from Definition 2.4 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \sigma, v_{h1}, v'_{h1}, \tau \sigma) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose  $W'$  as  $W_e$  and from SLIO\*-Sem-unlabel we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . And we already know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore from Lemma 2.21 we get  $(k - j, H_1, H_2) \triangleright W_e$

From SLIO\*-Sem-unlabel we know that  $v'_1, v'_2$  in (FB-U0) is  $v_{h1}, v'_{h1}$  respectively. And since from (FB-U1) we know that  $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \sigma, v_{h1}, v'_{h1}, \tau \sigma)$ . Therefore from Lemma 2.26 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, (\ell_i \sqcup \ell) \sigma, v_{h1}, v'_{h1}, \tau \sigma)$$

- (b)  $\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right.$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $\left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right):$

Case  $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W. \theta_l, H, j \text{ s.t } (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k$$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that

$$\forall m. (W. \theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W. \theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W. \theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.22 to get

$$(W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \sqcup \ell \tau) \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c < k. (\text{unlabel } e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\text{SLIO } \ell_i \ell_i \sqcup \ell \tau) \sigma]_V$$

This further means that given some  $c < k$  s.t  $(\text{unlabel } e')\gamma \downarrow_1 \downarrow_c v$ . From SLIO\*-Sem-val we know that  $c = 0$  and  $v = (\text{unlabel } e')\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \sqcup \ell \tau) \sigma]_V$$

From Definition 2.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K &\implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge & \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1) & \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

16. SLIO\*-tolabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{SLIO } \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

To prove:  $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \downarrow v'_{f1}$

From SLIO\*-Sem-val we know that  $v_{f1} = \text{toLabeled}(e') \gamma \downarrow_1$ ,  $v_{f2} = \text{toLabeled}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Let  $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$  and  $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 2.4 we are required to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. & \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies & \\ \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma) & \Big) \wedge \\ \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies & \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge & \end{aligned}$$



$$\begin{aligned}
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)
\end{aligned}$$

We need to prove:

$$\begin{aligned}
\text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\
& (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell_2, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma):
\end{aligned}$$

This means that we are given some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2, v'_1, v'_2, j < k$  s.t  
 $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell_o, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma) \quad (\text{FB-TL0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{SLIO } \ell_i \ell_o \tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, n-J, v_{h1}, v'_{h1}) \in [\text{SLIO } \ell_i \ell_o \tau \sigma]_V^A$$

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j (H'_1, v'_1)$  and  $(H_2, \text{toLabeled}(e')\gamma \downarrow_2) \Downarrow_j (H'_2, v'_2)$ . Therefore from SLIO\*-Sem-val we know that  $\exists J < j < k \leq n$  s.t  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly we also know that  $e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k-J, v_{h1}, v'_{h1}) \in [\text{SLIO } \ell_i \ell_o \tau \sigma]_V^A$$

From Definition 2.4 we know that

$$\begin{aligned}
& (\forall k_1 \leq (k-J), W_e'' \sqsupseteq W_e. \forall H_1'', H_2''. (k_1, H_1'', H_2'') \triangleright W_e'' \wedge \forall v''_1, v''_2, m. \\
& (H_1'', v_{h1}) \Downarrow_m^f (H'_1, v''_1) \wedge (H_2'', v'_{h1}) \Downarrow^f (H'_2, v''_2) \wedge m < k_1 \implies \\
& \exists W' \sqsupseteq W_e''. (k_1-m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1-m, \ell_o, v''_1, v''_2, \tau \sigma)) \wedge \\
& \forall \ell \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
& \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_l) \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \quad (\text{FB-TL1})
\end{aligned}$$

We instantiate  $W_e''$  with  $W_e$ ,  $H_1''$  with  $H_1$ ,  $H_2''$  with  $H_2$  and  $k_1$  with  $k$  in (FB-TL1).

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$ , therefore  $\exists m < j < k \leq n$  s.t  $(H_1, v_{h1}) \Downarrow_m^f (H'_1, v'_1) \wedge (H_2, v'_{h1}) \Downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W' \sqsupseteq W_e. (k-m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-m, \ell_o, v''_1, v''_2, \tau \sigma) \quad (\text{FB-TL2})$$

In order to prove (FB-TL0) we choose  $W'$  as  $W'$  from (FB-TL2). Since from SLIO\*-Sem-tolabeled we know that  $v'_1 = \text{Lb}_{\ell_o}(v''_1)$ ,  $v'_2 = \text{Lb}_{\ell_o}(v''_2)$  and  $j = m+1$ , therefore from Lemma 2.21 we get  $(k-j, H'_1, H'_2) \triangleright W'$ .

Since we have by assumption that  $\ell_i \sqsubseteq \ell_o$  therefore the following cases arise

i.  $\ell_i \sqsubseteq \ell_o \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$$(W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since  $v'_1 = \text{Lb}_{\ell_o}(v''_1)$  and  $v'_2 = \text{Lb}_{\ell_o}(v''_2)$ . Therefore from Definition 2.4 it suffices to prove that

$$\text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v''_1, v''_2, \tau \sigma)$$

We get this from (FB-TL2) and Lemma 2.26

ii.  $(\ell_i \sqsubseteq \ell_o) \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$$\forall m. (W', m, v'_1) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \text{ and } \forall m. (W', m, v'_2) \in [(\text{Labeled } \ell_o \tau) \sigma]_V$$

Since  $\ell_o \not\sqsubseteq \mathcal{A}$  therefore we get this from (FB-TL2), Definition 2.3 and Definition 2.6

iii.  $(\ell_i \sqsubseteq \mathcal{A} \sqsubseteq \ell_o)$ :

In this case from Definition 2.3 it suffices to prove that

$$(W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since  $v'_1 = \text{Lb}_{\ell_o}(v''_1)$  and  $v'_2 = \text{Lb}_{\ell_o}(v''_2)$ . Therefore from Definition 2.4 it suffices to prove that

$$\forall m. (W', m, v''_1) \in [\tau \sigma]_V \text{ and } \forall m. (W', m, v''_2) \in [\tau \sigma]_V$$

We obtain this directly from (FB-TL2) and Definition 2.3

$$\begin{aligned} \text{(b)} \quad \forall l \in \{1, 2\}. & \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) \right): \end{aligned}$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.22 to get

$$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$$

Instantiating  $c$  with 0 and from SLIO\*-Sem-val we know  $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$

From Definition 2.6 we have

$$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$$

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \llbracket \text{Labeled } \ell_o \tau \rrbracket_V \sigma \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

17. SLIO\*-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \text{SLIO } \ell_i \ell_i \tau}$$

To prove:  $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in \llbracket \text{SLIO } \ell_i \ell_i \tau \sigma \rrbracket_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} & \forall i < n. \text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ & (W, n - i, v_{f1}, v'_{f1}) \in \llbracket \text{SLIO } \ell_i \ell_i \tau \sigma \rrbracket_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From SLIO\*-Sem-val we know that  $v_{f1} = \text{ret}(e') \gamma \downarrow_1$ ,  $v_{f2} = \text{ret}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{ret}(e') \gamma \downarrow_1, \text{ret}(e') \gamma \downarrow_2) \in \llbracket \text{SLIO } \ell_i \ell_i \tau \sigma \rrbracket_V^A$$

Let  $v_1 = \text{ret}(e') \gamma \downarrow_1$  and  $v_2 = \text{ret}(e') \gamma \downarrow_2$

From Definition 2.4 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall v, i. (e_l \Downarrow_i v) \implies \right. \\ & \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \tau \rrbracket_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \end{aligned}$$

It suffices to prove:

$$\begin{aligned} & \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau): \end{aligned}$$

We are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

From SLIO\*-Sem-ret we know that  $H'_1 = H_1$  and  $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e.(k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau \sigma]_V^A$$

Since we know that  $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow^f (H_2, v'_2)$ , therefore  $\exists J < j < k$  s.t  $e' \gamma \downarrow_1 \downarrow_J v_{h1}$  and similarly  $e' \gamma \downarrow_2 \downarrow v'_{h1}$ .

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau \sigma]_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_e$  and from SLIO\*-Sem-ret we know that  $v'_1 = v_{h1}$  and  $v'_2 = v'_{h1}$ . We need to prove the following:

i.  $(k - j, H_1, H_2) \triangleright W_e$ :

Since we have  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.21 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_i, v'_1, v'_2, \tau)$ :

2 cases arise:

A.  $\ell_i \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in [\tau \sigma]_V^A$$

Since  $j = J + 1$  therefore we get this from (FB-R1) and Lemma 2.17

B.  $\ell_i \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in [\tau \sigma]_V \text{ and } \forall m. (W_e, m, v'_2) \in [\tau \sigma]_V$$

We get this From (FB-R1) and Lemma 2.15

$$\begin{aligned} \text{(b) } \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_1, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_1) \Downarrow_j^f (H', v'_1) \wedge j < k$

We need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.22 to get

$$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \tau) \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\text{SLIO } \ell_i \ell_i \tau) \sigma]_V$$

Instantiating  $c$  with 0 and from SLIO\*-Sem-val we know that  $v = (\text{ret } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\text{SLIO } \ell_i \ell_i \tau) \sigma]_V$

From Definition 2.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K &\implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \sigma \wedge & \\ (\forall a.H_1(a) \neq H'(a) \implies \exists \ell'.\theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e).\theta'(a) \searrow \ell_i \sigma) & \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

18. SLIO\*-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \text{SLIO } \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_b : \text{SLIO } \ell_o \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_l, x.e_b) : \text{SLIO } \ell_i \ell_o \tau'}$$

To prove:  $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From SLIO\*-Sem-val we know that  $v_{f1} = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$ ,  $v_{f2} = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b)\gamma \downarrow_1, \text{bind}(e_l, x.e_b)\gamma \downarrow_2) \in [\text{SLIO } \ell_i \ell_o \tau' \sigma]_V^A$$

Let  $v_1 = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$  and  $v_2 = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$

This means from Definition 2.4 we need to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. & \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies & \\ \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \Big) \wedge & \\ \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies & \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge & \\ (\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell_i) & \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad \forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. & \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies & \\ \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma): & \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\text{SLIO } \ell_i \ell \tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{SLIO } \ell_i \ell \tau \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{SLIO } \ell_i \ell \tau \sigma]_V^A$$

This means from Definition 2.4 we have

$$\begin{aligned} & (\forall K \leq (k - f), W'_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W'_e \wedge \forall v''_1, v''_2, J. \\ & (H''_1, v_{h1}) \Downarrow_J^f (H'_1, v''_1) \wedge (H''_2, v'_{h1}) \Downarrow^f (H'_2, v''_2) \wedge J < K \implies \\ & \exists W'' \sqsupseteq W'_e. (K - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell \sigma, v''_1, v''_2, \tau \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)) \end{aligned}$$

Instantiating  $K$  with  $(k - f)$ ,  $W'_e$  with  $W_e$ ,  $H''_1$  with  $H_1$  and  $H''_2$  with  $H_2$  in the first conjunct of the above equation. Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.21 we also have  $(k - f, H_1, H_2) \triangleright W_e$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J < j - f < k - f$  s.t  $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v''_1) \wedge (H_2, v'_{h1}) \Downarrow^f (H'_2, v''_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e. (k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell \sigma, v''_1, v''_2, \tau \sigma) \quad (\text{FB-B1})$$

From Definition 2.3 two cases arise:

i.  $\ell \sigma \sqsubseteq \mathcal{A}$ :

$$\text{In this case we know that } (W'', k - f - J, v''_1, v''_2) \in [\tau \sigma]_V^A$$

IH2:

$$(W'', k - f - J, e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}), e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\})) \in [\text{SLIO } \ell \ell_o \tau' \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall s < k - f - J. e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2} \implies (W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\text{SLIO } \ell \ell_o \tau' \sigma]_V^A$$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f - J < k - f - J$  s.t.  $e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\text{SLIO } \ell \ell_o \tau' \sigma]_V^A$$

This means from Definition 2.4 we know that

$$\left( \forall K_s \leq (k - f - J - s), W_s \sqsupseteq W'' . \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right.$$

$$(H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies$$

$$\left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_i, v'_1, v'_2, \tau' \sigma) \right) \wedge$$

$$\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$$

$$\left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right.$$

$$\left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_1 \sqsubseteq \ell') \wedge \right.$$

$$\left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right)$$

Instantiating  $K_s$  with  $(k - f - J - s)$ ,  $W_s$  with  $W''$ ,  $H_1$  with  $H'_1$  and  $H'_2$  with  $H_2$ . Since we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Lemma 2.21 we also have  $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J_s < j - f - J - s < k - f - J - s$  s.t.  $(H'_1, v'_1) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_2) \Downarrow^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s. (k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma) \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose  $W'$  as  $W'_s$ . From SLIO\*-Sem-bind we know that  $H'_1 = H'_{s1}$ ,  $H'_2 = H'_{s2}$ ,  $v'_1 = v'_{s1}$ ,  $v'_2 = v'_{s2}$  and  $j = f + J + s + J_s + 1$ . And we need to prove:

A.  $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$ :

Since from (FB-B2) we know that  $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$  therefore from Lemma 2.21 we get

$$(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$$

B.  $\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$ :

Since from (FB-B2) we know that  $\text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$  therefore from Lemma 2.26 we get

$$\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$$

ii.  $\ell \sigma \not\sqsubseteq \mathcal{A}$ :

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma)$$

Since  $\ell_i \sigma \sqsubseteq \ell \sigma \sqsubseteq \ell_o \sigma$  (by assumption) and  $\ell \sigma \not\sqsubseteq \mathcal{A}$  therefore we have  $\ell_o \sigma \not\sqsubseteq \mathcal{A}$

This means that from Definition 2.3 it suffices to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'. \theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge \forall m_{u2}. (W'. \theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V$$

This means given some  $m_{u1}, m_{u2}$  and we need to prove

$$\exists W' \sqsupseteq W_e.(k-j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V \quad (\text{FB-B01})$$

In this case we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in [\tau \sigma]_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in [\tau \sigma]_V \quad (\text{FB-B3})$$

Since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_1 \downarrow_j v'_1$  therefore  $\exists J_1 < j - f - J < k - f - J$  s.t  $(e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{J_1} v'_1$ . Similarly,  $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$  s.t  $(H'_1, v'_1) \downarrow_{J'_1}^f -$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  in the first conjunct of (FB-B3)  
 $(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in [\tau \sigma]_V$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that  
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  we get  $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 2.18 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V \quad (\text{FB-B4})$$

Now we can apply Theorem 2.22 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_V \quad (\text{FB-B5})$$

Instantiating  $c_1$  with  $J_1$  in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_V$$

From Definition 2.6 we have

$$\begin{aligned} \forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies \\ \exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) / \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Instantiating  $K$  with  $m_{u1} + 1 + J'_1$ ,  $\theta'_e$  with  $W''.\theta_1$ ,  $H_1$  with  $H'_1$  (from FB-B1) and  $J_2$  with  $J'_1$  we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1. (m_{u1} + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) / \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow v'_2$ . Say this reduction happens in  $t$  steps. Therefore  $\exists t_1 < t < k \leq n$  s.t  $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{t_1} v'_2$  and similarly  $\exists t_2 < t - t_1 < k - t_1$  s.t  $(H, v'_2)\gamma \downarrow_2 \downarrow_{t_2}^f (H''_2, v''_2)$

Again since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow_t v'_2$  therefore  $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$  s.t  $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{J_2} v'_2$ . Similarly  $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$  s.t  $(H'_2, v'_2) \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with  $m_{u2} + 1 + J_2 + J'_2$  we get  
 $(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v''_2) \in [\tau \sigma]_V$



Again since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that  
 $\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $m_{u_2} + 1 + J_2 + J'_2$  we get  $(W.\theta_2, m_{u_2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 2.18 we know that

$$(W''.\theta_2, m_{u_2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 2.22 to get

$$(W''.\theta_2, m_{u_2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v_2''\}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c_2 < (m_{u_2} + 1 + J_2 + J'_2).(e_b)\gamma \downarrow_2 \cup \{x \mapsto v_2''\} \downarrow_{c_2} v_{o_2} \implies (W''.\theta_2, m_{u_2} + 1 + J_2 - c_2, v_{o_2}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_V \quad (\text{FB-B8})$$

Instantiating  $c_2$  with  $J_2$  in (FB-B8) we get

$$(W''.\theta_2, m_{u_2} + 1 + J'_2, v_{o_2}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_V$$

From Definition 2.6 we have

$$\begin{aligned} \forall K \leq (m_{u_2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3.(K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o_2}) \downarrow_{J_3}^f (H_2'', v_2') \wedge J_3 < K \implies \\ \exists \theta'_2 \sqsupseteq \theta'_e.(K - J_3, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v_2') \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) / \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \end{aligned}$$

Instantiating  $K$  with  $m_{u_2} + 1 + J'_2$ ,  $\theta'_e$  with  $W''.\theta_2$ ,  $H_2$  with  $H_2'$  (from FB-B1) and  $J_3$  with  $J'_2$ , we get

$$\begin{aligned} \exists \theta'_2 \sqsupseteq W''.\theta_2.(m_{u_2} + 1, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, m_{u_2} + 1, v_2') \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) / \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose  $W'$  as  $W_n$  where  $W_n$  is defined as follows:

$$\begin{aligned} W_n.\theta_1 &= \theta'_1 \quad (\text{From (FB-B6)}) \\ W_n.\theta_2 &= \theta'_2 \quad (\text{From (FB-B9)}) \\ W_n.\hat{\beta} &= W''.\hat{\beta} \quad (\text{From (FB-B1)}) \end{aligned}$$

It suffices to prove

- $(k - j, H_1'', H_2'') \triangleright W_n$ :

From Definition 2.9 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'') \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$$

From (FB-B6) we know that  $(m_{u_1} + 1, H_1'') \triangleright \theta'_1$  therefore from Definition 2.8 we know that  $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'')$

Similarly from (FB-B9) we know that  $(m_{u_2} + 1, H_2'') \triangleright \theta'_2$  therefore from Definition 2.8 we know that  $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that  $(k - f - J, H_1', H_2') \triangleright W''$  therefore from Definition 2.9 we know that  $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq W_n.\theta_1$  and  $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

–  $\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k-j-1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A)$ :

4 cases arise for each  $(a_1, a_2) \in W_n.\hat{\beta}$

A.  $H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) = H_2''(a_2)$ :

To prove:

$\overline{W_n.\theta_1(a_1)} = W_n.\theta_2(a_2)$ :

We know from that  $(k - f - J, H_1', H_2') \triangleright W''$

Therefore from Definition 2.9 we have

$\forall (a_1', a_2') \in (W''.\hat{\beta}). W''.\theta_1(a_1') = W''.\theta_2(a_2')$

Since  $W_n.\hat{\beta} = W''.\hat{\beta}$  by construction therefore

$\forall (a_1', a_2') \in (W_n.\hat{\beta}). W''.\theta_1(a_1') = W''.\theta_2(a_2')$

From (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq \theta_1'$  and  $W''.\theta_2 \sqsubseteq \theta_2'$  respectively.

Therefore from Definition 2.1

$\forall (a_1', a_2') \in (W_n.\hat{\beta}). \theta_1'(a_1) = \theta_2'(a_2)$

To prove:

$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$ :

From (FB-B1) we know that  $(k - f - J, H_1', H_2') \triangleright^A W''$

This means from Definition 2.9 we know that

$\forall (a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) = W''.\theta_2(a_{i2}) \wedge$

$(W'', k - f - J - 1, H_1'(a_{i1}), H_2'(a_{i2})) \in \lceil W''.\theta_1(a_{i1}) \rceil_V^A$

Instantiating with  $a_1$  and  $a_2$  and since  $W'' \sqsubseteq W_n$  and  $k - j - 1 < k - f - J - 1$  (since  $j = f + J + J_1 + 1$  therefore from Lemma 2.17 we get

$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$

B.  $H_1'(a_1) \neq H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$ :

To prove:

$\overline{W_n.\theta_1(a_1)} = W_n.\theta_2(a_2)$

Same reasoning as in the previous case

To prove:

$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$

From (FB-B6) and (FB-B9) we know that

$(\forall a. H_1'(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$

$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$

This means we have

$\exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$  and

$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell' \not\sqsubseteq \mathcal{A}$ .

Also from (FB-B6) and (FB-B9),  $(m_{u1}+1, H_1'') \triangleright \theta_1'$  and  $(m_{u2}+1, H_2'') \triangleright \theta_2'$ .

Therefore from Definition 2.8 we have

$$\begin{aligned}(\theta'_1, m_{u1}, H''_1(a_1)) &\in \lfloor \theta'_1(a_1) \rfloor_V \text{ and} \\(\theta'_2, m_{u2}, H''_2(a_1)) &\in \lfloor \theta'_2(a_2) \rfloor_V\end{aligned}$$

Since  $m_{u1}$  and  $m_{u2}$  are arbitrary indices therefore from Definition 2.4 we get

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$C. H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) \neq H''_2(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

From (FB-B9) we know that

$$(\forall a. H'_2(a) \neq H''_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$$

Since  $\ell \sigma \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell' \not\sqsubseteq \mathcal{A}$ .

Since from (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \triangleright^A W''$  that means from Definition 2.9 that  $(W'', k - f - J - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W''.\theta_1(a_1) \rceil_V^A$ . Since  $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$  and since  $\ell' \not\sqsubseteq \mathcal{A}$  therefore from Definition 2.4 and Definition 2.3 we know that

Therefore

$$\forall m. (W''.\theta_1, m, H'_1(a_1)) \in W''.\theta_1(a_1) \quad (F)$$

Instantiating the (F) with  $m_{u1}$  and using Lemma 2.16 we get

$$(\theta'_1, m_{u1}, H'_1(a_1)) \in \theta'_1(a_1)$$

Since from (FB-B9) we know that  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$  therefore from Definition 2.8 we know that  $(\theta'_2, m_{u2}, H''_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', k - j - 1, H''_1(a_1), H''_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$D. H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

Symmetric reasoning as in the previous case

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V:$$

Case  $i = 1$

Given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$$

This further means that given some  $a_1 \in \text{dom}(W_n.\theta_1)$  we need to show

$$(W_n.\theta_1, m, H''_1(a_1)) \in \lfloor W_n.\theta_1(a_1) \rfloor_V$$

Since  $W_n.\theta_1 = \theta'_1$ , it suffices to prove

$$(\theta'_1, m, H''_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$

Like before we apply Theorem 2.22 on  $e_b \gamma \downarrow_1 \cup \{x \mapsto v''_1\}$  but this time at  $m + 1 + J_1 + J'_1$  to get

$$\begin{aligned} & \exists \theta'_1 \sqsupseteq W''.\theta_1.(m+1, H_1'') \triangleright \theta'_1 \wedge (\theta'_1, m_{u1}+1, v'_1) \in [\tau' \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) / \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Since we have  $(m+1, H_1'') \triangleright \theta'_1$  therefore from Definition 2.8 we get the desired.

Case  $i = 2$

Similar reasoning as in the  $i = 1$  case

- $(W'.\theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V$ :

We get this from (FB-B6), (FB-B9) and Lemma 2.16 we get the desired

19. SLIO\*-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new}(e') : \text{SLIO } \ell \ell (\text{ref } \ell' \tau)}$$

To prove:  $(W, n, \text{new}(e')(\gamma \downarrow_1), \text{new}(e')(\gamma \downarrow_2)) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} & \forall i < n. \text{new}(e')(\gamma \downarrow_1 \downarrow_i v_{f1}) \wedge \text{new}(e')(\gamma \downarrow_2 \downarrow v'_{f1}) \implies \\ & (W, n-i, v_{f1}, v'_{f1}) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{new}(e')(\gamma \downarrow_1 \downarrow_i v_{f1}) \wedge \text{new}(e')(\gamma \downarrow_2 \downarrow v'_{f1})$

From SLIO\*-Sem-val we know that  $v_{f1} = \text{new}(e')(\gamma \downarrow_1)$ ,  $v_{f2} = \text{new}(e')(\gamma \downarrow_2)$  and  $i = 0$

We are required to prove

$$(W, n, \text{new}(e')(\gamma \downarrow_1), \text{new}(e')(\gamma \downarrow_2)) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_V^A$$

Let  $v_1 = \text{new}(e')(\gamma \downarrow_1)$  and  $v_2 = \text{new}(e')(\gamma \downarrow_2)$

From Definition 2.4 we are required to prove

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma)) \wedge \\ & \forall \ell \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, v_\ell) \Downarrow_j^f (H', v'_\ell) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, v'_\ell) \in [(\text{ref } \ell' \tau)]_V \sigma \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)) \end{aligned}$$

This means we need to prove the following:

- (a)  $\forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2.$   
 $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies$   
 $\exists W' \sqsupseteq W_e.(k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma):$

This means we are given some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also we are given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e.(k-j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k-j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma) \quad (\text{FB-R0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ (W_e, k-f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e' \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k-f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_n$  where

$$W_n.\theta_1 = W_e.\theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n.\theta_2 = W_e.\theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n.\hat{\beta} = W_e.\hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

i.  $(k-j, H'_1, H'_2) \triangleright W_n$ :

From Definition 2.9 it suffices to prove:

$$\begin{aligned} \text{dom}(W_n.\theta_1) &\subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge \\ (W_n.\hat{\beta}) &\subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W_n.\hat{\beta}). &(W_n.\theta_1(a_1) = W_n.\theta_2(a_2)) \wedge \\ (W_n, (k-j) - 1, &H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau \sigma]_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in &\text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a_i)) \in [W_n.\theta_i(a_i)]_V \end{aligned}$$

This means we need to prove

- $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$ :

We know that  $\text{dom}(W_n.\theta_1) = \text{dom}(W_e.\theta_1) \cup \{a_1\}$  and  $\text{dom}(W_n.\theta_2) = \text{dom}(W_e.\theta_2) \cup \{a_2\}$

Also  $\text{dom}(H'_1) = \text{dom}(H_1) \cup \{a_1\}$  and  $\text{dom}(H'_2) = \text{dom}(H_2) \cup \{a_2\}$

Therefore from  $(k, H_1, H_2) \triangleright W_e$  and from construction of  $W_n$  we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2)) \wedge (W_n, k-j-1, H'_1(a'_1), H'_2(a'_2)) \in [\text{Labeled } \ell' \tau \sigma]_V^A$ :

$\forall (a'_1, a'_2) \in (W_n.\hat{\beta})$ .

A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

Since from (FB-R1) we know that  $(W_e, k-f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$

And since from SLIO\*-Sem-ref we know that  $H'_1(a_1) = v_{h1}$ ,  $H'_2(a_2) = v'_{h1}$

and  $j = f + 1$  therefore from Lemma 2.17 we get

$$(W_n, k-j-1, H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau \sigma]_V^A$$

B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise

C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise

D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.9

•  $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in [W_n.\theta_i(a'_i)]_V$ :

When  $i = 1$

Given some  $m$

$\forall a'_1 \in \text{dom}(W_n.\theta_1)$ .

– when  $a'_1 = a_1$ :

From construction

$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$

And from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$

Therefore from Lemma 2.15 get the desired

– Otherwise:

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.9

When  $i = 2$

Similar reasoning as with  $i = 1$

ii.  $\text{ValEq}(\mathcal{A}, W_n, k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma)$ :

From SLIO\*-Sem-ref we know that  $v'_1 = a_1$  and  $v'_2 = a_2$

2 cases arise:

A.  $\ell \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$(W_n, k - j, a_1, a_2) \in (\text{ref } \ell' \tau) \sigma$

From Definition 2.4 it suffices to prove

$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$

This holds from construction of  $W_n$

B.  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$\forall m. (W_n.\theta_1, m, a_1) \in (\text{ref } \ell' \tau) \sigma$  and  $(W_n.\theta_2, m, a_2) \in (\text{ref } \ell' \tau) \sigma$

From Definition 2.6 this means for any given  $m$  we need to prove that

$W_n.\theta_1(a_1) \in (\text{Labeled } \ell' \tau) \sigma$  and  $W_n.\theta_2(a_2) \in (\text{Labeled } \ell' \tau) \sigma$

This holds from construction of  $W_n$

(b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)$ :

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$

$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that  
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.22 to get

$$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_E$$

This means from Definition 2.7 we get

$$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$$

This further means that given some  $c < k$  s.t  $\text{ref } (e')\gamma \downarrow_1 \downarrow_c v$ . From SLIO\*-Sem-val we know that  $c = 0$  and  $v = \text{ref } (e')\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\text{SLIO } \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$$

From Definition 2.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau) \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

## 20. SLIO\*-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)}$$

To prove:  $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall i < n. !e' \gamma \downarrow_1 \downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $!e' \gamma \downarrow_1 \downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \downarrow v'_{f1}$

From SLIO\*-Sem-val we know that  $v_{f1} = !e' \gamma \downarrow_1$ ,  $v_{f2} = !e' \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, !e' \gamma \downarrow_1, !e' \gamma \downarrow_2) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A$$

Let  $v_1 = !e' \gamma \downarrow_1$  and  $v_2 = !e' \gamma \downarrow_2$

From Definition 2.4 it suffices to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)) \wedge \end{aligned}$$

$$\begin{aligned}
& \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma) \right)
\end{aligned}$$

This means we need to prove:

$$\begin{aligned}
\text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\
& (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma):
\end{aligned}$$

This means we are given is some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\begin{aligned}
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma) \\
& \text{(FB-D0)}
\end{aligned}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in \llbracket \text{ref } \ell \tau \rrbracket_E^A$$

This means from Definition 2.5 we need to prove:

$$\begin{aligned}
& \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\
& (W_e, k - f, v_{h1}, v'_{h1}) \in \llbracket \text{ref } \ell \tau \rrbracket_V^A
\end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in \llbracket \text{ref } \ell \tau \rrbracket_V^A \quad \text{(FB-D1)}$$

In order to prove (FB-D0) we choose  $W'$  as  $W_e$ . Also from SLIO\*-Sem-deref we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . Also we know that  $v_{h1} = a_1$  and  $v'_{h1} = a_2$ .

- $(k - j, H_1, H_2) \triangleright W_e$ :  
Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 2.21 we get  $(k - j, H_1, H_2) \triangleright W_e$
- $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)$ :  
From SLIO\*-Sem-ref we know that  $v'_1 = H_1(a_1)$  and  $v'_2 = H_2(a_2)$

2 cases arise:

–  $\ell' \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$$(W_e, k - j, v'_1, v'_2) \in \llbracket \text{Labeled } \ell \tau \rrbracket \sigma$$

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in \llbracket \text{ref } \ell \tau \rrbracket_V^A$

Therefore from Definition 2.4 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that  $(W_e, k, H_1(a_1), H_2(a_2)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V^A$

From Lemma 2.17 we get  $(W_e, k - j, H_1(a_1), H_2(a_2)) \in \llbracket (\text{Labeled } \ell \tau) \rrbracket_V^A$



–  $\ell' \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 it suffices to prove that

$\forall m. (W_e.\theta_1, m, H_1(a_1)) \in (\text{Labeled } \ell \tau) \sigma$  and  $(W_e.\theta_2, m, H_2(a_2)) \in (\text{Labeled } \ell \tau) \sigma$   
(FB-B2)

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau \sigma]_V^A$

Therefore from Definition 2.4 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that  $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau \sigma]_V^A$

Finally from Lemma 2.15 we get (FB-B2)

- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 2.24 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.22 to get

$(W.\theta_1, k, (!e'\gamma \downarrow_1) \in [(\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_E$

This means from Definition 2.7 we get

$\forall c < k. !e'\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$

Instantiating  $c$  with 0 and from SLIO\*-Sem-val we know that  $v = !e'\gamma \downarrow_1$

And we have  $(W.\theta_1, k, !e'\gamma \downarrow_1) \in [(\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$

From Definition 2.6 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$   
 $\exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma)$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

21. SLIO\*-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_l := e_r : \text{SLIO } \ell \ell \text{ unit}}$$

To prove:  $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in [\text{SLIO } \ell \ell \text{ unit } \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall i < n. (e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\text{SLIO } \ell \ell \text{ unit } \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $(e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1}$

From SLIO\*-Sem-val we know that  $v_{f1} = (e_l := e_r) \gamma \downarrow_1$ ,  $v_{f2} = (e_l := e_r) \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in [\text{SLIO } \ell \ell \text{ unit } \sigma]_V^A$$

Let  $e_1 = (e_l := e_r) \gamma \downarrow_1$  and  $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 2.4 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{unit}]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}): \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

And finally given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\begin{aligned} & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}) \\ & \text{(FB-A0)} \end{aligned}$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\text{ref } \ell' \tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} &\implies \\ (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_V^A & \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \Downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \Downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-A1})$$

IH2:

$$(W_e, k - f, e_r (\gamma \Downarrow_1), e_r (\gamma \Downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall s < k - f. e' \gamma \Downarrow_1 \Downarrow_s v_{h2} \wedge e' \gamma \Downarrow_2 \Downarrow v'_{h2} &\implies \\ (W_e, k - f - s, v_{h2}, v'_{h2}) &\in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A \end{aligned}$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f < k - f$  s.t  $e_r \gamma \Downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \Downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-A2})$$

In order to prove (FB-A0) we choose  $W'$  as  $W_e$ . Also from SLIO\*-Sem-assign we know that  $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$  and  $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$ , and  $j = f + s + 1$

We need to prove the following:

- i.  $(k - j, H'_1, H'_2) \triangleright W_e$ :

Say  $v_{h1} = a_1$  and  $v'_{h1} = a_2$

From Definition 2.9 it suffices to prove:

$$\begin{aligned} \text{dom}(W_e.\theta_1) &\subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge \\ (W_e.\hat{\beta}) &\subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W_e.\hat{\beta}). &(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) \wedge \\ (W_e, (k - j) - 1, &H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A) \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in &\text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A \end{aligned}$$

This means we need to prove

- $\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2))$ :

Since  $\text{dom}(H_1) = \text{dom}(H'_1)$  and  $\text{dom}(H_2) = \text{dom}(H'_2)$ , and also we know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore we obtain the desired directly from Definition 2.9

- $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A)$ :

$\forall (a'_1, a'_2) \in (W_e.\hat{\beta})$ .

- A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :

From (FB-A1) and from Definition 2.4 we get

$$(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau \sigma))_{\mathcal{V}}^A$$

Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A$

And since from SLIO\*-Sem-assign we know that  $H'_1(a_1) = v_{h2}$ ,  $H'_2(a_2) = v'_{h2}$  and  $j = f + s + 1$  therefore from Lemma 2.17 we get

$$(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau \sigma]_{\mathcal{V}}^A$$

- B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise

- C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise

D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.9

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a'_i)) \in \llbracket W_e.\theta_i(a'_i) \rrbracket_V$ :

When  $i = 1$

Given some  $m$

$\forall a'_1 \in \text{dom}(W_e.\theta_1)$ .

– when  $a'_1 = a_1$ :

From (FB-A1) and from Definition 2.4 we get

$$(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

Since from (FB-A2) we know that  $(W_e, k-f-s, v_{h2}, v'_{h2}) \in \llbracket \text{Labeled } \ell' \tau \sigma \rrbracket_V^A$

Therefore from Lemma 2.15 get the desired

– Otherwise:

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 2.9

When  $i = 2$

Similar reasoning as with  $i = 1$

- ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell, (), (), \text{unit})$ :

Holds directly from Definition 2.3 and Definition 2.4

- (b)  $\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \text{unit} \rrbracket_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell \sigma)$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket (\text{unit}) \sigma \rrbracket_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sigma \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell \sigma)$$

Since  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V^A$  therefore from Lemma 2.24 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in \llbracket \Gamma \rrbracket_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$

Now we can apply Theorem 2.22 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in \llbracket (\text{SLIO } \ell \ell (\text{unit})) \sigma \rrbracket_E$$

This means from Definition 2.7 we get

$$\forall c < k. (e_l := e_r)\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in \llbracket (\text{SLIO } \ell \ell (\text{unit})) \sigma \rrbracket_V$$

Instantiating  $c$  with 0 and from SLIO\*-Sem-val we know that  $v = (e_l := e_r)\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in \llbracket (\text{SLIO } \ell \ell (\text{unit})) \sigma \rrbracket_V$

From Definition 2.6 we have

$$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$$

$$\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V \wedge$$

$$(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma)$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

□

**Lemma 2.26** (SLIO\*: Equivalence of values).  $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies$$

$$\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$$

*Proof.* Given that  $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$ . From Definition 2.3 two cases arise

1.  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

2 cases arise

(a)  $\ell' \sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 2.17 we know that  $(W', j, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

And thus from Definition 2.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 2.15 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.16 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Given is  $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.16 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

**Lemma 2.27** (SLIO\*: Subtyping binary). *The following holds:*

$$\forall \Sigma, \Psi, \sigma, \tau, \tau'.$$

$$1. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\mathcal{A}} \subseteq [(\tau' \sigma)]_V^{\mathcal{A}}$$

$$2. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\mathcal{A}} \subseteq [(\tau' \sigma)]_E^{\mathcal{A}}$$

*Proof.* Proof of statement (1)

Proof by induction on the  $\tau <: \tau'$

1. SLIO\*sub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^A] \subseteq [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^A]$

IH1:  $[\lceil \tau'_1 \sigma \rceil_V^A] \subseteq [\lceil \tau_1 \sigma \rceil_V^A]$  (Statement 1)

$[\lceil \tau_2 \sigma \rceil_E^A] \subseteq [\lceil \tau'_2 \sigma \rceil_E^A]$  (Sub-A0 From Statement 2)

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^A]. (W, n, \lambda x.e_1, \lambda x.e_2) \in [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^A]$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^A]$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^A]$

From Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A &\implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V &\implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2 \sigma]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V &\implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 2.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A &\implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2 \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V &\implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V &\implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E) \end{aligned}$$

This means need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A) :$$

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (85)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 85 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using (Sub-A0) we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

$$(b) \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E) :$$

Given:  $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$

Since we are given  $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$  and since  $\tau'_1 \sigma <: \tau_1 \sigma$  therefore from Lemma 2.23 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (86)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_l, k, v'_1$  and  $v'_2$  we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E) \quad (87)$$

Therefore from Equation 86 and 87 we get  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E$

Since  $\tau_2 \sigma <: \tau'_2 \sigma$  therefore from Lemma 2.23 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$$

$$(c) \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E):$$

Similar reasoning as in the previous case

## 2. SLIO\*sub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $\lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A \subseteq \lceil ((\tau'_1 \times \tau'_2) \sigma) \rceil_V^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau'_1 \sigma) \rceil_V^A$  (Statement (1))

IH2:  $\lceil (\tau_2 \sigma) \rceil_V^A \subseteq \lceil (\tau'_2 \sigma) \rceil_V^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$$(W, n, v_1, v'_1) \in \lceil \tau_1 \sigma \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau_2 \sigma \rceil_V^A \quad (88)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in \lceil \tau'_1 \sigma \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau'_2 \sigma \rceil_V^A$$

Since from Equation 88 we know that  $(W, n, v_1, v'_1) \in \lceil \tau_1 \sigma \rceil_V^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in \lceil \tau'_1 \sigma \rceil_V^A$

Similarly since  $(W, n, v_2, v'_2) \in \lceil \tau_2 \sigma \rceil_V^A$  from Equation 88 therefore from IH2 we have  $(W, n, v_2, v'_2) \in \lceil \tau'_2 \sigma \rceil_V^A$

## 3. SLIO\*sub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A \subseteq \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau'_1 \sigma) \rceil_V^A$  (Statement (1))

IH2:  $\lceil (\tau_2 \sigma) \rceil_V^A \subseteq \lceil (\tau'_2 \sigma) \rceil_V^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A. (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_V^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_V^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s1} = \text{inl } v_{i2}$ :

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1 \sigma \rceil_V^A \quad (89)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_V^A$$

From Equation 89 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_V^A$$

(b)  $v_s = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2 \sigma \rceil_V^A \quad (90)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_V^A$$

From Equation 90 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_V^A$$

#### 4. SLIO\*sub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $\lceil ((\forall \alpha. \tau_1) \sigma) \rceil_V^A \subseteq \lceil ((\forall \alpha. \tau_2) \sigma) \rceil_V^A$

$\forall \sigma. \lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$  (Sub-F2, From Statement (2))

It suffices to prove:  $\forall (W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. \tau_1) \sigma) \rceil_V^A.$

$(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. \tau_2) \sigma) \rceil_V^A$

This means that given:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\tau_1)) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in \lceil \tau_1[\ell'/\alpha] \sigma \rceil_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in \lceil \tau_1[\ell'/\alpha] \rceil_E) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in \lceil \tau_1[\ell''/\alpha] \rceil_E) \quad (\text{Sub-F1}) \end{aligned}$$

And it suffices to prove:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. \tau_2) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:



$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge \\ \forall \theta'_i \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_i, k, e_1) \in [\tau_2[\ell''/\alpha]]_E) \wedge \\ \forall \theta'_i \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_i, k, e_2) \in [\tau_2[\ell''/\alpha]]_E) \end{aligned}$$

This means we are required to show:

$$(a) \quad \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A):$$

By instantiating the first conjunct of Sub-F1 with  $W''$ ,  $n''$  and  $\ell''$  we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from Sub-F2 instantiated at  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

$$(b) \quad \forall \theta'_i \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_i, k, e_1) \in [\tau_2[\ell''/\alpha]]_E):$$

By instantiating the second conjunct of Sub-F1 with  $\theta'_i$  and  $\ell''$  we know that the following holds

$$((\theta'_i, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E)$$

Since  $\tau_1 \sigma \cup \{\alpha \mapsto \ell''\} <: \tau_2 \sigma \cup \{\alpha \mapsto \ell''\}$  therefore from Lemma 2.23 we know that

$$((\theta'_i, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E)$$

$$(c) \quad \forall \theta'_i \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_i, k, e_2) \in [\tau_2[\ell''/\alpha]]_E):$$

Similar reasoning as in the previous case

##### 5. SLIO\*sub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2}$$

To prove:  $[\!(c_1 \implies \tau_1) \sigma]_{\mathcal{V}}^A \subseteq [\!(c_2 \implies \tau_2) \sigma]_{\mathcal{V}}^A$

$$[\!(\tau_1 \sigma) \sigma]_E^A \subseteq [\!(\tau_2 \sigma) \sigma]_E^A \text{ (Sub-C0, From Statement (2))}$$

It suffices to prove:  $\forall (W, n, \nu e_1, \nu e_2) \in [\!(c_1 \implies \tau_1) \sigma]_{\mathcal{V}}^A. (W, n, \nu e_1, \nu e_2) \in [\!(c_2 \implies \tau_2) \sigma]_{\mathcal{V}}^A$

This means that given:  $(W, n, \nu e_1, \nu e_2) \in [\!(c_1 \implies \tau_1) \sigma]_{\mathcal{V}}^A$

Therefore from Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma \implies (W', n', e_1, e_2) \in [\tau_1 \sigma]_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_1) \in [\tau_1 \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_2) \in [\tau_1 \sigma]_E \quad \text{(Sub-C1)} \end{aligned}$$

And it suffices to prove:  $(W, n, \nu e_1, \nu e_2) \in [\!(c_2 \implies \tau_2) \sigma]_{\mathcal{V}}^A$

Again from Definition 2.4, it suffices to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A \wedge \\ \forall \theta'_i \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 \implies (\theta'_i, j, e_1) \in [\tau_2 \sigma]_E \wedge \\ \forall \theta'_i \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_i, j, e_2) \in [\tau_2 \sigma]_E \end{aligned}$$

This means that we are required to show the following:

(a)  $\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A$ :

We are given  $W'' \sqsupseteq W, n'' < n$  also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with  $W''$  and  $n''$  we know that the following holds

$$(W'', n'', e_1, e_2) \in [\tau_1 \sigma]_E^A$$

Therefore from (Sub-C0) we get  $(W'', n'', e_1, e_2) \in [\tau_2 \sigma]_E^A$

(b)  $\forall \theta'_i \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \implies (\theta'_i, k, e_1) \in [\tau_2 \sigma]_E$ :

We are given some  $\theta'_i \sqsupseteq W.\theta_1, k$ , also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with  $\theta'_i$  we know that the following holds

$$(\theta'_i, k, e_1) \in [\tau_1 \sigma]_E$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  therefore from Lemma 2.23 we get

$$(\theta'_i, k, e_1) \in [\tau_2 \sigma]_E$$

(c)  $\forall \theta'_i \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_i, j, e_2) \in [\tau_2 \sigma]_E$ :

Similar reasoning as in the previous case

## 6. SLIO\*sub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $[((\text{Labeled } \ell \tau) \sigma)]_{\mathcal{V}}^A \subseteq [((\text{Labeled } \ell' \tau') \sigma)]_{\mathcal{V}}^A$

IH:  $[(\tau \sigma)]_{\mathcal{V}}^A \subseteq [(\tau' \sigma)]_{\mathcal{V}}^A$

It suffices to prove:  $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_{\mathcal{V}}^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_{\mathcal{V}}^A$

This means we are given  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_{\mathcal{V}}^A$

From Definition 2.4 it means we have  $\text{ValEq}(\mathcal{A}, W, \ell \sigma, n, v_1, v_2, \tau \sigma)$  (Sub-L0)

and it suffices to prove  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_{\mathcal{V}}^A$

Again from Definition 2.4 it means we need to prove that

$$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$$

Since we have (Sub-L0) and  $\ell \sigma \sqsubseteq \ell' \sigma$  therefore from Lemma 2.26 we have

$$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau \sigma)$$

2 cases arise:

(a)  $\ell' \sigma \sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 we know that  $(W, n, v_1, v_2) \in [\tau \sigma]_{\mathcal{V}}^A$

From IH we also know that  $(W, n, v_1, v_2) \in [\tau' \sigma]_{\mathcal{V}}^A$

And from Definition 2.4 we get  $\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

(b)  $\ell' \sigma \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 2.3 we know that  $\forall j. (W.\theta_1, j, v_1) \in [\tau \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau \sigma]_V$

Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.23 we get  $(W.\theta_1, j, v_1) \in [\tau' \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau' \sigma]_V$

And from Definition 2.4 we get  $ValEq(\mathcal{A}, W, \ell' \sigma, n, Lb(v_1), Lb_\ell(v_2), \tau' \sigma)$

7. SLIO\*sub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \text{SLIO } \ell_i \ell_o \tau <: \text{SLIO } \ell'_i \ell'_o \tau'}$$

To prove:  $\lceil ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rceil_V^A \subseteq \lceil ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rceil_V^A$

IH:  $\lceil (\tau \sigma) \rceil_V^A \subseteq \lceil (\tau' \sigma) \rceil_V^A$

It suffices to prove:  $\forall (W, n, e_1, e_2) \in \lceil ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rceil_V^A. (W, n, e_1, e_2) \in \lceil ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rceil_V^A$

This means we are given  $(W, n, e_1, e_2) \in \lceil ((\text{SLIO } \ell_i \ell_o \tau) \sigma) \rceil_V^A$

From Definition 2.4 it means we have

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \quad (\text{Sub-CG0}) \end{aligned}$$

And we need to prove

$$(W, n, e_1, e_2) \in \lceil ((\text{SLIO } \ell'_i \ell'_o \tau') \sigma) \rceil_V^A$$

Again from Definition 2.4 it means we need to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) \right) \end{aligned}$$

It means we need to prove:

$$\begin{aligned} & \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau' \sigma): \end{aligned}$$

This means we are given  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2, v'_1, v'_2, j < k$  s.t  
 $(k, H_1, H_2) \triangleright W_e$ ,  $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2)$

And we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma)$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1)  $\text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma)$

Therefore from Lemma 2.26 we get  $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau \sigma)$

- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma):$

Case  $l = 1$

Here we are given  $k, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

- i.  $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau' \sigma]_V:$

Instantiating the second conjunct of (Sub-CG0) with the given  $k, \theta_e, H, j$  to get  
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau \sigma]_V$

Since  $\tau \sigma <: \tau' \sigma$  therefore from Lemma 2.23 we get  $(\theta', k - j, v'_1) \in [\tau' \sigma]_V$

- ii.  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell')$ :

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell')$$

Since  $\ell'_i \sigma \sqsubseteq \ell_i \sigma$  therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell')$$

- iii.  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma):$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$$

Since  $\ell'_i \sigma \sqsubseteq \ell_i \sigma$  therefore we also get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$$

Case  $l = 2$

Symmetric reasoning as in the previous  $l = 1$  case

## 8. SLIO\*sub-base:

Trivial

Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [(\tau \sigma)]_E^A. (W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$$

This means given  $(W, n, e_1, e_2) \in [(\tau \sigma)]_E^A$

From Definition 2.5 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau \sigma]_{\mathcal{V}}^A \quad (\text{Sub-E0})$$

And it suffices to prove  $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

Again from Definition 2.5 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_{\mathcal{V}}^A$$

This means that given  $i < n$  s.t.  $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$  we need to prove  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_{\mathcal{V}}^A$

Instantiating (Sub-E0) with the given  $i$  we get  $(W, n - i, v_1, v_2) \in [\tau \sigma]_{\mathcal{V}}^A$

From Statement (1) we get  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_{\mathcal{V}}^A$  □

**Theorem 2.28** (SLIO\*: NI). *Say*  $\text{bool} = (\text{unit} + \text{unit})$

$\forall v_1, v_2, e, \tau, n.$

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \text{SLIO } \perp \perp \text{ bool} \wedge$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v_1') \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v_2') \implies v_1' = v_2'$$

*Proof.* Given some

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \text{SLIO } \perp \perp \text{ bool} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v_1') \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v_2')$

And we need to prove

$$v_1' = v_2'$$

From Theorem 2.25 we know that

$$\forall n. (\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \text{ bool}]_E^\perp$$

Similarly from Theorem 2.25 and Definition 2.14 we also get

$$\forall n. (\emptyset, n, e[v_1/x], e[v_2/x]) \in [\text{SLIO } \perp \perp \text{ bool}]_E^\perp$$

From Definition 2.5 we get

$$\forall n. \forall i < n. e[v_1/x] \Downarrow_i v_{11} \wedge e[v_2/x] \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\text{SLIO } \perp \perp \text{ bool}]_{\mathcal{V}}^\perp$$

Instantiating it with  $n' + 1$  and then with 0, from CG-val we have  $v_{11} = e[v_1/x]$  and  $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n' + 1, e[v_1/x], e[v_2/x]) \in [\text{SLIO } \perp \perp \text{ bool}]_{\mathcal{V}}^\perp$$

From Definition 2.6 we have

$$\left( \forall k \leq (n' + 1), W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right.$$

$$\forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow_j^f (H_2', v_2'') \wedge j < k \implies \exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \Big) \wedge$$

$$\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right.$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\mathbf{b}]_{\mathcal{V}} \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge$$

$$\left( \forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp \right)$$

Instantiating the first conjunct with  $n' + 1, \emptyset, \emptyset, \emptyset$ .

Since we know that

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v'_2)$$

Therefore we instantiate  $v''_1$  with  $v'_1$ ,  $v''_2$  with  $v'_2$ ,  $j$  with  $n'$  to get  
 $\exists W' \sqsupseteq \emptyset.(n - n', H'_1, H'_2) \triangleright W' \wedge ValEq(\perp, W', k - j, \perp, v'_1, v'_2, \text{bool})$

From Definition 2.3 and Definition 2.6 we get  $v'_1 = v'_2$

□

### 3 Translations between FG and SLIO\*

#### 3.1 Translation from SLIO\* to FG

##### 3.1.1 Type directed translation from SLIO\* to FG

SLIO\* types are translated into FG types by the following definition of  $\llbracket \cdot \rrbracket$

$$\begin{array}{ll}
\llbracket \mathbf{b} \rrbracket = \mathbf{b}^\perp & \llbracket \text{ref } \ell \ \tau \rrbracket = (\text{ref } (\llbracket \tau \rrbracket) + \text{unit})^\ell^\perp \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \text{SLIO } \ell_i \ \ell_o \ \tau \rrbracket = (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket) + \text{unit})^{\ell_o}^\perp \\
\llbracket \tau_1 \times \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp & \llbracket c \Rightarrow \tau \rrbracket = (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \\
\llbracket \tau_1 + \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp & \llbracket \forall \alpha. \tau \rrbracket = (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \\
\llbracket \text{Labeled } \ell \ \tau \rrbracket = (\llbracket \tau \rrbracket + \text{unit})^\ell &
\end{array}$$

The translation judgment for expressions is of the form  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$ . Its rules are shown below.

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{var} \\
\frac{\Sigma; \Psi; \Gamma, x : \tau \vdash e : \tau' \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \lambda x. e : \tau \rightarrow \tau' \rightsquigarrow \lambda x. e_F} \text{lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau \rightarrow \tau' \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash e_1 \ e_2 : \tau' \rightsquigarrow e_{F1} \ e_{F2}} \text{app} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2 \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{F1}, e_{F2})} \text{prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e) : \tau_1 \rightsquigarrow \text{fst}(e_F)} \text{fst} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{snd}(e) : \tau_1 \rightsquigarrow \text{snd}(e_F)} \text{snd} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inl}(e_F)} \text{inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inr}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inr}(e_F)} \text{inr} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 + \tau_2 \rightsquigarrow e_F \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{case}(e_F, x.e_{F1}, y.e_{F2})} \text{case} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \ \tau) \rightsquigarrow \text{inl}(e_F)} \text{label} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \text{SLIO } \ell_i \ (\ell_i \sqcup \ell) \ \tau \rightsquigarrow \lambda_. e_F} \text{unlabel} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{SLIO } \ell_i \ \ell_o \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \text{SLIO } \ell_i \ \ell_i \ (\text{Labeled } \ell_o \ \tau) \rightsquigarrow \lambda_. \text{inl}(e_F \ ())} \text{toLabeled} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \text{SLIO } \ell_i \ \ell_i \ \tau \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ret}
\end{array}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{SLIO } \ell_i \ell \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \text{SLIO } \ell \ell_o \tau' \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \text{SLIO } \ell_i \ell_o \tau' \rightsquigarrow \lambda\_.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{bind} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda\_.\text{inl}(\text{new } (e_F))} \text{ref} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash !e : \text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda\_.\text{inl}(e_F)} \text{deref} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \text{SLIO } \ell \ell \text{ unit} \rightsquigarrow \lambda\_.\text{inl}(e_{F1} := e_{F2})} \text{assign} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F} \text{sub} \\
\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \Lambda e : \forall \alpha. \tau \rightsquigarrow \Lambda e_F} \text{FI} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \rightsquigarrow e_F \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha] \rightsquigarrow e_F[]} \text{FE} \\
\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \nu e : c \Rightarrow \tau \rightsquigarrow \nu e_F} \text{CI} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau \rightsquigarrow e_F \bullet} \text{CE}
\end{array}$$

### 3.1.2 Type preservation for SLIO\* to FG translation

**Assumption 3.1.**  $\forall e, \tau, \Sigma, \Psi, \Gamma, \ell_i, \ell_o.$

$$\Sigma; \Psi; \Gamma \vdash e : \text{SLIO } \ell_i \ell_o \tau \implies \ell_i \sqsubseteq \ell_o$$

**Theorem 3.2** (SLIO\*  $\rightsquigarrow$  FG: Type preservation).  $\forall \Sigma, \Psi, \Gamma, e_C, \tau.$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \text{ is a valid typing derivation in SLIO}^* \implies$$

$$\exists e_F.$$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$$

$$\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau] \text{ is a valid typing derivation in FG}$$

*Proof.* Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\perp}} \text{FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\ell}} \text{FG-sub}$$

2. unlabel:

P1:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i \sqcup \ell \quad \overline{\Sigma; \Psi \vdash ([\tau] + \text{unit}) <: ([\tau] + \text{unit})} \text{Lemma 1.1}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell} <: ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{FGsub-label}$$



Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\top} e_F : ([\tau] + \mathbf{unit})^\ell} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top \quad P1}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_F : ([\tau] + \mathbf{unit})^{\ell_i \sqcup \ell}} \text{FG-sub}}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\top} \lambda_{-}.e_F : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^{\ell_i \sqcup \ell})^\perp} \text{FG-lam}}$$

3. toLabeled:

P2:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\top} e_F : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^{\ell_o})^\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_F : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^{\ell_o})^\perp} \text{FG-sub}}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} () : \mathbf{unit}} \text{P2} \quad \Sigma; \Psi \vdash \ell_i \sqcup \perp \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash ([\tau] + \mathbf{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_F() : ([\tau] + \mathbf{unit})^{\ell_o}} \text{FG-app}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} \mathbf{inl}(e_F()) : (([\tau] + \mathbf{unit})^{\ell_o} + \mathbf{unit})^{\ell_i}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\mathbf{inl}(e_F()) : (\mathbf{unit} \xrightarrow{\ell_i} (([\tau] + \mathbf{unit})^{\ell_o} + \mathbf{unit})^{\ell_i})^\perp} \text{FG-lam}}$$

4. ret:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\top} e_F : [\tau]} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_F : [\tau]} \text{FG-sub} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} \mathbf{inl}(e_F) : ([\tau] + \mathbf{unit})^{\ell_i}} \text{FG-sub, FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\mathbf{inl}(e_F) : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^{\ell_i})^\perp}$$

5. bind:

P1.1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\top} e_{F1} : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^\ell)^\perp} \text{IH1, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_{F1} : (\mathbf{unit} \xrightarrow{\ell_i} ([\tau] + \mathbf{unit})^\ell)^\perp} \text{FG-sub}}$$

P1:

$$\frac{P1.1 \quad \frac{}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} () : \mathbf{unit}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi \vdash (\ell_i \sqcup \perp) \sqsubseteq \ell_i \quad \frac{}{\Sigma; \Psi \vdash ([\tau] + \mathbf{unit})^\ell \searrow \perp} \text{FG-app}}{\Sigma; \Psi; [\Gamma], - : \mathbf{unit} \vdash_{\ell_i} e_{F1}() : ([\tau] + \mathbf{unit})^\ell} \text{FG-app}$$

P2.1:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}}{\Sigma; \Psi \vdash \ell \sqsubseteq \top} \text{IH2, Weakening}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}} \text{FG-sub}$$

P2:

$$\frac{\begin{array}{c} P2.1 \quad \frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi \vdash \perp \sqsubseteq \ell_o} \text{FG-var} \\ \Sigma; \Psi \vdash (\ell \sqcup \perp) \sqsubseteq \ell \quad \frac{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app} \end{array}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell_i \sqcup \ell} e_{F2}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app}$$

P3:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_o}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-sub, FG-inr}$$

Main derivation:

$$\frac{\begin{array}{c} P1 \quad P2 \quad P3 \quad \frac{\frac{\Sigma; \Psi; \Gamma \vdash e_2 : \text{SLIO } \ell \ell_o \tau}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_o} \text{Assumption 3.1}}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \ell} \text{FG-case} \end{array}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-lam, weak}$$

$$\frac{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\text{unit} \xrightarrow{\ell_i} ([\tau'] + \text{unit})^{\ell_o})^{\perp}}$$

6. ref:

P1:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : ([\tau] + \text{unit})^{\ell'} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_F : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}}{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_F : ([\tau] + \text{unit})^{\ell'} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow \ell} \text{FG-ref}} \text{FG-ref}$$

$$\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{new } e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp}}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(\text{new } e_F) : ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(\text{new } e_F) : (\text{unit} \xrightarrow{\ell} ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell})^{\perp}} \text{FG-lam}}$$

7. deref:

P2:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell' \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp}} \text{FG-sub}}$$

P1:

$$P2 \quad \frac{\frac{}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell} \text{Lemma 1.1} \quad \frac{}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell \searrow \perp}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} !e_F : ([\tau] + \text{unit})^\ell} \text{FG-deref}$$

Main derivation:

$$P1 \quad \frac{\frac{\Sigma; \Psi \vdash \perp \sqsubseteq \ell' \quad \frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'}} \text{FG-lam}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\ell'} (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'})^\perp}$$

8. assign:

P3:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{IH2, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}$$

P2:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{IH1, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{FG-sub}$$

P1:

$$P2 \quad P3 \quad \frac{\frac{\frac{}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{Given}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp)} \text{FG-assign}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\ell} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^\ell)^\perp} \text{FG-lam}$$

9. sub:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau']} \text{IH} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top \quad \frac{\Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi \vdash [\tau'] <: [\tau]} \text{Lemma 3.3}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{FG-sub}$$

10. FI:

$$\frac{\frac{}{\Sigma, \alpha; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp} \text{FG-FI}$$

11. FE:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (\forall \alpha. (\top, [\tau]))^{\perp}} \text{ IH}}{\text{FV}(\ell) \in \Sigma \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau[\ell/\alpha]] \searrow \perp} \text{ FG-FE}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F [] : [\tau][\ell/\alpha]} \text{ Lemma 3.6}$$

12. CI:

$$\frac{\overline{\Sigma; \Psi, c; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \nu e_F : (c \stackrel{\top}{\Rightarrow} [\tau])^{\perp}} \text{ FG-CI}$$

13. CE:

$$\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (c \stackrel{\top}{\Rightarrow} [\tau])^{\perp}} \text{ IH} \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau] \searrow \perp}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F \bullet : [\tau]} \text{ FG-CE}$$

□

**Lemma 3.3** (SLIO\*  $\rightsquigarrow$  FG: Subtyping). *For any SLIO\* types  $\tau$  and  $\tau'$ ,  $\Sigma$ , and  $\Psi$ , if  $\Sigma; \Psi \vdash \tau <: \tau'$ , then  $\Sigma; \Psi \vdash [\tau] <: [\tau']$ .*

*Proof.* Proof by induction on SLIO\*'s subtyping relation

1. SLIO\*sub-base:

$$\overline{\Sigma; \Psi \vdash [\tau] <: [\tau]} \text{ Lemma 1.1}$$

2. SLIO\*sub-arrow:

$$\frac{\overline{\Sigma; \Psi \vdash [\tau'_1] <: [\tau_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau'_2] <: [\tau_2]} \text{ IH2} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top}{\Sigma; \Psi \vdash ([\tau_1] \stackrel{\top}{\mapsto} [\tau_2])^{\perp} <: ([\tau'_1] \stackrel{\top}{\mapsto} [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \stackrel{\ell_s}{\mapsto} \tau_2)) <: ((\tau'_1 \stackrel{\ell'_s}{\mapsto} \tau'_2))} \text{ Definition of } [\cdot]$$

3. SLIO\*sub-prod:

$$\frac{\overline{\Sigma; \Psi \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\Sigma; \Psi \vdash ([\tau_1] \times [\tau_2])^{\perp} <: ([\tau'_1] \times [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \times \tau_2)) <: ((\tau'_1 \times \tau'_2))} \text{ Definition of } [\cdot]$$

4. SLIO\*sub-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket <: \llbracket \tau'_2 \rrbracket}}{\text{IH2}}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp <: (\llbracket \tau'_1 \rrbracket + \llbracket \tau'_2 \rrbracket)^\perp} \text{FGsub-arrow}}{\Sigma; \Psi \vdash \llbracket (\tau_1 + \tau_2) \rrbracket <: \llbracket (\tau'_1 + \tau'_2) \rrbracket} \text{Definition of } \llbracket \cdot \rrbracket$$

5. SLIO\*sub-labeled:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum}}{\frac{\frac{\overline{\text{Labeled } \ell_1 \tau_1 <: \text{Labeled } \ell'_1 \tau'_1}}{\text{Given}}}{\ell_1 \sqsubseteq \ell'_1} \text{By inversion}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_1}} \text{FGsub-arrow}}{\Sigma; \Psi \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket} \text{Definition of } \llbracket \cdot \rrbracket$$

6. SLIO\*sub-monad:

P3:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum}$$

P2:

$$\frac{P3 \quad \frac{\overline{\Sigma; \Psi \vdash \text{SLIO } \ell_i \ell_o \tau_1 <: \text{SLIO } \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o} \text{By inversion}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o}} \text{FGsub-label}$$

P1:

$$\frac{\overline{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \quad P2 \quad \frac{\overline{\Sigma; \Psi \vdash \text{SLIO } \ell_i \ell_o \tau_1 <: \text{SLIO } \ell'_i \ell'_o \tau'_1}}{\Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i} \text{Given}}{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o}) <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})} \text{FGsub-arrow}$$

Main derivation:

$$\frac{P1 \quad \overline{\Sigma; \Psi \vdash \perp \sqsubseteq \perp}}{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})^\perp} \text{FGsub-label}}{\Sigma; \Psi \vdash \llbracket \text{SLIO } \ell_i \ell_o \tau_1 \rrbracket <: \llbracket \text{SLIO } \ell'_i \ell'_o \tau'_1 \rrbracket} \text{Definition of } \llbracket \cdot \rrbracket$$

7. SLIO\*sub-forall:

P1:

$$\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket <: \llbracket \tau' \rrbracket} \text{IH, Weakening} \quad \overline{\Sigma, \alpha; \Psi \vdash \top \sqsubseteq \top}}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))} \text{FGsub-forall}$$

Main derivation:

$$\frac{P1 \quad \overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))^\perp} \text{FGsub-label} \\ \hline \Sigma; \Psi \vdash \llbracket \forall \alpha. \tau \rrbracket <: \llbracket \forall \alpha. \tau' \rrbracket$$

8. SLIO\*sub-constraint:

P1:

$$\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket <: \llbracket \tau' \rrbracket} \text{IH} \quad \frac{\overline{\Sigma; \Psi \vdash \top \sqsubseteq \top} \quad \frac{\overline{\Sigma; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given}}{\Sigma; \Psi \vdash c' \Longrightarrow c} \text{By inversion}}{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket) <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)} \text{FGsub-constra}}$$

Main derivation:

$$\frac{P1 \quad \overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket)^\perp <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)^\perp} \text{FGsub-label} \\ \hline \Sigma; \Psi \vdash \llbracket c \Rightarrow \tau \rrbracket <: \llbracket c' \Rightarrow \tau' \rrbracket$$

□

**Lemma 3.4** (SLIO\*  $\rightsquigarrow$  FG: Preservation of well-formedness).  $\forall \Sigma, \Psi, \tau.$

$$\Sigma; \Psi \vdash \tau \text{ WF} \Longrightarrow \Sigma; \Psi \vdash \llbracket \tau \rrbracket \text{ WF}$$

*Proof.* Proof by induction on the  $\tau$  WF relation.

1. SLIO\*-wff-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base}}{\Sigma; \Psi \vdash \mathbf{b}^\perp \text{ WF}} \text{FG-wff-label}$$

2. SLIO\*-wff-unit:

$$\overline{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit}$$

3. SLIO\*-wff-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \text{IH2}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket) \text{ WF}} \text{FG-wff-arrow}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \text{ WF}} \text{FG-wff-label}$$

4. SLIO\*-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \text{IH2}}{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket \text{ WF}} \text{FG-wff-prod}}{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket^\perp \text{ WF}} \text{FG-wff-label}$$

5. SLIO\*-wff-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket WF} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket WF} \text{ IH2}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket WF} \text{ FG-wff-prod}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp WF} \text{ FG-wff-label}$$

6. SLIO\*-wff-ref:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi \vdash \text{ref } \ell \tau WF} \text{ Given}}{\text{FV}(\tau) = \emptyset} \text{ By inversion}}{\text{FV}(\llbracket \tau \rrbracket) = \emptyset} \text{ Lemma 3.5}}{\frac{\frac{\overline{\Sigma; \Psi \vdash \text{ref } \ell \tau WF} \text{ Given}}{\text{FV}(\ell) = \emptyset} \text{ By inversion}}{\text{FV}(\text{unit}) = \emptyset}}{\Sigma; \Psi \vdash \text{FV}((\llbracket \tau \rrbracket + \text{unit})^\ell) = \emptyset} \text{ FG-wff-ref}}{\Sigma; \Psi \vdash \text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell WF} \text{ FG-wff-label}}{\Sigma; \Psi \vdash (\text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp WF} \text{ FG-wff-label}$$

7. SLIO\*-wff-forall:

$$\frac{\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket WF} \text{ IH}}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) WF} \text{ FG-wff-forall}}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp WF} \text{ SLIO*-wff-label}$$

8. SLIO\*-wff-constraint:

$$\frac{\frac{\overline{\Sigma; \Psi, c \vdash \llbracket \tau \rrbracket WF} \text{ IH}}{\Sigma; \Psi \vdash (c \xRightarrow{\top} \llbracket \tau \rrbracket) WF} \text{ FG-wff-constraint}}{\Sigma; \Psi \vdash (c \xRightarrow{\top} \llbracket \tau \rrbracket)^\perp WF} \text{ SLIO*-wff-label}$$

9. SLIO\*-wff-labeled:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket WF} \text{ IH} \quad \overline{\Sigma; \Psi \vdash \text{unit } WF} \text{ FG-wff-unit}}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) WF} \text{ FG-wff-sum}}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^\ell WF} \text{ SLIO*-wff-label}$$

10. SLIO\*-wff-monad:

P1:

$$\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket WF} \text{ IH} \quad \overline{\Sigma; \Psi \vdash \text{unit } WF} \text{ FG-wff-unit}}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) WF} \text{ FG-wff-sum}$$

Main derivation:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \mathbf{unit} \text{ } WF} \text{FG-wff-unit} \quad \frac{P1}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell_o} \text{ } WF} \text{FG-wff-label} \\
\hline
\Sigma; \Psi \vdash (\mathbf{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell_o}) \text{ } WF \quad \text{FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\mathbf{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell_o})^\perp \text{ } WF \quad \text{SLIO}^*\text{-wff-label}
\end{array}$$

□

**Lemma 3.5** (SLIO\*  $\rightsquigarrow$  FG: Free variable lemma).  $\forall \tau. FV(\llbracket \tau \rrbracket) \subseteq FV(\tau)$

*Proof.* Proof by induction on the SLIO\* types,  $\tau$

1.  $\tau = \mathbf{b}$ :

$$\begin{aligned}
& FV(\llbracket \mathbf{b} \rrbracket) \\
&= FV(\mathbf{b}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= \emptyset \\
&= FV(\mathbf{b})
\end{aligned}$$

2.  $\tau = \mathbf{unit}$ :

$$\begin{aligned}
& FV(\llbracket \mathbf{b} \rrbracket) \\
&= FV(\mathbf{unit}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= \emptyset \\
&= FV(\mathbf{unit})
\end{aligned}$$

3.  $\tau = \tau_1 \rightarrow \tau_2$ :

$$\begin{aligned}
& FV(\llbracket \tau_1 \rightarrow \tau_2 \rrbracket) \\
&= FV(\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\
&\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\
&= FV(\tau_1 \rightarrow \tau_2)
\end{aligned}$$

4.  $\tau = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& FV(\llbracket \tau_1 \times \tau_2 \rrbracket) \\
&= FV(\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\
&\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\
&= FV(\tau_1 \times \tau_2)
\end{aligned}$$

5.  $\tau = \tau_1 + \tau_2$ :

$$\begin{aligned}
& FV(\llbracket \tau_1 + \tau_2 \rrbracket) \\
&= FV(\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\
&\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\
&= FV(\tau_1 + \tau_2)
\end{aligned}$$

6.  $\tau = \text{ref } \ell_i \tau_i$ :

$$\begin{aligned}
& FV(\llbracket \text{ref } \ell_i \tau_i \rrbracket) \\
&= FV(\text{ref } (\llbracket \tau_i \rrbracket + \mathbf{unit})^{\ell_i})^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\
&= FV(\llbracket \tau_i \rrbracket) \cup FV(\ell_i) \\
&\subseteq FV(\tau_i) \cup FV(\ell_i) \quad \text{IH} \\
&= FV(\text{ref } \ell_i \tau_i)
\end{aligned}$$



7.  $\tau = \forall\alpha.\tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \forall\alpha.\tau_i \rrbracket) \\
&= \text{FV}(\forall\alpha.(\top, \llbracket \tau_i \rrbracket))^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket \tau_i \rrbracket) - \{\alpha\} \\
&\subseteq \text{FV}(\tau_i) - \{\alpha\} && \text{IH} \\
&= \text{FV}(\forall\alpha.\tau_i)
\end{aligned}$$

8.  $\tau = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket c \Rightarrow \tau_i \rrbracket) \\
&= \text{FV}(c \overset{\top}{\Rightarrow} \llbracket \tau_i \rrbracket)^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\llbracket \tau_i \rrbracket) \\
&\subseteq \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\tau_i) && \text{IH} \\
&= \text{FV}(c \Rightarrow \tau_i)
\end{aligned}$$

9.  $\tau = \text{Labeled } \ell_i \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \text{Labeled } \ell_i \tau_i \rrbracket) \\
&= \text{FV}(\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i} && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \\
&\subseteq \text{FV}(\tau_i) \cup \text{FV}(\ell_i) && \text{IH} \\
&= \text{FV}(\text{Labeled } \ell_i \tau_i)
\end{aligned}$$

10.  $\tau = \text{SLIO } \ell_i \ell_o \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \text{SLIO } \ell_i \ell_o \tau_i \rrbracket) \\
&= \text{FV}(\text{unit} \overset{\ell_i}{\rightarrow} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_o})^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) \\
&\subseteq \text{FV}(\tau_i) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) && \text{IH} \\
&= \text{FV}(\text{SLIO } \ell_i \ell_o \tau_i)
\end{aligned}$$

□

**Lemma 3.6** (SLIO\*  $\rightsquigarrow$  FG: Substitution lemma).  $\forall\tau. s.t \vdash \tau \text{ WF}$  the following holds:

$$\llbracket \tau \rrbracket[\ell/\alpha] = \llbracket \tau[\ell/\alpha] \rrbracket$$

*Proof.* Proof by induction on the SLIO\* types,  $\tau$

1.  $\tau = \mathbf{b}$ :

$$\begin{aligned}
& (\llbracket \mathbf{b} \rrbracket)[\ell/\alpha] \\
&= (\mathbf{b}^\perp)[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
&= (\mathbf{b}^\perp) \\
&= \llbracket \mathbf{b} \rrbracket \\
&= \llbracket (\mathbf{b}[\ell/\alpha]) \rrbracket
\end{aligned}$$

2.  $\tau = \text{unit}$ :

$$\begin{aligned}
& (\llbracket \text{unit} \rrbracket)[\ell/\alpha] \\
&= (\text{unit}^\perp)[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
&= (\text{unit}^\perp) \\
&= \llbracket \text{unit} \rrbracket \\
&= \llbracket (\text{unit}[\ell/\alpha]) \rrbracket
\end{aligned}$$

3.  $\tau = \tau_1 \rightarrow \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 \rightarrow \tau_2 \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
= & \llbracket (\tau_1[\ell/\alpha] \rightarrow \tau_2[\ell/\alpha]) \rrbracket \\
= & \llbracket (\tau_1 \rightarrow \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$

4.  $\tau = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 \times \tau_2 \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] \times \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] \times \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
= & \llbracket (\tau_1[\ell/\alpha] \times \tau_2[\ell/\alpha]) \rrbracket \\
= & \llbracket (\tau_1 \times \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$

5.  $\tau = \tau_1 + \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 + \tau_2 \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] + \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] + \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
= & \llbracket (\tau_1[\ell/\alpha] + \tau_2[\ell/\alpha]) \rrbracket \\
= & \llbracket (\tau_1 + \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$

6.  $\tau = \text{ref } \ell_i \tau_i$ :

$$\begin{aligned}
& (\llbracket \text{ref } \ell_i \tau_i \rrbracket)[\ell/\alpha] \\
= & (\text{ref } (\llbracket \tau_i \rrbracket + \text{unit}^{\ell_i})^\perp)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\text{ref } (\llbracket \tau_i \rrbracket + \text{unit}^{\ell_i})^\perp)^\perp && \text{Lemma 3.4} \\
= & \llbracket (\text{ref } \ell_i \tau_i) \rrbracket && \text{since } \vdash \tau \text{ } WF \\
= & \llbracket (\text{ref } \ell_i \tau_i)[\ell/\alpha] \rrbracket
\end{aligned}$$

7.  $\tau = \forall \alpha. \tau_i$ :

$$\begin{aligned}
& (\llbracket \forall \alpha. \tau_i \rrbracket)[\ell/\alpha] \\
= & (\forall \alpha. (\top, \llbracket \tau_i \rrbracket))^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\forall \alpha. (\top, \llbracket \tau_i \rrbracket[\ell/\alpha]))^\perp \\
= & (\forall \alpha. (\top, \llbracket \tau_i \rrbracket[\ell/\alpha]))^\perp && \text{IH} \\
= & (\forall \alpha. \tau_i[\ell/\alpha]) \\
= & (\forall \alpha. \tau_i)[\ell/\alpha]
\end{aligned}$$

8.  $\tau = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& (\llbracket c \Rightarrow \tau_i \rrbracket)[\ell/\alpha] \\
= & (c \overset{\top}{\Rightarrow} \llbracket \tau_i \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (c[\ell/\alpha] \overset{\top}{\Rightarrow} \llbracket \tau_i \rrbracket[\ell/\alpha])^\perp \\
= & (c[\ell/\alpha] \overset{\top}{\Rightarrow} \llbracket \tau_i \rrbracket[\ell/\alpha])^\perp && \text{IH} \\
= & (c[\ell/\alpha] \Rightarrow \tau_i[\ell/\alpha]) \\
= & (c \Rightarrow \tau_i)[\ell/\alpha]
\end{aligned}$$

9.  $\tau = \text{Labeled } \ell_i \tau_i$ :

$$\begin{aligned}
& \llbracket \text{Labeled } \ell_i \tau_i \rrbracket [\ell/\alpha] \\
= & \llbracket \tau_i \rrbracket + \text{unit}^{\ell_i} [\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \llbracket \tau_i \rrbracket [\ell/\alpha] + \text{unit}^{\ell_i} [\ell/\alpha] \\
= & \llbracket \tau_i [\ell/\alpha] \rrbracket + \text{unit}^{\ell_i} [\ell/\alpha] && \text{IH} \\
= & \llbracket (\text{Labeled } \ell_i [\ell/\alpha] \tau_i [\ell/\alpha]) \rrbracket \\
= & \llbracket (\text{Labeled } \ell_i \tau_i) [\ell/\alpha] \rrbracket
\end{aligned}$$

10.  $\tau = \text{SLIO } \ell_i \ell_o \tau_i$ :

$$\begin{aligned}
& \llbracket \text{SLIO } \ell_i \ell_o \tau_i \rrbracket [\ell/\alpha] \\
= & (\text{unit}^{\ell_i} \xrightarrow{\ell_i} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_o})^\perp [\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\text{unit}^{\ell_i} \xrightarrow{\ell_i} \llbracket \tau_i \rrbracket [\ell/\alpha] + \text{unit})^{\ell_o} [\ell/\alpha]^\perp \\
= & (\text{unit}^{\ell_i} \xrightarrow{\ell_i} \llbracket \tau_i [\ell/\alpha] \rrbracket + \text{unit})^{\ell_o} [\ell/\alpha]^\perp && \text{IH} \\
= & (\text{SLIO } \ell_i [\ell/\alpha] \ell_o [\ell/\alpha] \tau_i [\ell/\alpha]) \\
= & (\text{SLIO } \ell_i \ell_o \tau_i) [\ell/\alpha]
\end{aligned}$$

□

### 3.1.3 Model for SLIO\* to FG translation

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \leftrightarrow \text{Loc}))$

**Definition 3.7** (SLIO\*  $\rightsquigarrow$  FG:  ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$   
 $\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

**Definition 3.8** (SLIO\*  $\rightsquigarrow$  FG:  $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$   
 $\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

**Definition 3.9** (SLIO\*  $\rightsquigarrow$  FG: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\
\llbracket \text{unit} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\
& \quad ({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\
& \quad \{({}^s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'}, \\
& \quad \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \forall \alpha. \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \Lambda e_s, \Lambda e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket c \Rightarrow \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \nu e_s, \nu e_t) \mid \mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \text{ref } \ell \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s a, {}^t a) \mid {}^s\theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\
\llbracket \text{Labeled } \ell \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \\
& \quad \exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in \llbracket \tau \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \text{SLIO } \ell_1 \ell_2 \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
& \quad (k, H_s, H_t) \triangleright ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \\
& \quad \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge \\
& \quad \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \rrbracket_V^{\hat{\beta}''}\}
\end{aligned}$$

**Definition 3.10** (SLIO\*  $\rightsquigarrow$  FG: Unary expression relation).

$$\begin{aligned} \llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv.e_s \Downarrow_i {}^sv \implies \\ &\exists H'_t, {}^tv. (H_t, e_t) \Downarrow (H'_t, {}^tv) \wedge ({}^s\theta, n - i, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta\} \end{aligned}$$

**Definition 3.11** (SLIO\*  $\rightsquigarrow$  FG: Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a) \rrbracket_V^{\hat{\beta}} \end{aligned}$$

**Definition 3.12** (SLIO\*  $\rightsquigarrow$  FG: Label substitution).  $\sigma : Lvar \mapsto Label$

**Definition 3.13** (SLIO\*  $\rightsquigarrow$  FG: Value substitution to values).  $\delta^s : Var \mapsto Val, \delta^t : Var \mapsto Val$

**Definition 3.14** (SLIO\*  $\rightsquigarrow$  FG: Unary interpretation of  $\Gamma$ ).

$$\begin{aligned} \llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ &\forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\} \end{aligned}$$

### 3.1.4 Soundness proof for SLIO\* to FG translation

**Lemma 3.15** (SLIO\*  $\rightsquigarrow$  FG: Monotonicity).  $\forall {}^s\theta, {}^s\theta', n, {}^sv, {}^tv, n', \beta, \beta'$ .

$$({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}$$

*Proof.* Proof by induction on  $\tau$

1. Case b:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}}$  therefore from Definition 3.9 we know that  ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 3.9  ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$  we get the desired

2. Case unit:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}}$  therefore from Definition 3.9 we know that  ${}^sv \in [\mathbf{unit}] \wedge {}^tv \in [\mathbf{unit}]$

Therefore from Definition 3.9  ${}^sv \in [\mathbf{unit}] \wedge {}^tv \in [\mathbf{unit}]$  we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.9 we know that  ${}^sv = ({}^sv_1, {}^sv_2)$  and  ${}^tv = ({}^tv_1, {}^tv_2)$ .

We also know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\underline{\text{IH1:}} ({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\underline{\text{IH2:}} ({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 3.9, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.9 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

$$\underline{\text{IH:}} ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 3.9 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \rightarrow \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.9 we know that

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 3.9 we are required to prove

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some  ${}^s\theta'_1 \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$  s.t.  $({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}}$  and we are required to prove

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with  ${}^s\theta'_1, {}^sv_2, {}^tv_2, j, \hat{\beta}''$  since  ${}^s\theta'_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

## 6. Case $\forall\alpha.\tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\forall\alpha.\tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\forall\alpha.\tau]_V^{\hat{\beta}'}$$

From Definition 3.9 we know that  ${}^sv = \Lambda e'_s$  and  ${}^tv = \Lambda e'_t$ . And

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'' . ({}^s\theta'', j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''} \quad (\text{F0})$$

Similarly from Definition 3.9 we are required to prove

$$\forall {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

This means we are given some  ${}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

Instantiating (F0) with  ${}^s\theta''_1, j, \hat{\beta}''_1$  since  ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$  therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

## 7. Case $c \Rightarrow \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [c \Rightarrow \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \Rightarrow \tau]_{\mathcal{V}}^{\hat{\beta}'}$$

From Definition 3.9 we know that  ${}^s v = \nu(e'_s)$  and  ${}^t v = \nu(e'_t)$ . And

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \hat{\beta}' \sqsubseteq \hat{\beta}_1'' . ({}^s\theta'', j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}'} \quad (\text{C0})$$

Similarly from Definition 3.9 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta_1'' \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}_1'' . ({}^s\theta_1'', j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

This means we are given some  $\mathcal{L} \models c, {}^s\theta_1'' \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}_1''$

and we are required to prove

$$({}^s\theta_1'', j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

Since  $\mathcal{L} \models c$  and instantiating (C0) with  ${}^s\theta_1'', j, \hat{\beta}_1''$  since  ${}^s\theta_1'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}_1''$  therefore we get

$$({}^s\theta_1'', j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

8. Case ref  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \ell \tau]_{\mathcal{V}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \tau]_{\mathcal{V}}^{\hat{\beta}'}$$

From Definition 3.9 we know that  ${}^s v = {}^s a$  and  ${}^t v = {}^t a$ . We also know that

$${}^s\theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}$$

From Definition 3.9, Definition 3.7 and Definition 3.8 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \tau]_{\mathcal{V}}^{\hat{\beta}'}$$

9. Case Labeled  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\hat{\beta}'}$$

From Definition 3.9 it means

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_{\ell}({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, n, {}^s v', {}^t v') \in [\tau]_{\mathcal{V}}^{\hat{\beta}}$$

$$\underline{\text{IH:}} ({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_{\mathcal{V}}^{\hat{\beta}'}$$

Similarly from Definition 3.9 we need to prove that

$$\exists s v'', t v''. s v = \text{Lb}_\ell(s v'') \wedge t v = \text{inl } t v'' \wedge (s \theta', n', s v'', t v'') \in [\tau]_V^{\hat{\beta}}$$

We choose  $s v''$  as  $s v'$  and  $t v''$  as  $t v'$  and since from IH we know that  $(s \theta', n', s v', t v') \in [\tau]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we get

$$(s \theta', n', s v, t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

10. Case SLIO  $\ell_1 \ell_2 \tau$ :

Given:

$$(s \theta, n, s v, t v) \in [\text{SLIO } \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge s \theta \sqsubseteq s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$(s \theta', n', s v, t v) \in [\text{SLIO } \ell_1 \ell_2 \tau]_V^{\hat{\beta}'}$$

This means from Definition 3.9 we know that

$$\forall s \theta_e \sqsupseteq s \theta, H_s, H_t, i, s v', t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1.$$

$$(k, H_s, H_t) \triangleright^{\hat{\beta}_1} (s \theta_e) \wedge (H_s, s v) \Downarrow_i^f (H'_s, s v') \wedge i < k \implies$$

$$\exists t v'. (H_t, t v()) \Downarrow (H'_t, t v') \wedge \exists s \theta' \sqsupseteq s \theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}_2} s \theta' \wedge$$

$$\exists t v''. t v' = \text{inl } t v'' \wedge (s \theta', t \theta', k - i, s v', t v'') \in [\tau]_V^{\hat{\beta}_2} \wedge$$

$$(\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. s \theta_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(s \theta') / \text{dom}(s \theta_e). s \theta'(a) \searrow \ell_1) \quad (\text{CG0})$$

Similarly from Definition 3.9 we need to prove

$$\forall s \theta'_e \sqsupseteq s \theta', H'_s, H'_t, i', s v'', t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1.$$

$$(k', H'_s, H'_t) \triangleright^{\hat{\beta}'_1} (s \theta'_e) \wedge (H'_s, s v) \Downarrow_i^f (H''_s, s v'') \wedge (H'_t, t v()) \Downarrow (H''_t, t v'') \wedge i' < k' \implies$$

$$\exists t v''. (H'_t, t v()) \Downarrow (H''_t, t v'') \wedge \exists s \theta'' \sqsupseteq s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} s \theta'' \wedge$$

$$\exists t v'''. t v' = \text{inl } t v''' \wedge (s \theta', k' - i, s v', t v''') \in [\tau]_V^{\hat{\beta}'_2} \wedge$$

$$(\forall a. H'_s(a) \neq H''_s(a) \implies \exists \ell'. s \theta'_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(s \theta') / \text{dom}(s \theta_e). s \theta'(a) \searrow \ell_1)$$

This means we are given some  $s \theta'_e \sqsupseteq s \theta', H'_s, H'_t, i', s v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  s.t.  $(k', H'_s, H'_t) \triangleright (s \theta'_e) \wedge (H'_s, s v) \Downarrow_i^f (H''_s, s v'') \wedge i' < k'$

And we need to prove

$$\exists t v''. (H'_t, t v()) \Downarrow (H''_t, t v'') \wedge \exists s \theta'' \sqsupseteq s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} s \theta'' \wedge$$

$$\exists t v'''. t v' = \text{inl } t v''' \wedge (s \theta', k' - i, s v', t v''') \in [\tau]_V^{\hat{\beta}'_2} \wedge$$

$$(\forall a. H'_s(a) \neq H''_s(a) \implies \exists \ell'. s \theta'_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(s \theta') / \text{dom}(s \theta_e). s \theta'(a) \searrow \ell_1)$$

Instantiating (CG0) with  $s \theta'_e \sqsupseteq s \theta', H'_s, H'_t, i', s v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  we get the desired

□

**Lemma 3.16** (SLIO\*  $\rightsquigarrow$  FG: Unary monotonicity for  $\Gamma$ ).  $\forall s \theta, s \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'.$

$$(s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge s \theta \sqsubseteq s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (s \theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$



*Proof.* Given:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 3.14 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 3.14 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 3.15 we get

$$\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

□

**Lemma 3.17** (SLIO\*  $\rightsquigarrow$  FG: Unary monotonicity for  $H$ ).  $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$

From Definition 3.11 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 3.11 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta}' \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta}' \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 3.15 we get

$$\forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

□

**Theorem 3.18** (SLIO\*  $\rightsquigarrow$  FG: Fundamental theorem).  $\forall \Gamma, \tau, e, \delta^s, \delta^t, \sigma, {}^s\theta, n.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}} \\ & \implies \\ & ({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}} \end{aligned}$$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. CF-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\} \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 3.10 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $x \delta^s \Downarrow_i {}^s v$   
From SLIO\*-Sem-val we know that  $i = 0, {}^s v = x \delta^s$ .

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-V0})$$

From fg-val we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$ . So we are left with proving

$$({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we are given  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\} \sigma]_V^{\hat{\beta}}$ , therefore from Definition 3.14 we get

$$({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_V^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ in the context. So we are done.}$$

2. CF-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{lam}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 3.10 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n, {}^s v$  s.t  $(\lambda x.e_s) \delta^s \Downarrow_i {}^s v$

From SLIO\*-Sem-val and fg-val we know that  ${}^s v = (\lambda x.e_s) \delta^s, {}^t v = (\lambda x.e_t) \delta^t, H'_t = H_t$  and  $i = 0$

It suffices to prove that

$$({}^s \theta, n, (\lambda x.e_s) \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V'}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

We know  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\lambda x.e_s) \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V'}^{\hat{\beta}}$$

From Definition 3.9 it suffices to prove

$$\begin{aligned} \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'} \\ \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2 \sigma]_{E'}^{\hat{\beta}'} \end{aligned}$$

This means that we are given  ${}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'}$

And we need to prove

$$({}^s \theta', j, e_s[{}^s v/x] \delta^s, e_t[{}^t v/x] \delta^t) \in [\tau_2 \sigma]_{E'}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V'}^{\hat{\beta}}$  therefore from Lemma 3.16 we also have

$$({}^s \theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{V'}^{\hat{\beta}'}$$

IH:

$$({}^s \theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_1\}, e_t \cup \{x \mapsto {}^t v_1\}) \in [\tau_2 \sigma]_{E'}^{\hat{\beta}'} \text{ s.t}$$

$$({}^s \theta', j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'}$$

We get (F-L0) directly from IH

### 3. CF-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V'}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2 \sigma]_{E'}^{\hat{\beta}}$

This means from Definition 3.10 it suffices to prove

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1} \ e_{t2}) \ \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau_2 \ \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-A0)

IH1:

$$({}^s \theta, n, e_{s1} \ \delta^s, e_{t1} \ \delta^t) \in [(\tau_1 \rightarrow \tau_2) \ \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.10 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. e_{s1} \ \delta^s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \ \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \ \sigma]_{V}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1} \ e_{s2}) \ \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_{s1} \ \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \ \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \ \sigma]_{V}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-A1)

IH2:

$$({}^s \theta, n - j, e_{s2} \ \delta^s, e_{t2} \ \delta^t) \in [\tau_1 \ \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.10 it suffices to prove

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall k < n - j, {}^s v_2. e_{s2} \ \delta^s \Downarrow_i {}^s v_2 \implies$$

$$\exists H'_{t2}, {}^t v_2. (H_t, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \ \sigma]_{V}^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta'$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $(e_{s1} \ e_{s2}) \ \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s2} \ \delta^s \Downarrow_k {}^s v_2$ .

And we have

$$\exists H'_{t2}, {}^t v_2. (H_t, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \ \sigma]_{V}^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-A2)

Since from (F-A1) we know that  $({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \ \sigma]_{V}^{\hat{\beta}}$  where

$${}^s v_1 = \lambda x. e'_s \text{ and } {}^t v_1 = \lambda x. e'_t$$

From Definition 3.9 we have

$$\forall {}^s \theta'_3 \sqsupseteq {}^s \theta, {}^s v, {}^t v, l < n - j, \hat{\beta}_3 \sqsupseteq \hat{\beta}. ({}^s \theta'_3, l, {}^s v, {}^t v) \in [\tau_1 \ \sigma]_{V}^{\hat{\beta}_3}$$

$$\implies ({}^s \theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2 \ \sigma]_E^{\hat{\beta}_3}$$

Instantiating with  ${}^s \theta, {}^s v_2, {}^t v_2, n - j - k, \hat{\beta}$  we get

$$({}^s \theta, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \ \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (n - j - k, H_{s4}, H_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall k' < n - j - k, {}^s v_4. e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ & \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s \theta, n - j - k - k', {}^s v_4, {}^t v_4) \in [\tau_2 \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - j - k - k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t2}$ , from (F-A2) we know that  $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta$ . Instantiating  ${}^s v_4$  with  ${}^s v$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k' < i - j - k < n - j - k$  s.t  $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$ . therefore we have

$$\begin{aligned} & \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s \theta, n - j - k - k', {}^s v, {}^t v_4) \in [\tau_2 \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j - k - \\ & k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-A3}) \end{aligned}$$

Since from SLIO\*-Sem-app we know that  $i = j + k + k'$  and  $H'_t = H'_{t4}$ ,  ${}^t v = {}^t v_4$  therefore we get (F-A0) from (F-A3) and Lemma 3.15 and Lemma 3.17

#### 4. CF-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 3.10 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t, \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$  and given some  $i < n$  s.t  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \\ & (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t_1}, {}^t v_1. (H_{t_1}, e_{t_1} \delta^t) \Downarrow (H'_{t_1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t_1}) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-P1)

IH2:

$$({}^s \theta, n - j, e_{s_2} \delta^s, e_{t_2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\forall H_{s_2}, H_{t_2}. (n, H_{s_2}, H_{t_2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall k < n - j. e_{s_2} \delta^s \Downarrow_k {}^s v_2 \implies$$

$$\exists H'_{t_2}, {}^t v_2. (H_{t_2}, e_{t_2} \delta^t) \Downarrow (H'_{t_2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_{s_2}, H'_{t_2}) \triangleright^{\hat{\beta}} {}^s \theta$$

Instantiating with  $H_s, H'_{t_1}, \hat{\beta}'_1$  and since we know that  $(e_{s_1}, e_{s_2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s_2} \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t_2}, {}^t v_2. (H_{t_2}, e_{t_2} \delta^t) \Downarrow (H'_{t_2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t_2}) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-P2)

From SLIO\*-Sem-prod we know that  $i = j + k + 1$ ,  $H'_t = H'_{t_2}$  and  ${}^t v = ({}^t v_1, {}^t v_2)$  therefore from Definition 3.9 and Lemma 3.15 we get  $({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$

And since we have  $(n - j - k, H_s, H'_{t_2}) \triangleright^{\hat{\beta}} {}^s \theta$  therefore from Lemma 3.17 we also get

$$(n - i, H_s, H'_{t_2}) \triangleright^{\hat{\beta}} {}^s \theta$$

5. CF-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$  (F-F0)

This means from Definition 3.10 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n. {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

(F-F0)

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{E}}^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j ({}^s v_1, -) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and  ${}^s v_1$  with  ${}^s v$  since we know that  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j ({}^s v, -)$ .

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F1}) \end{aligned}$$

From SLIO\*-Sem-fst we know that  $i = j + 1$ ,  $H'_t = H'_{t1}$  and  ${}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}}$  therefore from Definition 3.9 and Lemma 3.15 we get

$$({}^s\theta, n - i, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}}$$

And since from (F-F1) we have  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 3.17 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

## 6. CF-snd:

Symmetric reasoning as in the CF-fst case

## 7. CF-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\tau_1 + \tau_2) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 3.10 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v.\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright {}^s \theta \quad (\text{F-IL0})$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright {}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright {}^s \theta$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright {}^s \theta \quad (\text{F-IL1})$$

From SLIO\*-Sem-inl we know that  $i = j + 1$  and  $H'_t = H'_{t1}, {}^t v = {}^t v_1$ . Since we know  $({}^s \theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from Definition 3.9 and Lemma 3.15 we get

$$({}^s \theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}$$

And since from (F-IL1) we have  $(n - j, H_s, H'_{t1}) \triangleright {}^s \theta$  therefore from Lemma 3.17 we get

$$(n - i, H_s, H'_{t1}) \triangleright {}^s \theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{case}(e_t, x.e_{t1}, y.e_{t2})} \text{ case}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 3.10 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright {}^s \theta \wedge \forall i < n, {}^s v. \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright {}^s \theta$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright {}^s \theta$  and given some  $i < n$  s.t  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$



And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-C0})$$

IH1:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_{V'}^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_{V'}^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-C1})$$

Two cases arise:

- (a)  ${}^s v_1 = \text{inl}({}^s v'_1)$  and  ${}^t v_1 = \text{inl}({}^t v'_1)$ :

IH2:

$$({}^s \theta, n-j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall k < n-j, {}^s v_2.e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n-j-k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}} \wedge (n-j-k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i-j < n-j$  s.t  $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n-j-k, {}^s v, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}} \wedge (n-j-k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$$

From SLIO\*-Sem-case1 we know that  $i = j+k+1$  and  $H'_t = H'_{t2}$ ,  ${}^t v = {}^t v_2$ . Since we know  $({}^s \theta, n-j-k, {}^s v, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}}$  therefore from Definition 3.9 and Lemma 3.15 we get

$$({}^s \theta, n-i, {}^s v, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}}$$

And since from (F-C2) we have  $(n-j-k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$  therefore from Lemma 3.17 we get  $(n-i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$

(b)  ${}^s v_1 = \text{inr}({}^s v'_1)$  and  ${}^t v_1 = \text{inr}({}^t v'_1)$ :

Symmetric reasoning as in the previous case

10. CF-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \Lambda e_s : \forall \alpha. \tau \rightsquigarrow \Lambda e_t} \text{FI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \Lambda e_s \delta^s, \Lambda e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 3.10 we know that

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \Lambda e_s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \Lambda e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\forall \alpha. \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t.  $(\Lambda e_s) \delta^s \Downarrow_i {}^s v$

From SLIO\*-Sem-val and fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s$ ,  ${}^t v = (\Lambda e_t) \delta^t$ ,  $i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$$

From Definition 3.9 it suffices to prove

$$\forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}'}$$

This means that we are given  ${}^s \theta' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}'} \quad (\text{F-FI0})$$

Since  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$  therefore from Lemma 3.16 we also have

$$({}^s \theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}'}$$

IH:

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_{\hat{E}}^{\hat{\beta}'}$$

We get (F-FI0) directly from IH

11. CF-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \forall \alpha. \tau \rightsquigarrow e_t \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e_s [] : \tau[\ell/\alpha] \rightsquigarrow e_t[]} \text{FE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s [] \delta^s, e_t [] \delta^t) \in [\tau[\ell/\alpha] \sigma]_E^{\hat{\beta}}$

From Definition 3.10 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s [] \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $e_s [] \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.10 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s []) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n, {}^s v_1$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-FE1})$$

From SLIO\*-Sem-FE we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_1 = \Lambda e'_t$

Therefore we have

$$({}^s\theta, n - j, \Lambda e'_s, \Lambda e'_t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}}$$

This means from Definition 3.9 we have

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s\theta', k, e'_s, e'_t) \in [\tau[\ell'/\alpha] \sigma]_E^{\hat{\beta}_2}$$

Instantiating  ${}^s\theta'$  with  ${}^s\theta$ ,  $k$  with  $n - j - 1$ ,  $\ell'$  with  $\ell$   $\sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$$({}^s\theta, n - j - 1, e'_s, e'_t) \in [\tau[\ell/\alpha] \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we get

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n-j-1, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}_2} s\theta'_1 \wedge \forall k < n-j-1, {}^s v_2. e'_s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n-j-1-k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n-j-1-k, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-FE1) we know that  $(n-j, H_s, H'_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta$  therefore from Lemma 3.17 we get  $(n-j-1, H_s, H'_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta$

Since we know that  $e_s \Downarrow \delta^s \Downarrow_i {}^s v$  and from SLIO\*-Sem-FE we know that  $i = j+k+1$  (for some  $k$ ) and  $i < n$  therefore we have  $k < n-j-1$  s.t.  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n-j-1-k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n-j-1-k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-FE2})$$

Since  $H'_t = H_{t2'}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-FE0) directly from (F-FE2)

12. CF-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \nu e_s : c \Rightarrow \tau \rightsquigarrow \nu e_t} \text{ CI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \nu e_s \delta^s, \nu e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_E^{\hat{\beta}}$

This means from Definition 3.10 we know that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall i < n. \nu e_s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \nu e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [(c \Rightarrow \tau) \hat{\beta} \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$  and given some  $i < n$  s.t.  $(\nu e_s) \delta^s \Downarrow_i {}^s v$

From SLIO\*-Sem-val and fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  ${}^t v = (\nu e_t) \delta^t$ ,  $i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$$

We know  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$$

From Definition 3.9 it suffices to prove

$$\mathcal{L} \models c \sigma \implies \forall s\theta' \sqsupseteq s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}'}$$

This means that we are given  $\mathcal{L} \models c \sigma$  and  ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'} \quad (\text{F-CI0})$$

Since  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$  therefore from Lemma 3.16 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}'}$$

And since we know that  $\mathcal{L} \models c \sigma$  therefore

$$\underline{\text{IH:}} ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

### 13. CF-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : c \Rightarrow \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e_s \bullet : \tau \rightsquigarrow e_t \bullet} \text{CE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s \bullet \delta^s, e_t \bullet \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 3.10 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s \bullet \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $e_s \bullet \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.10 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s \bullet) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t.  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE1})$$

From SLIO\*-Sem-CE we know that  ${}^s v_1 = \nu e'_s$  and  ${}^t v_1 = \nu e'_t$

Therefore we have

$$({}^s \theta, n - j, \nu e'_s, \nu e'_t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$$

This means from Definition 3.9 we have

$$\forall {}^s \theta' \sqsubseteq {}^s \theta'_1, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s \theta', k, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}_2}$$

Instantiating  ${}^s \theta'$  with  ${}^s \theta$ ,  $k$  with  $n - j - 1$ ,  $\ell'$  with  $\ell$   $\sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$$({}^s \theta, n - j - 1, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we get

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s \theta'_1 \wedge \forall k < n - j - 1. e'_s \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-CE1) we know that  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$  therefore from Lemma 3.17 we get  $(n - j - 1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$

Since we know that  $e_s \bullet \delta^s \Downarrow_i {}^s v$  and from SLIO\*-Sem-CE we know that  $i = j + k + 1$  (for some  $k$ ) and  $i < n$  therefore we have  $k < n - j - 1$  s.t  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-CE2})$$

Since  $H'_t = H_{t2'}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-CE0) directly from (F-CE2)

14. CF-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e_s) : \text{SLIO } \ell_i \ell_i \tau \rightsquigarrow \lambda_{\cdot} \text{inl}(e_t)} \text{ret}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot} \text{inl}(e_t) \delta^t) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_E^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n. {}^s v. \text{ret}(e_s) \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \lambda_{\cdot} \text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t  $\text{ret}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot} \text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From SLIO\*-ret and FG-lam we know that  $i = 0$ ,  ${}^s v = \text{ret}(e_s) \delta^s$ ,  ${}^t v = \lambda_{-}.\text{inl}(e_t) \delta^t$  and  $H'_t = H_t$ .

So we need to prove

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [\text{SLIO } \ell_i \ell_i \tau \sigma]_{V}^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k &\implies \exists H'_t, {}^t v'. (H_t, (\lambda_{-}.\text{inl}(e_t) ()) \delta^t) \Downarrow \\ (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k. \text{ Also from SLIO*}-\text{Sem-ret we know that } H'_s = H_s$$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v'. (H_t, (\lambda_{-}.\text{inl}(e_t) ()) \delta^t) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s, H'_t) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-R0}) \end{aligned}$$

**IH:**

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies \\ \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s \theta_e, k - f, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s1}$  with  $H_s$  and  $H_{t1}$  with  $H_t$ . And since we know that  $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ . Therefore we have

$$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s \theta_e, k - f, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'} \wedge (k - f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-R1})$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_{t1}$ ,  ${}^t v'$  as  $\text{inl}({}^t v)$ ,  ${}^s \theta'$  as  ${}^s \theta_e$ ,  $\hat{\beta}''$  as  $\hat{\beta}'$ . Since from SLIO\*-Sem-ret we know that  $i = f + 1$  therefore from (F-R1) and Lemma 3.17 we know that  $(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

Next we choose  ${}^t v''$  as  ${}^t v$  (from F-R1) and from Lemma 3.15 we get  $({}^s \theta_e, k - i, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'}$  (we know from SLIO\*-Sem-ret that  ${}^s v' = {}^s v$ )

15. CF-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \text{SLIO } \ell_i \ell \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_{s2} : \text{SLIO } \ell \ell_o \tau' \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \text{SLIO } \ell_i \ell_o \tau' \rightsquigarrow \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{bind}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\text{SLIO } \ell_i \ell_o \tau') \sigma]_E^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\text{SLIO } \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\text{SLIO } \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

From SLIO\*-Sem-val and fg-val we know that  $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s, {}^t v = \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\text{SLIO } \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\text{SLIO } \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'. (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'. (k - \\ & i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \quad (\text{F-B0}) \end{aligned}$$



IH1:

$$({}^s\theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_{E}^{\hat{\beta}}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\ \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k - \\ j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{SLIO } \ell_i \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k - \\ j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-B1.1})$$

From Definition 3.9 we know have

$$\begin{aligned} \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s\theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s\theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s\theta_e$  with  ${}^s\theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s\theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s\theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s\theta'', k - j - b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_{E}^{\hat{\beta}''}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}''} {}^s\theta \wedge \forall c < (k - j - b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s\theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_{V}^{\hat{\beta}''} \wedge \\ (k - j - b - c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s\theta'' \end{aligned}$$

Instantiating  $H_{s4}$  with  $H'_{s3}$  and  $H_{t4}$  with  $H'_{t3}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists c < i - j - b < k - j - b$  s.t  $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s\theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{SLIO } \ell \ell_o \tau') \sigma]_{V}^{\hat{\beta}''} \wedge \\ (k - j - b - c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s\theta'' \quad (\text{F-B2.1}) \end{aligned}$$

From Definition 3.9 we know have

$$\begin{aligned} \forall^s \theta_e \sqsupseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k - j - b - c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1. \\ (m, H_{s5}, H_{t5}) \hat{\triangleright}^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}''_2. (m - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} {}^s \theta''' \wedge \\ \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta''$ ,  $H_{s5}$  with  $H'_{s3}$ ,  $H_{t5}$  with  $H'_{t3}$ ,  $m$  with  $k - j - b - c$  and  $\hat{\beta}''_1$  with  $\hat{\beta}''$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists d < i - j - b - c < k - j - b - c$  s.t.  $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$ .

Therefore we have

$$\begin{aligned} \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}''_2. (k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} {}^s \theta''' \wedge \\ \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose  $H'_{t1}$  as  $H'_{t5}$  and  ${}^t v'$  as  ${}^t v'_{h2}$ . Next we choose  ${}^s \theta'$  as  ${}^s \theta'''$  and  $\hat{\beta}''$  as  $\hat{\beta}''_2$  (both chosen from (F-B2)). Also from SLIO\*-Sem-bind we know that in (F-B0)  $H'_{s1}$  will be  $H'_{s5}$ .

Since  $(k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} {}^s \theta'''$  therefore Lemma 3.15 we get  $(k - i, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} {}^s \theta'''$ . Also since from (F-B2) we have  $\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$

Since  $i = j + b + c + d + 1$  therefore from Lemma 3.15 we get

$$\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$$

16. CF-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{ label}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 3.10 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ (n - i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t.  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ .

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 3.10 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove  $({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}}$ . Since from SLIO\*-Sem-label we know that  $i = j + 1$ ,  ${}^s v = {}^s v_1$  and  ${}^t v = {}^t v_1$ . Therefore we get this from Definition 3.9, (F-LB1) and Lemma 3.15.

$$\text{From Lemma 3.15 we get } (n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

17. CF-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{SLIO } \ell_i \ell_o \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e_s) : \text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau) \rightsquigarrow \lambda_{\cdot} \text{inl}(e_t ())} \text{toLabeled}$$

$$\text{Also given is: } \mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \in [(\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_E^{\hat{\beta}}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t  $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

From SLIO\*-Sem-val and fg-val we know that  $i = 0$ ,  ${}^s v = \text{toLabeled}(e_s) \delta^s$ ,  
 ${}^t v = (\lambda\_.\text{inl } e_t()) \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda\_.\text{inl } e_t()) \delta^t) \in [(\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda\_.\text{inl } e_t()) \delta^t) \in [(\text{SLIO } \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{V}^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_.\text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge$$

$$\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{V}^{\hat{\beta}''}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda\_.\text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge$$

$$\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-TL0})$$

III:

$$({}^s \theta, k, e_s \delta^s, e_t \delta^t) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_E^{\hat{\beta}}$$

It means from Definition 3.10 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies$$

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k -$$

$$j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{SLIO } \ell_i \ell_o \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k -$$

$$j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-TL1.1})$$

From Definition 3.9 we know have

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies$$

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-TL1}) \end{aligned}$$

In order to prove (F-TL0) we choose  ${}^s \theta'$  as  ${}^s \theta''$  and  $\hat{\beta}'$  as  $\hat{\beta}''$  (both chosen from (F-TL2))

Also from SLIO\*-Sem-toLabeled and fg-inl, fg-app we know that  $H'_s = H'_{s3}$  and  $H'_t = H'_{t3}$ , and  ${}^s v' = {}^s v'_{h1}$ ,  ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 3.15

18. CF-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell \ \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e_s) : \text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \rightsquigarrow \lambda_{\cdot} e_t} \text{unlabel}$$

Also given is:  $\mathcal{L} \models \Psi \ \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \sigma]_E^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \\ & \text{From SLIO*-Sem-val and fg-val we know that } i = 0, {}^s v = \text{unlabel}(e_s) \delta^s, {}^t v = \lambda_{\cdot} e_t \delta^t, \\ & H'_t = H_t \end{aligned}$$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in [\text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\text{SLIO } \ell_i (\ell_i \sqcup \ell) \ \tau \sigma]_V^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\begin{aligned}
& \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^sv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
& (k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^sv') \wedge i < k \implies \\
& \exists H'_{t1}, {}^tv'. (H_{t1}, (\lambda_.e_t)() \delta^t) \Downarrow (H'_{t1}, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\
& \exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', k - i, {}^sv', {}^tv'') \in [\tau \sigma]_V^{\hat{\beta}''}
\end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^sv', {}^tv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^sv') \wedge i < k.$$

And we need to prove

$$\begin{aligned}
& \exists H'_{t1}, {}^tv'. (H_{t1}, (\lambda_.e_t)() \delta^t) \Downarrow (H'_{t1}, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\
& \exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', k - i, {}^sv', {}^tv'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-U0})
\end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned}
& \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^sv_h.e_s \delta^s \Downarrow_f {}^sv_h \implies \\
& \exists H'_{t2}, {}^tv_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^tv_h) \wedge ({}^s\theta_e, k - f, {}^sv_h, {}^tv_h) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - \\
& f, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e
\end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^sv')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^sv_h$ .

Therefore we have

$$\exists H'_{t2}, {}^tv_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^tv_h) \wedge ({}^s\theta_e, k - f, {}^sv_h, {}^tv_h) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-U1})$$

In order to prove (F-U0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^tv'$  as  ${}^tv_h$ ,  ${}^s\theta'$  as  ${}^s\theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From SLIO\*-Sem-unlabel and fg-app we also know that  $H'_{s1} = H_{s1}$  and  $H'_{t1} = H'_{t2}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e:$$

$$\text{Since from (F-U1) we know that } (k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$$

$$\text{Therefore from Lemma 3.17 we also get } (k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$$

$$(b) \exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta_e, k - i, {}^sv', {}^tv'') \in [\tau \sigma]_V^{\hat{\beta}'}$$

Since from (F-U1) we have

$$({}^s\theta_e, k - f, {}^sv_h, {}^tv_h) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}'}$$

This means from Definition 3.9 we know that

$$\exists {}^sv_i, {}^tv_i. {}^sv_h = \text{Lb}_\ell({}^sv_i) \wedge {}^tv_h = \text{inl } {}^tv_i \wedge ({}^s\theta_e, k - f - 1, {}^sv_i, {}^tv_i) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-U2})$$

Since we know that  ${}^t v' = {}^t v_h$  and since from (F-U2) we have  ${}^t v_h = \text{inl } {}^t v_i$ . Therefore from we choose  ${}^t v''$  as  ${}^t v_i$  to get the first conjunct

From SLIO\*-Sem-unlabel we know that  ${}^s v = {}^s v_i$  and since we know that  $({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$

Therefore from Lemma 3.15 we also get  $({}^s \theta_e, k - i, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$

19. CF-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e_s : \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{-} \text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-} \text{inl}(\text{new } (e_t)) \delta^t) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_E^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-} \text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From SLIO\*-Sem-val and fg-val we know that  $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda_{-} \text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-} \text{inl}(\text{new } (e_t)) \delta^t) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-} \text{inl}(\text{new } (e_t)) \delta^t) \in \llbracket \text{SLIO } \ell \ell (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\dots} \text{inl}(\text{new } e_t))(\delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k-i, {}^s v', {}^t v'') \in \llbracket (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-N0}) \end{aligned}$$

From SLIO\*-Sem-ref we know that  ${}^s v' = a_s$  and from fg-ref, fg-inl we know that  ${}^t v' = \text{inl } a_t$ .

**IIH:**

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h.e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k-f, \\ & H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-N1})$$

In order to prove (F-N0) we choose  $H'_{t1}$  as  $H'_{t2} \cup \{a_t \mapsto {}^t v_h\}$ ,  ${}^t v$  as  $a_t$ ,  ${}^s \theta'$  as  ${}^s \theta_n$  where  ${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$

And we choose  $\hat{\beta}''$  as  $\hat{\beta}_n$  where  $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From SLIO\*-Sem-ref and fg-ref we also know that  $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

We need to prove

$$(a) \ (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}_n} {}^s \theta_n:$$

From Definition 3.11 it suffices to prove that

- $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$ :

Since  $\text{dom}({}^s \theta_e) \subseteq \text{dom}(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\} \text{ and } H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$$

Therefore we get  $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (\text{dom}({}^s \theta_e) \times \text{dom}(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}, H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\} \text{ and } \hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$$

Therefore we get  $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s \theta_n, k-i-1, H'_{s1}(a_1), H'_{t1}(a_2)) \in \llbracket {}^s \theta_n(a) \rrbracket_V^{\hat{\beta}_n}$ :

$$\forall (a_1, a_2) \in \hat{\beta}_n$$



–  $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-N1) we know that  $({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'}$

From Lemma 3.15 we get  $({}^s\theta_n, k - i - 1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}_n}$

–  $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}' } {}^s\theta_e$  therefore  
from Definition 3.11 we get

$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$

From Lemma 3.15 we get

$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_n(a_1)]_{V}^{\hat{\beta}'}$

(b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}_n}$ :

We choose  ${}^t v''$  as  ${}^t v_h$  from (F-N1), fg-inl and fg-ref we know that  ${}^t v' = \text{inl } {}^t v_h$

In order to prove  $({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}_n}$ , from Definition 3.9 it suffices  
to prove that

${}^s\theta_n(a_s) = (\text{Labeled } \ell' \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}_n$

We get this by construction of  ${}^s\theta_n$  and  $\hat{\beta}_n$

20. CF-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{ref } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash !e_s : \text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{\cdot} \text{inl}(e_t)} \text{deref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e_s \delta^s, \lambda_{\cdot} \text{inl}(e_t) \delta^t) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_{E}^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot} \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  
 $!e_s \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot} \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

From SLIO\*-Sem-val and fg-val we know that  $i = 0, {}^s v = !e_s \delta^s, {}^t v = \lambda_{\cdot} \text{inl}(e_t) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{SLIO } \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, !e_s \delta^s, \lambda_{\cdot} \text{inl}(e_t) \delta^t) \in \llbracket \text{SLIO } \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \stackrel{\hat{\beta}'}{\triangleright} ({}^s\theta_e) \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} \text{inl}(e_t))(\delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}''}{\triangleright} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e \wedge (H_{s1}, !(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} \text{inl}(e_t))(\delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}''}{\triangleright} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-D0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'_1$  as  $H'_{t2}(a)$  (where  ${}^t v_h = a_t$  from fg-deref),  ${}^s\theta'$  as  ${}^s\theta_e$  and we choose  $\hat{\beta}''$  as  $\hat{\beta}'$ .

From SLIO\*-Sem-deref we also know that  $H'_{s1} = H_{s1}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e:$$

Since from (F-D1) we have  $(k - f, H_{s1}, H'_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e$  and since  $f < i$  thfore from Lemma 3.17 we get  $(k - i, H_{s1}, H'_{t2}) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

Since from SLIO\*-Sem-deref and fg-deref we know that  ${}^s v_h = a_s$  and  ${}^t v_h = a_t$ .

Therefore from (F-D1) and from Definition 3.9 we know that

$${}^s\theta_e(a_s) = (\text{Labeled } \ell \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that  $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$  which means from Definition 3.11 we know that

$$({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$
 (F-D2)

This means from Definition 3.9 we know that

$$\exists {}^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_\ell({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$$

We choose  ${}^t v''$  as  ${}^t v_i$  and we know that  ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$ . This proves the first conjunct.

Since from (F-D2) we have  $({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$  therefore from Lemma 3.15 we get

$$({}^s\theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

This proves the second conjunct.

## 21. CF-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_{s1} := e_{s2} : \text{SLIO } \ell \ell \text{ unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in \llbracket \text{SLIO } \ell \ell \text{ unit} \sigma \rrbracket_E^{\hat{\beta}}$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in \llbracket \text{SLIO } \ell \ell \text{ unit} \sigma \rrbracket_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in \llbracket \text{SLIO } \ell \ell \text{ unit} \sigma \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{s\theta}$$

From SLIO\*-Sem-val and fg-val we know that  $i = 0, {}^s v = (e_{s1} := e_{s2}) \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in \llbracket \text{SLIO } \ell \ell \text{ unit} \sigma \rrbracket_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{s\theta}$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  from the context so we are left with proving

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in \llbracket \text{SLIO } \ell \ell \text{ unit} \sigma \rrbracket_V^{\hat{\beta}}$$

From Definition 3.9 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl}(e_{t1} := e_{t2})() \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} \\ & {}^s\theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \_ . \text{inl}(e_{t1} := e_{t2})() \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ & (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s\theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - \\ & f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S1})$$

IH2:

$$({}^s\theta_e, k - f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 3.10 that we need to prove

$$\begin{aligned} & \forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k - f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} \implies \\ & \exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - \\ & f - l, H_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s3}$  with  $H_{s1}$  and  $H_{t3}$  with  $H'_{t2}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists l < i - f < k - f$  s.t  $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$ .

Therefore we have

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - f - l, H_{s1}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S2})$$

In order to prove (F-S0) we choose  $H'_{t1}$  as  $H'_{t3}[a_t \mapsto {}^t v_{h3}]$ ,  ${}^t v'$  as  $()$ ,  ${}^s \theta'$  as  ${}^s \theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$ .  
 From SLIO\*-Sem-assign and fg-assign we also know that  ${}^s v_{h2} = a_s$ ,  ${}^t v_{h2} = a_t$ ,  $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$  and  $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

(a)  $(k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ :

From Definition 3.11 it suffices to prove that

- $dom({}^s \theta_e) \subseteq dom(H'_{s1})$ :

Since  $dom({}^s \theta_e) \subseteq dom(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since  $dom(H_{s1}) = dom(H'_{s1})$  therefore we also get  
 $dom({}^s \theta_e) \subseteq dom(H'_{s1})$

- $\hat{\beta}' \subseteq (dom({}^s \theta_e) \times dom(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (dom({}^s \theta_e) \times dom(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since  $dom(H_{t1}) \subseteq dom(H'_{t1})$  therefore we also have  $\hat{\beta}' \subseteq (dom({}^s \theta_e) \times dom(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}'.({}^s \theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_V^{\hat{\beta}'}$   
 $\forall (a_1, a_2) \in \hat{\beta}'_n$

–  $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-S2) we know that  $({}^s \theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'}$

From Lemma 3.15 we get  $({}^s \theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'}$

–  $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$  therefore  
 from Definition 3.11 we get

$({}^s \theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_V^{\hat{\beta}'}$

From Lemma 3.15 we get

$({}^s \theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_V^{\hat{\beta}'}$

(b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}'_n}$ :

We choose  ${}^t v''$  as  $()$  from (F-S1), fg-inl and fg-assign we know that  ${}^t v' = \text{inl } ()$

To prove:  $({}^s \theta_n, k - i, (), ()) \in [\text{unit}]_V^{\hat{\beta}'_n}$ ,

We get this directly from Definition 3.9

□

**Lemma 3.19** (SLIO\*  $\rightsquigarrow$  FG: Subtyping). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

1.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$

2.  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. SLIO\*sub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^{\hat{\beta}}] \subseteq [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^{\hat{\beta}}]$

It suffices to prove:  $\forall ({}^s\theta, n, \lambda x.e_i) \in [\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^{\hat{\beta}}]. ({}^s\theta, n, \lambda x.e_i) \in [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^{\hat{\beta}}]$

This means that given some  ${}^s\theta, n$  and  $\lambda x.e_i$  s.t  $({}^s\theta, n, \lambda x.e_i) \in [\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^{\hat{\beta}}]$

Therefore from Definition 3.9 we are given:

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ ({}^s\theta', j, {}^sv, {}^tv) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{S-A0}) \end{aligned}$$

And it suffices to prove:  $({}^s\theta, n, \lambda x.e_i) \in [\lceil (\tau'_1 \rightarrow \tau'_2) \sigma \rceil_V^{\hat{\beta}}]$

Again from Definition 3.9 it suffices to prove:

$$\begin{aligned} \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1 \\ ({}^s\theta'_1, k, {}^sv_1, {}^tv_1) \in [\tau'_1]_V^{\hat{\beta}'_1} \implies ({}^s\theta'_1, k, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1} \end{aligned}$$

This means that given some  ${}^s\theta'_1 \sqsubseteq {}^s\theta, {}^sv_1, {}^tv_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t  $({}^s\theta'_1, k, {}^sv_1, {}^tv_1) \in [\tau'_1]_V^{\hat{\beta}'_1}$

And we are required to prove:  $({}^s\theta'_1, k, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

IH:  $[\lceil \tau'_1 \sigma \rceil_V^{\hat{\beta}'_1}] \subseteq [\lceil \tau_1 \sigma \rceil_V^{\hat{\beta}'_1}]$  (Statement (1))

$[\lceil \tau_2 \sigma \rceil_E^{\hat{\beta}'_1}] \subseteq [\lceil \tau'_2 \sigma \rceil_E^{\hat{\beta}'_1}]$  (Sub-A0, From Statement (2))

Instantiating (S-A0) with  ${}^s\theta'_1, {}^sv_1, {}^tv_1, k, \hat{\beta}'_1$

Since  $({}^s\theta'_1, k, {}^sv_1, {}^tv_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_1}$  therefore from IH1 we know that  $({}^s\theta'_1, k, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1}$

As a result we get

$$({}^s\theta'_1, k, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

From (Sub-A0), we know that

$$({}^s\theta'_1, k, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_1}$$

2. SLIO\*sub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[\lceil (\tau_1 \times \tau_2) \sigma \rceil_V^{\hat{\beta}}] \subseteq [\lceil (\tau'_1 \times \tau'_2) \sigma \rceil_V^{\hat{\beta}}]$

IH1:  $[\lceil \tau_1 \sigma \rceil_V^{\hat{\beta}}] \subseteq [\lceil \tau'_1 \sigma \rceil_V^{\hat{\beta}}]$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we are given:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s\theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 3.9, it suffices to prove:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$$

Since from (S-P0) we know that  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from IH1 we have  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$

Similarly since from (S-P0) we have  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$  therefore from IH2 we get  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}}$

### 3. SLIO\*sub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $[(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \subseteq [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

IH1:  $[(\tau_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_1 \sigma)]_V^{\hat{\beta}}$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:  $\forall ({}^s\theta, n, {}^s v, {}^t v) \in [((\tau_1 + \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in [((\tau'_1 + \tau'_2) \sigma)]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^s v, {}^t v) \in [((\tau_1 + \tau_2) \sigma)]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^s v, {}^t v) \in [((\tau'_1 + \tau'_2) \sigma)]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^s v = \text{inl } {}^s v_i$  and  ${}^t v = \text{inl } {}^t v_i$ :

From Definition 3.9 we are given:

$$({}^s\theta, n, {}^s v_i, {}^t v_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s\theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

(b)  ${}^s v = \text{inr } {}^s v_i$  and  ${}^t v = \text{inr } {}^t v_i$ :

Symmetric reasoning

4. SLIO\*sub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $[(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we are given:

$$\forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'} \quad (\text{S-F0})$$

And it suffices to prove:  $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

Again from Definition 3.9, it suffices to prove:

$$\forall {}^s \theta'_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

This means that given  ${}^s \theta_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:  $({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$

Instantiating (S-F0) with  ${}^s \theta_1, k, \ell'_1, \hat{\beta}'_1$  we get

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_1[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell']))]_E^{\hat{\beta}'_1} \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell']))]_E^{\hat{\beta}'_1} \quad (\text{Sub-F0, Statement (2)})$$

From (Sub-F0), we know that

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

5. SLIO\*sub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $[(c_1 \Rightarrow \tau_1) \sigma]_V^{\hat{\beta}} \subseteq [(c_2 \Rightarrow \tau_2) \sigma]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s \theta, n, \nu e_s, \nu e_t) \in [(c_1 \Rightarrow \tau_1) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \nu e_s, \nu e_t) \in [(c_2 \Rightarrow \tau_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s \theta, n, \nu e_s, \nu e_t) \in [(c_1 \Rightarrow \tau_1) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we are given:



$$\mathcal{L} \models c_1 \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in [\tau_1 \sigma]_E^{\hat{\beta}'} \quad (\text{S-C0})$$

And it suffices to prove:  $({}^s\theta, n, \nu e_s, \nu e_t) \in [((c_2 \Rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 3.9, it suffices to prove:

$$\mathcal{L} \models c_2 \sigma \implies \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, k, e_s, e_t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

This means that given  $\mathcal{L} \models c_2, {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

since we know that  $c_2 \implies c_1$  and since  $\mathcal{L} \models c_2 \sigma$  therefore  $\mathcal{L} \models c_1 \sigma$ . Next we instantiate (S-C0) with  ${}^s\theta'_1, k, \hat{\beta}'_1$  to get

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_1 \sigma]_E^{\hat{\beta}'_1}$$

$$[(\tau_1 \sigma)]_E^{\hat{\beta}'_1} \subseteq [(\tau_2 \sigma)]_E^{\hat{\beta}'_1} \quad (\text{Sub-C0, Statement (2)})$$

Therefore from (Sub-C0), we get

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

6. SLIO\*sub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $[((\text{Labeled } \ell \tau) \sigma)]_V^{\hat{\beta}} \subseteq [((\text{Labeled } \ell' \tau') \sigma)]_V^{\hat{\beta}}$

IH:  $[(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell \tau) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^{\hat{\beta}}$$

This means that given some  $({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell \tau) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we are given:

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in [\tau \sigma]_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^{\hat{\beta}}$$

From Definition 3.9 it suffices to prove

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in [\tau' \sigma]_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

7. SLIO\*sub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_1 \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell'_2}{\Sigma; \Psi \vdash \text{SLIO } \ell_1 \ell_2 \tau <: \text{SLIO } \ell'_1 \ell'_2 \tau'}$$

To prove:  $[\text{((SLIO } \ell_i \ell_2 \tau) \sigma)]_V^{\hat{\beta}} \subseteq [\text{((SLIO } \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$

It suffices to prove:

$$\forall {}^s\theta, n, {}^s v, {}^t v \in [\text{((SLIO } \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}. \quad {}^s\theta, n, {}^s v, {}^t v \in [\text{((SLIO } \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, {}^s v, {}^t v) \in [\text{((SLIO } \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 3.9 we are given:

$$\begin{aligned} \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}' \\ (k, H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \\ \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{S-M0}) \end{aligned}$$

And we are required to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{((SLIO } \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

So again from Definition 3.9 we need to prove

$$\begin{aligned} \forall {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1 \\ (k_1, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1 \implies \\ \exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''_1} {}^s\theta' \wedge \\ \exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1} \end{aligned}$$

This means we are given some  ${}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t.  $(k_1, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''_1} {}^s\theta' \wedge \\ \exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1} \end{aligned}$$

We instantiate (S-M0) with  ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$  we get

$$\begin{aligned} \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

IH:  $[(\tau \sigma)]_V^{\hat{\beta}''} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}''}$  (Statement (1))

Since we have  $({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''}$  therefore from IH we get  $({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\tau' \sigma)]_V^{\hat{\beta}''}$

8. SLIO\*sub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}. ({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means that we are given  $({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

From Definition 3.10 it means we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s \downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, e_t) \downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

From Definition 3.10 we need to prove

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

This further means that given  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $j < n, {}^s v_1$  s.t  $e_s \downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

Instantiating (Sub-E0) with the given  $H_{s1}, H_{t1}$  and  $j < n, {}^s v_1$ . We get

$$\exists H'_t, {}^t v. (H_{t1}, e_t) \downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we have  $({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}}$  therefore from Statement(1) we get  $({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau' \sigma]_V^{\hat{\beta}}$

□

**Theorem 3.20** (SLIO\*  $\rightsquigarrow$  FG: Deriving CG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s1}, H'_{s2}$ .

let  $\text{bool} = (\text{unit} + \text{unit})$ .

$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \text{SLIO } \perp \perp \text{ bool} \wedge$

$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s1}, {}^s v'_1) \wedge$

$(\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2}^f (H'_{s2}, {}^s v'_2)$

$\implies$

$${}^s v'_1 = {}^s v'_2$$

*Proof.* From the CG to FG translation we know that  $\exists e_t$  s.t

$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \text{SLIO } \perp \perp \text{ bool} \rightsquigarrow e_t$

Similarly we also know that  $\exists {}^t v_1, {}^t v_2$  s.t

$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1$  and  $\emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$

From type preservation theorem we know that

$$\emptyset, \emptyset, x : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \vdash_\top e_t : (\text{unit} \xrightarrow{\perp} ((\text{unit} + \text{unit})^\perp + \text{unit})^\perp)^\perp$$

$$\begin{aligned} \emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_1 : ((\mathbf{unit} + \mathbf{unit})^{\perp} + \mathbf{unit})^{\top} \\ \emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_2 : ((\mathbf{unit} + \mathbf{unit})^{\perp} + \mathbf{unit})^{\top} \quad (\text{NI-1}) \end{aligned}$$

Since we have  $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 3.18 we have (we choose  $n$  s.t  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\text{Labeled } \top \text{ bool}]_E^{\emptyset} \quad (\text{NI-2})$$

And therefore from Definition 3.14 and (NI-2) we have

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_1)) \in [x \mapsto \text{Labeled } \top \text{ bool}]_V^{\emptyset}$$

From (NI-0) we know that  $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \text{SLIO } \perp \perp \text{ bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 3.18 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\text{SLIO } \perp \perp \text{ bool}]_E^{\emptyset} \quad (\text{NI-3.1})$$

Applying Definition 3.10 on (NI-3.1) we get

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright} \emptyset \wedge \forall i < n. e_s[{}^s v_1/x] \Downarrow_i {}^s v \implies \\ \exists H'_{t2}. {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n - i, {}^s v, {}^t v) \in [\text{SLIO } \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset \end{aligned}$$

Instantiating with  $\emptyset, \emptyset$ . From SLIO\*-Sem-val we know that  $i = 0$  and  ${}^s v = e_s[{}^s v_1/x]$ .

Therefore we have

$$\exists H'_{t2}. {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n, {}^s v, {}^t v) \in [\text{SLIO } \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset$$

From translation and from (NI-1) we know that  ${}^t v = e_t[{}^t v_1/x] = \lambda_{\cdot}. e_{b1}$  and therefore from fg-val we have  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_1/x], \lambda_{\cdot}. e_{b1}) \in [\text{SLIO } \perp \perp \text{ bool}]_V^{\emptyset}$$

Expanding  $(\emptyset, n, e_s[{}^s v_1/x], \lambda_{\cdot}. e_{b1}) \in [\text{SLIO } \perp \perp \text{ bool}]_V^{\emptyset}$  using Definition 3.9 we get

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'. \\ (k, H_{s3}, H_{t3}) \hat{\triangleright}' ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_1/x]) \Downarrow_i^f (H'_{s1}, {}^s v'') \wedge i < k \implies \\ \exists H''_{t1}. {}^t v''. (H_{t3}, (\lambda_{\cdot}. e_{b1})()) \Downarrow (H''_{t1}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H''_{t1}) \hat{\triangleright}'' {}^s \theta' \wedge \exists {}^t v''' . {}^t v'' = \\ \text{inl } {}^t v''' \wedge ({}^s \theta', k - i, {}^s v'_1, {}^t v''') \in [\text{bool}]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_1, {}^s v'_1, n, \emptyset$  we get

$$\begin{aligned} \exists H''_{t1}. {}^t v''. (\emptyset, (\lambda_{\cdot}. e_{b1})()) \Downarrow (H''_{t1}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s1}, H''_{t1}) \hat{\triangleright}'' {}^s \theta' \wedge \exists {}^t v''' . {}^t v'' = \\ \text{inl } {}^t v''' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''') \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2}) \end{aligned}$$

Since we have  $\exists {}^t v''' . {}^t v'' = \text{inl } {}^t v''' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''') \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$ , therefore from Definition 3.9 we know that 2 cases arise

- ${}^s v'_1 = \text{inl}^s v'_{i1}$  and  ${}^t v''' = \text{inl}^t v'_{i1}$ :

And from Definition 3.9 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means  ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr}^s v'_{i1}$  and  ${}^t v''' = \text{inr}^t v'_{i1}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v'_1 = {}^t v''''_1$  (NI-3.3)

Similarly we can apply Theorem 3.18 with the other substitution to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\text{SLIO} \perp \perp \text{bool}]_E^{\hat{\theta}} \quad (\text{NI-4.1})$$

Applying Definition 3.10 on (NI-4.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\theta}} \emptyset \wedge \forall i < n, {}^s v_s.e_s[{}^s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\text{SLIO} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From SLIO\*-Sem-val we know that  $i = 0$  and  ${}^s v_s = e_s[{}^s v_2/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\text{SLIO} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Also from (NI-1) and from translation we know that  ${}^t v = e_t[{}^t v_2/x] = \lambda_.e_{b2}$  and therefore from fg-val we know that  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_.e_{b2}) \in [\text{SLIO} \perp \perp \text{bool}]_V^{\hat{\theta}}$$

Expanding  $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b2}) \in [\text{SLIO} \perp \perp \text{bool}]_V^{\hat{\theta}}$  using Definition 3.9 we get

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}' \\ (k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s2}, {}^s v''_2) \wedge i < k \implies \\ & \exists H''_{t2}, {}^t v'', (H_{t3}, (\lambda_.e_{b2})()) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v''''_2 . {}^t v''_2 = \\ & \text{inl } {}^t v''''_2 \wedge ({}^s \theta', k - i, {}^s v''_1, {}^t v''''_2) \in [\text{bool}]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$  we get

$$\begin{aligned} & \exists H''_{t2}, {}^t v''. (\emptyset, (\lambda_.e_{b2})()) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v''''_2 . {}^t v''_2 = \\ & \text{inl } {}^t v''''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''''_2) \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-4.2}) \end{aligned}$$

Since we have  $\exists {}^t v''''_2 . {}^t v''_2 = \text{inl } {}^t v''''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''''_2) \in [\text{bool}]_V^{\hat{\beta}''}$ , therefore from Definition 3.9 2 cases arise

- ${}^s v'_2 = \text{inl } {}^s v'_{i2}$  and  ${}^t v''''_2 = \text{inl } {}^t v'_{i2}$ :  
And from Definition 3.9 we know that  
 $({}^s \theta', n - n_1, {}^s v'_{i2}, {}^t v'_{i2}) \in [\text{unit}]_V^{\hat{\beta}''}$   
which means  ${}^s v'_{i2} = {}^t v'_{i2} = ()$
- ${}^s v'_2 = \text{inr } {}^s v'_{i2}$  and  ${}^t v''''_2 = \text{inr } {}^t v'_{i2}$ :  
Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v'_2 = {}^t v''''_2$  (NI-4.3)

From SLIO\* to FG translation we know that  $\exists {}^t v_{i1}. {}^t v_1 = \text{inl } {}^t v_{i1}$  and similarly  $\exists {}^t v_{i2}. {}^t v_2 = \text{inl } {}^t v_{i2}$

From (NI-1) since  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_1 : (\text{bool}^{\perp} + \text{unit})^{\top}$  therefore from SLIO\*-inl we know that  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\perp}$

And from SLIO\*sub-sum we know that  $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\top}$   
Therefore we also have  $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i1} : \text{bool}^{\top}$  (NI-5.1)

Similarly we also have  $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i2} : \text{bool}^{\top}$  (NI-5.2)

Next, let  $e_T = (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top} . \text{case}(e_t(), y.y, z.{}^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) :$   
 $\text{bool}^{\perp}$

where  $\text{true} = \text{inl } ()$  and  $\text{false} = \text{inr } ()$

We claim  $\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} e_T : \text{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{true} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.1:

$$\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} u : \text{bool}^{\top}}$$

P2:

$$\frac{\frac{P2.1 \quad P2.2 \quad P2.3 \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\top} \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : (\text{bool}^{\perp} + \text{unit})^{\top}}}$$

P1.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^{\perp} + \text{unit})^{\perp})^{\perp}}{\text{NI-1}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} () : \text{unit}}{\text{FG-unit}}}{\frac{\frac{}{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\perp} \searrow \perp}}{\text{FG-app}}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t() : (\text{bool}^{\perp} + \text{unit})^{\perp}}$$

P1.1:

$$\frac{\frac{P1.2 \quad \frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, y : \text{bool}^{\perp} \vdash_{\perp} y : \text{bool}^{\perp}}{\text{FG-var}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, z : \text{unit} \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-var}} \quad \frac{}{\emptyset, \emptyset \models \text{bool}^{\perp} \searrow \perp}}{\text{FG-case}}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z.{}^t v_b) : \text{bool}^{\perp}}$$

P1:

$$\frac{\frac{P1.1 \quad \frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z.{}^t v_b) : \text{bool}^{\perp}}{\text{FG-case}}}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top} . \text{case}(e_t(), y.y, z.{}^t v_b)) : ((\text{bool}^{\perp} + \text{unit})^{\top} \xrightarrow{\perp} \text{bool}^{\perp})^{\perp}}$$

Main derivation:

$$\frac{\begin{array}{c} P1 \quad P2 \\ \overline{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \overline{\emptyset, \emptyset \models \text{bool}^\perp \searrow \perp} \end{array}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -.inl \text{true}, -.inl \text{false})) : \text{bool}^\perp} \text{FG-app}$$

Assuming  $e_{b1}()$  reduces in  $n_{t1}$  steps in (NI-3.2) and  $e_{b2}()$  reduces in  $n_{t2}$  steps in (NI-4.2).

We instantiate Theorem 1.29 with  $e_T, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$  and  $\perp$  and therefore from (NI-3.3) and (NI-4.3) we get  ${}^t v_1''' = {}^t v_2'''$  and thus  ${}^s v_1' = {}^s v_2'$

□

## 3.2 Translation from FG to FG<sup>-</sup>

### 3.2.1 FG<sup>-</sup> typesystem

**Lemma 3.21** (FG<sup>-</sup>: Reflexivity of subtyping). *The following hold:*

1. For all  $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all  $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$ .

Proof of statement (1)

Let  $\tau = A^\ell$ . Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on  $A$ .

1.  $A = \mathbf{b}$ :

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2.  $A = \text{ref } \tau$ :

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}}$$

4.  $A = \tau_1 + \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6.  $A = \text{unit}$ :

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$



**Type system:**  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow pc}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau'} \text{FG}^- \text{-var} \quad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^{pc}} \text{FG}^- \text{-lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG}^- \text{-app} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^{pc}} \text{FG}^- \text{-prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG}^- \text{-fst} \quad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^{pc}} \text{FG}^- \text{-inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x. e_1, y. e_2) : \tau} \text{FG}^- \text{-case} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG}^- \text{-sub} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^{pc}} \text{FG}^- \text{-ref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG}^- \text{-deref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^{pc}} \text{FG}^- \text{-assign} \\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^{pc}} \text{FG}^- \text{-unitI} \quad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^{pc}} \text{FG}^- \text{-FI} \\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG}^- \text{-FE} \\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^{pc}} \text{FG}^- \text{-CI} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG}^- \text{-CE}
\end{array}$$

Figure 8: Type system for  $\text{FG}^-$

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A'^{\ell'}} \text{FG}^- \text{-sub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FG}^- \text{-sub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FG}^- \text{-sub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FG}^- \text{-sub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FG}^- \text{-sub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FG}^- \text{-sub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FG}^- \text{-sub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{FG}^- \text{-sub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{FG}^- \text{-sub-constraint}
\end{array}$$

Figure 9:  $\text{FG}^-$  subtyping

7.  $A = \forall \alpha. \tau_i$ :

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}$$

8.  $A = c \Rightarrow \tau_i$ :

$$\frac{\frac{}{\Sigma; \Psi \vdash c \implies c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}$$

□

### 3.2.2 Type translation

We define a translation of types, indexed by a label  $\ell$  (which represents a  $pc$  joined with all outer labels) below. This is written  $\llbracket \tau \rrbracket_\ell$ .

**Definition 3.22** ( $FG \rightsquigarrow FG^-$ : Type translation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_\ell &= \mathbf{b} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_\ell &= \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha) \\
\llbracket \tau_1 \times \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell \times \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \tau_1 + \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \text{ref } \tau \rrbracket_\ell &= \text{ref } \llbracket \tau \rrbracket_\perp \\
\llbracket \text{unit} \rrbracket_\ell &= \text{unit} \\
\llbracket \forall \gamma. (\ell_e, \tau) \rrbracket_\ell &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_\ell &= \forall \alpha. \alpha, (((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket A' \rrbracket_\ell &= (\llbracket A \rrbracket_{\ell \sqcup \ell'})^{\ell \sqcup \ell'}
\end{aligned}$$

Translation judgement:

$$\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau \rrbracket_{pc'}} \text{ where}$$

$pc' \sqsubseteq pc$  and  $\forall i \in 1 \dots n. \ell_i \sqsubseteq pc'$   
 $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$   
 $\Gamma' = x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$

### 3.2.3 Type preservation: $FG$ to $FG^-$

**Theorem 3.23** ( $FG \rightsquigarrow FG^-$ : Type preservation). *Suppose (1)  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$  and (2)  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$  in  $FG$ . Suppose  $\ell_1, \dots, \ell_n$  and  $pc'$  are arbitrary labels with free variables in  $\Sigma$  such that (3)  $\Sigma; \Psi \vdash pc' \sqsubseteq pc$  and (4) For each  $i \in [1, n]$ ,  $\Sigma; \Psi \vdash \ell_i \sqsubseteq pc'$ .*

*Let  $\Gamma'$  be the  $FG^-$  context  $x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$ . Then,  $\Sigma; \Psi; \Gamma' \vdash_{pc'} e : \llbracket \tau \rrbracket_{pc'}$  in  $FG^-$ .*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. var:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} x : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}} \text{var}$$

$$\frac{\llbracket \tau \rrbracket_{\ell_n} <: \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}}$$

2. lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma, x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{\ell'_e} e_m : \llbracket \tau_2 \rrbracket_{\ell'_e} \quad \ell_{n+1} \sqsubseteq \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : T_1} \text{lam}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.2} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Given}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Weakening}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma', x : \llbracket \tau_1 \rrbracket_{\beta} \vdash_{\alpha} e_m : \llbracket \tau_2 \rrbracket_{\alpha}} \text{IH}}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{P1}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma' \vdash_{\alpha} \lambda x. e_m : T_{1.3}} \text{FG}^{-}\text{-lam}}{\Sigma, \alpha, \beta; \Psi; \Gamma' \vdash_{\alpha} \nu(\lambda x. e_m) : T_{1.2}} \text{FG}^{-}\text{-CI}}{\Sigma, \alpha; \Psi; \Gamma' \vdash_{\alpha} \Lambda(\nu(\lambda x. e_m)) : T_{1.1}} \text{FG}^{-}\text{-FI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\Lambda(\nu(\lambda x. e_m))) : T_1} \text{FG}^{-}\text{-FI}$$

3. app:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^{\ell} \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau_2 \rrbracket_{pc'}} \text{app}}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_{\alpha})^{\alpha})^{\alpha})^{pc' \sqcup \ell}$$

$$T_{1.1} = (\forall \beta. (pc' \sqcup \ell), (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge \beta \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell))$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.4} = (\llbracket \tau_1 \rrbracket_{(pc')} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

P7:

$$\overline{pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P6:

$$\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}} \text{IH2}$$

P5:

$$\overline{\Sigma; \Psi \vdash T_{1.3} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\overline{\Sigma; \Psi \vdash T_{1.2} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\overline{\Sigma; \Psi \vdash T_{1.1} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P2:

$$\overline{pc' \sqcup pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P1:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1} \text{ IH1} \quad P2 \quad P3}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square : T_{1.1}} \text{ FG}^- \text{-FE} \quad P2 \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square \square : T_{1.2}} \text{ FG}^- \text{-FE}}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad \overline{\Sigma; \Psi \vdash c_1} \quad P2 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.3}} \text{ FG}^- \text{-CE} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.4}} \text{ FG}^- \text{-sub} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell}} \text{ FG}^- \text{-app}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 3.26}$$

4. prod:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ prod}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'}} \text{ IH1} \quad \overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ FG}^- \text{-prod}}$$

5. fst:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ fst}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-fst}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ Lemma 3.26}}$$

6. snd:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ snd}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-snd}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 3.26}}$$

7. inl:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}{\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}} \text{IH} \quad \text{FG}^- \text{-inl}$$

8. inr:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}{\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inr}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}}} \text{IH} \quad \text{FG}^- \text{-inr}$$

9. case:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{\ell_{n+2}} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, y.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{case}$$

P2:

$$\frac{\Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}}{\text{IH3 @ } pc' \sqcup \ell}$$

P1:

$$\frac{\Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}}{\text{IH2 @ } pc' \sqcup \ell}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, y.e_{m2}) : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH1} \quad \text{P1} \quad \text{P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, y.e_{m2}) : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{FG}^- \text{-case}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, y.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{Lemma 3.26}$$

10. sub:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc''} e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'} \quad \Sigma; \Psi \vdash_{pc} pc \sqsubseteq pc'' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{sub}$$

$$\frac{\frac{\frac{pc' \sqsubseteq pc \sqsubseteq pc''}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'}} \text{IH} \quad \frac{\tau' <: \tau}{\llbracket \tau' \rrbracket_{pc'} <: \llbracket \tau \rrbracket_{pc'}} \text{Lemma 3.24}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{IH}$$

11. ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{ref}$$

P1:

$$\frac{\frac{\text{Given} \quad \Sigma; \Psi \vdash pc' \sqsubseteq pc}{\Sigma; \Psi \vdash \tau \searrow pc}}{\Sigma; \Psi \vdash \tau \searrow pc'} \text{Lemma 3.29}}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp \searrow pc'}$$

Main derivation:

$$\frac{\frac{\frac{\text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_\perp} \text{Lemma 3.26} \quad P1}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{FG}^- \text{-new}$$

12. deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{deref}$$

$$\frac{\frac{\frac{\tau <: \tau'}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp <: \llbracket \tau' \rrbracket_{pc' \sqcup \ell}} \text{Lemma 3.24} \quad \frac{\text{IH1}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc'} \text{Definition of } \searrow}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc' \sqcup \ell}}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{Lemma 3.26}$$

13. assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{assign}$$

P1:

$$\frac{\frac{\text{IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp} \text{Lemma 3.26}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp} \text{Given} \quad \frac{\tau \searrow pc}{\tau \searrow pc'}}$$

Main derivation:

$$\frac{\frac{\text{IH1} \quad P1 \quad \frac{\tau \searrow (\ell \sqcup pc)}{\llbracket \tau \rrbracket_\perp \searrow \ell \sqcup pc'} \text{Lemma 3.29}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{FG}^- \text{-assign}}$$

14. unitI:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} () : \text{unit}^{pc'}} \text{unitI}$$

$$\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} () : \text{unit}^{pc'}} \text{FG}^- \text{-unitI}$$

15. FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma, \alpha; \Psi; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.2} = (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_\alpha$$

P1:

$$\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_\alpha e_m : T_{1.3}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_\alpha \Lambda e_m : T_{1.2}} \text{FG}^- \text{-FI}$$

Main derivation:

$$\frac{\frac{P1}{\Sigma, \alpha; \Psi; \Gamma' \vdash_\alpha \nu(\Lambda e_m) : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FG}^- \text{-FI}$$

16. CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma; \Psi, c; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \overset{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{CI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \overset{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi, c_1; \Gamma' \vdash_\alpha e_m : T_{1.2}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha; \Psi; \Gamma' \vdash_\alpha \nu e_m : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{FG}^- \text{-FI}$$



17. FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \gamma. (\ell_e, \tau))^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \text{FV}(\ell') \in \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\gamma] \quad \Sigma; \Psi \vdash \tau[\ell'/\gamma] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau[\ell'/\gamma] \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}} \text{FE}$$

$$T_1 = (\forall \alpha. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha) \alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}[\ell'/\gamma]$$

$$T_{1.31} = \llbracket \tau[\ell'/\gamma] \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.4} = \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}$$

P5:

$$\frac{}{T_{1.2} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\frac{}{T_{1.1} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\frac{}{(pc' \sqcup \ell) \sqsubseteq (pc \sqcup \ell) \sqsubseteq \ell_e} \text{Given}$$

P2:

$$\frac{\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{IH} \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.1}} \text{FG}^- \text{-FE}$$

P1:

$$\frac{P2 \quad P3 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.2}} \text{FG}^- \text{-CE}$$

P0:

$$\frac{P1 \quad \frac{\frac{}{\Sigma; \Psi \vdash T_{1.3} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket \quad P2}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.3}} \text{FG}^- \text{-FE}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.31}} \text{Lemma 3.28}$$

Main derivation:

$$\frac{P0 \quad \frac{}{\Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell} \text{Lemma 3.26}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.4}} \text{Lemma 3.26}$$

18. CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\rightsquigarrow} \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : \llbracket \tau \rrbracket_{pc'}} \text{CE}$$

$$T_1 = (\forall \alpha. \alpha, ((c \wedge (pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = ((c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} \llbracket \tau \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.3} = \llbracket \tau \rrbracket_{pc'}$$

$$c_1 = (c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

P3:

$$\frac{\overline{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_e} \text{ Given}}{\Sigma; \Psi \vdash (pc' \sqcup \ell) \sqsubseteq \ell_e}$$

P2:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)}$$

P1:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)}$$

P0:

$$\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{ IH} \quad P1}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.1}} \text{ FG}^- \text{-FE} \quad \frac{\overline{\Sigma; \Psi \vdash c} \text{ Given, Weakening} \quad P3}{\Sigma; \Psi \vdash c_1} \text{ P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.2}} \text{ FG}^- \text{-CE}$$

Main derivation:

$$\frac{P0.1 \quad \frac{\overline{\tau \searrow \ell} \text{ Given}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}} \text{ Lemma 3.26}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}}$$

□

**Lemma 3.24** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping).  $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$  and the following holds:

1.  $\forall \tau, \tau'.$

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_\ell <: \llbracket \tau' \rrbracket_{\ell'}$$

2.  $\forall A, A'.$

$$\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \llbracket A \rrbracket_\ell <: \llbracket A' \rrbracket_{\ell'}$$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash A_1 <: A_2} \text{ By inversion } P1}{\Sigma; \Psi \vdash (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2})} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion } \quad \overline{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{ Given}}{\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1})^{\ell \sqcup \ell_1} <: (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2})^{\ell' \sqcup \ell_2}}{\Sigma; \Psi \vdash \llbracket A_1^{\ell_1} \rrbracket_{\ell} <: \llbracket A_2^{\ell_2} \rrbracket_{\ell'}}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{ FG}^- \text{ sub-base}}{\Sigma; \Psi \vdash \llbracket \mathbf{b} \rrbracket_{\ell} <: \llbracket \mathbf{b} \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

2. FGsub-ref:

$$\frac{\overline{\Sigma; \Psi \vdash \text{ref } \llbracket \tau_i \rrbracket_{\perp} <: \text{ref } \llbracket \tau_i \rrbracket_{\perp}} \text{ FG}^- \text{ sub-ref}}{\Sigma; \Psi \vdash \llbracket \text{ref } \tau_i \rrbracket_{\ell} <: \llbracket \text{ref } \tau_i \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

3. FGsub-prod:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell} \times \llbracket \tau'_2 \rrbracket_{\ell'}}{\Sigma; \Psi \vdash \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \times \tau'_2 \rrbracket_{\ell'}} \text{ FG}^- \text{ sub-prod}} \text{ Definition of } \llbracket \cdot \rrbracket$$

4. FGsub-sum:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'} + \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ FG}^- \text{ sub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

5. FGsub-arrow:

$$T_1 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{1.0} = \forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.1} = ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.2} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$

$$T_2 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{2.0} = \forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{2.1} = ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{2.2} = (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha$$

$$c_2 = (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha)$$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\alpha <: \llbracket \tau'_2 \rrbracket_\alpha} \text{ IH(1) with } \ell = \ell' = \alpha$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau'_1 <: \tau_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{ IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.3} <: T_{2.3}} \text{ FG}^- \text{ sub-arrow}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell'_e \sqsubseteq \ell_e}}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e}}{\Sigma, \alpha, \beta; \Psi \vdash c_2 \implies c_1} \quad \text{P1}}{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.2} <: T_{2.2}}{\text{Weakening, FG}^- \text{-sub-label}}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}}} \text{FG}^- \text{-sub-constraint}$$

P0.1:

$$\frac{\text{P0}}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \text{FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{\frac{\frac{\text{P0.1}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{FG}^- \text{-sub-label}}{\Sigma; \Psi \vdash \left[ \left[ \tau_1 \xrightarrow{\ell_\xi} \tau_2 \right]_\ell <: \left[ \left[ \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2 \right]_{\ell'} \right]_{\ell'}} \text{Definition of } \llbracket \cdot \rrbracket$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FG}^- \text{-sub-unit}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_\ell <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{Definition of } \llbracket \cdot \rrbracket$$

7. FGsub-forall:

$$\begin{aligned}
T_1 &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
T_{1.0} &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
T_{1.1} &= (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
c_1 &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \\
T_{1.2} &= \llbracket \tau \rrbracket_\alpha \\
T_2 &= \forall \alpha. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha)^\alpha \\
T_{2.0} &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
T_{2.1} &= (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
c_2 &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \\
T_{2.2} &= \llbracket \tau' \rrbracket_\alpha
\end{aligned}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e}}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

P1:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}}{\text{IH}} \quad \text{FG}^- \text{-sub-forall} \quad \frac{}{\Sigma; \Psi \vdash c_2 \implies c_1} \text{P0}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \quad \text{FG}^- \text{-sub-constraint}}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \quad \text{FG}^- \text{-sub-forall}}{\Sigma; \Psi \vdash T_1 <: T_2} \quad \text{FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{}{\Sigma; \Psi \vdash \llbracket \forall \gamma. \tau_1 \rrbracket_\ell <: \llbracket \forall \gamma. \tau_2 \rrbracket_{\ell'}} \text{P0.1} \quad \text{Definition of } \llbracket \cdot \rrbracket$$

8. FGsub-constraint:

$$T_1 = \forall \alpha. \alpha, (((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.1} = ((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_\alpha$$

$$c_1 = (c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_2 = \forall \alpha. \alpha, (((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} \llbracket \tau' \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{2.1} = ((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} \llbracket \tau' \rrbracket_\alpha)^\alpha$$

$$T_{2.2} = \llbracket \tau' \rrbracket_\alpha$$

$$c_2 = (c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e)$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau_2$$

P1:

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash c' \implies c} \text{By inversion}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e} \text{P1}}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket_\alpha <: \llbracket \tau' \rrbracket_\alpha} \text{IH}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \text{FG}^- \text{-sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{Definition of } \llbracket \cdot \rrbracket_\ell} \text{P0}$$

□

**Lemma 3.25** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping with label). *If  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell'$ , then  $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau \rrbracket_{\ell'}$  in FG<sup>-</sup>.*

*Proof.* From Lemma 3.24 with  $\tau = \tau'$  and from Lemma 3.21 □

**Lemma 3.26** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Subtyping for  $\tau \searrow \ell$ ). *If  $\Sigma; \Psi \vdash \tau \searrow \ell$ , then  $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'}$  in FG<sup>-</sup>.*

*Proof.* Since  $\Sigma; \Psi \vdash \tau \searrow \ell$ , there exists  $\ell''$  such that  $\tau = \mathbf{A}^{\ell''}$  and  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$ . Now we have:

$$\begin{aligned} & \Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'} \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell \sqcup \ell'} <: \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell'} && (\tau = \mathbf{A}^{\ell''}) \\ = & \Sigma; \Psi \vdash (\llbracket \mathbf{A} \rrbracket_{\ell \sqcup \ell' \sqcup \ell''})^{\ell \sqcup \ell' \sqcup \ell''} <: (\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell''})^{\ell' \sqcup \ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell \sqcup \ell''} <: \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \end{aligned}$$

The last statement holds by Lemma 3.25, since  $\Sigma; \Psi \vdash \ell \sqcup \ell'' \sqsubseteq \ell''$  follows from our earlier assertion that  $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$ . □

**Lemma 3.27** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Lemma for protection relation).  $\forall \Sigma, \Psi, \alpha, \tau, \ell, \ell'$ .

$$\Sigma, \alpha; \Psi \vdash \tau \searrow \ell \implies \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell[\ell'/\alpha], \text{ where } FV(\ell') \in \Sigma$$

*Proof.* Say  $\tau = \mathbf{A}^{\ell_g}$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell_g} \text{By inversion on } \Sigma, \alpha; \Psi \vdash \tau \searrow \ell}}{\Sigma; \Psi \vdash \ell[\ell'/\alpha] \sqsubseteq \ell_g[\ell'/\alpha]} \text{Substitution over constraints}}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_g}[\ell'/\alpha] \searrow \ell[\ell'/\alpha]} \text{Definition of } \searrow$$

□

**Lemma 3.28** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Substitution lemma). *For all  $\ell, \ell'$  the following hold:*

1.  $\forall \tau. \llbracket \tau \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \tau[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$
2.  $\forall \mathbf{A}. \llbracket \mathbf{A} \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $\mathbf{A}$

Proof of statement (1)

$$\begin{aligned} & \text{Let } \tau = \mathbf{A}^{\ell_i} \\ & \llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell}[\ell'/\alpha] \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell})^{\ell_i \sqcup \ell}[\ell'/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell}[\ell'/\alpha])^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} \\ = & (\llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]})^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} && \text{IH(2) on } \mathbf{A} \\ = & \llbracket (\mathbf{A}[\ell'/\alpha])^{\ell_i[\ell'/\alpha]} \rrbracket_{\ell[\ell'/\alpha]} \\ = & \llbracket \mathbf{A}^{\ell_i}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \end{aligned}$$

Proof of statement (2)

We consider cases of  $\mathbf{A}$

1.  $\mathbf{A} = \mathbf{b}$ :

$$\begin{aligned} & \llbracket \mathbf{b} \rrbracket_{\ell}[\ell'/\alpha] \\ = & \mathbf{b}[\ell'/\alpha] && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \mathbf{b} && \alpha \notin FV(\mathbf{b}) \\ = & \llbracket \mathbf{b} \rrbracket_{\ell} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \llbracket \mathbf{b}[\ell'/\alpha] \rrbracket_{\ell} \end{aligned}$$

2.  $A = \text{ref } \tau_i$ :

$$\begin{aligned}
& \llbracket \text{ref } \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \text{ref } \llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \text{ref } (\llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha]) \\
&= \text{ref } \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\perp} && \text{IH(1) on } \tau_i \\
&= \llbracket \text{ref } \tau_i[\ell'/\alpha] \rrbracket_{\ell}
\end{aligned}$$

3.  $A = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] \times \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \times \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] \times \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 \times \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

4.  $A = \tau_1 + \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 + \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} + \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] + \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} + \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] + \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 + \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

5.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\begin{aligned}
& \llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell \sqsubseteq \beta_1 \sqsubseteq \ell_e \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta}[\ell'/\alpha] \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1}[\ell'/\alpha])^{\beta_1})^{\beta_1})^{\beta_1} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1[\ell'/\alpha] \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1} \\
& \quad \text{(IH1 on } \tau_1 \text{ and } \tau_2 \text{)} \\
&= \llbracket (\tau_1[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} \\
&= \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

6.  $A = \forall \gamma. \tau_i$ :

$$\begin{aligned}
& \llbracket \forall \gamma. \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta. \beta, ((\ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \quad \text{IH1 on } \tau_i \\
&= \llbracket \forall \beta. \ell_e[\ell'/\alpha], \tau_i[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

7.  $A = c \Rightarrow \tau_i$ :



$$\begin{aligned}
& \llbracket c \Rightarrow \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
= & \forall \beta. \beta, (((c \wedge \ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \text{IH1 on } \tau_i \\
= & \left[ \left[ (c[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_i[\ell'/\alpha]) \right] \right]_{\ell[\ell'/\alpha]} \\
= & \left[ (c \xrightarrow{\ell_e} \tau_i)[\ell'/\alpha] \right]_{\ell[\ell'/\alpha]}
\end{aligned}$$

□

**Lemma 3.29** (FG  $\rightsquigarrow$  FG<sup>-</sup>: Preservation of protection relation).  $\forall \tau, \ell, \ell'$ .

$$\tau \searrow \ell \implies \llbracket \tau \rrbracket_{\ell'} \searrow \ell$$

*Proof.* Let  $\tau = \mathbf{A}^{\ell_i}$

$$\begin{array}{c}
\frac{}{\tau \searrow \ell} \text{ Given} \\
\frac{}{\ell \sqsubseteq \ell_i} \text{ Given} \\
\frac{}{\ell \sqsubseteq (\ell' \sqcup \ell_i)} \text{ By inversion} \\
\frac{}{(\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell_i})^{\ell' \sqcup \ell_i} \searrow \ell} \text{ Definition of } \llbracket \cdot \rrbracket \\
\frac{}{\llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell'} \searrow \ell} \\
\hline
\llbracket \tau \rrbracket_{\ell'} \searrow \ell
\end{array}$$

□

### 3.3 Translation from FG to SLIO\*

#### 3.3.1 Type directed (direct) translation from FG to SLIO\*

**Definition 3.30** (FG  $\rightsquigarrow$  SLIO\*: Type translation).

$$\begin{aligned}
\langle \mathbf{b} \rangle_\ell &= \mathbf{b} \\
\langle \mathbf{unit} \rangle_\ell &= \mathbf{unit} \\
\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle_\ell &= \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \text{SLIO } \gamma \gamma \langle \tau_2 \rangle_\alpha \\
\langle \tau_1 \times \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell \\
\langle \tau_1 + \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell \\
\langle \text{ref } \mathbf{A}^{\ell'} \rangle_\ell &= \text{ref } \ell' \langle \mathbf{A} \rangle_{\ell'} \\
\langle \forall \alpha. (\ell_e, \tau) \rangle_\ell &= \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma \langle \tau \rangle_{\alpha'} \\
\langle c \xrightarrow{\ell_e} \tau \rangle_\ell &= \forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma \langle \tau \rangle_\alpha \\
\langle \mathbf{A}^{\ell'} \rangle_\ell &= \text{Labeled } (\ell \sqcup \ell') \langle \mathbf{A} \rangle_{\ell \sqcup \ell'}
\end{aligned}$$

For  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$  and  $\bar{\ell} = \ell_1, \dots, \ell_n$ , define  $\langle \Gamma \rangle_{\bar{\ell}} = x_1 : \langle \tau_1 \rangle_{\ell_1}, \dots, x_n : \langle \tau_n \rangle_{\ell_n}$ .

We use a coercion function defined as follows:

$$\begin{aligned}
\text{coerce\_taint} &: \text{SLIO } \gamma \alpha_c \tau' \rightarrow \text{SLIO } \gamma \gamma \tau' \quad \text{when } \tau' = \text{Labeled } \alpha'_c \tau \text{ and } \Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \\
\text{coerce\_taint} &\triangleq \lambda x. \text{toLabeled}(\text{bind}(x, y. \text{unlabel}(y)))
\end{aligned}$$

$$\begin{aligned}
&\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var} \\
&\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(\lambda x. e_{c1}))))} \text{FC-lam} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a. \text{bind}(e_{c2}, b. \text{bind}(\text{unlabel } a, c. (c \square \square \square \bullet) b))))} \text{FC-app} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a. \text{bind}(e_{c2}, b. \text{ret}(\text{Lb}(a, b))))} \text{FC-prod} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } (a), b. \text{ret}(\text{fst}(b))))} \text{FC-fst} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } (a), b. \text{ret}(\text{snd}(b))))} \text{FC-snd} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a. \text{ret}(\text{Lbinl}(a)))} \text{FC-inl} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a. \text{ret}(\text{Lbinr}(a)))} \text{FC-inr}
\end{aligned}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)))} \text{FC-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()))} \text{FC-assign} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda (\nu(e_c))))} \text{FC-FI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda (\nu(e_c))))} \text{FC-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-CE}
\end{array}$$

### 3.3.2 Type preservation for FG to SLIO\* translation

**Lemma 3.31** (Coercion lemma - typing).  $\forall \Sigma, \Psi, \Gamma, \alpha_c, \alpha'_c, \tau.$

$$\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \implies$$

$$\Sigma; \Psi; \Gamma \vdash \text{coerce\_taint} : \text{SLIO } \gamma \alpha_c \text{ Labeled } \alpha'_c \tau \rightarrow \text{SLIO } \gamma \gamma \text{ Labeled } \alpha'_c \tau$$

*Proof.*  $T_{c4} = \text{Labeled } \alpha'_c \tau$

$$T_{c3} = \text{SLIO } \alpha_c \alpha'_c \tau$$

$$T_{c2} = \text{SLIO } \gamma \alpha'_c \tau$$

$$T_{c1} = \text{SLIO } \gamma \gamma \text{ Labeled } \alpha'_c \tau$$

$$T_{c0} = \text{SLIO } \gamma \alpha_c \text{ Labeled } \alpha'_c \tau$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}}{\text{SLIO}^*\text{-var}} \quad \frac{\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c}{\text{Given}}}{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{SLIO}^*\text{-unlabel}$$

Pc1:

$$\frac{}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash x : T_{c0}} \text{SLIO}^* \text{-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{SLIO}^* \text{-bind}}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{SLIO}^* \text{-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\Sigma; \Psi; \Gamma \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{SLIO}^* \text{-lam}}{\Sigma; \Psi; \Gamma \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

□

**Theorem 3.32** (FG  $\rightsquigarrow$  SLIO\*: Type preservation). *Suppose  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$  in FG. Then, there exists  $e'$  such that  $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$  and for any  $\alpha', \beta', \gamma'$  with  $\beta' \sqcup \gamma' \sqsubseteq pc \sqcap \alpha'$ , there is a derivation of  $\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e' : \text{SLIO } \gamma' \gamma' (\tau)_{\alpha'}$  in SLIO\*.*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{\frac{\frac{}{(\Gamma)_{\beta''}(x) = (\tau)_{\beta''}} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\beta''} \vdash x : (\tau)_{\beta''}} \text{SLIO}^* \text{-var} \quad \frac{\frac{}{\Sigma; \Psi \vdash \beta'_o \sqcup \gamma'_o \sqsubseteq \alpha'_o \sqcap pc} \text{Given}}{\Sigma; \Psi \vdash \beta'_o \sqsubseteq \alpha'_o} \text{Lemma 3.33, SLIO}^* \text{-sub}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash x : (\tau)_{\alpha'_o}} \text{Lemma 3.33, SLIO}^* \text{-sub}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{ret } x : \text{SLIO } \gamma'_o \gamma'_o (\tau)_{\alpha'_o}} \text{FC-var}$$

2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1}))))} \text{FC-lam}$$

$$T_0 = \text{SLIO } \gamma'_j \gamma'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp)_{\alpha'_j} = \text{SLIO } \gamma'_j \gamma'_j \text{ Labeled } \alpha'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2)_{\alpha'_j})$$

$$T_1 = \text{SLIO } \gamma'_j \gamma'_j \text{ Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

$$T_{1.0} = \text{Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

$$T_{1.1} = \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

$$T_{1.2} = (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

$$T_{1.3} = (\tau_1)_{\beta_t} \rightarrow \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

$$T_{1.4} = \text{SLIO } \gamma_t \gamma_t (\tau_2)_{\alpha_t}$$

P3:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \gamma_j \sqsubseteq \alpha'_j \sqcap pc \quad \text{Given, Weakening}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqsubseteq \alpha'_j}$$

P2:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e \quad P3}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e}$$

P1:

$$\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}}, x : (\tau_1)_{\beta_t} \vdash e_{c1} : T_{1.4} \quad \text{IH}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x. e_{c1} : T_{1.3}} \quad P2}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x. e_{c1} : T_{1.3}} \quad \text{SLIO*}-\text{lam}$$

P0:

$$\frac{\Sigma; \Psi \vdash \overline{\beta'_j} \sqcup \gamma'_j \sqsubseteq \alpha'_j \quad \text{Given}}{\Sigma; \Psi \vdash \gamma_j \sqsubseteq \alpha_j}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \nu(\lambda x. e_{c1}) : T_{1.2} \quad P1}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1})) : T_{1.1}} \quad \text{SLIO*}-\text{CI}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1}))) : T_{1.0}} \quad \text{3 applications SLIO*}-\text{FI} \quad P0}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_{c1})))) : T_1} \quad \text{SLIO*}-\text{label}} \quad \text{SLIO*}-\text{ret}$$

3. FC-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c \square \square \square \bullet) b))))} \quad \text{FC-app}$$

$$\beta' = \bigcup_{\beta_i \in \overline{\beta'}} \beta_i$$

$$T_0 = \text{SLIO } \gamma' \gamma' ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell)_{\beta' \sqcup \gamma'} = \text{SLIO } \gamma' \gamma' \text{ Labeled } (\beta' \sqcup \gamma' \sqcup \ell) ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell)_{\beta' \sqcup \gamma' \sqcup \ell}$$

$$T_1 = \text{SLIO } \gamma' \gamma' \text{ Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.1} = \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.2} = \text{SLIO } \gamma' (\gamma' \sqcup (\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.3} = \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.4} = \forall \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.5} = \forall \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.6} = (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow T_{1.7}$$

$$T_{1.7} = (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.8} = \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.9} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.10} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i})_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.11} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{ Labeled } (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.12} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.13} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } (\ell_i \sqcup \beta' \sqcup \gamma') (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma')}$$

$$T_2 = \text{SLIO } (\gamma') (\gamma') (\tau_2)_{(\beta' \sqcup \gamma')}$$

$$T_3 = \text{SLIO } (\gamma') (\gamma') (\tau_1)_{(\beta' \sqcup \gamma')}$$

P8:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash b : (\tau_1)_{(\beta' \sqcup \gamma')}} \text{SLIO}^* \text{-var}$$

P7:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \alpha' \sqcap pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha' \sqcap pc \sqsubseteq \ell_e}$$

P6:

$$P7 \quad \frac{\frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash (\ell \sqcup \beta' \sqcup \gamma') \sqsubseteq \ell_e}$$

P5:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c : T_{1.3}} \text{SLIO}^* \text{-var}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c \square : T_{1.4}} \text{SLIO}^* \text{-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c \square \square : T_{1.5}} \text{SLIO}^* \text{-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c \square \square \square : T_{1.6}} \text{SLIO}^* \text{-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c \square \square \square \bullet : T_{1.7}} \text{SLIO}^* \text{-CE}$$

P4:

$$\frac{P5 \quad P8}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash (c \square \square \square \bullet) b : T_{1.8}} \text{SLIO}^* \text{-app}$$

P3:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')} \vdash a : T_{1.1}} \text{SLIO}^* \text{-var}$$

P2:

$$\frac{\frac{P3}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')} \vdash \text{unlabel } a : T_{1.2}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')} \vdash \text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b) : T_{1.9}} \text{SLIO}^*\text{-unlabel } P4}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1}, b : (\tau_1)_{(\beta' \sqcup \gamma')} \vdash \text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b) : T_{1.9}} \text{SLIO}^*\text{-bind}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash e_{c2} : T_3}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.9}} \text{IH2, Weakening } P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.9}} \text{SLIO}^*\text{-bind}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_{c1} : T_1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.9}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.10}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.11}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.12}} \text{IH1 with } (\beta' \sqcup \gamma'), \bar{\beta}', \gamma' \quad P1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_{1.13}} \text{Lemma 3.31}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b))) : T_2}} \text{SLIO}^*\text{-bind}$$

4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((\tau_1 \times \tau_2)^\perp)_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\tau_1 \times \tau_2))_{\alpha'}$$

$$T_3 = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_{3.1} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_4 = \text{SLIO } \gamma' \gamma' (\tau_1)_{\alpha'}$$

$$T_5 = \text{SLIO } \gamma' \gamma' (\tau_2)_{\alpha'}$$

P4:

$$\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{SLIO}^*\text{-var}$$

P3:

$$\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash b : (\tau_2)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash b : (\tau_2)_{\alpha'}} \text{SLIO}^*\text{-var}$$

P2:

$$\frac{\frac{\frac{P3 \quad P4}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash (a, b) : (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{Lb}(a, b) : T_{3.1}} \text{SLIO}^*\text{-prod}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{SLIO}^*\text{-label}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{SLIO}^*\text{-ret}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash e_{c2} : T_5 \quad \text{IH2} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{SLIO}^*\text{-bind}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_{c1} : T_4 \quad \text{IH1} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1} \text{Definition 3.30}$$

5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{FC-fst}$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau_1)_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\tau_1 \times \tau_2)^\ell)_{\alpha'}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup \alpha' ((\tau_1 \times \tau_2)_{\alpha' \sqcup \ell})$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.4} = (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.5} = \text{SLIO } (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_3 = \text{SLIO } (\gamma' \sqcup \alpha' \sqcup \ell) (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.1} = \text{SLIO } (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.2} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.3} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) (\mathbf{A}^{\ell_i})_{\alpha' \sqcup \ell}$$

$$T_{3.4} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) \text{ Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.5} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.6} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell_i}$$

$$T_{3.7} = \text{SLIO } (\gamma') (\gamma') (\mathbf{A}^{\ell_i})_{\alpha'}$$

P2:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4} \quad \text{SLIO}^*\text{-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : (\tau_1)_{\alpha' \sqcup \ell} \quad \text{SLIO}^*\text{-fst}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{SLIO}^*\text{-ret}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{SLIO}^*\text{-ret}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5} \quad \text{SLIO}^*\text{-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1} \quad \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1} \quad \text{SLIO}^*\text{-bind}} \text{SLIO}^*\text{-bind}$$



P0:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e_c : T_{2.2}} \text{IH} \quad P1 \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1} \quad \text{SLIO}^*\text{-bind} \\
\frac{}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.4} \quad \text{Definition 3.30} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.5} \quad \text{Lemma 3.31}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P0 \quad \frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{By inversion}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.6} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.7} \quad \text{Definition 3.30} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_1
\end{array}$$

6. FC-snd:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd}$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau_2)_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\tau_1 \times \tau_2)^\ell)_{\alpha'}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup \alpha' ((\tau_1 \times \tau_2)^\ell)_{\alpha' \sqcup \ell}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.4} = (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.5} = \text{SLIO } (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_3 = \text{SLIO } (\gamma' \sqcup \alpha' \sqcup \ell) (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{3.1} = \text{SLIO } (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{3.2} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{3.3} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) (\mathbf{A}^{\ell_i})_{\alpha' \sqcup \ell}$$

$$T_{3.4} = \text{SLIO } (\gamma') (\alpha' \sqcup \ell) \text{ Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.5} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.6} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell_i}$$

$$T_{3.7} = \text{SLIO } (\gamma') (\gamma') (\mathbf{A}^{\ell_i})_{\alpha'}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}}{\text{SLIO}^*\text{-var}}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : (\tau_2)_{\alpha' \sqcup \ell}} \text{SLIO}^*\text{-snd}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{SLIO}^*\text{-ret}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{unlabel}(a) : T_{2.5}} \text{SLIO}^*\text{-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))) : T_{3.1}} \text{SLIO}^*\text{-bind}$$

P0:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e_c : T_{2.2}} \text{IH} \quad P1}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.1}} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.2}}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.4}} \text{Lemma 3.31}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.5}} \text{Lemma 3.31}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad \frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{By inversion}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.6}} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.7}} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_1}$$

7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_{1.3} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' (\tau_1)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{inl}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{SLIO}^*\text{-inl}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{Lbinl}(a) : T_{1.3}} \text{SLIO}^*\text{-label}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{SLIO}^*\text{-ret}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_{1.2}} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_1} \text{Definition 3.30}$$

8. FC-inr:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\tau_1)_{\alpha'} + (\tau_2)_{\alpha'})$$

$$T_{1.3} = \text{Labeled } \alpha' ((\tau_1)_{\alpha'} + (\tau_2)_{\alpha'})$$

$$T_2 = \text{SLIO } \gamma' \gamma' (\tau_2)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{inr}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{SLIO}^*\text{-inr}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{Lbinr}(a) : T_{1.3}} \text{SLIO}^*\text{-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{ret}(\text{Lbinr}(a)) : T_{1.2}} \text{SLIO}^*\text{-ret}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{SLIO}^*\text{-var}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_{1.2}} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_1} \text{Definition 3.30}$$

9. FC-case:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce.taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case}}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau)_{(\alpha')}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\tau_1 + \tau_2)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1 + \tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$T_{2.3} = \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$T_{2.4} = \text{SLIO } \gamma' (\gamma' \sqcup (\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$T_{2.5} = ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$T_3 = \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_4 = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_5 = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i})_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{5.1} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{ Labeled } \ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{5.2} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{5.3} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell_i \sqcup \beta' \sqcup \gamma' (\mathbf{A})_{\ell_i \sqcup \beta' \sqcup \gamma'}$$

$$T_{5.4} = \text{SLIO } (\gamma') (\gamma') (\mathbf{A}^{\ell_i})_{\beta' \sqcup \gamma'}$$

$$T_{5.5} = \text{SLIO } (\gamma') (\gamma') (\tau)_{\beta' \sqcup \gamma'}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{SLIO}^*\text{-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5}, x : (\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} \vdash e_{c1} : T_3} \text{IH2, Weakening}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5}, y : (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell} \vdash e_{c2} : T_3} \text{IH3, Weakening}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{case}(b, x.e_{c1}, y.e_{c2}) : T_3} \text{SLIO}^*\text{-case}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{SLIO}^*\text{-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_4} \text{SLIO}^*\text{-bind}$$

P0:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_4} \text{SLIO}^*\text{-bind}$$

P0.1:

$$\frac{\frac{}{P0}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_5} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_{5.1}} \text{Lemma 3.31}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.2}} \text{Lemma 3.31}$$

Main derivation:

$$\begin{array}{c}
\text{P0.1} \quad \frac{\frac{\overline{\Sigma; \Psi \vdash A^{\ell_i} \searrow \ell} \text{ Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.3}} \\
\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.4}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.5}} \text{ Definition 3.30} \\
\frac{\overline{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq \alpha'} \text{ Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_1}
\end{array}$$

10. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})))} \text{ FC-ref}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((\text{ref } \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' ((\text{ref } A^{\ell_i})^\perp)_{\alpha'}$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\text{ref } A^{\ell_i}))_{\alpha'}$$

$$T_{1.3} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' \text{ ref } \ell_i (\mathbf{A})_{\ell_i}$$

$$T_2 = \text{SLIO } \gamma' \gamma' (\tau)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' (A^{\ell_i})_{(\beta' \sqcup \gamma')}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.4} = \text{SLIO } \gamma' \gamma' \text{ ref } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.5} = \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.51} = \text{Labeled } \alpha' \text{ ref } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.6} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' \text{ ref } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' \text{ ref } \ell_i (\mathbf{A})_{\ell_i}$$

P3:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash A^{\ell_i} \searrow pc} \text{ Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_i} \text{ By inversion} \quad \frac{\overline{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq pc} \text{ Given}}{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq \ell_i}}{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq \ell_i}$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}} \text{ SLIO}^*\text{-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}} \text{ SLIO}^*\text{-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{2.6}} \text{ SLIO}^*\text{-ret} \quad \text{P3} \\
\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{2.6}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{new } (a) : T_{2.4}}{\text{SLIO}^*\text{-new}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{L}b)) : T_{1.3}} \text{SLIO}^*\text{-bind}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{L}b))) : T_{1.3}} \text{SLIO}^*\text{-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{L}b))) : T_1} \text{Definition 3.30}$$

11. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau')_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' (A^{\ell'_i})_{\alpha'}$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell'_i \sqcup \alpha' (A^{\ell'_i})_{\ell'_i \sqcup \alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\text{ref } \tau)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) ((\text{ref } \tau)^\ell)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) ((\text{ref } A^{\ell'_i}))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \text{SLIO } \gamma' \gamma' \text{ Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i (A)_{\ell_i})$$

$$T_{2.4} = \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i (A)_{\ell_i})$$

$$T_{2.5} = \text{SLIO } \gamma' \beta' \sqcup \gamma' \sqcup \ell (\text{ref } \ell_i (A)_{\ell_i})$$

$$T_{2.6} = (\text{ref } \ell_i (A)_{\ell_i})$$

$$T_{2.7} = \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i (A)_{\ell_i})$$

$$T_{2.8} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i (A)_{\ell_i})$$

$$T_{2.9} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell'_i (A^{\ell'_i})_{\ell'_i})$$

$$T_{2.10} = \text{SLIO } (\gamma') (\gamma') (\text{Labeled } \beta' \sqcup \gamma' \sqcup \ell \sqcup \ell'_i (A^{\ell'_i})_{\ell'_i})$$

$$T_{2.11} = \text{SLIO } (\gamma') (\gamma') (\text{Labeled } \alpha \sqcup \ell'_i (A^{\ell'_i})_{\ell'_i})$$

P2:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}}{\text{SLIO}^*\text{-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{SLIO}^*\text{-deref}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}}{\text{SLIO}^*\text{-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{SLIO}^*\text{-bind}$$

P0:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_{2.3}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.8}} \text{SLIO}^*\text{-bind}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}} \quad P0}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{2.10}} \text{Lemma 3.33}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{SLIO}^*\text{-sub}}{\frac{\frac{\frac{}{\Sigma; \Psi \vdash A^{\ell_i} \searrow \ell} \quad \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \quad \text{By inversion}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha'} \quad \text{Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{SLIO}^*\text{-sub}}$$

12. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\text{unit})_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' \text{unit}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\text{ref } \tau)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') ((\text{ref } \tau))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') ((\text{ref } A^{\ell_i}))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \text{SLIO } \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (\mathbf{A})_{\ell_i}$$

$$T_{2.4} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (\mathbf{A})_{\ell_i}$$

$$T_{2.5} = \text{SLIO } \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (\mathbf{A})_{\ell_i}$$

$$T_{2.6} = \text{ref } \ell_i (\mathbf{A})_{\ell_i}$$

$$T_{2.7} = \text{SLIO } \ell \sqcup (\beta' \sqcup \gamma') \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.8} = \text{SLIO } \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.9} = \text{SLIO } \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_3 = \text{SLIO } \gamma' \gamma' (\tau)_{(\beta' \sqcup \gamma')}$$

$$T_{3.1} = \text{SLIO } \gamma' \gamma' (A^{\ell_i})_{(\beta' \sqcup \gamma')}$$

$$T_{3.2} = \text{SLIO } \gamma' \gamma' \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.4} = \text{Labeled } \ell_i (\mathbf{A})_{\ell_i}$$

P4:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{SLIO}^*\text{-var}$$

P5:

$$\begin{array}{c}
\frac{\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}}{\text{SLIO}^* \text{-var}}}{\Sigma; \Psi \vdash \tau = \mathbf{A}^{\ell_i} \searrow (pc \sqcup \ell)}{\text{Given}}}{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \ell_i} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.4}}
\end{array}$$

P3:

$$\frac{\frac{P4 \quad P5}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}}{\text{SLIO}^* \text{-assign}}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}}{\text{SLIO}^* \text{-unlabel}} \quad P3}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{SLIO}^* \text{-bind}}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash e_{c2} : T_{3.2}}{\text{IH2}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{SLIO}^* \text{-bind}}$$

P0:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash e_{c1} : T_{2.3}}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.9}} \text{SLIO}^* \text{-bind}}$$

P0.1:

$$\frac{\frac{P0}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}}{\text{SLIO}^* \text{-toLabeled}}$$

Main derivation:

$$\frac{\frac{P0.1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()) : T_{1.1}}{\text{SLIO}^* \text{-bind}}$$

13. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c))))} \text{FC-FI}$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((\forall \alpha. (\ell_e, \tau))^\perp)_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((\forall \alpha. (\ell_e, \tau))_{\alpha'})$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' \forall \alpha. \forall \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$



$$T_{2.1} = (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.2} = \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \alpha' \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi, (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma)_{\beta'} \vdash e_c : T_2}{\text{IH, Weakening}}}{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi; (\Gamma)_{\beta'} \vdash \nu(e_c) : T_{2.1}}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \Lambda\Lambda\Lambda(\nu(e_c)) : T_{2.2}}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{Lb}(\Lambda\Lambda\Lambda(\nu(e_c))) : T_{2.3}}}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_c)))) : T_{1.2}} \text{SLIO}^*\text{-CI}$$

$$\text{SLIO}^*\text{-FI}$$

$$\text{SLIO}^*\text{-label}$$

14. FC-FE:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c}{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash_{pc} \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[]\bullet)))} \text{FC-FE}$$

$$\beta' = \bigcup_{\beta_i \in \beta'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau[\ell''/\alpha])_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((\forall \alpha. (\ell_e, \tau))^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup (\beta' \sqcup \gamma') (\forall \alpha. (\ell_e, \tau))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.4} = \text{SLIO } \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.5} = \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.6} = \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}[\ell''/\alpha]$$

$$T_{2.7} = \forall \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.8} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.81} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.9} = \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.10} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.11} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau[\ell''/\alpha])_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.12} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i}[\ell''/\alpha])_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.13} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{ Labeled } \ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}[\ell''/\alpha])_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.14} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{ Labeled } \ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma' (\mathbf{A}[\ell''/\alpha])_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.15} = \text{SLIO } (\gamma') (\gamma') \text{ Labeled } \ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma' (\mathbf{A}[\ell''/\alpha])_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma')}$$

$T_{2.16} = \text{SLIO}(\gamma')(\gamma')(\mathbb{A}[\ell''/\alpha]^{\ell_i[\ell''/\alpha]})(\beta' \sqcup \gamma')$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e[\ell''/\alpha]} \text{Given} \quad \frac{}{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq pc \sqsubseteq \ell_e[\ell''/\alpha]} \text{Given}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq \ell_e[\ell''/\alpha])}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha])}}$$

P2:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}} \text{SLIO}^*\text{-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[] : T_{2.6}} \text{SLIO}^*\text{-FE}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][] : T_{2.7}} \text{SLIO}^*\text{-FE}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][][] : T_{2.81}} \text{SLIO}^*\text{-FE} \quad P3}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][][] \bullet : T_{2.9}} \text{SLIO}^*\text{-CE}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}} \text{SLIO}^*\text{-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[][][] \bullet) : T_{2.10}} \text{SLIO}^*\text{-bind}}$$

P0:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash e_c : T_{2.2}} \text{IH} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet)) : T_{2.10}} \text{SLIO}^*\text{-bind}}$$

P0.1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \mathbb{A}[\ell''/\alpha]^{\ell_i[\ell''/\alpha]} \searrow \ell} \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i[\ell''/\alpha]} \text{By inversion}}$$

P0.2:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet)) : T_{2.11}} \text{Lemma 3.36}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet)) : T_{2.12}} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet)) : T_{2.13}} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet)) : T_{2.14}} \text{Lemma 3.31}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce.taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet))) : T_{2.15}} \text{Lemma 3.31}$$

Main derivation:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce.taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet))) : T_1} \text{Definition 3.30}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce.taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][][] \bullet))) : T_1} \text{Definition 3.30}$$

15. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{FC-CI}$$

$$T_1 = \text{SLIO } \gamma' \gamma' ((c \xrightarrow{\ell_e} \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' ((c \xrightarrow{\ell_e} \tau))_{\alpha'}$$

$$T_{1.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.3} = \text{Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.4} = \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.5} = (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \text{SLIO } \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma, \alpha_i, \gamma_i; \Psi, (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma) \vdash e_c : T_2}{\Sigma; \Psi; \Gamma \vdash \nu(e_c) : T_{1.5}} \text{IH, Weakening}}{\Sigma; \Psi; \Gamma \vdash \Lambda\Lambda(\nu(e_c)) : T_{1.4}} \text{SLIO}^*\text{-CI}}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(\Lambda\Lambda(\nu(e_c))) : T_{1.3}} \text{SLIO}^*\text{-FI}}{\Sigma; \Psi; \Gamma \vdash \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) : T_{1.2}} \text{SLIO}^*\text{-label}$$

16. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b[][\bullet])))} \text{FC-CE}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \text{SLIO } \gamma' \gamma' (\tau)_{\alpha'}$$

$$T_2 = \text{SLIO } \gamma' \gamma' ((c \xrightarrow{\ell_e} \tau)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup (\beta' \sqcup \gamma') ((c \xrightarrow{\ell_e} \tau))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \text{SLIO } \gamma' \gamma' \text{ Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\ell \sqcup \alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\ell \sqcup \alpha_i}$$

$$T_{2.4} = \text{SLIO } \gamma' \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\ell \sqcup \alpha_i}$$

$$T_{2.5} = \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\ell \sqcup \alpha_i}$$

$$T_{2.6} = \forall \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma_i \gamma_i (\tau)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e \Rightarrow \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.71} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.8} = \text{SLIO } (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.9} = \text{SLIO } (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau)_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.10} = \text{SLIO}(\gamma')(\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i})_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.11} = \text{SLIO}(\gamma')(\beta' \sqcup \gamma' \sqcup \ell) \text{ Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.12} = \text{SLIO}(\gamma')(\gamma') \text{ Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.13} = \text{SLIO}(\gamma')(\gamma') \text{ Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.14} = \text{SLIO}(\gamma')(\gamma') (\mathbf{A}^{\ell_i})_{(\beta' \sqcup \gamma')}$$

P2:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{SLIO}^*\text{-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[] : T_{2.6}}{\text{SLIO}^*\text{-FE}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[][] : T_{2.71}}{\text{SLIO}^*\text{-FE}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.8}}{\text{SLIO}^*\text{-CE}}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{SLIO}^*\text{-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[][]\bullet) : T_{2.9}}{\text{SLIO}^*\text{-bind}}$$

P0:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet)) : T_{2.9}}{\text{SLIO}^*\text{-bind}}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{P0}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet)) : T_{2.10}}{\text{SLIO}^*\text{-bind}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet)) : T_{2.11}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet))) : T_{2.12}}{\text{Lemma 3.31}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet))) : T_{2.13}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet))) : T_{2.14}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][]\bullet))) : T_1}$$

□

**Lemma 3.33** (FG  $\rightsquigarrow$  SLIO\*: Subtyping preservation).  $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$  and the following holds:

1.  $\forall \tau, \tau'.$

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau' \rrbracket_{\ell'}$$

2.  $\forall \mathbf{A}, \mathbf{A}'.$

$$\Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}' \implies \Sigma; \Psi \vdash \llbracket \mathbf{A} \rrbracket_{\ell} <: \llbracket \mathbf{A}' \rrbracket_{\ell'}$$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash A_1 <: A_2} \text{ By inversion } P1}{\Sigma; \Psi \vdash (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2})} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion } \quad \frac{}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{ Given}}{\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \text{Labeled } \ell \sqcup \ell_1 (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: \text{Labeled } \ell' \sqcup \ell_2 (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2})} \text{ SLIO}^* \text{-sub-labeled}}{\Sigma; \Psi \vdash \llbracket A_1^{\ell_1} \rrbracket_{\ell} <: \llbracket A_2^{\ell_2} \rrbracket_{\ell'}}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\frac{\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{ SLIO}^* \text{-ref}}{\Sigma; \Psi \vdash \llbracket \mathbf{b} \rrbracket_{\ell} <: \llbracket \mathbf{b} \rrbracket_{\ell'}} \text{ Definition 3.30}$$

2. FGsub-ref:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{ref } \ell_i \llbracket A \rrbracket_{\ell_i} <: \text{ref } \ell_i \llbracket A \rrbracket_{\ell_i}} \text{ SLIO}^* \text{-ref}}{\Sigma; \Psi \vdash \llbracket \text{ref } A^{\ell_i} \rrbracket_{\ell} <: \llbracket \text{ref } A^{\ell_i} \rrbracket_{\ell'}} \text{ Definition 3.30}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell \times \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_\ell \times \llbracket \tau'_2 \rrbracket_{\ell'}}{\text{SLIO}^* \text{sub-prod}}}{\Sigma; \Psi \vdash \llbracket \tau_1 \times \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \times \tau'_2 \rrbracket_{\ell'}} \text{Definition 3.30}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}{\text{Given}}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}{\text{Given}}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_\ell + \llbracket \tau'_2 \rrbracket_{\ell'}}{\text{SLIO}^* \text{sub-prod}}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{Definition 3.30}$$

5. FGsub-arrow:

$$T_1 = \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.0} = \forall \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.2} = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.3} = (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha$$

$$c_1 = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.0} = \forall \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.2} = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.3} = (\tau'_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau'_2)_\alpha$$

$$c_2 = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2}{\text{Given}}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \text{SLIO } \gamma \gamma (\tau_2)_\alpha <: \text{SLIO } \gamma \gamma (\tau'_2)_\alpha} \text{IH(1) with } \ell = \ell' = \alpha, \text{SLIO}^* \text{sub-monad}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{Given}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau'_1 <: \tau_1} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma; \Psi \vdash T_{1.3} <: T_{2.3}} \text{SLIO*sub-arrow}$$

P0.1:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell \sqsubseteq \ell'} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha)} \implies (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha)}{\frac{\frac{}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \ell'_e)} \implies (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \ell_e)} \Sigma, \alpha, \beta, \gamma; \Psi \vdash c_2 \implies c_1$$

P0:

$$\frac{P0.1 \quad \frac{P1}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.3} <: T_{2.3}} \text{SLIO*sub-arrow}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}} \text{SLIO*sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{SLIO*sub-forall}$$

Main derivation:

$$\frac{P0}{\Sigma; \Psi \vdash \llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \xrightarrow{\ell'_e} \tau'_2 \rrbracket_{\ell'}} \text{Definition 3.30}$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{SLIO*sub-unit}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_\ell <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{Definition 3.30}$$

7. FGsub-forall:

$$T_1 = \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.0} = \forall \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.2} = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.3} = \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'}$$

$$c_1 = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.0} = \forall \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.2} = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.3} = \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}$$

$$c_2 = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\text{Given, Weakening}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \tau_1 <: \tau_2}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\tau_1)_{\alpha'} <: \tau_{2\alpha'}} \text{ IH(1) with } \ell = \ell' = \alpha'}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \text{SLIO } \gamma \gamma (\tau_1)_{\alpha'} <: \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}}$$

P2:

$$\frac{\frac{\text{Given}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell'_e \sqsubseteq \ell_e)}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \ell'_e) \Rightarrow (\ell \sqcup \gamma \sqsubseteq \ell_e)}$$

P1:

$$\frac{\frac{\text{Given}}{(\ell \sqsubseteq \ell')}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \alpha') \Rightarrow (\ell \sqcup \gamma \sqsubseteq \alpha')}$$

P0:

$$\frac{P1 \quad P2}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash c_2 \Rightarrow c_1}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad P3}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash T_{1.2} <: T_{2.2}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{ SLIO*sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{ SLIO*sub-forall}}{\Sigma; \Psi \vdash [\forall \alpha. \tau_1]_{\ell} <: [\forall \alpha. \tau_2]_{\ell'}} \text{ Definition 3.30}$$

8. FGsub-constraint:

$$T_1 = \forall \alpha, \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha}$$

$$T_{1.0} = \forall \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha}$$

$$T_{1.1} = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_1)_{\alpha}$$

$$T_{1.2} = \text{SLIO } \gamma \gamma (\tau_1)_{\alpha}$$

$$C_1 = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{2.0} = \forall \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{2.1} = (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{2.2} = \text{SLIO } \gamma \gamma (\tau_2)_{\alpha}$$

$$C_2 = (c_2 \wedge \ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$



P1:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash \tau_1 <: \tau_2} \text{Given, Weakening}}{\Sigma, \alpha, \gamma; \Psi \vdash (\tau_1)_\alpha <: \tau_{2\alpha}} \text{IH(1) with } \ell = \ell' = \alpha}}{\Sigma, \alpha, \gamma; \Psi \vdash \text{SLIO } \gamma \gamma (\tau_1)_\alpha <: \text{SLIO } \gamma \gamma (\tau_2)_\alpha}$$

P0:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash c_2 \implies c_1} \text{Given}}{\Sigma, \alpha, \gamma; \Psi \vdash c_2 \wedge (\ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \implies c_1 \wedge (\ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)} \text{Weakening, } \ell \sqsubseteq \ell', \ell'_e \sqsubseteq \ell_e}}{\Sigma, \alpha, \gamma; \Psi \vdash C_2 \implies C_1}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad P1}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.1} <: T_{2.1}} \text{SLIO*sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{SLIO*sub-forall}}{\Sigma; \Psi \vdash \left[ c_1 \xrightarrow{\ell_e} \tau_1 \right]_\ell <: \left[ c_2 \xrightarrow{\ell'_e} \tau_2 \right]_{\ell'}} \text{Definition 3.30}$$

□

**Lemma 3.34** (FG  $\rightsquigarrow$  SLIO\*: Preservation of well-formedness). *Forall  $\Sigma, \Psi$  and  $\ell$  s.t  $FV(\ell) \in \Sigma$  the following hold:*

1.  $\forall \tau. \Sigma; \Psi \vdash \tau \text{ WF} \implies \Sigma; \Psi \vdash (\tau)_\ell \text{ WF}$
2.  $\forall A. \Sigma; \Psi \vdash A \text{ WF} \implies \Sigma; \Psi \vdash (A)_\ell \text{ WF}$

*Proof.* Proof by simultaneous induction on the  $WF$  relation of FG

Proof of statement (1)

Let  $\tau = A^{\ell'}$

$$\frac{\frac{\frac{\frac{}{FV(\ell') \in \Sigma} \text{By inversion}}{FV(\ell' \sqcup \ell) \in \Sigma}}{\Sigma; \Psi \vdash (A)_{\ell' \sqcup \ell} \text{ WF}} \text{IH(2) on A}}{\Sigma; \Psi \vdash \text{Labeled } \ell' \sqcup \ell (A)_{\ell' \sqcup \ell} \text{ WF}} \text{SLIO* -wff-labeled}$$

Proof of statement (2)

We proceed by case analyzing the last rule of given  $WF$  judgment.

1. FG-wff-base:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{SLIO* -wff-base}$$

2. FG-wff-unit:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{SLIO* -wff-unit}$$

3. FG-wff-arrow:

P1:

$$\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash (\tau_2)_\alpha \text{ WF}} \text{ IH(1) on } \tau_2}{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \text{SLIO } \gamma \gamma (\tau_2)_\alpha \text{ WF}} \text{ SLIO}^*\text{-wff-monad}}$$

P0:

$$\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash (\tau_1)_\beta \text{ WF}} \text{ IH(1) on } \tau_1 \quad P1}{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash ((\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha) \text{ WF}} \text{ SLIO}^*\text{-wff-arrow}}$$

Main derivation:

$$\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi \vdash ((\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha) \text{ WF}} \text{ P0}}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha) \text{ WF}} \text{ SLIO}^*\text{-wff-constraint}}$$

4. FG-wff-prod:

$$\frac{\overline{\Sigma; \Psi \vdash (\tau_1)_\ell \text{ WF}} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash (\tau_2)_\ell \text{ WF}} \text{ IH(1) on } \tau_2}{\overline{\Sigma; \Psi \vdash (\tau_1)_\ell \times (\tau_2)_\ell \text{ WF}} \text{ SLIO}^*\text{-wff-prod}}$$

5. FG-wff-sum:

$$\frac{\overline{\Sigma; \Psi \vdash (\tau_1)_\ell \text{ WF}} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash (\tau_2)_\ell \text{ WF}} \text{ IH(1) on } \tau_2}{\overline{\Sigma; \Psi \vdash (\tau_1)_\ell + (\tau_2)_\ell \text{ WF}} \text{ SLIO}^*\text{-wff-prod}}$$

6. FG-wff-ref:

Let  $\tau = \mathbf{A}^{\ell'}$

$$\frac{\overline{\text{FV}(\mathbf{A}) = \emptyset} \text{ By inversion} \quad \overline{\text{FV}(\ell') = \emptyset} \text{ By inversion}}{\overline{\text{FV}((\mathbf{A})_{\ell'}) = \emptyset} \text{ Lemma 3.35}} \text{ SLIO}^*\text{-wff-ref}$$

$$\frac{}{\overline{\Sigma; \Psi \vdash \text{ref } \ell' (\mathbf{A})_{\ell'} \text{ WF}}}$$

7. FG-wff-forall:

$$\frac{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash (\tau)_{\alpha'} \text{ WF}} \text{ IH(1) on } \tau}{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash \text{SLIO } \gamma \gamma (\tau)_{\alpha'} \text{ WF}} \text{ SLIO}^*\text{-wff-monad}} \text{ SLIO}^*\text{-wff-constraint}$$

$$\frac{}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau)_{\alpha'}) \text{ WF}}}$$

8. FG-wff-constraint:

$$\frac{\frac{\frac{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash (\tau)_\alpha WF}{\text{IH(1) on } \tau}}{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \text{SLIO } \gamma \gamma (\tau)_\alpha WF}{\text{SLIO}^*\text{-wff-monad}}}{\Sigma, \alpha, \gamma; \Psi \vdash (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau)_\alpha WF}{\text{SLIO}^*\text{-wff-constraint}} \frac{}{\Sigma; \Psi \vdash (\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau)_\alpha WF)}$$

□

**Lemma 3.35** (FG  $\rightsquigarrow$  SLIO\*: Free variable lemma).  $\forall \Sigma, \ell. FV(\ell) \in \Sigma$ , the following hold

1.  $\forall \tau. FV((\tau)_\ell) \subseteq FV(\tau) \cup FV(\ell)$
2.  $\forall A. FV((A)_\ell) \subseteq FV(A) \cup FV(\ell)$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$

Proof for (1)

Let  $\tau = A^{\ell_i}$

$$\begin{aligned} & FV((A^{\ell_i})_\ell) \\ = & FV(\text{Labeled } \ell_i \sqcup \ell (A)_{\ell_i \sqcup \ell}) && \text{Definition 3.30} \\ = & FV(\ell_i) \cup FV(\ell) \cup FV((A)_{\ell_i \sqcup \ell}) \\ \subseteq & FV(\ell_i) \cup FV(\ell) \cup FV(A) && \text{IH(2) on } A \\ = & FV(A^{\ell_i}) \cup FV(\ell) \end{aligned}$$

Proof for (2)

1.  $A = \mathbf{b}$ :

$$\begin{aligned} & FV((\mathbf{b})_\ell) \\ = & FV(\mathbf{b}) && \text{Definition 3.30} \\ \subseteq & FV(\mathbf{b}) \cup FV(\ell) \end{aligned}$$

2.  $A = \text{unit}$ :

$$\begin{aligned} & FV((\text{unit})_\ell) \\ = & FV(\text{unit}) && \text{Definition 3.30} \\ \subseteq & FV(\text{unit}) \cup FV(\ell) \end{aligned}$$

3.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\begin{aligned} & FV((\tau_1 \xrightarrow{\ell_e} \tau_2)_\ell) \\ = & FV(\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_\alpha) && \text{Definition 3.30} \\ = & FV(\ell) \cup FV((\tau_1)_\beta) \cup FV(\ell_e) \cup FV((\tau_2)_\alpha) \\ \subseteq & FV(\tau_1) \cup FV(\ell_e) \cup FV(\tau_2) \cup FV(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \xrightarrow{\ell_e} \tau_2) \cup FV(\ell) \end{aligned}$$

4.  $A = \tau_1 \times \tau_2$ :

$$\begin{aligned} & FV((\tau_1 \times \tau_2)_\ell) \\ = & FV((\tau_1)_\ell \times (\tau_2)_\ell) && \text{Definition 3.30} \\ = & FV((\tau_1)_\ell) \cup FV((\tau_2)_\ell) \cup FV(\ell) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \cup FV(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \times \tau_2) \cup FV(\ell) \end{aligned}$$

5.  $A = \tau_1 + \tau_2$ :

$$\begin{aligned}
& \text{FV}(\langle \tau_1 + \tau_2 \rangle_\ell) \\
&= \text{FV}(\langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell) && \text{Definition 3.30} \\
&= \text{FV}(\langle \tau_1 \rangle_\ell) \cup \text{FV}(\langle \tau_2 \rangle_\ell) \cup \text{FV}(\ell) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 + \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$

6.  $A = \text{ref } \tau_i$ :

Let  $\tau_i = A_i^{\ell_i}$

$$\begin{aligned}
& \text{FV}(\langle \text{ref } \tau_i \rangle_\ell) \\
&= \text{FV}(\text{ref } \ell_i \langle A_i \rangle) && \text{Definition 3.30} \\
&= \text{FV}(\ell_i) \cup \text{FV}(\langle A_i \rangle) \\
&\subseteq \text{FV}(\ell_i) \cup \text{FV}(A_i) \cup \text{FV}(\ell) && \text{IH(2) on } A_i \\
&= \text{FV}(\text{ref } A_i^{\ell_i}) \cup \text{FV}(\ell) \\
&= \text{FV}(\text{ref } \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$

7.  $A = \forall \alpha. (\ell_e, \tau_i)$ :

$$\begin{aligned}
& \text{FV}(\langle \forall \alpha. (\ell_e, \tau_i) \rangle) \\
&= \text{FV}(\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma \langle \tau_i \rangle_{\alpha'}) && \text{Definition 3.30} \\
&= \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_i \rangle) \\
&\subseteq \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_i) && \text{IH(1) on } \tau_i \\
&= \text{FV}(\ell) \cup \text{FV}(\forall \alpha. (\ell_e, \tau_i))
\end{aligned}$$

8.  $A = c \xrightarrow{\ell_c} \tau_i$ :

$$\begin{aligned}
& \text{FV}(\langle c \xrightarrow{\ell_c} \tau_i \rangle) \\
&= \text{FV}(\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma \langle \tau \rangle_\alpha) && \text{Definition 3.30} \\
&= \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\langle \tau_i \rangle) \cup \text{FV}(\ell) \\
&\subseteq \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\tau_i) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_i \\
&= \text{FV}(c \xrightarrow{\ell_c} \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$

□

**Lemma 3.36** (FG  $\rightsquigarrow$  SLIO\*: Substitution lemma).  $\forall \tau, A, \ell$  s.t.  $\alpha \notin \text{FV}(\ell)$ ,  $\vdash \tau$  WF and  $\vdash A$  WF. The following holds

1.  $(\langle \tau \rangle_\ell[\ell'/\alpha]) = \langle \tau[\ell'/\alpha] \rangle_\ell$
2.  $(\langle A \rangle_\ell[\ell'/\alpha]) = \langle A[\ell'/\alpha] \rangle_\ell$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $A$

Proof for (1)

Let  $\tau = A_i^{\ell_i}$

$$\begin{aligned}
& (\langle A_i^{\ell_i} \rangle_\ell)[\ell'/\alpha] \\
&= (\text{Labeled } (\ell_i \sqcup \ell) \langle A \rangle_{\ell_i \sqcup \ell})[\ell'/\alpha] && \text{Definition 3.30} \\
&= (\text{Labeled } (\ell_i[\ell'/\alpha] \sqcup \ell) \langle A \rangle_{\ell_i[\ell'/\alpha] \sqcup \ell})[\ell'/\alpha] \\
&= (\text{Labeled } (\ell_i[\ell'/\alpha] \sqcup \ell) \langle A[\ell'/\alpha] \rangle_{\ell_i[\ell'/\alpha] \sqcup \ell}) && \text{IH(2)} \\
&= (\langle A[\ell'/\alpha] \rangle_{\ell_i[\ell'/\alpha] \sqcup \ell}) \\
&= (\langle A_i^{\ell_i}[\ell'/\alpha] \rangle)_\ell
\end{aligned}$$

Proof for (2)

1.  $A = \mathbf{b}$ :

$$\begin{aligned}
& ((\mathbf{b})_\ell)[\ell'/\alpha] \\
&= (\mathbf{b})[\ell'/\alpha] && \text{Definition 3.30} \\
&= \mathbf{b} \\
&= ((\mathbf{b}))_\ell \\
&= ((\mathbf{b}[\ell'/\alpha]))_\ell
\end{aligned}$$

2.  $A = \mathbf{unit}$ :

$$\begin{aligned}
& ((\mathbf{unit})_\ell)[\ell'/\alpha] \\
&= (\mathbf{unit})[\ell'/\alpha] && \text{Definition 3.30} \\
&= \mathbf{unit} \\
&= ((\mathbf{unit}))_\ell \\
&= ((\mathbf{unit}[\ell'/\alpha]))_\ell
\end{aligned}$$

3.  $A = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\begin{aligned}
& ((\tau_1 \xrightarrow{\ell_e} \tau_2)_\ell)[\ell'/\alpha] \\
&= (\forall \alpha', \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'})[\ell'/\alpha] && \text{Definition 3.30} \\
&= (\forall \alpha', \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow (\tau_1)_\beta[\ell'/\alpha] \rightarrow \text{SLIO } \gamma \gamma (\tau_2)_{\alpha'}[\ell'/\alpha]) \\
&= (\forall \alpha', \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow (\tau_1[\ell'/\alpha])_\beta \rightarrow \text{SLIO } \gamma \gamma (\tau_2[\ell'/\alpha])_{\alpha'}) && \text{IH(1)} \\
&= ((\tau_1[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_2[\ell'/\alpha]))_\ell \\
&= ((\tau_1 \xrightarrow{\ell_e} \tau_2)[\ell'/\alpha])_\ell
\end{aligned}$$

4.  $A = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& ((\tau_1 \times \tau_2)_\ell)[\ell'/\alpha] \\
&= ((\tau_1)_\ell \times (\tau_2)_\ell)[\ell'/\alpha] && \text{Definition 3.30} \\
&= ((\tau_1)_\ell[\ell'/\alpha] \times (\tau_2)_\ell[\ell'/\alpha]) \\
&= ((\tau_1[\ell'/\alpha])_\ell \times (\tau_2[\ell'/\alpha])_\ell) && \text{IH(1)} \\
&= ((\tau_1[\ell'/\alpha] \times \tau_2[\ell'/\alpha]))_\ell \\
&= ((\tau_1 \times \tau_2)[\ell'/\alpha])_\ell
\end{aligned}$$

5.  $A = \tau_1 + \tau_2$ :

$$\begin{aligned}
& ((\tau_1 + \tau_2)_\ell)[\ell'/\alpha] \\
&= ((\tau_1)_\ell + (\tau_2)_\ell)[\ell'/\alpha] && \text{Definition 3.30} \\
&= ((\tau_1)_\ell[\ell'/\alpha] + (\tau_2)_\ell[\ell'/\alpha]) \\
&= ((\tau_1[\ell'/\alpha])_\ell + (\tau_2[\ell'/\alpha])_\ell) && \text{IH(1)} \\
&= ((\tau_1[\ell'/\alpha] + \tau_2[\ell'/\alpha]))_\ell \\
&= ((\tau_1 + \tau_2)[\ell'/\alpha])_\ell
\end{aligned}$$

6.  $A = \text{ref } \tau_i$ :

Let  $\tau_i = A_i^{\ell_i}$

$$\begin{aligned}
& ((\text{ref } \tau_i)_\ell)[\ell'/\alpha] \\
&= (\text{ref } \ell_i (A_i))[\ell'/\alpha] && \text{Definition 3.30} \\
&= (\text{ref } \ell_i (A_i)) && \text{Lemma 3.34} \\
&= ((\text{ref } A_i^{\ell_i}))_\ell \\
&= ((\text{ref } A_i^{\ell_i})[\ell'/\alpha])_\ell && \text{Since } \vdash \text{ref } \tau_i \text{ WF} \\
&= ((\text{ref } \tau_i)[\ell'/\alpha])_\ell
\end{aligned}$$

7.  $A = \forall \alpha''. (\ell_e, \tau_i)$ :

$$\begin{aligned}
& ((\forall \alpha''. (\ell_e, \tau_i)))[\ell'/\alpha] \\
= & (\forall \alpha'', \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i)_{\alpha'})[\ell'/\alpha] && \text{Definition 3.30} \\
= & (\forall \alpha'', \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i)_{\alpha'}[\ell'/\alpha]) \\
= & (\forall \alpha'', \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i[\ell'/\alpha])_{\alpha'}) && \text{IH(1)} \\
= & ((\forall \alpha''. (\ell_e[\ell'/\alpha], \tau_i[\ell'/\alpha])))_{\ell} \\
= & ((\forall \alpha''. (\ell_e, \tau_i))[\ell'/\alpha])_{\ell}
\end{aligned}$$

8.  $A = c \xrightarrow{\ell_e} \tau_i$ :

$$\begin{aligned}
& ((c \xrightarrow{\ell_e} \tau_i))[\ell'/\alpha] \\
= & (\forall \alpha', \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i)_{\alpha'})[\ell'/\alpha] && \text{Definition 3.30} \\
= & (\forall \alpha', \gamma. (c[\ell'/\alpha] \wedge \ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i)_{\alpha'}[\ell'/\alpha]) \\
= & (\forall \alpha', \gamma. (c[\ell'/\alpha] \wedge \ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e[\ell'/\alpha]) \Rightarrow \text{SLIO } \gamma \gamma (\tau_i[\ell'/\alpha])_{\alpha'}) && \text{IH(1)} \\
= & ((c[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_i[\ell'/\alpha]))_{\ell} \\
= & ((c \xrightarrow{\ell_e} \tau_i)[\ell'/\alpha])_{\ell}
\end{aligned}$$

□

### 3.3.3 Model for FG to SLIO\* translation

**Definition 3.37** (FG  $\rightsquigarrow$  SLIO\*:  ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$

$$\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$$

**Definition 3.38** (FG  $\rightsquigarrow$  SLIO\*:  $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$

$$\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$$

**Definition 3.39** (FG  $\rightsquigarrow$  SLIO\*: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\
\llbracket \text{unit} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\
& \quad ({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\
& \quad \{({}^s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \lambda x. e_s, \Lambda \Lambda(\nu(\lambda x. e_t))) \mid \\
& \quad \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \implies \\
& \quad ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \forall \alpha. (\ell_e, \tau) \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \Lambda e_s, \Lambda \Lambda(\nu(e_t))) \mid \\
& \quad \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \nu e_s, \Lambda(\nu(e_t))) \mid \\
& \quad \mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \text{ref } \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, a_s, a_t) \mid {}^s\theta(a_s) = \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\
\llbracket \mathbf{A}' \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^s v, \text{Lb}({}^t v)) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \mathbf{A} \rrbracket_V^{\hat{\beta}}\}
\end{aligned}$$

**Definition 3.40** (FG  $\rightsquigarrow$  SLIO\*: Unary expression relation).

$$\begin{aligned} \llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv.(H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\ &\exists H'_t, {}^tv.(H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \\ &\wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}\} \end{aligned}$$

**Definition 3.41** (FG  $\rightsquigarrow$  SLIO\*: Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a_1) \rrbracket_V^{\hat{\beta}} \end{aligned}$$

**Definition 3.42** (FG  $\rightsquigarrow$  SLIO\*: Label substitution).  $\sigma : Lvar \mapsto Label$

**Definition 3.43** (FG  $\rightsquigarrow$  SLIO\*: Value substitution to values).  $\delta^s : Var \mapsto Val, \delta^t : Var \mapsto Val$

**Definition 3.44** (FG  $\rightsquigarrow$  SLIO\*: Unary interpretation of  $\Gamma$ ).

$$\begin{aligned} \llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ &\forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\} \end{aligned}$$

### 3.3.4 Soundness proof for FG to SLIO\* translation

**Lemma 3.45** (FG  $\rightsquigarrow$  SLIO\*: Monotonicity).  $\forall {}^s\theta, {}^s\theta', n, {}^sv, {}^tv, n', \beta, \beta'$ .

1.  $\forall A. ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket A \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsupseteq {}^s\theta' \wedge \hat{\beta} \sqsupseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^sv, {}^tv) \in \llbracket A \rrbracket_V^{\hat{\beta}'}$
2.  $\forall \tau. ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsupseteq {}^s\theta' \wedge \hat{\beta} \sqsupseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}$

*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We case analyze  $A$  in the last step

1. Case **b**:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsupseteq {}^s\theta' \wedge \hat{\beta} \sqsupseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}}$  therefore from Definition 3.39 we know that  ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$  and  ${}^sv = {}^tv$

Therefore from Definition 3.39 we get the desired

2. Case  $\text{unit}$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}}$  therefore from Definition 3.39 we know that  ${}^sv \in \llbracket \text{unit} \rrbracket \wedge {}^tv \in \llbracket \text{unit} \rrbracket$

Therefore from Definition 3.39 we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.39 we know that  ${}^sv = ({}^sv_1, {}^sv_2)$  and  ${}^tv = ({}^tv_1, {}^tv_2)$ .

We also know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1:  $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

IH2:  $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 3.39, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.39 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

IH:  $({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 3.39 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case



5. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}'}$$

From Definition 3.39 we know that

${}^sv$  is of the form  $\lambda x.e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\Lambda\Lambda\Lambda(\nu(\lambda x.e_t))$  (for some  $e_t$ ) s.t

$$({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'_1} \quad (\text{A0})$$

Similarly from Definition 3.39 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''} \implies \\ ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

Instantiating (A0) with  ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}''$  since

${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case  $\forall\alpha.\tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}'}$$

From Definition 3.39 we know that  ${}^sv = \Lambda e'_s$  and  ${}^tv = \Lambda\Lambda\Lambda(\nu(e_t))$  s.t

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1 . ({}^s\theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'_1} \quad (\text{F0})$$

Similarly from Definition 3.39 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

This means we are given  ${}^s\theta''_1 \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

Instantiating (F0) with  ${}^s\theta''_1, k, \hat{\beta}''$  since  ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

7. Case  $c \stackrel{\ell_\xi}{\Rightarrow} \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [c \stackrel{\ell_\xi}{\Rightarrow} \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \stackrel{\ell_\xi}{\Rightarrow} \tau]_V^{\hat{\beta}'}$$

From Definition 3.39 we know that  ${}^s v = \nu(e'_s)$  and  ${}^t v = \Lambda\Lambda(\nu(e_t))$ . And

$$\mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in [\tau]_E^{\hat{\beta}'} \quad (C0)$$

Similarly from Definition 3.39 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

This means we are given  $\mathcal{L} \models c, {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

Since  $\mathcal{L} \models c$  and instantiating (C0) with  ${}^s\theta''_1, k, \hat{\beta}''$  since  ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

8. Case ref  $\tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

From Definition 3.39 we know that  ${}^s v = a_s$  and  ${}^t v = a_t$ . We also know that

$${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$$

From Definition 3.39, Definition 3.37 and Definition 3.38 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

Proof of Statement (2)

Let  $\tau = \mathbf{A}^{\ell''}$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 3.39 we know that

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \text{ and } ({}^s\theta, n, {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}'}$$

This means from Definition 3.39 we need to prove

$$({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_V^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_V^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

**Lemma 3.46** (FG  $\rightsquigarrow$  SLIO\*: Unary monotonicity for  $\Gamma$ ).  $\forall {}^s\theta, {}^s\theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies ({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 3.44 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$$

And again from Definition 3.44 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 3.45 we get

$$\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

□

**Lemma 3.47** (FG  $\rightsquigarrow$  SLIO\*: Unary monotonicity for  $H$ ).  $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$

From Definition 3.41 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 3.41 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n' - 1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$  (given)

Therefore from Lemma 3.45 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n' - 1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$$

□

**Lemma 3.48** (Coercion lemma).  $\forall H, e, v.$

$$(H, e) \Downarrow_{-}^f (H', \text{Lb } v) \implies (H, \text{coerce\_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$$

*Proof.* Given:  $(H, e) \Downarrow_{-}^f (H', \text{Lb } v)$

To prove:  $(H, \text{coerce\_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and SLIO\*-Sem-app it suffices to prove that  $(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_{-}^f (H', \text{Lb } v)$

From SLIO\*-Sem-tolabeled it suffices to prove that  $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From SLIO\*-Sem-bind it suffices to prove that

1.  $(H, e) \Downarrow_{-}^f (H'_1, v_1)$ :

We are given that  $(H, e) \Downarrow_{-}^f (H', v)$  therefore we have  $H'_1 = H'$  and  $v'_1 = \text{Lb } v$

2.  $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$ :

It suffices to prove that

$$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v):$$

We get this directly from SLIO\*-Sem-unlabel

□

**Theorem 3.49** (FG  $\rightsquigarrow$  SLIO\*: Fundamental theorem).  $\forall \Sigma, \Psi, \Gamma, \tau, e_s, e_t, pc, \mathcal{L}, \delta^s, \delta^t, \sigma, {}^s\theta, n, \hat{\beta}.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}} \\ & \implies \\ & ({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}} \end{aligned}$$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x \delta^s, \text{ret}(x) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 3.40 it suffices to prove that

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{ret}(x) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge \\ ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  $i = 0, {}^s v = x \delta^s$ . Also from SLIO\*-Sem-ret we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$

And we are required to prove

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ : Given

(b)  $({}^s\theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}}$ :

Since we are given  $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_V^{\hat{\beta}}$ , therefore from Definition 3.44 we get  $({}^s\theta, n, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}}$

## 2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t))))} \text{FC-lam}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\lambda x. e_s) \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^{\hat{\beta}}$

From Definition 3.40 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t.  $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  ${}^s v = (\lambda x. e_s) \delta^s, H'_s = H_s$  and  $i = 0$ . Also from SLIO\*-Sem-ret, SLIO\*-Sem-label and SLIO\*-Sem-FI we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n, H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ : Given

(b)  $({}^s\theta, n, \lambda x.e_s \delta^s, (\mathbf{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_{V}^{\hat{\beta}}$ :

From Definition 3.39 it suffices to prove that

$$({}^s\theta, n, \lambda x.e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(\lambda x.e_t))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_{V}^{\hat{\beta}}$$

Again from Definition 3.39 it suffices to prove that

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'.({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'} \implies$$

$$({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'}$$

This further means that given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

And we are required to prove

$$({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since we are given  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$ , therefore from Definition 3.44 and Lemma 3.46 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\}) \sigma]_{V}^{\hat{\beta}'}$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'}$$

We get (F-L0) directly from IH

### 3. FC-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b))))} \text{FC-app}}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$$

This means from Definition 3.40 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v.(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t

$$(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c \square \square \square \square \bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-A0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \\ & ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - \\ & j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-A1.0}) \end{aligned}$$

Since we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 3.39 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 3.39 we know that  ${}^s v_1 = \lambda x. e'_s$  and  ${}^t v_i = \Lambda \Lambda \Lambda (\nu (\lambda x. e'_t))$  s.t

$$\begin{aligned} & \forall {}^s \theta''_1 \sqsupseteq {}^s \theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}''_1 \sqsubseteq \hat{\beta}'_1. \\ & ({}^s \theta''_1, l, {}^s v', {}^t v') \in [\tau_1 \sigma]_V^{\hat{\beta}''_1} \implies ({}^s \theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}''_1} \quad (\text{F-A1}) \end{aligned}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 3.40 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n - j, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}''_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - \\ & j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}''_2} \end{aligned}$$

We instantiate with  $H'_{s1}, H'_{t1}$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s1}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$ .

This means we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}''_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - \\ & j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}''_2} \quad (\text{F-A2}) \end{aligned}$$

We instantiate (F-A1) with  $\theta_1''$  as  $\theta_2'$ ,  ${}^s v'$  as  ${}^s v_2$ ,  ${}^t v'$  as  ${}^t v_2$ ,  $l$  as  $n - j - k$  and  $\hat{\beta}_1''$  as  $\hat{\beta}_2'$ . Therefore we get

$$({}^s \theta_2', n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}_2'}$$

From Definition 3.40 we have

$$\forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}_2'} {}^s \theta_2' \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s3}, {}^s v_3) \implies \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta_3' \sqsupseteq {}^s \theta_2', \hat{\beta}_3' \sqsupseteq \hat{\beta}_2'.$$

$$(n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}_3'} {}^s \theta_3' \wedge ({}^s \theta_3', n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_V^{\hat{\beta}_3'}$$

Instantiating with  $H'_{s2}, H'_{t2}$ . since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists a < i - j - k < n - j - k$  s.t  $(H'_{s2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s3}, {}^s v_3)$

Therefore we have

$$\exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta_3' \sqsupseteq {}^s \theta_2', \hat{\beta}_3' \sqsupseteq \hat{\beta}_2'.$$

$$(n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}_3'} {}^s \theta_3' \wedge ({}^s \theta_3', n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_V^{\hat{\beta}_3'} \quad (\text{F-A3})$$

Let  $\tau_2 \sigma = \mathbf{A}_2^{\ell_i}$ , since  $\tau_2 \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s \theta_3', n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_V^{\hat{\beta}_3'}$$

Therefore from Definition 3.39 we know that

$$({}^s \theta_3', n - j - k - a, {}^s v_3, \mathbf{Lb}^t v_{3i}) \in [\tau_2 \sigma]_V^{\hat{\beta}_3'} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose  $H'_t$  as  $H'_{t3}$  and  ${}^t v$  as  $\mathbf{Lb}^t v_{3i}$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \mathbf{Lb}^t v_{3i}):$$

From Lemma 3.48 it suffices to prove that

$$(H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \mathbf{Lb}^t v_{3i})$$

From SLIO\*-Sem-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \mathbf{Lb}^t v_{3i}):$

From SLIO\*-Sem-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \mathbf{Lb}^t v_{3i}):$

From SLIO\*-Sem-bind again it suffices to prove

- \*  $(H'_{t2}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \mathbf{Lb}^t v_i$

Therefore from SLIO\*-Sem-unlabel and (F-A1) we know that  $H'_{t31} = H'_{t2}$  and  ${}^t v_{t2} = {}^t v_i = \Lambda\Lambda\Lambda(\nu(\lambda x. e'_t))$



\*  $((c \square \square \square \bullet b)[{}^t v_2/b][{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}$ :

It suffices to prove that

$((\Lambda \Lambda \Lambda (\nu(\lambda x. e'_t))) \square \square \square \bullet {}^t v_2) \delta^t \Downarrow {}^t v_{t21}$

From SLIO\*-Sem-FE it suffices to prove that

$((\Lambda \Lambda (\nu(\lambda x. e'_t))) \square \square \bullet {}^t v_2) \delta^t \Downarrow {}^t v_{t21}$

Again from SLIO\*-Sem-FE applied two times it suffices to prove that

$(\nu(\lambda x. e'_t) \bullet {}^t v_2) \delta^t \Downarrow {}^t v_{t21}$

From SLIO\*-Sem-CE it suffices to prove that

$((\lambda x. e'_t) {}^t v_2) \delta^t \Downarrow {}^t v_{t21}$

From SLIO\*-Sem-app we know that

${}^t v_{t21} = e'_t[{}^t v_2/x] \delta^t$

\*  $(H'_{t2}, {}^t v_{21}) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i}))$ :

We get this from (F-A3) and (F-A3.1)

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_3$  and  $\hat{\beta}'$  as  $\hat{\beta}'_3$ . From fg-app we know that  $i = j + k + a + 1$ ,  ${}^s v = {}^s v_3$  and  $H'_s = H'_{s3}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t3}$  and  ${}^t v = \text{Lb}({}^t v_3)$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta'$  from (F-A3) and Lemma 3.47

Since  ${}^t v = \text{Lb}({}^t v_3)$  therefore from Definition 3.39 it suffices to prove that

$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}$

We get this directly from (F-A3) and Lemma 3.45

#### 4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \lfloor \Gamma \sigma \rfloor_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \in \lfloor (\tau_1 \times \tau_2)^\perp \sigma \rfloor_V^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \lfloor (\tau_1 \times \tau_2)^\perp \sigma \rfloor_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v_1, {}^s v_2$  s.t  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in \lfloor (\tau_1 \times \tau_2)^\perp \sigma \rfloor_V^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 3.40 we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists k < i - j < n - j$  s.t  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose  $H_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_1, {}^t v_2)$

(a)  $(H_t, (\text{bind}(e_{t1}, a. \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :

From SLIO\*-Sem-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$ :  
From (F-P1) we know that  $H'_{tb1} = H'_{t1}$  and  ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))) [{}^t v_1/a] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From SLIO\*-Sem-bind it suffices to prove that
  - $(H_t, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$ :  
From (F-P2) we know that  $H'_{tb2} = H'_{t2}$  and  ${}^t v_{tb2} = {}^t v_2$
  - $(H'_{t2}, \text{ret}(\text{Lb}(a, b))) [{}^t v_1/a] [{}^t v_2/b] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From SLIO\*-Sem-ret, (F-P1) and (F-P2)

(b)  $\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \hat{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'}$ :  
 We choose  ${}^s\theta'$  as  ${}^s\theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$  and since from fg-prod  $i = j + k + 1$  and  $H'_s = H'_{s2}$ .  
 Therefore from (F-P2) and Lemma 3.47 we get

$$(n - i, H'_s, H'_{t2}) \hat{\triangleright} {}^s\theta'$$

In order to prove  $({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'}$

From Definition 3.39 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_{V'}^{\hat{\beta}'_2}$$

Since  ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$  therefore we get the desired from (F-P1), (F-P2), Definition 3.39 and Lemma 3.45

5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V'}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t) \in [(\tau_1 \sigma)]_E^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \stackrel{\gamma, \hat{\beta}}{\triangleright} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1}$  therefore from Definition 3.39 we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 3.39 we know that  ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$  and  ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$  s.t

$$({}^s \theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$

Let  $\tau_1 \sigma = \mathbf{A}_1^{\ell_i}$ , since  $\tau_1 \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$\text{Since } ({}^s \theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1}$$

Therefore from Definition 3.39 we know that

$$({}^s \theta'_1, n - j, {}^s v_{i1}, \text{Lb}({}^t v_{i11})) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{i11})$  as we need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11})):$$

From Lemma 3.48 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11}))$$

From SLIO\*-Sem-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this from (F-F1)

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11})):$

Again from SLIO\*-Sem-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

Since  ${}^t v_1 = \text{Lb}({}^t v_{i1}, {}^t v_{i2})$  from (F-F1.1) and (F-F1.2) therefore we get the desired from SLIO\*-Sem-unlabel

So,  $H_{t21} = H'_{t1}$  and  ${}^t v_{t21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11})):$

We get this from SLIO\*-Sem-fst, SLIO\*-Sem-ret and (F-F1.2) and (F-F1.3)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1}:$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . And from fg-fst we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-F1) and Lemma 3.47 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$$

Since from fg-fst we know that  ${}^s v = {}^s v_{i1}$  therefore from (F-F1.2) and Lemma 3.45 we get

$$({}^s \theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1}$$

## 6. FC-snd:

Symmetric reasoning as in the FC-fst case

7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))} \text{inl}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 3.40 we have

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 3.40 we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_t, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-IL1}) \end{aligned}$$

In order to prove (F-IL0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $(\text{Lb inl}({}^t v_1))$  and we need to prove:

(a)  $(H'_{t1}, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1)))$ :

From SLIO\*-Sem-bind it suffices to prove that

i.  $(H'_{t1}, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :

From (F-IL1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

ii.  $(H'_{t1}, \text{ret}(\text{Lbinl}(a)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1)))$ :

We get this from SLIO\*-Sem-ret, (F-IL1)

(b)  $\exists^s \theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{V}}^{\hat{\beta}'}$ :

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . Since from fg-inl we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-IL1) and Lemma 3.47 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$$

Now we need to prove  $({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{V}}^{\hat{\beta}'}$

Since  ${}^s v = \text{inl } {}^s v_1$  and  ${}^t v = \text{Lb}(\text{inl}({}^t v_1))$  therefore from Definition 3.39 it suffices to prove that

$$({}^s\theta', n - i, \text{inl } {}^s v_1, \text{inl } {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}'}$$

Since from (F-IL1) we know that  $({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}'}$

Therefore from Lemma 3.45 and Definition 3.39 we get

$$({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}'}$$

8. FC-inr:

Symmetric reasoning as in the FC-inl case

9. FC-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

This means we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-C0})$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell \sigma]_{\hat{\beta}}^E$$

This means from Definition 3.40 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}.(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1.(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1.(H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_{\hat{\beta}'_1}^V \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1.(H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_{\hat{\beta}'_1}^V \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_{\hat{\beta}'_1}^V$  therefore from Definition 3.39 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2) \sigma]_{\hat{\beta}'_1}^V \quad (\text{F-C1.1})$$

2 cases arise

- (a)  ${}^s v_1 = \text{inl}({}^s v_{i1})$  and  ${}^t v_i = \text{inl}({}^t v_{i1})$ :

Also from Lemma 3.46 and Definition 3.44 we know that

$$({}^s\theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\}) \sigma]_{\hat{\beta}'_1}^V$$

IH2:

$$({}^s\theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau \sigma]_{\hat{\beta}'_1}^E$$

This means from Definition 3.40 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}.(n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2.(H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{\beta}'_2}^V \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{\beta}'_2}^V \quad (\text{F-C2}) \end{aligned}$$

Let  $\tau \sigma = \mathbf{A}_2^{\ell_i}$ , since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{\beta}'_2}^V$$

Therefore from Definition 3.39 we know that

$$({}^s\theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{i1})) \in [\tau \sigma]_{\hat{\beta}'_2}^V \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{i1})$

And we need to prove:

i.  $(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}^t v_{2i})$ :

From Lemma 3.48 it suffices to prove that

$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}^t v_{2i})$

From SLIO\*-Sem-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :

From (F-C1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{2i})$ :

From SLIO\*-Sem-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :

Since from (F-C1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  therefore from SLIO\*-Sem-unlabel we know that

$H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$

- $(\text{case}(b, x.e_{t1}, y.e_{t2}) [{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$ :

Since we know that in this case  ${}^t v_i = \text{inl}({}^t v_{i1})$

Therefore from SLIO\*-Sem-case we know that  ${}^t v_{t22} = e_{t1} [{}^t v_{i1}/x] \delta^t$

- $(H'_{t1}, e_{t1} [{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \text{Lb}^t v_{2i})$ :

We get this from (F-C2) and (F-C2.1)

ii.  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . Since from fg-case we know that  $i = j + k + 1$  and  $H'_s = H'_{s2}$  therefore from (F-C2) and Lemma 3.47 we get

$(n - i, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$

Now we need to prove  $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

Since  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  and since from (F-C2) we know that

$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

Therefore from Lemma 3.45 and Definition 3.39 we get

$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

(b)  ${}^s v_1 = \text{inr}({}^s v_{i1})$  and  ${}^t v_1 = \text{inr}({}^t v_{i1})$ :

Symmetric reasoning as in the previous case

## 10. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_s : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t))))} \text{FI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \lfloor \Gamma \sigma \rfloor_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \Lambda e_s \delta^s, \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \in \lfloor \tau \sigma \rfloor_E^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies$   
 $\exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$

$(n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'}$



This means given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \stackrel{\gamma, \hat{\beta}}{\triangleright} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\forall \alpha_g. (\ell_e, \tau))^\perp \sigma]_{\hat{V}}^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from SLIO\*-Sem-ret we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\forall \alpha_g. (\ell_e, \tau))^\perp \sigma]_{\hat{V}}^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \stackrel{\hat{\beta}}{\triangleright} {}^s\theta$ : Given

(b)  $({}^s\theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_t)))) \delta^t) \in [(\forall \alpha_g. (\ell_e, \tau))^\perp \sigma]_{\hat{V}}^{\hat{\beta}}$ :

From Definition 3.39 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(e_t))) \delta^t) \in [(\forall \alpha_g. (\ell_e, \tau)) \sigma]_{\hat{V}}^{\hat{\beta}}$$

Again from Definition 3.39 it suffices to prove that

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{E}}^{\hat{\beta}'_1}$$

This further means that given  ${}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we need to prove

$$({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g]]_{\hat{E}}^{\hat{\beta}'_1} \quad (\text{F-FI0})$$

$$\underline{\text{IH}}: ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha_g \mapsto \ell'\}]_{\hat{E}}^{\hat{\beta}'_1}$$

We get (F-FI0) directly from IH

## 11. FC-FE:

$$\frac{\begin{array}{c} \Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_t \\ \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash_{pc} \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell \end{array}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s [] : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[]\bullet)))} \text{FE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s [] \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[]\bullet))) \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \stackrel{\hat{\beta}}{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s []) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[]\bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \stackrel{\hat{\beta}'}{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-FE0})$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we have

$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies$   
 $\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V'}^{\hat{\beta}'_1}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-FE1})$$

Since from (F-FE1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V'}^{\hat{\beta}'_1}$  therefore from Definition 3.39 we know that

$$\exists {}^t v_i. {}^t v_i = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\forall \alpha_g. (\ell_e, \tau)) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-FE1.1})$$

Therefore from Definition 3.39 we have

$${}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \Lambda \nu e'_t$$

$$\forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \ell'' \in \mathcal{L}, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in [\tau[\ell''/\alpha_g] \sigma]_E^{\hat{\beta}'_1} \quad (\text{F-FE1.2})$$

We instantiate with  ${}^s \theta'_1, \ell', n - j - 1, \hat{\beta}'$  we get  $({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in [\tau[\ell'/\alpha_g] \sigma]_E^{\hat{\beta}'_1}$

From Definition 3.40 we have

$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies$   
 $\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1.$

$$(n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V'}^{\hat{\beta}''}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  and from fg-FE we know that  $i = j + k + 1 < n$  therefore we know that  $k < n - j - 1$  s.t  $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1.$$

$$(n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V'}^{\hat{\beta}''} \quad (\text{F-FE1.3})$$

Let  $\tau[\ell'/\alpha] \sigma = \Lambda^{\ell_i}$ , since  $\tau[\ell'/\alpha] \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{\beta}''}^V$$

Therefore from Definition 3.39 we know that

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}^t v_{2i}) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{\beta}''}^V \quad (\text{F-FE1.4})$$

In order to prove (F-FE0) we choose  $H'_t$  as  $H'_{t_2}$  and  ${}^t v$  as  $\text{Lb}^t v_{2i}$ . We need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[\square\square\square\bullet]))) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}):$$

From Lemma 3.48 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[\square\square\square\bullet]))) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i})$$

From SLIO\*-Sem-bind it suffices to prove that

$$\bullet (H_t, e_t \delta^t) \Downarrow^f (H'_{t_{11}}, {}^t v_{t_{11}}):$$

From (F-FE1) we know that  $H'_{t_{11}} = H'_{t_1}$  and  ${}^t v_{t_{11}} = {}^t v_1$

$$\bullet (H'_{t_1}, \text{bind}(\text{unlabel } a, b.b[\square\square\square\bullet]) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}):$$

Again from SLIO\*-Sem-bind it suffices to prove that

$$- (H'_{t_1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t_{12}}, {}^t v_{t_{12}}):$$

From (F-FE1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from SLIO\*-Sem-unlabel we have  $H'_{t_{12}} = H'_{t_1}$  and  ${}^t v_{t_{12}} = {}^t v_i$

$$- (b[\square\square\square\bullet] [{}^t v_i/b] \delta^t \Downarrow^f {}^t v_{t_{13}}):$$

From (F-FE1.2) we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_i = \Lambda\Lambda\Lambda\nu e'_t$

Therefore from SLIO\*-Sem-FE and SLIO\*-Sem-CE we know that  ${}^t v_{t_{13}} = e'_t$

$$- (H'_{t_1}, e'_t \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}))$$

From (F-FE1.3) and (F-FE1.4) we get the desired.

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{\beta}'}^V:$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}''$ . From fg-FE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v'_2$ ,  ${}^t v = {}^t v'_2$ ,  $H'_s = H'_{s_2}$  and  $H'_t = H'_{t_2}$ .

Therefore from (F-FE1.3) we get the  $(n - i, H'_{s_2}, H'_{t_2}) \triangleright^{\hat{\beta}''} {}^s\theta'_2$

$$\underline{\text{To prove:}} ({}^s\theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{\beta}''}^V$$

We get this directly from (F-FE1.3)

## 12. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{CI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^V$

To prove:  $({}^s\theta, n, \nu e \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{\hat{\beta}'}^V$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from SLIO\*-Sem-ret we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ : Given

(b)  $({}^s \theta, n, \nu e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V^{\hat{\beta}}$ :

From Definition 3.39 it suffices to prove that

$$({}^s \theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V^{\hat{\beta}}$$

Again from Definition 3.39 it suffices to prove that

$$\mathcal{L} \models c \sigma \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

This further means that given  $\mathcal{L} \models c \sigma$  and  ${}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'} \quad (\text{F-CI0})$$

$$\underline{\text{IH}}: ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

### 13. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b.\square\square\bullet)))} \text{CE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, e_s \bullet \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b.\square\square\bullet)))) \delta^t \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, e_s \square) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b.\square\square\bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$\exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-CE0})$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [(c \xrightarrow{\ell_\varepsilon} \tau)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we have

$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(c \xrightarrow{\ell_\varepsilon} \tau)^\ell \sigma]_V^{\hat{\beta}'_1}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(c \xrightarrow{\ell_\varepsilon} \tau)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-CE1})$$

Since from (F-CE1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(c \xrightarrow{\ell_\varepsilon} \tau)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 3.39 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(c \xrightarrow{\ell_\varepsilon} \tau) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-CE1.1})$$

Therefore from Definition 3.39 we have

$${}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \nu e'_t$$

$$\mathcal{L} \models c \sigma \implies \forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}'_1} \quad (\text{F-CE1.2})$$

Since we know that  $\mathcal{L} \models c \sigma$ , we instantiate with  ${}^s \theta'_1, n - j - 1, \hat{\beta}'$  to get

$$({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

From Definition 3.40 we have

$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$

$$(n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}''}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$  and since from fg-CE we know that  $i = j + k + 1 < n$  therefore we know that  $k < n - j - 1$  s.t  $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-CE1.3})$$

Let  $\tau \sigma = A^{\ell_i}$ , since  $\tau \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}''}$$

Therefore from Definition 3.39 we know that

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}^t v_{2i}) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-CE1.4})$$

In order to prove (F-CE0) we choose  $H'_t$  as  $H'_{t_2}$  and  ${}^t v$  as  $\text{Lb}^t v_{2i}$ . We need to prove

$$(a) \ (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[\ ][\ ]\bullet)))) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}):$$

From Lemma 3.48 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[\ ][\ ]\bullet)))) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i})$$

From SLIO\*-Sem-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t_{11}}, {}^t v_{t_{11}}):$

From (F-CE1) we know that  $H'_{t_{11}} = H'_{t_1}$  and  ${}^t v_{t_{11}} = {}^t v_1$

- $(H'_{t_1}, \text{bind}(\text{unlabel } a, b.b[\ ][\ ]\bullet)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}):$

Again from SLIO\*-Sem-bind it suffices to prove that

- $(H'_{t_1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t_{12}}, {}^t v_{t_{12}}):$

From (F-CE1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from SLIO\*-Sem-unlabel we have  $H'_{t_{12}} = H'_{t_1}$  and  ${}^t v_{t_{12}} = {}^t v_i$

- $(b[\ ][\ ]\bullet)[{}^t v_i/b] \delta^t \Downarrow^f {}^t v_{t_{13}}:$

From (F-CE1.2) we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_i = \Lambda \Lambda \nu e'_t$

Therefore from SLIO\*-Sem-FE and SLIO\*-Sem-CE we know that  ${}^t v_{t_{13}} = e'_t$

- $(H'_{t_1}, e'_t \Downarrow^f (H'_{t_2}, \text{Lb}^t v_{2i}))$

We get the desired from From (F-CE1.3) and (F-CE1.4)

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}''$ . From fg-CE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v'_2$ ,  ${}^t v = {}^t v'_2$ ,  $H'_s = H'_{s_2}$  and  $H'_t = H'_{t_2}$ .

Therefore from (F-CE1.3) we get the  $(n - i, H'_{s_2}, H'_{t_2}) \triangleright^{\hat{\beta}''} {}^s\theta'_2$

To prove:  $({}^s\theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau \sigma]_V^{\hat{\beta}''}$

From (F-CE1.3) we know that  $({}^s\theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}''}$

14. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb} b)))} \text{ref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{new } (e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b)) \delta^t) \delta^t) \in [(\text{ref } \tau)^\perp \sigma]_E^{\hat{\beta}}$

This means from Definition 3.40 we have

$$\begin{aligned}
& \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}
\end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ .

And we are required to prove

$$\begin{aligned}
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'} \quad (\text{F-R0})
\end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we have

$$\begin{aligned}
& \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\
& \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} s\theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1}
\end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore we know that  $\exists j < n$  s.t.  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

Therefore we have

$$\begin{aligned}
& \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} s\theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-R1})
\end{aligned}$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_1 \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^t v = \text{Lb}(a_t)$ ,  ${}^s \theta'$  as  ${}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From SLIO\*-Sem-bind it suffices to prove that

$$\bullet (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-R1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$

$$\bullet (H'_1, \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2}):$$

From SLIO\*-Sem-bind it suffices to prove that

$$i. (H'_1, \text{new } (a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2}):$$

From SLIO\*-Sem-new we know that  $H'_{t2} = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_{t2} = a_t$

$$ii. (H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_t):$$

From SLIO\*-Sem-ret we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_t = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}$$

From (F-R1) we know that  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} s\theta'_1$  and since  $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$ ,  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$

Therefore from Definition 3.41 and Lemma 3.47 we get  $(n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'}$

To prove:  $(s\theta', n - i, s v, t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}$

Since we know that  $s v = a_s$  and  $t v = \text{Lb } a_t$  therefore we need to prove

$(s\theta', n - i, a_s, \text{Lb}(a_t)) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}$

From Definition 3.39 it suffices to prove that

$(s\theta', n - i, a_s, a_t) \in [(\text{ref } \tau) \sigma]_V^{\hat{\beta}'}$

Again from Definition 3.39 it suffices to prove that

$s\theta'(a_s) = \tau \sigma \wedge (a_s, a_t) \in \hat{\beta}'$

We get this by construction

15. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\text{ref } \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{deref}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge (s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $(s\theta, n, !e \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \in [\tau' \sigma]_E^{\hat{\beta}}$

This means from Definition 3.40 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, s v. (H_s, !e_s) \Downarrow_i (H'_s, s v) \implies \\ & \exists H'_t, t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, t v) \wedge \exists s\theta' \sqsupseteq s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge (s\theta', n - i, s v, t v) \in [\tau' \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\gamma, \hat{\beta}} s\theta$ . Also given some  $i < n, s v$  s.t  $(H_s, !e_s) \Downarrow_i (H'_s, s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, t v) \wedge \exists s\theta' \sqsupseteq s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge (s\theta', n - i, s v, t v) \in [\tau' \sigma]_V^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

**IH:**

$$(s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \tau)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 3.40 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^{s\theta} \wedge \forall j < n, s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, s v_1) \implies \\ & \exists H'_{t1}, t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, t v_1) \wedge \exists s\theta'_1 \sqsupseteq s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{s\theta'_1} \wedge (s\theta'_1, n - j, s v_1, t v_1) \in [(\text{ref } \tau)^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, !e_s) \Downarrow_i (H'_s, s v)$  therefore  $\exists j < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, s v)$

Therefore we have



$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$

From Definition 3.39 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 3.39 we know that  ${}^s v_1 = a_s$  and  ${}^t v_i = a_t$

$${}^s \theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Let  $\tau' \sigma = A^{\ell_i}$ , since  $\tau' \sigma \searrow \ell \sigma$  therefore  $\ell \sigma \sqsubseteq \ell_i$  and

Let  $v_g = H_t(a_t)$  therefore from Definition 5.27 we have

$$({}^s \theta, n - 1, H_s(a_s), \text{Lb} v_{gi}) \in [\tau']_{\hat{V}}^{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $H'_{t1}(a_t) = v_g = \text{Lb} v_{gi}$

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb} v_{gi}):$$

From Lemma 3.48 it suffices to prove that

$$(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb} v_{gi})$$

From SLIO\*-Sem-bind it suffices to prove

$$i. (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$

$$ii. (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t2}):$$

From SLIO\*-Sem-bind it suffices to prove that

$$A. (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from SLIO\*-Sem-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$

$$B. (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_t, \text{Lb} v_{gi}):$$

Since from (F-DR1.2) we know that  ${}^t v_i = a_t$  therefore from SLIO\*-Sem-deref we know that  $H'_t = H'_{t1}$  and  ${}^t v = H'_{t1}(a_t) = v_g = \text{Lb} v_{gi}$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'_1}:$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$

Therefore from (F-DR1) we get  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$  and since  $i = j + 1$  therefore

from Lemma 3.47 we get  $(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$

Since from (F-DR1.2) we know that  $(a_s, a_t) \in \hat{\beta}'_1$  and  ${}^s \theta'_1(a_s) = \tau$ . Also from (F-

DR1) we have  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$ . Therefore from Definition 3.40 we have  $(n - j - 1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s \theta'_1(a_s)]_{\hat{V}}^{\hat{\beta}'_1}$

Since  $i = j + 1$ ,  ${}^s \theta'_1(a_s) = \tau \sigma$ ,  $H'_{s1}(a_s) = {}^s v$  and  $H'_{t1}(a_t) = {}^t v$

Therefore we get

$$({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'_1}$$

Finally from Lemma 3.50 we get

$$({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'}$$

16. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} := e_{s2} : \mathbf{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{assign}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \in \llbracket \mathbf{unit} \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 3.40 we are required to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}'} \quad (\text{F-AN0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 3.40 we are required to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t.  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$  therefore from Definition 3.39 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 3.39 this further means that

$${}^s \theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 3.40 we are required to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright_{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright_{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < n - j$  s.t.  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright_{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}'_2} \wedge \\ & (\text{F-AN2}) \end{aligned}$$

In order to prove (F-AN0) we choose  $H'_t$  as  $H'_{t2}[a_t \mapsto {}^s v_2]$ ,  ${}^t v$  as ()

We need to prove

$$(a) \ (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From SLIO\*-Sem-bind it suffices to prove that

$$- (H_t, (\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From SLIO\*-Sem-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$$

where  ${}^t v_T = \text{Lb}({}^t v_{Ti})$

From SLIO\*-Sem-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$

From (F-AN1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12}):$

From SLIO\*-Sem-bind it suffices to prove

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13}):$

From (F-AN2) we know that  $H'_{t13} = H'_{t2}$  and  ${}^t v_{t13} = {}^t v_2$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From SLIO\*-Sem-bind it suffices to prove that

- \*  $(H'_{t1}, \text{unlabel } a [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{V'}^{\hat{\beta}'_1}$$

Therefore from SLIO\*-Sem-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i = a_t$

\*  $(H'_{t1}, (c := b)[^t v_1/a][^t v_2/b][^t v_i/c] \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$ :

From SLIO\*-Sem-assign we know that  $H'_T = H'_{t1}[a_t \mapsto {}^t v_2]$  and  ${}^t v_{Ti} = ()$

Since  ${}^t v_{t12} = {}^t v_{Ti} = ()$  therefore  ${}^t v_T = \text{Lb}()$

-  $(H'_T, \text{ret}()[^t v_T/d]) \delta^t \Downarrow^f (H'_t, ())$ :

From SLIO\*-Sem-ret and SLIO\*-Sem-val

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$

In order to prove  $(n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$  it suffices to prove

- $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 3.41 we get  $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 3.41 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s \theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$ :

$\forall (a_1, a_2) \in \hat{\beta}'_2$ .

- $a_1 = a_s$  and  $a_1 = a_t$ :

Since from (F-AN2) we know that  $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 3.37 we know that  ${}^s \theta'_2(a_1) = \tau \sigma$

Therefore from Lemma 3.45 we get

$$({}^s \theta'_2, n - i - 1, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$  and  $a_1 \neq a_t$ :

From (F-AN2) since we know that  $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 3.41 we get

$$({}^s \theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1) \sigma]_{V'}^{\hat{\beta}'_2}$$

Since  $i = j + k + 1$  therefore from Lemma 3.45 we get

$$({}^s \theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1) \sigma]_{V'}^{\hat{\beta}'_2}$$

- $a_1 = a_s$  and  $a_1 \neq a_t$ :

This case cannot arise

- $a_1 \neq a_s$  and  $a_1 = a_t$ :

This case cannot arise

And in order to prove  $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

Since we know that  ${}^s v = ()$  and  ${}^t v = ()$  therefore from Definition 3.39 we get  $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

□

**Lemma 3.50** (FG  $\rightsquigarrow$  SLIO\*: Semantic Subtyping lemma). *The following holds:*

$$\forall \Sigma, \Psi, \sigma, \mathcal{L}, \hat{\beta}.$$

1.  $\forall A, A'$ .

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (A \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (A' \sigma) \rrbracket_V^{\hat{\beta}}$$

2.  $\forall \tau, \tau'$ .

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_V^{\hat{\beta}}$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

$$\text{IH1: } \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{It suffices to prove: } \forall ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda\Lambda(\nu(\lambda x.e_t))) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}.$$

$$({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda\Lambda(\nu(\lambda x.e_t))) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, m$  and  $\lambda x.e_s, \Lambda\Lambda\Lambda(\nu(\lambda x.e_t))$  s.t

$$({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda\Lambda(\nu(\lambda x.e_t))) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Therefore from Definition 3.39 we are given:

$$\begin{aligned} \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-L0}) \end{aligned}$$

$$\text{And it suffices to prove: } ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda\Lambda(\nu(\lambda x.e_t))) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Again from Definition 3.39, it suffices to prove:

$$\begin{aligned} \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L1}) \end{aligned}$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$  s.t  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}'_2}$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with  ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$ . Since we have  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}'_2}$  therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2 \sigma)]_E^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2}$$

## 2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

$$\text{To prove: } [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_1 \sigma)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{IH2: } [(\tau_2 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, n$  and  ${}^s v_1, {}^s v_2, {}^t v_1, {}^t v_2$  s.t

$$({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 3.39 we are given:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 3.39, it suffices to prove:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that  $({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from IH1 we have

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

Similarly since we have  $({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$  from (S-P0) therefore from IH2 we have

$$({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}}$$

### 3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \subseteq [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

IH1:  $[(\tau_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_1 \sigma)]_V^{\hat{\beta}}$  (Statement 2(a))

IH2:  $[(\tau_2 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:  $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^sv = \text{inl } {}^sv_i$  and  ${}^tv = \text{inl } {}^tv_i$ :

From Definition 3.39 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b)  ${}^sv = \text{inr } {}^sv_i$  and  ${}^tv = \text{inr } {}^tv_i$ :

Symmetric reasoning as in the previous case

### 4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove:  $[(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s\theta, n, \Lambda e_s, \Lambda \Lambda \Lambda(\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}. ({}^s\theta, n, \Lambda e_s, \Lambda \Lambda \Lambda(\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, \Lambda e_s, \Lambda \Lambda \Lambda(\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 3.39 we have:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-F0})$$

And we need to prove

$$({}^s\theta, n, \Lambda e_s, \Lambda \Lambda \Lambda(\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

Again from Definition 3.39 it means we need to prove

$$\forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha] \sigma]_{E}^{\hat{\beta}'_2}$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s\theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha]]_{E}^{\hat{\beta}'_2} \quad (\text{S-F1})$$

Instantiating (S-F0) with  ${}^s\theta'_2, k, \ell'', \hat{\beta}'_2$  and we get

$$({}^s\theta'_2, k, e_s, e_t) \in [\tau_1[\ell''/\alpha]]_{E}^{\hat{\beta}'_2}$$

$$\text{IH: } [(\tau_1 \sigma \cup \{\alpha \mapsto \ell''\})]_{E}^{\hat{\beta}'_2} \subseteq [(\tau_2 \sigma \cup \{\alpha \mapsto \ell''\})]_{E}^{\hat{\beta}'_2} \quad (\text{Statement 2(b)})$$

Therefore from IH we get the desired

#### 5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell'_e} \tau_1 <: c_2 \xrightarrow{\ell_e} \tau_2} \text{FGsub-constraint}$$

$$\text{To prove: } [((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^{\hat{\beta}} \subseteq [((c_2 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

It suffices to prove:

$$\forall ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in [((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in [((c_2 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{This means that given: } ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in [((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 3.39 we are given:

$$\mathcal{L} \models c_1 \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s, e_t) \in [\tau_1 \sigma]_{E}^{\hat{\beta}'_1} \quad (\text{S-C0})$$

And it suffices to prove:

$$({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in [((c_1 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 3.39 it means that we need to prove:

$$\mathcal{L} \models c_2 \sigma \implies \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, e_s, e_t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'_2}$$

This means that given that  $\mathcal{L} \models c_2 \sigma$  and  ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s\theta'_2, k, e_s, e_t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'_2} \quad (\text{S-C1})$$

Instantiating (S-C0) with  ${}^s\theta'_2, k, \hat{\beta}'_2$  we get  $({}^s\theta'_2, k, e_s, e_t) \in [\tau_1 \sigma]_{E}^{\hat{\beta}'_2}$

$$\text{IH: } [(\tau_1 \sigma)]_{E}^{\hat{\beta}'_2} \subseteq [(\tau_2 \sigma)]_{E}^{\hat{\beta}'_2} \quad (\text{Statement 2(b)})$$

Finally from IH we get  $({}^s\theta'_2, k, e_s, e_t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'_2}$



6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

We get this directly from Definition 3.39

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{b}) \sigma) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 3.39

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{unit}) \sigma) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 3.39

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell' \sqsubseteq \ell'' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^{\ell'} <: \mathbf{A}^{\ell''}} \text{FGsub-label}$$

To prove:  $\llbracket ((\mathbf{A}^{\ell'}) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{A}^{\ell''}) \sigma) \rrbracket_V^{\hat{\beta}}$

This means from Definition 3.39 we need to prove

$$\forall ({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in \llbracket \mathbf{A}^{\ell'} \sigma \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in \llbracket \mathbf{A}^{\ell''} \sigma \rrbracket_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in \llbracket \mathbf{A}^{\ell'} \sigma \rrbracket_V^{\hat{\beta}}$

From Definition 3.39 it further means that we are given

$$({}^s\theta, n, {}^sv, {}^tv_i) \in \llbracket \mathbf{A} \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-LB0})$$

And we need to prove

$$({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in \llbracket \mathbf{A}^{\ell''} \sigma \rrbracket_V^{\hat{\beta}}$$

Again from Definition 3.39 it suffices to prove that

$$({}^s\theta, n, {}^sv, {}^tv_i) \in \llbracket \mathbf{A}' \sigma \rrbracket_V^{\hat{\beta}}$$

Since  $\ell' \sqsubseteq \ell''$  and  $A' <: A''$  therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given:  $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove:  $\llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$

This means we need to prove that

$$\forall ({}^s\theta, n, e_s, e_t) \in \llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}}. ({}^s\theta, n, e_s, e_t) \in \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$$

This means given  $({}^s\theta, n, e_s, e_t) \in \llbracket (\tau \sigma) \rrbracket_E^{\hat{\beta}}$

This means from Definition 3.40 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{S-E0})$$

And it suffices to prove that  $({}^s\theta, n, e_s, e_t) \in \llbracket (\tau' \sigma) \rrbracket_E^{\hat{\beta}}$

Again from Definition 3.40 it means we need to prove

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'_1}$$

This means that given some  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} {}^s\theta$ . Also given some  $j < n, {}^s v_1$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

And we need to prove

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{S-E1})$$

Instantiating (S-E0) with  $H_{s1}, H_{t1}$  and with  $j, {}^s v_1$ . Then we get

$$\exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'_1}$$

Since we have  $\tau <: \tau'$ . Therefore from IH (Statement 2(a)) we get

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'_1}$$

□

**Theorem 3.51** (FG  $\rightsquigarrow$  SLIO\*: Deriving FG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, n_1, n_2, H'_{s1}, H'_{s2}, pc.$

Let  $\text{bool} = (\text{unit} + \text{unit})$

$$\emptyset, \emptyset, x : \text{bool}^\top \vdash_{pc} e_s : \text{bool}^\perp \wedge$$

$$\emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_1 : \text{bool}^\top \wedge \emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_2 : \text{bool}^\top \wedge$$

$$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^s v'_1) \wedge$$

$$(\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^s v'_2) \wedge$$

$\implies$

$${}^s v'_1 = {}^s v'_2$$

*Proof.* From the FG to CG translation we know that  $\exists e_t$  s.t

$$\emptyset, \emptyset, x : \text{bool}^\top \vdash e_s : \text{bool}^\perp \rightsquigarrow e_t$$

Similarly we also know that  $\exists t v_1, t v_2$  s.t

$$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow t v_1 \text{ and } \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{bool}^\top \rightsquigarrow t v_2 \quad (\text{NI-0})$$

From type preservation theorem (choosing  $\alpha = \gamma = \bar{\beta} = \perp$ ) we know that

$$\emptyset, \emptyset, x : \text{Labeled} \top \text{ bool} \vdash e_t : \text{SLIO} \perp \perp \text{ Labeled} \perp \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash t v_1 : \text{SLIO} \perp \perp \text{ Labeled} \top \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash t v_2 : \text{SLIO} \perp \perp \text{ Labeled} \top \text{ bool} \quad (\text{NI-1})$$

Since we have  $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow t v_1$

And since  ${}^s v_1$  and  $t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 3.49 we have (we choose  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, t v_1) \in \lfloor \text{bool}^\top \rfloor_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 3.40 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^\emptyset \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, t v_{11}. (H_t, t v_1) \Downarrow^f (H'_t, t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, t v_{11}) \in \lfloor \text{bool}^\top \sigma \rfloor_V^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, t v_{11}. (H_t, t v_1) \Downarrow^f (H'_t, t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, t v_{11}) \in \lfloor \text{bool}^\top \sigma \rfloor_V^{\hat{\beta}'} \quad (\text{NI-2.1})$$

From Definition 3.39 we know that

$$t v_{11} = \text{Lb}(t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, t v_{i11}) \in \lfloor (\text{unit} + \text{unit}) \sigma \rfloor_V^{\hat{\beta}'}$$

Again from Definition 3.39 we know that

Either a)  ${}^s v_1 = \text{inl}()$  and  $t v_{i11} = \text{inl}()$  or b)  ${}^s v_1 = \text{inr}()$  and  $t v_{i11} = \text{inr}()$

But in either case we have that  $\emptyset, \emptyset, \emptyset \vdash t v_{i11} : (\text{unit} + \text{unit}) \quad (\text{NI-2.2})$

As a result we have  $\emptyset, \emptyset, \emptyset \vdash t v_{11} : \text{Labeled} \top (\text{unit} + \text{unit}) \quad (\text{NI-2.3})$

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash t v_{i11} : (\text{unit} + \text{unit})} \quad (\text{NI-2.2})}{\emptyset, \emptyset, \emptyset \vdash \text{Lb}(t v_{i11}) : \text{Labeled} \top (\text{unit} + \text{unit})}$$

From Definition 3.44 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto t v_{11})) \in \lfloor x \mapsto \text{bool}^\top \rfloor_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 3.49 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[t v_{11}/x]) \in \lfloor \text{bool}^\perp \rfloor_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 3.40 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v''_1. (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v''_1) \implies \\ \exists H'_{t1}, t v''_1. (H_t, e_t[t v_{11}/x]) \Downarrow^f (H'_{t1}, t v''_1) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v''_1, t v''_1) \in \lfloor \text{bool}^\perp \sigma \rfloor_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_1, {}^s v'_1$  we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}' \\ (n - n_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \quad (\text{NI-2.5}) \end{aligned}$$

Since we have  $({}^s \theta', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''}$  therefore from Definition 3.39 we have

$$\exists {}^t v_{i1}. {}^t v'' = \mathbf{Lb}({}^t v_{i1}) \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [\mathbf{bool} \sigma]_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$  therefore from Definition 3.39 two cases arise

- ${}^s v'_1 = \mathbf{inl} {}^s v_{i11}$  and  ${}^t v_{i1} = \mathbf{inl} {}^t v_{i11}$ :

From Definition 3.39 we have

$$({}^s \theta', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v'_1 = \mathbf{inr} {}^s v_{i11}$  and  ${}^t v_{i1} = \mathbf{inr} {}^t v_{i11}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v'_1 = {}^t v_{i1}$

$$\text{Similarly with other substitution we have } (\emptyset, n, {}^s v_2, {}^t v_2) \in [\mathbf{bool}^\top]_E^\emptyset \quad (\text{NI-3})$$

Therefore from Definition 3.40 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^\emptyset \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'} \quad (\text{NI-3.1})$$

From Definition 3.39 we know that

$${}^t v_2 = \mathbf{Lb}({}^t v_{22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [(\mathbf{unit} + \mathbf{unit}) \sigma]_V^{\hat{\beta}'}$$

Again from Definition 3.39 we know that

Either a)  ${}^s v_2 = \mathbf{inl}()$  and  ${}^t v_{22} = \mathbf{inl}()$  or b)  ${}^s v_2 = \mathbf{inr}()$  and  ${}^t v_{22} = \mathbf{inr}()$

But in either case we have that  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-3.2})$

As a result we have  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : \text{Labeled } \top (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-3.3})$

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : (\mathbf{unit} + \mathbf{unit})} \quad (\text{NI-3.2})}{\emptyset, \emptyset, \emptyset \vdash \mathbf{Lb}({}^t v_{22}) : \text{Labeled } \top (\mathbf{unit} + \mathbf{unit})}$$

From Definition 3.44 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \mathbf{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 3.49 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\mathbf{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-3.4})$$

From Definition 3.40 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v_2''.(H_s, e_s[{}^s v_2/x]) \Downarrow_i (H'_{s2}, {}^s v_2'') \implies \\ \exists H'_{t2}, {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_2'', {}^t v_2'') \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_2, {}^s v_2'$  we get

$$\exists H'_{t2}, {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - n_1, H'_s, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_2'', {}^t v_2'') \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \quad (\text{NI-3.5})$$

Since we have  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_2'') \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''}$  therefore from Definition 3.39 we have

$$\exists {}^t v_{i2}. {}^t v_2'' = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [\mathbf{bool} \sigma]_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$  therefore from Definition 3.39 two cases arise

- ${}^s v_2' = \text{inl } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inl } {}^t v_{i22}$ :

From Definition 3.39 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v_2' = \text{inr } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inr } {}^t v_{i22}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v_2' = {}^t v_{i2}$

We know that  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{11} : \text{Labeled } \top \text{ bool}$  (NI-2.3)

Also we have  $\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : \text{Labeled } \top \text{ bool}$  (NI-3.3)

Let  $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that  $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{S}\text{LIO} \perp \perp \text{ bool}$  by giving a typing derivation P2:

$$\frac{\frac{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool} \quad \text{SLIO}^*\text{-var}}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{S}\text{LIO} \perp \perp \text{ bool} \quad \text{SLIO}^*\text{-unlabel}}}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{S}\text{LIO} \perp \perp \text{ bool}}$$

P1:

$$\frac{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{S}\text{LIO} \perp \perp \text{Labeled } \perp \text{ bool}}{\text{From (NI-1)}}$$

Main derivation:

$$\frac{\frac{\text{P1} \quad \text{P2}}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{S}\text{LIO} \perp \perp \text{ bool}}}$$

Say  $e_t[{}^t v_{11}/x]$  reduces in  $n_{t1}$  steps in (NI-2.5) and  $e_t[{}^t v_{22}/x]$  reduces in  $n_{t2}$  steps in (NI-3.5)

We instantiate Theorem 2.28 with  $e_T, {}^t v_{11}, {}^t v_{22}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$  and from (NI-2.5) and (NI-3.5) we have  ${}^t v_{i1} = {}^t v_{i2}$  and thus  ${}^s v_1' = {}^s v_2'$  □

## 4 New coarse-grained IFC enforcement (CG)

### 4.1 CG type system

**Term, type, constraint syntax:**

Expressions	$e ::=$	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid$ $\text{new } e \mid !e \mid e := e \mid () \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid$ $\text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::=$	$\perp \mid \top \mid l \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::=$	$\mathbf{b} \mid \text{unit} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{Labeled } \ell \tau \mid \mathbb{C} \ell_1 \ell_2 \tau \mid \forall \alpha. \tau \mid$ $c \Rightarrow \tau$

**Type system:**  $\boxed{\Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types  $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau, \text{unit}$  are included.)

$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{CG-label}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau} \text{CG-unlabel}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell \ell' \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell \perp (\text{Labeled } \ell' \tau)} \text{CG-toLabeled}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ell \ell' \tau} \text{CG-ret}$
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau'}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_4 \quad \Sigma; \Psi \vdash \ell_4 \sqsubseteq \ell'} \text{CG-bind}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau} \text{CG-sub}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)} \text{CG-ref}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell' \tau}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell' \tau)} \text{CG-deref}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}} \text{CG-assign}$	
$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \Lambda e : \forall \alpha. \tau} \text{CG-FI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha]} \text{CG-FE}$
$\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \nu e : c \Rightarrow \tau} \text{CG-CI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau} \text{CG-CE}$

Figure 10: Type system of CG.

### 4.2 CG semantics

Judgement:  $e \Downarrow_i v$  and  $(H, e) \Downarrow_i^f (H', v)$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \tau <: \tau} \text{CGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{CGsub-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{CGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{CGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{CGsub-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'} \text{CGsub-monad} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{CGsub-forall} \qquad \frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2} \text{CGsub-constraint}
\end{array}$$

Figure 11: CG subtyping

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ } WF} \text{CG-wff-base} \qquad \frac{}{\Sigma; \Psi \vdash \text{unit} \text{ } WF} \text{CG-wff-unit} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \rightarrow \tau_2) \text{ } WF} \text{CG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \times \tau_2) \text{ } WF} \text{CG-wff-times} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 + \tau_2) \text{ } WF} \text{CG-wff-sum} \\
\\
\frac{\text{FV}(\ell) = \emptyset \quad \text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \ell \tau) \text{ } WF} \text{CG-wff-ref} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\forall \alpha. \tau) \text{ } WF} \text{CG-wff-forall} \\
\\
\frac{\Sigma; \Psi, c \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (c \implies \tau) \text{ } WF} \text{CG-wff-constraint} \qquad \frac{\Sigma; \Psi \vdash \tau \text{ } WF \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi \vdash (\text{Labeled } \ell \tau) \text{ } WF} \text{CG-wff-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ } WF \quad \text{FV}(\ell_i) \in \Sigma \quad \text{FV}(\ell_o) \in \Sigma}{\Sigma; \Psi \vdash (\text{SLIO } \ell_i \ell_o \tau) \text{ } WF} \text{CG-wff-monad}
\end{array}$$

Figure 12: Well-formedness relation for CG

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x.e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 e_2 \Downarrow_{i+j+k+1} v_3} \text{cg-app} \qquad \frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{cg-prod} \\
\\
\frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{cg-fst} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{cg-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{cg-inl} \\
\\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{cg-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{cg-case1} \\
\\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{cg-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{cg-Lb} \\
\\
\frac{e \Downarrow_i \Lambda e_i \quad e_i \Downarrow_j v}{e[] \Downarrow_{i+j+1} v} \text{SLIO*}-\text{Sem-FE} \qquad \frac{e \Downarrow_i \nu e_i \quad e_i \Downarrow_j v}{e\bullet \Downarrow_{i+j+1} v} \text{SLIO*}-\text{Sem-CE} \\
\\
\frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-ret} \\
\\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{cg-bind} \\
\\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-unlabel} \qquad \frac{e \Downarrow_i v \quad (H, v) \Downarrow_j^f (H', v')}{(H, \text{toLabeled}(e)) \Downarrow_{i+j+1}^f (H', \text{Lb}(v'))} \text{cg-toLabeled} \\
\\
\frac{e \Downarrow_i \text{Lb}v \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1}^f (H[a \mapsto \text{Lb}v], a)} \text{cg-ref} \qquad \frac{e \Downarrow_i a}{(H, !e) \Downarrow_{i+1}^f (H, H(a))} \text{cg-deref} \\
\\
\frac{e_1 \Downarrow_i a \quad e_2 \Downarrow_j \text{Lb}v}{(H, e_1 := e_2) \Downarrow_{i+j+1}^f (H[a \mapsto \text{Lb}v], ())} \text{cg-assign} \\
\\
\frac{e \in \{x, \lambda y.-, \Lambda, \nu, \text{ret}-, \text{bind}(-, -), \text{unlabel}(-), \text{toLabeled}(-), \text{new } (-), !-, - := -\}}{e \Downarrow_0 e} \text{cg-val}
\end{array}$$

Figure 13: CG semantics



### 4.3 Model for CG

$W : ((Loc \mapsto Type) \times (Loc \mapsto Type) \times (Loc \leftrightarrow Loc))$

**Definition 4.1** ( $\theta_2$  extends  $\theta_1$ ).  $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

**Definition 4.2** ( $W_2$  extends  $W_1$ ).  $W_1 \sqsubseteq W_2 \triangleq$

1.  $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2.  $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

**Definition 4.3** (Value Equivalence).

$$ValEq(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_{\mathcal{V}} \wedge (W.\theta_2, j, v_2) \in [\tau]_{\mathcal{V}} & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

**Definition 4.4** (Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \rightarrow \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\
[\forall \alpha. \tau]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, \ell' \in \mathcal{L}. \\
&\quad ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E\} \\
[ c \Rightarrow \tau ]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n. \\
&\quad \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E\} \\
[\text{ref } \ell \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \text{Labeled } \ell \tau\} \\
[\text{Labeled } \ell \tau]_V^A &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\
[\mathbb{C} \ell_1 \ell_2 \tau]_V^A &\triangleq \{(W, n, v_1, v_2) \mid \\
&\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\
&\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
&\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\
&\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\}
\end{aligned}$$

**Definition 4.5** (Binary expression relation).

$$[\tau]_E^A \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A\}$$

**Definition 4.6** (Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta' \supseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \forall \alpha. \tau \rrbracket_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta', j < m. \forall \ell' \in \mathcal{L}. (\theta', j, e) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E\} \\
\llbracket c \Rightarrow \tau \rrbracket_V &\triangleq \{(\theta, m, \nu e) \mid \mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < m. (\theta', j, e) \in \llbracket \tau \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v)) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbb{C} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \supseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \supseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

**Definition 4.7** (Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

**Definition 4.8** (Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

**Definition 4.9** (Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

**Definition 4.10** (Binary substitution).  $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

**Definition 4.11** (Unary substitution).  $\delta : \text{Var} \mapsto \text{Val}$

**Definition 4.12** (Unary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

**Definition 4.13** (Binary interpretation of  $\Gamma$ ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

#### 4.4 Soundness proof for CG

**Lemma 4.14** (Binary value relation subsumes unary value relation).  $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

*Proof.* Proof by induction on  $\tau$

1. Case **b, unit**:

From Definition 4.6

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$  (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$  (P02)

From Definition 4.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$  (P1)

IH1a:  $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$  and

IH1b:  $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a:  $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$  and

IH2b:  $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some  $m$  we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some  $m$  we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given  $m$  from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$  and  $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 4.6, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given  $m$  from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$  and  $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 4.6, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl}(v_{i1})$  and  $v_2 = \text{inl}(v_{j1})$

Given:  $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$  (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$  (S02)

From Definition 4.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

$$\text{IH1: } \forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \text{ and}$$

$$\text{IH2: } \forall m_2. (W.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$$

From (S01) we know that given some  $m$  and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

Also from (S02) we know that given some  $m$  and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH1 with  $m$  from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$$

Therefore from Definition 4.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with  $m$  from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 4.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

$$(b) \ v_1 = \text{inr}(v_{i2}) \text{ and } v_2 = \text{inr}(v_{j2})$$

Symmetric reasoning as in the (a) case above

#### 4. Case $\tau_1 \rightarrow \tau_2$ :

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V^A$$

This means from Definition 4.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V \implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E) \quad (\text{L0}) \end{aligned}$$

To prove:

$$(a) \ \forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \rightarrow \tau_2]_V:$$

This means from Definition 4.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

This further means that we have some  $\theta', j$  and  $v$  s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

Instantiating  $\theta_l, i$  and  $v_c$  in the second conjunct of L0 with  $\theta', j$  and  $v$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

$$(b) \ \forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V:$$

Similar reasoning with  $e_2$

5. Case  $\forall\alpha.\tau$ :

Given:  $(W, n, \Lambda e_1, \Lambda e_2) \in [\forall\alpha.\tau]_V^A$

This means from Definition 4.4 we know that

$$\begin{aligned} &\forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E) \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E) \end{aligned} \quad (\text{F0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall\alpha.\tau]_V$ :

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$$

This further means that we are given some  $\theta', m'$  and  $\ell_u$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\ell_u \in \mathcal{L}$

And we need to prove:  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

Instantiating  $\theta_l, i$  and  $\ell''$  in the second conjunct of F0 with  $\theta', m'$  and  $\ell_u$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\ell_u \in \mathcal{L}$

Therefore we get  $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

(b)  $\forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall\alpha.\tau]_V$ :

Symmetric reasoning for  $e_2$

6. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, \nu e_1, \nu e_2) \in [c \Rightarrow \tau]_V^A$

This means from Definition 4.4 we know that

$$\begin{aligned} &\forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \\ &\wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E \end{aligned} \quad (\text{C0})$$

To prove:

(a)  $\forall m. (W.\theta_1, m, \nu e_1) \in [c \Rightarrow \tau]_V$ :

This means from Definition 4.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E$$

This further means that we are given some  $\theta'$  and  $m'$  s.t  $W.\theta_1 \sqsubseteq \theta', m' < m$  and  $\mathcal{L} \models c$

And we need to prove:  $(\theta', m', e_1) \in [\tau]_E$

Instantiating  $\theta_l, j$  in the second conjunct of C0 with  $\theta', m'$  respectively and since we know that  $W.\theta_1 \sqsubseteq \theta'$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta', m', e_1) \in [\tau]_E$

(b)  $\forall m. (W.\theta_2, m, \nu e_2) \in [c \Rightarrow \tau]_V$ :

Symmetric reasoning for  $e_2$

7. Case ref  $\ell \tau$ :

From Definition 4.4 and 4.6

8. Case Labeled  $\ell \tau$ :

Given  $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

From Definition 4.3 we know that

$(W, n, v_1, v_2) \in [\tau]_V^A$

Therefore from IH we get  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

Directly from Definition 4.3

9. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

This means from Definition 4.4 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 4.6 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case  $l = 1$

And given some  $m$  and  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with  $l = 1$  and the given  $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$  we get the desired.

Case  $l = 2$

Symmetric reasoning as in the previous case above

□

**Lemma 4.15** (Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$$(\theta, m, v) \in \lfloor \tau \rfloor_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \lfloor \tau \rfloor_V$$

*Proof.* Proof by induction on  $\tau$

1. case **b**, unit:

Directly from Definition 4.6

2. case  $\tau_1 \times \tau_2$ :

Given:  $(\theta, m, (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

To prove:  $(\theta', m', (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

This means from Definition 4.6 we know that

$$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V \wedge (\theta, m, v_2) \in \lfloor \tau_2 \rfloor_V$$

$$\text{IH1} : (\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$$

$$\text{IH2} : (\theta', m', v_2) \in \lfloor \tau_2 \rfloor_V$$

We get the desired from IH1, IH2 and Definition 4.6

3. case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v = \text{inl}(v_1)$ :

Given:  $(\theta, m, (\text{inl } v_1)) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

To prove:  $(\theta', m', \text{inl } v_1) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

This means from Definition 4.6 we know that

$$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V$$

$$\text{IH} : (\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$$

Therefore from IH and Definition 4.6 we get the desired

(b)  $v = \text{inr}(v_2)$

Symmetric case

4. case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(\theta, m, (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

To prove:  $(\theta', m', (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

This means from Definition 4.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E \quad (91)$$

Similarly from Definition 4.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V \implies (\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$$

This means that given some  $\theta''', k$  and  $v_1$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

And we are required to prove  $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$



Instantiating Equation 91 with  $\theta''', k$  and  $v_1$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get  $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

5. case ref  $\ell \tau$ :

From Definition 4.6 and Definition 4.1

6. case  $\forall\alpha.\tau$ :

Given:  $(\theta, m, (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

To prove:  $(\theta', m', (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

This means from Definition 4.6 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E \quad (92)$$

Similarly from Definition 4.6 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

Instantiating Equation 92 with  $\theta''', k$  and  $\ell_j$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\ell_j \in \mathcal{L}$

Therefore we get  $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

7. case  $c \Rightarrow \tau$ :

Given:  $(\theta, m, (\nu e_1)) \in [c \Rightarrow \tau]_V$

To prove:  $(\theta', m', (\nu e_1)) \in [c \Rightarrow \tau]_V$

This means from Definition 4.6 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \mathcal{L} \models c \implies (\theta'', j, e_1) \in [\tau]_E \quad (93)$$

Similarly from Definition 4.6 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \mathcal{L} \models c \implies (\theta''', k, e_1) \in [\tau]_E$$

This means that given some  $\theta''', k$  and  $\ell_j$  such that  $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove  $(\theta''', k, e_1) \in [\tau]_E$

Instantiating Equation 93 with  $\theta''', k$  and since we know that  $\theta' \sqsubseteq \theta'''$  and  $\theta \sqsubseteq \theta'$  therefore we have  $\theta \sqsubseteq \theta'''$ . Also, we know that  $k < m' < m$  and  $\mathcal{L} \models c$

Therefore we get  $(\theta''', k, e_1) \in [\tau]_E$

8. case Labeled  $\ell \tau$ :

Given:  $(\theta, m, (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

To prove:  $(\theta', m', (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

This means from Definition 4.6 we know that  $(\theta, m, v) \in \llbracket \tau \rrbracket_V$

IH:  $(\theta', m', v) \in \llbracket \tau \rrbracket_V$

Therefore from IH and Definition 4.6 we get the desired

9. case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(\theta, m, e) \in \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V$

To prove:  $(\theta', m', e) \in \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V$

This means from Definition 4.6 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 4.6 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta', k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t. } (k_1, H) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta', k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0),  $k$  with  $k_1$ ,  $\theta_e$  with  $\theta_{e1}$ ,  $H$  with  $H_1$  and  $j$  with  $j_1$ . We know that  $k_1 < m' < m$ ,  $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$ ,  $(k_1, H_1) \triangleright \theta_{e1}$ ,  $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$  and  $i_1 + j_1 < k_1$ . Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta', k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

**Lemma 4.16** (Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}}$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**, unit:

From Definition 4.4

2. Case  $\tau_1 \times \tau_2$ :

Given:  $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:  $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 4.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 :  $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 :  $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 4.4 we get the desired.

3. Case  $\tau_1 + \tau_2$ :

2 cases arise:

(a)  $v_1 = \text{inl } v_{i1}$  and  $v_2 = \text{inl } v_{i2}$ :

Given:  $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove:  $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 4.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH :  $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 4.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b)  $v_1 = \text{inr}(v_{i1})$  and  $v_2 = \text{inr}(v_{i2})$ :

Symmetric case

4. Case  $\tau_1 \rightarrow \tau_2$ :

Given:  $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

To prove:  $(\theta', n', (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

This means from Definition 4.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$   
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E)$  (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)$  (BM-A2)

Similarly from Definition 4.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', k < n'$  and  $v'_1, v'_2$  s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with  $W'', k$  and  $v'_1, v'_2$  we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $v'_c$  s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove:  $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with  $\theta'_l, k$  and  $v'_c$  we get

$$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$$

#### 5. Case ref $\ell \tau$ :

From Definition 4.4 and Definition 4.2

#### 6. Case $\forall \alpha. \tau$ :

Given:  $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall \alpha. \tau]_V^A$

To prove:  $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall \alpha. \tau]_V^A$

This means from Definition 4.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F2})$$

Similarly from Definition 4.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with  $W'', n''$  and  $\ell''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . Also since  $n'' < n'$  and  $n' < n$  therefore  $n'' < n$ . And finally since  $\ell'' \in \mathcal{L}$  therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_l \sqsupseteq W.\theta_1$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$$

(c)  $\forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\ell'' \in \mathcal{L}$

And we are required to prove:  $((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

Instantiating BM-F1 with  $\theta'_l, k$  and  $\ell''$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\ell'' \in \mathcal{L}$  therefore we get

$((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

7. Case  $c \Rightarrow \tau$ :

Given:  $(W, n, (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

To prove:  $(\theta', n', (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

This means from Definition 4.4 we know that the following holds

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in \lceil \tau \rceil_E^A$  (BM-C0)

$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in \lfloor \tau \rfloor_E$  (BM-C1)

$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in \lfloor \tau \rfloor_E$  (BM-C2)

Similarly from Definition 4.4 we know that we are required to prove

(a)  $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$ :

This means that we are given some  $W'' \sqsupseteq W', n'' < n'$  and  $\mathcal{L} \models c$

And we are required to prove:  $(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

Instantiating BM-C0 with  $W'', n''$ . And since  $W'' \sqsupseteq W'$  and  $W' \sqsupseteq W$  therefore  $W'' \sqsupseteq W$ . And since  $\mathcal{L} \models c$  therefore we get

$(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

(b)  $\forall \theta'_l \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_1, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_1$  and  $W' \sqsupseteq W$  therefore  $\theta'_1 \sqsupseteq W.\theta_1$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

(c)  $\forall \theta'_l \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$ :

This means that we are given some  $\theta'_l \sqsupseteq W'.\theta_2, k$  and  $\mathcal{L} \models c$

And we are required to prove:  $(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with  $\theta'_l, k$ . And since  $\theta'_l \sqsupseteq W'.\theta_2$  and  $W' \sqsupseteq W$  therefore  $\theta'_2 \sqsupseteq W.\theta_2$ . And since  $\mathcal{L} \models c$  therefore we get

$(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

8. Case Labeled  $\ell \tau$ :

Given:  $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

To prove:  $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

From Definition 4.4 2 cases arise:

(a)  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Therefore from IH we know that  $(W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Hence from Definition 4.4 we get  $(W', n', (\mathbf{Lb}v_1), (\mathbf{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

(b)  $\ell \not\sqsubseteq \mathcal{A}$ :

In this case we know that  $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$  and  $(W.\theta_2, m, v_2) \in [\tau]_V$

Since  $W.\theta_1 \sqsubseteq W'.\theta_1$  (from Definition 4.2). Therefore from Lemma 4.15 we know that

$\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since  $W.\theta_2 \sqsubseteq W'.\theta_2$  (from Definition 4.2). Therefore from Lemma 4.15 we know that

$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$

Finally from Definition 4.4 we get  $(W', n', (\mathbf{Lb}v_1), (\mathbf{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

9. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:  $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

To prove:  $(W', n', v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 4.4 we are given that

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{BM-M0}) \end{aligned}$$

Similarly from Definition 4.4 it suffices to prove that

$$\begin{aligned} & \text{(a) } \left( \forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right): \\ & \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\ & (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given  $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  and since we know that  $n' \leq n$  and  $W \sqsubseteq W'$  we get the desired

$$\begin{aligned} & \text{(b) } \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right): \end{aligned}$$

Similar reasoning as in the previous case but using Lemma 4.15

□

**Lemma 4.17** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n'$ .  
 $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$   
 To prove:  $(\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

From Definition 4.12 it is given that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

And again from Definition 4.12 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

- $dom(\Gamma) \subseteq dom(\delta)$ :  
 Given
- $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$ :  
 Since we know that  $\forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$  (given)  
 Therefore from Lemma 4.15 we get  
 $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

□

**Lemma 4.18** (Binary monotonicity for  $\Gamma$ ).  $\forall W, W', \delta, \Gamma, n, n'$ .  
 $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

*Proof.* Given:  $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W'$   
 To prove:  $(W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

From Definition 4.13 it is given that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

And again from Definition 4.12 we are required to prove that  
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

- $dom(\Gamma) \subseteq dom(\gamma)$ :  
 Given
- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$ :  
 Since we know that  $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$  (given)  
 Therefore from Lemma 4.16 we get  
 $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

□

**Lemma 4.19** (Unary monotonicity for  $H$ ).  $\forall \theta, H, n, n'$ .  
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

*Proof.* Given:  $(n, H) \triangleright \theta \wedge n' < n$

To prove:  $(n', H) \triangleright \theta$

From Definition 4.8 it is given that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

And again from Definition 4.12 we are required to prove that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$ :

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$ :

Since we know that  $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$  (given)

Therefore from Lemma 4.15 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

**Lemma 4.20** (Binary monotonicity for heaps).  $\forall W, H_1, H_2, n, n'$ .

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

*Proof.* Given:  $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove:  $(n', H_1, H_2) \triangleright W$

From Definition 4.9 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W.\hat{\beta}). & (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). & (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 4.9 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$ :

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$ :

Given

- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$  and  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

$$\forall (a_1, a_2) \in (W.\hat{\beta}).$$

–  $(W.\theta_1(a_1) = W.\theta_2(a_2))$ : Given

–  $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$ :

Given and from Lemma 4.16

- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$ :

Given



□

**Theorem 4.21** (Fundamental theorem unary).  $\forall \Sigma, \Psi, \Gamma, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash e : \tau \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge \\ & (\theta, n, \delta) \in [\Gamma \sigma]_V \implies \\ & (\theta, n, e \delta) \in [\tau \sigma]_E \end{aligned}$$

*Proof.* Proof by induction on *CG* typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, x \delta) \in [\tau \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some  $i < n$  s.t  $x \delta \Downarrow_i v$

(from cg-val we know that  $v = x \delta$  and  $i = 0$ )

It suffices to prove  $(\theta, n, x \delta) \in [\tau \sigma]_V$  (FU-V0)

Since  $(\theta, n, \delta) \in [\Gamma']_V$  where  $\Gamma' = \Gamma \cup \{x : \tau\}$ . Therefore from Definition 4.12 we know that  $(\theta, n, \delta(x)) \in [\Gamma'(x)]_V$

So we are done.

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e' : \tau_2}{\Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from cg-val we know that  $v = \lambda x. e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-L0})$$

From Definition 4.6 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

This means given some  $\theta'', v', j$  s.t  $\theta'' \sqsupseteq \theta, j < n$  and  $(\theta'', j, v') \in [\tau_1]_V$  (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 4.17 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2]_E, \text{ s.t } (\theta_i, j, v_x) \in [\tau_1]_V$$

Instantiating IH with  $\theta''$  and  $v'$  from (FU-L1) we get  $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$

3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2 \sigma]_V$$

This means that given some  $i < n$  s.t  $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2 \sigma]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have  $(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$

From cg-app we know that  $v_1 = \lambda x. e'$ . Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-P1})$$

This means from Definition 4.6 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2 \sigma]_E \quad (94)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1]_V$$

Since we know that  $(e_1 e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1]_V \quad (\text{FU-P2})$$

Instantiating Equation 94 with  $\theta, (n - j - k), v_2$  and since we know that  $(\theta, n - j - k, v_2) \in \lfloor \tau_1 \rfloor_V$  therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in \lfloor \tau_2 \sigma \rfloor_E$$

This means from Definition 4.7 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore we know that  $\exists J < i < n$  s.t  $i = j + k + J$  (since  $j + k + J < n$  therefore  $J < n - j - k$ ) and  $e'[v_2/x] \Downarrow_J v_f$

Therefore we have  $(\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$

Since we know that  $i = j + k + J$  and  $v = v_J$  therefore we get  $(\theta, n - i, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$  (so FU-P0 is proved)

#### 4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, (e_1, e_2) \delta) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V$$

Since we know that  $(e_1, e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e_1 \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

Since we know that  $(e_1 \ e_2) \delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_2 \delta \Downarrow_k v_2$ . This means we have

$$(\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from cg-prod we know that  $i = j + k + 1$  and  $v = (v_1, v_2)$  therefore from Definition 4.6 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau_1 \rfloor_V \text{ and } (\theta, n - j - k - 1, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

We get this from (FU-PA1) and Lemma 4.15 and from (FU-PA2) and Lemma 4.15

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{fst}(e') \delta) \in [\tau_1 \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_1 \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_1 \sigma]_V \quad (\text{FU-F0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$$

Since we know that  $\text{fst}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j (v_1, v_2)$ . This means we have

$$(\theta, n - j, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$$

From Definition 4.6 we know the following holds

$$(\theta, n - j, v_1) \in [\tau_1 \sigma]_V \text{ and } (\theta, n - j, v_2) \in [\tau_2 \sigma]_V \quad (\text{FU-F1})$$

From cg-fst we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-F0), we are required to prove

$$(\theta, n - j - 1, v_1) \in [\tau_1 \sigma]_V$$

We get this from (FU-F1) and Lemma 4.15

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{inl}(e') \delta) \in [(\tau_1 + \tau_2) \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 + \tau_2) \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [(\tau_1 + \tau_2) \sigma]_V \quad (\text{FU-LE0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [\tau_1 \sigma]_V$$

Since we know that  $\text{inl}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e' \delta \Downarrow_j v_1$ . This means we have

$$(\theta, n - j, v_1) \in [\tau_1 \sigma]_V \quad (\text{FU-LE1})$$

From cg-inl we know that  $v = v_1$  and  $i = j + 1$ . Therefore from (FU-LE0) we are required to prove

$$(\theta, n - j - 1, v_1) \in [(\tau_1 + \tau_2) \sigma]_V$$

From Definition 4.6 it suffices to prove

$$(\theta, n - j - 1, v_1) \in [\tau_1 \sigma]_V$$

We get this from (FU-LE1) and Lemma 4.15

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some  $i < n$  s.t.  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau \sigma]_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in [(\tau_1 + \tau_2) \sigma]_V$$

Since we know that  $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t.  $e_c \delta \Downarrow_j v_c$ . This means we have

$$(\theta, n - j, v_c) \in [(\tau_1 + \tau_2) \sigma]_V \quad (\text{FU-C1})$$

2 cases arise:

(a)  $v_c = \text{inl}(v_l)$ :

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V$$

Since we know that (case  $e_c, x.e_1, y.e_2$ )  $\delta \Downarrow_i v$  therefore  $\exists k < i - j$  (since  $i < n$  therefore  $i - j < n - j$ ) s.t  $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$ . This means we have

$$(\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C2})$$

From cg-case1 we know that  $i = j + k + 1$  and  $v = v_1$ . Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau \sigma \rfloor_V$$

We get this from (FU-C2) and Lemma 4.15

(b)  $v_c = \text{inr}(v_r)$ :

Symmetric reasoning as in the previous case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, \Lambda e' \delta) \in \lfloor (\forall \alpha. (\ell_e, \tau)) \sigma \rfloor_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \Lambda e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\forall \alpha. \tau) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $\lambda x. e' \delta \Downarrow_i v$

(from CG-Sem-val we know that  $v = \Lambda e' \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \Lambda e' \delta) \in \lfloor (\forall \alpha. \tau) \sigma \rfloor_V \quad (\text{FU-FI0})$$

From Definition 4.6 it further suffices to prove

$$\forall \theta'. \theta \sqsubseteq \theta', j < n. \forall \ell' \in \mathcal{L}. (\theta', j, e' \delta) \in \lfloor \tau[\ell'/\alpha] \rfloor_E$$

This means given some  $\theta', j, \ell' \in \mathcal{L}$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-FI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in \lfloor \tau[\ell'/\alpha] \sigma \rfloor_E \quad (\text{FU-FI2})$$

Since  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$  therefore from Lemma 4.17 we know that  $(\theta, j, \delta) \in \lfloor \Gamma \sigma \rfloor_V$  where  $j < n$  (from FU-L1)

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in \lfloor \tau \sigma \cup \{\alpha \mapsto \ell'\} \rfloor_E$$

(FU-FI2) is obtained directly from IH

11. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_E$

This means that from Definition 4.7 we need to prove

$\forall i < n. \nu e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(c \Rightarrow \tau) \sigma]_V$

This means that given some  $i < n$  s.t  $\nu e' \delta \Downarrow_i v$

(from CG-Sem-val we know that  $v = \nu e' \delta$  and  $i = 0$ )

It suffices to prove

$(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_V$  (FU-CI0)

From Definition 4.6 it further suffices to prove

$\mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < n. (\theta', j, e' \delta) \in [\tau]_E$

This means given  $\mathcal{L} \models c$  and some  $\theta', j$  s.t  $\theta' \sqsupseteq \theta, j < n$  (FU-CI1)

We are required to prove

$(\theta', j, (e' \delta)) \in [\tau \sigma]_E$  (FU-CI2)

Since  $(\theta, n, \delta) \in [\Gamma \sigma]_V$  therefore from Lemma 4.17 we know that  $(\theta, j, \delta) \in [\Gamma \sigma]_V$  where  $j < n$  (from FU-L1). Also we know that  $\mathcal{L} \models c \sigma$  therefore  $\mathcal{L} \models (\Sigma \cup \{c\}) \sigma$

IH:  $(\theta', j, e' \delta) \in [\tau \sigma]_E$

(FU-CI2) is obtained directly from IH

## 12. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' [] \delta) \in [\tau[\ell/\alpha] \sigma]_E$

This means that from Definition 4.7 we need to prove

$\forall i < n. e' [] \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V$

This means that given some  $i < n$  s.t  $e' [] \delta \Downarrow_i v$

It suffices to prove

$(\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V$  (FU-FE0)

IH:  $(\theta, n, e' \delta) \in [\forall \alpha. \tau]_E$

From Definition 4.7 we know that

$\forall h_1 < n. e' \delta \Downarrow_{h_1} \Lambda e_{h_1} \implies (\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$

Since  $e' [] \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \Lambda e_i$

Therefore we know that  $(\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$

From Definition 4.6 we know that

$$\forall \theta'' \sqsupseteq \theta, x < (n - h_1), \ell_h \in \mathcal{L}. (\theta'', x, e_{h_1}) \in [(\tau[\ell_h/\alpha]) \sigma]_E$$

Instantiating  $\theta''$  with  $\theta$ ,  $x$  with  $n - h_1 - 1$  and  $\ell_h$  with  $\ell$ . So, we get  
 $(\theta, n - h_1 - 1, e_{h_1}) \in [(\tau[\ell/\alpha]) \sigma]_E$

From Definition 4.7 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $e' \Downarrow \delta$  reduces in  $i$  steps therefore from CG-Sem-FE we know that ( $i = h_1 + h_2 + 1$ ) and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ . Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

Also given is  $\mathcal{L} \models \Psi \sigma \wedge$  and  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, e' \bullet \delta) \in [\tau \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. e' \bullet \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some  $i < n$  s.t  $e' \bullet \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau \sigma]_V \quad (\text{FU-CE0})$$

IH:  $(\theta, n, e' \delta) \in [c \Rightarrow \tau \sigma]_E$

From Definition 4.7 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \nu e_{h_1} \implies (\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$$

Since  $e' \bullet \delta$  reduces therefore we know that  $\exists h_1 < i < n$  such that  $e' \delta \Downarrow_{h_1} \nu e_{h_1}$

Therefore we know that  $(\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$

From Definition 4.6 we know that

$$\mathcal{L} \models c \sigma \implies \forall \theta'' \sqsupseteq \theta, x < (n - h_1). (\theta'', x, e_{h_1}) \in [\tau \sigma]_E$$

Since we know that  $\mathcal{L} \models c \sigma$  and then we instantiate  $\theta''$  with  $\theta$ ,  $x$  with  $n - h_1 - 1$ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [\tau \sigma]_E$$

From Definition 4.7 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$



Since  $e' \bullet \delta$  reduces in  $i$  steps therefore from CG-Sem-CE we know that  $(i = h_1 + h_2 + 1)$  and since we know that  $i < n$  therefore we have  $h_2 < n - h_1 - 1$  such that  $e_{h_1} \delta \Downarrow_{h_2} v$ . Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in \lfloor \tau \sigma \rfloor_V$$

Since we know that  $i = h_1 + h_2 + 1$  therefore we get

$$(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$$

14. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } (e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, \text{new } (e') \delta) \in \lfloor \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \sigma \rfloor_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \sigma \rfloor_V$$

This means that given some  $i < n$  s.t  $\text{new } (e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{new } (e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{new } (e') \delta) \in \lfloor \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \sigma \rfloor_V$$

From Definition 4.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor (\text{ref } \ell' \tau \sigma) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-ref we know that  $v' = a$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in \lfloor (\text{ref } \ell' \tau) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_E$$

From Definition 4.7 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$$

Since we know that  $(H, \text{new } (e')) \Downarrow_j^f (H', a)$  therefore from cg-ref we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h$$

Therefore we have

$$(\theta_e, n - l, v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V \quad (\text{FU-R2})$$

In order to prove (FU-R0) we choose  $\theta'$  as  $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_n$ :

From Definition 4.8 it suffices to prove that

$$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in [\theta_n(a)]_V$$

- $\text{dom}(\theta_n) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H) \cup \{a\}$

We know that  $\text{dom}(\theta_n) = \text{dom}(\theta_e) \cup \{a\}$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in [\theta_n(a)]_V$ :

Since from (FU-R2) we know that  $(\theta_h, n - l, v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V$

Since  $\theta_h \sqsubseteq \theta_n$  and  $k - j - 1 < n - l$  (since  $k < n$  and  $l < j$ ) therefore from

Lemma 4.15 we know that  $(\theta_n, k - j - 1, v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V$

(b)  $(\theta_n, k - j - 1, a) \in [(\text{ref } \ell' \tau) \sigma]_V$ :

From Definition 4.6 it suffices to prove that  $\theta_n(a) = \text{Labeled } \ell' \tau$

We get this by construction of  $\theta_n$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow_{\downarrow} \ell)$ :

From CG-ref we know that  $\ell \sqsubseteq \ell'$

15. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \tau)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (!e') \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau) \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. (!e') \delta \downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau) \sigma]_V$$

(From cg-val we know that  $v = !e' \delta$  and  $i = 0$ )

This means that given some  $i < n$  s.t.  $!e' \delta \downarrow_i !e' \delta$

It suffices to prove

$$(\theta, n, !e' \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau) \sigma]_V$$

From Definition 4.6 it suffices to prove

$$\forall k \leq n. \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow_{\downarrow} \top)$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-D0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau) \sigma]_E$$

From Definition 4.7 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau) \sigma]_V$$

Since we know that  $(H, !(e')) \Downarrow_j^f (H', a)$  therefore from cg-deref we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h, v_h = a$$

Therefore we have

$$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau) \sigma]_V \quad (\text{FU-D1})$$

In order to prove (FU-D0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 4.8 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

•  $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

And  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

And since  $H' = H$  (from cg-deref) so we are done

•  $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$ :

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since  $H' = H$  and from Lemma 4.15 we get

$$\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b)  $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau) \sigma]_V$ :

From cg-deref we know that  $H = H'$  and  $v' = H(a)$

From (FU-D1) and Definition 4.6 we know that  $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from cg-deref we know that  $j \geq 1$ . Therefore from Lemma 4.15 we get  $(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau) \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \top)$ :

Holds vacuously

16. CG-assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, (e_1 := e_2) \delta) \in [(\mathbb{C} \ell \perp \text{unit})]_E^{pc}$

This means that from Definition 4.7 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

This means that given some  $i < n$  s.t.  $(e_1 := e_2) \delta \Downarrow_i v$ .

It suffices to prove

$$(\theta, n - i, ()) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

From Definition 4.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-assign we know that  $v' = ()$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, ()) \in [\text{unit}]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0}) \end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists l < j < k$  s.t.  $e_1 \delta \Downarrow_l a$ . This means we have

$$(\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k - l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k - l - m, v_2) \in [\text{Labeled } \ell' \tau]_V$$

Since we know that  $(e_1 := e_2) \delta \Downarrow_j^f v$  therefore  $\exists m < j - l$  (since  $j < k$  therefore  $j - l < k - l$ ) s.t.  $e_2 \delta \Downarrow_m v_2$ . This means we have

$$(\theta, k - l - m, v_2) \in [(\text{Labeled } \ell' \tau)]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose  $\theta'$  as  $\theta_e$

Now we need to prove:

(a)  $(k - j, H') \triangleright \theta_e$ :

From Definition 4.8 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in \lfloor \theta_e(a) \rfloor_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$ :

We know that  $\text{dom}(H') = \text{dom}(H)$

And  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that  $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in \lfloor \theta_e(a) \rfloor_V$ :

$\forall a \in \text{dom}(\theta_e)$ .

- i.  $H(a) = H'(a)$ :

Since  $(k, H) \triangleright \theta_e$  therefore from Definition 4.8 we know that

$$(\theta_e, k - 1, H(a)) \in \lfloor \theta_e(a) \rfloor_V$$

Therefore from Lemma 4.15 we get

$$(\theta_e, k - 1 - j, H(a)) \in \lfloor \theta_e(a) \rfloor_V$$

- ii.  $H(a) \neq H'(a)$ :

From cg-assign we know that  $H'(a) = v_2$

From (FU-A1) we know that  $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that  $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k - l - m, v_2) \in \lfloor (\text{Labeled } \ell' \tau) \rfloor_V$$

Therefore we get

$$(\theta, k - j + 1, v_2) \in \lfloor (\text{Labeled } \ell' \tau) \rfloor_V$$

Therefore from Lemma 4.15 we get

$$(\theta, k - j - 1, v_2) \in \lfloor (\text{Labeled } \ell' \tau) \rfloor_V$$

(b)  $(\theta_e, k - j - 1, ()) \in \lfloor \text{unit} \rfloor_V$ :

From Definition 4.6

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$ :

From CG-assign we know that  $\ell \sqsubseteq \ell'$

(d)  $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$ :

Holds vacuously

17. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is  $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove:  $(\theta, n, \text{Lb}(e') \delta) \in \lfloor \text{Labeled } \ell \tau \sigma \rfloor_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \text{Labeled } \ell \tau \sigma \rfloor_V$$

This means we are given some  $i < n$  s.t  $\text{Lb}(e') \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in \lfloor \text{Labeled } \ell \tau \sigma \rfloor_V$$

Let  $v = \text{Lb}(v_i)$ . This means from Definition 4.6 we are required to prove

$$(\theta, n - i, v_i) \in \lfloor \tau \sigma \rfloor_V$$

IH:  $(\theta, n, e' \delta) \in [\tau \sigma]_E$

This means from Definition 4.7 we have

$$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau]_V$$

Since we know that  $\text{Lb}(e') \delta \Downarrow_i v$  therefore  $\exists j < i < n$  s.t  $e' \delta \Downarrow_j v_i$

Therefore we have  $(\theta, n - j, v_i) \in [\tau \sigma]_V$

From cg-label we know that  $i = j + 1$  therefore from Lemma 4.15 we have

$$(\theta, n - i, v_i) \in [\tau \sigma]_V$$

18. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \top \ell \tau) \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \top \ell \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{unlabel}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \top \ell \tau) \sigma]_V$$

From Definition 4.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ . Also from cg-unlabel we know that  $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-U0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell \tau) \sigma]_E$$

This means that from Definition 4.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from cg-unlabel we know that

$\exists h_1 < j < k$  s.t  $e' \delta \Downarrow_{h_1} \text{Lb } v'$

This means we have

$(\theta_e, k - h_1, \text{Lb } v') \in [(\text{Labeled } \ell \tau) \sigma]_V$

This means from Definition 4.6 we have

$(\theta_e, k - h_1, v') \in [\tau \sigma]_V$  (FU-U1)

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we a required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since have  $(k, H) \triangleright \theta_e$  therefore from Lemma 4.19 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-U1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 4.15 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)$ :

Holds vacuously

19. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell \ell' \tau}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau \sigma]_E$

This means that from Definition 4.7 we need to prove

$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau \sigma]_V$

This means we are given some  $i < n$  s.t  $\text{ret}(e') \delta \Downarrow_i v$  and we are required to prove

$(\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau \sigma]_V$

(from cg-val we know that  $v = \text{ret}(e') \delta$  and  $i = 0$ )

It suffices to prove

$(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau \sigma]_V$

From Definition 4.6 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

This means given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

Also from cg-ret we know that  $H' = H$

It suffices to prove

$$\begin{aligned}
& \exists \theta' \sqsupseteq \theta_e.(k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0})
\end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [\tau \sigma]_E$$

This means that from Definition 4.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [\tau \sigma]_V$$

Since we know that  $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$  therefore from cg-ret we know that

$$\exists h_1 < j < k \text{ s.t. } e' \delta \Downarrow_{h_1} v'$$

This means we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose  $\theta'$  as  $\theta_e$ . And we are required to prove:

(a)  $(k - j, H) \triangleright \theta_e$ :

Since we have  $(k, H) \triangleright \theta_e$  therefore from Lemma 4.19 we get  $(k - j, H) \triangleright \theta_e$

(b)  $(\theta', k - j, v') \in [\tau \sigma]_V$ :

Since from (FU-R1) we know that  $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since  $j = h_1 + 1$ , therefore from Lemma 4.15 we get  $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'')$ :

Holds vacuously

(d)  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$ :

Holds vacuously

20. CG-bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau'}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in [\mathbb{C} \ell \ell' \tau' \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau' \sigma]_V$$

This means we are given some  $i < n$  s.t.  $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$  and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau' \sigma]_V$$

(from cg-val we know that  $v = \text{bind}(e_1, x.e_2) \delta$  and  $i = 0$ )

Therefore we need to prove



$$(\theta, n, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \sigma \rfloor_V$$

From Definition 4.6 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in \lfloor \tau' \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$ .

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in \lfloor \tau' \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \sigma \rfloor_E$$

This means that from Definition 4.7 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \sigma \rfloor_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore from cg-bind we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \sigma \rfloor_V$$

From Definition 4.6 we know that

$$\begin{aligned} & \forall k_{h1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h1}) \wedge J < k_{h1} \implies \\ & \exists \theta'' \sqsupseteq \theta'_e.(k_{h1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h1} - J, v') \in \lfloor \tau \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $k_{h1}$  with  $k - h_1$ ,  $\theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 4.19 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} & \exists \theta'' \sqsupseteq \theta_e.(k_{h1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h1} - J, v') \in \lfloor \tau \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-B1}) \end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_E$$

This means that from Definition 4.7 we need to prove

$$\forall h_2 < k - h_1 - J.e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$  therefore from cg-bind we know that  $\exists h_2 < j - h_1 - J < k - h_1 - J$  s.t  $e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

From Definition 4.6 we know that

$$\begin{aligned} \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\ \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau']_V \wedge \\ (\forall a.H(a) \neq H''(a) \implies \exists \ell''.\theta'_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e).\theta'''(a) \searrow \ell_3) \end{aligned}$$

Since we know that  $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$  therefore  $\exists v_{h_2}, i$  s.t  $(v'' \Downarrow_i v_{h_2})$ . From cg-val we know that  $v_{h_2} = v''$  and  $i = 0$ . Instantiating  $k_{h_2}$  with  $k - h_1 - J - h_2$ ,  $\theta'_e$  with  $\theta''$ ,  $H$  with  $H'$  (from FU-B1) and  $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$  s.t  $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$ . And since we already know that  $(k - h_1, H') \triangleright \theta''$  therefore from Lemma 4.19 we get  $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned} \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau \sigma]_V \wedge \\ (\forall a.H(a) \neq H''(a) \implies \exists \ell''.\theta'_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e).\theta'''(a) \searrow \ell_3) \quad (\text{FU-B2}) \end{aligned}$$

We get (FU-B0) by choosing  $\theta'$  as  $\theta'''$  (from FU-B2)

21. CG-toLabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)}$$

Also given is  $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove:  $(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau) \sigma]_E$

This means that from Definition 4.7 we need to prove

$$\forall i < n.\text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau) \sigma]_V$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \delta \Downarrow_i v$

(from cg-val we know that  $v = \text{toLabeled}(e') \delta$  and  $i = 0$ )

It suffices to prove

$$(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau) \sigma]_V$$

From Definition 4.6 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell_1) \end{aligned}$$

And given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$ .  
Also from cg-tolabeled we know that  $H' = H$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \quad (\text{FU-TL0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E$$

This means that from Definition 4.7 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V$$

Since  $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$  therefore from cg-tolabeled we know that  $\exists h_1 < j < k$  s.t  $e' \delta \Downarrow_{h_1} v_1$

$$\text{Therefore we get } (\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V$$

From Definition 4.6 we know that

$$\begin{aligned} & \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J.(k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ & \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in [\tau \sigma]_V \wedge \\ & (\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $k_{h_1}$  with  $k - h_1, H_h$  with  $H, \theta'_e$  with  $\theta_e$ . Since we know that  $(H, \text{toLabeled}(e')) \Downarrow_j^f (H', v_1)$  therefore  $\exists J < j - h_1 < k - h_1$  s.t  $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$ . And since we already know that  $(k, H) \triangleright \theta_e$  therefore from Lemma 4.19 we get  $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} & \exists \theta'' \sqsupseteq \theta'_e.(k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-TL1}) \end{aligned}$$

In order to prove (FU-TL0) we choose  $\theta'$  as  $\theta''$ . Now we need to prove the following

(a)  $(k - j, H') \triangleright \theta''$ :

Since  $(k - h_1 - J, H') \triangleright \theta''$  and  $j = h_1 + J + 1$  therefore from Lemma 4.19 we get  $(k - j, H') \triangleright \theta''$

(b)  $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau)]_V$ :

From cg-tolabeled we know that  $v' = \text{toLabeled}(v_1)$

From Definition 4.4 it suffices to prove that  $(\theta'', k - j - 1, v_1) \in [\tau \sigma]_V$

We get this from (FU-TL1) and Lemma 4.15

(c)  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$ :

Directly from (FU-TL1)

(d)  $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$ :

Directly from (FU-TL1)

□

**Lemma 4.22** (Subtyping unary). *The following holds:*

$\forall \mathcal{L}, \tau, \tau'$ .

$$1. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$$

$$2. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E \subseteq [(\tau' \sigma)]_E$$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

IH1:  $[(\tau'_1 \sigma)]_V \subseteq [(\tau_1 \sigma)]_V$  (Statement (1))

$[(\tau_2)]_E \subseteq [(\tau'_2)]_E$  (Sub-A0, From Statement (2))

It suffices to prove:  $\forall (\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V. (\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

This means that given some  $\theta, n$  and  $\lambda x.e_i$  s.t  $(\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V$

Therefore from Definition 4.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in [\tau_1 \sigma]_V \implies (\theta_1, i, e_i[v/x]) \in [\tau_2 \sigma]_E \quad (95)$$

And it suffices to prove:  $(\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

Again from Definition 4.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in [\tau'_1 \sigma]_V \implies (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

This means that given some  $\theta_2, j < n, v$  s.t  $\theta \sqsubseteq \theta_2$  and  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$

And we are required to prove:  $(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$

Since  $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$  therefore from IH1 we know that  $(\theta_2, j, v) \in [\tau_1 \sigma]_V$

As a result from Equation 95 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2 \sigma]_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V$  (Statement (1))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:  $\forall (\theta, n, (v_1, v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V. (\theta, n, (v_1, v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V$

This means that given some  $\theta, n$  and  $(v_1, v_2)$   $(\theta, (v_1, v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V$

Therefore from Definition 4.6 we are given:

$$(\theta, n, v_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge (\theta, n, v_2) \in \llbracket \tau_2 \sigma \rrbracket_V \quad (96)$$

And it suffices to prove:  $(\theta, (v_1, v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V$

Again from Definition 4.6, it suffices to prove:

$$(\theta, n, v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V \wedge (\theta, n, v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V$$

Since from Equation 96 we know that  $(\theta, n, v_1) \in \llbracket \tau_1 \sigma \rrbracket_V$  therefore from IH1 we have  $(\theta, n, v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V$

Similarly since  $(\theta, n, v_2) \in \llbracket \tau_2 \sigma \rrbracket_V$  from Equation 96 therefore from IH2 we have  $(\theta, n, v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V$  (Statement (1))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:  $\forall (\theta, n, v_s) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V. (\theta, v_s) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

This means that given:  $(\theta, n, v_s) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V$

And it suffices to prove:  $(\theta, n, v_s) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V$

2 cases arise

(a)  $v_s = \text{inl } v_i$ :

From Definition 4.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_1 \sigma \rfloor_V \quad (97)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

From Equation 97 and IH1 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

(b)  $v_s = \text{inr } v_i$ :

From Definition 4.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_2 \sigma \rfloor_V \quad (98)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

From Equation 98 and IH2 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

#### 4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $\lfloor ((\forall \alpha. \tau_1) \sigma) \rfloor_V \subseteq \lfloor (\forall \alpha. \tau_2) \sigma \rfloor_V$

It suffices to prove:  $\forall (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. \tau_1) \sigma) \rfloor_V. (\theta, n, \Lambda e_i) \in \lfloor (\forall \alpha. \tau_2) \sigma \rfloor_V$

This means that given:  $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. \tau_1) \sigma) \rfloor_V$

Therefore from Definition 4.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E \quad (99)$$

And it suffices to prove:  $(\theta, n, \Lambda e_i) \in \lfloor (\forall \alpha. \tau_2) \sigma \rfloor_V$

Again from Definition 4.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E$$

This means that given some  $\theta_2, j < n, \ell' \in \mathcal{L}$  s.t  $\theta \sqsubseteq \theta_2$

And we are required to prove:  $(\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$  therefore from Equation 99 we have

$$(\theta_2, j, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E$$

$$\lfloor (\tau_1 (\sigma \cup [\alpha \mapsto \ell'])) \rfloor_E \subseteq \lfloor (\tau_2 (\sigma \cup [\alpha \mapsto \ell'])) \rfloor_E \text{ (Sub-F0, Statement (2))}$$

From (Sub-F0), we know that

$$(\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E$$

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $\llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V \subseteq \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V$

It suffices to prove:  $\forall (\theta, n, \nu e_i) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V. (\theta, n, \nu e_i) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V$

This means that given:  $(\theta, n, \nu e_i) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V$

Therefore from Definition 4.6 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in \llbracket \tau_1(\sigma) \rrbracket_E \quad (100)$$

And it suffices to prove:  $(\theta, n, \nu e_i) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V$

Again from Definition 4.6, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in \llbracket \tau_2(\sigma) \rrbracket_E$$

This means that given some  $\theta_2, j$  s.t  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

And we are required to prove:  $(\theta_2, j, e_i) \in \llbracket \tau_2(\sigma) \rrbracket_E$

Since we are given  $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$  and  $\mathcal{L} \models c_2 \sigma \implies c_1 \sigma$  therefore from Equation 100 we have

$$(\theta_2, j, e_i) \in \llbracket \tau_1(\sigma) \rrbracket_E$$

$$\llbracket (\tau_1 \sigma) \rrbracket_E \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E \text{ (Sub-C0, Statement (2))}$$

From (Sub-C0), we know that

$$(\theta_2, j, e_i) \in \llbracket \tau_2(\sigma) \rrbracket_E$$

6. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V \subseteq \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$

IH:  $\llbracket (\tau \sigma) \rrbracket_V \subseteq \llbracket (\tau' \sigma) \rrbracket_V$  (Statement (1))

It suffices to prove:

$$\forall (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V. (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V$$

This means that given some  $\theta, n$  and  $\text{Lb}(e_i)$  s.t  $(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V$

Therefore from Definition 4.6 we are given:

$$(\theta, n, v_i) \in \llbracket (\tau \sigma) \rrbracket_V \quad (\text{SL})$$

And we are required to prove that

$$(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V$$

From Definition 4.6 it suffices to prove

$$(\theta, n, v_i) \in [(\tau' \sigma)]_V$$

We get this directly from (SL) and IH

## 7. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $[((\mathbb{C} \ell_i \ell_o \tau)) ]_V \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$

IH:  $[(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$  (Statement (1))

It suffices to prove:

$$\forall (\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V. (\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$$

This means that given some  $\theta, n$  and  $e$  s.t  $(\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V$

Therefore from Definition 4.6 we are given:

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) & \quad (\text{SC0}) \end{aligned}$$

And we are required to prove

$$(\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau')) ]_V$$

So again from Definition 4.6 we need to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau' \sigma]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) & \end{aligned}$$

This means we are given some  $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$   
(SC1)

And we need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau' \sigma]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) & \end{aligned}$$

We instantiate (SC0) with  $k, \theta_e, H, j$  from (SC1) and we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) & \end{aligned}$$

Since  $\tau <: \tau'$  therefore from IH we get



$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \sigma \rfloor_V$$

And since  $\ell'_i \sqsubseteq \ell_i$  therefore we also have

$$\begin{aligned} (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) \end{aligned}$$

8. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E. (\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

This means that we are given  $(\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E$

From Definition 4.7 it means we have

$$\forall i < n.e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

From Definition 4.7 we need to prove

$$\forall i < n.e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

This further means that given some  $i < n$  s.t  $e \downarrow_i v$ , it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

Instantiating (Sub-E0) with the given  $i$  we get  $(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$

Finally from Statement(1) we get  $(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$

□

**Lemma 4.23** (Binary interpretation of  $\Gamma$  implies Unary interpretation of  $\Gamma$ ).  $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

*Proof.* Given:  $(W, n, \gamma) \in [\Gamma]_V^A$

To prove:  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 4.13 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case  $i = 1$

Given some  $m$  we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i)$ :

$$\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i)$$

Therefore,  $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i))$  (Given)

- $\forall x \in \text{dom}(\Gamma).(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$ :

We are given:  $\forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

Therefore from Lemma 4.14 we know that

$\forall m'.(W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Instantiating  $m'$  with  $m$  we get

$(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case  $i = 2$

Symmetric reasoning as in the  $i = 1$  case above

□

**Theorem 4.24** (Fundamental theorem binary).  $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$

$\Sigma; \Psi; \Gamma \vdash e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge$

$(W, n, \gamma) \in [\Gamma \sigma]_V^A \implies$

$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

*Proof.* Proof by induction on the typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

To prove:  $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau]_E^A$

Say  $e_1 = x(\gamma \downarrow_1)$  and  $e_2 = x(\gamma \downarrow_2)$

From Definition 4.5 it suffices to prove that

$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau \sigma]_V^A$

This means given some  $i < n$  s.t  $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2$

We are required to prove:  $(W, n - i, v'_1, v'_2) \in [\tau \sigma]_V^A$

From cg-val we know that  $x(\gamma \downarrow_1) \downarrow x(\gamma \downarrow_1)$  and  $x(\gamma \downarrow_2) \downarrow x(\gamma \downarrow_2)$

This means  $v'_1 = x(\gamma \downarrow_1)$  and  $v'_2 = x(\gamma \downarrow_2)$

Since  $(W, n, \gamma) \in [\tau \sigma]_V^A$ . Therefore from Definition 4.13 we know that

$(W, n, v'_1, v'_2) \in [\tau \sigma]_V^A$

From Lemma 4.16 we get

$(W, n - i, v'_1, v'_2) \in [\tau \sigma]_V^A$

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove:  $(W, n, \lambda x. e(\gamma \downarrow_1), \lambda x. e(\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

Say  $e_1 = \lambda x.e$  ( $\gamma \downarrow_1$ ) and  $e_2 = \lambda x.e$  ( $\gamma \downarrow_2$ )

From Definition of  $\lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_E^A$  it suffices to prove that

$$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in \lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^A$$

This means given some  $i < n$  s.t  $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From cg-val we know that  $v'_1 = (\lambda x.e_i)\gamma \downarrow_1$  and  $v'_2 = (\lambda x.e_i)\gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\lambda x.e_i)\gamma \downarrow_1, (\lambda x.e_i)\gamma \downarrow_2) \in \lceil (\tau_1 \rightarrow \tau_2) \sigma \rceil_V^A$$

From Definition 4.4 it suffices to prove

$$\forall W' \sqsupseteq W, j < n, v_1, v_2.$$

$$((W', j, v_1, v_2) \in \lceil \tau_1 \sigma \rceil_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in \lceil \tau_2 \sigma \rceil_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$$

$$((\theta_l, j, v_c) \in \lfloor \tau_1 \sigma \rfloor_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in \lfloor \tau_2 \sigma \rfloor_E) \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$$

$$((\theta_l, j, v_c) \in \lfloor \tau_1 \sigma \rfloor_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in \lfloor \tau_2 \sigma \rfloor_E) \quad (\text{FB-L0})$$

IH:

$$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in \lceil \tau_2 \sigma \rceil_E^A$$

s.t

$$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in \lceil \Gamma \rceil_V^A$$

In order to prove (FB-L0) we need to prove the following:

$$(a) \forall W' \sqsupseteq W, j < n, v_1, v_2.$$

$$((W', j, v_1, v_2) \in \lceil \tau_1 \sigma \rceil_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in \lceil \tau_2 \sigma \rceil_E^A):$$

This means given some  $W' \sqsupseteq W, j < n, v_1, v_2$  s.t.  $(W', j, v_1, v_2) \in \lceil \tau_1 \sigma \rceil_V^A$

We need to prove  $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in \lceil \tau_2 \sigma \rceil_E^A$

We get this by instantiating IH with  $W'$  and  $j$

$$(b) \forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$$

$$((\theta_l, j, v_c) \in \lfloor \tau_1 \sigma \rfloor_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in \lfloor \tau_2 \sigma \rfloor_E):$$

This means given some  $\theta_l \sqsupseteq W.\theta_1, v_c, j$  s.t  $(\theta_l, j, v_c) \in \lfloor \tau_1 \sigma \rfloor_V$

We need to prove:  $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in \lfloor \tau_2 \sigma \rfloor_E$

It is given to us that

$$(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$$

Therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$$

Instantiating  $m$  with  $j$  we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$$

From Lemma 4.18 we know that

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Since we know that  $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

Therefore we also have

$$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1 \sigma\}]_V$$

Therefore, we can apply Theorem 4.21 to obtain

$$(\theta_l, j, e[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_V$$

$$(c) \quad \forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E):$$

Similar reasoning as in the previous case

### 3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

To prove:  $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

This further means that given some  $i < n$  s.t  $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

$$\underline{\text{IH1}}: (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_1 \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e_1 \gamma \downarrow_2 \downarrow v_{h2}$

$$\text{This means we have } (W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$$

From cg-app we know that  $val_{h1} = \lambda x. e_{h1}$  and  $val_{h2} = \lambda x. e_{h2}$

From Definition 4.4 this further means

$$\forall W' \sqsupseteq W, J < (n - j), v_1, v_2.$$

$$((W', J, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', J, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2 \sigma]_E) \wedge$$

$$\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \quad (\text{FB-A1})$$

$$\underline{\text{IH2}}: (W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 4.5 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_j v_{h1'} \wedge e_2 \gamma \downarrow_2 \Downarrow v_{h2'} \implies (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t.  $e_2 \gamma \downarrow_1 \Downarrow_k v_{h1'}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_2 \gamma \downarrow_2 \Downarrow v_{h2'}$

$$\text{This means we have } (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-A2})$$

Instantiating the first conjunct of (FB-A1) as follows  $W'$  with  $W$ ,  $J$  with  $n - j - k$ ,  $v_1$  and  $v_2$  with  $v'_{h1}$  and  $v'_{h2}$  respectively, we obtain

$$(W, n - j - k, e_{h1}[v'_{h1}/x], e_{h2}[v'_{h2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 4.5

$$\forall l < n - j - k. (e_{h1}[v'_{h1}/x]) \gamma \downarrow_l \Downarrow v_{f1} \wedge e_{h2}[v'_{h2}/x] \Downarrow v_{f2} \implies (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since we know that  $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists l < i - j - k < n - j - k$  s.t.  $e_{h1}[v'_{h1}/x] \Downarrow_l v_{f1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_{h2}[v'_{h2}/x] \Downarrow v_{f2}$

$$\text{Therefore we have } (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since  $i = j + k + l$  therefore we are done

#### 4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

$$\text{To prove: } (W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall i < n. (e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2}) \implies (W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means that given some  $i < n$  s.t.  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2})$

We are required to prove

$$(W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A \quad (\text{FB-P0})$$

$$\underline{\text{IH1:}} (W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 4.5 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \Downarrow_j v_{f1} \wedge e_1 \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2})$ . Therefore  $\exists j < i < n$  s.t.  $e_1 \gamma \downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A \quad (\text{FB-P1})$$

$$\underline{\text{IH2:}} (W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$$

This means from Definition 4.5 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \Downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2})$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_2 \gamma \downarrow_1 \Downarrow_j v_{f_2}$ . Similarly since  $(e_1, e_2) \gamma \downarrow_2 \Downarrow v_{f_2}$  therefore  $e_2 \gamma \downarrow_2 \Downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 4.4 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathcal{V}}^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since  $i = j + k + 1$  therefore from (FB-P1) and (FB-P2) and from Lemma 4.16 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove:  $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

Since we know that  $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$  therefore  $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

From Definition 4.4 we know that

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Since from cg-fst  $i = j + 1$  therefore from Lemma 4.16 we get

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove:  $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f1}) \wedge \text{inl}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f1}) \implies \\ (W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f1}) \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_j v_{f1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ (W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

Since we know that  $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $\text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f1})$  therefore  $e' \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 4.4 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

From cg-inl since  $i = j + 1$  therefore from (FB-IL1) and Lemma 4.16 we get (FB-IL0)

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove:  $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2} \implies \\ (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

This means that given some  $i < n$  s.t  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2}$

We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_i v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_{\mathcal{V}}^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$  therefore  $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_{\mathcal{V}}^A \quad (\text{FB-C1})$$

2 cases arise

- (a)  $v_{h1} = \text{inl}(v_1)$  and  $v'_{h1} = \text{inl}(v'_1)$ :

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 4.5 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_i v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies \\ (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

Since we know that  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$ . Therefore  $\exists k < i - j < n - j$  s.t  $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$ . Similarly since  $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$  therefore  $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

From cg-case1 we know that  $i = j + k + 1$  therefore from Lemma 4.16 we get (FB-C0)

- (b)  $v_{h1} = \text{inr}(v_1)$  and  $v'_{h1} = \text{inr}(v'_1)$ :

Symmetric case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

To prove:  $(W, n, \Lambda e' (\gamma \downarrow_1), \Lambda e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 4.5 it suffices to prove that



$$\forall i < n. (\Lambda e')\gamma \downarrow_1 \downarrow_i v_{f_1} \wedge (\Lambda e')\gamma \downarrow_2 \downarrow v_{f_2} \implies (W, n - i, v_{f_1}, v_{f_2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

This means given some  $i < n$  s.t  $(\Lambda e')\gamma \downarrow_1 \downarrow_i v_{f_1} \wedge (\Lambda e')\gamma \downarrow_2 \downarrow v_{f_2}$

From CG-Sem-val we know that  $v_{f_1} = (\Lambda e')\gamma \downarrow_1$  and  $v_{f_2} = (\Lambda e')\gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\Lambda e')\gamma \downarrow_1, (\Lambda e')\gamma \downarrow_2) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Let  $e_1 = (\Lambda e')\gamma \downarrow_1$  and  $e_2 = (\Lambda e')\gamma \downarrow_2$

From Definition 4.4 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FI0}) \end{aligned}$$

$$\underline{\text{IH}}: \forall W, n. (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E^A$$

In order to prove (FB-FI0) we need to prove the following

$$(a) \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A):$$

This means given  $W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}$  and we are required to prove

$$(W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

$$(b) \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E:$$

This means given  $\theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j$  and we are required to prove

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

Since from Lemma 4.23

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 4.16 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 4.21 to get

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

$$(c) \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E:$$

Symmetric reasoning as before

## 11. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

To prove:  $(W, n, e' [] (\gamma \downarrow_1), e' [] (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 4.5 it suffices to prove that

$$\forall i < n. (e' [])\gamma \downarrow_1 \downarrow_i v_{f_1} \wedge (e' [])\gamma \downarrow_2 \downarrow v_{f_2} \implies (W, n - i, v_{f_1}, v_{f_2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

This means given some  $i < n$  s.t  $(e' [])\gamma \downarrow_1 \downarrow_i v_{f_1} \wedge (e' [])\gamma \downarrow_2 \downarrow v_{f_2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_{\mathcal{V}}^A \quad (\text{FB-FE0})$$

$$\underline{\text{IH}}: (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$$

From Definition 4.5 it suffices to prove that

$$\forall i < n. (e') \gamma \downarrow_1 \downarrow_i v_{h1} \wedge (e') \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e'[]) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e'[]) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e' \gamma \downarrow_2 \downarrow v_{h2}$

$$\text{This means we have } (W, n - j, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_{\mathcal{V}}^A$$

From CG-Sem-FE we know that  $v_{h1} = \Lambda e_{h1}$  and  $v_{h2} = \Lambda e_{h2}$

From Definition 4.4 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, k < (n - j), \ell' \in \mathcal{L}. ((W', k, e_{h1}, e_{h2}) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h1}) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h2}) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FE1}) \end{aligned}$$

Instantiating the first conjunct of (FB-FE1) with  $W, n - j - 1$  and  $\ell$  we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau[\ell/\alpha] \sigma]_E^A$$

This means from Definition 4.5 we know that

$$\forall l < n - j - 1. (e_{h1}) \downarrow_l v_{f1} \wedge e_{h2} \downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_{\mathcal{V}}^A$$

Since we know that  $(e'[]) \gamma \downarrow_1 \downarrow_i v_{f1}$  therefore from CG-Sem-FE we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t  $e_{h1} \gamma \downarrow_1 \downarrow_j v_{f1}$ . Similarly since  $(e'[]) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e_{h2} \gamma \downarrow_2 \downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_{\mathcal{V}}^A \quad (\text{FB-FE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-FE2) we get (FB-FE0)

## 12. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

To prove:  $(W, n, \nu e' (\gamma \downarrow_1), \nu e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 4.5 it suffices to prove that

$$\forall i < n. (\nu e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e') \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(c \Rightarrow \tau) \sigma]_{\mathcal{V}}^A$$

This means given some  $i < n$  s.t  $(\nu e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e') \gamma \downarrow_2 \downarrow v_{f2}$

From CG-Sem-val we know that  $v_{f1} = (\nu e') \gamma \downarrow_1$  and  $v_{f2} = (\nu e') \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\nu e') \gamma \downarrow_1, (\nu e') \gamma \downarrow_2) \in [(c \Rightarrow \tau) \sigma]_{\mathcal{V}}^A$$

Let  $e_1 = (\nu e') \gamma \downarrow_1$  and  $e_2 = (\nu e') \gamma \downarrow_2$

From Definition 4.4 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c &\implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c &\implies (\theta_l, j, e_1) \in [\tau \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c &\implies (\theta_l, j, e_2) \in [\tau \sigma]_E \quad (\text{FB-CI0}) \end{aligned}$$

$$\underline{\text{IH}}: \forall W, n. (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

In order to prove (FB-CI0) we need to prove the following

$$(a) \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \sigma \implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A:$$

This means given  $W' \sqsupseteq W, j < n, \mathcal{L} \models c \sigma$  and we are required to prove

$$(W', j, e_1, e_2) \in [\tau \sigma]_E^A$$

Instantiating IH with  $W'$  and  $j$  we get the desired

$$(b) \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma \implies (\theta_l, j, e_1) \in [\tau \sigma]_E:$$

This means given  $\theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma$  and we are required to prove

$$(\theta_l, j, e_1) \in [\tau \sigma]_E$$

$$\text{Since from Lemma 4.23 } (W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 4.16 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 4.21 to get

$$(\theta_l, j, e_1) \in [\tau \sigma]_E$$

$$(c) \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau \sigma]_E:$$

Symmetric reasoning as before

### 13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

$$\text{To prove: } (W, n, e' \bullet (\gamma \downarrow_1), e' \bullet (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$$

From Definition 4.5 it suffices to prove that

$$\forall i < n. (e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

This means given some  $i < n$  s.t  $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-CE0})$$

$$\underline{\text{IH}}: (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$$

From Definition 4.5 it suffices to prove that

$$\forall i < n. e' \gamma \downarrow_1 \downarrow_i v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(c \Rightarrow \tau) \sigma]_V^A$$

Since we know that  $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1}$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \downarrow_1 \downarrow_j v_{h1}$ . Similarly since  $(e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$  therefore  $e' \gamma \downarrow_2 \downarrow v_{h2}$

This means we have  $(W, n - j, v_{h1}, v_{h2}) \in \lceil (c \Rightarrow \tau) \sigma \rceil_V^A$

From CG-Sem-CE we know that  $v_{h1} = \nu e_{h1}$  and  $v_{h2} = \nu e_{h2}$

From Definition 4.4 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, k < n - j. \mathcal{L} \models c \sigma &\implies (W', k, e_1, e_2) \in \lceil \tau \sigma \rceil_E^A \wedge \\ \forall \theta_l \sqsupseteq W. \theta_l, k. \mathcal{L} \models c \sigma &\implies (\theta_l, k, e_1) \in \lfloor \tau \sigma \rfloor_E \wedge \\ \forall \theta_l \sqsupseteq W. \theta_l, k. \mathcal{L} \models c \sigma &\implies (\theta_l, k, e_2) \in \lfloor \tau \sigma \rfloor_E \quad (\text{FB-CE1}) \end{aligned}$$

Instantiating the first conjunct of (FB-CE1) with  $W, n - j - 1$  and since we know that  $\mathcal{L} \models c \sigma$  therefore we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in \lceil \tau \sigma \rceil_E^A$$

This means from Definition 4.5 we know that

$$\forall l < n - j - 1. (e_{h1}) \Downarrow_l v_{f1} \wedge e_{h2} \Downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in \lceil \tau \sigma \rceil_V^A$$

Since we know that  $(e' \bullet) \gamma \Downarrow_1 \Downarrow_i v_{f1}$  therefore from CG-Sem-CE we know that  $(i = j + l + 1)$  and since we know that  $i < n$  therefore we have  $l < n - j - 1$  s.t  $e_{h1} \gamma \Downarrow_1 \Downarrow_l v_{f1}$ . Similarly since  $(e' \bullet) \gamma \Downarrow_2 \Downarrow v_{f2}$  therefore  $e_{h2} \gamma \Downarrow_2 \Downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in \lceil \tau \sigma \rceil_V^A \quad (\text{FB-CE2})$$

Since we know that  $i = j + l + 1$  therefore from (FB-CE2) we get (FB-CE0)

14. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove:  $(W, n, \text{Lb}(e') (\gamma \Downarrow_1), \text{Lb}(e') (\gamma \Downarrow_2)) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_E^A$

This means from Definition 4.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1}) &\implies \\ (W, n - i, \text{Lb}(v_{f1}), \text{Lb}(v'_{f1})) &\in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_V^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$

We are required to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_V^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \Downarrow_1), e' (\gamma \Downarrow_2)) \in \lceil \tau \sigma \rceil_E^A$$

This means from Definition 4.5 we have:

$$\forall j < n. e' \gamma \Downarrow_1 \Downarrow_j v_{f1} \wedge e' \gamma \Downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, v_{f1}, v'_{f1}) \in \lceil \tau \sigma \rceil_V^A$$

Since we know that  $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1})$ . Therefore  $\exists j < i < n$  s.t  $e' \gamma \Downarrow_1 \Downarrow_j v_{f1}$ . Similarly since  $\text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$  therefore  $e' \gamma \Downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau \sigma]_{\mathcal{V}}^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 4.4 it suffices to prove that

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau \sigma]_{\mathcal{V}}^A$$

From cg-label we know that  $i = j + 1$ . Therefore we get the desired from (FB-LB1) and Lemma 4.16

15. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

To prove:  $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \top \ell \tau) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies \\ (W, n - i, v_{f_1}, v'_{f_1}) \in [(\mathbb{C} \top \ell \tau) \sigma]_{\mathcal{V}}^A \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

From cg-val we know that  $v_{f_1} = \text{unlabel}(e') \gamma \downarrow_1$  and  $v'_{f_1} = \text{unlabel}(e') \gamma \downarrow_2$ . Also  $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\mathbb{C} \top \ell \tau) \sigma]_{\mathcal{V}}^A$$

This means from Definition 4.4 we need to prove

Let  $e_1 = \text{unlabel}(e') \gamma \downarrow_1$  and  $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau']_{\mathcal{V}} \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right) \end{aligned}$$

We need to show

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau \sigma): \end{aligned}$$

Also given is some  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau \sigma) \quad (\text{FB-U0})$$

III:  $(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_E^A$

This means from Definition 4.5 we are given

$$\forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) \implies \\ (W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_V^A$$

Since we know that

$$(H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \text{ therefore} \\ \exists I < j < k \text{ s.t. } e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1})$$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_V^A$$

This means from Definition 4.4 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau \sigma) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose  $W'$  as  $W_e$  and from cg-unlabel we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . And we already know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore from Lemma 4.20 we get  $(k - j, H_1, H_2) \triangleright W_e$

From cg-unlabel we know that  $v'_1, v'_2$  in (FB-U0) is  $v_{h1}, v'_{h1}$  respectively. And since from (FB-U1) we know that  $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau \sigma)$ . Therefore from Lemma 4.25 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, \ell, v_{h1}, v'_{h1}, \tau \sigma)$$

$$(b) \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right):$$

Case  $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W.\theta_1, H, j \text{ s.t. } (k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1) \wedge j < k$$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)$$

Since  $(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in \lfloor \Gamma \rfloor_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$

Now we can apply Theorem 4.21 to get

$$(W.\theta_1, k, (\text{unlabel } e') \gamma \downarrow_1) \in \lfloor (\mathbb{C} \top \ell \tau) \sigma \rfloor_E$$

This means from Definition 4.7 we get

$$\forall c < k. (\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in \lfloor (\mathbb{C} \top \ell \tau) \sigma \rfloor_V$$

This further means that given some  $c < k$  s.t.  $(\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v$ . From cg-val we know that  $c = 0$  and  $v = (\text{unlabel } e') \gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in [(\mathbb{C} \top \ell \tau) \sigma]_V$

From Definition 4.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K &\implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau \sigma]_V \wedge & \\ (\forall a.H_1(a) \neq H'(a) \implies \exists \ell'.\theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e).\theta'(a) \searrow \top) & \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

16. CG-tolabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)}$$

To prove:  $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{toLabeled}(e') \gamma \downarrow_1$ ,  $v_{f2} = \text{toLabeled}(e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_V^A$$

Let  $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$  and  $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 4.4 we are required to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. & \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies & \\ \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell_2 \tau) \sigma)) \wedge & \\ \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies & \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge & \\ (\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell_1)) & \end{aligned}$$

We need to prove:

$$\begin{aligned} \text{(a)} \quad \forall k \leq n, W_e \sqsupseteq W.\forall H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. & \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies & \\ \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell_2 \tau) \sigma): & \end{aligned}$$

This means that we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t

$$(k, H_1, H_2) \triangleright W_e \text{ and } (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell_2 \tau) \sigma)$$

From Definition 4.3 it suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell_2 \tau) \sigma]_V^A$$

Further from Definition 4.4 it suffices to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v''_1, v''_2, \tau \sigma) \quad (\text{FB-TL0})$$

where  $v'_1 = \text{Lb} v''_1$  and  $v'_2 = \text{Lb} v''_2$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau \sigma]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v''_{h1} \implies (W_e, n - J, v_{h1}, v''_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau \sigma]_V^A$$

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j (H'_1, v'_1)$  and  $(H_2, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j (H'_2, v'_2)$ . Therefore from cg-val we know that  $\exists J < j < k \leq n$  s.t  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly we also know that  $e' \gamma \downarrow_2 \Downarrow v''_{h1}$

This means we have

$$(W_e, k - J, v_{h1}, v''_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau \sigma]_V^A$$

From Definition 4.4 we know that

$$\begin{aligned} & (\forall k_1 \leq (k - J), W_e'' \sqsupseteq W_e. \forall H''_1, H''_2. (k_1, H''_1, H''_2) \triangleright W_e'' \wedge \forall v''_1, v''_2, m. \\ & (H''_1, v_{h1}) \Downarrow_m^f (H'_1, v''_1) \wedge (H''_2, v''_{h1}) \Downarrow^f (H'_2, v''_2) \wedge m < k_1 \implies \\ & \exists W' \sqsupseteq W_e''. (k_1 - m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1 - m, \ell_2, v''_1, v''_2, \tau \sigma) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{FB-TL1}) \end{aligned}$$

We instantiate  $W_e''$  with  $W_e$ ,  $H''_1$  with  $H_1$ ,  $H''_2$  with  $H_2$  and  $k_1$  with  $k$  in (FB-TL1).

Since we know that  $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$ , therefore  $\exists m < j < k \leq n$  s.t  $(H_1, v_{h1}) \Downarrow_m^f (H'_1, v'_1) \wedge (H_2, v''_{h1}) \Downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W' \sqsupseteq W_e.(k - m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - m, \ell_2, v''_1, v''_2, \tau \sigma) \quad (\text{FB-TL2})$$

In order to prove (FB-TL0) we choose  $W'$  as  $W'$  from (FB-TL2). Since from cg-tolabeled we know that  $v'_1 = \text{Lb}(v''_1)$ ,  $v'_2 = \text{Lb}(v''_2)$  and  $j = m + 1$  (therefore from Lemma 4.20 we get  $(k - j, H'_1, H'_2) \triangleright W'$ ) and from (FB-TL2) and Lemma 4.25 we get  $\text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v''_1, v''_2, \tau \sigma)$

$$\begin{aligned} \text{(b)} \quad & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_2 \tau) \sigma]_V \wedge \right. \end{aligned}$$



$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)$ :

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_2 \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 4.21 to get

$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_E$

This means from Definition 4.7 we get

$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$

Instantiating  $c$  with 0 and from cg-val we know  $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$

From Definition 4.6 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$   
 $\exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\text{Labeled } \ell_2 \tau]_V \wedge$   
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1)$

Instantiating  $K$  with  $k, \theta'_e$  with  $\theta_e, H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

17. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_1 \ell_2 \tau}$$

To prove:  $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$\forall i < n. \text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies$   
 $(W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^A$

This means that given some  $i < n$  s.t.  $\text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{ret}(e')\gamma \downarrow_1, v_{f2} = \text{ret}(e')\gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{ret}(e')\gamma \downarrow_1, \text{ret}(e')\gamma \downarrow_2) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^A$$

Let  $v_1 = \text{ret}(e')\gamma \downarrow_1$  and  $v_2 = \text{ret}(e')\gamma \downarrow_2$

From Definition 4.4 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall v, i. (e_l \Downarrow_i v) \implies \right. \\ & \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \end{aligned}$$

It suffices to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau \sigma): \end{aligned}$$

We are given is some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2, v'_1, v'_2, j < k$  s.t  $(k, H_1, H_2) \triangleright W_e$  and  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

From cg-ret we know that  $H'_1 = H_1$  and  $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau \sigma) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau \sigma]_V^A$$

Since we know that  $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow_j^f (H_2, v'_2)$ , therefore  $\exists J < j < k$  s.t  $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$  and similarly  $e' \gamma \downarrow_2 \Downarrow v'_{h1}$ .

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau \sigma]_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_e$  and from cg-ret we know that  $v'_1 = v_{h1}$  and  $v'_2 = v'_{h1}$ . We need to prove the following:

- i.  $(k - j, H_1, H_2) \triangleright W_e$ :  
Since we have  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 4.20 we get  $(k - j, H_1, H_2) \triangleright W_e$
- ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_2, v'_1, v'_2, \tau \sigma)$ :  
2 cases arise:

A.  $\ell_2 \sqsubseteq \mathcal{A}$ :

In this case from Definition 4.3 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in [\tau \sigma]_V^A$$

Since  $j = J + 1$  therefore we get this from (FB-R1) and Lemma 4.16

B.  $\ell_2 \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 4.3 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in [\tau \sigma]_V \text{ and } \forall m. (W_e, m, v'_2) \in [\tau \sigma]_V$$

We get this From (FB-R1) and Lemma 4.14

$$\begin{aligned} \text{(b)} \quad \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) : \right. \end{aligned}$$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t.  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_o \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_o) \end{aligned}$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 4.21 to get

$$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E$$

This means from Definition 4.7 we get

$$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V$$

Instantiating  $c$  with 0 and from cg-val we know that  $v = (\text{ret } e')\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V$

From Definition 4.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

18. CG-bind:

$$\frac{\Gamma \vdash e_l : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_b : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_l, x.e_b) : \mathbb{C} \ell \ell' \tau'}$$

To prove:  $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in [(\mathbb{C} \ell \ell' \tau') \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell \ell' \tau') \sigma]_V^A$$

This means that given some  $i < n$  s.t  $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{bind}(e_l, x.e_b) \gamma \downarrow_1$ ,  $v_{f2} = \text{bind}(e_l, x.e_b) \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b) \gamma \downarrow_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \in [(\mathbb{C} \ell \ell' \tau') \sigma]_V^A$$

Let  $v_1 = \text{bind}(e_l, x.e_b) \gamma \downarrow_1$  and  $v_2 = \text{bind}(e_l, x.e_b) \gamma \downarrow_2$

This means from Definition 4.4 we need to prove

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau \sigma): \end{aligned}$$

This means we are given some  $k \leq n, W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau \sigma) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ (W_e, k - f, v_{h1}, v'_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^A$$

This means from Definition 4.4 we have

$$\begin{aligned}
& \left( \forall K \leq (k - f), W'_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W'_e \wedge \forall v''_1, v''_2, J. \right. \\
& (H''_1, v_{h1}) \Downarrow_J^f (H''_1, v''_1) \wedge (H''_2, v_{h1}) \Downarrow^f (H''_2, v''_2) \wedge J < K \implies \\
& \left. \exists W'' \sqsupseteq W'_e. (K - J, H''_1, H''_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell_2, v''_1, v''_2, \tau \sigma) \right) \wedge \\
& \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \tau \sigma \rrbracket_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right)
\end{aligned}$$

Instantiating  $K$  with  $(k - f)$ ,  $W'_e$  with  $W_e$ ,  $H''_1$  with  $H_1$  and  $H''_2$  with  $H_2$  in the first conjunct of the above equation. Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 4.20 we also have  $(k - f, H_1, H_2) \triangleright W_e$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J < j - f < k - f$  s.t  $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v''_1) \wedge (H_2, v_{h1}) \Downarrow^f (H'_2, v''_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e. (k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell_2, v''_1, v''_2, \tau \sigma) \quad (\text{FB-B1})$$

From Definition 4.3 two cases arise:

i.  $\ell_2 \sqsubseteq \mathcal{A}$ :

$$\text{In this case we know that } (W'', k - f - J, v''_1, v''_2) \in \llbracket \tau \sigma \rrbracket_V^A$$

IH2:

$$(W'', k - f - J, e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}), e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\})) \in \llbracket (\mathbb{C} \ell_3 \ell_4 \tau') \sigma \rrbracket_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall s < k - f - J. e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2} \implies (W'', k - f - J - s, v_{h2}, v'_{h2}) \in \llbracket (\mathbb{C} \ell_3 \ell_4 \tau') \sigma \rrbracket_V^A$$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f - J < k - f - J$  s.t  $e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in \llbracket (\mathbb{C} \ell_3 \ell_4 \tau') \sigma \rrbracket_V^A$$

This means from Definition 4.4 we know that

$$\begin{aligned}
& \left( \forall K_s \leq (k - f - J - s), W_s \sqsupseteq W''. \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right. \\
& (H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies \\
& \left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_4, v'_1, v'_2, \tau' \sigma) \right) \wedge \\
& \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \tau \sigma \rrbracket_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_3 \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_3) \right)
\end{aligned}$$

Instantiating  $K_s$  with  $(k - f - J - s)$ ,  $W_s$  with  $W''$ ,  $H_1$  with  $H'_1$  and  $H'_2$  with  $H_2$ . Since we know that  $(k - f - J, H'_1, H'_2) \triangleright W''$  therefore from Lemma 4.20 we also have  $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that  $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists J_s < j - f - J - s < k - f - J - s$  s.t.  $(H'_1, v'_1) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_2) \Downarrow^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s. (k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau' \sigma) \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose  $W'$  as  $W'_s$ . From cg-bind we know that  $H'_1 = H'_{s1}$ ,  $H'_2 = H'_{s2}$ ,  $v'_1 = v'_{s1}$ ,  $v'_2 = v'_{s2}$  and  $j = f + J + s + J_s + 1$ . And we need to prove:

A.  $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$ :

Since from (FB-B2) we know that  $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$  therefore from Lemma 4.20 we get

$$(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$$

B.  $\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau' \sigma)$ :

Since from (FB-B2) we know that  $\text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau' \sigma)$  therefore from Lemma 4.25 we get

$$\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau' \sigma)$$

ii.  $\ell_2 \not\sqsubseteq \mathcal{A}$ :

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma)$$

Since  $\ell_2 \sqsubseteq \ell_4 \sqsubseteq \ell'$  and  $\ell \not\sqsubseteq \mathcal{A}$  therefore we have  $\ell_4 \not\sqsubseteq \mathcal{A}$  and  $\ell' \not\sqsubseteq \mathcal{A}$

This means that from Definition 4.3 it suffices to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'.\theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge \forall m_{u2}. (W'.\theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V$$

This means given some  $m_{u1}, m_{u2}$  and we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in [\tau' \sigma]_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau' \sigma]_V \quad (\text{FB-B01})$$

In this case from (FB-B1) and Definition 4.3 we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in [\tau \sigma]_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in [\tau \sigma]_V \quad (\text{FB-B3})$$

Since  $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_j v'_1$  therefore  $\exists J_1 < j - f - J < k - f - J$  s.t.  $(e_b) \gamma \downarrow_1 \cup \{x \mapsto v''_1\} \Downarrow_{J_1} v'_1$ . Similarly,  $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$  s.t.  $(H'_1, v'_1) \Downarrow_{J'_1}^f -$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  in the first conjunct of (FB-B3)

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in [\tau \sigma]_V$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating  $m$  with  $m_{u1} + 1 + J_1 + J'_1$  we get  $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 4.17 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V \quad (\text{FB-B4})$$

Now we can apply Theorem 4.21 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_E$$

This means from Definition 4.7 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_V \quad (\text{FB-B5})$$

Instantiating  $c_1$  with  $J_1$  in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_V$$

From Definition 4.6 we have

$$\begin{aligned} \forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies \\ \exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \end{aligned}$$

Instantiating  $K$  with  $m_{u1} + 1 + J'_1$ ,  $\theta'_e$  with  $W''.\theta_1$ ,  $H_1$  with  $H'_1$  (from FB-B1) and  $J_2$  with  $J'_1$  we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1. (m_{u1} + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow v'_2$ . Say this reduction happens in  $t$  steps. Therefore  $\exists t_1 < t < k \leq n$  s.t.  $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{t_1} v_{l2}$  and similarly  $\exists t_2 < t - t_1 < k - t_1$  s.t.  $(H, v_{l2})\gamma \downarrow_2 \downarrow_{t_2}^f (H''_2, v''_2)$

Again since  $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow_t v'_2$  therefore  $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$  s.t.  $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{J_2} v'_2$ . Similarly  $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$  s.t.  $(H'_2, v'_2) \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with  $m_{u2} + 1 + J_2 + J'_2$  we get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v''_2) \in [\tau \sigma]_V$$

Again since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $m_{u2} + 1 + J_2 + J'_2$  we get  $(W.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 4.17 we know that

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 4.21 to get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_E$$

This means from Definition 4.7 we get

$$\forall c_2 < (m_{u2} + 1 + J_2 + J'_2). (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{c_2} v_{o2} \implies (W''.\theta_2, m_{u2} + 1 + J_2 - c_2, v_{o2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_V \quad (\text{FB-B8})$$

Instantiating  $c_2$  with  $J_2$  in (FB-B8) we get

$$(W''.\theta_2, m_{u2} + 1 + J'_2, v_{o2}) \in \llbracket (\mathbb{C} \ell_3 \ell_4 \tau') \sigma \rrbracket_V$$

From Definition 4.6 we have

$$\forall K \leq (m_{u2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3. (K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o2}) \Downarrow_{J_3}^f (H_2'', v'_2) \wedge J_3 < K \implies$$

$$\begin{aligned} & \exists \theta'_2 \sqsupseteq \theta'_e. (K - J_3, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v'_2) \in \llbracket \tau' \sigma \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \end{aligned}$$

Instantiating  $K$  with  $m_{u2} + 1 + J'_2$ ,  $\theta'_e$  with  $W''.\theta_2$ ,  $H_2$  with  $H_2'$  (from FB-B1) and  $J_3$  with  $J'_2$ , we get

$$\begin{aligned} & \exists \theta'_2 \sqsupseteq W''.\theta_2. (m_{u2} + 1, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, m_{u2} + 1, v'_2) \in \llbracket \tau' \sigma \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose  $W'$  as  $W_n$  where  $W_n$  is defined as follows:

$$\begin{aligned} W_n.\theta_1 &= \theta'_1 \quad (\text{From (FB-B6)}) \\ W_n.\theta_2 &= \theta'_2 \quad (\text{From (FB-B9)}) \\ W_n.\hat{\beta} &= W''.\hat{\beta} \quad (\text{From (FB-B1)}) \end{aligned}$$

It suffices to prove

- $(k - j, H_1'', H_2'') \triangleright W_n$ :

From Definition 4.9 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'') \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$$

From (FB-B6) we know that  $(m_{u1} + 1, H_1'') \triangleright \theta'_1$  therefore from Definition 4.8 we know that  $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'')$

Similarly from (FB-B9) we know that  $(m_{u2} + 1, H_2'') \triangleright \theta'_2$  therefore from Definition 4.8 we know that  $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that  $(k - f - J, H_1', H_2') \triangleright W''$  therefore from Definition 4.9 we know that  $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq W_n.\theta_1$  and  $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \llbracket W_n.\theta_1(a_1) \rrbracket_V^A):$$

4 cases arise for each  $(a_1, a_2) \in W_n.\hat{\beta}$

$$A. H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) = H_2''(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

We know from that  $(k - f - J, H_1', H_2') \triangleright W''$

Therefore from Definition 4.9 we have

$$\forall (a'_1, a'_2) \in (W''.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$



Since  $W_n.\hat{\beta} = W''.\hat{\beta}$  by construction therefore  
 $\forall(a'_1, a'_2) \in (W_n.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$

From (FB-B6) and (FB-B9) we know that  $W''.\theta_1 \sqsubseteq \theta'_1$  and  $W''.\theta_2 \sqsubseteq \theta'_2$  respectively.

Therefore from Definition 4.1

$$\forall(a'_1, a'_2) \in (W_n.\hat{\beta}).\theta'_1(a_1) = \theta'_2(a_2)$$

To prove:

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A:$$

From (FB-B1) we know that  $(k - f - J, H'_1, H'_2) \triangleright^A W''$

This means from Definition 4.9 we know that

$$\begin{aligned} \forall(a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) &= W''.\theta_2(a_{i2}) \wedge \\ (W'', k - f - J - 1, H'_1(a_{i1}), H'_2(a_{i2})) &\in [W''.\theta_1(a_{i1})]_V^A \end{aligned}$$

Instantiating with  $a_1$  and  $a_2$  and since  $W'' \sqsubseteq W_n$  and  $k - j - 1 < k - f - J - 1$  (since  $j = f + J + J_1 + 1$  therefore from Lemma 4.16 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

B.  $H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) \neq H''_2(a_2)$ :

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B6) and (FB-B9) we know that

$$\begin{aligned} (\forall a. H'_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell') \\ (\forall a. H'_2(a) \neq H''_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell') \end{aligned}$$

This means we have

$$\begin{aligned} \exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell' \text{ and} \\ \exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell' \end{aligned}$$

Since  $\ell_2 \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_3 \not\sqsubseteq \mathcal{A}$ .

Also from (FB-B6) and (FB-B9),  $(m_{u1} + 1, H''_1) \triangleright \theta'_1$  and  $(m_{u2} + 1, H''_2) \triangleright \theta'_2$ .

Therefore from Definition 4.8 we have

$$\begin{aligned} (\theta'_1, m_{u1}, H''_1(a_1)) \in [\theta'_1(a_1)]_V \text{ and} \\ (\theta'_2, m_{u2}, H''_2(a_1)) \in [\theta'_2(a_2)]_V \end{aligned}$$

Since  $m_{u1}$  and  $m_{u2}$  are arbitrary indices therefore from Definition 4.4 we get

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

C.  $H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) \neq H''_2(a_2)$ :

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B9) we know that

$$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell'$$

Since  $\ell_2 \not\sqsubseteq \mathcal{A}$ . Therefore,  $\ell_3 \not\sqsubseteq \mathcal{A}$ .

Since from (FB-B1) we know that  $(k - f - J, H_1', H_2') \triangleright^A W''$  that means from Definition 4.9 that  $(W'', k - f - J - 1, H_1'(a_1), H_2'(a_2)) \in [W''.\theta_1(a_1)]_V^A$ . Since  $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$  and since  $\ell' \not\sqsubseteq \mathcal{A}$  therefore from Definition 4.4 and Definition 4.3 we know that

Therefore

$$\forall m. (W''.\theta_1, m, H_1'(a_1)) \in W''.\theta_1(a_1) \quad (\text{F})$$

Instantiating the (F) with  $m_{u1}$  and using Lemma 4.15 we get

$$(\theta_1', m_{u1}, H_1'(a_1)) \in \theta_1'(a_1)$$

Since from (FB-B9) we know that  $(m_{u2} + 1, H_2'') \triangleright \theta_2'$  therefore from Definition 4.8 we know that  $(\theta_2', m_{u2}, H_2''(a_2)) \in \theta_2'(a_2)$

Therefore from Definition 4.4 we get

$$(W', k - j - 1, H_1''(a_1), H_2''(a_2)) \in [\theta_1'(a_1)]_V^A$$

D.  $H_1'(a_1) \neq H_1''(a_1) \wedge H_2'(a_2) = H_2''(a_2)$ :

Symmetric reasoning as in the previous case

–  $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i''(a_i)) \in [W_n.\theta_i(a_i)]_V$ :

Case  $i = 1$

Given some  $m$  we need to prove

$$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i''(a_i)) \in [W_n.\theta_i(a_i)]_V$$

This further means that given some  $a_1 \in \text{dom}(W_n.\theta_1)$  we need to show

$$(W_n.\theta_1, m, H_1''(a_1)) \in [W_n.\theta_1(a_1)]_V$$

Since  $W_n.\theta_1 = \theta_1'$ , it suffices to prove

$$(\theta_1', m, H_1''(a_1)) \in [\theta_1'(a_1)]_V$$

Like before we apply Theorem 4.21 on  $e_b \gamma \downarrow_1 \cup \{x \mapsto v_1''\}$  but this time at  $m + 1 + J_1 + J_1'$  to get

$$\begin{aligned} & \exists \theta_1' \sqsupseteq W''.\theta_1.(m + 1, H_1'') \triangleright \theta_1' \wedge (\theta_1', m_{u1} + 1, v_1') \in [\tau']_V \wedge \\ & (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_1') \setminus \text{dom}(\theta_e'). \theta_1'(a) \searrow \ell_3) \end{aligned}$$

Since we have  $\ell \sqsubseteq \ell_3$  and  $(m + 1, H_1'') \triangleright \theta_1'$  therefore from Definition 4.8 we get the desired.

Case  $i = 2$

Similar reasoning as in the  $i = 1$  case

•  $(W'.\theta_1, m_{u1}, v_1') \in [\tau']_V \wedge (W'.\theta_2, m_{u2}, v_2') \in [\tau' \sigma]_V$ :

We get this from (FB-B6), (FB-B9) and Lemma 4.15 we get the desired

19. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } (e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

To prove:  $(W, n, \text{new } (e') (\gamma \downarrow_1), \text{new } (e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\begin{aligned} \forall i < n. \text{new } (e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_V^A & \end{aligned}$$

This means that given some  $i < n$  s.t  $\text{new } (e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = \text{new } (e') \gamma \downarrow_1$ ,  $v_{f2} = \text{new } (e') \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, \text{new } (e') \gamma \downarrow_1, \text{new } (e') \gamma \downarrow_2) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_V^A$$

Let  $v_1 = \text{new } (e') \gamma \downarrow_1$  and  $v_2 = \text{new } (e') \gamma \downarrow_2$

From Definition 4.4 we are required to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau)) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove the following:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma): \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also we are given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma)$$

Further from Definition 4.3 it suffices to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau) \sigma]_V^A \quad (\text{FB-R0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e' \gamma \Downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \Downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \sigma \rceil_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose  $W'$  as  $W_n$  where

$$W_n.\theta_1 = W_e.\theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\theta_2 = W_e.\theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\hat{\beta} = W_e.\hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

- i.  $(k - j, H'_1, H'_2) \triangleright W_n$ :

From Definition 4.9 it suffices to prove:

$$\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge$$

$$(W_n, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A) \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$$

This means we need to prove

- $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$ :

We know that  $\text{dom}(W_n.\theta_1) = \text{dom}(W_e.\theta_1) \cup \{a_1\}$  and  $\text{dom}(W_n.\theta_2) = \text{dom}(W_e.\theta_2) \cup \{a_2\}$

Also  $\text{dom}(H'_1) = \text{dom}(H_1) \cup \{a_1\}$  and  $\text{dom}(H'_2) = \text{dom}(H_2) \cup \{a_2\}$

Therefore from  $(k, H_1, H_2) \triangleright W_e$  and from construction of  $W_n$  we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2) \wedge (W_n, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_n.\theta_1(a'_1) \rceil_V^A)$ :

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}).$$

- A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

Since from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$

And since from cg-ref we know that  $H'_1(a_1) = v_{h1}$ ,  $H'_2(a_2) = v'_{h1}$  and

$j = f + 1$  therefore from Lemma 4.16 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

- B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise

- C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise

- D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 4.9

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in \lfloor W_n.\theta_i(a'_i) \rfloor_V$ :

When  $i = 1$

Given some  $m$

$$\forall a'_1 \in \text{dom}(W_n.\theta_1).$$

– when  $a'_1 = a_1$ :

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

And from (FB-R1) we know that  $(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau]_V^A$

Therefore from Lemma 4.14 get the desired

– Otherwise:

Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 4.9

When  $i = 2$

Similar reasoning as with  $i = 1$

ii.  $(W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau) \sigma]_V^A$ :

From cg-ref we know that  $v'_1 = a_1$  and  $v'_2 = a_2$

From Definition 4.4 it suffices to prove

$$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau)$$

This holds from construction of  $W_n$

(b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 4.21 to get

$$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_E$$

This means from Definition 4.7 we get

$$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$$

This further means that given some  $c < k$  s.t  $\text{ref } (e')\gamma \downarrow_1 \downarrow_c v$ . From cg-val we know that  $c = 0$  and  $v = \text{ref } (e')\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$$

From Definition 4.6 we have

$$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$$

$$\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau)]_V \wedge$$

$$(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell)$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

20. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \tau)}$$

To prove:  $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_E^A$

This means from Definition 4.5 we need to prove:

$$\forall i < n. !e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^A$$

This means that given some  $i < n$  s.t  $!e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = !e' \gamma \downarrow_1$ ,  $v_{f2} = !e' \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, !e' \gamma \downarrow_1, !e' \gamma \downarrow_2) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^A$$

Let  $v_1 = !e' \gamma \downarrow_1$  and  $v_2 = !e' \gamma \downarrow_2$

From Definition 4.4 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \tau)) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \tau)): \end{aligned}$$

This means we are given is some  $k \leq n$ ,  $W_e \sqsupseteq W, H_1, H_2$  s.t  $(k, H_1, H_2) \triangleright W_e$

Also given some  $v'_1, v'_2, j < k$  s.t  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \tau))$$

This means from Definition 4.3 it suffices to prove  $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \tau)]_V^A$  (FB-D0)

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{ref } \ell \tau)]_E^A$$

This means from Definition 4.5 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} \implies \\ (W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A$$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A \quad (\text{FB-D1})$$

In order to prove (FB-D0) we choose  $W'$  as  $W_e$ . Also from cg-deref we know that  $H'_1 = H_1$  and  $H'_2 = H_2$ . Also we know that  $v_{h_1} = a_1$  and  $v'_{h_1} = a_2$ .

- $(k - j, H_1, H_2) \triangleright W_e$ :

Since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Lemma 4.20 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

- $(W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \tau)]_V^A$ :

Since from (FB-D1) we know that  $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau]_V^A$

Therefore from Definition 4.4 we know that  $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau$

And since we know that  $(k, H_1, H_2) \triangleright W_e$  therefore from Definition we know that  $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau]_V^A$

Also from cg-ref we know that  $v'_1 = H_1(a_1)$  and  $v'_2 = H_2(a_2)$

From Lemma 4.16 we get  $(W', k - j, H_1(a_1), H_2(a_2)) \in [(\text{Labeled } \ell \tau)]_V^A$

- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)):$

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$$

Since  $(W, n, \gamma) \in [\Gamma]_V^A$  therefore from Lemma 4.23 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 4.21 to get

$$(W.\theta_1, k, (!e'\gamma \downarrow_1) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_E$$

This means from Definition 4.7 we get

$$\forall c < k. !e'\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$$

Instantiating  $c$  with 0 and from cg-val we know that  $v = !e'\gamma \downarrow_1$

And we have  $(W.\theta_1, k, !e'\gamma \downarrow_1) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$

From Definition 4.6 we have

$$\begin{aligned} & \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ & \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \llbracket \text{Labeled } \ell \tau \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \top) \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

21. CG-assign:

$$\frac{\Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_l := e_r : \mathbb{C} \ell \perp \text{unit}}$$

To prove:  $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in \llbracket \mathbb{C} \ell \perp \text{unit } \sigma \rrbracket_E^A$

This means from Definition 4.5 we need to prove:

$$\begin{aligned} & \forall i < n. (e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ & (W, n - i, v_{f1}, v'_{f1}) \in \llbracket \mathbb{C} \ell \perp \text{unit} \rrbracket_V^A \end{aligned}$$

This means that given some  $i < n$  s.t.  $(e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that  $v_{f1} = (e_l := e_r) \gamma \downarrow_1$ ,  $v_{f2} = (e_l := e_r) \gamma \downarrow_2$  and  $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in \llbracket \mathbb{C} \ell \ell \text{unit} \rrbracket_V^A$$

Let  $e_1 = (e_l := e_r) \gamma \downarrow_1$  and  $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 4.4 it suffices to prove

$$\begin{aligned} & \left( \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}) \right) \wedge \\ & \forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \text{unit} \rrbracket_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} & \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}): \end{aligned}$$

This means we are given some  $k \leq n$ ,  $W_e \sqsupseteq W$ ,  $H_1, H_2$  s.t.  $(k, H_1, H_2) \triangleright W_e$

And finally given some  $v'_1, v'_2, j < k$  s.t.  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:



$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit})$   
(FB-A0)

IH1:

$(W_e, k, e_l(\gamma \downarrow_1), e_l(\gamma \downarrow_2)) \in [\text{ref } \ell' \tau]_E^A$

This means from Definition 4.5 we need to prove:

$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies$   
 $(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists f < j < k$  s.t  
 $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$  (FB-A1)

IH2:

$(W_e, k - f, e_r(\gamma \downarrow_1), e_r(\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau]_E^A$

This means from Definition 4.5 we need to prove:

$\forall s < k - f. e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2} \implies$   
 $(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$

Since we know that  $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$  therefore  $\exists s < j - f < k - f$  s.t  $e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$  (FB-A2)

In order to prove (FB-A0) we choose  $W'$  as  $W_e$ . Also from cg-assign we know that  $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$  and  $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$ , and  $j = f + s + 1$

We need to prove the following:

i.  $(k - j, H'_1, H'_2) \triangleright W_e$ :

Say  $v_{h1} = a_1$  and  $v'_{h1} = a_2$

From Definition 4.9 it suffices to prove:

$\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge$

$(W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge$

$\forall (a_1, a_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a_1) = W_e.\theta_2(a_2)) \wedge$

$(W_e, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [\text{ref } \ell' \tau]_V^A \wedge$

$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [W_e.\theta_i]_V$

This means we need to prove

- $\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2))$ :

Since  $\text{dom}(H_1) = \text{dom}(H'_1)$  and  $\text{dom}(H_2) = \text{dom}(H'_2)$ , and also we know that  $(k, H_1, H_2) \triangleright W_e$ . Therefore we obtain the desired directly from Definition 4.9

- $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).(W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_e.\theta_1(a'_1) \rceil_V^A)$   
 $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).$ 
  - A. When  $a'_1 = a_1$  and  $a'_2 = a_2$ :  
From (FB-A1) and from Definition 4.4 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$   
Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$   
And since from cg-assign we know that  $H'_1(a_1) = v_{h2}$ ,  $H'_2(a_2) = v'_{h2}$  and  $j = f + s + 1$  therefore from Lemma 4.16 we get  
 $(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_e.\theta_1(a_1) \rceil_V^A$
  - B. When  $a'_1 = a_1$  and  $a'_2 \neq a_2$ : This case cannot arise
  - C. When  $a'_1 \neq a_1$  and  $a'_2 = a_2$ : This case cannot arise
  - D. When  $a'_1 \neq a_1$  and  $a'_2 \neq a_2$ :  
Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 4.9
- $\forall i \in \{1, 2\}.\forall m.\forall a'_i \in \text{dom}(W_e.\theta_i).(W_e.\theta_i, m, H_i(a'_i)) \in \lfloor W_e.\theta_i(a'_i) \rfloor_V$ :  
When  $i = 1$   
Given some  $m$   
 $\forall a'_1 \in \text{dom}(W_e.\theta_1).$ 
  - when  $a'_1 = a_1$ :  
From (FB-A1) and from Definition 4.4 we get  
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$   
Since from (FB-A2) we know that  $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$   
Therefore from Lemma 4.14 get the desired
  - Otherwise:  
Since  $(k, H_1, H_2) \triangleright W_e$  therefore the desired is obtained directly from Definition 4.9
- When  $i = 2$   
Similar reasoning as with  $i = 1$

ii.  $\text{ValEq}(\mathcal{A}, W_e, k - j, \perp, (), (), \text{unit})$ :

Holds directly from Definition 4.3 and Definition 4.4

- (b)  $\forall l \in \{1, 2\}.\left(\forall k, \theta_e \sqsupseteq \theta, H, j, (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$   
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \text{unit} \rfloor_V \wedge$   
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$ :

Case  $l = 1$

Given some  $k, \theta_e \sqsupseteq W.\theta_l, H, j$  s.t  $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor (\text{unit}) \rfloor_V \wedge$   
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$

Since  $(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$  therefore from Lemma 4.23 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$  and  $(W.\theta_2, m, \gamma \downarrow_2) \in \lfloor \Gamma \rfloor_V$

Instantiating  $m$  with  $k$  we get  $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 4.21 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{unit}))]_E$$

This means from Definition 4.7 we get

$$\forall c < k.(e_l := e_r)\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{unit}))]_V$$

Instantiating  $c$  with 0 and from cg-val we know that  $v = (e_l := e_r)\gamma \downarrow_1$

And we have  $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit}))]_V$

From Definition 4.6 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a.H_1(a) \neq H'(a) \implies \exists \ell'.\theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e).\theta'(a) \searrow \ell') \end{aligned}$$

Instantiating  $K$  with  $k$ ,  $\theta'_e$  with  $\theta_e$ ,  $H_1$  with  $H$  and  $J$  with  $j$  we get the desired

Case  $l = 2$

Symmetric reasoning as in the  $l = 1$  case above

□

**Lemma 4.25.**  $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies$$

$$\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$$

*Proof.* Given that  $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$ . From Definition 4.3 two cases arise

1.  $\ell \sqsubseteq \mathcal{A}$ :

In this case we know that  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

2 cases arise

(a)  $\ell' \sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 4.16 we know that  $(W', j, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

And thus from Definition 4.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

Since  $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$  therefore from Lemma 4.14 we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 4.15 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 4.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2.  $\ell \not\sqsubseteq \mathcal{A}$ :

Given is  $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that  $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 4.15 we know that  $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 4.3 we know that  $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

**Lemma 4.26** (Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

$$1. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$$

$$2. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^A \subseteq [(\tau' \sigma)]_E^A$$

*Proof.* Proof of statement (1)

Proof by induction on the  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

IH1:  $[(\tau'_1 \sigma)]_V^A \subseteq [(\tau_1 \sigma)]_V^A$  (Statement 1)

$[(\tau_2 \sigma)]_E^A \subseteq [(\tau'_2 \sigma)]_E^A$  (Sub-A0 From Statement 2)

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

This means that given:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A$

And it suffices to prove:  $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

From Definition 4.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2 \sigma]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 4.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2 \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E) \end{aligned}$$

This means need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A) :$$

Given:  $W'' \sqsupseteq W, k < n$  and  $v'_1, v'_2$ . We are also given  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove:  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with  $W'', k, v'_1$  and  $v'_2$  we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (101)$$

Since  $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$  therefore from IH1 we know that  $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 101 we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using (Sub-A0) we get  $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

(b)  $\forall \theta'_i \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_i, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E)$ :

Given:  $\theta'_i \sqsupseteq W.\theta_1, k, v'_c$ . We are also given  $(\theta'_i, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove:  $(\theta'_i, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$

Since we are given  $(\theta'_i, k, v'_c) \in [\tau'_1 \sigma]_V$  and since  $\tau'_1 <: \tau_1$  therefore from Lemma 4.22 we get

$$(\theta'_i, k, v'_c) \in [\tau_1 \sigma]_V \quad (102)$$

Instantiating the second conjunct of Sub-A1 with  $\theta'_i, k, v'_1$  and  $v'_2$  we get

$$((\theta'_i, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_i, e_1[v'_c/x]) \in [\tau_2 \sigma]_E) \quad (103)$$

Therefore from Equation 102 and 103 we get  $(\theta'_i, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E$

Since  $\tau_2 <: \tau'_2$  therefore from Lemma 4.22 we get

$$(\theta'_i, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$$

(c)  $\forall \theta'_i \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_i, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E)$ :

Similar reasoning as in the previous case

## 2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove:  $[((\tau_1 \times \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

IH1:  $[(\tau_1 \sigma)]_V^A \subseteq [(\tau'_1 \sigma)]_V^A$  (Statement (1))

IH2:  $[(\tau_2 \sigma)]_V^A \subseteq [(\tau'_2 \sigma)]_V^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

This means that given:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A$

Therefore from Definition 4.4 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1 \sigma]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2 \sigma]_V^A \quad (104)$$

And it suffices to prove:  $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

Again from Definition 4.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_V^A \wedge (W, n, v_2, v'_2) \in [\tau'_2 \sigma]_V^A$$

Since from Equation 104 we know that  $(W, n, v_1, v'_1) \in [\tau_1 \sigma]_V^A$  therefore from IH1 we have  $(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_V^A$

Similarly since  $(W, n, v_2, v'_2) \in [\tau_2 \sigma]_V^A$  from Equation 104 therefore from IH2 we have  $(W, n, v_2, v'_2) \in [\tau'_2 \sigma]_V^A$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $\lceil ((\tau_1 + \tau_2) \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_{\mathcal{V}}^A$

IH1:  $\lceil (\tau_1 \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil (\tau'_1 \sigma) \rceil_{\mathcal{V}}^A$  (Statement (1))

IH2:  $\lceil (\tau_2 \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil (\tau'_2 \sigma) \rceil_{\mathcal{V}}^A$  (Statement (1))

It suffices to prove:  $\forall (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_{\mathcal{V}}^A$

This means that given:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2) \sigma) \rceil_{\mathcal{V}}^A$

And it suffices to prove:  $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2) \sigma) \rceil_{\mathcal{V}}^A$

2 cases arise

(a)  $v_{s1} = \text{inl } v_{i1}$  and  $v_{s2} = \text{inl } v_{i2}$ :

From Definition 4.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1 \sigma \rceil_{\mathcal{V}}^A \tag{105}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_{\mathcal{V}}^A$$

From Equation 105 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_1 \sigma \rceil_{\mathcal{V}}^A$$

(b)  $v_{s1} = \text{inr } v_{i1}$  and  $v_{s2} = \text{inr } v_{i2}$ :

From Definition 4.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2 \sigma \rceil_{\mathcal{V}}^A \tag{106}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_{\mathcal{V}}^A$$

From Equation 106 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau'_2 \sigma \rceil_{\mathcal{V}}^A$$

4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $\lceil ((\forall \alpha. \tau_1) \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil (\forall \alpha. \tau_2) \sigma \rceil_{\mathcal{V}}^A$

$\forall \sigma. \lceil (\tau_1 \sigma) \rceil_{\mathcal{E}}^A \subseteq \lceil (\tau_2 \sigma) \rceil_{\mathcal{E}}^A$  (Sub-F2, From Statement (2))

It suffices to prove:  $\forall (W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. \tau_1) \sigma) \rceil_{\mathcal{V}}^A.$

$$(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. \tau_2) \sigma) \rceil_{\mathcal{V}}^A$$

This means that given:  $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\tau_1)) \sigma) \rceil_{\mathcal{V}}^A$

Therefore from Definition 4.4 we are given:

$$\begin{aligned}
& \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E) \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E) \quad (\text{Sub-F1})
\end{aligned}$$

And it suffices to prove:  $(W, n, \Lambda e_1, \Lambda e_2) \in [((\forall \alpha. \tau_2) \sigma)]_V^A$

Again from Definition 4.4, it suffices to prove:

$$\begin{aligned}
& \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge \\
& \forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E) \wedge \\
& \forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E)
\end{aligned}$$

This means we are required to show:

$$(a) \quad \forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A):$$

By instantiating the first conjunct of Sub-F1 with  $W''$ ,  $n''$  and  $\ell''$  we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from Sub-F2 instantiated at  $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

$$(b) \quad \forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E):$$

By instantiating the second conjunct of Sub-F1 with  $\theta'_l$  and  $\ell''$  we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E)$$

Since  $\tau_1 \sigma \cup \{\alpha \mapsto \ell''\} <: \tau_2 \sigma \cup \{\alpha \mapsto \ell''\}$  therefore from Lemma 4.22 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E)$$

$$(c) \quad \forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E):$$

Similar reasoning as in the previous case

## 5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove:  $[((c_1 \Rightarrow \tau_1) \sigma)]_V^A \subseteq [((c_2 \Rightarrow \tau_2)) \sigma]_V^A$

$$[(\tau_1 \sigma)]_E^A \subseteq [(\tau_2 \sigma)]_E^A \text{ (Sub-C0, From Statement (2))}$$

It suffices to prove:  $\forall (W, n, \nu e_1, \nu e_2) \in [((c_1 \Rightarrow \tau_1) \sigma)]_V^A. (W, n, \nu e_1, \nu e_2) \in [((c_2 \Rightarrow \tau_2) \sigma)]_V^A$

This means that given:  $(W, n, \nu e_1, \nu e_2) \in [((c_1 \Rightarrow \tau_1) \sigma)]_V^A$

Therefore from Definition 4.4 we are given:

$$\begin{aligned}
& \forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma \implies (W', n', e_1, e_2) \in [\tau_1 \sigma]_E^A \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_1) \in [\tau_1 \sigma]_E \wedge \\
& \forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_2) \in [\tau_1 \sigma]_E \quad (\text{Sub-C1})
\end{aligned}$$

And it suffices to prove:  $(W, n, \nu e_1, \nu e_2) \in [((c_2 \Rightarrow \tau_2) \sigma)]_V^A$

Again from Definition 4.4, it suffices to prove:

$$\begin{aligned}
\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma &\implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A \wedge \\
\forall \theta'_i \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 &\implies (\theta'_i, j, e_1) \in \lfloor \tau_2 \sigma \rfloor_E \wedge \\
\forall \theta'_i \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 &\implies (\theta'_i, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E
\end{aligned}$$

This means that we are required to show the following:

$$(a) \quad \forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A:$$

We are given  $W'' \sqsupseteq W, n'' < n$  also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with  $W''$  and  $n''$  we know that the following holds

$$(W'', n'', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A$$

Therefore from (Sub-C0) we get  $(W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$

$$(b) \quad \forall \theta'_i \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \implies (\theta'_i, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E:$$

We are given some  $\theta'_i \sqsupseteq W.\theta_1, k$ , also we know that  $\mathcal{L} \models c_2 \sigma$  and  $c_2 \sigma \implies c_1 \sigma$  therefore we also know that  $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with  $\theta'_i$  we know that the following holds

$$(\theta'_i, k, e_1) \in \lfloor \tau_1 \sigma \rfloor_E$$

Since  $\tau_1 \sigma <: \tau_2 \sigma$  therefore from Lemma 4.22 we get

$$(\theta'_i, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E$$

$$(c) \quad \forall \theta'_i \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_i, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E:$$

Similar reasoning as in the previous case

## 6. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\lceil ((\text{Labeled } \ell \tau) \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil ((\text{Labeled } \ell' \tau') \sigma) \rceil_{\mathcal{V}}^A$

IH:  $\lceil (\tau \sigma) \rceil_{\mathcal{V}}^A \subseteq \lceil (\tau' \sigma) \rceil_{\mathcal{V}}^A$

It suffices to prove:  $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil ((\text{Labeled } \ell \tau) \sigma) \rceil_{\mathcal{V}}^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil ((\text{Labeled } \ell' \tau') \sigma) \rceil_{\mathcal{V}}^A$

This means we are given  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil ((\text{Labeled } \ell \tau) \sigma) \rceil_{\mathcal{V}}^A$

From Definition 4.4 it means we have  $\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau \sigma)$  (Sub-L0)

and it suffices to prove  $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil ((\text{Labeled } \ell' \tau') \sigma) \rceil_{\mathcal{V}}^A$

Again from Definition 4.4 it means we need to prove that

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$$

Since we have (Sub-L0) and  $\ell \sqsubseteq \ell'$  therefore from Lemma 4.25 we have

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}(v_2), \tau \sigma)$$

2 cases arise:



(a)  $\ell' \sqsubseteq \mathcal{A}$ :

In this case from Definition 4.3 we know that  $(W, n, v_1, v_2) \in [\tau \sigma]_V^{\mathcal{A}}$

From IH we also know that  $(W, n, v_1, v_2) \in [\tau' \sigma]_V^{\mathcal{A}}$

And from Definition 4.4 we get  $ValEq(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau' \sigma)$

(b)  $\ell' \not\sqsubseteq \mathcal{A}$ :

In this case from Definition 4.3 we know that  $\forall j. (W.\theta_1, j, v_1) \in [\tau \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau \sigma]_V$

Since  $\tau <: \tau'$  therefore from Lemma 4.22 we get  $(W.\theta_1, j, v_1) \in [\tau' \sigma]_V$  and  $(W.\theta_2, j, v_2) \in [\tau' \sigma]_V$

And from Definition 4.4 we get  $ValEq(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau' \sigma)$

7. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell'_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove:  $[((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^{\mathcal{A}} \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^{\mathcal{A}}$

IH:  $[(\tau \sigma)]_V^{\mathcal{A}} \subseteq [(\tau' \sigma)]_V^{\mathcal{A}}$

It suffices to prove:  $\forall (W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^{\mathcal{A}}. (W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^{\mathcal{A}}$

This means we are given  $(W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^{\mathcal{A}}$

From Definition 4.4 it means we have

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \quad (\text{Sub-CG0}) \end{aligned}$$

And we need to prove

$$(W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^{\mathcal{A}}$$

Again from Definition 4.4 it means we need to prove

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau' \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)) \end{aligned}$$

It means we need to prove:

- (a)  $\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j.$   
 $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$   
 $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma):$

This means we are given  $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$  s.t  
 $(k, H_1, H_2) \triangleright W_e, (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2)$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau' \sigma)$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1)  $\text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma)$

Therefore from Lemma 4.25 we get  $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau \sigma)$

- (b)  $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$   
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \sigma \rfloor_V \wedge$   
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge$   
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i):$

Case  $l = 1$

Here we are given  $k, \theta_e \sqsupseteq \theta, H, j < k$  s.t  $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

- i.  $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau' \sigma \rfloor_V:$

Instantiating the second conjunct of (Sub-CG0) with the given  $k, \theta_e, H, j$  to get  
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau \sigma \rfloor_V$

Since  $\tau <: \tau'$  therefore from Lemma 4.22 we get  $(\theta', k - j, v'_1) \in \lfloor \tau' \sigma \rfloor_V$

- ii.  $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')::$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell')$$

Since  $\ell'_i \sqsubseteq \ell_i$  therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')$$

- iii.  $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)::$

Instantiating the second conjunct of (Sub-CG0) with the given  $v, i, k, \theta_e, H, j$  to get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)$$

Since  $\ell'_i \sqsubseteq \ell_i$  therefore we also get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

Case  $l = 2$

Symmetric reasoning as in the previous  $l = 1$  case

## 8. CGsub-base:

Trivial

### Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [(\tau \sigma)]_E^A. (W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$$

This means given  $(W, n, e_1, e_2) \in [(\tau \sigma)]_E^A$

From Definition 4.5 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau \sigma]_V^A \quad (\text{Sub-E0})$$

And it suffices to prove  $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

Again from Definition 4.5 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$$

This means that given  $i < n$  s.t  $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$  we need to prove  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$

Instantiating (Sub-E0) with the given  $i$  we get  $(W, n - i, v_1, v_2) \in [\tau \sigma]_V^A$

From Statement (1) we get  $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$  □

**Theorem 4.27** (NI for CG). *Say*  $\text{bool} = (\text{unit} + \text{unit})$

$$\forall v_1, v_2, e, n'.$$

$$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$$

$$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2) \implies v'_1 = v'_2$$

*Proof.* Given some

$$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$$

$$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2)$$

And we need to prove

$$v'_1 = v'_2$$

From Theorem 4.24 we know that

$$\forall n. (\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \text{ bool}]_E^\perp$$

Similarly from Theorem 4.24 and Definition 4.13 we also get

$$\forall n. (\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\perp$$

From Definition 4.5 we get

$$\forall n. \forall i < n. e[v_1/x] \Downarrow_i v_{11} \wedge e[v_2/x] \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\perp$$

Instantiating it with  $n' + 1$  and then with 0, from CG-val we have  $v_{11} = e[v_1/x]$  and  $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n' + 1, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\perp$$

From Definition 4.6 we have

$$\left( \forall k \leq n' + 1, W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right.$$

$$\forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow^f (H_2', v_2'') \wedge j < k \implies$$

$$\exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \Big) \wedge$$

$$\forall l \in \{1, 2\}. \left( \forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right.$$

$$\begin{aligned}
& \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_i) \in [b]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp)
\end{aligned}$$

Instantiating the first conjunct with  $n' + 1, \emptyset, \emptyset, \emptyset$ . And then with  $v'_1, v'_2, n'$  we get  $\exists W' \sqsupseteq \emptyset.(1, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\perp, W', 1, \perp, v'_1, v'_2, \text{bool})$

From Definition 4.3 and Definition 4.6 we get  $v'_1 = v'_2$

□

## 5 Translations between FG and CG

### 5.1 CG to FG translation

#### 5.1.1 Type directed translation from CG to FG

CG types are translated into FG types by the following definition of  $\llbracket \cdot \rrbracket$

$$\begin{array}{ll}
 \llbracket \mathbf{b} \rrbracket = \mathbf{b}^\perp & \llbracket \text{ref } \ell \tau \rrbracket = (\text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp \\
 \llbracket \tau_1 \rightarrow \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket = (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp \\
 \llbracket \tau_1 \times \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp & \llbracket c \Rightarrow \tau \rrbracket = (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \\
 \llbracket \tau_1 + \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp & \llbracket \forall \alpha. \tau \rrbracket = (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \\
 \llbracket \text{Labeled } \ell \tau \rrbracket = (\llbracket \tau \rrbracket + \text{unit})^\ell &
 \end{array}$$

The translation judgment for expressions is of the form  $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$ .

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \lambda x. e : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_F} \text{ lambda} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau \rightsquigarrow e_{F1} e_{F2}} \text{ app} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : \tau_1 \times \tau_2 \rightsquigarrow (e_{F1}, e_{F2})} \text{ prod} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e) : \tau_1 \rightsquigarrow \text{fst}(e_F)} \text{ fst} \qquad \frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{snd}(e) : \tau_2 \rightsquigarrow \text{snd}(e_F)} \text{ snd} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inl}(e_F)} \text{ inl} \qquad \frac{\Sigma; \Psi; \Gamma \vdash e : \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inr}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inr}(e_F)} \text{ inr} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : (\tau_1 + \tau_2) \rightsquigarrow e_F \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{case}(e_F, x.e_{F1}, y.e_{F2})} \text{ case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{ label} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_. e_F} \text{ unlabel} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_. \text{inl}(e_F ())} \text{ toLabeled} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ ret} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_4 \quad \Sigma; \Psi \vdash \ell_4 \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_. \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{ bind} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_. \text{inl}(\text{new } (e_F))} \text{ ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_. \text{inl}(e_{F1} := e_{F2})} \text{ assign}
\end{array}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F} \text{sub} \qquad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \Lambda e : \forall \alpha. \tau \rightsquigarrow \Lambda e_F} \text{FI} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \rightsquigarrow e_F \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e \llbracket : \tau[\ell/\alpha] \rightsquigarrow e_F \rrbracket} \text{FE} \qquad \frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \nu e : c \Rightarrow \tau \rightsquigarrow \nu e_F} \text{CI} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau \rightsquigarrow e_F \bullet} \text{CE}
\end{array}$$

### 5.1.2 Type preservation for CG to FG translation

**Theorem 5.1** (Type preservation,  $\text{CG} \rightsquigarrow \text{FG}$ ).  $\forall \Sigma; \Psi; \Gamma, e_C, \tau$ .

$\Gamma \vdash e_C : \tau$  is a valid typing derivation in CG  $\implies$

$\exists e_F$ .

$\Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$

$\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket$  is a valid typing derivation in FG

*Proof.* Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\begin{array}{c}
\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{Lb}_{\ell}(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{label} \\
\frac{\frac{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket}{\llbracket \Gamma \rrbracket \vdash_{\top} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\perp}} \text{IH}}{\llbracket \Gamma \rrbracket \vdash_{\top} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell}} \text{FG-inl} \\
\text{FG-sub}
\end{array}$$

2. unlabel:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{\cdot} e_F} \text{unlabel}$$

Main derivation:

$$\frac{\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \lambda_{\cdot} e_F : (\text{unit} \xrightarrow{\top} (\llbracket \tau \rrbracket + \text{unit})^{\ell})^{\perp}} \text{IH}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \lambda_{\cdot} e_F : (\text{unit} \xrightarrow{\top} (\llbracket \tau \rrbracket + \text{unit})^{\ell})^{\perp}} \text{FG-lam}$$

3. toLabeled:

$$\frac{\Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_{\cdot} \text{inl}(e_F ())} \text{toLabeled}$$

P2:

$$\frac{\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{FG-sub}$$

P1:

$$P2 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} () : \text{unit}} \quad \mathcal{L} \vdash \ell_1 \sqcup \perp \sqsubseteq \ell_1 \quad \mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{FG-app}$$

Main derivation:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F()) : ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp} \text{FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(e_F()) : (\text{unit} \xrightarrow{\ell_1} ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp)^\perp} \text{FG-lam}}$$

4. ret:

$$\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{-}.\text{inl}(e_F)} \text{ret}$$

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : \llbracket \tau \rrbracket} \text{IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : \llbracket \tau \rrbracket} \text{FG-sub} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{FG-sub, FG-inl}$$

$$\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(e_F) : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{FG-lam}$$

5. bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{bind}$$

P1.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{FG-sub}$$

P1:

$$P1.1 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} () : \text{unit}} \text{FG-var} \quad \mathcal{L} \vdash (\ell \sqcup \perp) \sqsubseteq \ell_1 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_{F1}() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{FG-app}$$

P2.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^\perp} \text{IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqcup \ell_2 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^\perp} \text{FG-sub}$$



P2:

$$\begin{array}{c}
P2.1 \quad \frac{}{\llbracket \Gamma \rrbracket, - : \mathbf{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} () : \mathbf{unit}} \text{FG-var} \\
\mathcal{L} \vdash (\ell \sqcup \ell_2 \sqcup \perp) \sqsubseteq \ell_3 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell_4} \searrow \perp} \\
\hline
\llbracket \Gamma \rrbracket, - : \mathbf{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2}() : (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell_4} \text{FG-app}
\end{array}$$

P3:

$$\frac{\frac{}{\llbracket \Gamma \rrbracket, - : \mathbf{unit}, y : \mathbf{unit} \vdash_{\ell \sqcup \ell_2} () : \mathbf{unit}} \text{FG-var} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\llbracket \Gamma \rrbracket, - : \mathbf{unit}, y : \mathbf{unit} \vdash_{\ell \sqcup \ell_2} \mathbf{inr}() : (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell_4}} \text{FG-sub, FG-inr}$$

Main derivation:

$$\begin{array}{c}
P1 \quad P2 \quad P3 \quad \frac{\frac{}{\mathcal{L} \vdash \ell_2 \sqsubseteq \ell_4} \text{Given}}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell_4} \searrow \ell_2} \quad \frac{}{\ell_4 \sqsubseteq \ell'} \text{Given} \\
\hline
\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\ell} \mathbf{case}(e_{F1}(), x.e_{F2}(), y.\mathbf{inr}()) : (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell'} \text{FG-case, FG-sub} \\
\hline
\llbracket \Gamma \rrbracket \vdash_{\top} \lambda\_.\mathbf{case}(e_{F1}(), x.e_{F2}(), y.\mathbf{inr}()) : (\mathbf{unit} \xrightarrow{\ell} (\llbracket \tau' \rrbracket + \mathbf{unit})^{\ell'})^{\perp} \text{FG-lam}
\end{array}$$

6. ref:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \mathbf{new } e : \mathbb{C} \ell \perp (\mathbf{ref } \ell' \tau) \rightsquigarrow \lambda\_.\mathbf{inl}(\mathbf{new } (e_F))} \text{ref}$$

P1:

$$\begin{array}{c}
\frac{\frac{}{\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\top} e_F : (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'}}{\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\ell} e_F : (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'}} \text{IH, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\ell} e_F : (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'}} \text{FG-sub} \\
\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'} \searrow \ell} \\
\hline
\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\ell} \mathbf{new } e_F : (\mathbf{ref}(\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'})^{\perp} \text{FG-ref}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P1 \\
\frac{\frac{}{\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\ell} \mathbf{inl}(\mathbf{new } e_F) : ((\mathbf{ref}(\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'})^{\perp} + \mathbf{unit})^{\perp}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda\_.\mathbf{inl}(\mathbf{new } e_F) : (\mathbf{unit} \xrightarrow{\ell} ((\mathbf{ref}(\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'})^{\perp} + \mathbf{unit})^{\perp})^{\perp}} \text{FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda\_.\mathbf{inl}(\mathbf{new } e_F) : (\mathbf{unit} \xrightarrow{\ell} ((\mathbf{ref}(\llbracket \tau \rrbracket + \mathbf{unit})^{\ell'})^{\perp} + \mathbf{unit})^{\perp})^{\perp}} \text{FG-lam}
\end{array}$$

7. deref:

$$\frac{\Gamma \vdash e : \mathbf{ref } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda\_.\mathbf{inl}(e_F)} \text{deref}$$

P2:

$$\frac{}{\llbracket \Gamma \rrbracket, - : \mathbf{unit} \vdash_{\top} e_F : (\mathbf{ref}(\llbracket \tau \rrbracket + \mathbf{unit})^{\ell})^{\perp}} \text{IH}$$

P1:

$$\frac{P2 \quad \frac{\mathcal{L} \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell \quad \text{Lemma 1.1} \quad \mathcal{L} \vdash ([\tau] + \text{unit})^\ell \searrow \perp}{\mathcal{L} \vdash ([\tau] + \text{unit})^\ell \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} !e_F : ([\tau] + \text{unit})^\ell} \text{FG-deref}$$

Main derivation:

$$\frac{\frac{P1 \quad \frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^\perp}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^\perp} \text{FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\top} (([\tau] + \text{unit})^\ell + \text{unit})^\perp)} \text{FG-lam}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^\perp} \text{FG-lam}}$$

8. assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{F1} := e_{F2})} \text{assign}$$

P3:

$$\frac{\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F2} : ([\tau] + \text{unit})^{\ell'} \quad \text{IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}$$

P2:

$$\frac{\frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp \quad \text{IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{FG-sub}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{FG-sub}$$

P1:

$$\frac{P2 \quad P3 \quad \frac{\mathcal{L} \vdash \ell \sqsubseteq \ell' \quad \text{Given}}{\mathcal{L} \vdash ([\tau] + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp)} \text{FG-assign}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}} \text{FG-assign}$$

Main derivation:

$$\frac{\frac{P1 \quad \frac{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\perp}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\perp} \text{FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^\perp)} \text{FG-lam}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\perp} \text{FG-lam}}$$

9. sub:

$$\frac{\frac{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : [\tau'] \quad \text{IH} \quad \mathcal{L} \vdash \top \sqsubseteq \top \quad \frac{\mathcal{L} \vdash \tau' <: \tau}{\mathcal{L} \vdash [\tau'] <: [\tau]} \text{Lemma 5.2}}{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : [\tau]} \text{FG-sub}}{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : [\tau]} \text{FG-sub}$$

10. FI:

$$\frac{\frac{\Sigma, \alpha; \Psi; \llbracket \Gamma \rrbracket \vdash_{\top} e_F : [\tau]}{\Sigma, \alpha; \Psi; \llbracket \Gamma \rrbracket \vdash_{\top} e_F : [\tau]} \text{IH}}{\Sigma; \Psi; \llbracket \Gamma \rrbracket \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp} \text{FG-FI}}{\Sigma; \Psi; \llbracket \Gamma \rrbracket \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp} \text{FG-FI}}$$

11. FE:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (\forall \alpha. (\top, [\tau]))^{\perp}} \text{ IH}}{\text{FV}(\ell) \in \Sigma \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau[\ell/\alpha]] \searrow \perp} \text{ FG-FE}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F [] : [\tau][\ell/\alpha]} \text{ Lemma 5.5}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F [] : [\tau][\ell/\alpha]}$$

12. CI:

$$\frac{\overline{\Sigma; \Psi, c; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \nu e_F : (c \Rightarrow [\tau])^{\perp}} \text{ FG-CI}$$

13. CE:

$$\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (c \Rightarrow [\tau])^{\perp}} \text{ IH} \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau] \searrow \perp}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F \bullet : [\tau]} \text{ FG-CE}$$

□

**Lemma 5.2** (Subtyping type preservation: CG to FG). *For any CG types  $\tau$  and  $\tau'$ ,  $\Sigma$ , and  $\Psi$ , if  $\mathcal{L} \vdash \tau <: \tau'$ , then  $\mathcal{L} \vdash [\tau] <: [\tau']$ .*

*Proof.* Proof by induction on CG's subtyping relation

1. CGsub-base:

$$\overline{\mathcal{L} \vdash [\tau] <: [\tau]} \text{ Lemma 1.1}$$

2. CGsub-arrow:

$$\frac{\overline{\mathcal{L} \vdash [\tau'_1] <: [\tau_1]} \text{ IH1} \quad \overline{\mathcal{L} \vdash [\tau'_2] <: [\tau'_2]} \text{ IH2} \quad \mathcal{L} \vdash \top \sqsubseteq \top}{\mathcal{L} \vdash ([\tau_1] \xrightarrow{\top} [\tau_2])^{\perp} <: ([\tau'_1] \xrightarrow{\top} [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\mathcal{L} \vdash ([\tau_1] \xrightarrow{\ell_{\xi}} \tau_2) <: ([\tau'_1] \xrightarrow{\ell'_{\xi}} \tau'_2)} \text{ Definition of } [\cdot]$$

3. CGsub-prod:

$$\frac{\overline{\mathcal{L} \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\mathcal{L} \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\mathcal{L} \vdash ([\tau_1] \times [\tau_2])^{\perp} <: ([\tau'_1] \times [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\mathcal{L} \vdash [(\tau_1 \times \tau_2)] <: [(\tau'_1 \times \tau'_2)]} \text{ Definition of } [\cdot]$$

4. CGsub-sum:

$$\frac{\overline{\mathcal{L} \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\mathcal{L} \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\mathcal{L} \vdash ([\tau_1] + [\tau_2])^{\perp} <: ([\tau'_1] + [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\mathcal{L} \vdash [(\tau_1 + \tau_2)] <: [(\tau'_1 + \tau'_2)]} \text{ Definition of } [\cdot]$$

5. CGsub-labeled:

$$\begin{array}{c}
\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket} \text{ IH1} \quad \overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{ FGsub-unit}}{\overline{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})}} \text{ FGsub-sum} \\
\frac{\overline{\text{Labeled } \ell_1 \tau_1 <: \text{Labeled } \ell'_1 \tau'_1} \text{ Given}}{\ell_1 \sqsubseteq \ell'_1} \text{ By inversion} \\
\hline
\overline{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_1}} \text{ FGsub-arrow} \\
\hline
\overline{\mathcal{L} \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket} \text{ Definition of } \llbracket \cdot \rrbracket
\end{array}$$

6. CGsub-monad:

P3:

$$\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket} \text{ IH} \quad \overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{ FGsub-unit}}{\overline{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})}} \text{ FGsub-sum}$$

P2:

$$\begin{array}{c}
P3 \quad \frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1} \text{ Given}}{\mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o} \text{ By inversion} \\
\hline
\overline{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o}} \text{ FGsub-label}
\end{array}$$

P1:

$$\frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \quad P2 \quad \frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1} \text{ Given}}{\mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i}}{\overline{\mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o}) <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})}} \text{ FGsub-arrow}$$

Main derivation:

$$\frac{P1 \quad \overline{\mathcal{L} \vdash \perp \sqsubseteq \perp}}{\overline{\mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})^\perp}} \text{ FGsub-label} \\
\hline
\overline{\mathcal{L} \vdash \llbracket \mathbb{C} \ell_i \ell_o \tau_1 \rrbracket <: \llbracket \mathbb{C} \ell'_i \ell'_o \tau'_1 \rrbracket} \text{ Definition of } \llbracket \cdot \rrbracket$$

7. SLIO\*sub-forall:

P1:

$$\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket <: \llbracket \tau' \rrbracket} \text{ IH, Weakening} \quad \overline{\Sigma, \alpha; \Psi \vdash \top \sqsubseteq \top}}{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))}} \text{ FGsub-forall}$$

Main derivation:

$$\frac{P1 \quad \overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))^\perp}} \text{ FGsub-label} \\
\hline
\overline{\Sigma; \Psi \vdash \llbracket \forall \alpha. \tau \rrbracket <: \llbracket \forall \alpha. \tau' \rrbracket}$$

8. SLIO\*sub-constraint:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket} <: \llbracket \tau' \rrbracket}}{\text{IH}} \quad \frac{\overline{\Sigma; \Psi \vdash \top \sqsubseteq \top} \quad \frac{\overline{\Sigma; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'}}{\text{Given}} \quad \text{By inversion}}{\overline{\Sigma; \Psi \vdash c' \Rightarrow c}} \quad \text{FGsub-constra}}{\overline{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket) <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)}} \quad \text{FGsub-constra}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}}}{\overline{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket)^\perp <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)^\perp}} \quad \text{FGsub-label}}{\overline{\Sigma; \Psi \vdash \llbracket c \Rightarrow \tau \rrbracket <: \llbracket c' \Rightarrow \tau' \rrbracket}} \quad \text{FGsub-label}$$

□

**Lemma 5.3** (CG  $\rightsquigarrow$  FG: Preservation of well-formedness).  $\forall \Sigma, \Psi, \tau$ .

$$\Sigma; \Psi \vdash \tau \text{ WF} \implies \Sigma; \Psi \vdash \llbracket \tau \rrbracket \text{ WF}$$

*Proof.* Proof by induction on the  $\tau$  WF relation.

1. CG-wff-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \quad \text{FG-wff-base}}{\overline{\Sigma; \Psi \vdash \mathbf{b}^\perp \text{ WF}} \quad \text{FG-wff-label}}$$

2. CG-wff-unit:

$$\overline{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \quad \text{FG-wff-unit}$$

3. CG-wff-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \quad \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \quad \text{IH2}}{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket) \text{ WF}} \quad \text{FG-wff-arrow}}}{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \text{ WF}} \quad \text{FG-wff-label}}$$

4. CG-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \quad \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \quad \text{IH2}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket \text{ WF}} \quad \text{FG-wff-prod}}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket^\perp \text{ WF}} \quad \text{FG-wff-label}}$$

5. CG-wff-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket \text{ WF}} \quad \text{IH1} \quad \overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket \text{ WF}} \quad \text{IH2}}{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket \rrbracket \text{ WF}} \quad \text{FG-wff-prod}}}{\overline{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket) \rrbracket^\perp \text{ WF}} \quad \text{FG-wff-label}}$$

6. CG-wff-ref:

$$\begin{array}{c}
\frac{\overline{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ Given}}{\text{FV}(\tau) = \emptyset} \text{ By inversion} \\
\hline
\text{FV}(\llbracket \tau \rrbracket) = \emptyset \text{ Lemma 5.4} \\
\hline
\frac{\overline{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ Given}}{\text{FV}(\ell) = \emptyset} \text{ By inversion} \\
\hline
\text{FV}(\text{unit}) = \emptyset \\
\hline
\Sigma; \Psi \vdash \text{FV}((\llbracket \tau \rrbracket + \text{unit})^\ell) = \emptyset \\
\hline
\Sigma; \Psi \vdash \text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF \text{ FG-wff-ref} \\
\hline
\Sigma; \Psi \vdash (\text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp \ WF \text{ FG-wff-label}
\end{array}$$

7. CG-wff-forall:

$$\begin{array}{c}
\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH}}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) \ WF} \text{ FG-wff-forall} \\
\hline
\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \ WF \text{ CG-wff-label}
\end{array}$$

8. CG-wff-constraint:

$$\begin{array}{c}
\frac{\overline{\Sigma; \Psi, c \vdash \llbracket \tau \rrbracket \ WF} \text{ IH}}{\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket) \ WF} \text{ FG-wff-constraint} \\
\hline
\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \ WF \text{ CG-wff-label}
\end{array}$$

9. CG-wff-labeled:

$$\begin{array}{c}
\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \overline{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit}}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF} \text{ FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF \text{ CG-wff-label}
\end{array}$$

10. CG-wff-monad:

P1:

$$\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \overline{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit}}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF} \text{ FG-wff-sum}$$

Main derivation:

$$\begin{array}{c}
\frac{\overline{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \quad \frac{P1}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \ WF} \text{ FG-wff-label}}{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}) \ WF} \text{ FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp \ WF \text{ CG-wff-label}
\end{array}$$

□

**Lemma 5.4** (CG  $\rightsquigarrow$  FG: Free variable lemma).  $\forall \tau. FV(\llbracket \tau \rrbracket) \subseteq FV(\tau)$

*Proof.* Proof by induction on the CG types,  $\tau$

1.  $\tau = \mathbf{b}$ :

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ = & FV(\mathbf{b}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & \emptyset \\ = & FV(\mathbf{b}) \end{aligned}$$

2.  $\tau = \mathbf{unit}$ :

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ = & FV(\mathbf{unit}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & \emptyset \\ = & FV(\mathbf{unit}) \end{aligned}$$

3.  $\tau = \tau_1 \rightarrow \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 \rightarrow \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \rightarrow \tau_2) \end{aligned}$$

4.  $\tau = \tau_1 \times \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 \times \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 \times \tau_2) \end{aligned}$$

5.  $\tau = \tau_1 + \tau_2$ :

$$\begin{aligned} & FV(\llbracket \tau_1 + \tau_2 \rrbracket) \\ = & FV(\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ \subseteq & FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ = & FV(\tau_1 + \tau_2) \end{aligned}$$

6.  $\tau = \mathbf{ref} \ell_i \tau_i$ :

$$\begin{aligned} & FV(\llbracket \mathbf{ref} \ell_i \tau_i \rrbracket) \\ = & FV(\mathbf{ref} (\llbracket \tau_i \rrbracket + \mathbf{unit})^{\ell_i})^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_i \rrbracket) \cup FV(\ell_i) \\ \subseteq & FV(\tau_i) \cup FV(\ell_i) \quad \text{IH} \\ = & FV(\mathbf{ref} \ell_i \tau_i) \end{aligned}$$

7.  $\tau = \forall \alpha. \tau_i$ :

$$\begin{aligned} & FV(\llbracket \forall \alpha. \tau_i \rrbracket) \\ = & FV(\forall \alpha. (\top, \llbracket \tau_i \rrbracket))^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ = & FV(\llbracket \tau_i \rrbracket) - \{\alpha\} \\ \subseteq & FV(\tau_i) - \{\alpha\} \quad \text{IH} \\ = & FV(\forall \alpha. \tau_i) \end{aligned}$$

8.  $\tau = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket c \Rightarrow \tau_i \rrbracket) \\
&= \text{FV}(c \xrightarrow{\top} \llbracket \tau_i \rrbracket)^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\llbracket \tau_i \rrbracket) \\
&\subseteq \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\tau_i) && \text{IH} \\
&= \text{FV}(c \Rightarrow \tau_i)
\end{aligned}$$

9.  $\tau = \text{Labeled } \ell_i \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \text{Labeled } \ell_i \tau_i \rrbracket) \\
&= \text{FV}(\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i} && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \\
&\subseteq \text{FV}(\tau_i) \cup \text{FV}(\ell_i) && \text{IH} \\
&= \text{FV}(\text{Labeled } \ell_i \tau_i)
\end{aligned}$$

10.  $\tau = \text{SLIO } \ell_1 \ell_2 \tau_i$ :

$$\begin{aligned}
& \text{FV}(\llbracket \text{SLIO } \ell_1 \ell_2 \tau_i \rrbracket) \\
&= \text{FV}(\text{unit} \xrightarrow{\ell_1} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_2})^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
&= \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_1) \cup \text{FV}(\ell_2) \\
&\subseteq \text{FV}(\tau_i) \cup \text{FV}(\ell_1) \cup \text{FV}(\ell_2) && \text{IH} \\
&= \text{FV}(\text{SLIO } \ell_1 \ell_2 \tau_i)
\end{aligned}$$

□

**Lemma 5.5** (CG  $\rightsquigarrow$  FG: Substitution lemma).  $\forall \tau. s.t \vdash \tau \text{ WF}$  the following holds:

$$\llbracket \tau \rrbracket[\ell/\alpha] = \llbracket \tau[\ell/\alpha] \rrbracket$$

*Proof.* Proof by induction on the CG types,  $\tau$

1.  $\tau = \mathbf{b}$ :

$$\begin{aligned}
& (\llbracket \mathbf{b} \rrbracket)[\ell/\alpha] \\
&= (\mathbf{b}^\perp)[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
&= (\mathbf{b}^\perp) \\
&= \llbracket \mathbf{b} \rrbracket \\
&= \llbracket (\mathbf{b}[\ell/\alpha]) \rrbracket
\end{aligned}$$

2.  $\tau = \text{unit}$ :

$$\begin{aligned}
& (\llbracket \text{unit} \rrbracket)[\ell/\alpha] \\
&= (\text{unit}^\perp)[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
&= (\text{unit}^\perp) \\
&= \llbracket \text{unit} \rrbracket \\
&= \llbracket (\text{unit}[\ell/\alpha]) \rrbracket
\end{aligned}$$

3.  $\tau = \tau_1 \rightarrow \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 \rightarrow \tau_2 \rrbracket)[\ell/\alpha] \\
&= (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
&= (\llbracket \tau_1 \rrbracket[\ell/\alpha] \xrightarrow{\top} \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
&= (\llbracket \tau_1[\ell/\alpha] \rrbracket \xrightarrow{\top} \llbracket \tau_2[\ell/\alpha] \rrbracket)^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell/\alpha] \rightarrow \tau_2[\ell/\alpha]) \rrbracket \\
&= \llbracket (\tau_1 \rightarrow \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$



4.  $\tau = \tau_1 \times \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 \times \tau_2 \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] \times \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
= & (\llbracket \tau_1[\ell/\alpha] \rrbracket \times \llbracket \tau_2[\ell/\alpha] \rrbracket)^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
= & \llbracket (\tau_1[\ell/\alpha] \times \tau_2[\ell/\alpha]) \rrbracket \\
= & \llbracket (\tau_1 \times \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$

5.  $\tau = \tau_1 + \tau_2$ :

$$\begin{aligned}
& (\llbracket \tau_1 + \tau_2 \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_1 \rrbracket[\ell/\alpha] + \llbracket \tau_2 \rrbracket[\ell/\alpha])^\perp \\
= & (\llbracket \tau_1[\ell/\alpha] \rrbracket + \llbracket \tau_2[\ell/\alpha] \rrbracket)^\perp && \text{IH on } \tau_1 \text{ and } \tau_2 \\
= & \llbracket (\tau_1[\ell/\alpha] + \tau_2[\ell/\alpha]) \rrbracket \\
= & \llbracket (\tau_1 + \tau_2)[\ell/\alpha] \rrbracket
\end{aligned}$$

6.  $\tau = \text{ref } \ell_i \tau_i$ :

$$\begin{aligned}
& (\llbracket \text{ref } \ell_i \tau_i \rrbracket)[\ell/\alpha] \\
= & (\text{ref } (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i})^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\text{ref } (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i})^\perp && \text{Lemma 5.3} \\
= & \llbracket (\text{ref } \ell_i \tau_i) \rrbracket && \text{since } \vdash \tau \text{ WF} \\
= & \llbracket (\text{ref } \ell_i \tau_i)[\ell/\alpha] \rrbracket
\end{aligned}$$

7.  $\tau = \forall \alpha. \tau_i$ :

$$\begin{aligned}
& (\llbracket \forall \alpha. \tau_i \rrbracket)[\ell/\alpha] \\
= & (\forall \alpha. (\top, \llbracket \tau_i \rrbracket))^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\forall \alpha. (\top, \llbracket \tau_i \rrbracket[\ell/\alpha]))^\perp \\
= & (\forall \alpha. (\top, \llbracket \tau_i[\ell/\alpha] \rrbracket))^\perp && \text{IH} \\
= & (\forall \alpha. \tau_i[\ell/\alpha]) \\
= & (\forall \alpha. \tau_i)[\ell/\alpha]
\end{aligned}$$

8.  $\tau = c \Rightarrow \tau_i$ :

$$\begin{aligned}
& (\llbracket c \Rightarrow \tau_i \rrbracket)[\ell/\alpha] \\
= & (c \stackrel{\top}{\Rightarrow} \llbracket \tau_i \rrbracket)^\perp[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (c[\ell/\alpha] \stackrel{\top}{\Rightarrow} \llbracket \tau_i \rrbracket[\ell/\alpha])^\perp \\
= & (c[\ell/\alpha] \stackrel{\top}{\Rightarrow} \llbracket \tau_i[\ell/\alpha] \rrbracket)^\perp && \text{IH} \\
= & (c[\ell/\alpha] \Rightarrow \tau_i[\ell/\alpha]) \\
= & (c \Rightarrow \tau_i)[\ell/\alpha]
\end{aligned}$$

9.  $\tau = \text{Labeled } \ell_i \tau_i$ :

$$\begin{aligned}
& (\llbracket \text{Labeled } \ell_i \tau_i \rrbracket)[\ell/\alpha] \\
= & (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i}[\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\llbracket \tau_i \rrbracket[\ell/\alpha] + \text{unit})^{\ell_i}[\ell/\alpha] \\
= & (\llbracket \tau_i[\ell/\alpha] \rrbracket + \text{unit})^{\ell_i}[\ell/\alpha] && \text{IH} \\
= & \llbracket (\text{Labeled } \ell_i[\ell/\alpha] \tau_i[\ell/\alpha]) \rrbracket \\
= & \llbracket (\text{Labeled } \ell_i \tau_i)[\ell/\alpha] \rrbracket
\end{aligned}$$

10.  $\tau = \mathbb{C} \ell_1 \ell_2 \tau_i$ :

$$\begin{aligned}
& (\llbracket \mathbb{C} \ell_1 \ell_2 \tau_i \rrbracket) [\ell/\alpha] \\
= & (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_2})^\perp [\ell/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\
= & (\text{unit} \xrightarrow{\ell_1[\ell/\alpha]} (\llbracket \tau_i \rrbracket [\ell/\alpha] + \text{unit})^{\ell_2[\ell/\alpha]})^\perp \\
= & (\text{unit} \xrightarrow{\ell_1[\ell/\alpha]} (\llbracket \tau_i[\ell/\alpha] \rrbracket + \text{unit})^{\ell_2[\ell/\alpha]})^\perp && \text{IH} \\
= & (\mathbb{C} \ell_1[\ell/\alpha] \ell_2[\ell/\alpha] \tau_i[\ell/\alpha]) \\
= & (\mathbb{C} \ell_1 \ell_2 \tau_i) [\ell/\alpha]
\end{aligned}$$

□

### 5.1.3 Model for CG to FG translation

**Definition 5.6** ( ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$   
 $\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

**Definition 5.7** ( $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$   
 $\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

**Definition 5.8** (Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket \wedge {}^sv = {}^tv\} \\
\llbracket \text{unit} \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \text{unit} \rrbracket \wedge {}^tv \in \llbracket \text{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \mid \\
& \quad ({}^s\theta, m, {}^sv_1, {}^tv_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \text{inl } {}^sv, \text{inl } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\
& \quad \{({}^s\theta, m, \text{inr } {}^sv, \text{inr } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \\
& \quad \implies ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \forall \alpha. \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \Lambda e_s, \Lambda e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket c \Rightarrow \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, \nu e_s, \nu e_t) \mid \mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \text{ref } \ell \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^sa, {}^ta) \mid {}^s\theta({}^sa) = \text{Labeled } \ell \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}\} \\
\llbracket \text{Labeled } \ell \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid \\
& \quad \exists {}^sv', {}^tv'. {}^sv = \text{Lb}({}^sv') \wedge {}^tv = \text{inl } {}^tv' \wedge ({}^s\theta, m, {}^sv', {}^tv') \in \llbracket \tau \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V^{\hat{\beta}} & \triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^sv', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
& \quad (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^sv) \Downarrow_i^f (H_s', {}^sv') \wedge i < k \implies \\
& \quad \exists H_t', {}^tv'. (H_t, {}^tv) \Downarrow (H_t', {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s', H_t') \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\
& \quad \exists {}^tv''. {}^tv = \text{inl } {}^tv'' \wedge ({}^s\theta', k - i, {}^sv', {}^tv'') \in \llbracket \tau \rrbracket_V^{\hat{\beta}''}\}
\end{aligned}$$

**Definition 5.9** (Unary expression relation).

$$\begin{aligned}
\llbracket \tau \rrbracket_E^{\hat{\beta}} & \triangleq \{({}^s\theta, n, e_s, e_t) \mid \\
& \quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv.e_s \Downarrow_i {}^sv \implies \\
& \quad \exists H_t', {}^tv. (H_t, e_t) \Downarrow (H_t', {}^tv) \wedge ({}^s\theta, n - i, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H_t') \triangleright^{\hat{\beta}} {}^s\theta\}
\end{aligned}$$

**Definition 5.10** (Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta &\triangleq \text{dom}(s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\hat{\beta} \subseteq (\text{dom}(s\theta) \times \text{dom}(H_t)) \wedge \\ &\forall (a_1, a_2) \in \hat{\beta}. (s\theta, n-1, H_s(a_1), H_t(a_2)) \in [s\theta(a)]_{\hat{\beta}}^V \end{aligned}$$

**Definition 5.11** (Value substitution).  $\delta^s : \text{Var} \mapsto \text{Val}$ ,  $\delta^t : \text{Var} \mapsto \text{Val}$

**Definition 5.12** (Unary interpretation of  $\Gamma$ ).

$$\begin{aligned} [\Gamma]_{\hat{\beta}}^V &\triangleq \{ (s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ &\forall x \in \text{dom}(\Gamma). (s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_{\hat{\beta}}^V \} \end{aligned}$$

#### 5.1.4 Soundness proof for CG to FG translation

**Lemma 5.13** (Monotonicity).  $\forall s\theta, s\theta', n, s_v, t_v, n', \beta, \beta'$ .

$$(s\theta, n, s_v, t_v) \in [\tau]_{\hat{\beta}}^V \wedge s\theta \sqsubseteq s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies (s\theta', n', s_v, t_v) \in [\tau]_{\hat{\beta}'}^V$$

*Proof.* Proof by induction on  $\tau$

1. Case **b**:

Given:

$$(s\theta, n, s_v, t_v) \in [\mathbf{b}]_{\hat{\beta}}^V \wedge s\theta \sqsubseteq s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$(s\theta', n', s_v, t_v) \in [\mathbf{b}]_{\hat{\beta}'}^V$$

Since  $(s\theta, n, s_v, t_v) \in [\mathbf{b}]_{\hat{\beta}}^V$  therefore from Definition 5.8 we know that  $s_v \in \llbracket \mathbf{b} \rrbracket \wedge t_v \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 5.8  $s_v \in \llbracket \mathbf{b} \rrbracket \wedge t_v \in \llbracket \mathbf{b} \rrbracket$  we get the desired

2. Case **unit**:

Given:

$$(s\theta, n, s_v, t_v) \in [\mathbf{unit}]_{\hat{\beta}}^V \wedge s\theta \sqsubseteq s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$(s\theta', n', s_v, t_v) \in [\mathbf{unit}]_{\hat{\beta}'}^V$$

Since  $(s\theta, n, s_v, t_v) \in [\mathbf{unit}]_{\hat{\beta}}^V$  therefore from Definition 5.8 we know that  $s_v \in \llbracket \mathbf{unit} \rrbracket \wedge t_v \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 5.8  $s_v \in \llbracket \mathbf{unit} \rrbracket \wedge t_v \in \llbracket \mathbf{unit} \rrbracket$  we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 5.8 we know that  ${}^s v = ({}^s v_1, {}^s v_2)$  and  ${}^t v = ({}^t v_1, {}^t v_2)$ .

We also know that  $({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\underline{\text{IH1:}} ({}^s\theta', n', {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\underline{\text{IH2:}} ({}^s\theta', n', {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 5.8, IH1 and IH2 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 5.8 two cases arise

(a)  ${}^s v = \text{inl}({}^s v')$  and  ${}^t v = \text{inl}({}^t v')$ :

$$\underline{\text{IH:}} ({}^s\theta', n', {}^s v', {}^t v') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 5.8 and IH we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b)  ${}^s v = \text{inr}({}^s v')$  and  ${}^t v = \text{inr}({}^t v')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \rightarrow \tau_2$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 5.8 we know that

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 5.8 we are required to prove

$$\forall^s \theta'_1 \sqsupseteq {}^s \theta', {}^s v_2, {}^t v_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s \theta'_1, j, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s \theta'_1, j, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some  ${}^s \theta'_1 \sqsupseteq {}^s \theta', {}^s v_2, {}^t v_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$  s.t  $({}^s \theta'_1, j, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}}$  and we are required to prove

$$({}^s \theta'_1, j, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with  ${}^s \theta'_1, {}^s v_2, {}^t v_2, j, \hat{\beta}''$  since  ${}^s \theta'_1 \sqsupseteq {}^s \theta' \sqsupseteq {}^s \theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s \theta'_1, j, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case  $\forall \alpha. \tau$ :

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\forall \alpha. \tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\forall \alpha. \tau]_V^{\hat{\beta}'}$$

From Definition 5.8 we know that  ${}^s v = \Lambda e'_s$  and  ${}^t v = \Lambda e'_t$ . And

$$\forall^s \theta'' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'' . ({}^s \theta'', j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''} \quad (\text{F0})$$

Similarly from Definition 5.8 we are required to prove

$$\forall^s \theta''_1 \sqsupseteq {}^s \theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s \theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

This means we are given some  ${}^s \theta''_1 \sqsupseteq {}^s \theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s \theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

Instantiating (F0) with  ${}^s \theta''_1, j, \hat{\beta}''_1$  since  ${}^s \theta''_1 \sqsupseteq {}^s \theta' \sqsupseteq {}^s \theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$  therefore we get

$$({}^s \theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}''_1}$$

7. Case  $c \Rightarrow \tau$ :

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [c \Rightarrow \tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [c \Rightarrow \tau]_V^{\hat{\beta}'}$$

From Definition 5.8 we know that  ${}^s v = \nu(e'_s)$  and  ${}^t v = \nu(e'_t)$ . And

$$\mathcal{L} \models c \implies \forall^s \theta'' \sqsupseteq {}^s \theta, j < n, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s \theta'', j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 5.8 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}_1''. ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

This means we are given some  $\mathcal{L} \models c, {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}_1''$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

Since  $\mathcal{L} \models c$  and instantiating (C0) with  ${}^s\theta''_1, j, \hat{\beta}_1''$  since  ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}_1''$  therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_E^{\hat{\beta}_1''}$$

8. Case ref  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \ell \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \tau]_V^{\hat{\beta}'}$$

From Definition 5.8 we know that  ${}^s v = {}^s a$  and  ${}^t v = {}^t a$ . We also know that

$${}^s\theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}$$

From Definition 5.8, Definition 5.6 and Definition 5.7 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \tau]_V^{\hat{\beta}'}$$

9. Case Labeled  $\ell \tau$ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}'}$$

From Definition 5.8 it means

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, n, {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}$$

$$\text{IH: } ({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}$$

Similarly from Definition 5.8 we need to prove that

$$\exists {}^s v'', {}^t v''. {}^s v = \text{Lb}_\ell({}^s v'') \wedge {}^t v = \text{inl } {}^t v'' \wedge ({}^s\theta', n', {}^s v'', {}^t v'') \in [\tau]_V^{\hat{\beta}'}$$

We choose  ${}^s v''$  as  ${}^s v'$  and  ${}^t v''$  as  ${}^t v'$  and since from IH we know that  $({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}'}$

Therefore from Definition 5.8 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}'}$$

10. Case  $\mathbb{C} \ell_1 \ell_2 \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}'}$$

This means from Definition 5.8 we know that

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^sv', {}^tv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1. \\ & (k, H_s, H_t) \stackrel{\hat{\beta}_1}{\triangleright} ({}^s\theta_e) \wedge (H_s, {}^sv) \Downarrow_i^f (H_s', {}^sv') \wedge i < k \implies \\ & \exists {}^tv'. (H_t, {}^tv()) \Downarrow (H_t', {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H_s', H_t') \stackrel{\hat{\beta}_2}{\triangleright} {}^s\theta' \wedge \\ & \exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', {}^t\theta', k - i, {}^sv', {}^tv'') \in [\tau]_V^{\hat{\beta}_2} \wedge \\ & (\forall a. H_s(a) \neq H_s'(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \quad (\text{CG0}) \end{aligned}$$

Similarly from Definition 5.8 we need to prove

$$\begin{aligned} & \forall {}^s\theta'_e \sqsupseteq {}^s\theta', H_s', H_t', i', {}^sv'', {}^tv'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1. \\ & (k', H_s', H_t') \stackrel{\hat{\beta}'_1}{\triangleright} ({}^s\theta'_e) \wedge (H_s', {}^sv) \Downarrow_i^f (H_s'', {}^sv'') \wedge (H_t', {}^tv()) \Downarrow (H_t'', {}^tv'') \wedge i' < k' \implies \\ & \exists {}^tv''. (H_t', {}^tv()) \Downarrow (H_t'', {}^tv'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H_s'', H_t'') \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'' \wedge \\ & \exists {}^tv'''. {}^tv' = \text{inl } {}^tv''' \wedge ({}^s\theta', k' - i', {}^sv', {}^tv''') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\ & (\forall a. H_s(a) \neq H_s'(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \end{aligned}$$

This means we are given some  ${}^s\theta'_e \sqsupseteq {}^s\theta', H_s', H_t', i', {}^sv'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  s.t.  $(k', H_s', H_t') \triangleright ({}^s\theta'_e) \wedge (H_s', {}^sv) \Downarrow_i^f (H_s'', {}^sv'') \wedge i' < k'$

And we need to prove

$$\begin{aligned} & \exists {}^tv''. (H_t', {}^tv()) \Downarrow (H_t'', {}^tv'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H_s'', H_t'') \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'' \wedge \\ & \exists {}^tv'''. {}^tv' = \text{inl } {}^tv''' \wedge ({}^s\theta', k' - i', {}^sv', {}^tv''') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\ & (\forall a. H_s(a) \neq H_s'(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with  ${}^s\theta'_e \sqsupseteq {}^s\theta', H_s', H_t', i', {}^sv'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  we get the desired

□

**Lemma 5.14** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove:  $(\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 5.12 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 5.12 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_{V}^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$ :

Given

- $\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_{V}^{\hat{\beta}'}$ :

Since we know that  $\forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_{V}^{\hat{\beta}}$  (given)

Therefore from Lemma 5.13 we get

$$\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_{V}^{\hat{\beta}'}$$

□

**Lemma 5.15** (Unary monotonicity for  $H$ ).  $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$

From Definition 5.10 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$$

And again from Definition 5.10 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$ :

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$  (given)

Therefore from Lemma 5.13 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{V}^{\hat{\beta}}$$

□

**Theorem 5.16** (Fundamental theorem).  $\forall \Gamma, \tau, e, \delta^s, \delta^t, \sigma, {}^s\theta, n$ .

$$\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge$$

$$\mathcal{L} \models \Psi \sigma \wedge$$

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$$

$\implies$

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$$



*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. CF-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\}]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 5.9 it suffices to prove that

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $x \delta^s \Downarrow_i {}^s v$

From cg-val we know that  $i = 0, {}^s v = x \delta^s$ .

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-V0})$$

From fg-val we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$ . So we are left with proving

$$({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we are given  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau \sigma\}]_V^{\hat{\beta}}$ , therefore from Definition 5.12 we get

$$({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_V^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ in the context. So we are done.}$$

2. CF-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{lam}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 5.9 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(\lambda x. e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  ${}^s v = (\lambda x. e_s) \delta^s, {}^t v = (\lambda x. e_t) \delta^t, H'_t = H_t$  and  $i = 0$

It suffices to prove that

$$({}^s\theta, n, (\lambda x.e_s) \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context. So, we are only left to prove

$$({}^s\theta, n, (\lambda x.e_s) \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}}$$

From Definition 5.8 it suffices to prove

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'} \\ \implies ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \end{aligned}$$

This means that we are given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s\theta', j, {}^sv, {}^tv) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

And we need to prove

$$({}^s\theta', j, e_s[{}^sv/x] \delta^s, e_t[{}^tv/x] \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}}$  therefore from Lemma 5.14 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^sv_1\}, e_t \cup \{x \mapsto {}^tv_1\}) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \text{ s.t}$$

$$({}^s\theta', j, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$$

We get (F-L0) directly from IH

### 3. CF-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}}$

This means from Definition 5.9 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^sv \implies \\ \exists H'_t, {}^tv. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^tv) \wedge ({}^s\theta, n - i, {}^sv, {}^tv) \in [\tau_2 \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^sv$  s.t  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^sv$

And we need to prove

$$\exists H'_t, {}^tv. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^tv) \wedge ({}^s\theta, n - i, {}^sv, {}^tv) \in [\tau_2 \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.9 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A1})$$

IH2:

$$({}^s\theta, n-j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.9 it suffices to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n-j, {}^s v_2. e_{s2} \Downarrow_i {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n-j-k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n-j-k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta' \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i-j < n-j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

And we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n-j-k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n-j-k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A2})$$

Since from (F-A1) we know that  $({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}}$  where

$${}^s v_1 = \lambda x. e'_s \text{ and } {}^t v_1 = \lambda x. e'_t$$

From Definition 5.8 we have

$$\begin{aligned} \forall {}^s\theta'_3 \sqsupseteq {}^s\theta, {}^s v, {}^t v, l < n-j, \hat{\beta}_3 \sqsupseteq \hat{\beta}. ({}^s\theta'_3, l, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}_3} \\ \implies ({}^s\theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}_3} \end{aligned}$$

Instantiating with  ${}^s\theta, {}^s v_2, {}^t v_2, n-j-k, \hat{\beta}$  we get

$$({}^s\theta, n-j-k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} \forall H_{s4}, H_{t4}. (n-j-k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k' < n-j-k, {}^s v_4. e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n-j-k-k', {}^s v_4, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge \\ (n-j-k-k', H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t_2}$ , from (F-A2) we know that  $(n-j-k, H_s, H'_{t_2}) \hat{\triangleright}^s \theta$ . Instantiating  ${}^s v_4$  with  ${}^s v$  and since we know that  $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k' < i - j - k < n - j - k$  s.t  $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$ . therefore we have

$$\exists H'_{t_4}, {}^t v_4. (H_{t_4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t_4}, {}^t v_4) \wedge ({}^s \theta, n - j - k - k', {}^s v, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k - k', H_{s_4}, H'_{t_4}) \hat{\triangleright}^s \theta \quad (\text{F-A3})$$

Since from cg-app we know that  $i = j + k + k'$  and  $H'_t = H'_{t_4}$ ,  ${}^t v = {}^t v_4$  therefore we get (F-A0) from (F-A3) and Lemma 5.13 and Lemma 5.15

#### 4. CF-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{prod}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 5.9 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n. {}^s v. (e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n$  s.t  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta' \quad (\text{F-P0})$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists j < i < n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-P1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j. e_{s2} \delta^s \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}, \hat{\beta}'_1$  and since we know that  $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s2} \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-P2})$$

From cg-prod we know that  $i = j + k + 1$ ,  $H'_t = H'_{t2}$  and  ${}^t v = ({}^t v_1, {}^t v_2)$  therefore from Definition 5.8 and Lemma 5.13 we get  $({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$

And since we have  $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 5.15 we also get

$$(n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$$

5. CF-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}} \quad (\text{F-F0})$

This means from Definition 5.9 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n. {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j ({}^s v_1, -) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s \theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and  ${}^s v_1$  with  ${}^s v$  since we know that  $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j ({}^s v, -)$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s \theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-F1}) \end{aligned}$$

From cg-fst we know that  $i = j + 1$ ,  $H'_t = H'_{t1}$  and  ${}^t v = {}^t v_1$ . Since we know  $({}^s \theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}}$  therefore from Definition 5.8 and Lemma 5.13 we get  $({}^s \theta, n - i, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}}$

And since from (F-F1) we have  $(n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta$  therefore from Lemma 5.15 we get

$$(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$$

#### 6. CF-snd:

Symmetric reasoning as in the CF-fst case

#### 7. CF-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{CF-inl}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{inl}(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\tau_1 + \tau_2) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 5.9 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^s \theta \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL1})$$

From cg-inl we know that  $i = j + 1$  and  $H'_t = H'_{t1}, {}^t v = {}^t v_1$ . Since we know  $({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from Definition 5.8 and Lemma 5.13 we get

$$({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}$$

And since from (F-IL1) we have  $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 5.15 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{case}(e_t, x.e_{t1}, y.e_{t2})} \text{CF-case}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 5.9 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-C0})$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-C1})$$

Two cases arise:

- (a)  ${}^s v_1 = \text{inl}({}^s v'_1)$  and  ${}^t v_1 = \text{inl}({}^t v'_1)$ :

IH2:

$$({}^s\theta, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j, {}^s v_2. e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - \\ j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$  and since we know that  $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$  therefore  $\exists k < i - j < n - j$  s.t  $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$$

From cg-case1 we know that  $i = j + k + 1$  and  $H'_t = H'_{t2}$ ,  ${}^t v = {}^t v_2$ . Since we know  $({}^s\theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$  therefore from Definition 5.8 and Lemma 5.13 we get  $({}^s\theta, n - i, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$

And since from (F-C2) we have  $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Lemma 5.15 we get  $(n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$

- (b)  ${}^s v_1 = \text{inr}({}^s v'_1)$  and  ${}^t v_1 = \text{inr}({}^t v'_1)$ :

Symmetric reasoning as in the previous case

## 10. CF-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \Lambda e_s : \forall \alpha. \tau \rightsquigarrow \Lambda e_t} \text{FI}$$



Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \Lambda e_s \delta^s, \Lambda e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_E^{\hat{\beta}}$

This means from Definition 5.9 we know that

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \Lambda e_s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \Lambda e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $(\Lambda e_s) \delta^s \Downarrow_i {}^s v$

From CG-Sem-val and fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s$ ,  ${}^t v = (\Lambda e_t) \delta^t$ ,  $i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s\theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context. So, we are only left to prove

$$({}^s\theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 5.8 it suffices to prove

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'}$$

This means that we are given  ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'} \quad (\text{F-FI0})$$

Since  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$  therefore from Lemma 5.14 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E^{\hat{\beta}'}$$

We get (F-FI0) directly from IH

#### 11. CF-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \forall \alpha. \tau \rightsquigarrow e_t \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e_s [] : \tau[\ell/\alpha] \rightsquigarrow e_t []} \text{FE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s [] \delta^s, e_t [] \delta^t) \in [\tau[\ell/\alpha] \sigma]_E^{\hat{\beta}}$

From Definition 5.9 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v.e_s \sqcap \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t \sqcap) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n, {}^s v$  s.t  $e_s \sqcap \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t \sqcap) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall\alpha.\tau) \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.9 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\forall\alpha.\tau) \sigma]_{V}^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s \sqcap) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n, {}^s v_1$  s.t  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\forall\alpha.\tau) \sigma]_{V}^{\hat{\beta}} \wedge (n-j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-FE1})$$

From CG-Sem-FE we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_1 = \Lambda e'_t$

Therefore we have

$$({}^s\theta, n-j, \Lambda e'_s, \Lambda e'_t) \in [(\forall\alpha.\tau) \sigma]_{V}^{\hat{\beta}}$$

This means from Definition 5.8 we have

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, k < n-j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s\theta', k, e'_s, e'_t) \in [\tau[\ell'/\alpha] \sigma]_E^{\hat{\beta}_2}$$

Instantiating  ${}^s\theta'$  with  ${}^s\theta$ ,  $k$  with  $n-j-1$ ,  $\ell'$  with  $\ell$   $\sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$$({}^s\theta, n-j-1, e'_s, e'_t) \in [\tau[\ell/\alpha] \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we get

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n-j-1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s\theta'_1 \wedge \forall k < n-j-1, {}^s v_2.e'_s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n-j-1-k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n-j-1-k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-FE1) we know that  $(n-j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} s\theta$  therefore from Lemma 5.15 we get  $(n-j-1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} s\theta$

Since we know that  $e_s \sqcap \delta^s \Downarrow_i {}^s v$  and from CG-Sem-FE we know that  $i = j+k+1$  (for some k) and  $i < n$  therefore we have  $k < n-j-1$  s.t  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j - 1 - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-FE2})$$

Since  $H'_t = H_{t2}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-FE0) directly from (F-FE2)

12. CF-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \nu e_s : c \Rightarrow \tau \rightsquigarrow \nu e_t} \text{ CI}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \nu e_s \delta^s, \nu e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 5.9 we know that

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n. \nu e_s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \nu e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(c \Rightarrow \tau) \hat{\beta} \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  and given some  $i < n$  s.t  $(\nu e_s) \delta^s \Downarrow_i {}^s v$

From CG-Sem-val and fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  ${}^t v = (\nu e_t) \delta^t$ ,  $i = 0$  and  $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$  from the context. So, we are only left to prove

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$$

From Definition 5.8 it suffices to prove

$$\mathcal{L} \models c \sigma \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}'}$$

This means that we are given  $\mathcal{L} \models c \sigma$  and  ${}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}'} \quad (\text{F-CI0})$$

Since  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$  therefore from Lemma 5.14 we also have

$$({}^s \theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}'}$$

And since we know that  $\mathcal{L} \models c \sigma$  therefore

$$\underline{\text{IH:}} ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

13. CF-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : c \Rightarrow \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e_s \bullet : \tau \rightsquigarrow e_t \bullet} \text{CE}$$

Also given is:  $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, e_s \bullet \delta^s, e_t \bullet \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

From Definition 5.9 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s \bullet \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $e_s \bullet \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.9 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(e_s \bullet) \delta^s \Downarrow_i {}^s v$  therefore  $\exists j < i < n$  s.t.  $e_s \delta^s \Downarrow_j {}^s v_1$ .

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE1})$$

From CG-Sem-CE we know that  ${}^s v_1 = \nu e'_s$  and  ${}^t v_1 = \nu e'_t$

Therefore we have

$$({}^s\theta, n - j, \nu e'_s, \nu e'_t) \in [(c \Rightarrow \tau) \sigma]_V^{\hat{\beta}}$$

This means from Definition 5.8 we have

$$\forall {}^s\theta' \sqsupseteq {}^s\theta'_1, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s\theta', k, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}_2}$$

Instantiating  ${}^s\theta'$  with  ${}^s\theta$ ,  $k$  with  $n - j - 1$ ,  $\ell'$  with  $\ell$   $\sigma$  and  $\hat{\beta}_2$  with  $\hat{\beta}$  and we get

$$({}^s\theta, n - j - 1, e'_s, e'_t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we get

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \hat{\triangleright}^{\beta_2} s\theta'_1 \wedge \forall k < n - j - 1. e'_s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright}^{\beta} s\theta \end{aligned}$$

Instantiating with  $H_s, H'_{t1}$ . Since from (F-CE1) we know that  $(n - j, H_s, H'_{t1}) \hat{\triangleright}^{\beta} s\theta$  therefore from Lemma 5.15 we get  $(n - j - 1, H_s, H'_{t1}) \hat{\triangleright}^{\beta} s\theta$

Since we know that  $e_s \bullet \delta^s \Downarrow_i {}^s v$  and from CG-Sem-CE we know that  $i = j + k + 1$  (for some  $k$ ) and  $i < n$  therefore we have  $k < n - j - 1$  s.t.  $e'_s \delta^s \Downarrow_k {}^s v_2$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_{t2}) \hat{\triangleright}^{\beta} s\theta \quad (\text{F-CE2})$$

Since  $H'_t = H_{t2'}$ ,  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  therefore we get (F-CE0) directly from (F-CE2)

14. CF-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e_s) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{-}.\text{inl}(e_t)} \text{ret}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\beta} s\theta \wedge \forall i < n, {}^s v.\text{ret}(e_s) \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\beta} s\theta \end{aligned}$$

This means that given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \hat{\triangleright}^{\beta} s\theta$  and given some  $i < n$  s.t.  $\text{ret}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\beta} s\theta$$

From CG-ret and FG-lam we know that  $i = 0$ ,  ${}^s v = \text{ret}(e_s) \delta^s$ ,  ${}^t v = \lambda_{-}.\text{inl}(e_t) \delta^t$  and  $H'_t = H_t$ .

So we need to prove

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{\beta} s\theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^{\beta} s\theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ .

$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \exists H'_t, {}^t v'. (H_t, (\lambda_{-}.\text{inl}(e_t) ()) \delta^t) \Downarrow$   
 $(H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge$   
 $\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''}$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k$ . Also from cg-ret we know that  $H'_s = H_s$

And we need to prove

$\exists H'_t, {}^t v'. (H_t, (\lambda_{-}.\text{inl}(e_t) ()) \delta^t) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge$   
 $\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-R0})$

**IH:**

$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'}$

It means from Definition 5.9 that we need to prove

$\forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies$   
 $\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

Instantiating  $H_{s1}$  with  $H_s$  and  $H_{t1}$  with  $H_t$ . And since we know that  $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ . Therefore we have

$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-R1})$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_{t1}$ ,  ${}^t v'$  as  $\text{inl}({}^t v)$ ,  ${}^s\theta'$  as  ${}^s\theta_e$ ,  $\hat{\beta}''$  as  $\hat{\beta}'$ . Since from cg-ret we know that  $i = f + 1$  therefore from (F-R1) and Lemma 5.15 we know that  $(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

Next we choose  ${}^t v''$  as  ${}^t v$  (from F-R1) and from Lemma 5.13 we get  $({}^s\theta_e, k - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$  (we know from cg-ret that  ${}^s v' = {}^s v$ )

15. CF-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_{s2} : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_3 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell_4 \quad \Sigma; \Psi \vdash \ell_4 \sqsubseteq \ell_o}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_{-}.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{bind}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \rrbracket_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_{-}.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in \llbracket (\mathbb{C} \ell_i \ell_o \tau') \sigma \rrbracket_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned}
& \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\
& \exists H'_t, {}^t v. (H_t, \lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())) \delta^t \Downarrow (H'_t, {}^t v) \wedge \\
& ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{s\theta}
\end{aligned}$$

This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^s v$  s.t.  $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned}
& \exists H'_t, {}^t v. (H_t, \lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())) \delta^t \Downarrow (H'_t, {}^t v) \wedge \\
& ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{s\theta}
\end{aligned}$$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s,$   
 ${}^t v = \lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{s\theta}$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  from the context so we are left with proving

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_{V}^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\begin{aligned}
& \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
& (k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\
& \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))()) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \\
& \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_{V}^{\hat{\beta}''}
\end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned}
& \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_. \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))()) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - \\
& i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-B0})
\end{aligned}$$

IH1:

$$({}^s \theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_{E}^{\hat{\beta}}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned}
& \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^{s\theta} \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\
& \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k - \\
& j, H_{s2}, H'_{t2}) \hat{\triangleright}^{s\theta}
\end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_V^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-B1.1})$$

From Definition 5.8 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s \theta_e$  with  ${}^s \theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s \theta'', k - j - b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_E^{\hat{\beta}''}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta \wedge \forall c < (k - j - b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_V^{\hat{\beta}''} \wedge \\ (k - j - b - c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \end{aligned}$$

Instantiating  $H_{s4}$  with  $H'_{s3}$  and  $H_{t4}$  with  $H'_{t3}$ . And since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists c < i - j - b < k - j - b$  s.t  $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau') \sigma]_V^{\hat{\beta}''} \wedge \\ (k - j - b - c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \quad (\text{F-B2.1}) \end{aligned}$$

From Definition 5.8 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k - j - b - c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1. \\ (m, H_{s5}, H_{t5}) \triangleright^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsupseteq {}^s \theta_e, \hat{\beta}''_1 \sqsubseteq \hat{\beta}''_2 . (m - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s \theta''' \wedge \\ \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_V^{\hat{\beta}''_2} \end{aligned}$$



Instantiating  ${}^s\theta_e$  with  ${}^s\theta''$ ,  $H_{s5}$  with  $H'_{s3}$ ,  $H_{t5}$  with  $H'_{t3}$ ,  $m$  with  $k - j - b - c$  and  $\hat{\beta}'_1$  with  $\hat{\beta}''$ . Since we know that  $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists d < i - j - b - c < k - j - b - c$  s.t.  $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s\theta''' \sqsupseteq {}^s\theta_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_2. (k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta''' \wedge \\ & \exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{V'}^{\hat{\beta}''_2} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose  $H'_{t1}$  as  $H'_{t5}$  and  ${}^t v'$  as  ${}^t v'_{h2}$ . Next we choose  ${}^s\theta'$  as  ${}^s\theta'''$  and  $\hat{\beta}''$  as  $\hat{\beta}''_2$  (both chosen from (F-B2)). Also from cg-bind we know that in (F-B0)  $H'_{s1}$  will be  $H'_{s5}$ .

Since  $(k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta'''$  therefore Lemma 5.13 we get  $(k - i, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta'''$ . Also since from (F-B2) we have  $\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{V'}^{\hat{\beta}''_2}$

Since  $i = j + b + c + d + 1$  therefore from Lemma 5.13 we get

$$\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{V'}^{\hat{\beta}''_2}$$

16. CF-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{ label}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_E^{\hat{\beta}}$

From Definition 5.9 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some  $H_s, H_t, \hat{\beta}$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t.  $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ .

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 5.9 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1.e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $\mathbf{Lb}_\ell(e_s) \delta^s \Downarrow_i \mathbf{Lb}_\ell({}^s v)$  therefore  $\exists j < i < n$  s.t  $e_s \delta^s \Downarrow_j {}^s v$ .

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove  $({}^s \theta, n - i, \mathbf{Lb}_\ell({}^s v), \mathbf{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}}$ . Since from cg-label we know that  $i = j + 1$ ,  ${}^s v = {}^s v_1$  and  ${}^t v = {}^t v_1$ . Therefore we get this from Definition 5.8, (F-LB1) and Lemma 5.13.

From Lemma 5.13 we get  $(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$

17. CF-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e_s) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_. \mathbf{inl}(e_t ())} \text{toLabeled}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \mathbf{inl} e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda_. \mathbf{inl} e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  and given some  $i < n$  s.t  $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda_. \mathbf{inl} e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = \text{toLabeled}(e_s) \delta^s$ ,

$${}^t v = (\lambda_. \mathbf{inl} e_t()) \delta^t, H'_t = H_t$$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \mathbf{inl} e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright}^s \theta$  from the context so we are left with proving

$$({}^s\theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_2 \tau) \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_2 \tau) \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \quad (\text{F-TL0}) \end{aligned}$$

IH:

$$({}^s\theta, k, e_s \delta^s, e_t \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_E^{\hat{\beta}}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies \\ \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} \\ {}^s\theta \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists j < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_j {}^s v_{h1}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} \\ {}^s\theta \quad (\text{F-TL1.1}) \end{aligned}$$

From Definition 5.8 we know have

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (m, H_{s3}, H_{t3}) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s\theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \end{aligned}$$

Instantiating  ${}^s\theta_e$  with  ${}^s\theta$ ,  $H_{s3}$  with  $H_{s1}$ ,  $H_{t3}$  with  $H'_{t2}$ ,  $m$  with  $k - j$  and  $\hat{\beta}'$  with  $\hat{\beta}$ . Since we know that  $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists b < i - j < k - j$  s.t  $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \hat{\triangleright} {}^s \theta'' \wedge \\ & \exists {}^t v'' . {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-TL1}) \end{aligned}$$

In order to prove (F-TL0) we choose  ${}^s \theta'$  as  ${}^s \theta''$  and  $\hat{\beta}'$  as  $\hat{\beta}''$  (both chosen from (F-TL2))

Also from cg-toLabeled and fg-inl, fg-app we know that  $H'_s = H'_{s3}$  and  $H'_t = H'_{t3}$ , and  ${}^s v' = {}^s v'_{h1}$ ,  ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 5.13

18. CF-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e_s) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{-}.e_t} \text{unlabel}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-}.e_t \delta^t) \in [(\mathbb{C} \top (\ell) \tau) \sigma]_{E}^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s \theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \top (\ell) \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright} {}^s \theta$  and given some  $i < n, {}^s v$  s.t  $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \top (\ell) \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta$$

From cg-val and fg-val we know that  $i = 0$ ,  ${}^s v = \text{unlabel}(e_s) \delta^s$ ,  ${}^t v = \lambda_{-}.e_t \delta^t$ ,  $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in [(\mathbb{C} \top (\ell) \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s \theta$$

Since we already know  $(n, H_s, H_t) \hat{\triangleright} {}^s \theta$  from the context so we are left with proving

$$({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-}.e_t \delta^t) \in [(\mathbb{C} \top (\ell) \tau) \sigma]_{V}^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright} ({}^s \theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.e_t)()) \delta^t \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright} {}^s \theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-U0}) \end{aligned}$$

IH:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h.e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - \\ & f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-U1})$$

In order to prove (F-U0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'$  as  ${}^t v_h$ ,  ${}^s \theta'$  as  ${}^s \theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From cg-unlabel and fg-app we also know that  $H'_{s1} = H_{s1}$  and  $H'_{t1} = H'_{t2}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e:$$

Since from (F-U1) we know that  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

Therefore from Lemma 5.15 we also get  $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

$$(b) \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}:$$

Since from (F-U1) we have

$$({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

This means from Definition 5.8 we know that

$$\exists {}^s v_i, {}^t v_i. {}^s v_h = \text{Lb}_\ell({}^s v_i) \wedge {}^t v_h = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-U2})$$

Since we know that  ${}^t v' = {}^t v_h$  and since from (F-U2) we have  ${}^t v_h = \text{inl } {}^t v_i$ . Therefore from we choose  ${}^t v''$  as  ${}^t v_i$  to get the first conjunct

From cg-unlabel we know that  ${}^s v = {}^s v_i$  and since we know that  $({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$

Therefore from Lemma 5.13 we also get  $({}^s \theta_e, k - i, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$

19. CF-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e_s : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{-}.\text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma ]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{new } e_s \delta^s, \lambda_{-}\text{inl}(\text{new } (e_t)) \delta^t) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}\text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that  $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda_{-}\text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, \text{new } e_s \delta^s, \lambda_{-}\text{inl}(\text{new } (e_t)) \delta^t) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, \text{new } e_s \delta^s, \lambda_{-}\text{inl}(\text{new } (e_t)) \delta^t) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \sigma]_V^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}''} \quad (\text{F-N0}) \end{aligned}$$

From cg-ref we know that  ${}^s v' = a_s$  and from fg-ref, fg-inl we know that  ${}^t v' = \text{inl } a_t$ .

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h. e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'} \wedge (k - \\ & f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, \text{new}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_{V'}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright {}^s \theta_e \quad (\text{F-N1})$$

In order to prove (F-N0) we choose  $H'_{t1}$  as  $H'_{t2} \cup \{a_t \mapsto {}^t v_h\}$ ,  ${}^t v$  as  $a_t$ ,  ${}^s \theta'$  as  ${}^s \theta_n$  where  ${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}$

And we choose  $\hat{\beta}''$  as  $\hat{\beta}_n$  where  $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From cg-ref and fg-ref we also know that  $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

We need to prove

(a)  $(k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}_n} {}^s \theta_n$ :

From Definition 5.10 it suffices to prove that

- $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$ :

Since  $\text{dom}({}^s \theta_e) \subseteq \text{dom}(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\} \text{ and } H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$$

Therefore we get  $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (\text{dom}({}^s \theta_e) \times \text{dom}(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ )

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}, H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\} \text{ and } \hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$$

Therefore we get  $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s \theta_n, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s \theta_n(a)]_{V'}^{\hat{\beta}_n}$ :

$$\forall (a_1, a_2) \in \hat{\beta}_n$$

- $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-N1) we know that  $({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V'}^{\hat{\beta}'}$

From Lemma 5.13 we get  $({}^s \theta_n, k - i - 1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V'}^{\hat{\beta}_n}$

- $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$  therefore

from Definition 5.10 we get

$$({}^s \theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_{V'}^{\hat{\beta}'}$$

From Lemma 5.13 we get

$$({}^s \theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_n(a_1)]_{V'}^{\hat{\beta}'}$$

(b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V'}^{\hat{\beta}_n}$ :

We choose  ${}^t v''$  as  ${}^t v_h$  from (F-N1), fg-inl and fg-ref we know that  ${}^t v' = \text{inl } {}^t v_h$

In order to prove  $({}^s \theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_{V'}^{\hat{\beta}_n}$ , from Definition 5.8 it suffices to prove that

$${}^s\theta_n(a_s) = (\text{Labeled } \ell' \tau) \wedge (a_s, a_t) \in \hat{\beta}_n$$

We get this by construction of  ${}^s\theta_n$  and  $\hat{\beta}_n$

20. CF-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{ref } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash !e_s : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{-}.\text{inl}(e_t)} \text{deref}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, !e_s \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n$  s.t  $!e_s \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

From cg-val and fg-val we know that  $i = 0, {}^s v = !e_s \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_t) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, !e_s \delta^s, \lambda_{-}.\text{inl}(e_t) \delta^t) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau)) \sigma]_V^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\begin{aligned} \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove



$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\dots} \text{inl}(e_t))(\delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-D0}) \end{aligned}$$

III:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\text{ref } \ell \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket (\text{ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_h$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket (\text{ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose  $H'_{t1}$  as  $H'_{t2}$ ,  ${}^t v'_1$  as  $H'_{t2}(a)$  (where  ${}^t v_h = a_t$  from fg-deref),  ${}^s \theta'$  as  ${}^s \theta_e$  and we choose  $\hat{\beta}''$  as  $\hat{\beta}'$ .

From cg-deref we also know that  $H'_{s1} = H_{s1}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e:$$

Since from (F-D1) we have  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$  and since  $f < i$  threfore from Lemma 5.15 we get  $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

Since from cg-deref and fg-deref we know that  ${}^s v_h = a_s$  and  ${}^t v_h = a_t$ .

Therefore from (F-D1) and from Definition 5.8 we know that

$${}^s \theta_e(a_s) = (\text{Labeled } \ell \tau) \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that  $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$  which means from Definition 5.10 we know that

$$({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-D2})$$

This means from Definition 5.8 we know that

$$\exists {}^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_{\ell}({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$$

We choose  ${}^t v''$  as  ${}^t v_i$  and we know that  ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$ . This proves the first conjunct.

Since from (F-D2) we have  $({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$  therefore from Lemma 5.13 we get

$$({}^s \theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in \llbracket (\text{Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

This proves the second conjunct.

21. CF-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_{s1} := e_{s2} : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_E^{\hat{\beta}}$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

From cg-val and fg-val we know that  $i = 0, {}^s v = (e_{s1} := e_{s2}) \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  from the context so we are left with proving

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}}$$

From Definition 5.8 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} \\ & {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some  ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ & (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s\theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} &\implies \\ \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} &{}^s\theta_e \end{aligned}$$

Instantiating  $H_{s2}$  with  $H_{s1}$  and  $H_{t2}$  with  $H_{t1}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists f < i < k \leq n$  s.t  $e_s \delta^s \Downarrow_f {}^s v_{h1}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} &{}^s\theta_e \quad (\text{F-S1}) \end{aligned}$$

IH2:

$$({}^s\theta_e, k-f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 5.9 that we need to prove

$$\begin{aligned} \forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k-f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} &\implies \\ \exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'} &{}^s\theta_e \end{aligned}$$

Instantiating  $H_{s3}$  with  $H_{s1}$  and  $H_{t3}$  with  $H'_{t2}$ . And since we know that  $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$  therefore  $\exists l < i-f < k-f$  s.t  $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$ .

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s1}, H'_{t3}) \triangleright^{\hat{\beta}'} &{}^s\theta_e \quad (\text{F-S2}) \end{aligned}$$

In order to prove (F-S0) we choose  $H'_{t1}$  as  $H'_{t3}[a_t \mapsto {}^t v_{h3}]$ ,  ${}^t v'$  as  $(\ )$ ,  ${}^s\theta'$  as  ${}^s\theta_e$  and  $\hat{\beta}''$  as  $\hat{\beta}'$

From cg-assign and fg-assign we also know that  ${}^s v_{h2} = a_s$ ,  ${}^t v_{h2} = a_t$ ,  $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$  and  $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

$$(a) \ (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

From Definition 5.10 it suffices to prove that

- $dom({}^s\theta_e) \subseteq dom(H'_{s1})$ :

Since  $dom({}^s\theta_e) \subseteq dom(H_{s1})$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ )

And since  $dom(H_{s1}) = dom(H'_{s1})$  therefore we also get  $dom({}^s\theta_e) \subseteq dom(H'_{s1})$

- $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H'_{t1}))$ :

Since  $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H_{t1}))$  (given that we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ )

And since  $\text{dom}(H_{t1}) \subseteq \text{dom}(H'_{t1})$  therefore we also have  $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$ :  
 $\forall (a_1, a_2) \in \hat{\beta}'_n$

–  $(a_1, a_2) = (a_s, a_t)$ :

Since from (F-S2) we know that  $({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

From Lemma 5.13 we get  $({}^s\theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

–  $(a_1, a_2) \neq (a_s, a_t)$ :

Since we have  $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$  therefore from Definition 5.10 we get

$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$

From Lemma 5.13 we get

$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$

- (b)  $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}'_n}$ :

We choose  ${}^t v''$  as () from (F-S1), fg-inl and fg-assgn we know that  ${}^t v' = \text{inl } ()$

To prove:  $({}^s\theta_n, k - i, (), ()) \in [\text{unit}]_V^{\hat{\beta}'_n}$ ,

We get this directly from Definition 5.8

□

**Lemma 5.17** (Subtyping). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$ .

$$1. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}'}$$

$$2. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}'}$$

*Proof.* Proof of Statement (1)

Proof by induction on  $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove:  $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}'}$

It suffices to prove:  $\forall ({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}'}$

This means that given some  ${}^s\theta, n$  and  $\lambda x.e_i$  s.t  $({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:

$$\begin{aligned} \forall^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}'} \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \end{aligned} \quad (\text{S-A0})$$

And it suffices to prove:  $({}^s \theta, n, \lambda x. e_i) \in [((\tau_1' \rightarrow \tau_2') \sigma)]_V^{\hat{\beta}}$

Again from Definition 5.8 it suffices to prove:

$$\begin{aligned} \forall^s \theta'_1 \sqsupseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta}' \sqsubseteq \hat{\beta}'_1 \\ ({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1' \sigma]_V^{\hat{\beta}'_1} \implies ({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2' \sigma]_E^{\hat{\beta}'_1} \end{aligned}$$

This means that given some  ${}^s \theta'_1 \sqsubseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta}' \sqsubseteq \hat{\beta}'_1$  s.t  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1' \sigma]_V^{\hat{\beta}'}$

And we are required to prove:  $({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2' \sigma]_E^{\hat{\beta}'_1}$

$$\text{IH: } [(\tau_1' \sigma)]_V^{\hat{\beta}'_1} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}'_1} \text{ (Statement (1))}$$

$$[(\tau_2 \sigma)]_E^{\hat{\beta}'_1} \subseteq [(\tau_2' \sigma)]_E^{\hat{\beta}'_1} \text{ (Sub-A0, From Statement (2))}$$

Instantiating (S-A0) with  ${}^s \theta'_1, {}^s v_1, {}^t v_1, k, \hat{\beta}'_1$

Since  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1' \sigma]_V^{\hat{\beta}'}$  therefore from IH1 we know that  $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$

As a result we get

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

From (Sub-A0), we know that

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2' \sigma]_E^{\hat{\beta}'_1}$$

## 2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau_1' \quad \Sigma; \Psi \vdash \tau_2 <: \tau_2'}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1' \times \tau_2'}$$

To prove:  $[((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau_1' \times \tau_2') \sigma)]_V^{\hat{\beta}}$

$$\text{IH1: } [(\tau_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau_1' \sigma)]_V^{\hat{\beta}} \text{ (Statement (1))}$$

$$\text{IH2: } [(\tau_2 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau_2' \sigma)]_V^{\hat{\beta}} \text{ (Statement (1))}$$

It suffices to prove:

$$\forall ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1' \times \tau_2') \sigma)]_V^{\hat{\beta}}$$

This means that given  $({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s \theta, n, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s \theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1' \times \tau_2') \sigma)]_V^{\hat{\beta}}$

Again from Definition 5.8, it suffices to prove:

$$({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$$

Since from (S-P0) we know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from IH1 we have  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$

Similarly since from (S-P0) we have  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$  therefore from IH2 we get  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}}$

### 3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove:  $[\tau_1 + \tau_2 \sigma]_V^{\hat{\beta}} \subseteq [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

IH1:  $[\tau_1 \sigma]_V^{\hat{\beta}} \subseteq [\tau'_1 \sigma]_V^{\hat{\beta}}$  (Statement (1))

IH2:  $[\tau_2 \sigma]_V^{\hat{\beta}} \subseteq [\tau'_2 \sigma]_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:  $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2) \sigma]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^sv = \text{inl } {}^sv_i$  and  ${}^tv = \text{inl } {}^tv_i$ :

From Definition 5.8 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

(b)  ${}^sv = \text{inr } {}^sv_i$  and  ${}^tv = \text{inr } {}^tv_i$ :

Symmetric reasoning

### 4. SLIO\*sub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove:  $[(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}. ({}^s\theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, \Lambda e_s, \Lambda e_t) \in [((\forall\alpha.\tau_1) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'} \quad (\text{S-F0})$$

And it suffices to prove:  $({}^s\theta, n, \Lambda e_s, \Lambda e_t) \in [((\forall\alpha.\tau_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 5.8, it suffices to prove:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

This means that given  ${}^s\theta_1 \sqsupseteq {}^s\theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:  $({}^s\theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$

Instantiating (S-F0) with  ${}^s\theta_1, k, \ell'_1, \hat{\beta}'_1$  we get

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_1[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \quad (\text{Sub-F0, Statement (2)})$$

From (Sub-F0), we know that

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

#### 5. SLIO\*sub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2}$$

To prove:  $[((c_1 \implies \tau_1) \sigma)]_V^{\hat{\beta}} \subseteq [((c_2 \implies \tau_2) \sigma)]_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, n, \nu e_s, \nu e_t) \in [((c_1 \implies \tau_1) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \nu e_t) \in [((c_2 \implies \tau_2) \sigma)]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, \nu e_s, \nu e_t) \in [((c_1 \implies \tau_1) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:

$$\mathcal{L} \models c_1 \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in [\tau_1 \sigma]_E^{\hat{\beta}'} \quad (\text{S-C0})$$

And it suffices to prove:  $({}^s\theta, n, \nu e_s, \nu e_t) \in [((c_2 \implies \tau_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 5.8, it suffices to prove:

$$\mathcal{L} \models c_2 \sigma \implies \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, k, e_s, e_t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

This means that given  $\mathcal{L} \models c_2, {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:

$$({}^s\theta'_1, k, e_s, e_t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

since we know that  $c_2 \implies c_1$  and since  $\mathcal{L} \models c_2 \sigma$  therefore  $\mathcal{L} \models c_1 \sigma$ . Next we instantiate (S-C0) with  ${}^s\theta'_1, k, \hat{\beta}'_1$  to get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

$$\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_1} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_1} \text{ (Sub-C0, Statement (2))}$$

Therefore from (Sub-C0), we get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove:  $\llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$

IH:  $\llbracket (\tau \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some  $({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:

$$\exists {}^sv', {}^tv'. {}^sv = \text{Lb}_\ell({}^sv') \wedge {}^tv = \text{inl } {}^tv' \wedge ({}^s\theta, m, {}^sv', {}^tv') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$$

From Definition 5.8 it suffices to prove

$$\exists {}^sv', {}^tv'. {}^sv = \text{Lb}_\ell({}^sv') \wedge {}^tv = \text{inl } {}^tv' \wedge ({}^s\theta, m, {}^sv', {}^tv') \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_1 \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell'_2}{\Sigma; \Psi \vdash \mathbb{C} \ell_1 \ell_2 \tau <: \mathbb{C} \ell'_1 \ell'_2 \tau'}$$

To prove:  $\llbracket ((\mathbb{C} \ell_1 \ell_2 \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\mathbb{C} \ell_1 \ell_2 \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given  $({}^s\theta, n, {}^sv, {}^tv) \in \llbracket ((\mathbb{C} \ell_1 \ell_2 \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 5.8 we are given:



$$\begin{aligned}
& \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}' \\
& (k, H_s, H_t) \stackrel{\hat{\beta}'}{\triangleright} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \\
& \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \stackrel{\hat{\beta}''}{\triangleright} {}^s\theta' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{S-M0})
\end{aligned}$$

And we are required to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

So again from Definition 5.8 we need to prove

$$\begin{aligned}
& \forall {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1 \\
& (k_1, H_{s1}, H_{t1}) \stackrel{\hat{\beta}'_1}{\triangleright} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1 \implies \\
& \exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}''_1}{\triangleright} {}^s\theta' \wedge \\
& \exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}
\end{aligned}$$

This means we are given some  ${}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$  s.t.  $(k_1, H_{s1}, H_{t1}) \stackrel{\hat{\beta}'_1}{\triangleright} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1$

And we need to prove

$$\begin{aligned}
& \exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}''_1}{\triangleright} {}^s\theta' \wedge \\
& \exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}
\end{aligned}$$

We instantiate (S-M0) with  ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$  we get

$$\begin{aligned}
& \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \stackrel{\hat{\beta}''}{\triangleright} {}^s\theta' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''}
\end{aligned}$$

$$\text{IH: } [(\tau \sigma)]_V^{\hat{\beta}''} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}' \hat{\beta}''} \text{ (Statement (1))}$$

Since we have  $({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''}$  therefore from IH we get  $({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\tau' \sigma)]_V^{\hat{\beta}''}$

## 8. CGsub-base:

Trivial

### Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}. ({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means that we are given  $({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

From Definition 5.9 it means we have

$$\forall H_s, H_t. (n, H_s, H_t) \stackrel{\hat{\beta}}{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \stackrel{\hat{\beta}}{\triangleright} {}^s\theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

From Definition 5.9 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $j < n, {}^s v_1$  s.t  $e_s \Downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

Instantiating (Sub-E0) with the given  $H_{s1}, H_{t1}$  and  $j < n, {}^s v_1$ . We get

$$\exists H'_t, {}^t v. (H_{t1}, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we have  $({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}}$  therefore from Statement(1) we get  $({}^s\theta, n - j, {}^s v_1, {}^t v) \in [\tau' \sigma]_V^{\hat{\beta}}$

□

**Theorem 5.18** (Deriving CG NI via compilation).  $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s1}, H'_{s2}$ .

let  $\mathbf{bool} = (\mathbf{unit} + \mathbf{unit})$ .

$x : \mathbf{Labeled} \top \mathbf{bool} \vdash e_s : \mathbb{C} \perp \perp \mathbf{bool} \wedge$

$\emptyset \vdash {}^s v_1 : \mathbf{Labeled} \top \mathbf{bool} \wedge \emptyset \vdash {}^s v_2 : \mathbf{Labeled} \top \mathbf{bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s1}, {}^s v'_1) \wedge$

$(\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2}^f (H'_{s2}, {}^s v'_2)$

$\implies$

${}^s v'_1 = {}^s v'_2$

*Proof.* From the CG to FG translation we know that  $\exists e_t$  s.t

$x : \mathbf{Labeled} \top \mathbf{bool} \vdash e_s : \mathbb{C} \perp \perp \mathbf{bool} \rightsquigarrow e_t$

Similarly we also know that  $\exists {}^t v_1, {}^t v_2$  s.t

$\emptyset \vdash {}^s v_1 : \mathbf{Labeled} \top \mathbf{bool} \rightsquigarrow {}^t v_1$  and  $\emptyset \vdash {}^s v_2 : \mathbf{Labeled} \top \mathbf{bool} \rightsquigarrow {}^t v_2$  (NI-0)

From type preservation theorem we know that

$x : ((\mathbf{unit} + \mathbf{unit})^\perp + \mathbf{unit})^\top \vdash_\top e_t : (\mathbf{unit} \xrightarrow{\perp} ((\mathbf{unit} + \mathbf{unit})^\perp + \mathbf{unit})^\perp)^\perp$

$\emptyset \vdash_\top {}^t v_1 : ((\mathbf{unit} + \mathbf{unit})^\perp + \mathbf{unit})^\top$

$\emptyset \vdash_\top {}^t v_2 : ((\mathbf{unit} + \mathbf{unit})^\perp + \mathbf{unit})^\top$  (NI-1)

Since we have  $\emptyset \vdash {}^s v_1 : \mathbf{Labeled} \top \mathbf{bool} \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 5.16 we have (we choose  $n$  s.t  $n > n_1$  and  $n > n_2$ )

$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\mathbf{Labeled} \top \mathbf{bool}]_E^\emptyset$  (NI-2)

And therefore from Definition 5.12 and (NI-2) we have

$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_1)) \in [x \mapsto \mathbf{Labeled} \top \mathbf{bool}]_V^\emptyset$

From (NI-0) we know that  $x : \mathbf{Labeled} \top \mathbf{bool} \vdash e_s : \mathbb{C} \perp \perp \mathbf{bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 5.16 to get

$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\mathbb{C} \perp \perp \mathbf{bool}]_E^\emptyset$  (NI-3.1)

Applying Definition 5.9 on (NI-3.1) we get

$\forall H_{s2}, H_{t2}.(n, H_{s2}, H_{t2}) \hat{\triangleright} \emptyset \wedge \forall i < n. e_s[s v_1/x] \Downarrow_i {}^s v \implies$   
 $\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n - i, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset$   
 Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  ${}^s v = e_s[s v_1/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset$$

From translation and from (NI-1) we know that  ${}^t v = e_t[{}^t v_1/x] = \lambda_. e_{b1}$  and therefore from fg-val we have  $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[s v_1/x], \lambda_. e_{b1}) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\emptyset}$$

Expanding  $(\emptyset, n, e_s[s v_1/x], \lambda_. e_{b1}) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\emptyset}$  using Definition 5.8 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$(k, H_{s3}, H_{t3}) \hat{\triangleright} ({}^s \theta_e) \wedge (H_{s3}, e_s[s v_1/x]) \Downarrow_i^f (H'_{s1}, {}^s v'_1) \wedge i < k \implies$

$$\exists H''_{t1}, {}^t v'', (H_{t3}, (\lambda_. e_{b1})()) \Downarrow (H''_{t1}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H''_{t1}) \hat{\triangleright} {}^s \theta' \wedge \exists {}^t v''' . {}^t v'' = \text{inl } {}^t v''' \wedge ({}^s \theta', k - i, {}^s v'_1, {}^t v''') \in [\mathbf{bool}]_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_1, {}^s v'_1, n, \emptyset$  we get

$$\exists H''_{t1}, {}^t v'' . (\emptyset, (\lambda_. e_{b1})()) \Downarrow (H''_{t1}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s1}, H''_{t1}) \hat{\triangleright} {}^s \theta' \wedge \exists {}^t v''' . {}^t v'' = \text{inl } {}^t v''' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''') \in [\mathbf{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2})$$

Since we have  $\exists {}^t v''' . {}^t v'' = \text{inl } {}^t v''' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''') \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$ , therefore from Definition 5.8 we know that 2 cases arise

- ${}^s v'_1 = \text{inl } {}^s v'_{i1}$  and  ${}^t v''' = \text{inl } {}^t v'_{i1}$ :

And from Definition 5.8 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means  ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr } {}^s v'_{i1}$  and  ${}^t v''' = \text{inr } {}^t v'_{i1}$ :

Same reasoning as in the previous case

Thus no matter which case occurs we have  ${}^s v'_1 = {}^t v'''$  (NI-3.3)

Similarly we can apply Theorem 5.16 with the other substitution to get

$$(\emptyset, n, e_s[s v_2/x], e_t[{}^t v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{bool}]_E^{\emptyset} \quad (\text{NI-4.1})$$

Applying Definition 5.9 on (NI-4.1) we get

$$\forall H_{s2}, H_{t2}.(n, H_{s2}, H_{t2}) \hat{\triangleright} \emptyset \wedge \forall i < n. {}^s v_s. e_s[s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset$$

Instantiating with  $\emptyset, \emptyset$ . From cg-val we know that  $i = 0$  and  ${}^s v_s = e_s[s v_2/x]$ .

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \hat{\triangleright} \emptyset$$

Also from (NI-1) and from translation we know that  ${}^t v = e_t[{}^t v_2/x] = \lambda_{\cdot} e_{b_2}$  and therefore from fg-val we know that  $H'_{t_2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_{\cdot} e_{b_2}) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\emptyset}$$

Expanding  $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b_2}) \in [\mathbb{C} \perp \perp \mathbf{bool}]_V^{\emptyset}$  using Definition 5.8 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s_3}, H_{t_3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s_3}, H_{t_3}) \hat{\beta}' \triangleright ({}^s \theta_e) \wedge (H_{s_3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s_2}, {}^s v'') \wedge i < k \implies$$

$$\exists H''_{t_2}, {}^t v'', (H_{t_3}, (\lambda_{\cdot} e_{b_2})()) \Downarrow (H''_{t_2}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' \cdot (k - i, H'_{s_2}, H''_{t_2}) \hat{\beta}'' \triangleright {}^s \theta' \wedge \exists {}^t v_2'''. {}^t v_2'' = \text{inl } {}^t v_2''' \wedge ({}^s \theta', k - i, {}^s v_1'', {}^t v_2''') \in [\mathbf{bool}]_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$  we get

$$\exists H''_{t_2}, {}^t v''. (\emptyset, (\lambda_{\cdot} e_{b_2})()) \Downarrow (H''_{t_2}, {}^t v'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' \cdot (n - n_1, H'_{s_2}, H''_{t_2}) \hat{\beta}'' \triangleright {}^s \theta' \wedge \exists {}^t v_2'''. {}^t v_2'' = \text{inl } {}^t v_2''' \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_2''') \in [\mathbf{bool}]_V^{\hat{\beta}''} \quad (\text{NI-4.2})$$

Since we have  $\exists {}^t v_2'''. {}^t v_2'' = \text{inl } {}^t v_2''' \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_2''') \in [\mathbf{bool}]_V^{\hat{\beta}''}$ , therefore from Definition 5.8 2 cases arise

- ${}^s v'_2 = \text{inl } {}^s v'_{i_2}$  and  ${}^t v_2''' = \text{inl } {}^t v'_{i_2}$ :

And from Definition 5.8 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i_2}, {}^t v'_{i_2}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means  ${}^s v'_{i_2} = {}^t v'_{i_2} = ()$

- ${}^s v'_2 = \text{inr } {}^s v'_{i_2}$  and  ${}^t v_2''' = \text{inr } {}^t v'_{i_2}$ :

Same reasoning as in the previous case

$$\text{Thus no matter which case occurs we have } {}^s v'_2 = {}^t v_2''' \quad (\text{NI-4.3})$$

From CG to FG translation we know that  $\exists {}^t v_{i_1}. {}^t v_1 = \text{inl } {}^t v_{i_1}$  and similarly  $\exists {}^t v_{i_2}. {}^t v_2 = \text{inl } {}^t v_{i_2}$

From (NI-1) since  $\emptyset \vdash_{\top} {}^t v_1 : (\mathbf{bool}^{\perp} + \mathbf{unit})^{\top}$  therefore from CG-inl we know that  $\emptyset \vdash_{\top} {}^t v_{i_1} : \mathbf{bool}^{\perp}$

And from CGsub-sum we know that  $\emptyset \vdash_{\top} {}^t v_{i_1} : \mathbf{bool}^{\top}$

$$\text{Therefore we also have } \emptyset \vdash_{\perp} {}^t v_{i_1} : \mathbf{bool}^{\top} \quad (\text{NI-5.1})$$

$$\text{Similarly we also have } \emptyset \vdash_{\perp} {}^t v_{i_2} : \mathbf{bool}^{\top} \quad (\text{NI-5.2})$$

Next, let  $e_T = (\lambda x : (\mathbf{bool}^{\perp} + \mathbf{unit})^{\top} . \text{case}(e_t(), y.y, z.{}^t v_b)) (\text{case}(u, -.\text{inl } \text{true}, -.\text{inl } \text{false})) : \mathbf{bool}^{\perp}$

where  $\text{true} = \text{inl } ()$  and  $\text{false} = \text{inr } ()$

We claim  $u : \mathbf{bool}^{\top} \vdash_{\perp} e_T : \mathbf{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{u : \mathbf{bool}^{\top}, - \vdash_{\perp} \text{false} : \mathbf{bool}^{\perp}}{\text{FG-inl}}}{u : \mathbf{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\mathbf{bool}^{\perp} + \mathbf{unit})^{\perp}}{\text{FG-inl}}}{u : \mathbf{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\mathbf{bool}^{\perp} + \mathbf{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{u : \text{bool}^\top, - \vdash_\perp \text{true} : \text{bool}^\perp} \text{FG-inl}}{u : \text{bool}^\top, - \vdash_\perp \text{inl true} : (\text{bool}^\perp + \text{unit})^\perp} \text{FG-inl}}{u : \text{bool}^\top, - \vdash_\perp \text{inl true} : (\text{bool}^\perp + \text{unit})^\top} \text{FGSub-base}$$

P2.1:

$$\frac{}{u : \text{bool}^\top \vdash_\perp u : \text{bool}^\top}$$

P2:

$$\frac{\frac{P2.1 \quad P2.2 \quad P2.3 \quad \frac{}{\mathcal{L} \models (\text{bool}^\perp + \text{unit})^\top \searrow \perp}}{u : \text{bool}^\top \vdash_\perp (\text{case}(u, -. \text{inl true}, -. \text{inl false})) : (\text{bool}^\perp + \text{unit})^\top}}$$

P1.2:

$$\frac{\frac{\frac{\frac{}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^\perp + \text{unit})^\perp} \text{NI-1}}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp () : \text{unit}} \text{FG-unit}}{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp \quad \mathcal{L} \models (\text{bool}^\perp + \text{unit})^\perp \searrow \perp}}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp e_t() : (\text{bool}^\perp + \text{unit})^\perp} \text{FG-app}}$$

P1.1:

$$\frac{\frac{P1.2 \quad \frac{}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top, y : \text{bool}^\perp \vdash_\perp y : \text{bool}^\perp} \text{FG-var}}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top, z : \text{unit} \vdash_\perp \text{false} : \text{bool}^\perp} \text{FG-var} \quad \frac{}{\mathcal{L} \models \text{bool}^\perp \searrow \perp}}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp} \text{FG-case}}$$

P1:

$$\frac{\frac{P1.1}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) : ((\text{bool}^\perp + \text{unit})^\top \xrightarrow{\perp} \text{bool}^\perp)^\perp}}$$

Main derivation:

$$\frac{\frac{P1 \quad P2 \quad \frac{}{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\mathcal{L} \models \text{bool}^\perp \searrow \perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -. \text{inl true}, -. \text{inl false})) : \text{bool}^\perp} \text{FG-app}}$$

Assuming  $e_{b1}()$  reduces in  $n_{t1}$  steps in (NI-3.2) and  $e_{b2}()$  reduces in  $n_{t2}$  steps in (NI-4.2).

We instantiate Theorem 5.38 with  $e_T, ^t v_{i1}, ^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$  and  $\perp$  and therefore from (NI-3.3) and (NI-4.3) we get  $^t v_1''' = ^t v_2'''$  and thus  $^s v_1' = ^s v_2'$

□

## 5.2 FG to CG translation

### 5.2.1 Type directed (direct) translation from FG to CG

**Definition 5.19.**

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket &= \mathbf{b} \\
\llbracket \text{unit} \rrbracket &= \text{unit} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket &= \llbracket \tau_1 \rrbracket \rightarrow \mathbb{C} \ell_e \perp \llbracket \tau_2 \rrbracket \\
\llbracket \forall \alpha. (\ell_e, \tau) \rrbracket &= \forall \alpha. \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket &= c \Rightarrow \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket \\
\llbracket \tau_1 \times \tau_2 \rrbracket &= \llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket \\
\llbracket \tau_1 + \tau_2 \rrbracket &= \llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket \\
\llbracket \text{ref } \mathbf{A}^\ell \rrbracket &= \text{ref } \ell \llbracket \mathbf{A} \rrbracket \\
\llbracket \mathbf{A}^\ell \rrbracket &= \text{Labeled } (\ell) \llbracket \mathbf{A} \rrbracket
\end{aligned}$$

For  $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$ , define  $\llbracket \Gamma \rrbracket = x_1 : \llbracket \tau_1 \rrbracket, \dots, x_n : \llbracket \tau_n \rrbracket$ .

We use a coercion function defined as follows:

$ \begin{aligned} \text{coerce\_taint} &: \mathbb{C} \text{ pc } \ell_c \tau' \rightarrow \mathbb{C} \text{ pc } \perp \tau' \quad \text{when } \tau' = \text{Labeled } \ell'_c \tau \text{ and } \ell_c \sqsubseteq \ell'_c \\ \text{coerce\_taint} &\triangleq \lambda x. \text{toLabeled}(\text{bind}(x, y. \text{unlabel}(y))) \end{aligned} $
--

$$\begin{aligned}
&\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{\text{pc}} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var} \\
&\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\lambda x. e_{c1}))} \text{FC-lam} \\
&\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} \Lambda e : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda e_c))} \text{FC-FI} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e : (\forall \alpha. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash_{\text{pc}} \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e [] : \tau[\ell'/\alpha] \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } a, b. (b[]))))} \text{FG-FE} \\
&\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\nu e_c))} \text{FG-CI} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash_{\text{pc}} \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } a, b. (b\bullet))))} \text{FG-CE} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{\text{pc}} e_2 : \tau_2 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup_{\text{pc}} \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a. \text{bind}(e_{c2}, b. \text{bind}(\text{unlabel } a, c. (c b)))))} \text{FC-app} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{\text{pc}} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a. \text{bind}(e_{c2}, b. \text{ret}(\text{Lb}(a, b))))} \text{FC-prod} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{\text{pc}} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } (a), b. \text{ret}(\text{fst}(b)))))} \text{FC-fst}
\end{aligned}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a)))} \text{FC-inl} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a)))} \text{FC-inr} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new}(e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb})))} \text{FC-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.!b)))} \text{FC-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel}(a), c.c := b))))), d.\text{ret}())} \text{FC-assig}
\end{array}$$

## 5.2.2 Type preservation for FG to CG translation

**Theorem 5.20** (Type preservation: FG to CG). *If  $\Gamma \vdash_{pc} e : \tau$  in FG then there exists  $e'$  such that  $\Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$  such that there is a derivation of  $(\Gamma) \vdash e' : \mathbb{C} pc \perp (\tau)$  in CG.*

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\begin{array}{c}
\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var} \\
\\
\frac{}{(\Gamma), x : (\tau) \vdash x : (\tau)} \text{CG-var} \\
\frac{}{(\Gamma), x : (\tau) \vdash \text{ret } x : \mathbb{C} pc \perp (\tau)} \text{CG-ret}
\end{array}$$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\lambda x.e_{c1}))} \text{FC-lam}$$

$$T_0 = \mathbb{C} pc \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) = \mathbb{C} pc \perp \text{Labeled} \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2))$$

$$T_1 = \mathbb{C} pc \perp \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1,0} = \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \ell_e \perp (\tau_2)$$

P1:

$$\frac{\frac{P2}{(\Gamma), x : (\tau_1) \vdash e_{c1} : T_{1.2}} \text{IH}}{(\Gamma) \vdash \lambda x. e_{c1} : T_{1.1}} \text{CG-lam}$$

Main derivation:

$$\frac{\frac{P1}{(\Gamma) \vdash (\text{Lb}(\lambda x. e_{c1})) : T_{1.0}} \text{CG-label}}{(\Gamma) \vdash \text{ret}(\text{Lb}(\lambda x. e_{c1})) : T_1} \text{CG-ret}$$

3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))))} \text{FC-app}$$

$$T_0 = \mathbb{C} pc \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell) = \mathbb{C} pc \perp \text{Labeled } \ell ((\tau_1 \xrightarrow{\ell_e} \tau_2))$$

$$T_1 = \mathbb{C} pc \perp \text{Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = \text{Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \top \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.3} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.4} = \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.5} = \mathbb{C} \ell_e \ell (\tau_2)$$

$$T_{1.6} = \mathbb{C} pc \ell (\mathbb{A}^{\ell_i})$$

$$T_{1.7} = \mathbb{C} pc \ell \text{Labeled } (\ell_i) (\mathbb{A})$$

$$T_{1.9} = \mathbb{C} pc \perp \text{Labeled } \ell_i (\mathbb{A})$$

$$T_{1.10} = \mathbb{C} pc \perp (\tau_2)$$

$$T_2 = \mathbb{C} pc \perp (\tau_1)$$

$$T_{c4} = \text{Labeled } \ell_i (\mathbb{A})$$

$$T_{c3} = \mathbb{C} \top \ell_i (\mathbb{A})$$

$$T_{c2} = \mathbb{C} pc \ell_i (\mathbb{A})$$

$$T_{c1} = \mathbb{C} pc \perp \text{Labeled } \ell_i (\mathbb{A})$$

$$T_{c0} = \mathbb{C} pc \ell \text{Labeled } \ell_i (\mathbb{A})$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}}$$

Pc1:

$$\frac{}{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$



Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{P0}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P6:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash b : (\tau_1)} \text{CG-var}$$

P5:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash c : T_{1.3}} \text{CG-var}$$

P4:

$$\frac{\frac{P5 \quad P6}{(\Gamma), a : T_{1.1}, b : (\tau_2), c : T_{1.3} \vdash c b : T_{1.4}} \text{CG-app}}{(\Gamma), a : T_{1.1}, b : (\tau_2), c : T_{1.3} \vdash c b : T_{1.5}} \text{CGSub-monad}$$

P3:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1) \vdash a : T_{1.1}} \text{CG-var}$$

P2:

$$\frac{\frac{P3}{(\Gamma), a : T_{1.1}, b : (\tau_1) \vdash \text{unlabel } a : T_{1.2}} \text{CG-unlabel} \quad P4}{(\Gamma), a : T_{1.1}, b : (\tau_1) \vdash \text{bind}(\text{unlabel } a, c.(c b)) : T_{1.6}} \text{CG-bind}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{1.1} \vdash e_{c2} : T_2} \text{IH2, Weakening} \quad P2}{(\Gamma), a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b))) : T_{1.6}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Main derivation:

$$\frac{Pc \quad \frac{\frac{}{(\Gamma) \vdash e_{c1} : T_1} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.7}} \text{CG-bind}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.9}} \text{CG-app}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.10}} \text{Definition 5.19}$$

4. FC-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda e_c))} \text{FC-FI}$$

$$T_0 = \mathbb{C} \text{ pc } \perp \langle (\forall \alpha. (\ell_e, \tau))^\perp \rangle = \mathbb{C} \text{ pc } \perp \text{Labeled } \perp \langle (\forall \alpha. (\ell_e, \tau)) \rangle$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\text{Labeled } \perp (\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle))$$

$$T_{1.0} = \text{Labeled } \perp (\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle)$$

$$T_{1.1} = \forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle$$

P1:

$$\frac{\frac{P2}{\Sigma, \alpha; \Psi; \langle \Gamma \rangle \vdash e_c : \langle \tau \rangle} \text{IH}}{\Sigma; \Psi; \langle \Gamma \rangle \vdash \Lambda e_c : T_{1.1}} \text{CG-lam}$$

Main derivation:

$$\frac{\frac{P1}{\Sigma; \Psi; \langle \Gamma \rangle \vdash \text{Lb}(\Lambda e_c) : T_{1.0}} \text{CG-label}}{\Sigma; \Psi; \langle \Gamma \rangle \vdash \text{ret}(\text{Lb}(\Lambda e_c)) : T_1} \text{CG-ret, CG-sub}$$

5. FC-FE:

$$\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash_{pc} \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha] \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel } a, b. (b[]))))} \text{FG-FE}$$

$$T_0 = \mathbb{C} \text{ pc } \perp \langle (\forall \alpha. (\ell_e, \tau))^\ell \rangle = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\forall \alpha. (\ell_e, \tau)) \rangle$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\text{Labeled } \ell (\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle))$$

$$T_{1.1} = (\text{Labeled } \ell (\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle))$$

$$T_{1.9} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i[\ell'/\alpha] \langle \mathbf{A} \rangle[\ell'/\alpha]$$

$$T_{1.10} = \mathbb{C} \text{ pc } \perp \langle \tau[\ell'/\alpha] \rangle$$

$$T_2 = \mathbb{C} \top \ell (\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle)$$

$$T_{2.1} = \forall \alpha. \mathbb{C} \ell_e \perp \langle \tau \rangle$$

$$T_{2.2} = (\mathbb{C} \ell_e \perp \langle \tau \rangle)[\ell'/\alpha]$$

$$T_{2.3} = \mathbb{C} \ell_e[\ell'/\alpha] \perp \langle \tau \rangle[\ell'/\alpha]$$

$$T_{2.4} = \mathbb{C} \text{ pc } \ell \langle \mathbf{A}^{\ell_i} \rangle[\ell'/\alpha]$$

$$T_{2.5} = \mathbb{C} \text{ pc } \ell \text{Labeled } (\ell_i[\ell'/\alpha]) \langle \mathbf{A} \rangle[\ell'/\alpha]$$

$$T_{c4} = \text{Labeled } \ell_i \langle \mathbf{A} \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle \mathbf{A} \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle \mathbf{A} \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle \mathbf{A} \rangle$$

$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \text{ (A)}$

$T_c = T_{c0} \rightarrow T_{c1}$

Pc2:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}}{\Sigma; \Psi; (\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}$$

Pc1:

$$\frac{}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{P0}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\Sigma; \Psi; (\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P4:

$$\frac{}{\Sigma; \Psi; (\Gamma), a : T_{1.1}, b : T_{2.1} \vdash b[] : T_{2.3}} \text{CG-FE}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma), a : T_{1.1} \vdash \text{unlabel } a : T_2} \text{P2}}{\Sigma; \Psi; (\Gamma), a : T_{1.1} \vdash \text{bind}(\text{unlabel } a, b.(b[])) : T_{2.5}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Main derivation:

$$\frac{\frac{Pc \quad \frac{\frac{}{\Sigma; \Psi; (\Gamma) \vdash e_c : T_1} \text{IH1} \quad P1}{\Sigma; \Psi; (\Gamma) \vdash (\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b[])))} : T_{2.5}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b[])))} : T_{1.9}} \text{CG-app}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b[])))} : T_{1.10}} \text{Lemma 5.24 and Def 5.19}$$

6. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\nu e_c))} \text{FG-CI}$$

$$T_0 = \mathbb{C} \text{ pc } \perp \llbracket (c \xrightarrow{\ell_e} \tau)^\perp \rrbracket = \mathbb{C} \text{ pc } \perp (\text{Labeled } \perp \llbracket (c \xrightarrow{\ell_e} \tau) \rrbracket)$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\text{Labeled } \perp (c \Rightarrow \mathbb{C} \ell_e \perp (\tau)))$$

$$T_{1.0} = \text{Labeled } \perp (c \Rightarrow \mathbb{C} \ell_e \perp (\tau))$$

$$T_{1.1} = c \Rightarrow \mathbb{C} \ell_e \perp (\tau)$$

P1:

$$\frac{\frac{P2}{\Sigma; \Psi, c; (\Gamma) \vdash e_c : (\tau)} \text{IH}}{\Sigma; \Psi; (\Gamma) \vdash \nu e_c : T_{1.1}} \text{CG-CI}$$

Main derivation:

$$\frac{\frac{P1}{\Sigma; \Psi; (\Gamma) \vdash \text{Lb}(\nu e_c) : T_{1.0}} \text{CG-label}}{\Sigma; \Psi; (\Gamma) \vdash \text{ret}(\text{Lb}(\nu e_c)) : T_1} \text{CG-ret,CG-sub}$$

7. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b\bullet)))} \text{FG-CE}$$

$$T_0 = \mathbb{C} \text{ pc } \perp ((c \xrightarrow{\ell_e} \tau)^\ell) = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell ((c \xrightarrow{\ell_e} \tau))$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\text{Labeled } \ell (c \Rightarrow \mathbb{C} \ell_e \perp (\tau)))$$

$$T_{1.1} = (\text{Labeled } \ell (c \Rightarrow \mathbb{C} \ell_e \perp (\tau)))$$

$$T_{1.9} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i (\mathbf{A})$$

$$T_{1.10} = \mathbb{C} \text{ pc } \perp (\tau)$$

$$T_2 = \mathbb{C} \top \ell (c \Rightarrow \mathbb{C} \ell_e \perp (\tau))$$

$$T_{2.1} = c \Rightarrow \mathbb{C} \ell_e \perp (\tau)$$

$$T_{2.2} = \mathbb{C} \ell_e \perp (\tau)$$

$$T_{2.4} = \mathbb{C} \text{ pc } \ell (\mathbf{A}^{\ell_i})$$

$$T_{2.5} = \mathbb{C} \text{ pc } \ell \text{Labeled } (\ell_i) (\mathbf{A})$$

$$T_{c4} = \text{Labeled } \ell_i (\mathbf{A})$$

$$T_{c3} = \mathbb{C} \top \ell_i (\mathbf{A})$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i (\mathbf{A})$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i (\mathbf{A})$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i (\mathbf{A})$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{\Sigma; \Psi; (\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}}{\Sigma; \Psi; (\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}}{\Sigma; \Psi; (\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}$$

Pc1:

$$\frac{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash x : T_{c0}}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{P0}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\Sigma; \Psi; (\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P4:

$$\frac{}{\Sigma; \Psi; (\Gamma), a : T_{1.1}, b : T_{2.1} \vdash b \bullet : T_{2.2}} \text{CG-CE}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma), a : T_{1.1} \vdash \text{unlabel } a : T_2} P2}{\Sigma; \Psi; (\Gamma), a : T_{1.1} \vdash \text{bind}(\text{unlabel } a, b.(b \bullet)) : T_{2.5}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Main derivation:

$$\frac{Pc \quad \frac{\frac{}{\Sigma; \Psi; (\Gamma) \vdash e_c : T_1} \text{IH1} \quad P1}{\Sigma; \Psi; (\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b \bullet))) : T_{2.5}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b \bullet)))) : T_{1.9}} \text{CG-app}}{\Sigma; \Psi; (\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.(b \bullet)))) : T_{1.10}} \text{Definition 5.19}$$

8. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \mathbb{C} \text{ pc } \perp ((\tau_1 \times \tau_2)^\perp)$$

$$T_2 = \mathbb{C} \text{ pc } \perp \text{Labeled } \perp ((\tau_1 \times \tau_2))$$

$$T_3 = \mathbb{C} \text{ pc } \perp \text{Labeled } \perp ((\tau_1) \times ((\tau_2)))$$

$$T_{3.1} = \text{Labeled } \perp ((\tau_1) \times ((\tau_2)))$$

$$T_4 = \mathbb{C} \text{ pc } \perp ((\tau_1))$$

$$T_5 = \mathbb{C} \text{ pc } \perp ((\tau_2))$$

P4:

$$\frac{}{(\Gamma), a : ((\tau_1)), b : ((\tau_1)) \vdash a : ((\tau_1))} \text{CG-var}$$

P3:

$$\frac{}{(\Gamma), a : (\tau_1), b : (\tau_1) \vdash b : (\tau_2)} \text{CG-var}$$

P2:

$$\frac{\frac{\frac{P3 \quad P4}{(\Gamma), a : (\tau_1), b : (\tau_1) \vdash (a, b) : (\tau_1) \times (\tau_2)}{\text{CG-prod}}}{(\Gamma), a : (\tau_1), b : (\tau_2) \vdash \text{Lb}(a, b) : T_{3.1}} \text{CG-label}}{(\Gamma), a : (\tau_1), b : (\tau_2) \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{}{(\Gamma), a : (\tau_1) \vdash e_{c2} : T_5} \text{IH2} \quad P2}{(\Gamma), a : (\tau_1) \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_{c1} : T_4} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1} \text{Definition 5.19}$$

9. FC-fst:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))} \text{FC-fst}$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\tau_1)$$

$$T_2 = \mathbb{C} \text{ pc } \perp ((\tau_1 \times \tau_2)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell ((\tau_1 \times \tau_2))$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell (\tau_1) \times (\tau_2)$$

$$T_{2.3} = \text{Labeled } \ell (\tau_1) \times (\tau_2)$$

$$T_{2.4} = (\tau_1) \times (\tau_2)$$

$$T_{2.5} = \mathbb{C} \top \ell (\tau_1) \times (\tau_2)$$

$$T_3 = \mathbb{C} \top \ell (\tau_1)$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell (\tau_1)$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell (A^{\ell_i})$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i (A)$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i (A)$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp (A^{\ell_i})$$

$$T_{c4} = \text{Labeled } \ell_i (A)$$

$$T_{c3} = \mathbb{C} \top \ell_i (A)$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i (A)$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i (A)$$

$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \text{ (A)}$

$T_c = T_{c0} \rightarrow T_{c1}$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau_1 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{ CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{ From Definition of coerce\_taint}$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{ CG-var}}{\frac{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : \langle \tau_1 \rangle}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{ CG-ret}} \text{ CG-fst}$$

P1:

$$\frac{\overline{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{ CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1}} \text{ CG-bind}$$

P0:

$$\frac{\overline{(\Gamma) \vdash e_c : T_{2.2}} \text{ IH} \quad P1}{\overline{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1}} \text{ CG-bind}}{\frac{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2}}{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3}} \text{ Definition 5.19}}$$

Main derivation:

$$\frac{\overline{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.5}} \text{ CG-app}}{\overline{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.6}} \text{ Definition 5.19}}{\overline{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_1}}$$

10. FC-snd:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau_2 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\tau_1 \times \tau_2) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.3} = \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.4} = \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} \top \ell \langle \tau_2 \rangle$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell \langle \tau_2 \rangle$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell \langle A^{\ell_i} \rangle$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle A \rangle$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle A \rangle$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp \langle A^{\ell_i} \rangle$$

$$T_{c4} = \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle A \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle A \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle A \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \quad \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Pc2:

$$\frac{\overline{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \quad \text{CG-var}}{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}$$

Pc1:

$$\overline{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\overline{\overline{\overline{Pc1} \quad Pc2} \quad \frac{Pg}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i}} \quad \text{CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-tolabeled}} \langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}$$



Pc:

$$\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x. \text{toLabeled}(\text{bind}(x, y. \text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P2:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : \langle \tau_2 \rangle} \text{CG-snd}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{CG-ret}}$$

P1:

$$\frac{\frac{\frac{}{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b))) : T_{3.1}} \text{CG-bind}}$$

P0:

$$\frac{\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_{2.2}} \text{IH} \quad P1}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.1}} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.2}}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 5.19}$$

Main derivation:

$$\frac{\frac{\frac{\frac{Pc \quad P0}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.5}} \text{CG-app}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.6}} \text{Definition 5.19}}{(\Gamma) \vdash \text{coerce\_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_1}}$$

11. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a. \text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \mathbb{C} \text{ pc} \perp ((\tau_1 + \tau_2)^\perp)$$

$$T_{1.1} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp ((\tau_1 + \tau_2))$$

$$T_{1.2} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp (\langle \tau_1 \rangle) + (\langle \tau_2 \rangle)$$

$$T_{1.3} = \text{Labeled} \perp (\langle \tau_1 \rangle) + (\langle \tau_2 \rangle)$$

$$T_2 = \mathbb{C} \text{ pc} \perp (\langle \tau_1 \rangle)$$

P1:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : \langle \tau_1 \rangle \vdash a : \langle \tau_1 \rangle} \text{CG-var}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{inl}(a) : \langle \tau_1 \rangle + \langle \tau_2 \rangle} \text{CG-inl}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{Lbinl}(a) : T_{1.3}} \text{CG-label}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_2} \text{IH} \quad P1}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_{1.2}} \text{CG-bind}}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_1} \text{Definition 5.19}}$$

12. FC-inr:

$$\frac{\Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \mathbb{C} \text{ pc } \perp ((\tau_1 + \tau_2)^\perp)$$

$$T_{1.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp ((\tau_1 + \tau_2))$$

$$T_{1.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_{1.3} = \mathbf{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_2 = \mathbb{C} \text{ pc } \perp (\tau_2)$$

P1:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : (\tau_2) \vdash a : (\tau_2)} \text{CG-var}}{(\Gamma), a : (\tau_2) \vdash \mathbf{inr}(a) : (\tau_1) + (\tau_2)} \text{CG-inr}}{(\Gamma), a : (\tau_2) \vdash \mathbf{Lbinr}(a) : T_{1.3}} \text{CG-label}}{(\Gamma), a : (\tau_2) \vdash \mathbf{ret}(\mathbf{Lbinr}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_2} \text{IH} \quad P1}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_{1.2}} \text{CG-bind}}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_1} \text{Definition 5.19}}$$

13. FC-case:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \mathbf{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_c, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.\mathbf{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case}$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\tau)$$

$$T_2 = \mathbb{C} \text{ pc } \perp ((\tau_1 + \tau_2)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ (\tau_1 + \tau_2)$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.3} = \mathbf{Labeled} \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.4} = \mathbb{C} \top \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.5} = (\tau_1) + (\tau_2)$$

$$T_3 = \mathbb{C} (\text{pc } \sqcup \ell) \perp (\tau)$$

$$\begin{aligned}
T_4 &= \mathbb{C} (pc \sqcup \ell) \ell (\tau) \\
T_5 &= \mathbb{C} (pc) \ell (\mathbb{A}^{\ell_i}) \\
T_{5.1} &= \mathbb{C} (pc) \ell \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{5.3} &= \mathbb{C} (pc) (\perp) \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{5.4} &= \mathbb{C} (pc) (\perp) (\mathbb{A}^{\ell_i}) \\
T_{c4} &= \text{Labeled } \ell_i (\mathbb{A}) \\
T_{c3} &= \mathbb{C} \top \ell_i (\mathbb{A}) \\
T_{c2} &= \mathbb{C} pc \ell_i (\mathbb{A}) \\
T_{c1} &= \mathbb{C} pc \perp \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{c0} &= \mathbb{C} pc \ell \text{ Labeled } \ell_i (\mathbb{A}) \\
T_c &= T_{c0} \rightarrow T_{c1}
\end{aligned}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \overline{\mathcal{L} \models \ell \sqsubseteq \ell_i} \text{ Pg}}{\overline{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{ Pc0} \text{ CG-lam}}{(\Gamma) \vdash \text{coerce\_taint} : T_c} \text{ From Definition of coerce\_taint}$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}} \text{ CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.5}, x : (\tau_1) \vdash e_{c1} : T_3} \text{ IH2, Weakening} \\
\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.5}, y : (\tau_2) \vdash e_{c2} : T_3} \text{ IH3, Weakening}}{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash \text{case}(b, x.e_{c1}, y.e_{c2}) : T_3} \text{ CG-case}$$

$$\begin{array}{c}
\text{P1:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_3} \text{CG-bind}}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_4} \text{CG-sub}
\end{array}$$

$$\begin{array}{c}
\text{P0:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.2}}{\text{IH1}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_5} \text{CG-bind}
\end{array}$$

$$\begin{array}{c}
\text{P0.2:} \\
\frac{P0}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_{5.1}} \text{Definition 5.19}
\end{array}$$

$$\begin{array}{c}
\text{P0.1:} \\
\frac{\frac{\frac{P_c \quad P0.2}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.3}}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.4}} \text{Definition 5.19}
\end{array}$$

Main derivation:

$$\frac{P0.1}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_1}$$

14. FC-ref:

$$\frac{\Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow_{pc}}{\Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})))} \text{FC-ref}$$

$$\begin{aligned}
T_1 &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } \tau)^\perp \rangle \\
T_{1.1} &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } A^{\ell_i})^\perp \rangle \\
T_{1.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \langle (\text{ref } A^{\ell_i}) \rangle \\
T_{1.3} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle \\
T_2 &= \mathbb{C} \text{ } pc \perp \langle \tau \rangle \\
T_{2.1} &= \mathbb{C} \text{ } pc \perp \langle A^{\ell_i} \rangle \\
T_{2.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled } \ell_i \langle A \rangle \\
T_{2.3} &= \text{Labeled } \ell_i \langle A \rangle \\
T_{2.4} &= \mathbb{C} \text{ } pc \perp \text{ref } \ell_i \langle A \rangle \\
T_{2.5} &= \text{ref } \ell_i \langle A \rangle \\
T_{2.51} &= \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle
\end{aligned}$$

$$\begin{array}{c}
\text{P2:} \\
\frac{\frac{\frac{}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}} \text{CG-label}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}} \text{CG-ret}
\end{array}$$

P1:

$$\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{new}(a) : T_{2.4}} \text{CG-new} \quad P2}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{CG-bind}}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{CG-bind}}$$

Main derivation:

$$\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{IH} \quad P1}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{CG-bind}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb}))) : T_{1.3}} \text{CG-bind}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lbb}))) : T_1} \text{Definition 5.19}}$$

15. FC-deref:

$$\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\begin{aligned} T_1 &= \mathbb{C} \text{ pc } \perp \langle \tau' \rangle \\ T_{1.1} &= \mathbb{C} \text{ pc } \perp \langle A'^{\ell'_i} \rangle \\ T_{1.2} &= \mathbb{C} \text{ pc } \perp \text{Labeled } \ell'_i \langle A' \rangle \\ T_2 &= \mathbb{C} \text{ pc } \perp \langle (\text{ref } \tau)^\ell \rangle \\ T_{2.1} &= \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } \tau) \rangle \\ T_{2.2} &= \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } A^{\ell_i}) \rangle \\ T_{2.3} &= \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle \\ T_{2.4} &= \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle \\ T_{2.5} &= \mathbb{C} \top \ell \langle \text{ref } \ell_i \langle A \rangle \rangle \\ T_{2.6} &= \text{ref } \ell_i \langle A \rangle \\ T_{2.7} &= \mathbb{C} \top \perp \langle \text{Labeled } \ell_i \langle A \rangle \rangle \\ T_{2.8} &= \mathbb{C} \top \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle \\ T_{2.9} &= \mathbb{C} \text{ pc } \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle \\ T_{c4} &= \text{Labeled } \ell_i \langle A \rangle \\ T_{c3} &= \mathbb{C} \top \ell_i \langle A \rangle \\ T_{c2} &= \mathbb{C} \text{ pc } \ell_i \langle A \rangle \\ T_{c1} &= \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle A \rangle \\ T_{c0} &= \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle A \rangle \\ T_c &= T_{c0} \rightarrow T_{c1} \end{aligned}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau' = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}}$$

Pc1:

$$\frac{}{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\langle \Gamma \rangle \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\langle \Gamma \rangle \vdash \text{coerce\_taint} : T_c} \text{From Definition of coerce\_taint}$$

P2:

$$\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}} \text{CG-var}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{CG-deref}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.8}} \text{CG-sub, Lemma 5.21}$$

P1:

$$\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}} \text{CG-unlabel} \quad P2}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.3}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{Pc \quad P0}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.2}} \text{CG-app}}{\langle \Gamma \rangle \vdash \text{coerce\_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{Definition 5.19}$$

16. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$T_1 = \mathbb{C} \text{ } pc \perp (\text{unit})$$

$$T_{1.1} = \mathbb{C} \text{ } pc \perp \text{unit}$$

$$T_2 = \mathbb{C} \text{ } pc \perp ((\text{ref } \tau)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } \tau))$$

$$T_{2.2} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } A^{\ell_i}))$$

$T_{2.3} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.4} = \text{Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.5} = \mathbb{C} \top (\ell) \text{ ref } \ell_i \text{ (A)}$

$T_{2.6} = \text{ref } \ell_i \text{ (A)}$

$T_{2.7} = \mathbb{C} (\text{pc } \sqcup \ell) \perp \text{unit}$

$T_{2.71} = \mathbb{C} (\text{pc } \sqcup \ell) \ell \text{unit}$

$T_{2.8} = \mathbb{C} \text{ pc } (\ell) \text{unit}$

$T_{2.9} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \text{unit}$

$T_3 = \mathbb{C} \text{ pc } \perp (\tau)$

$T_{3.1} = \mathbb{C} \text{ pc } \perp (\text{A}^{\ell_i})$

$T_{3.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \text{ (A)}$

$T_{3.3} = \text{Labeled } \ell_i \text{ (A)}$

P4:

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{CG-var}$$

P5:

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}} \text{CG-var}$$

P3:

$$\frac{\begin{array}{c} \text{Given} \\ \mathcal{L} \vdash \tau \searrow (\text{pc } \sqcup \ell) \\ \text{By inversion} \\ \mathcal{L} \vdash (\text{pc } \sqcup \ell) \sqsubseteq \ell_i \end{array}}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}} \text{CG-assign} \quad \text{CGsub-monad}$$

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.71}}$$

P2:

$$\frac{\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}} \text{CG-unlabel} \quad P3}{(\Gamma), a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{CG-bind}}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{2.4} \vdash e_{c2} : T_{3.2}} \text{IH2} \quad P2}{(\Gamma), a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{CG-bind}}$$

P0:

$$\frac{\frac{}{(\Gamma) \vdash e_{c1} : T_{2.3}} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.8}} \text{CG-bind}}$$

P0.1:

$$\frac{P0}{(\Gamma) \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}} \text{CG-toLabeled}}$$

Main derivation:

$$\frac{P0.1 \quad \frac{}{(\Gamma), d : \text{Labeled } \ell \text{ unit} \vdash \text{ret}() : T_{1.1}}{(\Gamma) \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) , d.\text{ret}()) : T_{1.1}} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) , d.\text{ret}()) : T_{1.1}}$$

□

**Lemma 5.21** (Subtyping - Type preservation).  $\forall \Sigma; \Psi$ .

The following holds:

1.  $\forall \tau, \tau'$ .

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \Sigma; \Psi \vdash \langle \tau \rangle <: \langle \tau' \rangle$$

2.  $\forall A, A'$ .

$$\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \langle A \rangle <: \langle A' \rangle$$

*Proof.* Proof by simultaneous induction on  $\tau <: \tau$  and  $A <: A$

Proof of statement (1)

Let  $\tau = A_1^{\ell_1}$  and  $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\frac{}{A_1^{\ell_1} <: A_2^{\ell_2}}{\text{Given}}}{\Sigma; \Psi \vdash A_1 <: A_2} \text{By inversion} \quad P1}{\Sigma; \Psi \vdash \langle A_1 \rangle <: \langle A_2 \rangle} \text{IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\frac{}{A_1^{\ell_1} <: A_2^{\ell_2}}{\text{Given}}}{\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2} \text{By inversion}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \text{Labeled } \ell_1 \langle A_1 \rangle <: \text{Labeled } \ell_2 \langle A_2 \rangle} \text{CGsub-labeled}}{\Sigma; \Psi \vdash \langle A_1^{\ell_1} \rangle <: \langle A_2^{\ell_2} \rangle}$$

Proof of statement (2)

We proceed by cases on  $A <: A$

1. FGsub-base:

$$\frac{\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{CG-refl}}{\Sigma; \Psi \vdash \langle \mathbf{b} \rangle <: \langle \mathbf{b} \rangle} \text{Definition 5.19}$$

2. FGsub-ref:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{ref } \ell_i \langle A \rangle <: \text{ref } \ell_i \langle A \rangle} \text{CG-refl}}{\Sigma; \Psi \vdash \langle \text{ref } A^{\ell_i} \rangle <: \langle \text{ref } A^{\ell_i} \rangle} \text{Definition 5.19}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle <: \langle \tau'_1 \rangle} \text{IH(1) on } \tau_1 <: \tau'_1$$



P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash (\tau_2) <: (\tau'_2)} \text{ By inversion} \quad \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash (\tau_1) \times (\tau_2) <: (\tau'_1) \times (\tau'_2)} \text{ CGsub-prod}}{\Sigma; \Psi \vdash (\tau_1 \times \tau_2) <: (\tau'_1 \times \tau'_2)} \text{ Definition 5.19}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ Given}}{\Sigma; \Psi \vdash (\tau_1) <: (\tau'_1)} \text{ By inversion} \quad \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash (\tau_2) <: (\tau'_2)} \text{ By inversion} \quad \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash (\tau_1) + (\tau_2) <: (\tau'_1) + (\tau'_2)} \text{ CGsub-prod}}{\Sigma; \Psi \vdash (\tau_1 + \tau_2) <: (\tau'_1 + \tau'_2)} \text{ Definition 5.19}$$

5. FGsub-arrow:

$$T_1 = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_2 = (\tau'_1) \rightarrow \mathbb{C} \ell'_e \perp (\tau'_2)$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\zeta} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\zeta} \tau'_2}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion, Weakening} \quad \frac{\overline{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\zeta} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\zeta} \tau'_2}}{\Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{ Given}}{\Sigma; \Psi \vdash \mathbb{C} \ell_e \perp (\tau_2) <: \mathbb{C} \ell'_e \perp (\tau'_2)} \text{ IH(1), CGsub-monad}$$

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\zeta} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\zeta} \tau'_2}}{\Sigma; \Psi \vdash \tau'_1 <: \tau_1} \text{ Given}}{\Sigma; \Psi \vdash (\tau'_1) <: (\tau_1)} \text{ By inversion, Weakening} \quad \text{IH(1)}$$

Main derivation:

$$\frac{P1 \quad P2}{\Sigma; \Psi \vdash (\tau_1 \xrightarrow{\ell_e} \tau_2) <: (\tau'_1 \xrightarrow{\ell'_e} \tau'_2)} \text{ Definition 5.19}$$

6. FGsub-forall:

P1:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau) <: \forall \alpha. (\ell'_e, \tau') \text{ Given}}{\Sigma, \alpha; \Psi \vdash \tau <: \tau'} \text{ By inversion}}{\Sigma, \alpha; \Psi \vdash (\tau) <: (\tau')} \text{ IH(1)}}{\frac{\frac{\frac{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau) <: \forall \alpha. (\ell'_e, \tau') \text{ Given}}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{ By inversion}}{\Sigma, \alpha; \Psi \vdash \mathbb{C} \ell_e \perp (\tau) <: \mathbb{C} \ell'_e \perp (\tau')} \text{ CGsub-monad}}}$$

Main derivation:

$$\frac{P1}{\frac{\Sigma; \Psi \vdash \forall \alpha. \mathbb{C} \ell_e \perp (\tau) <: \forall \alpha. \mathbb{C} \ell'_e \perp (\tau')}{\Sigma; \Psi \vdash (\forall \alpha. (\ell_e, \tau)) <: (\forall \alpha. (\ell'_e, \tau'))} \text{ Definition 5.19}}$$

7. FGsub-constraint:

P1:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash c \xrightarrow{\ell_e} \tau <: c' \xrightarrow{\ell'_e} \tau'}{\Sigma; \Psi \vdash \tau <: \tau'} \text{ By inversion}}{\Sigma; \Psi \vdash (\tau) <: (\tau')} \text{ IH(1)}}{\frac{\frac{\frac{\Sigma; \Psi \vdash c \xrightarrow{\ell_e} \tau <: c' \xrightarrow{\ell'_e} \tau'}{\Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{ By inversion}}{\Sigma; \Psi \vdash \mathbb{C} \ell_e \perp (\tau) <: \mathbb{C} \ell'_e \perp (\tau')} \text{ CGsub-monad}}}$$

P0:

$$\frac{\frac{\Sigma; \Psi \vdash c \xrightarrow{\ell_e} \tau <: c' \xrightarrow{\ell'_e} \tau'}{\Sigma; \Psi \vdash c' \Rightarrow c} \text{ By inversion}}{\text{ Given}}$$

Main derivation:

$$\frac{P0 \quad P1}{\frac{\Sigma; \Psi \vdash c \Rightarrow \mathbb{C} \ell_e \perp (\tau) <: c' \Rightarrow \mathbb{C} \ell'_e \perp (\tau')}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_e} \tau) <: (c' \xrightarrow{\ell'_e} \tau')} \text{ Definition 5.19}}$$

8. FGsub-unit:

$$\frac{\text{CGsub-unit}}{\frac{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}{\Sigma; \Psi \vdash (\text{unit}) <: (\text{unit})} \text{ Definition 5.19}}$$

□

**Lemma 5.22** (FG  $\rightsquigarrow$  CG: Preservation of well-formedness). *For all  $\Sigma, \Psi$  the following hold:*

$$1. \forall \tau. \Sigma; \Psi \vdash \tau \text{ WF} \implies \Sigma; \Psi \vdash \langle \tau \rangle \text{ WF}$$

$$2. \forall A. \Sigma; \Psi \vdash A \text{ WF} \implies \Sigma; \Psi \vdash \langle A \rangle \text{ WF}$$

*Proof.* Proof by simultaneous induction on the  $WF$  relation of FG

Proof of statement (1)

Let  $\tau = A^{\ell'}$

$$\frac{\frac{\Sigma; \Psi \vdash \langle A \rangle \text{ WF} \quad \text{IH(2) on } A \quad \frac{\text{FV}(\ell') \in \Sigma}{\text{By inversion}}}{\Sigma; \Psi \vdash \text{Labeled } \ell' \langle A \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle A^{\ell'} \rangle \text{ WF}} \text{CG-wff-labeled}$$

Proof of statement (2)

We proceed by case analyzing the last rule of given  $WF$  judgment.

1. FG-wff-base:

$$\frac{}{\Sigma; \Psi \vdash b \text{ WF}} \text{CG-wff-base}$$

2. FG-wff-unit:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{CG-wff-unit}$$

3. FG-wff-arrow:

P0:

$$\frac{\frac{\Sigma; \Psi \vdash \langle \tau_2 \rangle \text{ WF} \quad \text{IH(1) on } \tau_2 \quad \frac{\text{FV}(\ell_e) \in \Sigma}{\text{By inversion}}}{\Sigma; \Psi \vdash \mathbb{C} \ell_e \perp \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \mathbb{C} \ell_e \perp \langle \tau_2 \rangle \text{ WF}} \text{CG-wff-monad}$$

Main derivation:

$$\frac{\frac{\Sigma; \Psi \vdash \langle \tau_1 \rangle \text{ WF} \quad \text{IH(1) on } \tau_1 \quad P0}{\Sigma; \Psi \vdash \langle \tau_1 \rangle \rightarrow \mathbb{C} \ell_e \perp \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle \rightarrow \mathbb{C} \ell_e \perp \langle \tau_2 \rangle \text{ WF}} \text{CG-wff-arrow}$$

4. FG-wff-prod:

$$\frac{\frac{\Sigma; \Psi \vdash \langle \tau_1 \rangle \text{ WF} \quad \text{IH(1) on } \tau_1 \quad \frac{\Sigma; \Psi \vdash \langle \tau_2 \rangle \text{ WF} \quad \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \langle \tau_1 \rangle \times \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle \times \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle \times \langle \tau_2 \rangle \text{ WF}} \text{CG-wff-prod}$$

5. FG-wff-sum:

$$\frac{\frac{\Sigma; \Psi \vdash \langle \tau_1 \rangle \text{ WF} \quad \text{IH(1) on } \tau_1 \quad \frac{\Sigma; \Psi \vdash \langle \tau_2 \rangle \text{ WF} \quad \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \langle \tau_1 \rangle + \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle + \langle \tau_2 \rangle \text{ WF}}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle + \langle \tau_2 \rangle \text{ WF}} \text{CG-wff-prod}$$

6. FG-wff-ref:

Let  $\tau = \mathbf{A}^{\ell'}$

$$\frac{\frac{\overline{FV(\mathbf{A}) = \emptyset} \text{ By inversion} \quad \overline{FV(\ell') = \emptyset} \text{ By inversion}}{\overline{FV(\llbracket \mathbf{A} \rrbracket) = \emptyset} \text{ Lemma 5.23}}}{\Sigma; \Psi \vdash \text{ref } \ell' \llbracket \mathbf{A} \rrbracket \text{ WF}} \text{ CG-wff-ref}$$

7. FG-wff-forall:

$$\frac{\frac{\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket \text{ WF}} \text{ IH(1) on } \tau \quad \overline{FV(\ell_e) \in \Sigma \cup \{\alpha\}} \text{ By inversion}}{\overline{\Sigma, \alpha; \Psi \vdash \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket \text{ WF}} \text{ CG-wff-monad}}}{\overline{\Sigma; \Psi \vdash (\forall \alpha. \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket) \text{ WF}} \text{ CG-wff-forall}}$$

8. FG-wff-constraint:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi, c \vdash \llbracket \tau \rrbracket \text{ WF}} \text{ IH(1) on } \tau \quad \overline{FV(\ell_e) \in \Sigma} \text{ By inversion}}{\overline{\Sigma; \Psi, c \vdash \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket \text{ WF}} \text{ CG-wff-monad}}}{\overline{\Sigma; \Psi \vdash c \Rightarrow \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket \text{ WF}} \text{ CG-wff-constraint}}$$

□

**Lemma 5.23** (FG  $\rightsquigarrow$  CG: Free variable lemma).  $\forall \tau, \mathbf{A}$ . *The following hold*

1.  $FV(\llbracket \tau \rrbracket) \subseteq FV(\tau)$
2.  $FV(\llbracket \mathbf{A} \rrbracket) \subseteq FV(\mathbf{A})$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $\mathbf{A}$

Proof for (1)

Let  $\tau = \mathbf{A}^{\ell_i}$

$$\begin{aligned} & FV(\llbracket \mathbf{A}^{\ell_i} \rrbracket) \\ &= FV(\text{Labeled } \ell_i \llbracket \mathbf{A} \rrbracket) \quad \text{Definition 5.19} \\ &= FV(\ell_i) \cup FV(\llbracket \mathbf{A} \rrbracket) \\ &\subseteq FV(\ell_i) \cup FV(\mathbf{A}) \quad \text{IH(2) on } \mathbf{A} \\ &= FV(\mathbf{A}^{\ell_i}) \end{aligned}$$

Proof for (2)

1.  $\mathbf{A} = \mathbf{b}$ :

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ &= FV(\mathbf{b}) \quad \text{Definition 5.19} \\ &\subseteq FV(\mathbf{b}) \end{aligned}$$

2.  $\mathbf{A} = \text{unit}$ :

$$\begin{aligned} & FV(\llbracket \text{unit} \rrbracket) \\ &= FV(\text{unit}) \quad \text{Definition 5.19} \\ &\subseteq FV(\text{unit}) \end{aligned}$$

$$\begin{aligned}
3. \mathbf{A} &= \tau_1 \xrightarrow{\ell_e} \tau_2: \\
& \text{FV}(\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle) \\
&= \text{FV}(\langle \tau_1 \rangle \rightarrow \mathbb{C} \ell_e \perp \langle \tau_2 \rangle) && \text{Definition 5.19} \\
&= \text{FV}(\langle \tau_1 \rangle) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_2 \rangle) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_2) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 \xrightarrow{\ell_e} \tau_2)
\end{aligned}$$

$$\begin{aligned}
4. \mathbf{A} &= \tau_1 \times \tau_2: \\
& \text{FV}(\langle \tau_1 \times \tau_2 \rangle) \\
&= \text{FV}(\langle \tau_1 \rangle \times \langle \tau_2 \rangle) && \text{Definition 5.19} \\
&= \text{FV}(\langle \tau_1 \rangle) \cup \text{FV}(\langle \tau_2 \rangle) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\tau_2) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 \times \tau_2)
\end{aligned}$$

$$\begin{aligned}
5. \mathbf{A} &= \tau_1 + \tau_2: \\
& \text{FV}(\langle \tau_1 + \tau_2 \rangle) \\
&= \text{FV}(\langle \tau_1 \rangle + \langle \tau_2 \rangle) && \text{Definition 5.19} \\
&= \text{FV}(\langle \tau_1 \rangle) \cup \text{FV}(\langle \tau_2 \rangle) \\
&\subseteq \text{FV}(\tau_1) \cup \text{FV}(\tau_2) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \text{FV}(\tau_1 + \tau_2)
\end{aligned}$$

$$\begin{aligned}
6. \mathbf{A} &= \text{ref } \tau_i: \\
& \text{Let } \tau_i = \mathbf{A}_i^{\ell_i} \\
& \text{FV}(\langle \text{ref } \tau_i \rangle) \\
&= \text{FV}(\text{ref } \ell_i \langle \mathbf{A}_i \rangle) && \text{Definition 5.19} \\
&= \text{FV}(\ell_i) \cup \text{FV}(\langle \mathbf{A}_i \rangle) \\
&\subseteq \text{FV}(\ell_i) \cup \text{FV}(\mathbf{A}_i) && \text{IH(2) on } \mathbf{A}_i \\
&= \text{FV}(\text{ref } \mathbf{A}_i^{\ell_i}) \\
&= \text{FV}(\text{ref } \tau_i)
\end{aligned}$$

$$\begin{aligned}
7. \mathbf{A} &= \forall \alpha. (\ell_e, \tau_i): \\
& \text{FV}(\langle \forall \alpha. (\ell_e, \tau_i) \rangle) \\
&= \text{FV}(\forall \alpha. \mathbb{C} \ell_e \perp \langle \tau_i \rangle) && \text{Definition 5.19} \\
&= \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_i \rangle) \\
&\subseteq \text{FV}(\ell_e) \cup \text{FV}(\tau_i) && \text{IH(1) on } \tau_i \\
&= \text{FV}(\forall \alpha. (\ell_e, \tau_i))
\end{aligned}$$

$$\begin{aligned}
8. \mathbf{A} &= c \xrightarrow{\ell_e} \tau_i: \\
& \text{FV}(\langle c \xrightarrow{\ell_e} \tau_i \rangle) \\
&= \text{FV}(c) \cup \text{FV}(\mathbb{C} \ell_e \perp \langle \tau_i \rangle) && \text{Definition 5.19} \\
&= \text{FV}(c) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_i \rangle) \\
&\subseteq \text{FV}(c) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_i) && \text{IH(1) on } \tau_i \\
&= \text{FV}(c \xrightarrow{\ell_e} \tau_i)
\end{aligned}$$

□

**Lemma 5.24** (FG  $\rightsquigarrow$  CG: Substitution lemma).  $\forall \tau, \mathbf{A} \text{ s.t. } \vdash \tau \text{ WF and } \vdash \mathbf{A} \text{ WF. The following hold}$

$$1. \llbracket \tau \rrbracket[\ell/\alpha] = \llbracket (\tau[\ell/\alpha]) \rrbracket$$

$$2. \llbracket \mathbf{A} \rrbracket[\ell/\alpha] = \llbracket \mathbf{A}[\ell/\alpha] \rrbracket$$

*Proof.* Proof by simultaneous induction on  $\tau$  and  $\mathbf{A}$

Proof for (1)

$$\begin{aligned} \text{Let } \tau &= \mathbf{A}^{\ell_i} \\ & \llbracket (\mathbf{A}^{\ell_i}) \rrbracket[\ell/\alpha] \\ &= \llbracket \text{Labeled } \ell_i \ (\mathbf{A}) \rrbracket[\ell/\alpha] && \text{Definition 5.19} \\ &= \llbracket \text{Labeled } \ell_i[\ell/\alpha] \ (\mathbf{A}) \rrbracket[\ell/\alpha] \\ &= \llbracket \text{Labeled } \ell_i[\ell/\alpha] \ (\mathbf{A}[\ell/\alpha]) \rrbracket && \text{IH(2) on } \mathbf{A} \\ &= \llbracket \mathbf{A}[\ell/\alpha]^{\ell_i[\ell/\alpha]} \rrbracket \\ &= \llbracket \mathbf{A}^{\ell_i}[\ell/\alpha] \rrbracket \end{aligned}$$

Proof for (2)

1.  $\mathbf{A} = \mathbf{b}$ :

$$\begin{aligned} & \llbracket (\mathbf{b}) \rrbracket[\ell/\alpha] \\ &= \llbracket (\mathbf{b})[\ell/\alpha] \rrbracket && \text{Definition 5.19} \\ &= \llbracket \mathbf{b} \rrbracket \\ &= \llbracket \mathbf{b} \rrbracket \\ &= \llbracket \mathbf{b}[\ell/\alpha] \rrbracket \end{aligned}$$

2.  $\mathbf{A} = \text{unit}$ :

$$\begin{aligned} & \llbracket (\text{unit}) \rrbracket[\ell/\alpha] \\ &= \llbracket (\text{unit})[\ell/\alpha] \rrbracket && \text{Definition 5.19} \\ &= \llbracket \text{unit} \rrbracket \\ &= \llbracket \text{unit} \rrbracket \\ &= \llbracket \text{unit}[\ell/\alpha] \rrbracket \subseteq \llbracket \text{unit} \rrbracket \end{aligned}$$

3.  $\mathbf{A} = \tau_1 \xrightarrow{\ell_e} \tau_2$ :

$$\begin{aligned} & \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2) \rrbracket[\ell/\alpha] \\ &= \llbracket (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2) \rrbracket[\ell/\alpha] && \text{Definition 5.19} \\ &= \llbracket (\tau_1)[\ell/\alpha] \rightarrow \mathbb{C} \ell_e[\ell/\alpha] \perp (\tau_2)[\ell/\alpha] \rrbracket \\ &= \llbracket (\tau_1[\ell/\alpha]) \rightarrow \mathbb{C} \ell_e[\ell/\alpha] \perp (\tau_2[\ell/\alpha]) \rrbracket && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\ &= \llbracket (\tau_1[\ell/\alpha] \xrightarrow{\ell_e[\ell/\alpha]} \tau_2[\ell/\alpha]) \rrbracket \\ &= \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2)[\ell/\alpha] \rrbracket \end{aligned}$$

4.  $\mathbf{A} = \tau_1 \times \tau_2$ :

$$\begin{aligned} & \llbracket (\tau_1 \times \tau_2) \rrbracket[\ell/\alpha] \\ &= \llbracket (\tau_1)[\ell/\alpha] \times (\tau_2)[\ell/\alpha] \rrbracket && \text{Definition 5.19} \\ &= \llbracket (\tau_1[\ell/\alpha]) \times (\tau_2[\ell/\alpha]) \rrbracket && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\ &= \llbracket (\tau_1[\ell/\alpha] \times \tau_2[\ell/\alpha]) \rrbracket \\ &= \llbracket (\tau_1 \times \tau_2)[\ell/\alpha] \rrbracket \end{aligned}$$

5.  $\mathbf{A} = \tau_1 + \tau_2$ :

$$\begin{aligned} & \llbracket (\tau_1 + \tau_2) \rrbracket[\ell/\alpha] \\ &= \llbracket (\tau_1)[\ell/\alpha] + (\tau_2)[\ell/\alpha] \rrbracket && \text{Definition 5.19} \\ &= \llbracket (\tau_1[\ell/\alpha]) + (\tau_2[\ell/\alpha]) \rrbracket && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\ &= \llbracket (\tau_1[\ell/\alpha] + \tau_2[\ell/\alpha]) \rrbracket \\ &= \llbracket (\tau_1 + \tau_2)[\ell/\alpha] \rrbracket \end{aligned}$$

6.  $A = \text{ref } \tau_i$ :

$$\begin{aligned}
& \text{Let } \tau_i = A_i^{\ell_i} \\
& \quad (\llbracket \text{ref } \tau_i \rrbracket) [\ell/\alpha] \\
& = \quad (\text{ref } \ell_i \llbracket A_i \rrbracket) [\ell/\alpha] \quad \text{Definition 5.19} \\
& = \quad (\text{ref } \ell_i \llbracket A_i \rrbracket) \quad \text{Lemma 5.22} \\
& = \quad (\text{ref } A_i^{\ell_i}) \quad \text{since } \vdash \text{ref } \tau_i \text{ WF} \\
& = \quad (\llbracket \text{ref } \tau_i [\ell/\alpha] \rrbracket) \\
& = \quad (\llbracket \text{ref } \tau_i \rrbracket [\ell/\alpha])
\end{aligned}$$

7.  $A = \forall \alpha. (\ell_e, \tau_i)$ :

$$\begin{aligned}
& \quad (\llbracket \forall \alpha. (\ell_e, \tau_i) \rrbracket) [\ell/\alpha] \\
& = \quad (\forall \alpha. \mathbb{C} \ell_e \perp \llbracket \tau_i \rrbracket) [\ell/\alpha] \quad \text{Definition 5.19} \\
& = \quad (\forall \alpha. \mathbb{C} \ell_e [\ell/\alpha] \perp \llbracket \tau_i \rrbracket [\ell/\alpha]) \\
& = \quad (\forall \alpha. \mathbb{C} \ell_e [\ell/\alpha] \perp \llbracket \tau_i [\ell/\alpha] \rrbracket) \quad \text{IH(1) on } \tau_i \\
& = \quad (\llbracket \forall \alpha. (\ell_e [\ell/\alpha], \tau_i [\ell/\alpha]) \rrbracket) \\
& = \quad (\llbracket \forall \alpha. (\ell_e, \tau_i) \rrbracket) [\ell/\alpha]
\end{aligned}$$

8.  $A = c \stackrel{\ell_c}{\Rightarrow} \tau_i$ :

$$\begin{aligned}
& \quad (\llbracket c \stackrel{\ell_c}{\Rightarrow} \tau_i \rrbracket) [\ell/\alpha] \\
& = \quad (c \Rightarrow \mathbb{C} \ell_e \perp \llbracket \tau \rrbracket) [\ell/\alpha] \quad \text{Definition 5.19} \\
& = \quad c [\ell/\alpha] \Rightarrow (\mathbb{C} \ell_e [\ell/\alpha] \perp \llbracket \tau \rrbracket [\ell/\alpha]) \\
& = \quad c [\ell/\alpha] \Rightarrow (\mathbb{C} \ell_e [\ell/\alpha] \perp \llbracket \tau [\ell/\alpha] \rrbracket) \quad \text{IH(1) on } \tau_i \\
& = \quad (\llbracket c \stackrel{\ell_c}{\Rightarrow} \tau_i \rrbracket) [\ell/\alpha]
\end{aligned}$$

□

### 5.2.3 Model for FG to CG translation

**Definition 5.25** ( ${}^s\theta_2$  extends  ${}^s\theta_1$ ).  ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$

$$\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$$

**Definition 5.26** ( $\hat{\beta}_2$  extends  $\hat{\beta}_1$ ).  $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$

$$\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$$

**Definition 5.27** (Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\
[\mathbf{unit}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\
&\quad (s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 + \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \mathbf{inl} \, {}^s v, \mathbf{inl} \, {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}}\} \cup \\
&\quad \{(s\theta, m, \mathbf{inr} \, {}^s v, \mathbf{inr} \, {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \\
&\quad \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \implies \\
&\quad ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}'}\} \\
[\forall \alpha. (\ell_e, \tau)]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \Lambda e_s, \Lambda e_t) \mid \\
&\quad \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}', \ell' \in \mathcal{L}. ({}^s \theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'}\} \\
[c \xrightarrow{\ell_e} \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \nu e_s, \nu e_t) \mid \\
&\quad \mathcal{L} \models c \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau]_E^{\hat{\beta}'}\} \\
[\mathbf{ref} \, \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, a_s, a_t) \mid {}^s \theta(a_s) = \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\
[\mathbf{A}^{\ell'}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, \mathbf{Lb}({}^t v)) \mid ({}^s \theta, m, {}^s v, {}^t v) \in [\mathbf{A}]_V^{\hat{\beta}}\}
\end{aligned}$$

**Definition 5.28** (Unary expression relation).

$$\begin{aligned}
[\tau]_E^{\hat{\beta}} &\triangleq \{(s\theta, n, e_s, e_t) \mid \\
&\quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \\
&\quad \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \\
&\quad \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}\}
\end{aligned}$$

**Definition 5.29** (Unary heap well formedness).

$$\begin{aligned}
(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta &\triangleq \text{dom}({}^s \theta) \subseteq \text{dom}(H_s) \wedge \\
&\quad \hat{\beta} \subseteq (\text{dom}({}^s \theta) \times \text{dom}(H_t)) \wedge \\
&\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s \theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s \theta(a_1)]_V^{\hat{\beta}}
\end{aligned}$$

**Definition 5.30** (Value substitution).  $\delta^s : \text{Var} \mapsto \text{Val}$ ,  $\delta^t : \text{Var} \mapsto \text{Val}$

**Definition 5.31** (Unary interpretation of  $\Gamma$ ).

$$\begin{aligned}
[\Gamma]_V^{\hat{\beta}} &\triangleq \{(s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\
&\quad \forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}\}
\end{aligned}$$

#### 5.2.4 Soundness proof for FG to CG translation

**Lemma 5.32** (Monotonicity).  $\forall {}^s \theta, {}^s \theta', n, {}^s v, {}^t v, n', \beta, \beta'$ .

1.  $\forall \mathbf{A}. ({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{A}]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\mathbf{A}]_V^{\hat{\beta}'}$
2.  $\forall \tau. ({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$



*Proof.* Proof by simultaneous induction on  $A$  and  $\tau$

Proof of statement (1)

We case analyze  $A$  in the last step

1. Case  $b$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [b]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [b]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [b]_V^{\hat{\beta}}$  therefore from Definition 5.27 we know that  ${}^sv \in \llbracket b \rrbracket \wedge {}^tv \in \llbracket b \rrbracket$  and  ${}^sv = {}^tv$

Therefore from Definition 5.27 we get the desired

2. Case  $\text{unit}$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'}$$

Since  $({}^s\theta, n, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}}$  therefore from Definition 5.27 we know that  ${}^sv \in \llbracket \text{unit} \rrbracket \wedge {}^tv \in \llbracket \text{unit} \rrbracket$

Therefore from Definition 5.27 we get the desired

3. Case  $\tau_1 \times \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 5.27 we know that  ${}^sv = ({}^sv_1, {}^sv_2)$  and  ${}^tv = ({}^tv_1, {}^tv_2)$ .

We also know that  $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$  and  $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1:  $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$  (From Statement (2))

IH2:  $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$  (From Statement (2))

Therefore from Definition 5.27, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case  $\tau_1 + \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 5.27 two cases arise

(a)  ${}^sv = \text{inl}({}^sv')$  and  ${}^tv = \text{inl}({}^tv')$ :

$$\text{IH: } ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{From Statement (2)})$$

Therefore from Definition 5.27 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_{\hat{\beta}'}^{\hat{\beta}'}$$

(b)  ${}^sv = \text{inr}({}^sv')$  and  ${}^tv = \text{inr}({}^tv')$ :

Symmetric reasoning as in the previous case

5. Case  $\tau_1 \xrightarrow{\ell_e} \tau_2$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 5.27 we know that

${}^sv$  is of the form  $\lambda x.e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\lambda x.e_t$  (for some  $e_t$ ) s.t

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}' . ({}^s\theta', j, {}^sv_1, {}^tv_1) \in [\tau_1]_{\hat{\beta}'}^{\hat{\beta}'} \implies \\ ({}^s\theta', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{A0}) \end{aligned}$$

Similarly from Definition 5.27 we are required to prove

$$\begin{aligned} \forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{\hat{\beta}''}^{\hat{\beta}''} \implies \\ ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{\hat{\beta}''}^{\hat{\beta}''} \end{aligned}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{\hat{\beta}''}^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{\hat{\beta}''}^{\hat{\beta}''}$$

Instantiating (A0) with  ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}''$  since

${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta$ ,  $k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{\hat{\beta}''}^{\hat{\beta}''}$$

6. Case  $\forall\alpha.(\ell_e, \tau)$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}'}$$

From Definition 5.27 we know that

${}^sv$  is of the form  $\Lambda e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\Lambda e_t$  (for some  $e_t$ ) s.t

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}', \ell' \in \mathcal{L}. ({}^s\theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'_1} \quad (\text{F0})$$

Similarly from Definition 5.27 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'', \ell'' \in \mathcal{L}. ({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'', \ell'' \in \mathcal{L}$$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

Instantiating (F0) with  ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}'', \ell''$  since

${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

7. Case  $c \stackrel{\ell_e}{\Rightarrow} \tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}'}$$

From Definition 5.27 we know that

${}^sv$  is of the form  $\nu e_s$  (for some  $e_s$ ) and  ${}^tv$  is of the form  $\nu e_t$  (for some  $e_t$ ) s.t

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}', \mathcal{L} \models c \implies ({}^s\theta', j, e_s, e_t) \in [\tau]_E^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 5.27 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'', \mathcal{L} \models c \implies ({}^s\theta'', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t } \mathcal{L} \models c$$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

Instantiating (C0) with  ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}''$  since  
 ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta$ ,  $k < n' < n$  and  $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$  therefore we get  
 $({}^s\theta'', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$

8. Case ref  $\tau$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{ref } \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

From Definition 5.27 we know that  ${}^sv = a_s$  and  ${}^tv = a_t$ . We also know that  
 ${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$

From Definition 5.27, Definition 5.25 and Definition 5.26 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

Proof of Statement (2)

Let  $\tau = \mathbf{A}^{\ell''}$ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 5.27 we know that

$$\exists {}^tv_i. {}^tv = \text{Lb}({}^tv_i) \text{ and } ({}^s\theta, n, {}^sv, {}^tv_i) \in [\mathbf{A}]_V^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}'}$$

This means from Definition 5.27 we need to prove

$$({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_V^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_V^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

**Lemma 5.33** (Unary monotonicity for  $\Gamma$ ).  $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$ .

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

*Proof.* Given:  $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

$$\text{To prove: } (\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

From Definition 5.31 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$$

And again from Definition 5.31 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

- $dom(\Gamma) \subseteq dom(\delta^s) \wedge dom(\Gamma) \subseteq dom(\delta^t)$ :

Given

- $\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall x_i \in dom(\Gamma). (^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 5.32 we get

$$\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

□

**Lemma 5.34** (Unary monotonicity for  $H$ ).  $\forall ^s\theta, H_s, H_t, n, n', \hat{\beta}$ .

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta$$

*Proof.* Given:  $(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n$

To prove:  $(n', H_s, H_t) \triangleright^{\hat{\beta}'} ^s\theta$

From Definition 5.29 it is given that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 5.29 we are required to prove that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

- $dom(^s\theta) \subseteq dom(H_S)$ :

Given

- $\hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t))$ :

Given

- $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$ :

Since we know that  $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$  (given)

Therefore from Lemma 5.32 we get

$$\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

□

**Lemma 5.35** (Coercion lemma).  $\forall H, e, v$ .

$$(H, e) \Downarrow_-^f (H', \text{Lb } v) \implies (H, \text{coerce\_taint } e) \Downarrow_-^f (H', \text{Lb } v)$$

*Proof.* Given:  $(H, e) \Downarrow_-^f (H', \text{Lb } v)$

To prove:  $(H, \text{coerce\_taint } e) \Downarrow_-^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and cg-app it suffices to prove that

$$(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_-^f (H', \text{Lb } v)$$

From cg-tolabeled it suffices to prove that  
 $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From cg-bind it suffices to prove that

1.  $(H, e) \Downarrow_{-}^f (H'_1, v_1)$ :

We are given that  $(H, e) \Downarrow_{-}^f (H', v)$  therefore we have  $H'_1 = H'$  and  $v'_1 = \text{Lb } v$

2.  $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$ :

It suffices to prove that

$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v)$ :

We get this directly from cg-unlabel

□

**Theorem 5.36** (Fundamental theorem).  $\forall \Sigma, \Psi, \Gamma, \tau, e_s, e_t, pc, \mathcal{L}, \delta^s, \delta^t, \sigma, {}^s\theta, n, \hat{\beta}$ .

$\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge$

$\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

$\implies$

$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

*Proof.* Proof by induction on the  $\rightsquigarrow$  relation

1. FC-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_{V}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, x \delta^s, \text{ret}(x) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

From Definition 5.28 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(x) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge \\ & ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  $i = 0, {}^s v = x \delta^s$ . Also from cg-ret we know that  ${}^t v = x \delta^t$  and  $H'_t = H_t$

And we are required to prove

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ : Given

(b)  $({}^s\theta, n, {}^sv, {}^tv) \in [\tau \sigma]_V^{\hat{\beta}}$ :

Since we are given  $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_V^{\hat{\beta}}$ , therefore from Definition 5.31 we get  $({}^s\theta, n, {}^sv, {}^tv) \in [\tau \sigma]_V^{\hat{\beta}}$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \mathbf{ret}(\mathbf{Lb} \lambda x. e_t)} \text{ FC-lam}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\lambda x. e_s) \delta^s, \mathbf{ret}(\mathbf{Lb} \lambda x. e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^{\hat{\beta}}$

From Definition 5.28 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^sv. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv) \implies \\ & \exists H'_t, {}^tv. (H_t, \mathbf{ret}(\mathbf{Lb}(\lambda x. e_t))) \delta^t \Downarrow^f (H'_t, {}^tv) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$  and given some  $i < n, {}^sv$  s.t  $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv)$

From fg-val we know that  ${}^sv = (\lambda x. e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret, cg-label and cg-FI we know that  $H'_t = H_t$  and  ${}^tv = (\mathbf{Lb}(\lambda x. e_t)) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ : Given

(b)  $({}^s\theta, n, \lambda x. e_s \delta^s, \mathbf{Lb}(\lambda x. e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}}$ :

From Definition 5.27 it suffices to prove that

$$({}^s\theta, n, \lambda x. e_s \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_V^{\hat{\beta}}$$

Again from Definition 5.27 it suffices to prove that

$$\begin{aligned} & \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'} \implies \\ & ({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \end{aligned}$$

This further means that given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$

And we are required to prove

$$({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \quad (\text{F-L0})$$

Since we are given  $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'}$ , therefore from Definition 5.31 and Lemma 5.33 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\}) \sigma]_{V'}^{\hat{\beta}'}$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2 \sigma]_E^{\hat{\beta}'}$$

We get (F-L0) directly from IH

### 3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))))} \text{FC-app}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-A0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have



$$\exists H'_{t_1}, {}^t v_1. (H_{t_1}, e_{t_1}) \Downarrow^f (H'_{t_1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s_1}, H'_{t_1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-A1.0})$$

Since we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 5.27 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 5.27 we know that  ${}^s v_1 = \lambda x. e'_s$  and  ${}^t v_i = \lambda x. e'_t$  s.t

$$\begin{aligned} \forall {}^s \theta''_1 \sqsupseteq {}^s \theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1. \\ ({}^s \theta''_1, l, {}^s v', {}^t v') \in [\tau_1 \sigma]_V^{\hat{\beta}''_1} \implies ({}^s \theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}''_1} \end{aligned} \quad (\text{F-A1})$$

IH2:

$$({}^s \theta'_1, n - j, e_{s_2} \delta^s, e_{t_2} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s_2}, H_{t_2}. (n - j, H_{s_2}, H_{t_2}) \triangleright^{\hat{\beta}'_1} {}^s \theta \wedge \forall k < n - j, {}^s v_2. (H_{s_2}, e_{s_2} \delta^s) \Downarrow_j (H'_{s_2}, {}^s v_2) \implies \\ \exists H'_{t_2}, {}^t v_2. (H_{t_2}, e_{t_2}) \Downarrow^f (H'_{t_2}, {}^t v_2 \delta^t) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s_2}, H'_{t_2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

We instantiate with  $H'_{s_1}, H'_{t_1}$ . And since we know that  $(H_s, (e_{s_1} e_{s_2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s_1}, e_{s_2} \delta^s) \Downarrow_k (H'_{s_2}, {}^s v_2)$ .

This means we have

$$\exists H'_{t_2}, {}^t v_2. (H_{t_2}, e_{t_2}) \Downarrow^f (H'_{t_2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s_2}, H'_{t_2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-A2})$$

We instantiate (F-A1) with  $\theta''_1$  as  $\theta'_2$ ,  ${}^s v'$  as  ${}^s v_2$ ,  ${}^t v'$  as  ${}^t v_2$ ,  $l$  as  $n - j - k$  and  $\hat{\beta}''_1$  as  $\hat{\beta}'_2$ . Therefore we get

$$({}^s \theta'_2, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'_2}$$

From Definition 5.28 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s_3}, {}^s v_3) \implies \\ \exists H'_{t_3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t_3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s_3}, H'_{t_3}) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_V^{\hat{\beta}'_3} \end{aligned}$$

Instantiating with  $H'_{s_2}, H'_{t_2}$ . since we know that  $(H_s, (e_{s_1} e_{s_2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists a < i - j - k < n - j - k$  s.t  $(H'_{s_2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s_3}, {}^s v_3)$

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_t) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3} \quad (\text{F-A3}) \end{aligned}$$

Let  $\tau_2 = A_2^{\ell_i}$ , since  $\tau_2 \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}$$

Therefore from Definition 5.27 we know that

$$({}^s \theta'_3, n - j - k - a, {}^s v_3, \text{Lb}({}^t v_{3i})) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose  $H'_t$  as  $H'_{t3}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{3i})$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$$

From Lemma 5.35 it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_3))$$

From cg-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.c \ b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind again it suffices to prove

- \*  $(H'_{t2}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel and (F-A1) we know that  $H'_{t31} = H'_{t2}$  and  ${}^t v_{t2} = {}^t v_i = \lambda x. e'_t$

- \*  $((c \ b) [{}^t v_2/b] [{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}:$

It suffices to prove that

$$((\lambda x. e'_t) {}^t v_2 \delta^t) \Downarrow {}^t v_{t21}$$

From cg-app we know that

$${}^t v_{t21} = e'_t [{}^t v_2/x] \delta^t$$

- \*  $(H'_{t2}, {}^t v_{21}) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From (F-A3) and (F-A3.1) we get the desired

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_3$  and  $\hat{\beta}'$  as  $\hat{\beta}'_3$ . From fg-app we know that  $i = j + k + a + 1$ ,  ${}^s v = {}^s v_3$  and  $H'_s = H'_{s3}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t3}$  and  ${}^t v = \text{Lb}({}^t v_3)$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3$  from (F-A3) and Lemma 5.34

Since  ${}^t v = \text{Lb}({}^t v_3)$  therefore from Definition 5.27 it suffices to prove that

$$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}$$

We get this directly from (F-A3) and Lemma 5.32

4. FC-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_s : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda e_t))} \text{FC-FI}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, (\Lambda e_s) \delta^s, \text{ret}(\text{Lb} \Lambda e_t) \delta^t) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{\hat{\beta}}^{\hat{\beta}}$

From Definition 5.28 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (\Lambda e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda e_t))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{\hat{\beta}'}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (\Lambda e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  ${}^s v = (\Lambda e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret, cg-label and cg-val we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\Lambda e_t)) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$ : Given

(b)  $({}^s\theta, n, \Lambda e_s \delta^s, \text{Lb}(\Lambda e_t) \delta^t) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{\hat{\beta}}^{\hat{\beta}}$ :

From Definition 5.27 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

Again from Definition 5.27 it suffices to prove that

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}', \ell' \in \mathcal{L}. ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha] \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

This further means that given  ${}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}', \ell' \in \mathcal{L}$

And we are required to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha] \sigma]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-F0})$$

We get (F-F0) directly from IH

5. FC-FE:

$$\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\forall \alpha. (\ell_e, \tau))^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s [] : \tau[\ell'/\alpha] \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b))))} \text{FG-FE}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_s \ \square]) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b[\ ])))) \delta^t) \in [\tau[\ell'/\alpha] \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_s \ \square]) \delta^s \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b[\ ])))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \\ \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau[\ell'/\alpha] \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (e_s \ \square]) \delta^s \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b[\ ])))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \\ \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau[\ell'/\alpha] \sigma]_V^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \ \delta^s, e_t \ \delta^t) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_s \ \square]) \delta^s \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n, H'_{s1}$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - \\ j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.0}) \end{aligned}$$

Since we know that  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 5.27 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 5.27 we know that  ${}^s v_1 = \Lambda e'_s$  and  ${}^t v_i = \Lambda e'_t$  s.t

$$\forall {}^s\theta''_1 \sqsupseteq {}^s\theta'_1, l < (n - j), \hat{\beta}''_1 \sqsupseteq \hat{\beta}'_1, \ell'' \in \mathcal{L}. ({}^s\theta''_1, l, e'_s, e'_t) \in [\tau[\ell''/\alpha] \sigma]_E^{\hat{\beta}''_1} \quad (\text{F-F1})$$

Therefore we instantiate (F-F1) with  $\theta''_1$  as  $\theta'_1$ ,  $l$  as  $(n - j - 1)$ ,  $\hat{\beta}''_1$  as  $\hat{\beta}'_1$  and  $\ell''$  as  $\ell'$ . Therefore we get

$$({}^s\theta'_1, n - j - 1, e'_s, e'_t) \in [\tau[\ell'/\alpha] \sigma]_{\hat{E}}^{\hat{\beta}'_2}$$

From Definition 5.28 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - 1, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'_1 \wedge \forall a < n - j - 1, {}^s v. (H_s, e'_s) \Downarrow_a (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_t, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_2. \\ (n - j - 1 - a, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - 1 - a, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}'_2} \end{aligned}$$

Since we know that  $(H_s, (e_s \ []]) \delta^s \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k = i - j - 1$  s.t.  $(H_{s1}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . We know that  $k = i - j - 1 < n - j - 1$ . Therefore instantiating with  $H'_{s1}, H'_{t1}, k$  we get

$$\begin{aligned} \exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_2. \\ (n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - 1 - a, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}'_2} \quad (\text{F-F3}) \end{aligned}$$

Let  $\tau[\ell'/\alpha] = \mathbf{A}_2^{\ell_i}$ , since  $\tau[\ell'/\alpha] \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}'_2}$$

Therefore from Definition 5.27 we know that

$$({}^s\theta'_2, n - j - 1 - k, {}^s v_2, \mathbf{Lb}^t v_{2i}) \in [\tau[\ell'/\alpha] \sigma]_{\hat{V}}^{\hat{\beta}'_2} \quad (\text{F-F3.1})$$

In order to prove (F-F0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\mathbf{Lb}^t v_{2i}$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b[])))) \delta^t) \Downarrow^f (H'_{t2}, \mathbf{Lb}^t v_{2i}):$$

From Lemma 5.35 it suffices to prove that

$$(H_t, \text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b[]))) \delta^t) \Downarrow^f (H'_{t2}, \mathbf{Lb}^t v_{2i})$$

From cg-bind it further suffices to show that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1)$ :

We get this directly from (F-F1.0)

- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.(b[])) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \mathbf{Lb}^t v_{2i})$ :

From cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t2})$ :

Since from (F-F1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \mathbf{Lb}^t v_i$

Therefore from cg-unlabel and (F-F1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t2} = {}^t v_i = \Lambda e'_t$

- $((b \ []) [{}^t v_{t2}/b] \delta^t) \Downarrow {}^t v_{t21}$ :

It suffices to prove that

$$((\Lambda e'_t) \ [] \delta^t) \Downarrow {}^t v_{t21}$$

From cg-FE and cg-val we know that

$${}^t v_{t21} = e'_t \delta^t$$

- $(H'_{t1}, {}^t v_{21}) \Downarrow^f (H'_{t2}, \mathbf{Lb}^t v_{2i})$ :

From (F-F3) we get the desired

(b)  $\exists^s \theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge (s\theta', n - i, {}^s v, {}^t v) \in [\tau[l'/\alpha] \sigma]_{V}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . From fg-FE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v_2$  and  $H'_s = H'_{s2}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t2}$  and  ${}^t v = \text{Lb}({}^t v_{2i})$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta'$  from (F-F3) and Lemma 5.34

Since  ${}^t v = {}^t v_2 = \text{Lb}({}^t v_{2i})$  therefore from Definition 5.27 it suffices to prove that  $({}^s \theta'_3, n - j - k - 1, {}^s v_2, {}^t v_2) \in [\tau[l'/\alpha] \sigma]_{V}^{\hat{\beta}'_3}$

We get this directly from (F-F3) and Lemma 5.32

## 6. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e_s : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\nu e_t))} \text{FG-CI}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (\nu e_s) \delta^s, \text{ret}(\text{Lb}(\nu e_t)) \delta^t) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{E}^{\hat{\beta}}$

From Definition 5.28 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, (\nu e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\nu e_t))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s \theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge (s\theta', n - i, {}^s v, {}^t v) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (\nu e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that  ${}^s v = (\nu e_s) \delta^s$ ,  $H'_s = H_s$  and  $i = 0$ . Also from cg-ret, cg-label and cg-val we know that  $H'_t = H_t$  and  ${}^t v = (\text{Lb}(\nu e_t)) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}'} s\theta' \wedge (s\theta', n, {}^s v, {}^t v) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose  ${}^s \theta'$  as  ${}^s \theta$  and  $\hat{\beta}'$  as  $\hat{\beta}$

(a)  $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ : Given

(b)  $({}^s \theta, n, \nu e_s \delta^s, \text{Lb}(\nu e_t) \delta^t) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}}$ :

From Definition 5.27 it suffices to prove that

$$({}^s \theta, n, \nu e_s \delta^s, (\nu e_t) \delta^t) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_{V}^{\hat{\beta}}$$

Again from Definition 5.27 it suffices to prove that

$$\forall {}^s \theta' \sqsupseteq^s \theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. \mathcal{L} \models c \implies ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}'}$$

This further means that given  ${}^s \theta' \sqsupseteq^s \theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$  s.t  $\mathcal{L} \models c \implies$

And we are required to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'} \quad (\text{F-C0})$$

We get (F-C0) directly from IH

## 7. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet))))} \text{FG-CE}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_s \bullet) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet)))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

This means from Definition 5.28 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_s \bullet) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \\ \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This further means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  and given some  $i < n, {}^s v$  s.t  $(H_s, (e_s \bullet) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \\ \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-C0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with  $H_s, H_t$ . And since we know that  $(H_s, (e_s \bullet) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n, H'_{s1}$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

This means we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - \\ j, {}^s v_1, {}^t v_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-C1.0}) \end{aligned}$$

Since we know that  $({}^s\theta'_1, n-j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\varepsilon} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$  therefore from Definition 5.27 we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s\theta'_1, n-j, {}^s v_1, {}^t v_i) \in \llbracket (c \xrightarrow{\ell_\varepsilon} \tau) \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

From Definition 5.27 we know that  ${}^s v_1 = \nu e'_s$  and  ${}^t v_i = \nu e'_t$  s.t

$$\forall {}^s\theta''_1 \sqsupseteq {}^s\theta'_1, l < (n-j), \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1, \ell'' \in \mathcal{L}. ({}^s\theta''_1, l, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}''_1} \quad (\text{F-C1})$$

Therefore we instantiate (F-C1) with  $\theta''_1$  as  $\theta'_1$ ,  $l$  as  $(n-j-1)$ ,  $\hat{\beta}''_1$  as  $\hat{\beta}'_1$  and  $\ell''$  as  $\ell'$ . Therefore we get

$$({}^s\theta'_1, n-j-1, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1}$$

From Definition 5.28 we have

$$\begin{aligned} \forall H_s, H_t. (n-j-1, H_s, H_t) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \forall a < n-j-1, {}^s v. (H_s, e'_s) \Downarrow_a (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_t, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \end{aligned}$$

$$(n-j-1-a, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n-j-1-a, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'_2}$$

Since we know that  $(H_s, (e_s \bullet) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k = i-j-1$  s.t  $(H_{s1}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$ . We know that  $k = i-j-1 < n-j-1$ . Therefore instantiating with  $H'_{s1}, H'_{t1}, k$  we get

$$\exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1.$$

$$(n-j-1-k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n-j-1-a, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'_2} \quad (\text{F-C3})$$

Let  $\tau = \mathbf{A}_2^{\ell_i}$ , since  $\tau \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s\theta'_2, n-j-1-k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'_2}$$

Therefore from Definition 5.27 we know that

$$({}^s\theta'_2, n-j-1-k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'_2} \quad (\text{F-C3.1})$$

In order to prove (F-C0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_{2i})$ . We need to prove:

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet)))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i})):$$

From Lemma 5.35 it suffices to prove that

$$(H_t, \text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.(b\bullet))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it further suffices to show that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1)$ :  
We get this directly from (F-C1.0)
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.(b\bullet))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :  
From cg-bind it suffices to prove that



- $(H'_{t1}, (\text{unlabel } a)[^t v_1/a] \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t2})$ :  
Since from (F-C1.1) we know that  $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$   
Therefore from cg-unlabel and (F-C1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t2} = {}^t v_i = \nu e'_t$
- $((b \bullet)[^t v_{t2}/b] \delta^t) \Downarrow {}^t v_{t21}$ :  
It suffices to prove that  $((\nu e'_t) \bullet \delta^t) \Downarrow {}^t v_{t21}$   
From cg-CE and cg-val we know that  ${}^t v_{t21} = e'_t \delta^t$
- $(H'_{t1}, {}^t v_{21}) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$ :  
From (F-C3) we get the desired

(b)  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . From fg-CE we know that  $i = j + k + 1$ ,  ${}^s v = {}^s v_2$  and  $H'_s = H'_{s2}$ . Also from the termination proof (previous point) we know that  $H'_t = H'_{t2}$  and  ${}^t v = \text{Lb}({}^t v_{2i})$

We get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta'$  from (F-C3) and Lemma 5.34

Since  ${}^t v = {}^t v_2 = \text{Lb}({}^t v_{2i})$  therefore from Definition 5.27 it suffices to prove that

$$({}^s \theta'_3, n - j - k - 1, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_3}$$

We get this directly from (F-C3) and Lemma 5.32

## 8. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{V'}^{\hat{\beta}}$

To prove:  $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \in [(\tau_1 \times \tau_2)^\perp \sigma]_{E'}^{\hat{\beta}}$

This means from Definition 5.28 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ . Also given some  $i < n, {}^s v_1, {}^s v_2$  s.t  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 5.28 we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$  therefore  $\exists k < i - j < n - j$  s.t  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose  $H_t$  as  $H'_{i2}$  and  ${}^t v$  as  $\text{Lb}({}^t v_1, {}^t v_2)$

$$(a) (H_t, (\text{bind}(e_{t1}, a. \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2)):$$

From cg-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$ :  
From (F-P1) we know that  $H'_{tb1} = H'_{t1}$  and  ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
From cg-bind it suffices to prove that
  - $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$ :  
From (F-P2) we know that  $H'_{tb2} = H'_{i2}$  and  ${}^t v_{tb2} = {}^t v_2$
  - $(H'_{t2}, \text{ret}(\text{Lb}(a, b))) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$ :  
We get this from cg-ret, (F-P1) and (F-P2)

(b)  $\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'}$ :  
 We choose  ${}^s\theta'$  as  ${}^s\theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$  and since from fg-prod  $i = j + k + 1$  and  $H'_s = H'_{s2}$ .  
 Therefore from (F-P2) and Lemma 5.34 we get

$$(n - i, H'_s, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta'$$

In order to prove  $({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{V'}^{\hat{\beta}'}$

From Definition 5.27 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s\theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_{V'}^{\hat{\beta}'_2}$$

Since  ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$  therefore we get the desired from (F-P1), (F-P2), Definition 5.27 and Lemma 5.32

9. FC-fst:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t) \in [(\tau_1 \sigma)]_E^{\hat{\beta}}$

This means from Definition 5.28 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'} \end{aligned}$$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1}$  therefore from Definition 5.27 we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 5.27 we know that  ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$  and  ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$  s.t

$$({}^s \theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$

Let  $\tau_1 = \mathbf{A}_1^{\ell_i}$ , since  $\tau_1 \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$  and

$$({}^s \theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\mathbf{A}_1^{\ell_i}]_{V'}^{\hat{\beta}'_1}$$

Therefore from Definition 5.27 we know that

$$({}^s \theta'_1, n - j, {}^s v_{i1}, \text{Lb}({}^t v_{i11})) \in [\mathbf{A}_1]_{V'}^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  ${}^t v_{i1} (= \text{Lb}({}^t v_{i11}))$  as we need to prove

$$(a) (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11})):$$

From Lemma 5.35 it suffices to prove that

$$(H_t, \text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :

From (F-F1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1 = \text{Lb}({}^t v_i)$

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11}))$ :

Again from cg-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :

Since  ${}^t v_1 = \text{Lb}({}^t v_{i1}, {}^t v_{i2})$  from (F-F1.1) and (F-F1.2) therefore we get the desired from cg-unlabel

So,  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b)))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{i11}))$ :

We get the desired from cg-fst and cg-ret and (F-F1.3)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{i1}) \in [\tau_1]_{V'}^{\hat{\beta}'_1}$$

We choose  ${}^s \theta'$  as  ${}^s \theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . And from fg-fst we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-F1) and Lemma 5.34 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$$

Since from fg-fst we know that  ${}^s v = {}^s v_{i1}$  therefore from (F-F1.2) and Lemma 5.32 we get

$$({}^s \theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'_1}$$

## 10. FC-snd:

Symmetric reasoning as in the FC-fst case

11. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))} \text{inl}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove:  $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))\delta^t) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))\delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 5.28 we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_t, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-IL1}) \end{aligned}$$

In order to prove (F-IL0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $(\text{Lb inl}({}^t v_1))$  and we need to prove:

(a)  $(H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1)))$ :

From cg-bind it suffices to prove that

i.  $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :

From (F-IL1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$

ii.  $(H'_{t1}, \text{ret}(\text{Lbinl}(a)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1)))$ :

From cg-ret and (F-IL1)

(b)  $\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$ :

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$ . Since from fg-inl we know that  $i = j + 1$  and  $H'_s = H'_{s1}$  therefore from (F-IL1) and Lemma 5.34 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$$

Now we need to prove  $({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$

Since  ${}^sv = \text{inl } {}^sv_1$  and  ${}^tv = \text{Lb}(\text{inl}({}^tv_1))$  therefore from Definition 5.27 it suffices to prove that

$$({}^s\theta', n - i, \text{inl } {}^sv_1, \text{inl } {}^tv_1) \in [(\tau_1 + \tau_2) \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$$

Since from (F-IL1) we know that  $({}^s\theta', n - j, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$

Therefore from Lemma 5.32 and Definition 5.27 we get

$$({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 + \tau_2) \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$$

12. FC-inr:

Symmetric reasoning as in the FC-inl case

13. FC-case:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}}$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau \sigma]_{\hat{\beta}'_E}^{\hat{\beta}'}$$

This means from Definition 5.28 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta \wedge \forall i < n, {}^sv.(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^sv) \implies \exists H'_t, {}^tv. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^tv) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [\tau \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'}$$

This means we are given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}'} {}^s\theta$ . Also given some  $i < n, {}^sv$  s.t  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^sv)$

And we need to prove

$$\exists H'_t, {}^tv. (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^tv) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [\tau \sigma]_{\hat{\beta}'_V}^{\hat{\beta}'} \quad (\text{F-C0})$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}.(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1.(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1.(H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < i < n$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1.(H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$  therefore from Definition 5.27 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

2 cases arise

- (a)  ${}^s v_1 = \text{inl}({}^s v_{i1})$  and  ${}^t v_i = \text{inl}({}^t v_{i1})$ :

Also from Lemma 5.33 and Definition 5.31 we know that

$$({}^s\theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\})]_V^{\hat{\beta}'_1}$$

**IH2:**

$$({}^s\theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}.(n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2.(H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < i - j < n - j$  s.t  $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_2.(H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-C2}) \end{aligned}$$

Let  $\tau = A^{\ell_i}$  and since we know that  $\tau \searrow \ell$  therefore we have  $\ell \sqsubseteq \ell_i$

Since we have  $({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}'_2}$

Therefore from Definition 5.27 we have

$$({}^s\theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [A^{\ell_i}]_V^{\hat{\beta}'_2} \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose  $H'_t$  as  $H'_{t2}$  and  ${}^t v$  as  ${}^t v_2 = \text{Lb}({}^t v_{2i})$

And we need to prove:

- i.  $(H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}^t v_{2i})$ :  
 From Lemma 5.35 it suffices to prove that  
 $(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}^t v_{2i})$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :  
 From (F-C1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}^t v_{2i})$ :  
 From cg-bind it suffices to prove that
  - $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :  
 Since from (F-C1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$  therefore from cg-unlabel we know that  
 $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$
  - $(\text{case}(b, x.e_{t1}, y.e_{t2}) [{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$ :  
 Since we know that in this case  ${}^t v_i = \text{inl}({}^t v_{i1})$   
 Therefore from cg-case we know that  ${}^t v_{t22} = e_{t1} [{}^t v_{i1}/x] \delta^t$
  - $(H'_{t1}, e_{t1} [{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \text{Lb}^t v_{2i})$ :  
 From (F-C2) and (F-C2.1) we get the desired

- ii.  $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'}$ :  
 We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$ . Since from fg-case we know that  $i = j + k + 1$  and  $H'_s = H'_{s2}$  therefore from (F-C2) and Lemma 5.34 we get

$$(n - i, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$$

Now we need to prove  $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

Since  ${}^s v = {}^s v_2$  and  ${}^t v = {}^t v_2$  and since from (F-C2) we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

Therefore from Lemma 5.32 and Definition 5.27 we get

$$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

- (b)  ${}^s v_1 = \text{inr}({}^s v_{i1})$  and  ${}^t v_1 = \text{inr}({}^t v_{i1})$ :

Symmetric reasoning as in the previous case

#### 14. FC-ref:

$$\frac{\Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new}(e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lb}b)))} \text{ref}$$

Also given is:  $({}^s \theta, n, \delta^s, \delta^t) \in \lfloor \Gamma \rfloor_V^{\hat{\beta}}$

To prove:  $({}^s \theta, n, \text{new}(e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lb}b))) \delta^t) \delta^t \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_E^{\hat{\beta}}$

This means from Definition 5.28 we have

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{new}(e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new}(a), b.\text{ret}(\text{Lb}b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'}$$



This means that given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, \text{new}(e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ .

And we are required to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a. \text{bind}(\text{new}(a), b. \text{ret}(\text{Lb} b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 5.28 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, \text{new}(e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore we know that  $\exists j < n$  s.t.  $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$ .

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-R1}) \end{aligned}$$

In order to prove (F-R0) we choose  $H'_t$  as  $H'_1 \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^t v = \text{Lb}(a_t)$ ,  ${}^s \theta'$  as  ${}^s \theta'_1 \cup \{a_s \mapsto \tau\}$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new}(a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1})$ :  
From (F-R1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{new}(a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v)$ :  
From cg-bind it suffices to prove that
  - i.  $(H'_{t1}, \text{new}(a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v_{t2})$ :  
From cg-new we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v = a_t$
  - ii.  $(H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_t)$ :  
From cg-ret we know that  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$  and  ${}^t v_t = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}$$

From (F-R1) we know that  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$  and since  $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$ ,  $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ ,  ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau\}$

Therefore from Definition 5.29 and Lemma 5.34 we get  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'$

To prove:  $({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_V^{\hat{\beta}'}$

Since we know that  ${}^s v = a_s$  and  ${}^t v = \text{Lb } a_t$  therefore we need to prove

$$({}^s\theta', n - i, a_s, \mathbf{Lb}(a_t)) \in \llbracket (\mathbf{ref} \ \tau)^\perp \sigma \rrbracket_V^{\hat{\beta}'}$$

From Definition 5.27 it suffices to prove that

$$({}^s\theta', n - i, a_s, a_t) \in \llbracket (\mathbf{ref} \ \tau) \sigma \rrbracket_V^{\hat{\beta}'}$$

Again from Definition 5.27 it suffices to prove that

$${}^s\theta'(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'$$

We get this by construction

15. FC-deref:

$$\frac{\Gamma \vdash_{pc} e_s : (\mathbf{ref} \ \tau)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \mathbf{coerce\_taint}(\mathbf{bind}(e_t, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.!b)))} \text{deref}$$

$$\text{Also given is: } ({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \rrbracket_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s\theta, n, !e \ \delta^s, \mathbf{coerce\_taint}(\mathbf{bind}(e_t, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.!b)))) \delta^t \in \llbracket \tau' \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 5.28 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, !e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \mathbf{coerce\_taint}(\mathbf{bind}(e_t, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that we are given some  $H_s, H_t$  s.t.  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t.  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \mathbf{coerce\_taint}(\mathbf{bind}(e_t, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau' \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

IIH:

$$({}^s\theta, n, e_s \ \delta^s, e_t \ \delta^t) \in \llbracket (\mathbf{ref} \ \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 5.28 we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \ \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\mathbf{ref} \ \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t.  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v)$

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \ \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\mathbf{ref} \ \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have  $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{\beta}'_1}$

From Definition 5.27 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 5.27 we know that  ${}^s v_1 = a_s$  and  ${}^t v_i = a_t$

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Since we are given that  $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$  therefore from Definition 5.29 we know that

$$({}^s\theta, n - 1, H_s(a_s), H_t(a_t)) \in [{}^s\theta(a_s)]_{\hat{\beta}}$$

which means we have

$$({}^s\theta, n - 1, H_s(a_s), H_t(a_t)) \in [\tau \sigma]_{\hat{\beta}}$$

From Lemma 5.37 we know that

$$({}^s\theta, n - 1, H_s(a_s), H_t(a_t)) \in [\tau' \sigma]_{\hat{\beta}}$$

Let  $\tau' = A^{\ell_i}$  since  $\tau' \searrow \ell$  therefore  $\ell \sqsubseteq \ell_i$

Let  $v_g = H_t(a_t)$  therefore from Definition 5.27 we have

$$({}^s\theta, n - 1, H_s(a_s), \text{Lb } v_{gi}) \in [\tau' \sigma]_{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose  $H'_t$  as  $H'_{t1}$  and  ${}^t v$  as  $H'_{t1}(a_t) = v_g = \text{Lb } v_{gi}$

$$(a) \ (H_t, \text{coerce\_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From Lemma 5.35 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi})$$

From cg-bind it suffices to prove

$$i. \ (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t1} = {}^t v_1$

$$ii. \ (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From cg-bind it suffices to prove that

$$A. \ (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that  ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i$

$$B. \ (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

We get the desired from CG-deref, (F-DR1.2) and (F-DR1.3)

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, \text{Lb } v_{gi}) \in [\tau' \sigma]_{\hat{\beta}'}$$

We choose  ${}^s\theta'$  as  ${}^s\theta'_1$  and  $\hat{\beta}'$  as  $\hat{\beta}'_1$

Therefore from (F-DR1) we get  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$  and since  $i = j + 1$  therefore

from Lemma 5.34 we get  $(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$

Since from (F-DR1.2) we know that  $(a_s, a_t) \in \hat{\beta}'_1$  and  ${}^s\theta'_1(a_s) = \tau$ . Also from (F-DR1) we have  $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$ . Therefore from Definition 5.28 we have  $(n - j - 1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s\theta'_1(a_s)]_{V'}^{\hat{\beta}'_1}$

Since  $i = j + 1$ ,  ${}^s\theta'_1(a_s) = \tau$ ,  $H'_{s1}(a_s) = {}^s v$  and  $H'_{t1}(a_t) = {}^t v_g = \text{Lb } v_g i$

Therefore we get  $({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau']_{V'}^{\hat{\beta}'_1}$

from (F-DR1.3) and Lemma 5.32

16. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_{s1} := e_{s2} : \text{unit} \rightsquigarrow} \text{ assign}$$

$$\text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())$$

Also given is:  $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V'}^{\hat{\beta}}$

To prove:

$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \in [\text{unit}]_E^{\hat{\beta}}$

This means from Definition 5.28 we are required to prove

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies$   
 $\exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f$   
 $(H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

This means that given some  $H_s, H_t$  s.t  $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$ . Also given some  $i < n, {}^s v$  s.t  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$\exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f$   
 $(H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$  (F-AN0)

IH1:

$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \tau)^\ell \sigma]_E^{\hat{\beta}}$

This means from Definition 5.28 we are required to prove

$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies$   
 $\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$   
 $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{V'}^{\hat{\beta}'_1}$

Instantiating with  $H_s, H_t$  and since we know that  $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists j < n$  s.t  $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that  $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$  therefore from Definition 5.27 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 5.27 this further means that

$${}^s \theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 5.28 we are required to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with  $H'_{s1}, H'_{t1}$  and since we know that  $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$  therefore  $\exists k < n - j$  s.t  $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'_2} \quad (\text{F-AN2}) \end{aligned}$$

In order to prove (F-AN0) we choose  $H'_t$  as  $H'_{t2}[a_t \mapsto {}^s v_2]$ ,  ${}^t v$  as ()

We need to prove

$$(a) \ (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

$$- (H_t, \text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From cg-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$$

where  ${}^t v_T = \text{Lb}({}^t v_{Ti})$

From cg-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$ :  
From (F-AN1) we know that  $H'_{t11} = H'_{t1}$  and  ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12})$ :  
From cg-bind it suffices to prove
  - $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13})$ :  
From (F-AN2) we know that  $H'_{t13} = H'_{t2}$  and  ${}^t v_{t13} = {}^t v_2$

–  $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b)[^t v_1/a][^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v_{t12})$ :

From cg-bind it suffices to prove that

\*  $(H'_{t1}, \text{unlabel } a[^t v_1/a][^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$ :

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{V'}^{\hat{\beta}'_1}$$

Therefore from cg-unlabel we know that  $H'_{t21} = H'_{t1}$  and  ${}^t v_{t21} = {}^t v_i = a_t$

\*  $(H'_{t1}, (c := b)[^t v_1/a][^t v_2/b][^t v_i/c] \delta^t) \Downarrow^f (H'_t, {}^t v)$ :

From cg-assign we know that  $H'_t = H'_{t1}[a_t \mapsto {}^t v_2]$  and  ${}^t v_{t12} = ()$

Since  ${}^t v_{t12} = {}^t v_{Ti} = ()$  therefore  ${}^t v_T = \text{Lb}()$

–  $(H'_T, \text{ret}()[^t v_T/d]) \Downarrow^f (H'_t, ())$ :

From cg-ret and cg-val

(b)  $\exists {}^s \theta' \sqsubseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'}$ :

We choose  ${}^s \theta'$  as  ${}^s \theta'_2$  and  $\hat{\beta}'$  as  $\hat{\beta}'_2$

In order to prove  $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  it suffices to prove

- $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 5.29 we get  $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$ :

Since from (F-AN2) we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 5.29 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s \theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'}$ :

$\forall (a_1, a_2) \in \hat{\beta}'_2$ .

- $a_1 = a_s$  and  $a_1 = a_t$ :

Since from (F-AN2) we know that  $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 5.25 we know that  ${}^s \theta'_2(a_1) = \tau$

Therefore from Lemma 5.32 we get

$$({}^s \theta'_2, n - i - 1, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$  and  $a_1 \neq a_t$ :

From (F-AN2) since we know that  $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$  therefore from Definition 5.29 we get

$$({}^s \theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

Since  $i = j + k + 1$  therefore from Lemma 5.32 we get

$$({}^s \theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

- $a_1 = a_s$  and  $a_1 \neq a_t$ :

This case cannot arise

- $a_1 \neq a_s$  and  $a_1 = a_t$ :

This case cannot arise

And in order to prove  $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

Since we know that  ${}^s v = ()$  and  ${}^t v = ()$  therefore from Definition 5.27 we get  $({}^s \theta', n -$

$i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

□

**Lemma 5.37** (Subtyping lemma). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \mathcal{L}, \hat{\beta}$ .

1.  $\forall A, A'$ .

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies [(A \sigma)]_V^{\hat{\beta}} \subseteq [(A' \sigma)]_V^{\hat{\beta}}$$

2.  $\forall \tau, \tau'$ .

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$$

*Proof.* Proof by simultaneous induction on  $A <: A'$  and  $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of  $A <: A'$  in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } [((\tau_1 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \xrightarrow{\ell_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau'_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{It suffices to prove: } \forall ({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in [((\tau_1 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^{\hat{\beta}}.$$

$$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in [((\tau'_1 \xrightarrow{\ell_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, m$  and  $\lambda x.e_s, (\lambda x.e_t)$  s.t

$$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in [((\tau_1 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 5.27 we are given:

$$\begin{aligned} \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-L0}) \end{aligned}$$

$$\text{And it suffices to prove: } ({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 5.27, it suffices to prove:

$$\begin{aligned} \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L1}) \end{aligned}$$

$$\text{This means that given } {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2 \text{ s.t } ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with  ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$ . Since we have  $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$  therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2 \sigma)]_E^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2}$$

## 2. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

$$\text{To prove: } [(\forall \alpha. (\ell_e, \tau_1) \sigma)]_V^{\hat{\beta}} \subseteq [(\forall \alpha. (\ell'_e, \tau_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{It suffices to prove: } \forall ({}^s\theta, m, \Lambda e_s, (\Lambda e_t)) \in [(\forall \alpha. (\ell_e, \tau_1) \sigma)]_V^{\hat{\beta}}.$$

$$({}^s\theta, m, \Lambda e_s, (\Lambda e_t)) \in [(\forall \alpha. (\ell'_e, \tau_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, m$  and  $\Lambda e_s, (\Lambda e_t)$  s.t

$$({}^s\theta, m, \Lambda e_s, (\Lambda e_t)) \in [(\forall \alpha. (\ell_e, \tau_1) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 5.27 we are given:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1, \ell'_1 \in \mathcal{L}. ({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau_1[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-F0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, \Lambda e_s, (\Lambda e_t)) \in [(\forall \alpha. (\ell'_e, \tau_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 5.27, it suffices to prove:

$$\forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2, \ell'_2 \in \mathcal{L}. ({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in [\tau_2[\ell'_2/\alpha] \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-F1})$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2, \ell'_2 \in \mathcal{L}$

And we need to prove

$$({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in [\tau_2[\ell'_2/\alpha] \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-F2})$$

Instantiating (S-F0) with  ${}^s\theta'_2, k, \hat{\beta}'_2, \ell'_2$  we get

$$({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in [\tau_1[\ell'_2/\alpha] \sigma]_E^{\hat{\beta}'_2}$$

$$\text{IH: } [(\tau_1[\ell'_2/\alpha] \sigma)]_E^{\hat{\beta}'_2} \subseteq [(\tau_2[\ell'_2/\alpha] \sigma)]_E^{\hat{\beta}'_2} \text{ (Statement 2(b))}$$

Finally using IH we get the desired.



3. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove:  $\llbracket (c_1 \xrightarrow{\ell_e} \tau_1) \sigma \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (c_2 \xrightarrow{\ell'_e} \tau_2) \sigma \rrbracket_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s\theta, m, \nu e_s, (\nu e_t)) \in \llbracket (c_1 \xrightarrow{\ell_e} \tau_1) \sigma \rrbracket_V^{\hat{\beta}}$ .

$({}^s\theta, m, \nu e_s, (\nu e_t)) \in \llbracket (c_2 \xrightarrow{\ell'_e} \tau_2) \sigma \rrbracket_V^{\hat{\beta}}$

This means that given some  ${}^s\theta, m$  and  $\nu e_s, (\nu e_t)$  s.t

$({}^s\theta, m, \nu e_s, (\nu e_t)) \in \llbracket (c_1 \xrightarrow{\ell_e} \tau_1) \sigma \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 5.27 we are given:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. \mathcal{L} \models c_1 \implies ({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-C0})$$

And it suffices to prove:  $({}^s\theta, m, \nu e_s, (\nu e_t)) \in \llbracket (c_2 \xrightarrow{\ell'_e} \tau_2) \sigma \rrbracket_V^{\hat{\beta}}$

Again from Definition 5.27, it suffices to prove:

$$\forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. \mathcal{L} \models c_2 \implies ({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-C1})$$

This means that given  ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$  s.t  $\mathcal{L} \models c_2$

And we need to prove

$$({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-C2})$$

Instantiating (S-C0) with  ${}^s\theta'_2, k, \hat{\beta}'_2$  and since we know that  $\mathcal{L} \models c_2 \sigma \implies c_1 \sigma$  therefore we get

$$({}^s\theta'_2, k, e_s \delta^s, e_t \delta^t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_2}$$

IH:  $\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_2} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_2}$  (Statement 2(b))

Finally using IH we get the desired.

4. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove:  $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1:  $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

IH2:  $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}. \quad ({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some  ${}^s\theta, n$  and  ${}^sv_1, {}^sv_2, {}^tv_1, {}^tv_2$  s.t

$$({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 5.27 we are given:

$$({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove:  $({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 5.27, it suffices to prove:

$$({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that  $({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$  therefore from IH1 we have  $({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$

Similarly since we have  $({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau_2 \sigma]_V^{\hat{\beta}}$  from (S-P0) therefore from IH2 we have  $({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau'_2 \sigma]_V^{\hat{\beta}}$

## 5. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove:  $[((\tau_1 + \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 + \tau'_2) \sigma)]_V^{\hat{\beta}}$

IH1:  $[((\tau_1) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1) \sigma)]_V^{\hat{\beta}}$  (Statement 2(a))

IH2:  $[((\tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_2) \sigma)]_V^{\hat{\beta}}$  (Statement 2(a))

It suffices to prove:  $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2) \sigma)]_V^{\hat{\beta}}. \quad ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2) \sigma)]_V^{\hat{\beta}}$

This means that given:  $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2) \sigma)]_V^{\hat{\beta}}$

And it suffices to prove:  $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2) \sigma)]_V^{\hat{\beta}}$

2 cases arise

(a)  ${}^sv = \text{inl } {}^sv_i$  and  ${}^tv = \text{inl } {}^tv_i$ :

From Definition 5.27 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b)  ${}^s v = \text{inr } {}^s v_i$  and  ${}^t v = \text{inr } {}^t v_i$ :

Symmetric reasoning as in the previous case

6. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove:  $\llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:  $\forall ({}^s \theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s \theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

We get this directly from Definition 5.27

7. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove:  $\llbracket ((\mathbf{b})) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{b})) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 5.27

8. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove:  $\llbracket ((\text{unit})) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{unit})) \rrbracket_V^{\hat{\beta}}$

Directly from Definition 5.27

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell' \sqsubseteq \ell'' \quad \mathcal{L} \vdash \mathbf{A} <: \mathbf{A}'}{\mathcal{L} \vdash \mathbf{A}^{\ell'} <: \mathbf{A}^{\ell''}} \text{FGsub-label}$$

To prove:  $\llbracket ((\mathbf{A}^{\ell'})) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\mathbf{A}^{\ell''})) \rrbracket_V^{\hat{\beta}}$

This means from Definition 5.27 we need to prove

$\forall ({}^s \theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell'} \rrbracket_V^{\hat{\beta}}. ({}^s \theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell''} \rrbracket_V^{\hat{\beta}}$

This means that given  $({}^s \theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell'} \rrbracket_V^{\hat{\beta}}$

From Definition 5.27 it further means that we are given

$({}^s \theta, n, {}^s v, {}^t v_i) \in \llbracket \mathbf{A} \rrbracket_V^{\hat{\beta}} \quad (\text{S-LB0})$

And we need to prove

$({}^s \theta, n, {}^s v, \text{Lb}({}^t v_i)) \in \llbracket \mathbf{A}^{\ell''} \rrbracket_V^{\hat{\beta}}$

Again from Definition 5.27 it suffices to prove that

$$({}^s\theta, n, {}^sv, {}^tv_i) \in [A']_{V}^{\hat{\beta}}$$

Since  $\ell' \sqsubseteq \ell''$  and  $A' <: A''$  therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given:  $\mathcal{L} \vdash \tau <: \tau'$

To prove:  $[(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$

This means we need to prove that

$$\forall (\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}. (\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means given  $(\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

This means from Definition 5.28 we have

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\ \exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{S-E0})$$

And it suffices to prove that  $({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$

Again from Definition 5.28 it means we need to prove

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^sv_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^sv_1) \implies \\ \exists H'_{t1}, {}^tv_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^tv_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^sv_1, {}^tv_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1}$$

This means that given some  $H_{s1}, H_{t1}$  s.t  $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} {}^s\theta$ . Also given some  $j < n, {}^sv_1$  s.t  $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^sv_1)$

And we need to prove

$$\exists H'_{t1}, {}^tv_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^tv_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^sv_1, {}^tv_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1} \quad (\text{S-E1})$$

Instantiating (S-E0) with  $H_{s1}, H_{t1}$  and with  $j, {}^sv_1$ . Then we get

$$\exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^sv_1, {}^tv_1) \in [\tau \sigma]_V^{\hat{\beta}'_1}$$

Since we have  $\tau <: \tau'$ . Therefore from IH (Statement 2(a)) we get

$$\exists H'_{t1}, {}^tv_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^tv_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^sv_1, {}^tv_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1}$$

□

**Theorem 5.38** (Deriving FG NI via compilation).  $\forall e_s, {}^sv_1, {}^sv_2, n_1, n_2, H'_{s1}, H'_{s2}, \perp$ .

Let  $\text{bool} = (\text{unit} + \text{unit})$

$$x : \text{bool}^\top \vdash_\perp e_s : \text{bool}^\perp \wedge$$

$$\emptyset \vdash_\perp {}^sv_1 : \text{bool}^\top \wedge \emptyset \vdash_\perp {}^sv_2 : \text{bool}^\top \wedge$$

$$(\emptyset, e_s[{}^sv_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^sv'_1) \wedge$$

$$(\emptyset, e_s[{}^sv_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^sv'_2) \wedge$$

$\implies$

$${}^sv'_1 = {}^sv'_2$$

*Proof.* From the FG to CG translation we know that  $\exists e_t$  s.t

$$x : \mathbf{bool}^\top \vdash e_s : \mathbf{bool}^\perp \rightsquigarrow e_t$$

Similarly we also know that  $\exists^t v_1, {}^t v_2$  s.t

$$\emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1 \text{ and } \emptyset \vdash {}^s v_2 : \mathbf{bool}^\top \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$$

From type preservation theorem (choosing  $\alpha = \bar{\beta} = \perp$ ) we know that

$$x : \mathbf{Labeled} \top \mathbf{bool} \vdash e_t : \mathbb{C} \perp \perp \mathbf{Labeled} \perp \mathbf{bool}$$

$$\emptyset \vdash {}^t v_1 : \mathbb{C} \perp \perp \mathbf{Labeled} \top \mathbf{bool}$$

$$\emptyset \vdash {}^t v_2 : \mathbb{C} \perp \perp \mathbf{Labeled} \top \mathbf{bool} \quad (\text{NI-1})$$

Since we have  $\emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1$

And since  ${}^s v_1$  and  ${}^t v_1$  are closed terms (from given and NI-1)

Therefore from Theorem 5.36 we have (we choose  $n > n_1$  and  $n > n_2$ )

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in \lfloor \mathbf{bool}^\top \rfloor_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 5.28 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \triangleright {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{11}) \in \lfloor \mathbf{bool}^\top \rfloor_V^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \triangleright {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{11}) \in \lfloor \mathbf{bool}^\top \rfloor_V^{\hat{\beta}'} \quad (\text{NI-2.1})$$

From Definition 5.27 we know that

$${}^t v_{11} = \mathbf{Lb}({}^t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i11}) \in \lfloor (\mathbf{unit} + \mathbf{unit}) \rfloor_V^{\hat{\beta}'}$$

Again from Definition 5.27 we know that

Either a)  ${}^s v_1 = \mathbf{inl}()$  and  ${}^t v_{i11} = \mathbf{inl}()$  or b)  ${}^s v_1 = \mathbf{inr}()$  and  ${}^t v_{i11} = \mathbf{inr}()$

But in either case we have that  $\emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.2})$

As a result we have  $\emptyset \vdash {}^t v_{11} : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.3})$

We give it typing derivation

$$\frac{\overline{\emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit})} \quad (\text{NI-2.2})}{\emptyset \vdash \mathbf{Lb}({}^t v_{i11}) : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit})}$$

From Definition 5.31 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_{11})) \in \lfloor x \mapsto \mathbf{bool}^\top \rfloor_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 5.36 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_{11}/x]) \in \lfloor \mathbf{bool}^\perp \rfloor_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 5.28 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v''_1. (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v''_1) \implies \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta'' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s1}, H'_{t1}) \triangleright {}^s \theta'' \wedge ({}^s \theta'', n - i, {}^s v''_1, {}^t v''_1) \in \lfloor \mathbf{bool}^\perp \rfloor_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_1, {}^s v'_1$  we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ (n - n_1, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp]_V^{\hat{\beta}''} \quad (\text{NI-2.5}) \end{aligned}$$

Since we have  $({}^s \theta', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp]_V^{\hat{\beta}''}$  therefore from Definition 5.27 we have

$$\exists {}^t v_{i1}. {}^t v'' = \mathbf{Lb}({}^t v_{i1}) \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [\mathbf{bool}]_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$  therefore from Definition 5.27 two cases arise

- ${}^s v'_1 = \mathbf{inl} {}^s v_{i11}$  and  ${}^t v_{i1} = \mathbf{inl} {}^t v_{i11}$ :

From Definition 5.27 we have

$$({}^s \theta', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v'_1 = \mathbf{inr} {}^s v_{i11}$  and  ${}^t v_{i1} = \mathbf{inr} {}^t v_{i11}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v'_1 = {}^t v_{i1}$

$$\text{Similarly with other substitution we have } (\emptyset, n, {}^s v_2, {}^t v_2) \in [\mathbf{bool}^\top]_E^\emptyset \quad (\text{NI-3})$$

Therefore from Definition 5.28 we have

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^\emptyset \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'}$$

Instantiating with  $\emptyset, \emptyset$  and from fg-val we know that  $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$ . Therefore we have

$$\exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'} \quad (\text{NI-3.1})$$

From Definition 5.27 we know that

$${}^t v_2 = \mathbf{Lb}({}^t v_{22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}'}$$

Again from Definition 5.27 we know that

Either a)  ${}^s v_2 = \mathbf{inl}()$  and  ${}^t v_{22} = \mathbf{inl}()$  or b)  ${}^s v_2 = \mathbf{inr}()$  and  ${}^t v_{22} = \mathbf{inr}()$

But in either case we have that  $\emptyset \vdash {}^t v_{22} : (\mathbf{unit} + \mathbf{unit})$  (NI-3.2)

As a result we have  $\emptyset \vdash {}^t v_{22} : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit})$  (NI-3.3)

We give it typing derivation

$$\frac{\overline{\emptyset \vdash {}^t v_{22} : (\mathbf{unit} + \mathbf{unit})} \quad (\text{NI-3.2})}{\emptyset \vdash \mathbf{Lb}({}^t v_{22}) : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit})}$$

From Definition 5.31 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \mathbf{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 5.36 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\mathbf{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-3.4})$$

From Definition 5.28 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v_2'' . (H_s, e_s[{}^s v_2/x]) \Downarrow_i (H'_{s2}, {}^s v_2'') \implies \\ \exists H'_{t2}, {}^t v_2'' . (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}' .$$

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_2'', {}^t v_2'') \in [\mathbf{bool}^\perp]_V^{\hat{\beta}''}$$

Instantiating with  $\emptyset, \emptyset, n_2, {}^s v_2'$  we get

$$\exists H'_{t2}, {}^t v_2'' . (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}' .$$

$$(n - n_1, H'_s, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_2'', {}^t v_2'') \in [\mathbf{bool}^\perp]_V^{\hat{\beta}''} \quad (\text{NI-3.5})$$

Since we have  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_2'') \in [\mathbf{bool}^\perp]_V^{\hat{\beta}''}$  therefore from Definition 5.27 we have

$$\exists {}^t v_{i2} . {}^t v_2'' = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [\mathbf{bool}]_V^{\hat{\beta}''}$$

Since  $({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$  therefore from Definition 5.27 two cases arise

- ${}^s v_2' = \text{inl } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inl } {}^t v_{i22}$ :

From Definition 5.27 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have  ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v_2' = \text{inr } {}^s v_{i22}$  and  ${}^t v_{i2} = \text{inr } {}^t v_{i22}$ :

Symmetric reasoning as in the previous case

So no matter which case arise we have  ${}^s v_2' = {}^t v_{i2}$

We know that  $\emptyset \vdash {}^t v_{11} : \text{Labeled } \top \text{ bool}$  (NI-2.3)

Also we have  $\emptyset \vdash {}^t v_{22} : \text{Labeled } \top \text{ bool}$  (NI-3.3)

Let  $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that  $x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{C} \perp \perp \text{ bool}$  by giving a typing derivation P2:

$$\frac{\frac{}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool}}{\text{CG-var}}}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{C} \perp \perp \text{ bool}} \text{CG-unlabel}$$

P1:

$$\frac{}{x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \text{ bool}} \text{From (NI-1)}$$

Main derivation:

$$\frac{\frac{}{P1} \quad \frac{}{P2}}{x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{C} \perp \perp \text{ bool}}$$

Say  $e_t[{}^t v_{11}/x]$  reduces in  $n_{t1}$  steps in (NI-2.5) and  $e_t[{}^t v_{22}/x]$  reduces in  $n_{t2}$  steps in (NI-3.5)

We instantiate Theorem 5.18 with  $e_T, {}^t v_{11}, {}^t v_{22}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$  and from (NI-2.5) and (NI-3.5) we have  ${}^t v_{i1} = {}^t v_{i2}$  and thus  ${}^s v_1' = {}^s v_2'$

□